

Lecture 11: Singly Linked Lists

CS 62

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Writing code

- No complex code ever works the first time.
- Neither the "If I just fix this last thing..." attitude.
- Think about testing before you write the code.
- Comment before you write the code.
- Never write more than a method without testing it.
- We will talk about JUnits in lab next week.
- Once satisfied, commit your work.

```
public abstract class AbstractList<E>{  
    public AbstractList() { }  
    public boolean isEmpty() { return size() == 0; }  
    public void addFirst(E value) {add(0,value); }  
    public void addLast(E value) {add(size(),value);}  
    public E getFirst() { return get(0);}  
    public E getLast() {return get(size()-1);}  
    public E removeFirst() { return remove(0); }  
    public E removeLast() { return remove(size()-1); }  
    public void add(E value) { addLast(value);}  
    public E remove() { return removeLast();}  
    public E get() { return getLast(); }  
    public boolean contains(E value) {  
        return -1 != indexOf(value); }  
}
```

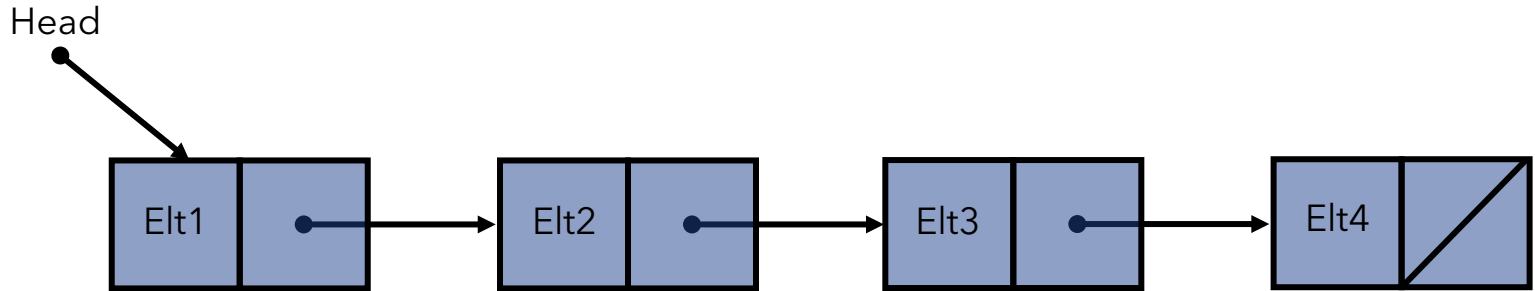
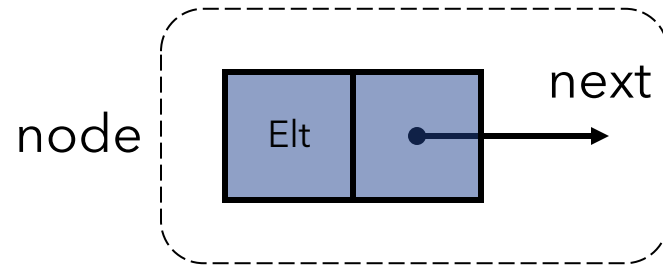
Linked List



- A data structure consisting of a linear sequence of nodes
 - Think of them as snap-lock beads
 - Varying length
- Alternate implementation of an abstract list
 - `Vector` (`ArrayList`) in `structure5` also extends `AbstractList`
- Trade-offs in complexity:
 - In `ArrayList`, adding elements is expensive at beginning of list
 - For Linked lists, it is inexpensive to add elements in early positions
 - However, accessing the i -th element is slow

Singly Linked List

- A linked list consisting of a sequence of nodes, starting from a head pointer
- Each node stores
 - Element
 - Link to the next node



```
public class Node<E>{
    protected E data; // value stored in this element
    protected Node<E> nextElement; // ref to next

    public Node(E v, Node<E> next) {
        data = v;
        nextElement = next; // construct the new head of a singly linked list
    }
    public Node(E v) {
        this(v,null); // constructs a new tail of a list with value v
    }
    public Node<E> next() {
        return nextElement; // returns reference to next value in list
    }
    public void setNext(Node<E> next) {
        nextElement = next; // sets reference to new next value
    }
    public E value() {
        return data; // returns value associated with this element
    }
    public void setValue(E value) {
        data = value; // sets value associated with this element
    }
}
```

