

# BALANCED SEARCH TREES

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CS 62 – Spring 2021

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## Admin

Last day for “normal” mentor hours, Friday (5/7)

More on mentor hours next week

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## Binary Search Trees

BST – A binary tree where each node has a key, and every node's key is:

- Larger than all keys in its left subtree. (everything left is smaller)
- Smaller than all keys in its right subtree. (everything right is larger)

```

graph TD
    12((12)) --- 8((8))
    12 --- 14((14))
    8 --- 5((5))
    8 --- 9((9))
    14 --- 20((20))
  
```

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## Operations

Search – Does the key exist in the tree

Insert – Insert the key into tree

Delete – Delete the key from the tree

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## Height of the tree

Most of the operations take time  $O(\text{height})$

We said trees built from random data have height  $O(\log n)$ , which is asymptotically tight

Two problems:

- We can't always insure random data
- What happens when we delete nodes and insert others after building a tree?

Worst case height for binary search trees is  $O(n)$  ☹

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## Balanced trees

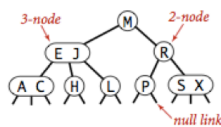
Make sure that the trees remain balanced!

- Red-black trees
- AVL trees
- 2-3 trees
- 2-3-4 trees
- B-trees
- ...

Height is guaranteed to be  $O(\log n)$

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## 2-3 trees



Anatomy of a 2-3 search tree

2-node: one key and two children (left and right)

- everything in left is smaller than key
- everything right is greater than (or equal to) key

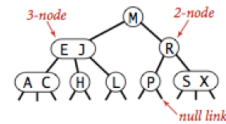
3-node: two keys ( $k_1, k_2$ ) and three children, left, middle and right

- $k_1 < k_2$
- everything in left is less than  $k_1$
- everything in middle is between  $k_1$  and  $k_2$  (greater than or equal to  $k_1$  and less than  $k_2$ )
- everything in right is greater than (or equal to)  $k_2$

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## Search

How do we search for a key?



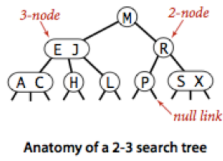
Anatomy of a 2-3 search tree

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### Search

Almost identical to BST search

Only difference: sometimes we have two keys

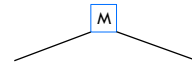


Anatomy of a 2-3 search tree

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### Search

Search(H)

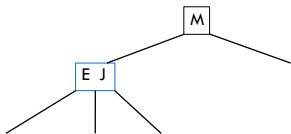


Which child?

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### Search

Search(H)

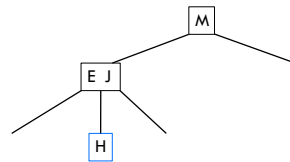


Which child?

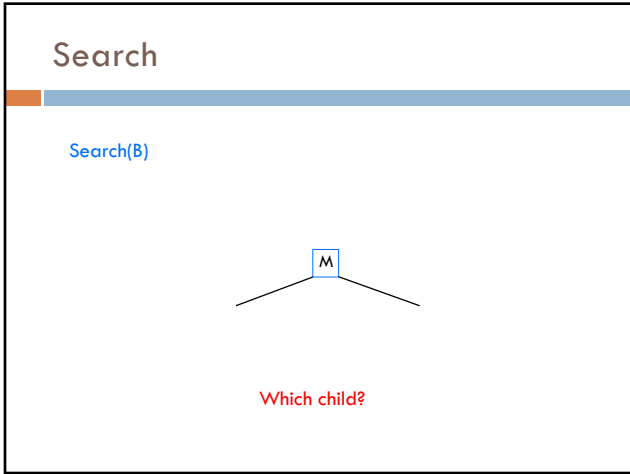
11

### Search

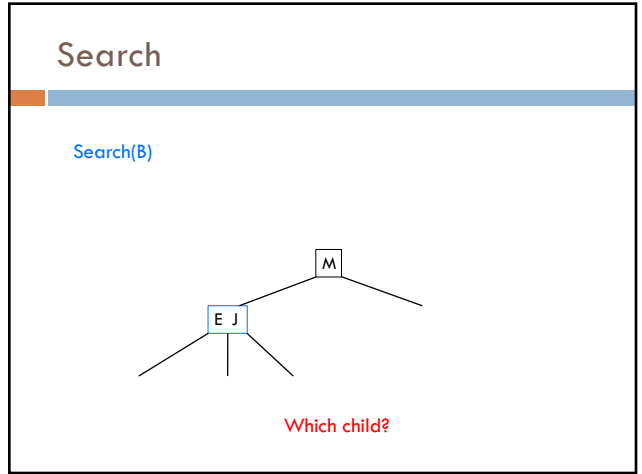
Search(H)



