

CS062

DATA STRUCTURES AND ADVANCED PROGRAMMING

13: Merge Sort, Quick Sort



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he/him/his

MERGESORT

	input	M	E	R	G	E	S	O	R	T	E	X	A	M	P	L	E
	sort left half	E	E	G	M	O	R	R	S	T	E	X	A	M	P	L	E
	sort right half	E	E	G	M	O	R	R	S	A	E	E	L	M	P	T	X
	merge results	A	E	E	E	E	G	L	M	M	O	P	R	R	S	T	X

Basics

[Mergesort overview](#)

- ▶ Invented by John von Neumann in 1945

- ▶ Algorithm sketch:

- ▶ Divide array into two halves.
- ▶ Recursively sort each half.
- ▶ Merge the two halves



<https://en.wikipedia.org/wiki/File:JohnvonNeumann-LosAlamos.gif>

Merging two already sorted halves into one sorted array

```
private static void merge(Comparable[] a, Comparable[] aux, int lo, int mid, int hi) {  
    for (int k = lo; k <= hi; k++)  
        aux[k] = a[k];  
  
    int i = lo, j = mid+1;  
    for (int k = lo; k <= hi; k++) {  
        if (i > mid) //ran out of elements in the left subarray  
            a[k] = aux[j++];  
        else if (j > hi) //ran out of elements in the right subarray  
            a[k] = aux[i++];  
        else if (less(aux[j], aux[i]))  
            a[k] = aux[j++];  
        else  
            a[k] = aux[i++];  
    }  
}
```

Merging Example - 1st step - copying to auxiliary array

Array aux

A G L O R H I M S T

0	1	2	3	4	5	6	7	8	9
---	---	---	---	---	---	---	---	---	---

lo

mid

hi

Array a

A G L O R H I M S T

0	1	2	3	4	5	6	7	8	9
---	---	---	---	---	---	---	---	---	---

```
private static void merge(Comparable[] a, Comparable[] aux, int lo, int mid, int hi) {  
    for (int k = lo; k <= hi; k++)  
        aux[k] = a[k]; // Copy a to aux array  
  
    int i = lo, j = mid+1;  
    for (int k = lo; k <= hi; k++) {  
        if (i > mid) //ran out of elements in the left subarray  
            a[k] = aux[j++];  
        else if (j > hi) //ran out of elements in the right subarray  
            a[k] = aux[i++];  
        else if (less(aux[j], aux[i]))  
            a[k] = aux[j++];  
        else  
            a[k] = aux[i++];  
    }  
}
```



<http://algs4.cs.princeton.edu>

2.2 MERGING DEMO

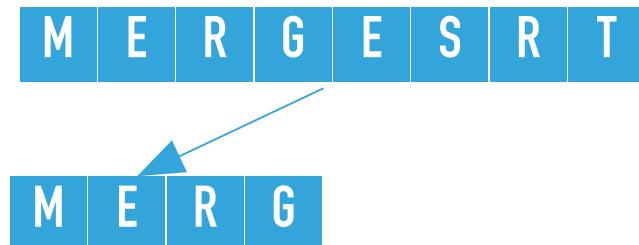
Mergesort - the quintessential example of divide-and-conquer

```
private static void sort(Comparable[] a, Comparable[] aux, int lo, int hi) {  
    if (hi <= lo)  
        return;  
    int mid = lo + (hi - lo) / 2;  
    sort(a, aux, lo, mid);  
    sort(a, aux, mid+1, hi);  
    merge(a, aux, lo, mid, hi);  
}  
  
public static void sort(Comparable[] a) {  
    Comparable[] aux = new Comparable[a.length];  
    sort(a, aux, 0, a.length - 1);  
}
```

```
private static void sort(Comparable[] a, Comparable[] aux, int lo, int hi) {  
    if (hi <= lo)  
        return;  
    int mid = lo + (hi - lo) / 2;  
    sort(a, aux, lo, mid);  
    sort(a, aux, mid+1, hi);  
    merge(a, aux, lo, mid, hi);  
}  
  
public static void sort(Comparable[] a) {  
    Comparable[] aux = new Comparable[a.length];  
    sort(a, aux, 0, a.length - 1);  
}
```

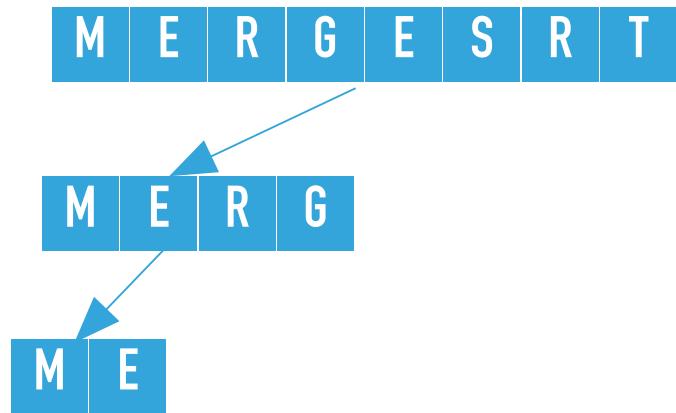
sort([M, E, R, G, E, S, R, T]) calls

sort([M, E, R, G, E, S, R, T], [null, null, null, null, null, null, null, null], 0, 7) where the array of nulls is the auxiliary array, lo = 0 and hi = 7.



```
private static void sort(Comparable[] a, Comparable[] aux, int lo, int hi) {  
    if (hi <= lo)  
        return;  
    int mid = lo + (hi - lo) / 2;  
    sort(a, aux, lo, mid);  
    sort(a, aux, mid+1, hi);  
    merge(a, aux, lo, mid, hi);  
}
```

sort([M, E, R, G, E, S, R, T], [null, null, null, null, null, null, null, null], 0, 7) calculates the `mid = 3` and calls recursively `sort` on the left subarray, that is `sort([M, E, R, G, E, S, R, T], [null, null, null, null, null, null, null, null], 0, 3)`, where `lo = 0, hi = 3`



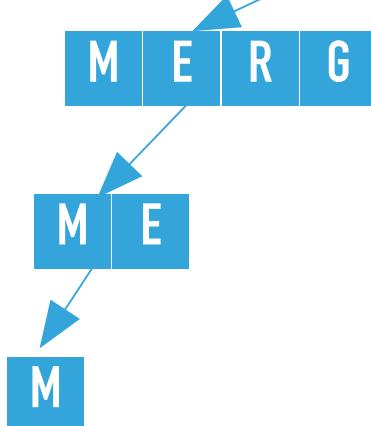
```

private static void sort(Comparable[] a, Comparable[] aux, int lo, int hi) {
    if (hi <= lo)
        return;
    int mid = lo + (hi - lo) / 2;
    sort(a, aux, lo, mid);
    sort(a, aux, mid+1, hi);
    merge(a, aux, lo, mid, hi);
}

```

sort([M, E, R, G, E, S, R, T], [null, null, null, null, null, null, null, null], 0, 3) calculates the `mid = 1` and calls recursively `sort` on the left subarray, that is `sort([M, E, R, G, E, S, R, T], [null, null, null, null, null, null, null, null], 0, 1)`, where `lo = 0, hi = 1`

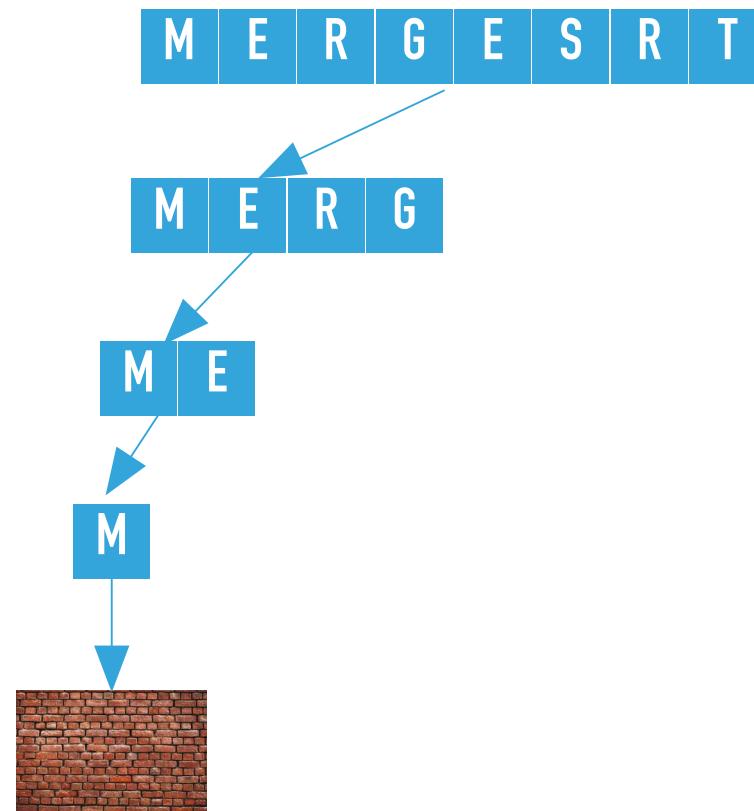
M E R G E S R T



10

```
private static void sort(Comparable[] a, Comparable[] aux, int lo, int hi) {  
    if (hi <= lo)  
        return;  
    int mid = lo + (hi - lo) / 2;  
    sort(a, aux, lo, mid);  
    sort(a, aux, mid+1, hi);  
    merge(a, aux, lo, mid, hi);  
}
```

sort([M, E, R, G, E, S, R, T], [null, null, null, null, null, null, null, null], 0, 1) calculates the $mid = 0$ and calls recursively `sort` on the left subarray, that is `sort([M, E, R, G, E, S, R, T], [null, null, null, null, null, null, null, null], 0, 0)`, where $lo = 0, hi = 0$

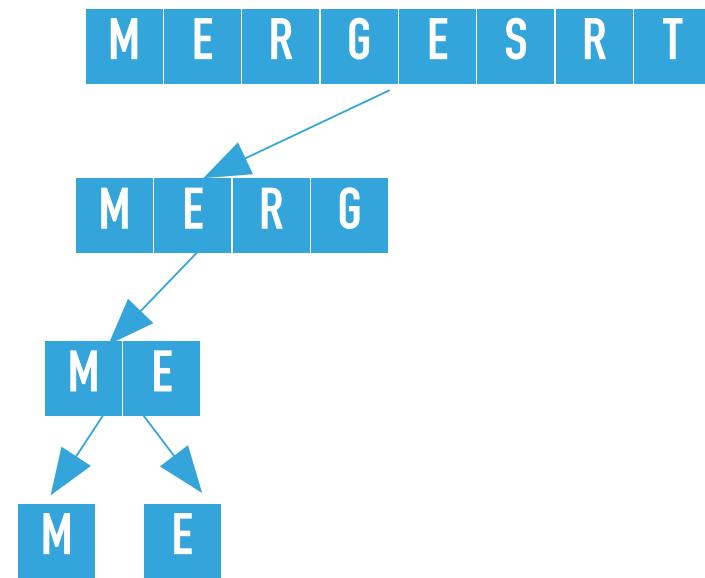


```

private static void sort(Comparable[] a, Comparable[] aux, int lo, int hi) {
    if (hi <= lo)
        return;
    int mid = lo + (hi - lo) / 2;
    sort(a, aux, lo, mid);
    sort(a, aux, mid+1, hi);
    merge(a, aux, lo, mid, hi);
}

```

sort([M, E, R, G, E, S, R, T], [null, null, null, null, null, null, null, null], 0, 0) finds $hi \leq lo$ and returns.

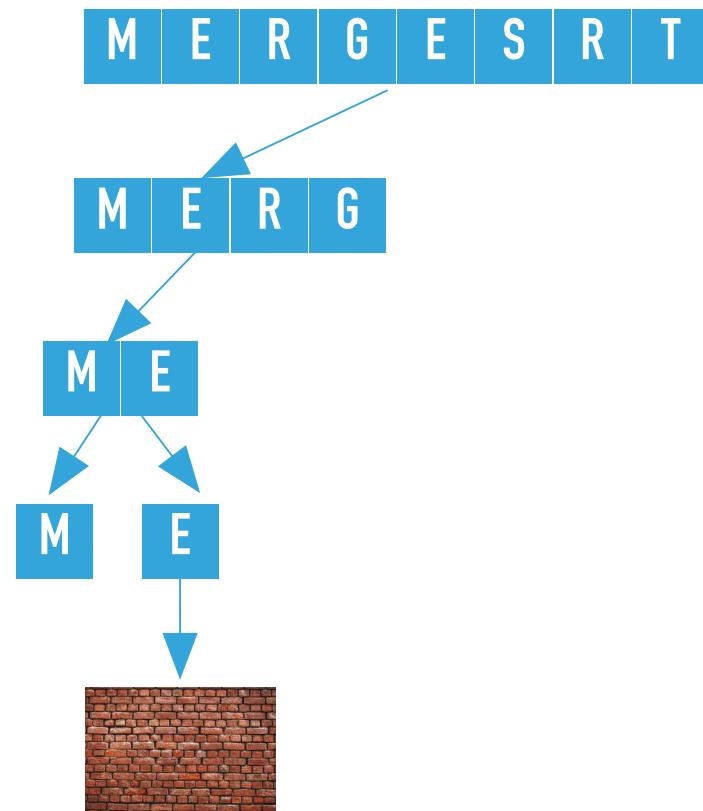


```

private static void sort(Comparable[] a, Comparable[] aux, int lo, int hi) {
    if (hi <= lo)
        return;
    int mid = lo + (hi - lo) / 2;
    sort(a, aux, lo, mid);
    sort(a, aux, mid+1, hi);
    merge(a, aux, lo, mid, hi);
}

```

sort([M, E, R, G, E, S, R, T], [null, null, null, null, null, null, null, null], 0, 1) calls recursively sort on the right subarray, that is sort([M, E, R, G, E, S, R, T], [null, null, null, null, null, null, null, null], 1, 1), where lo = 1, hi = 1

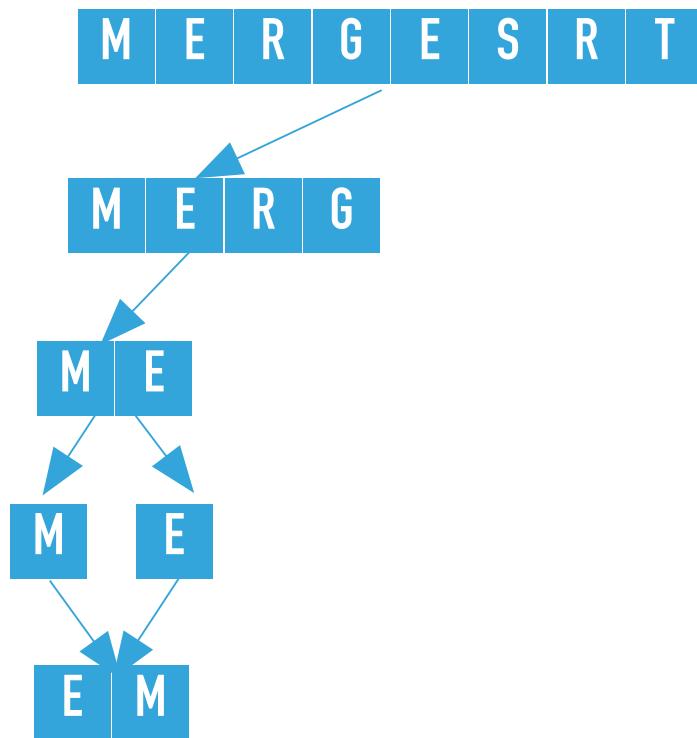


```

private static void sort(Comparable[] a, Comparable[] aux, int lo, int hi) {
    if (hi <= lo)
        return;
    int mid = lo + (hi - lo) / 2;
    sort(a, aux, lo, mid);
    sort(a, aux, mid+1, hi);
    merge(a, aux, lo, mid, hi);
}

```

sort([M, E, R, G, E, S, R, T], [null, null, null, null, null, null, null, null], 1, 1) finds $hi \leq lo$ and returns.

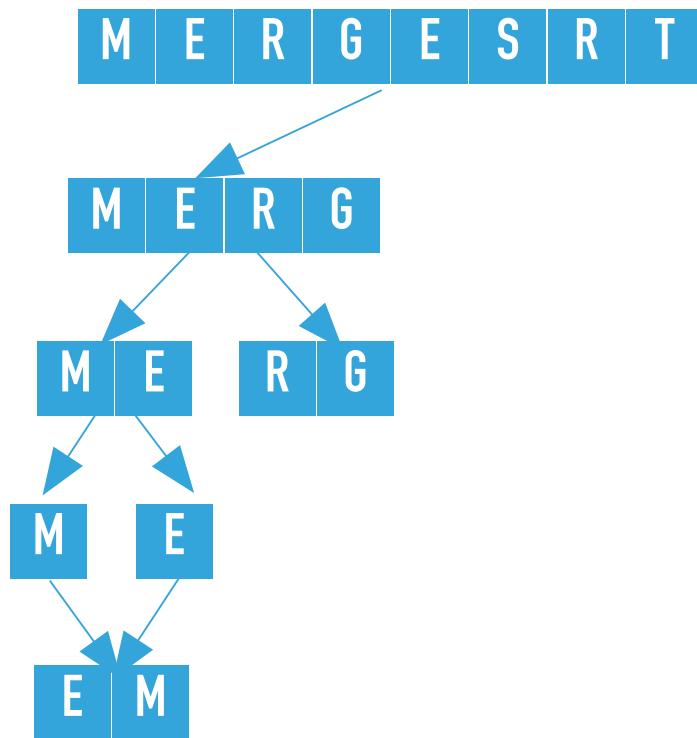


```

private static void sort(Comparable[] a, Comparable[] aux, int lo, int hi) {
    if (hi <= lo)
        return;
    int mid = lo + (hi - lo) / 2;
    sort(a, aux, lo, mid);
    sort(a, aux, mid+1, hi);
    merge(a, aux, lo, mid, hi);
}

```

sort([M, E, R, G, E, S, R, T], [null, null, null, null, null, null, null, null], 0, 1) merges the two subarrays that is calls merge([M, E, R, G, E, S, R, T], [null, null, null, null, null, null, null, null], 0, 0, 1), where lo = 0, mid = 0, and hi = 1. The resulting partially sorted array is [E, M, R, G, E, S, R T].

