

GRAPHS: SHORTEST PATHS

David Kauchak
CS 62 – Spring 2021

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Admin

TextGenerator assignment

Lab tomorrow

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Graph code

<https://github.com/pomonacs622021sp/LectureCode/tree/master/GraphExamples>

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Searching on graphs

```

treeSearch( toVisit )
while !toVisit.empty()
  v = toVisit.remove()
  // visit v, e.g., print it out
  for c in v.getChildren()
    toVisit.add(c)

graphSearch( toVisit )
while !toVisit.empty()
  v = toVisit.remove()
  if !visited[v]
    visited[v] = true
    for c in v.getAdjacent()
      if !visited[c]
        toVisit.add(c)
  
```

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Searching on graphs

```

graphBFS( start )
q = new Queue()
q.add(start)
treeSearch(q)

graphDFS( start )
s = new Stack()
s.add(start)
treeSearch(s)

graphSearch( toVisit )
while !toVisit.empty()
v = toVisit.remove()
if !visited[v]
visited[v] = true
for c in v.getAdjacent()
if !visited[c]
toVisit.add(c)
    
```

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DFS

```

graphDFS( start )
s = new Stack()
s.add(start)
treeSearch(s)
    
```

```

toVisit-stack: A
visited:
    
```

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DFS

```

graphSearch( toVisit )
while !toVisit.empty()
v = toVisit.remove()
if !visited[v]
visited[v] = true
for c in v.getAdjacent()
if !visited[c]
toVisit.add(c)
    
```

```

toVisit-stack: A
visited:
    
```

What order will the nodes get printed out?
Assume edges are traversed alphabetically.

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DFS

```

graphSearch( toVisit )
while !toVisit.empty()
v = toVisit.remove()
if !visited[v]
visited[v] = true
for c in v.getAdjacent()
if !visited[c]
toVisit.add(c)
    
```

```

toVisit-stack: A
visited:
    
```

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DFS

```

graphSearch( toVisit )
while !toVisit.empty()
  v = toVisit.remove()
  if !visited[v]
    visited[v] = true
    for c in v.getAdjacent()
      if !visited[c]
        toVisit.add(c)
    
```

toVisit-stack: B D E
visited: A

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DFS

```

graphSearch( toVisit )
while !toVisit.empty()
  v = toVisit.remove()
  if !visited[v]
    visited[v] = true
    for c in v.getAdjacent()
      if !visited[c]
        toVisit.add(c)
    
```

toVisit-stack: B D E
visited: A

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DFS

```

graphSearch( toVisit )
while !toVisit.empty()
  v = toVisit.remove()
  if !visited[v]
    visited[v] = true
    for c in v.getAdjacent()
      if !visited[c]
        toVisit.add(c)
    
```

toVisit-stack: B D C D F
visited: A E

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DFS

```

graphSearch( toVisit )
while !toVisit.empty()
  v = toVisit.remove()
  if !visited[v]
    visited[v] = true
    for c in v.getAdjacent()
      if !visited[c]
        toVisit.add(c)
    
```

toVisit-stack: B D C D F
visited: A E

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DFS

```

graphSearch( toVisit )
while !toVisit.empty()
  v = toVisit.remove()
  if !visited[v]
    visited[v] = true
    for c in v.getAdjacent()
      if !visited[c]
        toVisit.add(c)
  
```

toVisit-stack: B D C D G
visited: A E F

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DFS

```

graphSearch( toVisit )
while !toVisit.empty()
  v = toVisit.remove()
  if !visited[v]
    visited[v] = true
    for c in v.getAdjacent()
      if !visited[c]
        toVisit.add(c)
  
```

toVisit-stack: B D C D G
visited: A E F

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DFS

```

graphSearch( toVisit )
while !toVisit.empty()
  v = toVisit.remove()
  if !visited[v]
    visited[v] = true
    for c in v.getAdjacent()
      if !visited[c]
        toVisit.add(c)
  
```

toVisit-stack: B D C D
visited: A E F G

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DFS

```

graphSearch( toVisit )
while !toVisit.empty()
  v = toVisit.remove()
  if !visited[v]
    visited[v] = true
    for c in v.getAdjacent()
      if !visited[c]
        toVisit.add(c)
  