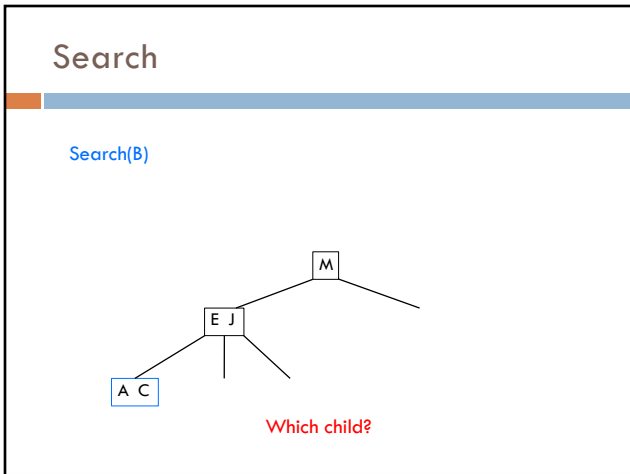
12



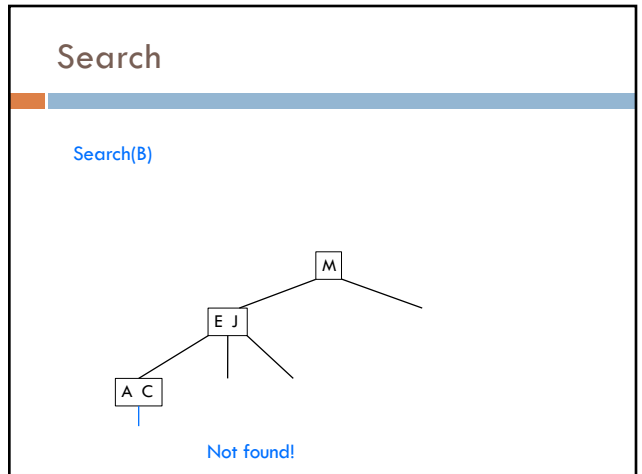
13



14

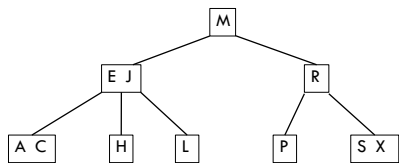


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## Search



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## Insertion

Like BST, insert always happens at a leaf

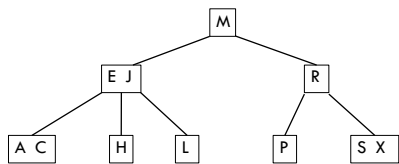
If the leaf is a 2-node, just insert it directly

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## Insertion

If the leaf is a 2-node, just insert it directly

Insert(F)



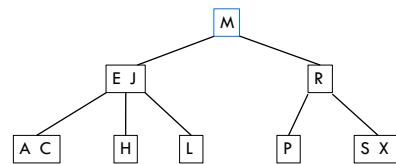
Where should it go?

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## Insertion

If the leaf is a 2-node, just insert it directly

Insert(F)

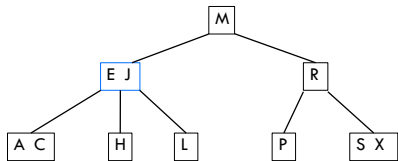


20

## Insertion

If the leaf is a 2-node, just insert it directly

Insert(F)

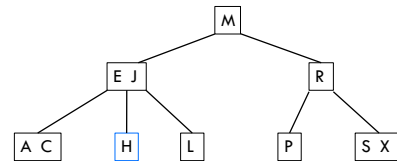


21

## Insertion

If the leaf is a 2-node, just insert it directly

Insert(F)

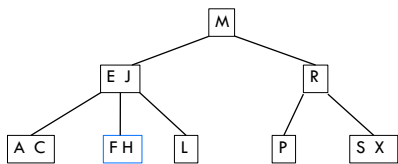


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## Insertion

If the leaf is a 2-node, just insert it directly

Insert(F)



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## Insertion

Like BST, insert always happens at a leaf

If the leaf is a 2-node, just insert it directly

If the leaf is a 3-node:

- We now have three values at this leaf
- Send the middle value up a node
- Make new 2-nodes out of the smallest and largest

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### Insertion

If the leaf is a 3-node:

- We now have three values at this leaf
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Insert(T)

Where should it go?