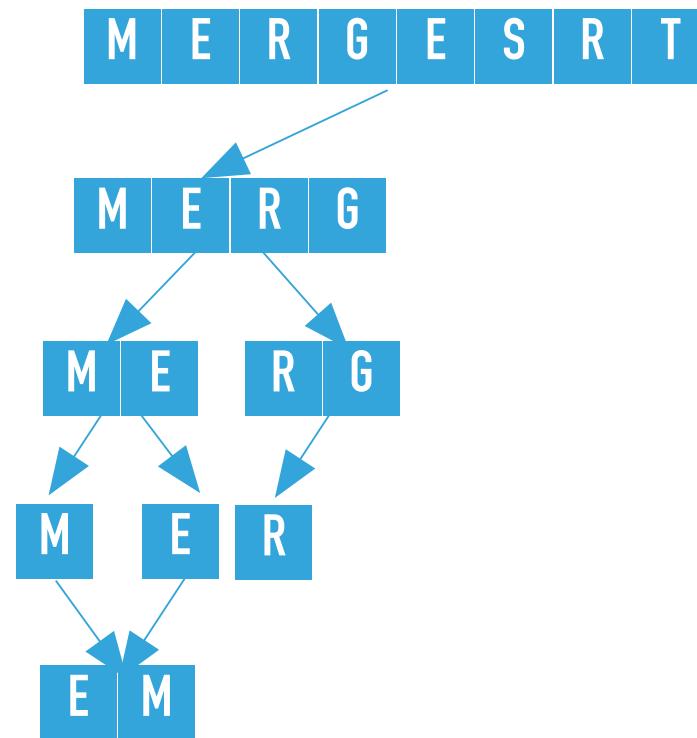


```

private static void sort(Comparable[] a, Comparable[] aux, int lo, int hi) {
    if (hi <= lo)
        return;
    int mid = lo + (hi - lo) / 2;
    sort(a, aux, lo, mid);
    sort(a, aux, mid+1, hi);
    merge(a, aux, lo, mid, hi);
}

```

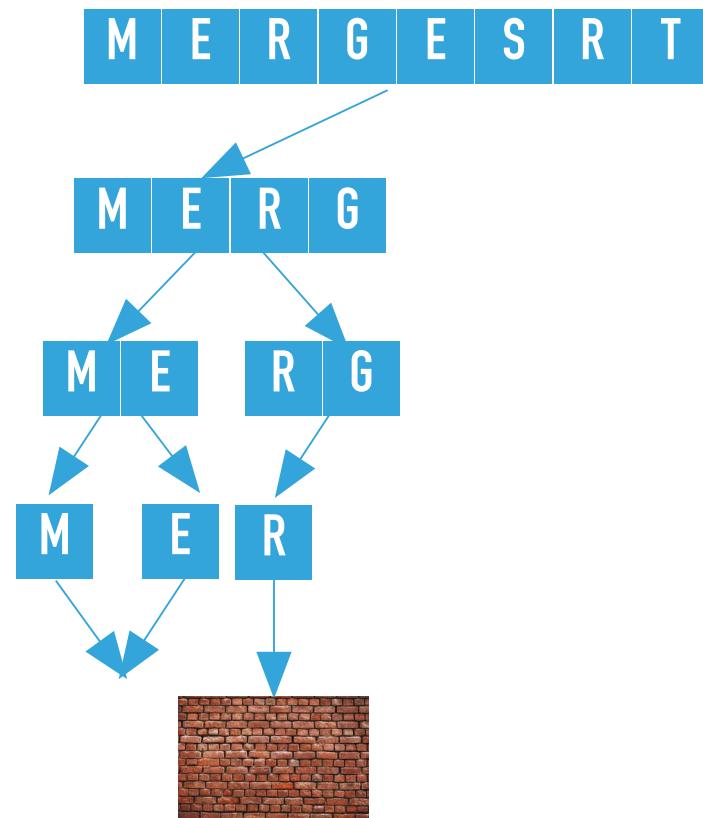
sort([E, M, R, G, E, S, R, T], [M, E, null, null, null, null, null, null], 0, 3) calls recursively `sort` on the right subarray, that is `sort([E, M, R, G, E, S, R, T], [M, E, null, null, null, null, null], 2, 3)`, where `lo = 2, hi = 3`



```

private static void sort(Comparable[] a, Comparable[] aux, int lo, int hi) {
    if (hi <= lo)
        return;
    int mid = lo + (hi - lo) / 2;
    sort(a, aux, lo, mid);
    sort(a, aux, mid+1, hi);
    merge(a, aux, lo, mid, hi);
}
  
```

sort([E, M, R, G, E, S, R, T], [M, E, null, null, null, null, null, null], 2, 3)
 calculates the `mid = 2` and calls recursively `sort` on the left subarray, that is `sort([E, M, R, G, E, S, R, T], [M, E, null, null, null, null, null, null], 2, 2)`, where `lo = 2, hi = 2`

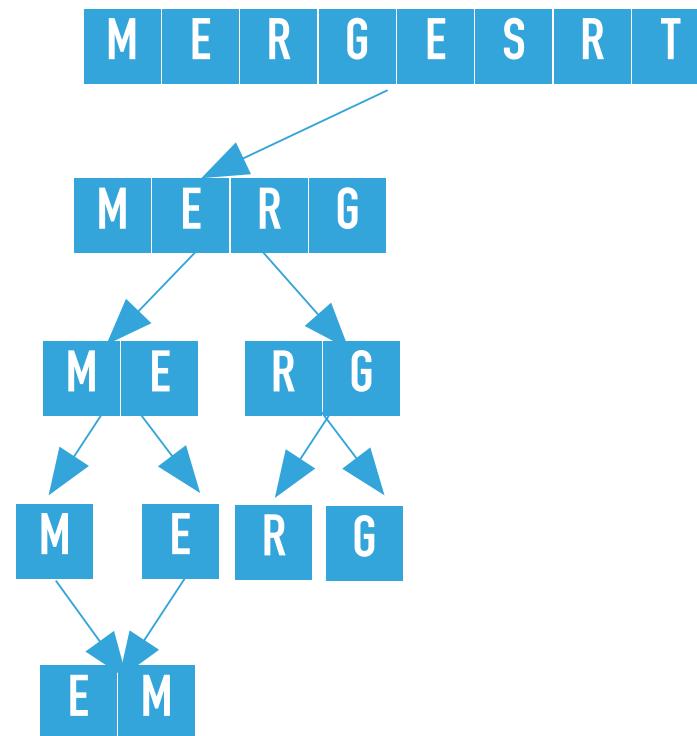


```

private static void sort(Comparable[] a, Comparable[] aux, int lo, int hi) {
    if (hi <= lo)
        return;
    int mid = lo + (hi - lo) / 2;
    sort(a, aux, lo, mid);
    sort(a, aux, mid+1, hi);
    merge(a, aux, lo, mid, hi);
}

```

sort([E, M, R, G, E, S, R, T], [M, E, null, null, null, null, null, null], 2, 2) finds $hi \leq lo$ and returns.

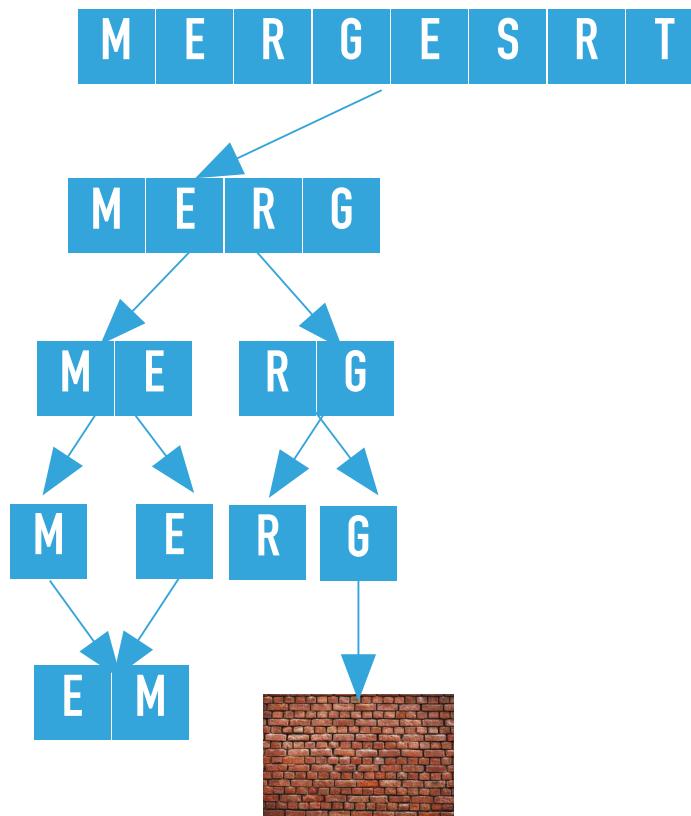


```

private static void sort(Comparable[] a, Comparable[] aux, int lo, int hi) {
    if (hi <= lo)
        return;
    int mid = lo + (hi - lo) / 2;
    sort(a, aux, lo, mid);
    sort(a, aux, mid+1, hi);
    merge(a, aux, lo, mid, hi);
}

```

sort([E, M, R, G, E, S, R, T], [M, E, null, null, null, null, null, null], 2, 3) calls recursively `sort` on the right subarray, that is `sort([E, M, R, G, E, S, R, T], [M, E, null, null, null, null, null], 3, 3)`, where `lo = 3, hi = 3`

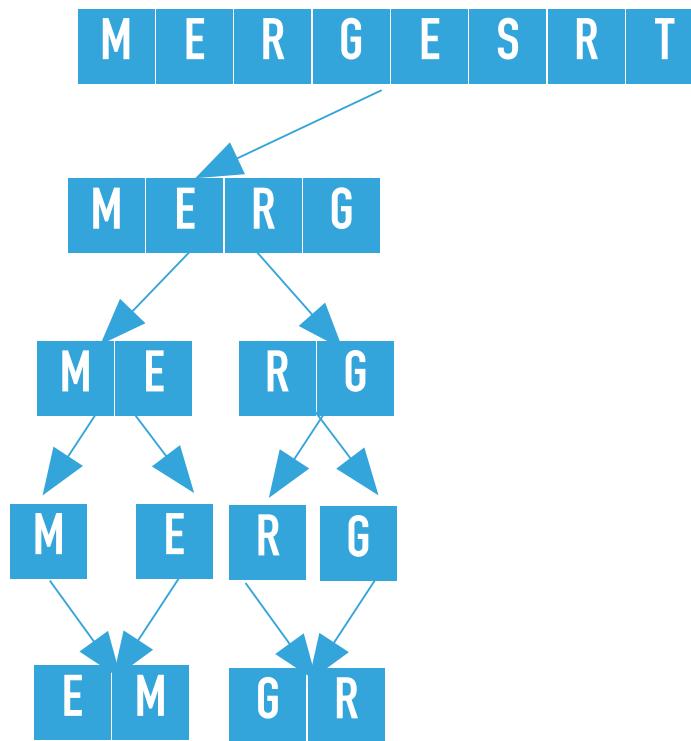


```

private static void sort(Comparable[] a, Comparable[] aux, int lo, int hi) {
    if (hi <= lo)
        return;
    int mid = lo + (hi - lo) / 2;
    sort(a, aux, lo, mid);
    sort(a, aux, mid+1, hi);
    merge(a, aux, lo, mid, hi);
}

```

sort([E, M, R, G, E, S, R, T], [M, E, null, null, null, null, null, null], 3, 3) finds $hi \leq lo$ and returns.

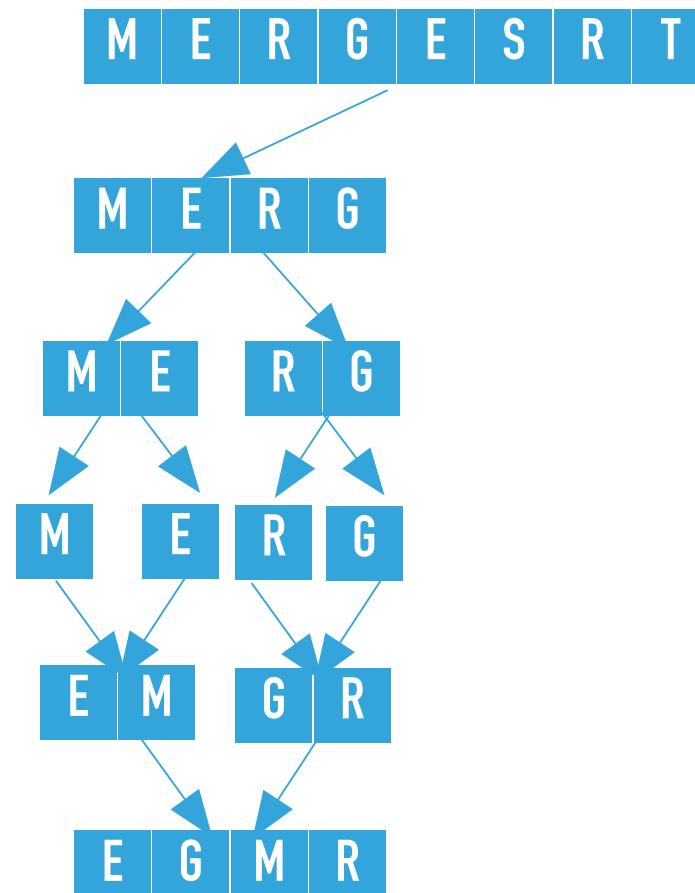


```

private static void sort(Comparable[] a, Comparable[] aux, int lo, int hi) {
    if (hi <= lo)
        return;
    int mid = lo + (hi - lo) / 2;
    sort(a, aux, lo, mid);
    sort(a, aux, mid+1, hi);
    merge(a, aux, lo, mid, hi);
}

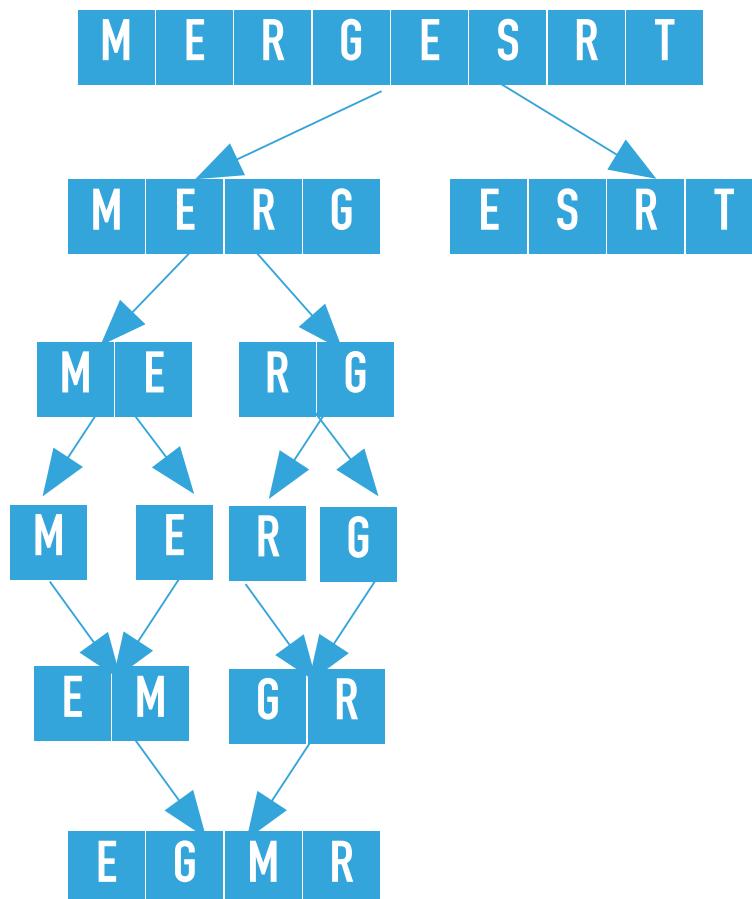
```

sort([E, M, R, G, E, S, R, T], [M, E, null, null, null, null, null, null], 2, 3)
merges the two subarrays that is calls merge([E, M, R, G, E, S, R, T], [M, E, null, null, null, null, null], 2, 2, 3), where lo = 2, mid = 2, and hi = 3. The resulting partially sorted array is [E, M, G, R, E, S, R T].



```
private static void sort(Comparable[] a, Comparable[] aux, int lo, int hi) {
    if (hi <= lo)
        return;
    int mid = lo + (hi - lo) / 2;
    sort(a, aux, lo, mid);
    sort(a, aux, mid+1, hi);
    merge(a, aux, lo, mid, hi);
}
sort([E, M, G, R, E, S, R, T], [M, E, R, G, null, null, null, null], 0, 3)
```

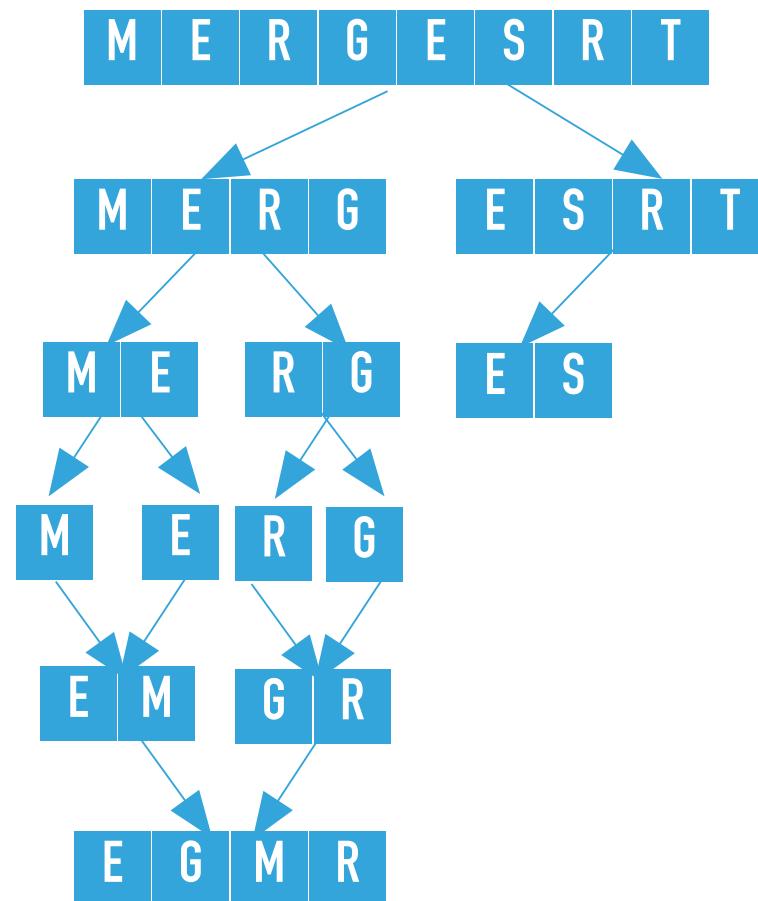
merges the two subarrays that is calls merge([E, M, G, R, E, S, R, T], [M, E, R, G, null, null, null, null], 0, 1, 3), where `lo = 0`, `mid = 1`, and `hi = 3`. The resulting partially sorted array is [E, G, M, R, E, S, R T].



```

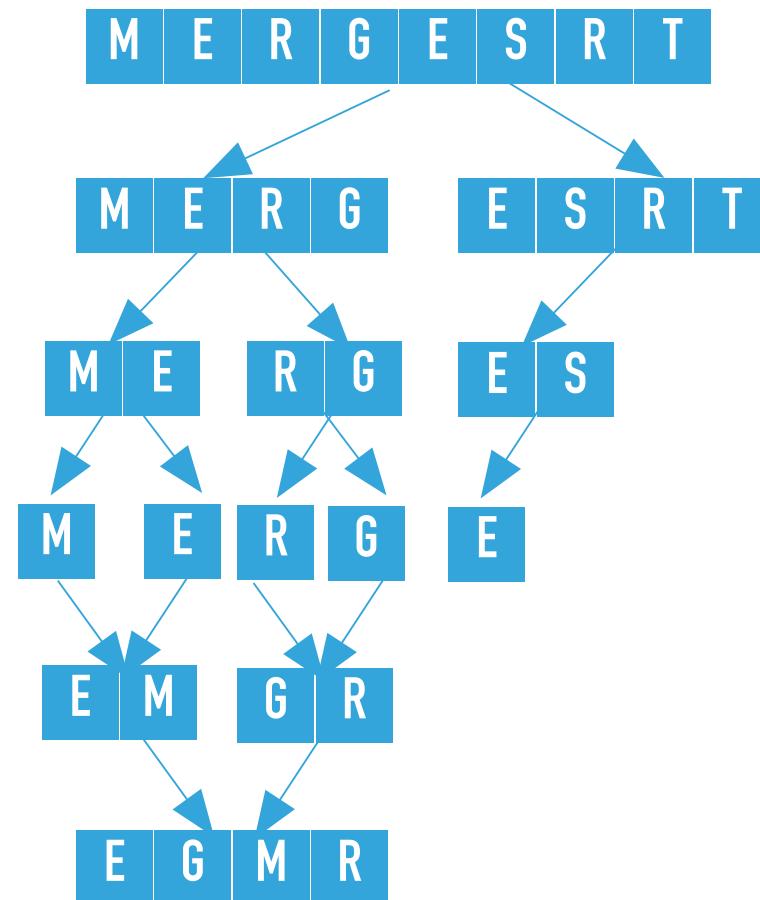
private static void sort(Comparable[] a, Comparable[] aux, int lo, int hi) {
    if (hi <= lo)
        return;
    int mid = lo + (hi - lo) / 2;
    sort(a, aux, lo, mid);
    sort(a, aux, mid+1, hi);
    merge(a, aux, lo, mid, hi);
}
  
```

sort([E, G, M, R, E, S, R, T], [E, M, G, R, null, null, null, null], 0, 7) calls recursively `sort` on the right subarray, that is `sort([E, G, M, R, E, S, R, T], [E, M, G, R, null, null, null, null], 4, 7)`, where `lo = 4, hi = 7`



```
private static void sort(Comparable[] a, Comparable[] aux, int lo, int hi) {
    if (hi <= lo)
        return;
    int mid = lo + (hi - lo) / 2;
    sort(a, aux, lo, mid);
    sort(a, aux, mid+1, hi);
    merge(a, aux, lo, mid, hi);
}
```

