CS062
DATA STRUCTURES AND ADVANCED PROGRAMMING

20: Left-Leaning Red-Black Trees

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Lecture 20: Left-leaning Red-Black Trees

- Introduction
- Elementary red-black BST operations
- Insertion
- Mathematical analysis
- Historical context

Some slides adopted from Algorithms 4th Edition or COS226
Left-leaning red-black BSTs correspond 1-1 with 2-3 trees

- Start with standard BSTs which are made up of 2-nodes.
- Add extra information to encode 3-nodes. We will introduce two types of links.
- **Red links**: bind together two 2-nodes to represent a 3-node.
  - Specifically, 3-nodes are represented as two 2-nodes connected by a single red link that leans left (one of the 2-nodes is the left child of the other).
- **Black links**: bind together the 2-3 tree.
- Advantage: Can use BST code with minimal modification.
Left-leaning red-black BSTs correspond 1-1 with 2-3 trees
Definition

- A left-leaning red-black tree is a BST such that:
  - No node has two red links connected to it.
  - Red links lean left.
  - Every path from root to leaves has the same number of black links (perfect black balance).
INTRODUCTION

Search

- Exactly the same as for elementary BSTs (we ignore the color).
- But runs faster because of better balance.

```java
public Value get(Key key) {
    if (key == null) throw new IllegalArgumentException("argument to get() is null");
    return get(root, key);
}
```

```java
private Value get(Node x, Key key) {
    while (x != null) {
        int cmp = key.compareTo(x.key);
        if (cmp < 0) x = x.left;
        else if (cmp > 0) x = x.right;
        else return x.val;
    }
    return null;
}
```

- Operations such as floor, iteration, rank, selection are also identical.
Representation

- Each node is pointed to by one node, its parent. We can use this to encode the color of the links in nodes.

- True if the link from the parent is red and false if it is black. Null links are black.

```java
private static final boolean RED = true;
private static final boolean BLACK = false;

private Node root; // root of the BST

// BST helper node data type
private class Node {
    private Key key;    // key
    private Value val;  // associated data
    private Node left, right; // links to left and right subtrees
    private boolean color; // color of parent link
    private int size;     // subtree count

    private boolean isRed(Node x) {
        if (x == null) return false;
        return x.color == RED;
    }
```
Practice Time

Which of the following are legal LLRB trees?
Answer

- Which of the following are legal LLRB trees?
- iii and iv
  - i is not balanced and ii is also not in symmetrical order
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**Left rotation**: Orient a (temporarily) right-leaning red link to lean left
Right rotation: Orient a left-leaning red link to a (temporarily) lean right
Color flip: Recolor to split a (temporary) 4-node
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Basic strategy: Maintain 1-1 correspondence with 2-3 trees

- During internal operations, maintain:
  - symmetric order
  - perfect black balance.

- But we might violate color invariants. For example:
  - Right-leaning red link.
  - Two red children (temporary 4-node).
  - Left-left red (temporary 4-node).
  - Left-right red (temporary 4-node).

- To restore color invariant we will be performing rotations and color flips.
Insertion into a LLRB

- Do standard BST insertion and color the new link red.
- Repeat until color invariants restored:
  - Both children red? Flip colors.
  - Right link red? Rotate left.
  - Two left reds in a row? Rotate right.
Red-black BST construction demo

red-black BST
Implementation

- Only three cases:
  - Right child red; left child black: rotate left.
  - Left child red; left-left grandchild red: rotate right.
  - Both children red: flip colors.

```java
// insert the key-value pair in the subtree rooted at h
private Node put(Node h, Key key, Value val) {
    if (h == null) return new Node(key, val, RED, 1);

    int cmp = key.compareTo(h.key);
    if (cmp < 0) h.left = put(h.left, key, val);
    else if (cmp > 0) h.right = put(h.right, key, val);
    else h.val = val;

    // fix-up any right-leaning links
    if (isRed(h.right) && !isRed(h.left)) h = rotateLeft(h);
    if (isRed(h.left) && isRed(h.left.left)) h = rotateRight(h);
    if (isRed(h.left) && isRed(h.right)) flipColors(h);
    h.size = size(h.left) + size(h.right) + 1;

    return h;
}
```
Visualization of insertion into a LLRB tree

- 255 insertions in ascending order.
Visualization of insertion into a LLRB tree

- 255 insertions in descending order.
Visualization of insertion into a LLRB tree

- 255 insertions in random order.
Practice Time - Worksheet #20

- Draw the LLRB tree that results when you insert the keys 10, 18, 7, 15, 16, 30, 25, 40, 60 in that order in an initially empty tree.
ANSWER - Worksheet #20

- Draw the LLRB tree that results when you insert the keys 10, 18, 7, 15, 16, 30, 25, 40, 60 in that order in an initially empty tree.
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Balance in LLRB trees

- Height of LLRB trees is \( \leq 2 \log n \) in the worst case.

- Worst case is a 2-3 tree that is all 2-nodes except that the left-most path is made up of 3-nodes.

- All ordered operations (min, max, floor, ceiling) etc. are also \( O(\log n) \).
### Summary for dictionary operations

<table>
<thead>
<tr>
<th></th>
<th>Worst case</th>
<th></th>
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<th>Average case</th>
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<th></th>
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</thead>
<tbody>
<tr>
<td></td>
<td>Search</td>
<td>Insert</td>
<td>Delete</td>
<td>Search</td>
<td>Insert</td>
<td>Delete</td>
</tr>
<tr>
<td>BST</td>
<td>$n$</td>
<td>$n$</td>
<td>$n$</td>
<td>$\log n$</td>
<td>$\log n$</td>
<td>$\sqrt{n}$</td>
</tr>
<tr>
<td>2-3 search tree</td>
<td>$\log n$</td>
<td>$\log n$</td>
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</tr>
<tr>
<td>Red-black BSTs</td>
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</tr>
</tbody>
</table>
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Red-black trees

- Why red-black? Xerox PARC had a laser printer and red and black had the best contrast...
- Left-leaning red-black trees [Sedgewick, 2008]
  - Inspired by difficulties in proper implementation of RB BSTs.
- RB BSTs have been involved in lawsuit because of improper implementation.
Balanced trees in the wild

- Red-black trees are widely used as system dictionaries.
  - e.g., Java: `java.util.TreeMap` and `java.util.TreeSet`.
- Other balanced BSTs: AVL, splay, randomized.
- 2-3 search trees are a subset of b-trees.
  - See recommended textbook for more.
- B-trees are widely used for file systems and databases.
Readings:

- Recommended Textbook: Chapter 3.3 (Pages 424-447)
- Website:
  - https://algs4.cs.princeton.edu/33balanced/
- Visualization:
  - https://algs4.cs.princeton.edu/GrowingTree/ (for LLRB trees)

Worksheet:

- Lecture 20 worksheet
Problem 1

- Draw the left-leaning red-black BST that results when you insert items with the keys E, A, S, Y, Q, U, T, I, O, N in that order into an initially empty tree.
ANSWER 1

- Draw the left-leaning red-black BST that results when you insert items with the keys E, A, S, Y, Q, U, T, I, O, N in that order into an initially empty tree.