```

toVisit-stack: B D C D
visited: A E F G

Frontier: Go as far down one path as possible

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DFS

```
graphSearch( toVisit )
while !toVisit.empty()
  v = toVisit.remove()
  if !visited[v]
    visited[v] = true
    for c in v.getAdjacent()
      if !visited[c]
        toVisit.add(c)
```

toVisit-stack: B D C D
visited: A E F G

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DFS

```
graphSearch( toVisit )
while !toVisit.empty()
  v = toVisit.remove()
  if !visited[v]
    visited[v] = true
    for c in v.getAdjacent()
      if !visited[c]
        toVisit.add(c)
```

toVisit-stack: B D C B
visited: A E F G D

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DFS

```
graphSearch( toVisit )
while !toVisit.empty()
  v = toVisit.remove()
  if !visited[v]
    visited[v] = true
    for c in v.getAdjacent()
      if !visited[c]
        toVisit.add(c)
```

toVisit-stack: B D C B
visited: A E F G D

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DFS

```
graphSearch( toVisit )
while !toVisit.empty()
  v = toVisit.remove()
  if !visited[v]
    visited[v] = true
    for c in v.getAdjacent()
      if !visited[c]
        toVisit.add(c)
```

toVisit-stack: B D C
visited: A E F G D B

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DFS

```

graphSearch( toVisit )
while !toVisit.empty()
  v = toVisit.remove()
  if !visited[v]
    visited[v] = true
    for c in v.getAdjacent()
      if !visited[c]
        toVisit.add(c)
  
```

toVisit-stack: B D C
visited: A E F G D B

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DFS

```

graphSearch( toVisit )
while !toVisit.empty()
  v = toVisit.remove()
  if !visited[v]
    visited[v] = true
    for c in v.getAdjacent()
      if !visited[c]
        toVisit.add(c)
  
```

toVisit-stack: B D
visited: A E F G D B C

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DFS

```

graphSearch( toVisit )
while !toVisit.empty()
  v = toVisit.remove()
  if !visited[v]
    visited[v] = true
    for c in v.getAdjacent()
      if !visited[c]
        toVisit.add(c)
  
```

toVisit-stack: B D
visited: A E F G D B C

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DFS

```

graphSearch( toVisit )
while !toVisit.empty()
  v = toVisit.remove()
  if !visited[v]
    visited[v] = true
    for c in v.getAdjacent()
      if !visited[c]
        toVisit.add(c)
  
```

toVisit-stack:
visited: A E F G D B C

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graphSearch run-time

```

graphSearch( toVisit )
while !toVisit.empty()
  v = toVisit.remove()
  if !visited[v]
    visited[v] = true
    for c in v.getAdjacent()
      if !visited[c]
        toVisit.add(c)

```

What is the big-O run-time of graphSearch?

Assume all of the stack/queue operations are constant.

How many times do we visit each vertex?

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graphSearch run-time

```

graphSearch( toVisit )
while !toVisit.empty()
  v = toVisit.remove()
  if !visited[v]
    visited[v] = true
    for c in v.getAdjacent()
      if !visited[c]
        toVisit.add(c)

```

How many times do we visit each vertex? Exactly once

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graphSearch run-time

```

graphSearch( toVisit )
while !toVisit.empty()
  v = toVisit.remove()
  if !visited[v]
    visited[v] = true
    for c in v.getAdjacent()
      if !visited[c]
        toVisit.add(c)

```

What is the cost to traverse all of the edges adjacent to a vertex?

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graphSearch run-time

```

graphSearch( toVisit )
while !toVisit.empty()
  v = toVisit.remove()
  if !visited[v]
    visited[v] = true
    for c in v.getAdjacent()
      if !visited[c]
        toVisit.add(c)

```

What is the cost to traverse all of the edges adjacent to a vertex?

Depends on the graph representation!

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graphSearch run-time

```

graphSearch( toVisit )
while !toVisit.empty()
  v = toVisit.remove()
  if !visited[v]
    visited[v] = true
    for c in v.getAdjacent()
      if !visited[c]
        toVisit.add(c)

```

What is the cost to traverse all of the edges adjacent to a vertex?

Depends on the graph representation!

Adjacency matrix?

Adjacency list?

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graphSearch run-time

```

graphSearch( toVisit )
while !toVisit.empty()
  v = toVisit.remove()
  if !visited[v]
    visited[v] = true
    for c in v.getAdjacent()
      if !visited[c]
        toVisit.add(c)

```

What is the cost to traverse all of the edges adjacent to a vertex?

Depends on the graph representation!

Adjacency matrix? V – we have to traverse the whole row

Adjacency list? Like with trees, we traverse each edge exactly once

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graphSearch run-time

```

graphSearch( toVisit )
while !toVisit.empty()
  v = toVisit.remove()
  if !visited[v]
    visited[v] = true
    for c in v.getAdjacent()
      if !visited[c]
        toVisit.add(c)

```

Overall runtime: Depends on the graph representation!