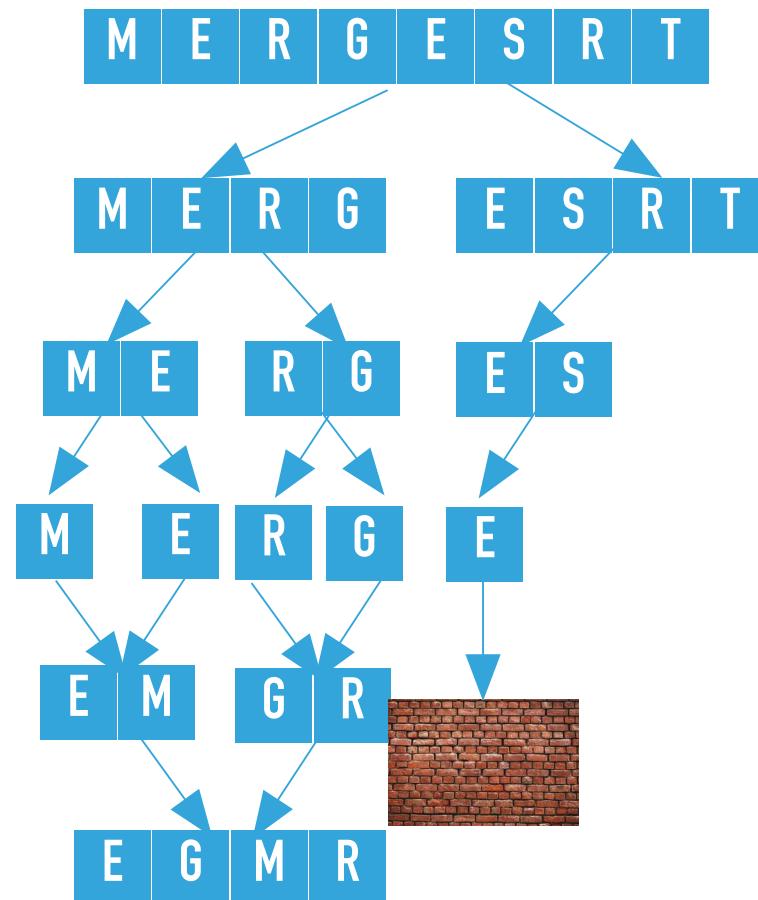
sort([E, G, M, R, E, S, R, T], [E, M, G, R, null, null, null, null], 4, 7) calculates the $mid = 5$ and calls recursively `sort` on the left subarray, that is `sort([E, G, M, R, E, S, R, T], [E, M, G, R, null, null, null, null], 4, 5)`, where $lo = 4, hi = 5$.



```

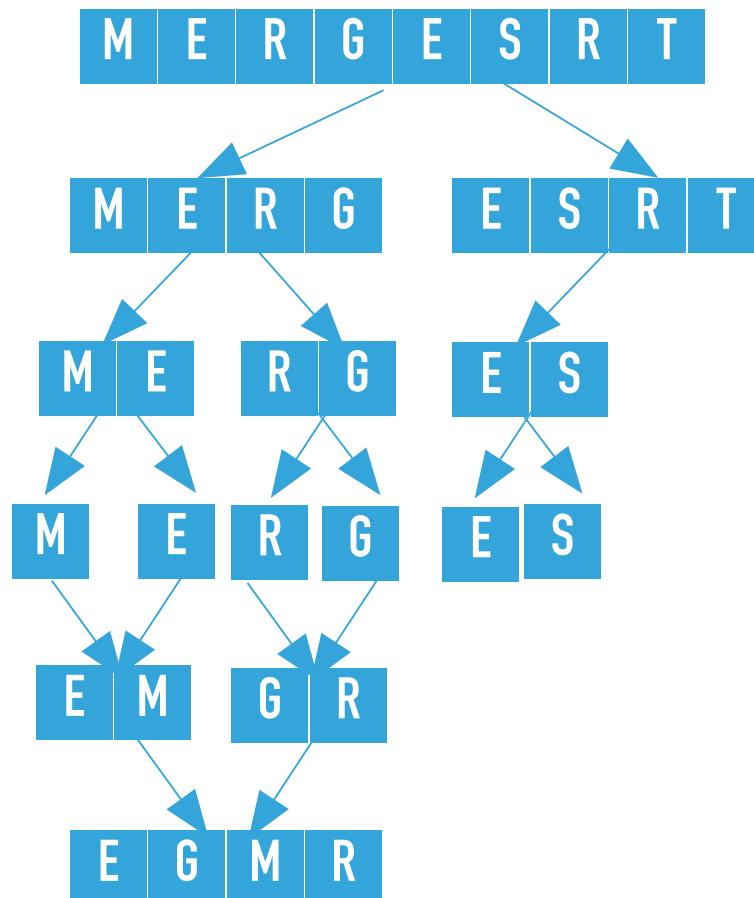
private static void sort(Comparable[] a, Comparable[] aux, int lo, int hi) {
    if (hi <= lo)
        return;
    int mid = lo + (hi - lo) / 2;
    sort(a, aux, lo, mid);
    sort(a, aux, mid+1, hi);
    merge(a, aux, lo, mid, hi);
}
  
```

sort([E, G, M, R, E, S, R, T], [E, M, G, R, null, null, null, null], 4, 5) calculates the $mid = 4$ and calls recursively `sort` on the left subarray, that is `sort([E, G, M, R, E, S, R, T], [E, M, G, R, null, null, null, null], 4, 4)`, where $lo = 4, hi = 4$.



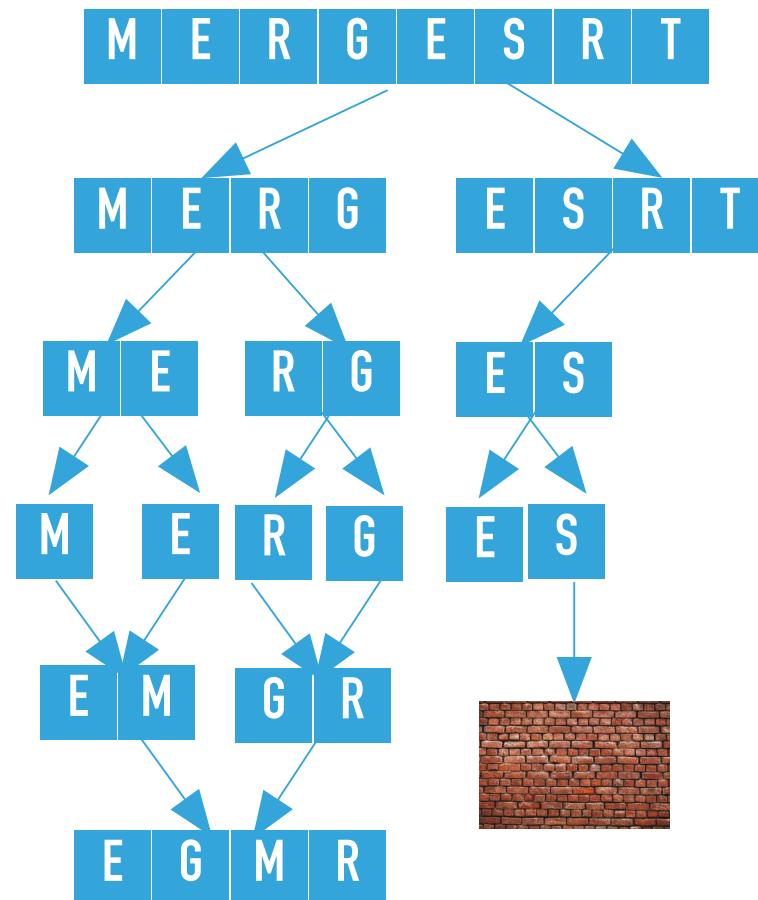
```
private static void sort(Comparable[] a, Comparable[] aux, int lo, int hi) {
    if (hi <= lo)
        return;
    int mid = lo + (hi - lo) / 2;
    sort(a, aux, lo, mid);
    sort(a, aux, mid+1, hi);
    merge(a, aux, lo, mid, hi);
}
```

sort([E, G, M, R, E, S, R, T], [E, M, G, R, null, null, null, null], 4, 4) finds $hi \leq lo$ and returns.



```
private static void sort(Comparable[] a, Comparable[] aux, int lo, int hi) {
    if (hi <= lo)
        return;
    int mid = lo + (hi - lo) / 2;
    sort(a, aux, lo, mid);
    sort(a, aux, mid+1, hi);
    merge(a, aux, lo, mid, hi);
}
```

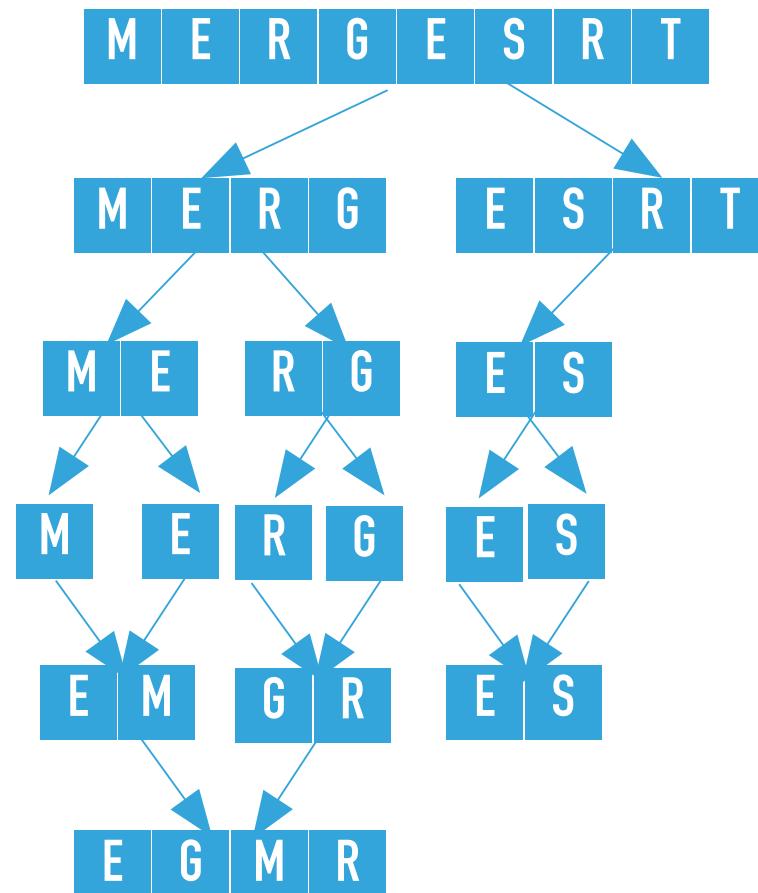
sort([E, G, M, R, E, S, R, T], [E, M, G, R, null, null, null, null], 4, 5) calls recursively `sort` on the right subarray, that is `sort([E, G, M, R, E, S, R, T], [E, M, G, R, null, null, null, null], 5, 5)`, where `lo = 5, hi = 5`



```

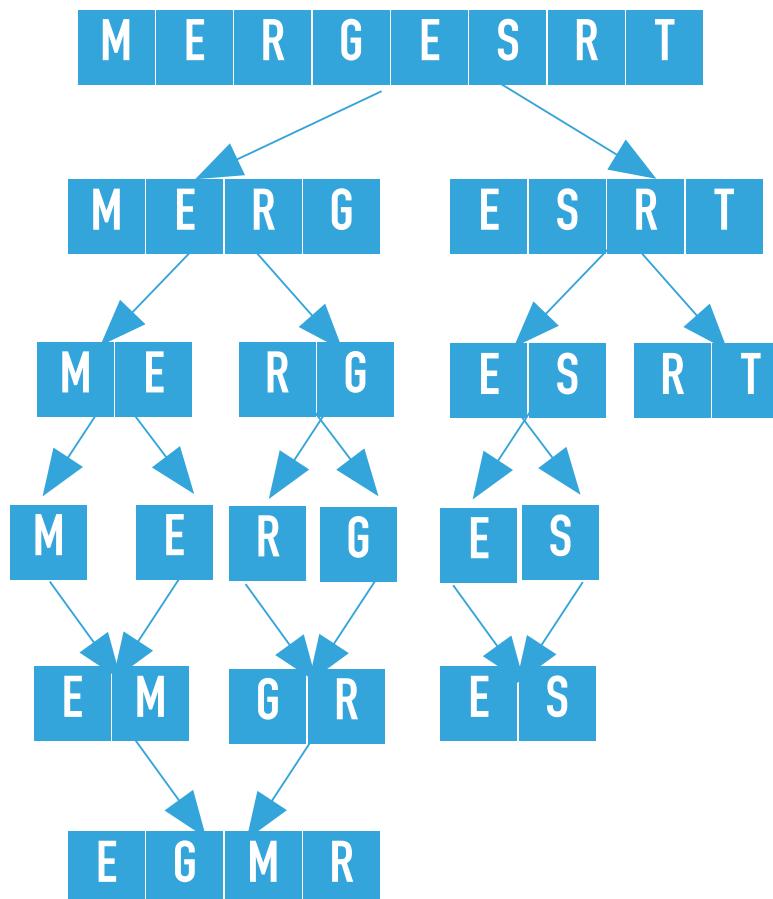
private static void sort(Comparable[] a, Comparable[] aux, int lo, int hi) {
    if (hi <= lo)
        return;
    int mid = lo + (hi - lo) / 2;
    sort(a, aux, lo, mid);
    sort(a, aux, mid+1, hi);
    merge(a, aux, lo, mid, hi);
}
  
```

sort([E, G, M, R, E, S, R, T], [E, M, G, R, null, null, null, null], 5, 5) finds
hi <= lo and returns.



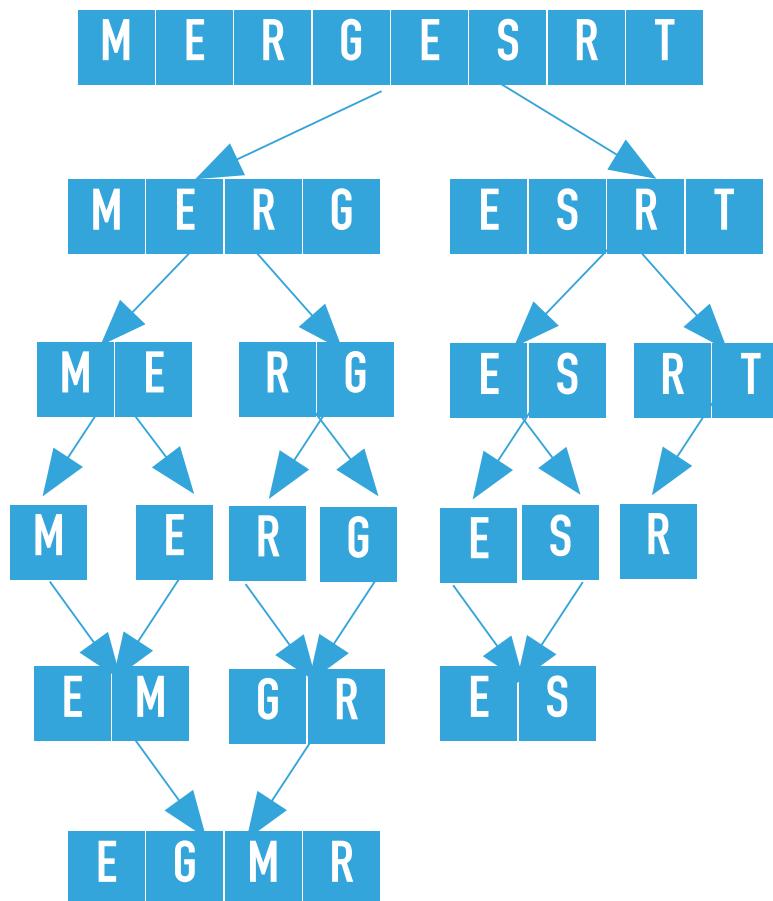
```
private static void sort(Comparable[] a, Comparable[] aux, int lo, int hi) {
    if (hi <= lo)
        return;
    int mid = lo + (hi - lo) / 2;
    sort(a, aux, lo, mid);
    sort(a, aux, mid+1, hi);
    merge(a, aux, lo, mid, hi);
}
```

sort([E, G, M, R, E, S, R, T], [E, M, G, R, null, null, null, null], 4, 5) merges the two subarrays that is calls merge([E, G, M, R, E, S, R, T], [E, M, G, R, null, null, null, null], 4, 4, 5), where lo = 4, mid = 4, and hi = 5. The resulting partially sorted array is [E, G, M, R, E, S, R, T].



```
private static void sort(Comparable[] a, Comparable[] aux, int lo, int hi) {
    if (hi <= lo)
        return;
    int mid = lo + (hi - lo) / 2;
    sort(a, aux, lo, mid);
    sort(a, aux, mid+1, hi);
    merge(a, aux, lo, mid, hi);
}
```

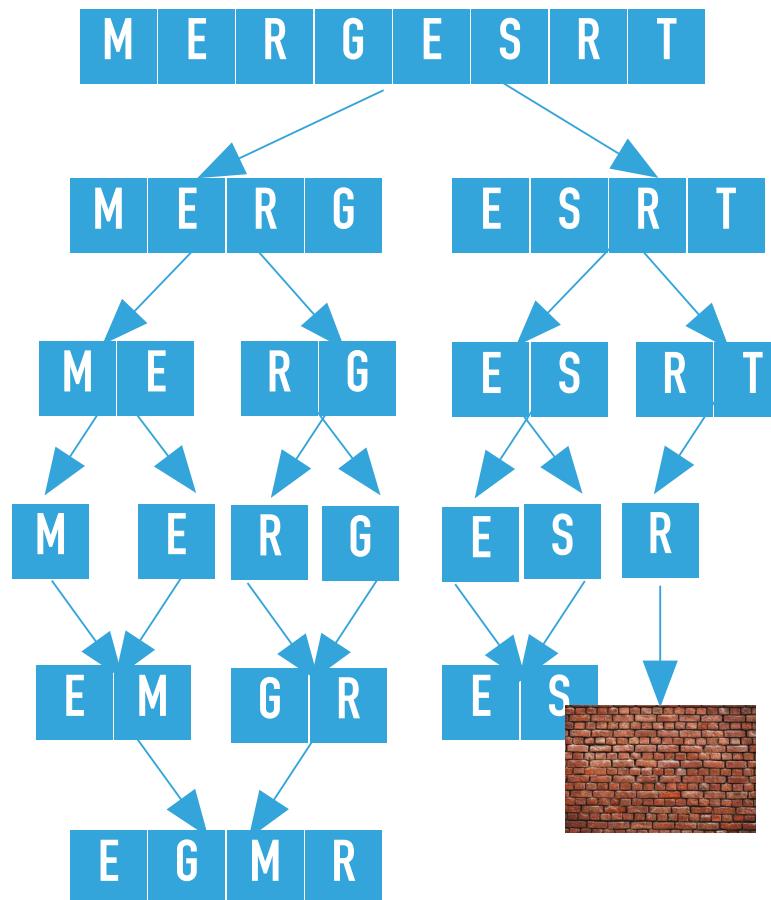
sort([E, G, M, R, E, S, R, T], [E, M, G, R, E, S, null, null], 4, 7) calls recursively sort on the right subarray, that is sort([E, G, M, R, E, S, R, T], [E, M, G, R, E, S, null, null], 6, 7), where lo = 6, hi = 7



```

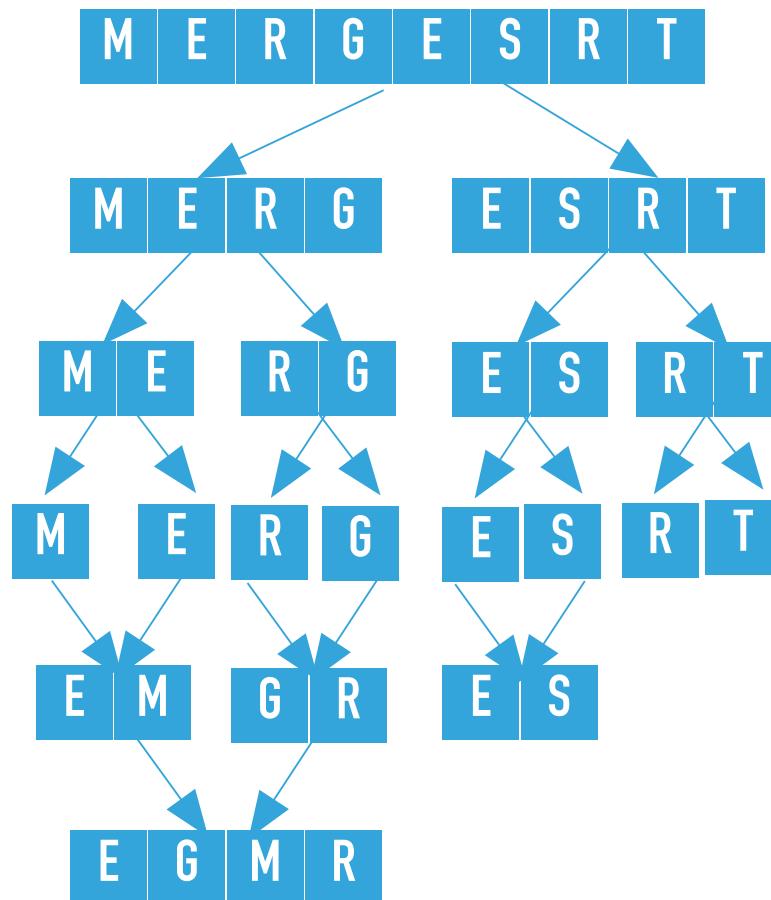
private static void sort(Comparable[] a, Comparable[] aux, int lo, int hi) {
    if (hi <= lo)
        return;
    int mid = lo + (hi - lo) / 2;
    sort(a, aux, lo, mid);
    sort(a, aux, mid+1, hi);
    merge(a, aux, lo, mid, hi);
}
  
```

sort([E, G, M, R, E, S, R, T], [E, M, G, R, E, S, null, null], 6, 7) calculates the `mid` = 6 and calls recursively `sort` on the left subarray, that is `sort([E, G, M, R, E, S, R, T], [E, M, G, R, E, S, null, null], 6, 6)`, where `lo = 6, hi = 6`.



