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## Representing graphs

Adjacency list - Each vertex $u \in V$ contains an adjacency list of the set of vertices $v$ such that there exists an edge $(u, v) \in E$

$$
\begin{array}{ll}
\hline \mathrm{A}: & \mathrm{B} \\
\mathrm{~B}: & \mathrm{D} \\
\mathrm{~A} & \mathrm{D} \\
\mathrm{C}: & \mathrm{D} \\
\mathrm{D}: & \mathrm{A} \\
\mathrm{E}: & \mathrm{B} \\
\mathrm{D} & \mathrm{C} \\
\end{array}
$$



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## Representing graphs

Adjacency matrix - $\mathrm{A}|\mathrm{V}| \mathrm{x}|\mathrm{V}|$ matrix A such that:
$a_{i j}= \begin{cases}1 & \text { if }(i, j) \in E \\ 0 & \text { otherwise }\end{cases}$
ABCDE


A 0101010
B 100010
C 0001010
D $\begin{array}{llllll}1 & 1 & 1 & 0 & 1\end{array}$
E 00010


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## Representing graphs

Adjacency matrix - A $|\mathrm{V}| \mathrm{x}|\mathrm{V}|$ matrix A such that:

$$
a_{i j}= \begin{cases}1 & \text { if }(i, j) \in E \\ 0 & \text { otherwise }\end{cases}
$$

ABCDE


| A | 0 | 1 | 0 | 1 | 0 |
| :--- | :--- | :--- | :--- | :--- | :--- |

$\begin{array}{llllll}\text { B } & 1 & 0 & 0 & 1 & 0 \\ \text { C } & 0 & 0 & 0 & 1 & 0\end{array}$
D $111 \begin{array}{llll}1 & 0 & 1\end{array}$
E 00010

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## Representing graphs

Adjacency matrix - $\mathrm{A}|\mathrm{V}| \mathrm{x}|\mathrm{V}|$ matrix A such that:
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| Adjacency list vs. adjacency matrix |  |
| :---: | :---: |
| Adjacency list | Adjacency matrix |
| Sparse graphs (e.g. web) <br> Space efficient <br> Must traverse the adjacency list to discover is an edge exists | Dense graphs <br> Constant time lookup to discover if an edge exists Simple to implement For non-weighted graphs, |

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| Adjacency list vs. adjacency matrix |  |
| :---: | :---: |
| Adjacency list | Adjacency matrix |
| Pros/Cons? |  |



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