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### Insertion

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Insert(T)

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Insert(T)

27

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Insert(T)

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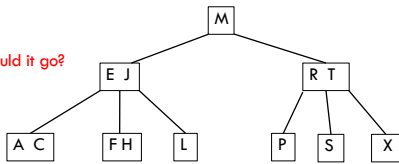
## Insertion

If the leaf is a 3-node:

- We now have three values at this leaf
- Send the middle value up a node
- Make new 2-nodes out of the smallest and largest

Insert(I)

Where should it go?



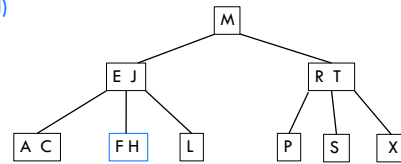
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## Insertion

If the leaf is a 3-node:

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Insert(I)



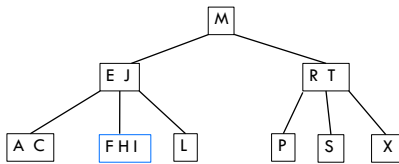
30

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Insert(I)



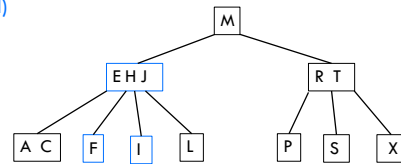
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## Insertion

If the leaf is a 3-node:

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Insert(I)



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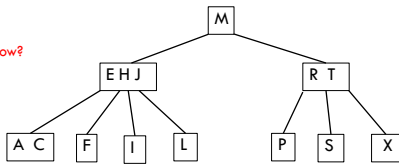
## Insertion

If the leaf is a 3-node:

- We now have three values at this leaf
- Send the middle value up a node
- Make new 2-nodes out of the smallest and largest

Insert(I)

What now?



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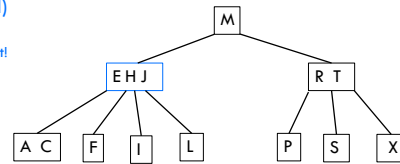
## Insertion

If the leaf is a 3-node:

- We now have three values at this leaf
- Send the middle value up a node
- Make new 2-nodes out of the smallest and largest

Insert(I)

Repeat!



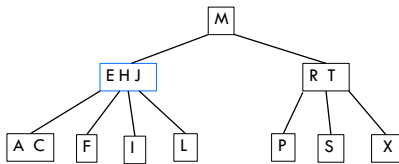
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## Insertion

If the leaf is a 3-node:

- We now have three values at this leaf
- Send the middle value up a node
- Make new 2-nodes out of the smallest and largest

Insert(I)



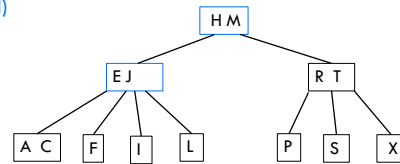
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If the leaf is a 3-node:

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Insert(I)



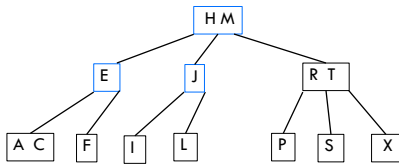
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## Insertion

If the leaf is a 3-node:

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Insert(I)



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## Insertion

If the leaf is a 2-node, just insert it directly

If the leaf is a 3-node:

- We now have three values at this leaf
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When will the height of the tree change?

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## Insertion

If the leaf is a 2-node, just insert it directly

If the leaf is a 3-node:

- We now have three values at this leaf
- Send the middle value up a node
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Only when the root is a 3-node and we insert into a path that is all 3-nodes!

Effect: The tree can hold quite a few values before having to increase the height

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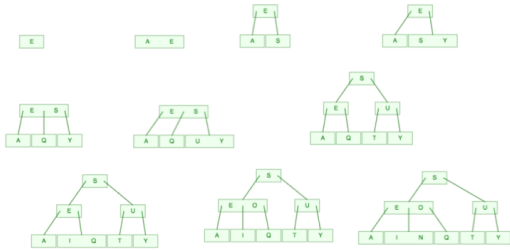
## Practice

Draw the 2-3 tree that results when you insert the keys: EASYQUATION in that order in an initially empty tree.

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## Practice

Draw the 2-3 tree that results when you insert the keys:  
E A S Y Q U T I O N in that order in an initially empty tree.



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## Running time

Worst case height:  $O(\log n)$

What does that mean?

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## Running time

Worst case height:  $O(\log n)$

Insert, search and delete are all  $O(\log n)$

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## 2-3 search trees in practice

A pain to implement

Overhead can often make slower than standard BST

Other balanced trees exist that provide the same worst case guarantee, but are faster (e.g, red-black trees)

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## Red-black tree high-level

<https://www.cs.usfca.edu/~galles/visualization/RedBlack.html>