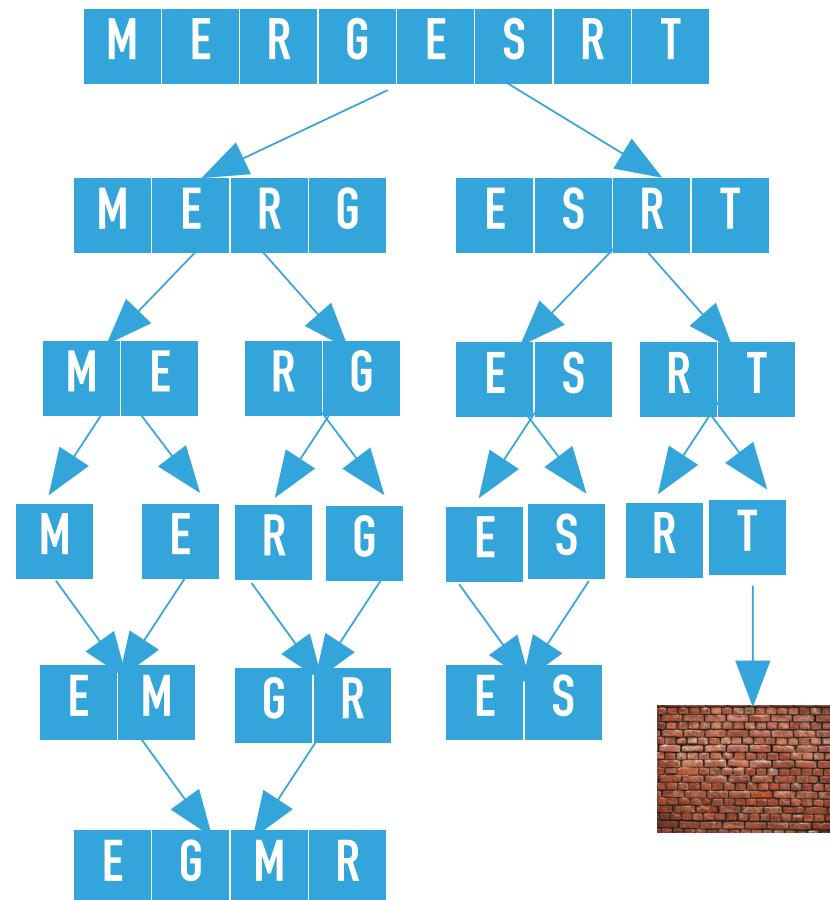
```
private static void sort(Comparable[] a, Comparable[] aux, int lo, int hi) {
    if (hi <= lo)
        return;
    int mid = lo + (hi - lo) / 2;
    sort(a, aux, lo, mid);
    sort(a, aux, mid+1, hi);
    merge(a, aux, lo, mid, hi);
}
```

sort([E, G, M, R, E, S, R, T], [E, M, G, R, E, S, null, null], 6, 6) finds `hi <= lo` and returns.



```
private static void sort(Comparable[] a, Comparable[] aux, int lo, int hi) {
    if (hi <= lo)
        return;
    int mid = lo + (hi - lo) / 2;
    sort(a, aux, lo, mid);
    sort(a, aux, mid+1, hi);
    merge(a, aux, lo, mid, hi);
}
```

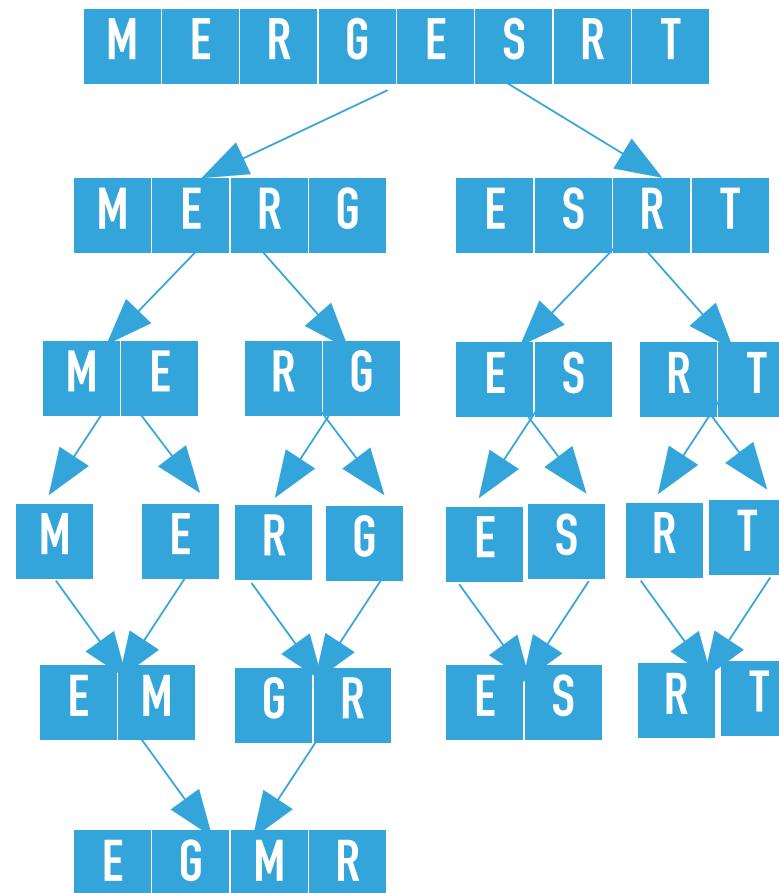
sort([E, G, M, R, E, S, R, T], [E, M, G, R, E, S, null, null], 6, 7) calls recursively sort on the right subarray, that is sort([E, G, M, R, E, S, R, T], [E, M, G, R, E, S, null, null], 7, 7), where lo = 7, hi = 7



```

private static void sort(Comparable[] a, Comparable[] aux, int lo, int hi) {
    if (hi <= lo)
        return;
    int mid = lo + (hi - lo) / 2;
    sort(a, aux, lo, mid);
    sort(a, aux, mid+1, hi);
    merge(a, aux, lo, mid, hi);
}
  
```

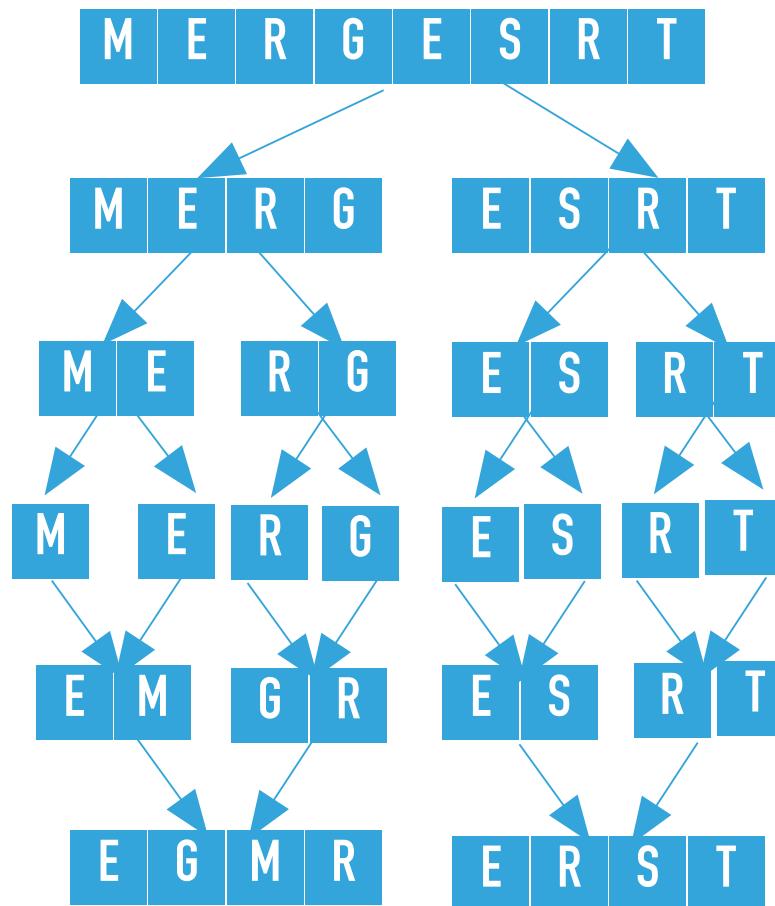
sort([E, G, M, R, E, S, R, T], [E, M, G, R, E, S, null, null], 7, 7) finds $hi \leq lo$ and returns.



```

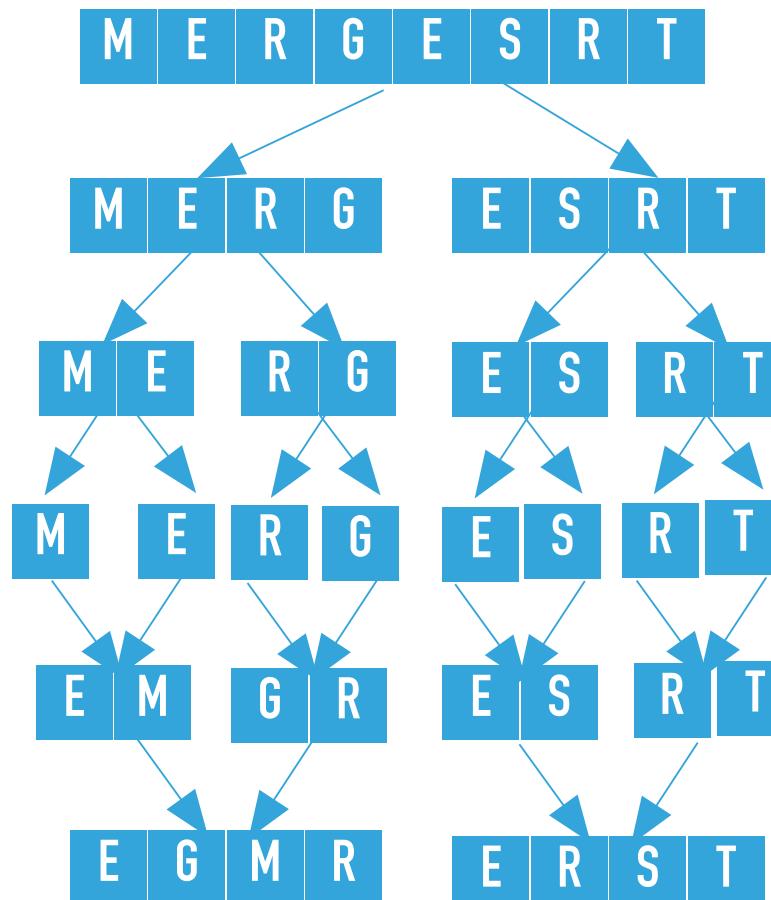
private static void sort(Comparable[] a, Comparable[] aux, int lo, int hi) {
    if (hi <= lo)
        return;
    int mid = lo + (hi - lo) / 2;
    sort(a, aux, lo, mid);
    sort(a, aux, mid+1, hi);
    merge(a, aux, lo, mid, hi);
}
  
```

sort([E, G, M, R, E, S, R, T], [E, M, G, R, E, S, null, null], 6, 7) merges the two subarrays that is calls merge([E, G, M, R, E, S, R, T], [E, M, G, R, E, S, null, null], 6, 6, 7), where `lo = 6`, `mid = 6`, and `hi = 7`. The resulting partially sorted array is [E, G, M, R, E, S, R, T].



```
private static void sort(Comparable[] a, Comparable[] aux, int lo, int hi) {
    if (hi <= lo)
        return;
    int mid = lo + (hi - lo) / 2;
    sort(a, aux, lo, mid);
    sort(a, aux, mid+1, hi);
    merge(a, aux, lo, mid, hi);
}
```

sort([E, G, M, R, E, S, R, T], [E, M, G, R, E, S, R, T], 4, 7) merges the two subarrays that is calls merge([E, G, M, R, E, S, R, T], [E, M, G, R, E, S, R, T], 4, 5, 7), where lo = 4, mid = 5, and hi = 7. The resulting partially sorted array is [E, G, M, R, E, R, S, T].



```
private static void sort(Comparable[] a, Comparable[] aux, int lo, int hi) {
    if (hi <= lo)
        return;
    int mid = lo + (hi - lo) / 2;
    sort(a, aux, lo, mid);
    sort(a, aux, mid+1, hi);
    merge(a, aux, lo, mid, hi);
}
```



sort([E, G, M, R, E, R, S, T], [E, M, G, R, E, S, R, T], 0, 7) merges the two subarrays that is calls merge([E, G, M, R, E, R, S, T], [E, M, G, R, E, S, R, T], 0, 3, 7), where lo = 0, mid = 3, and hi = 7. The resulting sorted array is [E, E, G, M, R, R, S, T].

Practice time

Which of the following subarray lengths will occur when running mergesort on an array of length 10?

- A. { 1, 2, 3, 5, 10 }
- B. { 2, 4, 6, 8, 10 }
- C. { 1, 2, 5, 10 }
- D. { 1,2,3,4,5,10}

Answer

Which of the following subarray lengths will occur when running mergesort on an array of length 10?

- A. { 1, 2, 3, 5, 10 }

Good algorithms are better than supercomputers

- ▶ Your laptop executes 10^8 comparisons per second
- ▶ A supercomputer executes 10^{12} comparisons per second

	Insertion sort			Mergesort		
Computer	Thousand inputs	Million inputs	Billion inputs	Thousand inputs	Million inputs	Billion inputs
Home	Instant	2 hours	300 years	instant	1 sec	15 min
Supercomputer	Instant	1 second	1 week	instant	instant	instant

Analysis of comparisons

- ▶ We will assume that n is a power of 2 ($n = 2^k$, where $k = \log_2 n$) and the number of comparisons $T(n)$ to sort an array of length n with merge sort satisfies the recurrence:
 - ▶ $T(n) = T(n/2) + T(n/2) + (n - 1) = O(n \log n)$
- ▶ Number of array accesses (rather than exchanges, here) is also $O(n \log n)$.

Mergesort uses $\leq n \log n$ compares to sort an array of length n

If $n = 4$, 2 levels

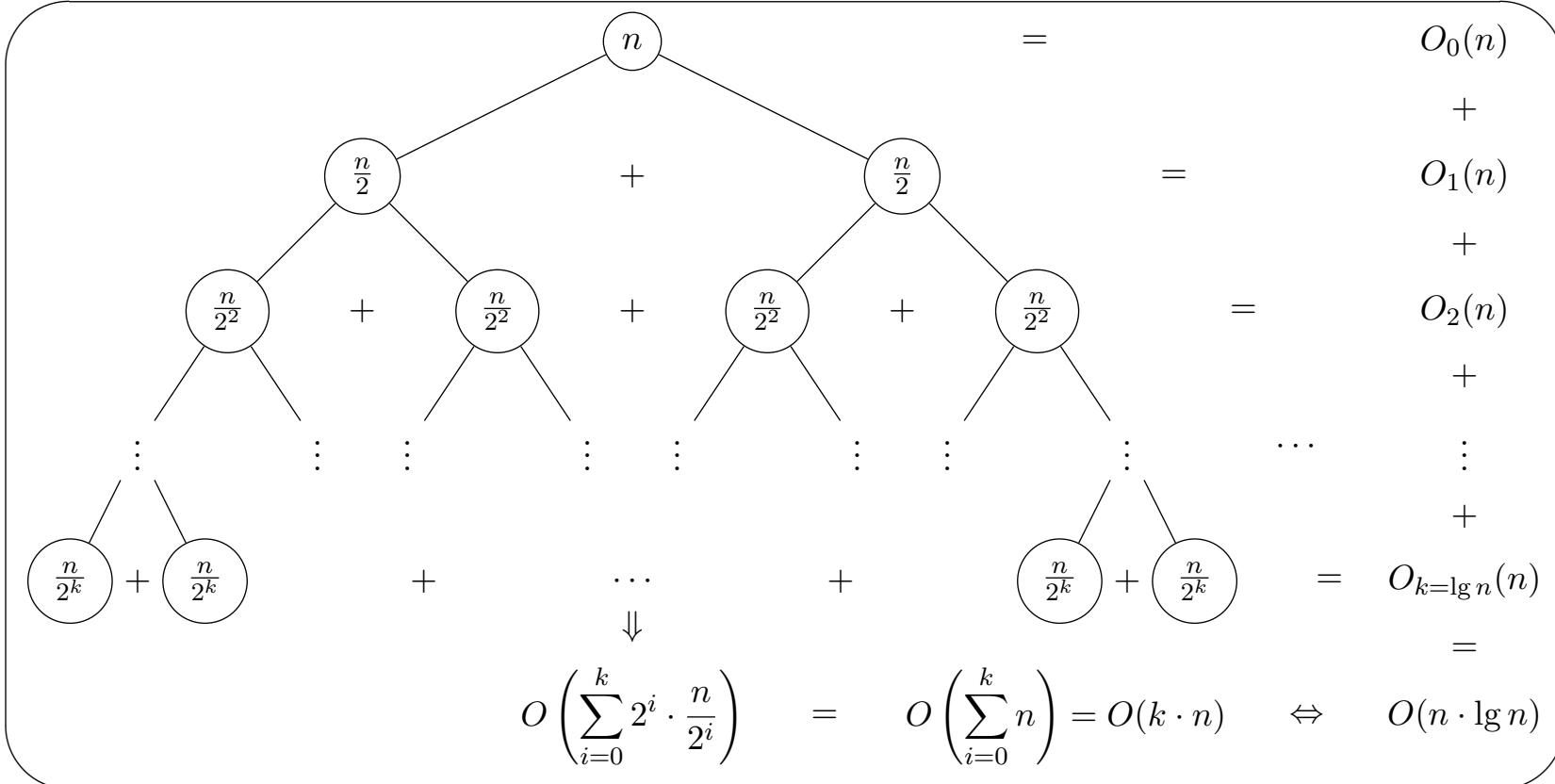
If $n = 8$, 3 levels

If $n = 16$, 4 levels

...

If $n = 2^k$, k levels,

or $k = \log_2 n$



Any algorithm with the same structure takes $n \log n$ time

```
public static void f(int n) {  
    if (n == 0)  
        return;  
    f(n/2);  
    f(n/2);  
    linear(n);  
}
```

Mergesort basics

- ▶ Auxiliary memory is proportional to n , as `aux[]` needs to be of length n for the last merge.
- ▶ At its simplest form, merge sort is **not an in-place algorithm**.
- ▶ There are modifications for halting the size of the auxiliary array but in-place merge is very hard.
- ▶ **Stable:** Look into `merge()`, if equal keys, it takes them from the left subarray.
 - ▶ So is insertion sort, but not selection sort.

Practical improvements for Mergesort

- ▶ Use insertion sort for small subarrays.
- ▶ Stop if already sorted.
- ▶ Eliminate the copy to the auxiliary array by saving time (not space).

```
private static void sort(Comparable[] src, Comparable[] dst, int lo, int hi) {  
    if (hi <= lo + 7) {  
        insertionSort(dst, lo, hi);  
        return;  
    }  
  
    int mid = lo + (hi - lo) / 2;  
    sort(dst, src, lo, mid);  
    sort(dst, src, mid+1, hi);  
  
    if (!less(src[mid+1], src[mid])) {  
        for (int i = lo; i <= hi; i++) dst[i] = src[i];  
        return;  
    }  
  
    merge(src, dst, lo, mid, hi);  
}
```

- ▶ For years, Java used this version to sort Collections of objects.

