CS62 Lab3: Timing ArrayLists

Basic Data Structures

size	increase by 1	increase by 10	double
1000	0.000002	0.000000	0.000000
2000	0.000002	0.000001	0.000000
4000	0.000004	0.000000	0.000000
8000	0.000007	0.000001	0.000000
16000	0.000013	0.000001	0.000000
32000	0.000024	0.000002	0.000000
64000	0.000047	0.000005	0.000000
128000	0.000094	0.000010	0.000000

Lab 2 agenda

- Quiz
- Lab

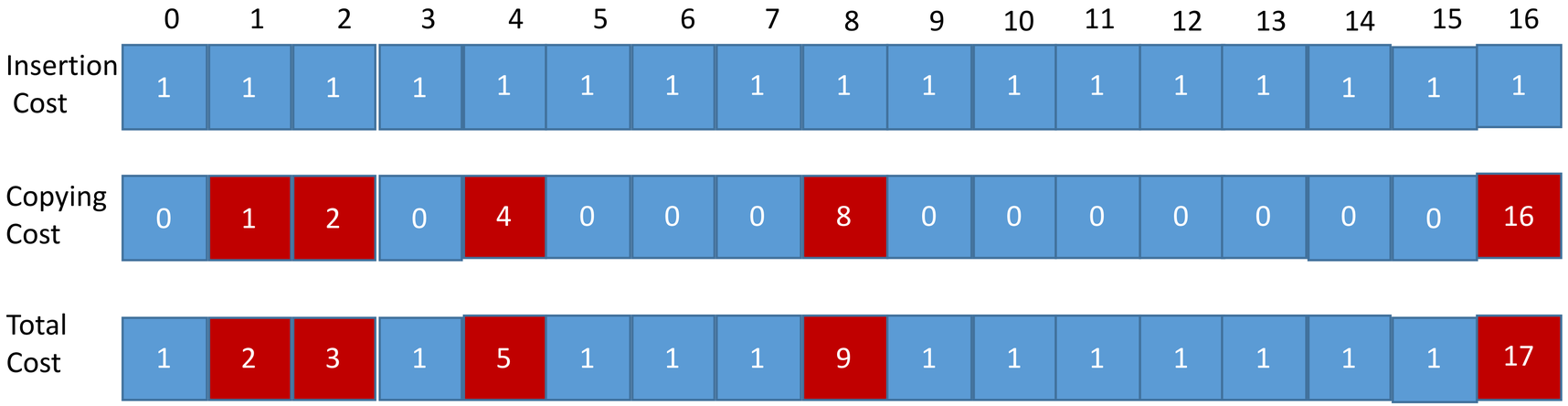
Standard Operations of `ArrayList<E>` class

- `ArrayList()`: Constructs an empty ArrayList with an initial capacity of 2 (can vary across implementations, another common initial capacity is 10).
- `ArrayList(int capacity)`: Constructs an empty ArrayList with the specified initial capacity.
- `isEmpty()`: Returns true if the ArrayList contains no elements.
- `size()`: Returns the number of elements in the ArrayList.
- `get(int index)`: Returns the element at the specified index.
- `add(E element)`: Appends the element to the end of the ArrayList.
- `add(int index, E element)`: Inserts the element at the specified index and shifts the element currently at that position (if any) and any subsequent elements to the right (adds one to their indices).
- `E remove()`: Removes and returns the element at the end of the ArrayList.
- `E remove(int index)`: Removes and returns the element at the specified index. Shifts any subsequent elements to the left (subtracts one from their indices).
- `E set(int index, E element)`: Replaces the element at the specified index with the specified element and returns the olde element.
- `clear()`: Removes all elements.

Lab structure

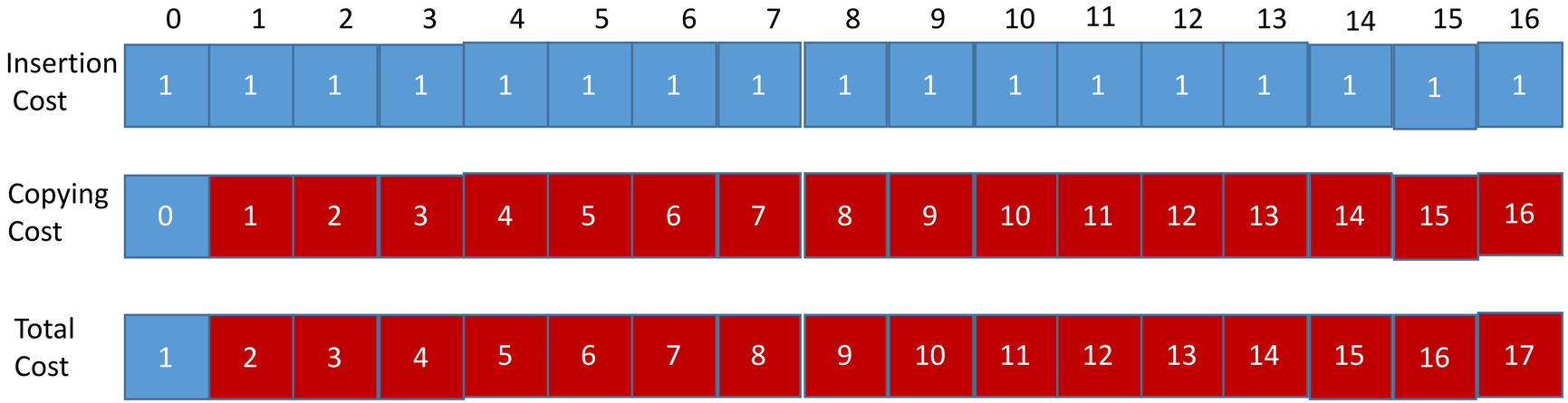
- We’ll answer the question, why do we double the size of the ArrayList? Why not just increase it by like 10 each time? Why is doubling better?
- You’ll need to create your own .java file for this one
- What we’re grading on Gradescope is answers.txt

Doubling analysis (we saw this in lecture)



- As the ArrayList increases, doubling happens *half as often* but *costs twice as much.*
- O(total cost)= $\sum$(“cost of insertions”) + $\sum$(“cost of copying”)
- $\sum$(“cost of insertions”) $= n$.
- $\sum$(“cost of copying”) = $1 + 2 + 2^2 + \ldots + 2^{\log_2 n - 1} \leq 2n$.
- O(total cost) $\leq 3n$, therefore amortized cost is $\leq \frac{3n}{n} = 3 = O^+(1)$, but “lumpy”.

Amortized analysis for n add() operations when increasing ArrayList by 1.



- $\sum$(“cost of insertions”) $= n$.
- $\sum$(“cost of copying”) $= 0 + 1 + 2 + 3 + \ldots + n - 1 = n(n-1)/2$.
- O(total cost) $= n + n(n-1)/2 = n(n+1)/2$, therefore amortized cost is $(n+1)/2$ or $O^+(n)$.
- Same idea when increasing ArrayList size by a constant (like 10).
  - This is why increasing the capacity by 1 is the slowest, and 10 the second slowest, and doubling the fastest in this lab.