

# Lecture 23: Blockchains

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CS 181S

Spring 2024

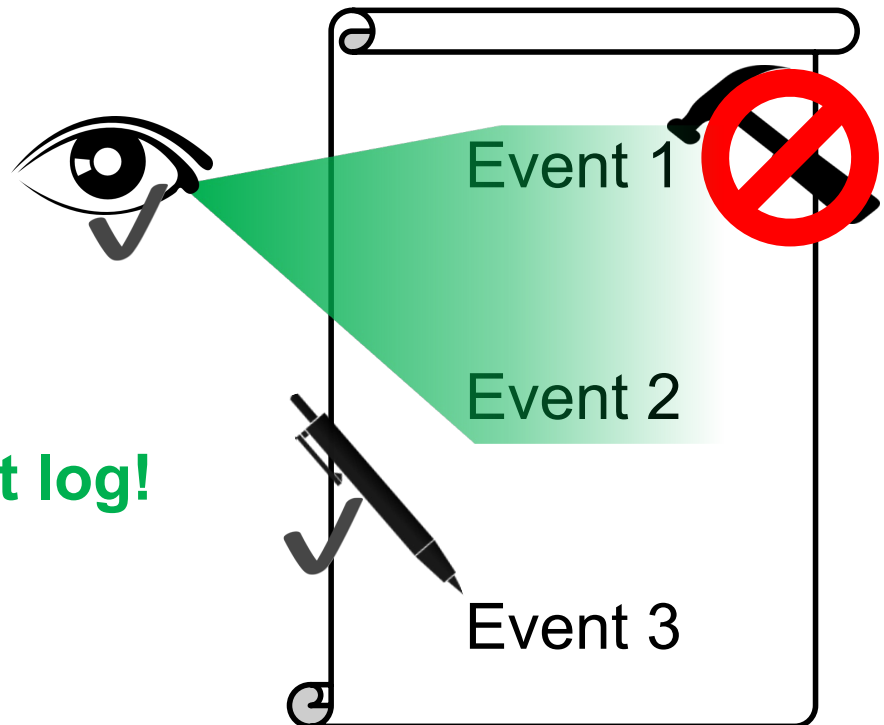
# Blockchain: A public tamper-proof log

Publicly visible

Publicly writable

Unmodifiable

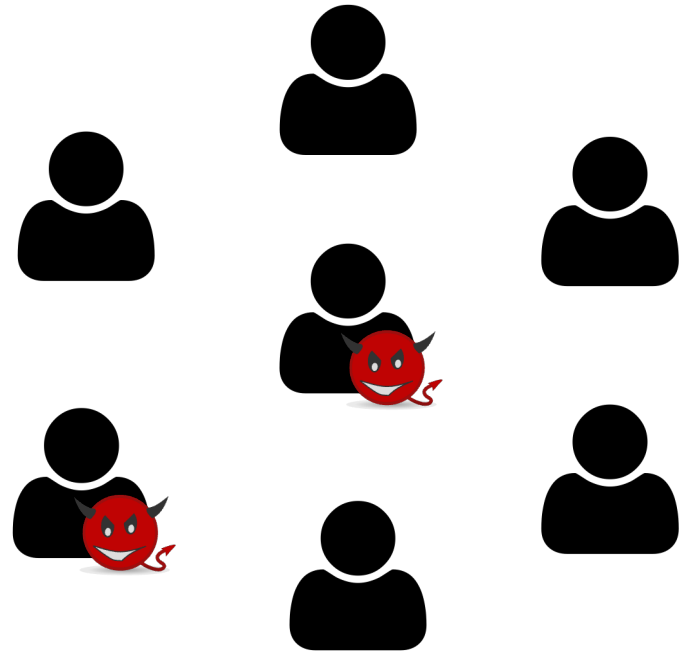