SinglyLinkedList

```
public class SinglyLinkedList<E> extends AbstractList<E> {
    protected int count; // list
    protected Node<E> head; // ref. to first element

    //construct an empty list
    public SinglyLinkedList() {
        head = null;
        count = 0;
    }
    public E getFirst() {
        return head.value(); //returns first value in list
    }
    public int size() {
        return count;
    }
}
```

Possible operations

A node can be added/removed:

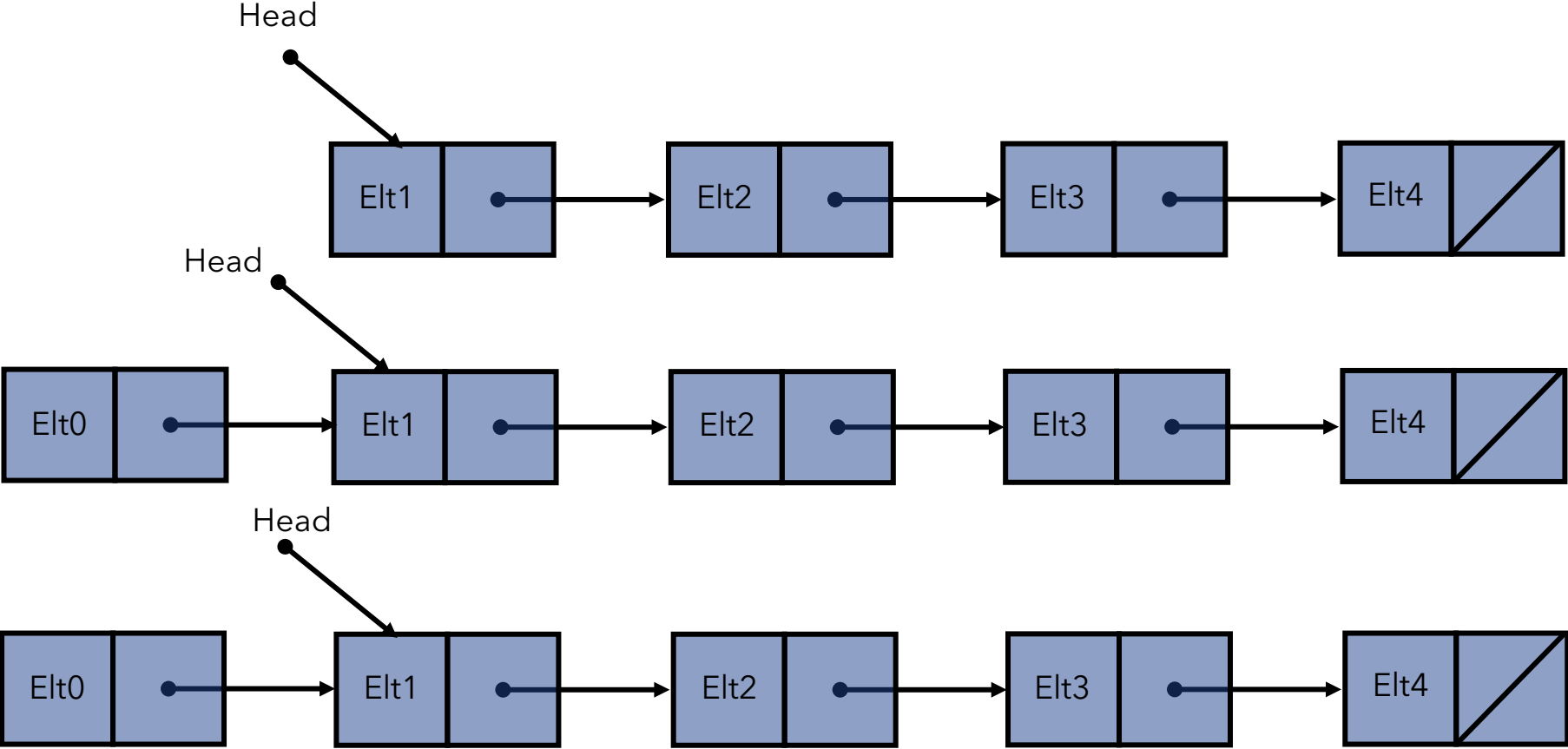
- At the beginning of the list

- At the end of the list

- Between nodes

- In an empty list (addition only)