Sorting: the story so far

	In place	Stable	Best	Average	Worst	Remarks
Selection	X		$O(n^2)$	$O(n^2)$	$O(n^2)$	n exchanges
Insertion	X	X	$O(n)$	$O(n^2)$	$O(n^2)$	Use for small arrays or partially
Merge sort		X	$O(n \log n)$	$O(n \log n)$	$O(n \log n)$	Guaranteed performance; stable

Lecture 13: Mergesort

- ▶ Mergesort

Readings:

- ▶ Textbook:
 - ▶ Chapter 2.2 (pages 270-277)
- ▶ Website:
 - ▶ Mergesort: <https://algs4.cs.princeton.edu/22mergesort/>
 - ▶ Code: <https://algs4.cs.princeton.edu/22mergesort/Merge.java.html>

Practice Problems:

- ▶ 2.2.1-2.2.2, 2.2.11

TEXT

Lecture 14: Quicksort

- ▶ This week's assignment
- ▶ Quicksort

Basics

- ▶ Divide-and-conquer method
- ▶ Invented by [Sir Tony Hoare](https://en.wikipedia.org/wiki/Tony_Hoare) in 1959.
 - ▶ Wanted to sort words before looking them up in dictionary.
 - ▶ Came up with quicksort but did not know how to implement it.
 - ▶ Learned Algol 60 and recursion and implemented it.
 - ▶ Won the 1980 Turing Award (also invented null - and regretted it).
- ▶ [Bob Sedgewick](#) (author of our textbook) refined and analyzed many versions of quicksort.



https://en.wikipedia.org/wiki/Tony_Hoare

Algorithm sketch

- ▶ **Shuffle** the array.
- ▶ **Partition** so that, for some pivot j :
 - ▶ Entry $a[j]$ is in place.
 - ▶ There is no larger entry to the left of j .
 - ▶ No smaller entry to the right of j .
- ▶ **Sort** each subarray recursively.



input	Q	U	I	C	K	S	O	R	T	E	X	A	M	P	L	E
shuffle	K	R	A	T	E	L	E	P	U	I	M	Q	C	X	O	S
partition	E	C	A	I	E	K	L	P	U	T	M	Q	R	X	O	S
	not greater					partitioning element					not less					
sort left	A	C	E	E	I	K	L	P	U	T	M	Q	R	X	O	S
sort right	A	C	E	E	I	K	L	M	O	P	Q	R	S	T	U	X
result	A	C	E	E	I	K	L	M	O	P	Q	R	S	T	U	X

Quicksort overview

Quicksort Code

```
// quicksort the subarray from a[lo] to a[hi]
private static void sort(Comparable[] a, int lo, int hi) {
    if (hi <= lo) return;
    int j = partition(a, lo, hi);
    sort(a, lo, j-1);
    sort(a, j+1, hi);
}

/**
 * Rearranges the array in ascending order, using the natural order.
 * @param a the array to be sorted
 */
public static void sort(Comparable[] a) {
    StdRandom.shuffle(a);
    sort(a, 0, a.length - 1);
}
```

Hoare's partition scheme (faster than Lomuto's you will see in CS140)

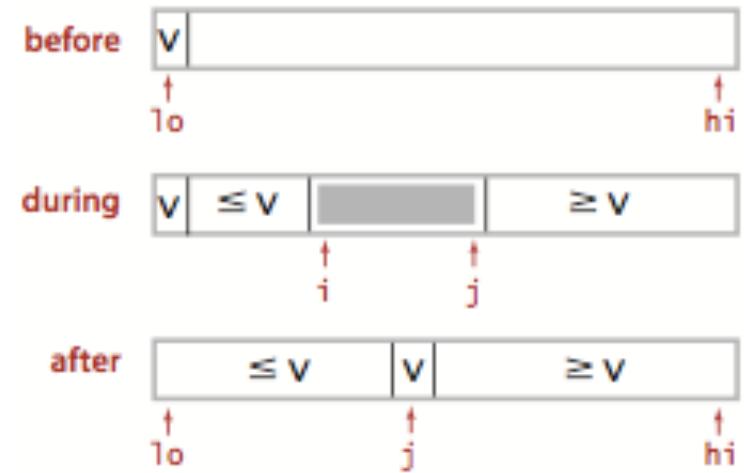
- ▶ Partition the subarray $a[lo..hi]$ so that
 $a[lo..j-1] \leq a[j] \leq a[j+1..hi]$
- ▶ Start with pivot at lo , pointer i at lo and pointer j at $hi+1$.
- ▶ Repeat the following until pointers i and j cross:
 - ▶ Scan i from left to right as long as $a[lo] > a[i]$.
 - ▶ (i.e. stop when you find a key that is greater than (or equal to) the pivot)
 - ▶ Scan j from right to left as long as $a[lo] < a[j]$.
 - ▶ (i.e. stop when you find a key that is less than (or equal to) the pivot)
 - ▶ Exchange $a[i]$ with $a[j]$.
- ▶ Exchange $a[lo]$ and $a[j]$. Return j .

Partition Code

```

// partition the subarray a[lo..hi] so that a[lo..j-1] <= a[j] <= a[j+1..hi] and return the index j.
private static int partition(Comparable[] a, int lo, int hi) {
    int i = lo;
    int j = hi + 1;
    Comparable v = a[lo];
    while (true) {
        // find item greater than (or equal to) v to swap
        while (less(a[++i], v)) {
            if (i == hi) break;
        }
        // find item smaller than (or equal to) v to swap
        while (less(v, a[--j])) {
            if (j == lo) break; // redundant since a[lo] acts as sentinel
        }
        // check if pointers cross
        if (i >= j) break;
        exch(a, i, j);
    }
    // put partitioning item v at a[j]
    exch(a, lo, j);
    // now, a[lo .. j-1] <= a[j] <= a[j+1 .. hi]
    return j;
}

```



Quicksort partitioning overview

Quicksort Example - Sort [Q,U,I,C,K,S,O,R,T,E,X,A,M,P,L,E]

```
public static void sort(Comparable[] a) {  
    StdRandom.shuffle(a);  
    sort(a, 0, a.length - 1);  
}  
  
// quicksort the subarray from a[lo] to a[hi]  
private static void sort(Comparable[] a, int lo, int hi) {  
    if (hi <= lo) return;  
    int j = partition(a, lo, hi);  
    sort(a, lo, j-1);  
    sort(a, j+1, hi);  
}
```

Shuffling resulted to:

K,R,A,T,E,L,E,P,U,I,M,Q,C,X,O,S

Call partition with $\text{lo}=0, \text{hi}=15$

K R A T E L E P U I M Q C X O S

0 1 2 3 4 5 6 7 8 9 10 11 12 13 14 15

lo hi

Partition Example, lo=0, hi=15

K	R	A	T	E	L	E	P	U	I	M	Q	C	X	O	S
0	1	2	3	4	5	6	7	8	9	10	11	12	13	14	15
lo															hi
i															j

```
// partition the subarray a[lo..hi] so that a[lo..j-1] <= a[j] <= a[j+1..hi] and return the index j.
private static int partition(Comparable[] a, int lo, int hi) {
    int i = lo;
    int j = hi + 1;
    Comparable v = a[lo];
```

Partition Example, $lo=0$, $hi=15$

K	R	A	T	E	L	E	P	U	I	M	Q	C	X	O	S
0	1	2	3	4	5	6	7	8	9	10	11	12	13	14	15
lo															hi
i															j

```
// find greater than (or equal to) v to swap
while (less(a[++i], v)) {
    if (i == hi) break;
}
```

Partition Example, $\text{lo}=0, \text{hi}=15$

K	R	A	T	E	L	E	P	U	I	M	Q	C	X	O	S
0	1	2	3	4	5	6	7	8	9	10	11	12	13	14	15
lo															hi
i															j

```
// find item smaller than (or equal to) v to swap
while (less(v, a[--j])) {
    if (j == lo) break; // redundant since a[lo] acts as sentinel
}
```

Partition Example, $\text{lo}=0, \text{hi}=15$

K	R	A	T	E	L	E	P	U	I	M	Q	C	X	O	S
0	1	2	3	4	5	6	7	8	9	10	11	12	13	14	15
lo															hi
i															j

```
// find item smaller than (or equal to) v to swap
while (less(v, a[--j])) {
    if (j == lo) break; // redundant since a[lo] acts as sentinel
}
```

Partition Example, $lo=0$, $hi=15$

K	R	A	T	E	L	E	P	U	I	M	Q	C	X	O	S
0	1	2	3	4	5	6	7	8	9	10	11	12	13	14	15
lo															hi
i															j

```
// find item smaller than (or equal to) v to swap
while (less(v, a[--j])) {
    if (j == lo) break; // redundant since a[lo] acts as sentinel
}
```

Partition Example, $\text{lo}=0, \text{hi}=15$

K	R	A	T	E	L	E	P	U	I	M	Q	C	X	O	S
0	1	2	3	4	5	6	7	8	9	10	11	12	13	14	15
lo															hi
i															j

```
// find item smaller than (or equal to) v to swap
while (less(v, a[--j])) {
    if (j == lo) break; // redundant since a[lo] acts as sentinel
}
```

Partition Example, $lo=0$, $hi=15$

K	R	A	T	E	L	E	P	U	I	M	Q	C	X	O	S
0	1	2	3	4	5	6	7	8	9	10	11	12	13	14	15
lo															hi
i											j				

```
exch(a, i, j);
```

Partition Example, $lo=0$, $hi=15$

K	C	A	T	E	L	E	P	U	I	M	Q	R	X	O	S
0	1	2	3	4	5	6	7	8	9	10	11	12	13	14	15
lo															hi
i											j				

```
exch(a, i, j);
```

Partition Example, lo=0, hi=15

K	C	A	T	E	L	E	P	U	I	M	Q	R	X	O	S
0	1	2	3	4	5	6	7	8	9	10	11	12	13	14	15
lo															hi
	i											j			

```
// find greater than (or equal to) v to swap
while (less(a[++i], v)) {
    if (i == hi) break;
}
```

Partition Example, $lo=0$, $hi=15$

K	C	A	T	E	L	E	P	U	I	M	Q	R	X	O	S
0	1	2	3	4	5	6	7	8	9	10	11	12	13	14	15
lo															hi
	i										j				

```
// find greater than (or equal to) v to swap
while (less(a[++i], v)) {
    if (i == hi) break;
}
```

Partition Example, lo=0, hi=15

```
// find item smaller than (or equal to) v to swap
while (less(v, a[--j])) {
    if (j == lo) break; // redundant since a[lo] acts as sentinel
}
```

Partition Example, $lo=0$, $hi=15$

K	C	A	T	E	L	E	P	U	I	M	Q	R	X	O	S
0	1	2	3	4	5	6	7	8	9	10	11	12	13	14	15
lo															hi

```
// find item smaller than (or equal to) v to swap
while (less(v, a[--j])) {
    if (j == lo) break; // redundant since a[lo] acts as sentinel
}
```

Partition Example, $\text{lo}=0, \text{hi}=15$

K	C	A	T	E	L	E	P	U	I	M	Q	R	X	O	S
0	1	2	3	4	5	6	7	8	9	10	11	12	13	14	15
lo															hi
	i							j							

```
// find item smaller than (or equal to) v to swap
while (less(v, a[--j])) {
    if (j == lo) break; // redundant since a[lo] acts as sentinel
}
```

Partition Example, $lo=0$, $hi=15$

K	C	A	T	E	L	E	P	U	I	M	Q	R	X	O	S
0	1	2	3	4	5	6	7	8	9	10	11	12	13	14	15
lo															hi
	i							j							

```
exch(a, i, j);
```

Partition Example, $lo=0$, $hi=15$

K	C	A	I	E	L	E	P	U	T	M	Q	R	X	O	S
0	1	2	3	4	5	6	7	8	9	10	11	12	13	14	15
lo														hi	

`exch(a, i, j);`

Partition Example, $lo=0$, $hi=15$

K	C	A	I	E	L	E	P	U	T	M	Q	R	X	O	S
0	1	2	3	4	5	6	7	8	9	10	11	12	13	14	15
lo														hi	

```
// find greater than (or equal to) v to swap
while (less(a[++i], v)) {
    if (i == hi) break;
}
```

Partition Example, $lo=0$, $hi=15$

K	C	A	I	E	L	E	P	U	T	M	Q	R	X	O	S
0	1	2	3	4	5	6	7	8	9	10	11	12	13	14	15
lo															hi

```
// find greater than (or equal to) v to swap
while (less(a[++i], v)) {
    if (i == hi) break;
}
```

Partition Example, $lo=0$, $hi=15$

K	C	A	I	E	L	E	P	U	T	M	Q	R	X	O	S
0	1	2	3	4	5	6	7	8	9	10	11	12	13	14	15
lo															hi

```
// find greater than (or equal to) v to swap
while (less(v, a[--j])) {
    if (j == lo) break; // redundant since a[lo] acts as sentinel
}
```

Partition Example, $lo=0$, $hi=15$

K	C	A	I	E	L	E	P	U	T	M	Q	R	X	O	S
0	1	2	3	4	5	6	7	8	9	10	11	12	13	14	15
lo														hi	
														i	j

```
// find greater than (or equal to) v to swap
while (less(v, a[--j])) {
    if (j == lo) break; // redundant since a[lo] acts as sentinel
}
```

Partition Example, $lo=0$, $hi=15$

K	C	A	I	E	L	E	P	U	T	M	Q	R	X	O	S
0	1	2	3	4	5	6	7	8	9	10	11	12	13	14	15
lo					i	j									hi

```
// find greater than (or equal to) v to swap
while (less(v, a[--j])) {
    if (j == lo) break; // redundant since a[lo] acts as sentinel
}
```

Partition Example, $lo=0$, $hi=15$

K	C	A	I	E	L	E	P	U	T	M	Q	R	X	O	S
0	1	2	3	4	5	6	7	8	9	10	11	12	13	14	15
lo														hi	
														i	j

```
exch(a, i, j);
```

Partition Example, $lo=0$, $hi=15$

K	C	A	I	E	E	L	P	U	T	M	Q	R	X	O	S
0	1	2	3	4	5	6	7	8	9	10	11	12	13	14	15
lo														hi	
														i	j

```
exch(a, i, j);
```

Partition Example, $lo=0$, $hi=15$

K	C	A	I	E	E	L	P	U	T	M	Q	R	X	O	S
0	1	2	3	4	5	6	7	8	9	10	11	12	13	14	15
lo														hi	

```
// find greater than (or equal to) v to swap
while (less(a[++i], v)) {
    if (i == hi) break;
}
```

Partition Example, $lo=0$, $hi=15$

K	C	A	I	E	E	L	P	U	T	M	Q	R	X	O	S
0	1	2	3	4	5	6	7	8	9	10	11	12	13	14	15
lo														hi	
														j	i

```
// find item smaller than (or equal to) v to swap
while (less(v, a[--j])) {
    if (j == lo) break; // redundant since a[lo] acts as sentinel
}
```

Partition Example, lo=0, hi=15

K	C	A	I	E	E	L	P	U	T	M	Q	R	X	O	S
0	1	2	3	4	5	6	7	8	9	10	11	12	13	14	15
lo														hi	
														j	i

```
// check if pointers cross
if (i >= j) break;
```

Partition Example, lo=0, hi=15

K	C	A	I	E	E	L	P	U	T	M	Q	R	X	O	S
0	1	2	3	4	5	6	7	8	9	10	11	12	13	14	15
lo														hi	

```
// put partitioning item v at a[j]
exch(a, lo, j);

// now, a[lo .. j-1] <= a[j] <= a[j+1 .. hi]
return j;
```

Partition Example, lo=0, hi=15

```
// put partitioning item v at a[j]
exch(a, lo, j);

// now, a[lo .. j-1] <= a[j] <= a[j+1 .. hi]
return j;
```

Partition for lo=0, hi=15 is complete. j=5

Call sort recursively on left subarray with lo=0, hi =4

```
private static void sort(Comparable[] a,  
int lo, int hi) {  
    if (hi <= lo) return;  
    int j = partition(a, lo, hi);  
    sort(a, lo, j-1);  
    sort(a, j+1, hi);  
}
```

```
private static void sort(Comparable[] a, int lo, int hi) {
    if (hi <= lo) return;
    int j = partition(a, lo, hi);
    sort(a, lo, j-1);
    sort(a, j+1, hi);
}
```

Quicksort Trace

	lo	j	hi	0	1	2	3	4	5	6	7	8	9	10	11	12	13	14	15	
initial values				Q	U	I	C	K	S	O	R	T	E	X	A	M	P	L	E	
random shuffle				K	R	A	T	E	L	E	P	U	I	M	Q	O	C	X	O	S
	0	5	15	E	C	A	I	E	K	L	P	U	T	M	Q	R	X	O	S	
	0	3	4	E	C	A	E	I	K	L	P	U	I	M	Q	R	X	O	S	
	0	2	2	A	C	E	E	I	K	L	P	U	T	M	Q	R	X	O	S	
	0	0	1	A	C	E	E	I	K	L	P	U	T	M	Q	R	X	O	S	
	1		1	A	C	E	E	I	K	L	P	U	T	M	Q	R	X	O	S	
	4		4	A	C	E	E	I	K	L	P	U	T	M	Q	R	X	O	S	
	6	6	15	A	C	E	E	I	K	L	P	U	T	M	Q	R	X	O	S	
	7	9	15	A	C	E	E	I	K	L	M	O	P	T	Q	R	X	U	S	
	7	7	8	A	C	E	E	I	K	L	M	O	P	T	Q	R	X	U	S	
	8		8	A	C	E	E	I	K	L	M	O	P	T	Q	R	X	U	S	
	10	13	15	A	C	E	E	I	K	L	M	O	P	S	Q	R	T	U	X	
	10	12	12	A	C	E	E	I	K	L	M	O	P	R	Q	S	T	U	X	
	10	11	11	A	C	E	E	I	K	L	M	O	P	Q	R	S	T	U	X	
	10		10	A	C	E	E	I	K	L	M	O	P	Q	R	S	T	U	X	
	14	14	15	A	C	E	E	I	K	L	M	O	P	Q	R	S	T	U	X	
	15		15	A	C	E	E	I	K	L	M	O	P	Q	R	S	T	U	X	
result				A	C	E	E	I	K	L	M	O	P	Q	R	S	T	U	X	

Partition Example, lo=0, hi=4

E	C	A	I	E	K	L	P	U	T	M	Q	R	X	O	S
0	1	2	3	4	5	6	7	8	9	10	11	12	13	14	15
lo	hi														
i	j														

```
// partition the subarray a[lo..hi] so that a[lo..j-1] <= a[j] <= a[j+1..hi] and return the index j.
private static int partition(Comparable[] a, int lo, int hi) {
    int i = lo;
    int j = hi + 1;
    Comparable v = a[lo];
```

Partition Example, $lo=0$, $hi=4$

E	C	A	I	E	K	L	P	U	T	M	Q	R	X	O	S
0	1	2	3	4	5	6	7	8	9	10	11	12	13	14	15
lo				hi											
i					j										

```
// find greater than (or equal to) v to swap
while (less(a[++i], v)) {
    if (i == hi) break;
}
```

Partition Example, $lo=0$, $hi=4$

E	C	A	I	E	K	L	P	U	T	M	Q	R	X	O	S
0	1	2	3	4	5	6	7	8	9	10	11	12	13	14	15
lo				hi											
	i				j										

```
// find greater than (or equal to) v to swap
while (less(a[++i], v)) {
    if (i == hi) break;
}
```

Partition Example, $lo=0$, $hi=4$

E	C	A	I	E	K	L	P	U	T	M	Q	R	X	O	S
0	1	2	3	4	5	6	7	8	9	10	11	12	13	14	15
lo				hi											
	i			j											

```
// find greater than (or equal to) v to swap
while (less(a[++i], v)) {
    if (i == hi) break;
}
```

Partition Example, $\text{lo}=0, \text{hi}=4$

E	C	A	I	E	K	L	P	U	T	M	Q	R	X	O	S
0	1	2	3	4	5	6	7	8	9	10	11	12	13	14	15
lo				hi											
	i	j													

```
// find item smaller than (or equal to) v to swap
while (less(v, a[--j])) {
    if (j == lo) break; // redundant since a[lo] acts as sentinel
}
```

Partition Example, $lo=0$, $hi=4$

E	C	A	I	E	K	L	P	U	T	M	Q	R	X	O	S
0	1	2	3	4	5	6	7	8	9	10	11	12	13	14	15
lo	hi														

```
exch(a, i, j);
```

Partition Example, $lo=0$, $hi=4$

E	C	A	E	I	K	L	P	U	T	M	Q	R	X	O	S
0	1	2	3	4	5	6	7	8	9	10	11	12	13	14	15
lo	hi														

exch(a , i , j);

Partition Example, $lo=0$, $hi=4$

E	C	A	E	I	K	L	P	U	T	M	Q	R	X	O	S
0	1	2	3	4	5	6	7	8	9	10	11	12	13	14	15
lo				hi											
				i,j											

```
// find greater than (or equal to) v to swap
while (less(a[++i], v)) {
    if (i == hi) break;
}
```

Partition Example, $lo=0$, $hi=4$

E	C	A	E	I	K	L	P	U	T	M	Q	R	X	O	S
0	1	2	3	4	5	6	7	8	9	10	11	12	13	14	15
lo				hi											
	j	i													

```
// find item smaller than (or equal to) v to swap
while (less(v, a[--j])) {
    if (j == lo) break; // redundant since a[lo] acts as sentinel
}
```