**Useful for an audit log!**



# Preventing Tampering

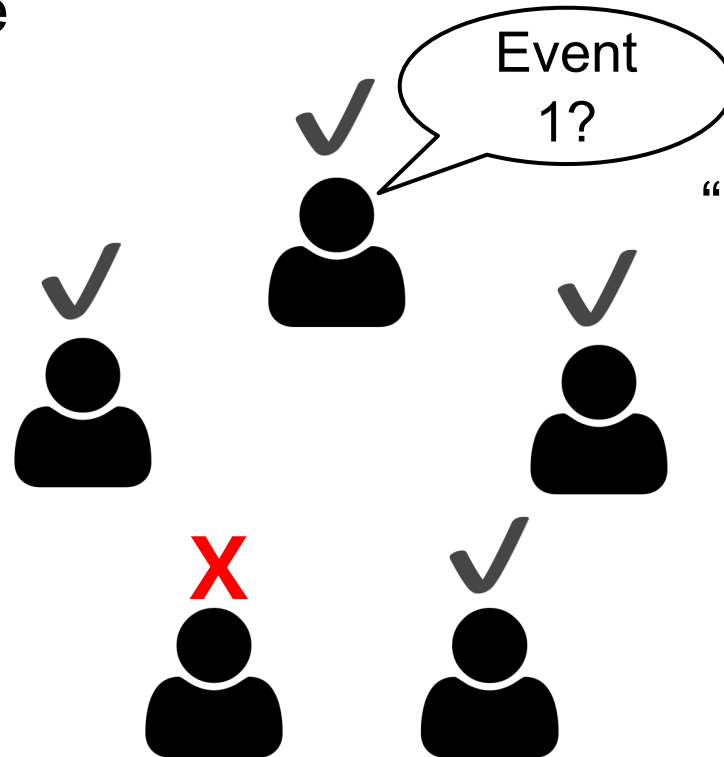


VS



# Traditional Consensus

Members Vote



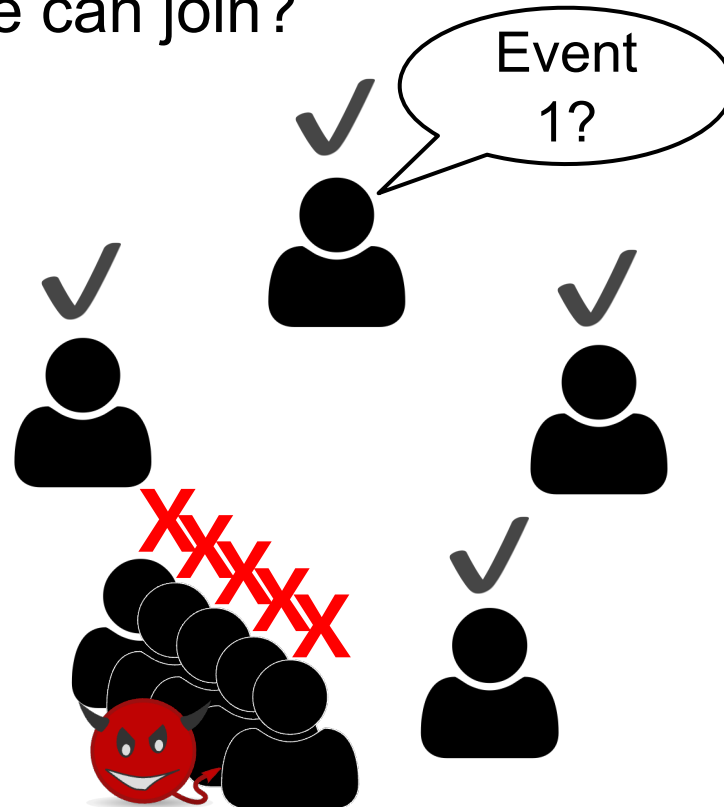
“Byzantine fault-tolerant  
(BFT) consensus”

Tolerates  $< 1/3$  faulty

Must know who  
everyone is!

# Sybil Attacks

What if anyone can join?