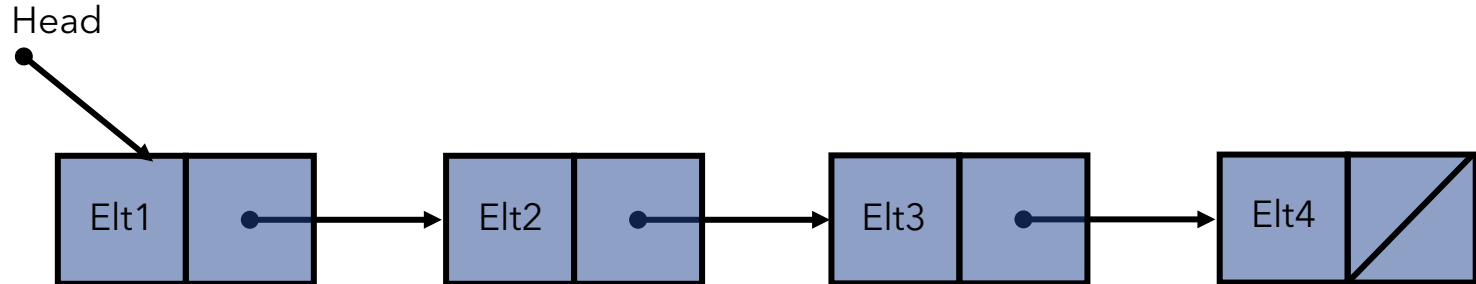
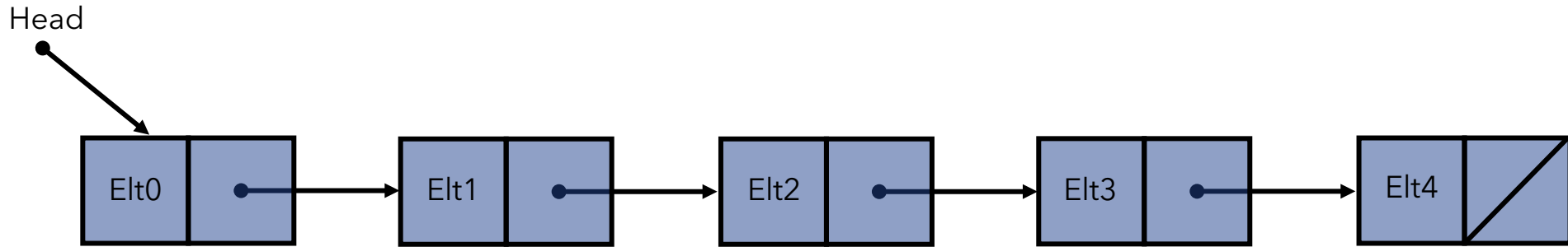
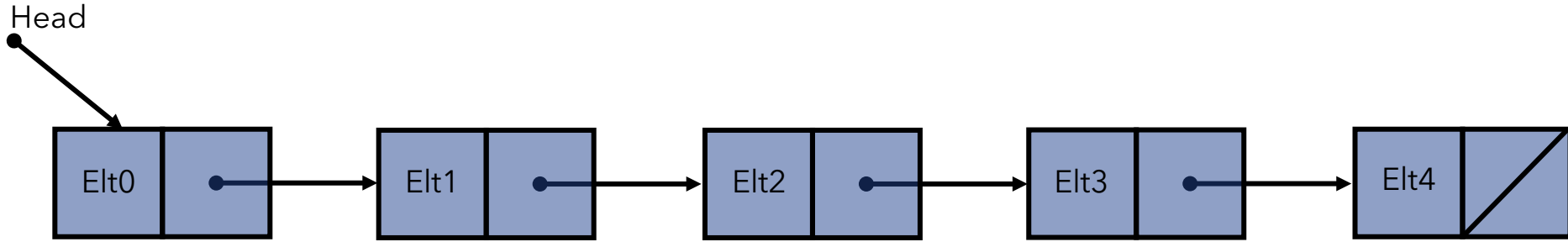
Adding at the head