Partition Example, $lo=0$, $hi=4$

E	C	A	E	I	K	L	P	U	T	M	Q	R	X	O	S
0	1	2	3	4	5	6	7	8	9	10	11	12	13	14	15
lo				hi											
		j	i												

```
// check if pointers cross
if (i >= j) break;
```

Partition Example, $\text{lo}=0, \text{hi}=4$

E	C	A	E	I	K	L	P	U	T	M	Q	R	X	O	S
0	1	2	3	4	5	6	7	8	9	10	11	12	13	14	15
lo	hi														

```
// put partitioning item v at a[j]
exch(a, lo, j);

// now, a[lo .. j-1] <= a[j] <= a[j+1 .. hi]
return j;
```

Partition Example, lo=0, hi=4

E	C	A	E	I	K	L	P	U	T	M	Q	R	X	O	S
0	1	2	3	4	5	6	7	8	9	10	11	12	13	14	15
lo	hi														

```
private static void sort(Comparable[] a, int lo, int hi) {
    if (hi <= lo) return;
    int j = partition(a, lo, hi);
    sort(a, lo, j-1);
    sort(a, j+1, hi);
}
```

Quicksort Trace

	lo	j	hi	0	1	2	3	4	5	6	7	8	9	10	11	12	13	14	15	
initial values				Q	U	I	C	K	S	O	R	T	E	X	A	M	P	L	E	
random shuffle				K	R	A	T	E	L	E	P	U	I	M	Q	C	X	O	S	
	0	5	15	F	C	A	T	E	K	L	P	U	T	M	O	R	X	O	S	
	0	3	4	E	C	A	E	I	K	L	P	U	T	M	Q	R	X	O	S	
	0	2	2	A	C	E	E	I	K	L	P	U	T	M	Q	R	X	O	S	
	0	0	1	A	C	E	E	I	K	L	P	U	T	M	Q	R	X	O	S	
		1		1	A	C	E	E	I	K	L	P	U	T	M	Q	R	X	O	S
		4		4	A	C	E	E	I	K	L	P	U	T	M	Q	R	X	O	S
	6	6	15	A	C	E	E	I	K	L	P	U	T	M	Q	R	X	O	S	
	7	9	15	A	C	E	E	I	K	L	M	O	P	T	Q	R	X	U	S	
	7	7	8	A	C	E	E	I	K	L	M	O	P	T	Q	R	X	U	S	
	8	8		A	C	E	E	I	K	L	M	O	P	T	Q	R	X	U	S	
	10	13	15	A	C	E	E	I	K	L	M	O	P	S	Q	R	T	U	X	
	10	12	12	A	C	E	E	I	K	L	M	O	P	R	Q	S	T	U	X	
	10	11	11	A	C	E	E	I	K	L	M	O	P	Q	R	S	T	U	X	
	10	10		A	C	E	E	I	K	L	M	O	P	Q	R	S	T	U	X	
	14	14	15	A	C	E	E	I	K	L	M	O	P	Q	R	S	T	U	X	
	15			A	C	E	E	I	K	L	M	O	P	Q	R	S	T	U	X	
result				A	C	E	E	I	K	L	M	O	P	Q	R	S	T	U	X	

Partition Example, lo=0, hi=2

E	C	A	E	I	K	L	P	U	T	M	Q	R	X	O	S
0	1	2	3	4	5	6	7	8	9	10	11	12	13	14	15
lo	hi														
i		j													

```
// partition the subarray a[lo..hi] so that a[lo..j-1] <= a[j] <= a[j+1..hi] and return the index j.
private static int partition(Comparable[] a, int lo, int hi) {
    int i = lo;
    int j = hi + 1;
    Comparable v = a[lo];
```

Partition Example, $lo=0$, $hi=2$

E	C	A	E	I	K	L	P	U	T	M	Q	R	X	O	S
0	1	2	3	4	5	6	7	8	9	10	11	12	13	14	15
lo	hi														
i		j													

```
// find greater than (or equal to) v to swap
while (less(a[++i], v)) {
    if (i == hi) break;
}
```

Partition Example, lo=0, hi=2

E	C	A	E	I	K	L	P	U	T	M	Q	R	X	O	S
0	1	2	3	4	5	6	7	8	9	10	11	12	13	14	15
lo	hi														
	i	j													

```
// find greater than (or equal to) v to swap
while (less(a[++i], v)) {
    if (i == hi) break;
}
```

Partition Example, $\text{lo}=0, \text{hi}=2$

E	C	A	E	I	K	L	P	U	T	M	Q	R	X	O	S
0	1	2	3	4	5	6	7	8	9	10	11	12	13	14	15
lo	hi														
		i,j													

```
// find item smaller than (or equal to) v to swap
while (less(v, a[--j])) {
    if (j == lo) break; // redundant since a[lo] acts as sentinel
}
```

Partition Example, lo=0, hi=2

E	C	A	E	I	K	L	P	U	T	M	Q	R	X	O	S
0	1	2	3	4	5	6	7	8	9	10	11	12	13	14	15
lo	hi														
	i,j														

```
// check if pointers cross
if (i >= j) break;
```

Partition Example, lo=0, hi=2

E	C	A	E	I	K	L	P	U	T	M	Q	R	X	O	S
0	1	2	3	4	5	6	7	8	9	10	11	12	13	14	15
lo	hi														
		i,j													

```
// put partitioning item v at a[j]
exch(a, lo, j);

// now, a[lo .. j-1] <= a[j] <= a[j+1 .. hi]
return j;
```

Partition Example, lo=0, hi=2

A	C	E	E	I	K	L	P	U	T	M	Q	R	X	O	S
0	1	2	3	4	5	6	7	8	9	10	11	12	13	14	15
lo	hi														
		i,j													

Partition for lo=0, hi=2 is complete. j=2

Call sort recursively on left subarray with lo=0, hi =1

```
// put partitioning item v at a[j]
exch(a, lo, j);

// now, a[lo .. j-1] <= a[j] <= a[j+1 .. hi]
return j;
```

```
private static void sort(Comparable[] a,
int lo, int hi) {
    if (hi <= lo) return;
    int j = partition(a, lo, hi);
    sort(a, lo, j-1);
    sort(a, j+1, hi);
}
```

```
private static void sort(Comparable[] a, int lo, int hi) {
    if (hi <= lo) return;
    int j = partition(a, lo, hi);
    sort(a, lo, j-1);
    sort(a, j+1, hi);
}
```

Quicksort Trace

	lo	j	hi	0	1	2	3	4	5	6	7	8	9	10	11	12	13	14	15	
initial values				Q	U	I	C	K	S	O	R	T	E	X	A	M	P	L	E	
random shuffle				K	R	A	T	E	L	E	P	U	I	M	Q	C	X	O	S	
	0	5	15	E	C	A	I	E	K	L	P	U	T	M	Q	R	X	O	S	
	0	3	4	E	C	A	E	I	K	L	P	U	T	M	O	R	X	O	S	
	0	2	2	A	C	E	E	I	K	L	P	U	T	M	Q	R	X	O	S	
	0	0	1	A	C	E	E	I	K	L	P	U	T	M	Q	R	X	O	S	
		1		1	A	C	E	E	I	K	L	P	U	T	M	Q	R	X	O	S
		4		4	A	C	E	E	I	K	L	P	U	T	M	Q	R	X	O	S
	6	6	15	A	C	E	E	I	K	L	P	U	T	M	Q	R	X	O	S	
	7	9	15	A	C	E	E	I	K	L	M	O	P	T	Q	R	X	U	S	
	7	7	8	A	C	E	E	I	K	L	M	O	P	T	Q	R	X	U	S	
	8	8		A	C	E	E	I	K	L	M	O	P	T	Q	R	X	U	S	
	10	13	15	A	C	E	E	I	K	L	M	O	P	S	Q	R	T	U	X	
	10	12	12	A	C	E	E	I	K	L	M	O	P	R	Q	S	T	U	X	
	10	11	11	A	C	E	E	I	K	L	M	O	P	Q	R	S	T	U	X	
	10	10		A	C	E	E	I	K	L	M	O	P	Q	R	S	T	U	X	
	14	14	15	A	C	E	E	I	K	L	M	O	P	Q	R	S	T	U	X	
	15			A	C	E	E	I	K	L	M	O	P	Q	R	S	T	U	X	
result				A	C	E	E	I	K	L	M	O	P	Q	R	S	T	U	X	

Partition Example, lo=0, hi=1

A	C	E	E	I	K	L	P	U	T	M	Q	R	X	O	S
0	1	2	3	4	5	6	7	8	9	10	11	12	13	14	15
lo	hi														
i	j														

```
// partition the subarray a[lo..hi] so that a[lo..j-1] <= a[j] <= a[j+1..hi] and return the index j.
private static int partition(Comparable[] a, int lo, int hi) {
    int i = lo;
    int j = hi + 1;
    Comparable v = a[lo];
```

Partition Example, $lo=0, hi=1$

A	C	E	E	I	K	L	P	U	T	M	Q	R	X	O	S
0	1	2	3	4	5	6	7	8	9	10	11	12	13	14	15
lo	hi														
i	j														

```
// find greater than (or equal to) v to swap
while (less(a[++i], v)) {
    if (i == hi) break;
}
```

Partition Example, $\text{lo}=0, \text{hi}=1$

A	C	E	E	I	K	L	P	U	T	M	Q	R	X	O	S
0	1	2	3	4	5	6	7	8	9	10	11	12	13	14	15
lo	hi														
i,j															

```
// find item smaller than (or equal to) v to swap
while (less(v, a[--j])) {
    if (j == lo) break; // redundant since a[lo] acts as sentinel
}
```

Partition Example, $\text{lo}=0, \text{hi}=1$

A	C	E	E	I	K	L	P	U	T	M	Q	R	X	O	S
0	1	2	3	4	5	6	7	8	9	10	11	12	13	14	15
lo	hi														
j	i														

```
// find item smaller than (or equal to) v to swap
while (less(v, a[--j])) {
    if (j == lo) break; // redundant since a[lo] acts as sentinel
}
```

Partition Example, $\text{lo}=0, \text{hi}=1$

A	C	E	E	I	K	L	P	U	T	M	Q	R	X	O	S
0	1	2	3	4	5	6	7	8	9	10	11	12	13	14	15
lo	hi														
j	i														

```
// check if pointers cross
if (i >= j) break;
```

Partition Example, $\text{lo}=0, \text{hi}=1$

A	C	E	E	I	K	L	P	U	T	M	Q	R	X	O	S
0	1	2	3	4	5	6	7	8	9	10	11	12	13	14	15
lo	hi														
j	i														

```
// put partitioning item v at a[j]
exch(a, lo, j);

// now, a[lo .. j-1] <= a[j] <= a[j+1 .. hi]
return j;
```

Partition Example, lo=0, hi=1

A	C	E	E	I	K	L	P	U	T	M	Q	R	X	O	S
0	1	2	3	4	5	6	7	8	9	10	11	12	13	14	15
lo	hi														
j	i														

```
// put partitioning item v at a[j]
exch(a, lo, j);

// now, a[lo .. j-1] <= a[j] <= a[j+1 .. hi]
return j;
```

Partition for lo=0, hi=1 is complete. j=0

No partition for subarrays of size 1.

Call sort recursively on right subarray with lo=6, hi =15

```
private static void sort(Comparable[] a, int lo, int hi) {

    if (hi <= lo) return;
    int j = partition(a, lo, hi);
    sort(a, lo, j-1);
    sort(a, j+1, hi);
}
```

```
private static void sort(Comparable[] a, int lo, int hi) {
    if (hi <= lo) return;
    int j = partition(a, lo, hi);
    sort(a, lo, j-1);
    sort(a, j+1, hi);
}
```

Quicksort Trace

	lo	j	hi	0	1	2	3	4	5	6	7	8	9	10	11	12	13	14	15
initial values				Q	U	I	C	K	S	O	R	T	E	X	A	M	P	L	E
random shuffle				K	R	A	T	E	L	E	P	U	I	M	Q	C	X	O	S
	0	5	15	E	C	A	I	E	K	L	P	U	T	M	Q	R	X	O	S
	0	3	4	E	C	A	E	I	K	L	P	U	T	M	Q	R	X	O	S
	0	2	2	A	C	F	F	T	K	L	P	U	T	M	Q	R	X	O	S
	0	0	1	A	C	E	E	I	K	L	P	U	T	M	Q	R	X	O	S
<td data-kind="ghost"></td>																			
<td data-kind="ghost"></td>																			

no partition
for subarrays
of size 1

Partition Example, lo=6, hi=15

A	C	E	E	I	K	L	P	U	T	M	Q	R	X	O	S
0	1	2	3	4	5	6	7	8	9	10	11	12	13	14	15
						lo									hi
						i									j

```
// partition the subarray a[lo..hi] so that a[lo..j-1] <= a[j] <= a[j+1..hi] and return the index j.
private static int partition(Comparable[] a, int lo, int hi) {
    int i = lo;
    int j = hi + 1;
    Comparable v = a[lo];
```

Partition Example, $lo=6$, $hi=15$

A	C	E	E	I	K	L	P	U	T	M	Q	R	X	O	S
0	1	2	3	4	5	6	7	8	9	10	11	12	13	14	15
lo												hi			
i												j			

```
// find greater than (or equal to) v to swap
while (less(a[++i], v)) {
    if (i == hi) break;
}
```

Partition Example, $lo=6$, $hi=15$

A	C	E	E	I	K	L	P	U	T	M	Q	R	X	O	S
0	1	2	3	4	5	6	7	8	9	10	11	12	13	14	15
lo												hi			
i												j			

```
// find item smaller than (or equal to) v to swap
while (less(v, a[--j])) {
    if (j == lo) break; // redundant since a[lo] acts as sentinel
}
```

Partition Example, $lo=6$, $hi=15$

A	C	E	E	I	K	L	P	U	T	M	Q	R	X	O	S
0	1	2	3	4	5	6	7	8	9	10	11	12	13	14	15
lo														hi	
<i>i</i>														<i>j</i>	

```
// find item smaller than (or equal to) v to swap
while (less(v, a[--j])) {
    if (j == lo) break; // redundant since a[lo] acts as sentinel
}
```

Partition Example, $lo=6$, $hi=15$

A	C	E	E	I	K	L	P	U	T	M	Q	R	X	O	S
0	1	2	3	4	5	6	7	8	9	10	11	12	13	14	15
lo														hi	
i														j	

```
// find item smaller than (or equal to) v to swap
while (less(v, a[--j])) {
    if (j == lo) break; // redundant since a[lo] acts as sentinel
}
```

Partition Example, $lo=6$, $hi=15$

A	C	E	E	I	K	L	P	U	T	M	Q	R	X	O	S
0	1	2	3	4	5	6	7	8	9	10	11	12	13	14	15
lo														hi	

```
// find item smaller than (or equal to) v to swap
while (less(v, a[--j])) {
    if (j == lo) break; // redundant since a[lo] acts as sentinel
}
```

Partition Example, $lo=6$, $hi=15$

A	C	E	E	I	K	L	P	U	T	M	Q	R	X	O	S
0	1	2	3	4	5	6	7	8	9	10	11	12	13	14	15
lo														hi	

```
// find item smaller than (or equal to) v to swap
while (less(v, a[--j])) {
    if (j == lo) break; // redundant since a[lo] acts as sentinel
}
```

Partition Example, $lo=6$, $hi=15$

A	C	E	E	I	K	L	P	U	T	M	Q	R	X	O	S
0	1	2	3	4	5	6	7	8	9	10	11	12	13	14	15
lo														hi	

```
// find item smaller than (or equal to) v to swap
while (less(v, a[--j])) {
    if (j == lo) break; // redundant since a[lo] acts as sentinel
}
```

Partition Example, $lo=6$, $hi=15$

A	C	E	E	I	K	L	P	U	T	M	Q	R	X	O	S
0	1	2	3	4	5	6	7	8	9	10	11	12	13	14	15
lo														hi	

```
// find item smaller than (or equal to) v to swap
while (less(v, a[--j])) {
    if (j == lo) break; // redundant since a[lo] acts as sentinel
}
```

Partition Example, $lo=6$, $hi=15$

A	C	E	E	I	K	L	P	U	T	M	Q	R	X	O	S
0	1	2	3	4	5	6	7	8	9	10	11	12	13	14	15
lo														hi	

```
// find item smaller than (or equal to) v to swap
while (less(v, a[--j])) {
    if (j == lo) break; // redundant since a[lo] acts as sentinel
}
```

Partition Example, lo=6, hi=15

A	C	E	E	I	K	L	P	U	T	M	Q	R	X	O	S
0	1	2	3	4	5	6	7	8	9	10	11	12	13	14	15
lo														hi	
i,j															

```
// find item smaller than (or equal to) v to swap
while (less(v, a[--j])) {
    if (j == lo) break; // redundant since a[lo] acts as sentinel
}
```

Partition Example, $lo=6$, $hi=15$

A	C	E	E	I	K	L	P	U	T	M	Q	R	X	O	S
0	1	2	3	4	5	6	7	8	9	10	11	12	13	14	15
lo														hi	
<i>j</i> <i>i</i>															

```
// find item smaller than (or equal to) v to swap
while (less(v, a[--j])) {
    if (j == lo) break; // redundant since a[lo] acts as sentinel
}
```

Partition Example, $lo=6$, $hi=15$

A	C	E	E	I	K	L	P	U	T	M	Q	R	X	O	S
0	1	2	3	4	5	6	7	8	9	10	11	12	13	14	15
lo														hi	

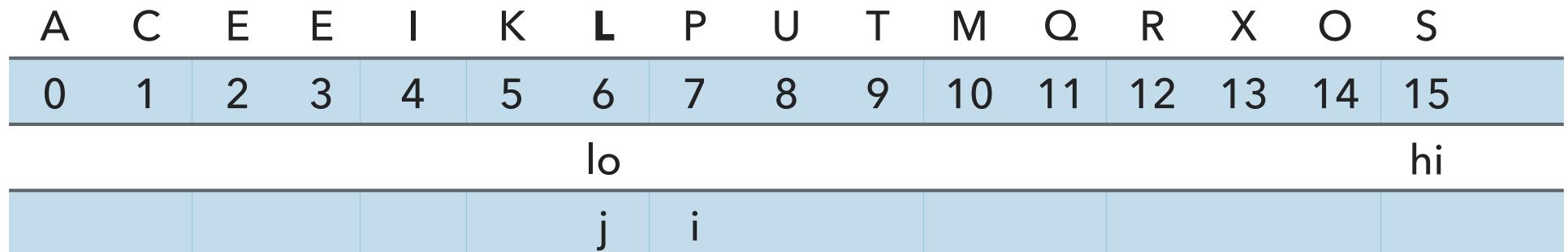
```
// check if pointers cross
if (i >= j) break;
```

Partition Example, $lo=6$, $hi=15$

A	C	E	E	I	K	L	P	U	T	M	Q	R	X	O	S
0	1	2	3	4	5	6	7	8	9	10	11	12	13	14	15
lo														hi	

```
// put partitioning item v at a[j]
exch(a, lo, j);

// now, a[lo .. j-1] <= a[j] <= a[j+1 .. hi]
return j;
```

Partition Example, $lo=6$, $hi=15$ 

```
// put partitioning item v at a[j]
exch(a, lo, j);

// now, a[lo .. j-1] <= a[j] <= a[j+1 .. hi]
return j;
```

Partition for $lo=6$, $hi=15$ is complete. $j=6$

No partition for subarrays of size 1.