# Defending against Sybil

## Need a **scarce resource**

- BFT consensus uses identity – you only get one
- What else can we use?
  - Money (Proof of Stake)
  - Computational power (Proof of Work)

# COMPUTATION AS A SCARCE RESOURCE: PROOF OF WORK

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# Proof of Work: The basics

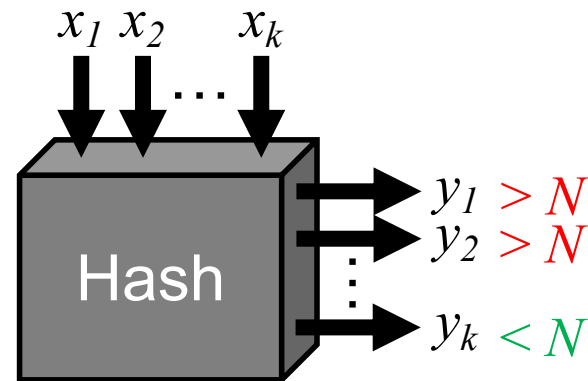
Find  $x$  such that  $\text{Hash}(x) < N$

This could take a while...

What about replays?

Add a nonce  $r$

Look for  $\text{Hash}(r \parallel x) < N$





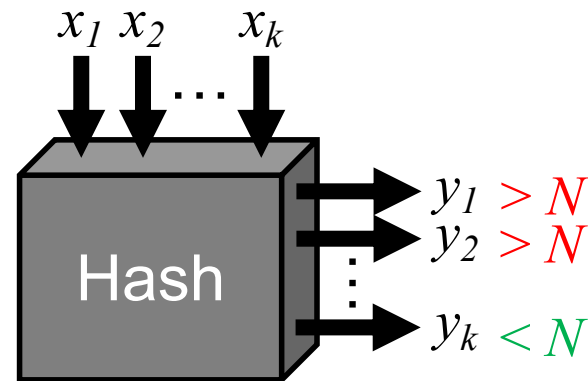
# Proof of Work: Building a log

Make the nonce useful

Use a message digest!