Adjacency matrix? $O(V^2)$

Adjacency list? $O(V+E)$

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graphSearch run-time

```

graphSearch( toVisit )
while !toVisit.empty()
  v = toVisit.remove()
  if !visited[v]
    visited[v] = true
    for c in v.getAdjacent()
      if !visited[c]
        toVisit.add(c)

```

How many times do we visit each vertex? Exactly once

How many times do we traverse each edge (via the for loop)? Exactly once

What is the big-O run-time of treeSearch? $O(|V| + |E|)$. Linear algorithm.

Nothing changes from treeSearch!
(assuming we're using an adjacency list representation!)

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Search

BFS: breadth first search

- Explores vertices in increasing distance (wrt number of edges) from the starting vertex
- Uses a queue to keep track of vertices to explore

DFS: depth first search:

- Goes as far down a path first and then works its way back
- Two versions: stack and recursive version

Run-time: $O(V + E)$ – assuming adjacency list representation

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BFS and DFS in java

<https://github.com/pomonacs622021/sp/LectureCode/tree/master/GraphExamples>

Look at Graph interface

Look at AdjacencyGraph and MatrixGraph

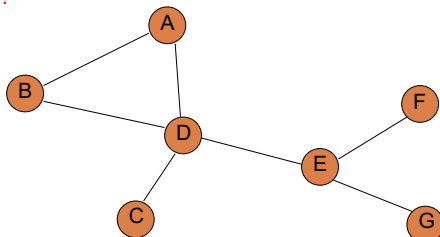
Look at BFS and DFS in Search code

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Connectedness

Connected – every pair of vertices is connected by a path

Algorithm?



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Connectedness

Connected – every pair of vertices is connected by a path

Pick any starting vertex u
Run DFS/BFS from u

Why does this work?

For each vertex v :
if !visited[v]
return false

If we can get from u to every vertex then we know a path exists between all vertices.

Path from a to b : $a - u - b$

return true

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Strongly connected

Strongly connected (directed graphs) –
Every two vertices are reachable by a path

```

    graph TD
      B((B)) --> A((A))
      B((B)) --> D((D))
      C((C)) --> D((D))
      D((D)) --> A((A))
      D((D)) --> E((E))
      E((E)) --> F((F))
      E((E)) --> G((G))
  
```

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Strongly connected

Strongly connected (directed graphs) –
Every two vertices are reachable by a path

Pick any starting vertex u
Run DFS/BFS from u

Does this work?

For each vertex v :
if !visited[v]
return false

return true

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Strongly connected

Strongly connected (directed graphs) –
Every two vertices are reachable by a path

Pick any starting vertex u
Run DFS/BFS from u

Does this work?

No!

For each vertex v :
if !visited[v]
return false

return true

Path from a to b : $a - u - b$

We know we can get from u to b ,
but we don't know that we can get
from a to s (directed graph!)

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Reverse of a graph

Given a graph G , we can calculate the reverse of a graph G^R by reversing the direction of all the edges

```

    graph TD
      subgraph G
        B((B)) --> A((A))
        B((B)) --> D((D))
        C((C)) --> D((D))
        D((D)) --> A((A))
        D((D)) --> E((E))
        E((E)) --> F((F))
        E((E)) --> G((G))
      end
      subgraph GR
        A((A)) --> B((B))
        D((D)) --> B((B))
        D((D)) --> C((C))
        A((A)) --> D((D))
        E((E)) --> D((D))
        F((F)) --> E((E))
        G((G)) --> E((E))
      end
  
```

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Strongly connected

Strongly-Connected(G)

- Run BFS/DFS from some node u
- If not all nodes are visited:
return false
- Create graph G^R
- Run BFS/DFS on G^R from node u
- If not all nodes are visited:
return false
- return true

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Is it correct?

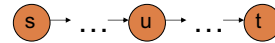
What do we know after the first search?

- Starting at u , we can reach every node

What do we know after the second search (reverse graph)?

- All nodes can reach u . Why?
- We can get from u to every node in G^R , therefore, if we reverse the edges (i.e. G), then we have a path from every node to u

Which means that any node can reach any other node! Given any two nodes s and t we can create a path through u



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Run-times?

Connectedness

Pick any starting vertex u
Run DFS/BFS from u

For each vertex v :

if !visited[v]
return false

return true

What is the run-time?

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