Call sort recursively on right subarray with $lo=7$, $hi = 15$

```
private static void sort(Comparable[] a, int lo, int hi) {

    if (hi <= lo) return;
    int j = partition(a, lo, hi);
    sort(a, lo, j-1);
    sort(a, j+1, hi);
}
```

```
private static void sort(Comparable[] a, int lo, int hi) {
    if (hi <= lo) return;
    int j = partition(a, lo, hi);
    sort(a, lo, j-1);
    sort(a, j+1, hi);
}
```

Quicksort Trace

	lo	j	hi	0	1	2	3	4	5	6	7	8	9	10	11	12	13	14	15	
initial values				Q	U	I	C	K	S	O	R	T	E	X	A	M	P	L	E	
random shuffle				K	R	A	T	E	L	E	P	U	I	M	Q	C	X	O	S	
	0	5	15	E	C	A	I	E	K	L	P	U	T	M	Q	R	X	O	S	
	0	3	4	E	C	A	E	I	K	L	P	U	T	M	Q	R	X	O	S	
	0	2	2	A	C	E	E	I	K	L	P	U	T	M	Q	R	X	O	S	
	0	0	1	A	C	E	E	I	K	L	P	U	T	M	Q	R	X	O	S	
	1			1	A	C	E	I	K	L	P	U	T	M	Q	R	X	O	S	
	4			4	A	C	F	F	T	K	I	P	U	T	M	Q	R	X	O	S
	6	6	15	A	C	E	E	I	K	L	P	U	T	M	Q	R	X	O	S	
<i>no partition for subarrays of size 1</i>	7	9	15	A	C	E	E	I	K	L	M	O	P	T	Q	R	X	U	S	
	7	7	8	A	C	E	E	I	K	L	M	O	P	T	Q	R	X	U	S	
	8			8	A	C	E	E	I	K	L	M	O	P	T	Q	R	X	U	S
	10	13	15	A	C	E	E	I	K	L	M	O	P	S	Q	R	T	U	X	
	10	12	12	A	C	E	E	I	K	L	M	O	P	R	Q	S	T	U	X	
	10	11	11	A	C	E	E	I	K	L	M	O	P	Q	R	S	T	U	X	
	10	10		A	C	E	E	I	K	L	M	O	P	Q	R	S	T	U	X	
	14	14	15	A	C	E	E	I	K	L	M	O	P	Q	R	S	T	U	X	
	15			A	C	E	E	I	K	L	M	O	P	Q	R	S	T	U	X	
result				A	C	E	E	I	K	L	M	O	P	Q	R	S	T	U	X	

Partition Example, lo=7, hi=15

A	C	E	E	I	K	L	P	U	T	M	Q	R	X	O	S
0	1	2	3	4	5	6	7	8	9	10	11	12	13	14	15
							lo								hi

							i								j
--	--	--	--	--	--	--	---	--	--	--	--	--	--	--	---

```
// partition the subarray a[lo..hi] so that a[lo..j-1] <= a[j] <= a[j+1..hi] and return the index j.
private static int partition(Comparable[] a, int lo, int hi) {
    int i = lo;
    int j = hi + 1;
    Comparable v = a[lo];
```

Partition Example, $lo=7$, $hi=15$

A	C	E	E	I	K	L	P	U	T	M	Q	R	X	O	S
0	1	2	3	4	5	6	7	8	9	10	11	12	13	14	15
							lo								hi

Below the array, the indices i and j are shown, indicating the current partitioning step.

```
// find greater than (or equal to) v to swap
while (less(a[++i], v)) {
    if (i == hi) break;
}
```

Partition Example, $lo=7$, $hi=15$

A	C	E	E	I	K	L	P	U	T	M	Q	R	X	O	S
0	1	2	3	4	5	6	7	8	9	10	11	12	13	14	15
lo														hi	
i														j	

```
// find item smaller than (or equal to) v to swap
while (less(v, a[--j])) {
    if (j == lo) break; // redundant since a[lo] acts as sentinel
}
```

Partition Example, $lo=7$, $hi=15$

A	C	E	E	I	K	L	P	U	T	M	Q	R	X	O	S
0	1	2	3	4	5	6	7	8	9	10	11	12	13	14	15
lo														hi	
i														j	

```
// find item smaller than (or equal to) v to swap
while (less(v, a[--j])) {
    if (j == lo) break; // redundant since a[lo] acts as sentinel
}
```

Partition Example, $lo=7$, $hi=15$

A	C	E	E	I	K	L	P	U	T	M	Q	R	X	O	S
0	1	2	3	4	5	6	7	8	9	10	11	12	13	14	15
lo														hi	
i														j	

```
exch(a, i, j);
```

Partition Example, $lo=7$, $hi=15$

A	C	E	E	I	K	L	P	O	T	M	Q	R	X	U	S
0	1	2	3	4	5	6	7	8	9	10	11	12	13	14	15
lo														hi	
i														j	

```
exch(a, i, j);
```

Partition Example, $lo=7$, $hi=15$

A	C	E	E	I	K	L	P	O	T	M	Q	R	X	U	S
0	1	2	3	4	5	6	7	8	9	10	11	12	13	14	15
lo														hi	
														i	j

```
// find greater than (or equal to) v to swap
while (less(a[++i], v)) {
    if (i == hi) break;
}
```

Partition Example, $lo=7$, $hi=15$

A	C	E	E	I	K	L	P	O	T	M	Q	R	X	U	S
0	1	2	3	4	5	6	7	8	9	10	11	12	13	14	15
lo														hi	
i														j	

```
// find item smaller than (or equal to) v to swap
while (less(v, a[--j])) {
    if (j == lo) break; // redundant since a[lo] acts as sentinel
}
```

Partition Example, $lo=7$, $hi=15$

A	C	E	E	I	K	L	P	O	T	M	Q	R	X	U	S
0	1	2	3	4	5	6	7	8	9	10	11	12	13	14	15
lo														hi	
i														j	

```
// find item smaller than (or equal to) v to swap
while (less(v, a[--j])) {
    if (j == lo) break; // redundant since a[lo] acts as sentinel
}
```

Partition Example, $lo=7$, $hi=15$

A	C	E	E	I	K	L	P	O	T	M	Q	R	X	U	S
0	1	2	3	4	5	6	7	8	9	10	11	12	13	14	15
lo														hi	

```
// find item smaller than (or equal to) v to swap
while (less(v, a[--j])) {
    if (j == lo) break; // redundant since a[lo] acts as sentinel
}
```

Partition Example, $lo=7$, $hi=15$

A	C	E	E	I	K	L	P	O	T	M	Q	R	X	U	S
0	1	2	3	4	5	6	7	8	9	10	11	12	13	14	15
lo														hi	
i j															

```
// find item smaller than (or equal to) v to swap
while (less(v, a[--j])) {
    if (j == lo) break; // redundant since a[lo] acts as sentinel
}
```

Partition Example, $lo=7$, $hi=15$

A	C	E	E	I	K	L	P	O	T	M	Q	R	X	U	S
0	1	2	3	4	5	6	7	8	9	10	11	12	13	14	15
lo														hi	
														i	j

```
exch(a, i, j);
```

Partition Example, $lo=7$, $hi=15$

A	C	E	E	I	K	L	P	O	M	T	Q	R	X	U	S
0	1	2	3	4	5	6	7	8	9	10	11	12	13	14	15
lo														hi	
i j															

```
exch(a, i, j);
```

Partition Example, $lo=7$, $hi=15$

A	C	E	E	I	K	L	P	O	M	T	Q	R	X	U	S
0	1	2	3	4	5	6	7	8	9	10	11	12	13	14	15
lo														hi	
i, j															

```
// find greater than (or equal to) v to swap
while (less(a[++i], v)) {
    if (i == hi) break;
}
```

Partition Example, $lo=7$, $hi=15$

A	C	E	E	I	K	L	P	O	M	T	Q	R	X	U	S
0	1	2	3	4	5	6	7	8	9	10	11	12	13	14	15
lo														hi	

```
// find item smaller than (or equal to) v to swap
while (less(v, a[--j])) {
    if (j == lo) break; // redundant since a[lo] acts as sentinel
}
```

Partition Example, $lo=7$, $hi=15$

A	C	E	E	I	K	L	P	O	M	T	Q	R	X	U	S
0	1	2	3	4	5	6	7	8	9	10	11	12	13	14	15
							lo								hi

The diagram shows an array of 16 elements with indices from 0 to 15. The elements are: A, C, E, E, I, K, L, P, O, M, T, Q, R, X, U, S. The index labels are: 0, 1, 2, 3, 4, 5, 6, 7, 8, 9, 10, 11, 12, 13, 14, 15. Below the array, two pointers are indicated: 'lo' at index 7 and 'hi' at index 15. The partitioning process has reached a stage where the pointers 'j' and 'i' have crossed, as shown by the code snippet below.

```
// check if pointers cross
if (i >= j) break;
```

Partition Example, $lo=7$, $hi=15$

A	C	E	E	I	K	L	P	O	M	T	Q	R	X	U	S
0	1	2	3	4	5	6	7	8	9	10	11	12	13	14	15
lo														hi	

```
// put partitioning item v at a[j]
exch(a, lo, j);

// now, a[lo .. j-1] <= a[j] <= a[j+1 .. hi]
return j;
```

Partition Example, $lo=7$, $hi=15$

A	C	E	E	I	K	L	M	O	P	T	Q	R	X	U	S
0	1	2	3	4	5	6	7	8	9	10	11	12	13	14	15
lo														hi	
j i															

Partition for $lo=6$, $hi=15$ is complete. $j=9$

Call sort recursively on left subarray with $lo=7$, $hi = 8$

```
// put partitioning item v at a[j]
exch(a, lo, j);

// now, a[lo .. j-1] <= a[j] <= a[j+1 .. hi]
return j;
```

```
private static void sort(Comparable[] a, int lo, int hi) {
    if (hi <= lo) return;
    int j = partition(a, lo, hi);
    sort(a, lo, j-1);
    sort(a, j+1, hi);
}
```

```
private static void sort(Comparable[] a, int lo, int hi) {
    if (hi <= lo) return;
    int j = partition(a, lo, hi);
    sort(a, lo, j-1);
    sort(a, j+1, hi);
}
```

Quicksort Trace

	lo	j	hi	0	1	2	3	4	5	6	7	8	9	10	11	12	13	14	15
initial values				Q	U	I	C	K	S	O	R	T	E	X	A	M	P	L	E
random shuffle				K	R	A	T	E	L	E	P	U	I	M	Q	C	X	O	S
	0	5	15	E	C	A	I	E	K	L	P	U	T	M	Q	R	X	O	S
	0	3	4	E	C	A	E	I	K	L	P	U	T	M	Q	R	X	O	S
	0	2	2	A	C	E	E	I	K	L	P	U	T	M	Q	R	X	O	S
	0	0	1	A	C	E	E	I	K	L	P	U	T	M	Q	R	X	O	S
	1	1	1	A	C	E	E	I	K	L	P	U	T	M	Q	R	X	O	S
	4	4	4	A	C	E	E	I	K	L	P	U	T	M	Q	R	X	O	S
	6	6	15	A	C	E	E	T	K	L	P	U	T	M	Q	R	X	O	S
<i>no partition for subarrays of size 1</i>	7	9	15	A	C	E	E	I	K	L	M	O	P	T	Q	R	X	U	S
	7	7	8	A	C	E	E	I	K	L	M	O	P	T	Q	R	X	U	S
	8	8	8	A	C	E	E	I	K	L	M	O	P	T	Q	R	X	U	S
	10	13	15	A	C	E	E	I	K	L	M	O	P	S	Q	R	T	U	X
	10	12	12	A	C	E	E	I	K	L	M	O	P	R	Q	S	T	U	X
	10	11	11	A	C	E	E	I	K	L	M	O	P	Q	R	S	T	U	X
	10	10	10	A	C	E	E	I	K	L	M	O	P	Q	R	S	T	U	X
	14	14	15	A	C	E	E	I	K	L	M	O	P	Q	R	S	T	U	X
result				A	C	E	E	I	K	L	M	O	P	Q	R	S	T	U	X

Partition Example, lo=7, hi=8

A	C	E	E	I	K	L	M	O	P	T	Q	R	X	U	S
0	1	2	3	4	5	6	7	8	9	10	11	12	13	14	15
lo hi															
i j															

```
// partition the subarray a[lo..hi] so that a[lo..j-1] <= a[j] <= a[j+1..hi] and return the index j.
private static int partition(Comparable[] a, int lo, int hi) {
    int i = lo;
    int j = hi + 1;
    Comparable v = a[lo];
```

Partition Example, $lo=7, hi=8$

A	C	E	E	I	K	L	M	O	P	T	Q	R	X	U	S
0	1	2	3	4	5	6	7	8	9	10	11	12	13	14	15
							lo	hi							

i j

```
// find greater than (or equal to) v to swap
while (less(a[++i], v)) {
    if (i == hi) break;
}
```

Partition Example, $lo=7, hi=8$

A	C	E	E	I	K	L	M	O	P	T	Q	R	X	U	S
0	1	2	3	4	5	6	7	8	9	10	11	12	13	14	15
							lo	hi							
								i,j							

```
// find item smaller than (or equal to) v to swap
while (less(v, a[--j])) {
    if (j == lo) break; // redundant since a[lo] acts as sentinel
}
```

Partition Example, $lo=7$, $hi=8$

A	C	E	E	I	K	L	M	O	P	T	Q	R	X	U	S
0	1	2	3	4	5	6	7	8	9	10	11	12	13	14	15
$lo \quad hi$															
$j \quad i$															

```
// find item smaller than (or equal to) v to swap
while (less(v, a[--j])) {
    if (j == lo) break; // redundant since a[lo] acts as sentinel
}
```

Partition Example, $lo=7$, $hi=8$

A	C	E	E	I	K	L	M	O	P	T	Q	R	X	U	S
0	1	2	3	4	5	6	7	8	9	10	11	12	13	14	15
$lo \quad hi$															
$j \quad i$															

```
// check if pointers cross
if (i >= j) break;
```

Partition Example, lo=7, hi=8

A	C	E	E	I	K	L	M	O	P	T	Q	R	X	U	S
0	1	2	3	4	5	6	7	8	9	10	11	12	13	14	15
lo hi															
j i															

```
// put partitioning item v at a[j]
exch(a, lo, j);

// now, a[lo .. j-1] <= a[j] <= a[j+1 .. hi]
return j;
```

Partition Example, $lo=7$, $hi=8$

A	C	E	E	I	K	L	M	O	P	T	Q	R	X	U	S
0	1	2	3	4	5	6	7	8	9	10	11	12	13	14	15
							lo	hi							
							j	i							

Partition for $lo=7$, $hi=8$ is complete. $j=7$

No partition for subarrays of size 1.

Call sort recursively on right subarray with $lo=10$, $hi = 15$

```
// put partitioning item v at a[j]
exch(a, lo, j);

// now, a[lo .. j-1] <= a[j] <= a[j+1 .. hi]
return j;
```

```
private static void sort(Comparable[] a, int lo, int hi) {

    if (hi <= lo) return;
    int j = partition(a, lo, hi);
    sort(a, lo, j-1);
    sort(a, j+1, hi);
}
```

```
private static void sort(Comparable[] a, int lo, int hi) {
    if (hi <= lo) return;
    int j = partition(a, lo, hi);
    sort(a, lo, j-1);
    sort(a, j+1, hi);
}
```

Quicksort Trace

	lo	j	hi	0	1	2	3	4	5	6	7	8	9	10	11	12	13	14	15	
initial values				Q	U	I	C	K	S	O	R	T	E	X	A	M	P	L	E	
random shuffle				K	R	A	T	E	L	E	P	U	I	M	Q	C	X	O	S	
	0	5	15	E	C	A	I	E	K	L	P	U	T	M	Q	R	X	O	S	
	0	3	4	E	C	A	E	I	K	L	P	U	T	M	Q	R	X	O	S	
	0	2	2	A	C	E	E	I	K	L	P	U	T	M	Q	R	X	O	S	
	0	0	1	A	C	E	E	I	K	L	P	U	T	M	Q	R	X	O	S	
	1			1	A	C	E	E	I	K	L	P	U	T	M	Q	R	X	O	S
	4			4	A	C	E	E	I	K	L	P	U	T	M	Q	R	X	O	S
	6	6	15	A	C	E	E	I	K	L	P	U	T	M	Q	R	X	O	S	
	7	9	15	A	C	E	E	I	K	L	M	O	P	T	Q	R	X	U	S	
<i>no partition for subarrays of size 1</i>	7	7	8	A	C	E	E	I	K	L	M	O	P	T	Q	R	X	U	S	
	8			A	C	E	E	I	K	L	M	O	P	T	Q	R	X	U	S	
	10	13	15	A	C	E	E	I	K	L	M	O	P	S	Q	R	T	U	X	
	10	12	12	A	C	E	E	I	K	L	M	O	P	R	Q	S	T	U	X	
	10	11	11	A	C	E	E	I	K	L	M	O	P	Q	R	S	T	U	X	
	10			A	C	E	E	I	K	L	M	O	P	Q	R	S	T	U	X	
	14	14	15	A	C	E	E	I	K	L	M	O	P	Q	R	S	T	U	X	
	15			A	C	E	E	I	K	L	M	O	P	Q	R	S	T	U	X	
result				A	C	E	E	I	K	L	M	O	P	Q	R	S	T	U	X	

Partition Example, lo=10, hi=15

A	C	E	E	I	K	L	M	O	P	T	Q	R	X	U	S
0	1	2	3	4	5	6	7	8	9	10	11	12	13	14	15
										lo					hi

										i					j
--	--	--	--	--	--	--	--	--	--	---	--	--	--	--	---

```
// partition the subarray a[lo..hi] so that a[lo..j-1] <= a[j] <= a[j+1..hi] and return the index j.
private static int partition(Comparable[] a, int lo, int hi) {
    int i = lo;
    int j = hi + 1;
    Comparable v = a[lo];
```

Partition Example, $lo=10$, $hi=15$

A	C	E	E	I	K	L	M	O	P	T	Q	R	X	U	S
0	1	2	3	4	5	6	7	8	9	10	11	12	13	14	15
										lo					hi
										i					j

```
// find greater than (or equal to) v to swap
while (less(a[++i], v)) {
    if (i == hi) break;
}
```

Partition Example, lo=10, hi=15

A	C	E	E	I	K	L	M	O	P	T	Q	R	X	U	S
0	1	2	3	4	5	6	7	8	9	10	11	12	13	14	15
										lo					hi

											i				j
--	--	--	--	--	--	--	--	--	--	--	---	--	--	--	---

```
// find greater than (or equal to) v to swap
while (less(a[++i], v)) {
    if (i == hi) break;
}
```

Partition Example, $lo=10$, $hi=15$

A	C	E	E	I	K	L	M	O	P	T	Q	R	X	U	S
0	1	2	3	4	5	6	7	8	9	10	11	12	13	14	15
										lo					hi
											i				j

```
// find greater than (or equal to) v to swap
while (less(a[++i], v)) {
    if (i == hi) break;
}
```

Partition Example, $lo=10$, $hi=15$

A	C	E	E	I	K	L	M	O	P	T	Q	R	X	U	S
0	1	2	3	4	5	6	7	8	9	10	11	12	13	14	15
lo														hi	
i														j	

```
// find item smaller than (or equal to) v to swap
while (less(v, a[--j])) {
    if (j == lo) break; // redundant since a[lo] acts as sentinel
}
```

Partition Example, $lo=10$, $hi=15$

A	C	E	E	I	K	L	M	O	P	T	Q	R	X	U	S
0	1	2	3	4	5	6	7	8	9	10	11	12	13	14	15
lo														hi	
													i	j	

```
exch(a, i, j);
```

Partition Example, $lo=10$, $hi=15$

A	C	E	E	I	K	L	M	O	P	T	Q	R	S	U	X
0	1	2	3	4	5	6	7	8	9	10	11	12	13	14	15
lo														hi	
														i	j

```
exch(a, i, j);
```

Partition Example, $lo=10$, $hi=15$

A	C	E	E	I	K	L	M	O	P	T	Q	R	S	U	X
0	1	2	3	4	5	6	7	8	9	10	11	12	13	14	15
lo														hi	
														i	j

```
// find greater than (or equal to) v to swap
while (less(a[++i], v)) {
    if (i == hi) break;
}
```

Partition Example, $lo=10$, $hi=15$

A	C	E	E	I	K	L	M	O	P	T	Q	R	S	U	X
0	1	2	3	4	5	6	7	8	9	10	11	12	13	14	15
lo														hi	
														i,j	

```
// find item smaller than (or equal to) v to swap
while (less(v, a[--j])) {
    if (j == lo) break; // redundant since a[lo] acts as sentinel
}
```

Partition Example, $\text{lo}=10, \text{hi}=15$

A	C	E	E	I	K	L	M	O	P	T	Q	R	S	U	X
0	1	2	3	4	5	6	7	8	9	10	11	12	13	14	15
lo														hi	

```
// find item smaller than (or equal to) v to swap
while (less(v, a[--j])) {
    if (j == lo) break; // redundant since a[lo] acts as sentinel
}
```

Partition Example, lo=10, hi=15

A	C	E	E	I	K	L	M	O	P	T	Q	R	S	U	X
0	1	2	3	4	5	6	7	8	9	10	11	12	13	14	15
lo														hi	
														j	i

```
// check if pointers cross
if (i >= j) break;
```

Partition Example, $lo=10$, $hi=15$

A	C	E	E	I	K	L	M	O	P	T	Q	R	S	U	X
0	1	2	3	4	5	6	7	8	9	10	11	12	13	14	15
lo														hi	
														j	i

```
// put partitioning item v at a[j]
exch(a, lo, j);

// now, a[lo .. j-1] <= a[j] <= a[j+1 .. hi]
return j;
```

Partition Example, $lo=10$, $hi=15$

A	C	E	E	I	K	L	M	O	P	S	Q	R	T	U	X
0	1	2	3	4	5	6	7	8	9	10	11	12	13	14	15
										lo					hi

j i

```
// put partitioning item v at a[j]
exch(a, lo, j);

// now, a[lo .. j-1] <= a[j] <= a[j+1 .. hi]
return j;
```