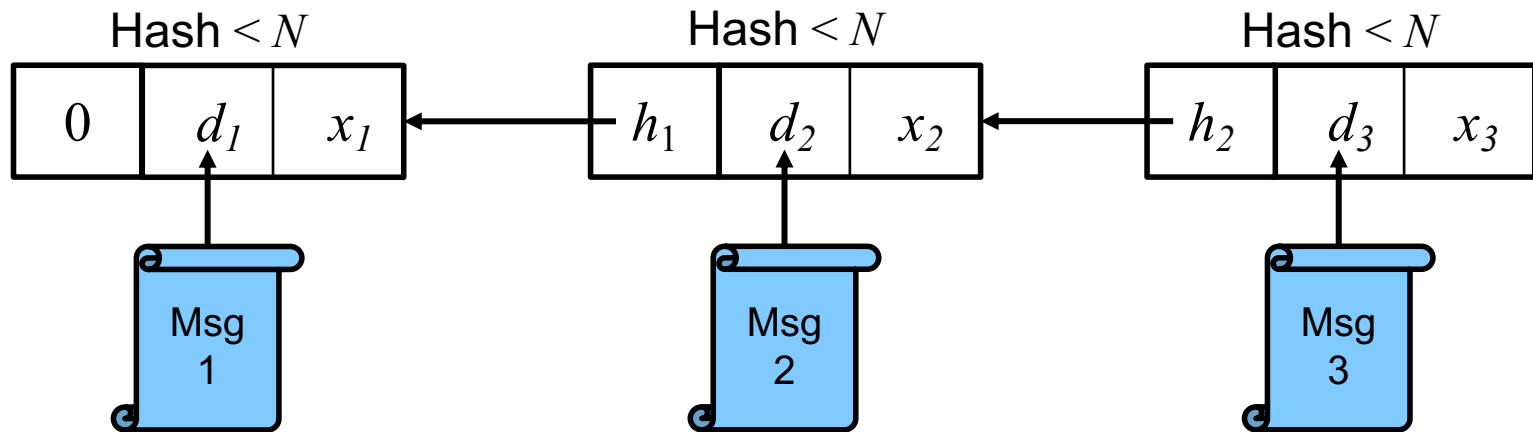
$d = \text{Digest}(m)$

Find  $x$  such that  $\text{Hash}(d \parallel x) < N$



# Proof of Work: Building a log

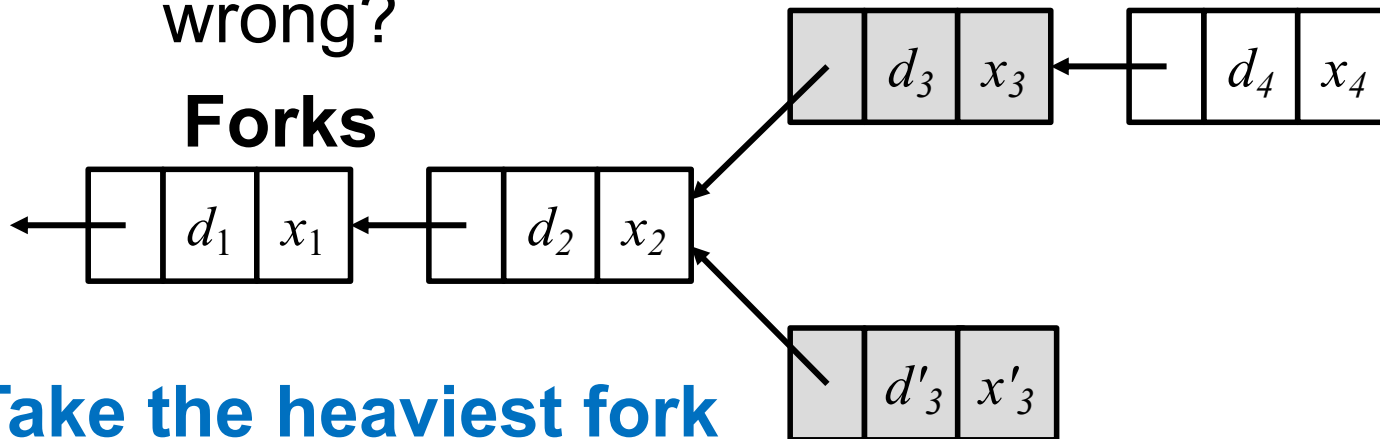
1. To add a message, generate a proof of work with that message
2. Connect each message to previous



# Proof of Work: Coming to consensus

What can go wrong?

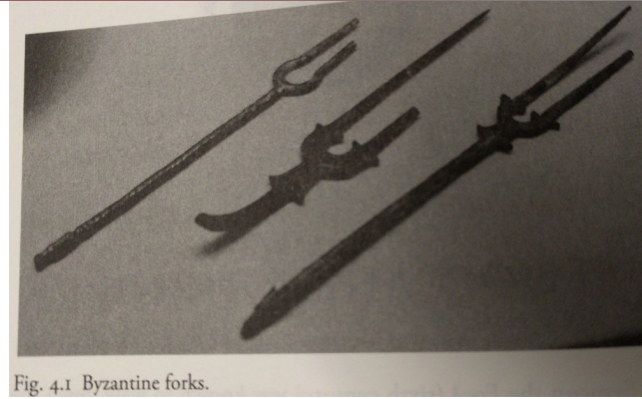
**Forks**



**Take the heaviest fork  
(one with the most work)**

# Nakamoto Consensus

- If majority of computation is honest, honest parties will agree (eventually)
- Log is tamper-proof
  - It would require redoing all of the work to tamper



# Blockchains for Audit

- **Individual accountability**
  - Everything is visible. Everyone is accountable.
- **Event reconstruction**
  - All of the events are there. Easy to reconstruct.
- **Real-time intelligence**
  - Miners can verify everything ~~as~~ it goes on the log.  
**before!**

# Not just a log!

## **Authoritative record**

- Instead of logging events elsewhere, the blockchain can record the definition of events (e.g. transactions)
- Online validation can prevent illegal events from ever happening!