Partition for $lo=6$, $hi=15$ is complete. $j=13$

Call sort recursively on left subarray with $lo=10$, $hi = 12$

```
private static void sort(Comparable[] a, int lo, int hi) {

    if (hi <= lo) return;
    int j = partition(a, lo, hi);
    sort(a, lo, j-1);
    sort(a, j+1, hi);
}
```

```
private static void sort(Comparable[] a, int lo, int hi) {
    if (hi <= lo) return;
    int j = partition(a, lo, hi);
    sort(a, lo, j-1);
    sort(a, j+1, hi);
}
```

Quicksort Trace

	lo	j	hi	0	1	2	3	4	5	6	7	8	9	10	11	12	13	14	15
initial values				Q	U	I	C	K	S	O	R	T	E	X	A	M	P	L	E
random shuffle				K	R	A	T	E	L	E	P	U	I	M	Q	C	X	O	S
	0	5	15	E	C	A	I	E	K	L	P	U	T	M	Q	R	X	O	S
	0	3	4	E	C	A	E	I	K	L	P	U	T	M	Q	R	X	O	S
	0	2	2	A	C	E	E	I	K	L	P	U	T	M	Q	R	X	O	S
	0	0	1	A	C	E	E	I	K	L	P	U	T	M	Q	R	X	O	S
		1		A	C	E	E	I	K	L	P	U	T	M	Q	R	X	O	S
		4		A	C	E	E	I	K	L	P	U	T	M	Q	R	X	O	S
	6	6	15	A	C	E	E	I	K	L	P	U	T	M	Q	R	X	O	S
	7	9	15	A	C	E	E	I	K	L	M	O	P	T	Q	R	X	U	S
	7	7	8	A	C	E	E	I	K	L	M	O	P	T	Q	R	X	U	S
	8	8	A	C	F	F	T	K	L	M	O	P	T	Q	R	X	U	S	
<i>no partition for subarrays of size 1</i>																			
	10	13	15	A	C	E	E	I	K	L	M	O	P	S	Q	R	T	U	X
	10	12	12	A	C	E	E	I	K	L	M	O	P	R	Q	S	T	U	X
	10	11	11	A	C	E	E	I	K	L	M	O	P	Q	R	S	T	U	X
	10	10	10	A	C	E	E	I	K	L	M	O	P	Q	R	S	T	U	X
	14	14	15	A	C	E	E	I	K	L	M	O	P	Q	R	S	T	U	X
	15			A	C	E	E	I	K	L	M	O	P	Q	R	S	T	U	X
result				A	C	E	E	I	K	L	M	O	P	Q	R	S	T	U	X

Partition Example, lo=10, hi=12

A	C	E	E	I	K	L	M	O	P	S	Q	R	T	U	X
0	1	2	3	4	5	6	7	8	9	10	11	12	13	14	15
										lo		hi			
										i		j			

```
// partition the subarray a[lo..hi] so that a[lo..j-1] <= a[j] <= a[j+1..hi] and return the index j.
private static int partition(Comparable[] a, int lo, int hi) {
    int i = lo;
    int j = hi + 1;
    Comparable v = a[lo];
```

Partition Example, lo=10, hi=12

A	C	E	E	I	K	L	M	O	P	S	Q	R	T	U	X
0	1	2	3	4	5	6	7	8	9	10	11	12	13	14	15
										lo		hi			
										i		j			

```
// find greater than (or equal to) v to swap
while (less(a[++i], v)) {
    if (i == hi) break;
}
```

Partition Example, lo=10, hi=12

A	C	E	E	I	K	L	M	O	P	S	Q	R	T	U	X
0	1	2	3	4	5	6	7	8	9	10	11	12	13	14	15
										lo	hi			i	j

```
// find greater than (or equal to) v to swap
while (less(a[++i], v)) {
    if (i == hi) break;
}
```

Partition Example, $lo=10$, $hi=12$

A	C	E	E	I	K	L	M	O	P	S	Q	R	T	U	X
0	1	2	3	4	5	6	7	8	9	10	11	12	13	14	15
										lo		hi			
											i,j				

```
// find item smaller than (or equal to) v to swap
while (less(v, a[--j])) {
    if (j == lo) break; // redundant since a[lo] acts as sentinel
}
```

Partition Example, lo=10, hi=12

A	C	E	E	I	K	L	M	O	P	S	Q	R	T	U	X
0	1	2	3	4	5	6	7	8	9	10	11	12	13	14	15
										lo	hi				

 i,j

```
// check if pointers cross
if (i >= j) break;
```

Partition Example, lo=10, hi=12

A	C	E	E	I	K	L	M	O	P	S	Q	R	T	U	X
0	1	2	3	4	5	6	7	8	9	10	11	12	13	14	15
										lo		hi			
											i,j				

```
// put partitioning item v at a[j]
exch(a, lo, j);

// now, a[lo .. j-1] <= a[j] <= a[j+1 .. hi]
return j;
```

Partition Example, $lo=10$, $hi=12$

A	C	E	E	I	K	L	M	O	P	R	Q	S	T	U	X
0	1	2	3	4	5	6	7	8	9	10	11	12	13	14	15
										lo		hi			

i,j

```
// put partitioning item v at a[j]
exch(a, lo, j);

// now, a[lo .. j-1] <= a[j] <= a[j+1 .. hi]
return j;
```

Partition for $lo=10$, $hi=12$ is complete. $j=12$

Call sort recursively on left subarray with $lo=10$, $hi = 11$

```
private static void sort(Comparable[] a, int lo, int hi) {

    if (hi <= lo) return;
    int j = partition(a, lo, hi);
    sort(a, lo, j-1);
    sort(a, j+1, hi);
}
```

```
private static void sort(Comparable[] a, int lo, int hi) {
    if (hi <= lo) return;
    int j = partition(a, lo, hi);
    sort(a, lo, j-1);
    sort(a, j+1, hi);
}
```

Quicksort Trace

	lo	j	hi	0	1	2	3	4	5	6	7	8	9	10	11	12	13	14	15	
initial values				Q	U	I	C	K	S	O	R	T	E	X	A	M	P	L	E	
random shuffle				K	R	A	T	E	L	E	P	U	I	M	Q	C	X	O	S	
	0	5	15	E	C	A	I	E	K	L	P	U	T	M	Q	R	X	O	S	
	0	3	4	E	C	A	E	I	K	L	P	U	T	M	Q	R	X	O	S	
	0	2	2	A	C	E	E	I	K	L	P	U	T	M	Q	R	X	O	S	
	0	0	1	A	C	E	E	I	K	L	P	U	T	M	Q	R	X	O	S	
	1			1	A	C	E	E	I	K	L	P	U	T	M	Q	R	X	O	S
	4			4	A	C	E	E	I	K	L	P	U	T	M	Q	R	X	O	S
	6	6	15	A	C	E	E	I	K	L	P	U	T	M	Q	R	X	O	S	
	7	9	15	A	C	E	E	I	K	L	M	O	P	T	Q	R	X	U	S	
	7	7	8	A	C	E	E	I	K	L	M	O	P	T	Q	R	X	U	S	
	8			8	A	C	E	E	I	K	L	M	O	P	T	Q	R	X	U	S
	10	13	15	A	C	F	F	T	K	L	M	O	P	S	O	R	T	U	X	
	10	12	12	A	C	E	E	I	K	L	M	O	P	R	Q	S	T	U	X	
	10	11	11	A	C	E	E	I	K	L	M	O	P	Q	R	S	T	U	X	
	10			10	A	C	E	E	I	K	L	M	O	P	Q	R	S	T	U	X
	14	14	15	A	C	E	E	I	K	L	M	O	P	Q	R	S	T	U	X	
	15			15	A	C	E	E	I	K	L	M	O	P	Q	R	S	T	U	X
result				A	C	E	E	I	K	L	M	O	P	Q	R	S	T	U	X	

no partition
for subarrays
of size 1



Partition Example, lo=10, hi=11

A	C	E	E	I	K	L	M	O	P	R	Q	S	T	U	X
0	1	2	3	4	5	6	7	8	9	10	11	12	13	14	15
lo hi															
										i	j				

```
// partition the subarray a[lo..hi] so that a[lo..j-1] <= a[j] <= a[j+1..hi] and return the index j.
private static int partition(Comparable[] a, int lo, int hi) {
    int i = lo;
    int j = hi + 1;
    Comparable v = a[lo];
```

Partition Example, lo=10, hi=11

A	C	E	E	I	K	L	M	O	P	R	Q	S	T	U	X
0	1	2	3	4	5	6	7	8	9	10	11	12	13	14	15
lo hi															
										i	j				

```
// find greater than (or equal to) v to swap
while (less(a[++i], v)) {
    if (i == hi) break;
}
```

Partition Example, $lo=10$, $hi=11$

A	C	E	E	I	K	L	M	O	P	R	Q	S	T	U	X
0	1	2	3	4	5	6	7	8	9	10	11	12	13	14	15
lo hi															
i,j															

```
// find item smaller than (or equal to) v to swap
while (less(v, a[--j])) {
    if (j == lo) break; // redundant since a[lo] acts as sentinel
}
```

Partition Example, $lo=10$, $hi=11$

A	C	E	E	I	K	L	M	O	P	R	Q	S	T	U	X
0	1	2	3	4	5	6	7	8	9	10	11	12	13	14	15
lo														hi	
i, j															

```
// check if pointers cross
if (i >= j) break;
```

Partition Example, lo=10, hi=11

A	C	E	E	I	K	L	M	O	P	R	Q	S	T	U	X
0	1	2	3	4	5	6	7	8	9	10	11	12	13	14	15
lo hi															
i,j															

```
// put partitioning item v at a[j]
exch(a, lo, j);

// now, a[lo .. j-1] <= a[j] <= a[j+1 .. hi]
return j;
```

Partition Example, lo=10, hi=11

A	C	E	E	I	K	L	M	O	P	Q	R	S	T	U	X
0	1	2	3	4	5	6	7	8	9	10	11	12	13	14	15
										lo	hi				
										ij					

Partition for lo=10, hi=11 is complete. j=11

No partition for subarrays of size 1.

Call sort recursively on right subarray with lo=14, hi = 15

private static void sort(Comparable[] a, int lo, int hi) {

```
if (hi <= lo) return;
int j = partition(a, lo, hi);
sort(a, lo, j-1);
sort(a, j+1, hi);
```

}

```
// put partitioning item v at a[j]
exch(a, lo, j);

// now, a[lo .. j-1] <= a[j] <= a[j+1 .. hi]
return j;
```

```
private static void sort(Comparable[] a, int lo, int hi) {
    if (hi <= lo) return;
    int j = partition(a, lo, hi);
    sort(a, lo, j-1);
    sort(a, j+1, hi);
}
```

Quicksort Trace

	lo	j	hi	0	1	2	3	4	5	6	7	8	9	10	11	12	13	14	15
initial values				Q	U	I	C	K	S	O	R	T	E	X	A	M	P	L	E
random shuffle				K	R	A	T	E	L	E	P	U	I	M	Q	C	X	O	S
	0	5	15	E	C	A	I	E	K	L	P	U	T	M	Q	R	X	O	S
	0	3	4	E	C	A	E	I	K	L	P	U	T	M	Q	R	X	O	S
	0	2	2	A	C	E	E	I	K	L	P	U	T	M	Q	R	X	O	S
	0	0	1	A	C	E	E	I	K	L	P	U	T	M	Q	R	X	O	S
	1			A	C	E	E	I	K	L	P	U	T	M	Q	R	X	O	S
	4			A	C	E	E	I	K	L	P	U	T	M	Q	R	X	O	S
	6	6	15	A	C	E	E	I	K	L	P	U	T	M	Q	R	X	O	S
	7	9	15	A	C	E	E	I	K	L	M	O	P	T	Q	R	X	U	S
	7	7	8	A	C	E	E	I	K	L	M	O	P	T	Q	R	X	U	S
	8			A	C	E	E	I	K	L	M	O	P	T	Q	R	X	U	S
	10	13	15	A	C	E	E	I	K	L	M	O	P	S	Q	R	T	U	X
	10	12	12	A	C	F	F	T	K	I	M	O	P	R	O	S	T	U	X
	10	11	11	A	C	E	E	I	K	L	M	O	P	Q	R	S	T	U	X
	10			A	C	E	E	I	K	L	M	O	P	Q	R	S	T	U	X
	14	14	15	A	C	E	E	I	K	L	M	O	P	Q	R	S	T	U	X
	15			A	C	E	E	I	K	L	M	O	P	Q	R	S	T	U	X
result				A	C	E	E	I	K	L	M	O	P	Q	R	S	T	U	X

Partition Example, lo=14, hi=15

A	C	E	E	I	K	L	M	O	P	Q	R	S	T	U	X
0	1	2	3	4	5	6	7	8	9	10	11	12	13	14	15
														lo	hi
														i	j

```
// partition the subarray a[lo..hi] so that a[lo..j-1] <= a[j] <= a[j+1..hi] and return the index j.
private static int partition(Comparable[] a, int lo, int hi) {
    int i = lo;
    int j = hi + 1;
    Comparable v = a[lo];
```

Partition Example, $lo=14$, $hi=15$

A	C	E	E	I	K	L	M	O	P	Q	R	S	T	U	X	
0	1	2	3	4	5	6	7	8	9	10	11	12	13	14	15	
															lo	hi
															i	j

```
// find greater than (or equal to) v to swap
while (less(a[++i], v)) {
    if (i == hi) break;
}
```

Partition Example, $\text{lo}=14$, $\text{hi}=15$

A	C	E	E	I	K	L	M	O	P	Q	R	S	T	U	X
0	1	2	3	4	5	6	7	8	9	10	11	12	13	14	15
lo hi														i,j	

```
// find item smaller than (or equal to) v to swap
while (less(v, a[--j])) {
    if (j == lo) break; // redundant since a[lo] acts as sentinel
}
```

Partition Example, $\text{lo}=14$, $\text{hi}=15$

A	C	E	E	I	K	L	M	O	P	Q	R	S	T	U	X	
0	1	2	3	4	5	6	7	8	9	10	11	12	13	14	15	
															lo	hi
															j	i

```
// find item smaller than (or equal to) v to swap
while (less(v, a[--j])) {
    if (j == lo) break; // redundant since a[lo] acts as sentinel
}
```

Partition Example, $\text{lo}=14$, $\text{hi}=15$

A	C	E	E	I	K	L	M	O	P	Q	R	S	T	U	X	
0	1	2	3	4	5	6	7	8	9	10	11	12	13	14	15	
															lo	hi
															j	i

```
// put partitioning item v at a[j]
exch(a, lo, j);

// now, a[lo .. j-1] <= a[j] <= a[j+1 .. hi]
return j;
```

Partition Example, $lo=14$, $hi=15$

A	C	E	E	I	K	L	M	O	P	Q	R	S	T	U	X
0	1	2	3	4	5	6	7	8	9	10	11	12	13	14	15
														lo	hi

~~j i~~

Partition for $lo=14$, $hi=15$ is complete. $j=14$

No partition for subarrays of size 1.

Sorting complete

```
// put partitioning item v at a[j]
exch(a, lo, j);

// now, a[lo .. j-1] <= a[j] <= a[j+1 .. hi]
return j;
```

```
private static void sort(Comparable[] a, int lo, int hi) {

    if (hi <= lo) return;
    int j = partition(a, lo, hi);
    sort(a, lo, j-1);
    sort(a, j+1, hi);
}
```

```
private static void sort(Comparable[] a, int lo, int hi) {
    if (hi <= lo) return;
    int j = partition(a, lo, hi);
    sort(a, lo, j-1);
    sort(a, j+1, hi);
}
```

Quicksort Trace

	lo	j	hi	0	1	2	3	4	5	6	7	8	9	10	11	12	13	14	15
initial values				Q	U	I	C	K	S	O	R	T	E	X	A	M	P	L	E
random shuffle				K	R	A	T	E	L	E	P	U	I	M	Q	C	X	O	S
	0	5	15	E	C	A	I	E	K	L	P	U	T	M	Q	R	X	O	S
	0	3	4	E	C	A	E	I	K	L	P	U	T	M	Q	R	X	O	S
	0	2	2	A	C	E	E	I	K	L	P	U	T	M	Q	R	X	O	S
	0	0	1	A	C	E	E	I	K	L	P	U	T	M	Q	R	X	O	S
	1			A	C	E	E	I	K	L	P	U	T	M	Q	R	X	O	S
	4			A	C	E	E	I	K	L	P	U	T	M	Q	R	X	O	S
	6	6	15	A	C	E	E	I	K	L	P	U	T	M	Q	R	X	O	S
	7	9	15	A	C	E	E	I	K	L	M	O	P	T	Q	R	X	U	S
	7	7	8	A	C	E	E	I	K	L	M	O	P	T	Q	R	X	U	S
	8			A	C	E	E	I	K	L	M	O	P	T	Q	R	X	U	S
	10	13	15	A	C	E	E	I	K	L	M	O	P	S	Q	R	T	U	X
	10	12	12	A	C	E	E	I	K	L	M	O	P	R	Q	S	T	U	X
	10	11	11	A	C	E	E	I	K	L	M	O	P	Q	R	S	T	U	X
	10			A	C	E	E	I	K	L	M	O	P	Q	R	S	T	U	X
	14	14	15	A	C	E	E	I	K	L	M	O	P	Q	R	S	T	U	X
	15			A	C	E	E	I	K	L	M	O	P	Q	R	S	T	U	X
result				A	C	E	E	I	K	L	M	O	P	Q	R	S	T	U	X

Great algorithms are better than good ones

- ▶ Your laptop executes 10^8 comparisons per second
- ▶ A supercomputer executes 10^{12} comparisons per second

	Insertion sort			Mergesort			Quicksort		
Computer	Thousands of inputs	Millions of inputs	Billions of inputs	Thousands of inputs	Millions of inputs	Billions of inputs	Thousands of inputs	Millions of inputs	Billions of inputs
Home	Instant	2 hours	300 years	instant	1 sec	15 min	Instant	0.5 sec	10 min
Supercomputer	Instant	1 second	1 week	instant	instant	instant	instant	instant	Instant

Quicksort analysis: best case

- ▶ Quicksort divides everything exactly in half.
- ▶ Similar to merge sort.
- ▶ Number of compares is $\sim n \log n$.

Quicksort analysis: worst case

- ▶ Data are already sorted or we pick the smallest or largest key as pivot.
- ▶ Number of compares is $\sim n^2$ - quadratic!
- ▶ Extremely unlikely (less likely than the probably that your computer is struck by lightning) if we shuffle and our shuffling is not broken.

Things to remember about quick sort

- ▶ 39% more compares than merge sort but in practice it is faster because it does not move data much.
 - ▶ If good implementation, even in sorted arrays it can be linearithmic. If not, we end up with quadratic.
- ▶ $O(n \log n)$ average, $O(n^2)$ worst, in practice faster than mergesort.
- ▶ **In-place** sorting.
- ▶ **Not stable.**

Quicksort practical improvements

- ▶ Use insertion sort for small subarrays.
- ▶ Best choice of pivot is the median of a small sample.
- ▶ For years, Java used quicksort for collections of primitives and mergesort for collections of objects due to stability.
- ▶ Has moved to dual-pivot quick sort (Yaroslavskiy, Bentley, and Bloch, 2009) and timsort (Peters, 1993), respectively.

Sorting: the story so far

Which Sort	In place	Stable	Best	Average	Worst	Remarks
Selection	X		$O(n^2)$	$O(n^2)$	$O(n^2)$	n exchanges
Insertion	X	X	$O(n)$	$O(n^2)$	$O(n^2)$	Use for small arrays or partially ordered
Merge		X	$O(n \log n)$	$O(n \log n)$	$O(n \log n)$	Guaranteed performance; stable
Quick	X		$O(n \log n)$	$O(n \log n)$	$O(n^2)$	$n \log n$ probabilistic guarantee; fastest in practice

Sorting based on comparisons

- ▶ All sorting algorithms we have seen so far and we will see in this class are compare-based.
- ▶ No compare-based sorting algorithm can sort n elements in less than $O(n \log n)$ time in the worst case.
 - ▶ Proof and proper notation in CS140.

Lecture 14: Quicksort

- ▶ This week's assignment
- ▶ Quicksort

Readings:

- ▶ Textbook:
 - ▶ Chapter 2.3 (Pages 288-296)
- ▶ Website:
 - ▶ Quicksort: <https://algs4.cs.princeton.edu/23quicksort/>
 - ▶ Code: <https://algs4.cs.princeton.edu/23quicksort/Quick.java.html>

Practice Problems:

- ▶ 2.3.1-2.3.4