# What restrictions make sense?

## Transaction Processing System

- Each block has a limited number of transactions (1 MB)
- Transactions cannot create money
  - Except coinbase transaction to reward miner
- Coins can only be spent once (spending creates new unspent coins)
- To spend a coin conditions must be met (e.g., owner authorizes)

**Bitcoin**



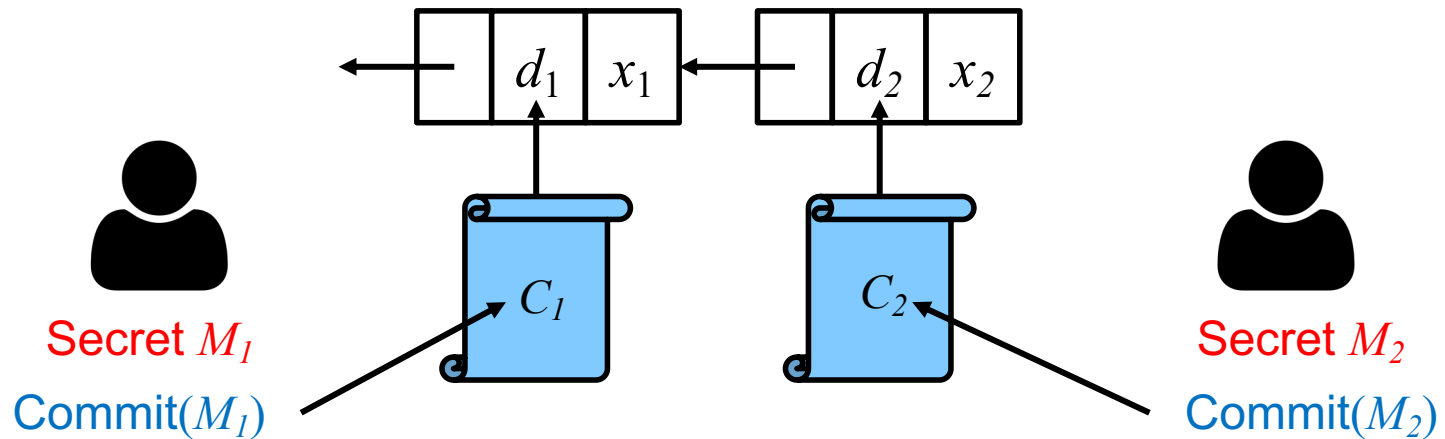
# BLOCKCHAINS AND CONFIDENTIALITY

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# What do we do with private data?

Cannot put it on the blockchain – everything is public  
Only publish commitments



# Commitment Schemes

- A commitment scheme  $\text{Com}$  is a two-phase, two-party protocol such that:
  - Secrecy: receiver does not learn anything about  $x$  from  $\text{Com}(x)$
  - Binding: sender cannot produce alternative  $x'$  such that  $\text{Com}(x) = \text{Com}(x')$
- Example Protocol
  1. B  $\rightarrow$  A:  $r$
  2. A: choose random bit  $b$ . If  $b=0$ ,  $\text{Com}(x) = \text{Hash}(x)$  else  $\text{Com}(x) = \text{Hash}(x) \oplus r$
  3. A  $\rightarrow$  B:  $\text{Com}(x)$
  - ...
  1. A  $\rightarrow$  B:  $x$

# What do we do with private data?

Cannot put it on the blockchain – everything is public

Only publish commitments

- Still tamper-proof ✓
- No longer able to see actions
  - Cannot reconstruct events ✗
  - Cannot perform online validation ✗ 😞



ethereum

# Doing better with private data

Verify data validity without leaking secrets

Ongoing research with two main tools

1. Heavy-duty cryptographic constructs
  - Complex zero-knowledge proofs
2. Trusted hardware
  - Places trust in hardware instead of crypto or a large group

# Zero-Knowledge Proof

- A zero-knowledge proof is a protocol that satisfies:
  1. **Completeness**: if the statement is true, a verifier will be convinced of this fact.
  2. **Soundness**: if the statement is false, no cheating prover can convince an honest verifier that it is true (except with some small probability).
  3. **Zero-knowledge**: if the statement is true, no verifier learns anything other than the fact that the statement is true.

# Cryptographic Example

ZKP gives strong publicly verifiable integrity guarantees

- Sender authorized transaction
- Sender had money to send
- Transaction value was not negative
- Transaction was processed correctly

Can (provably) furnish transaction details to external auditor



# Trusted Hardware

## Special machine instructions

Isolate process from the surrounding system

Can remotely attest that they're running specific code

Uses (literally) hard-wired keys in the CPU

Trustworthy code can operate on secret data and attest to correctness

Examples:

- Intel Software Guard eXtensions (SGX)
- ARM TrustZone