Adding at the head is $O(1)$

```
public void addFirst(E value){  
    // note order that things happen:  
    // head is parameter, then assigned  
    head = new Node<E>(value, head);  
    count++;  
}
```

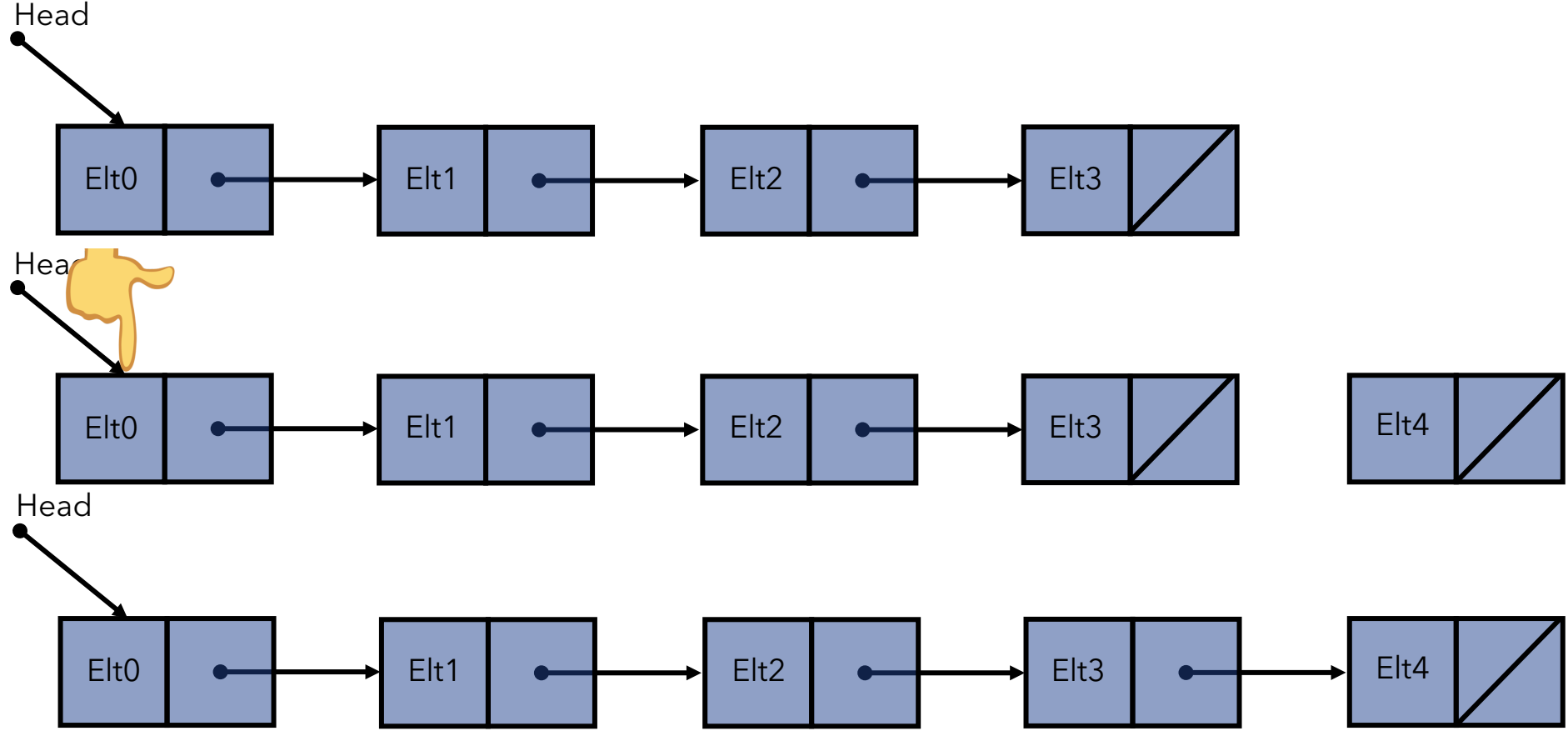
Removing at the head



Removing at the head is $O(1)$

```
public E removeFirst(){
    Node<E> temp = head;
    head = head.next(); // move head down list
    count--;
    return temp.value();
}
```

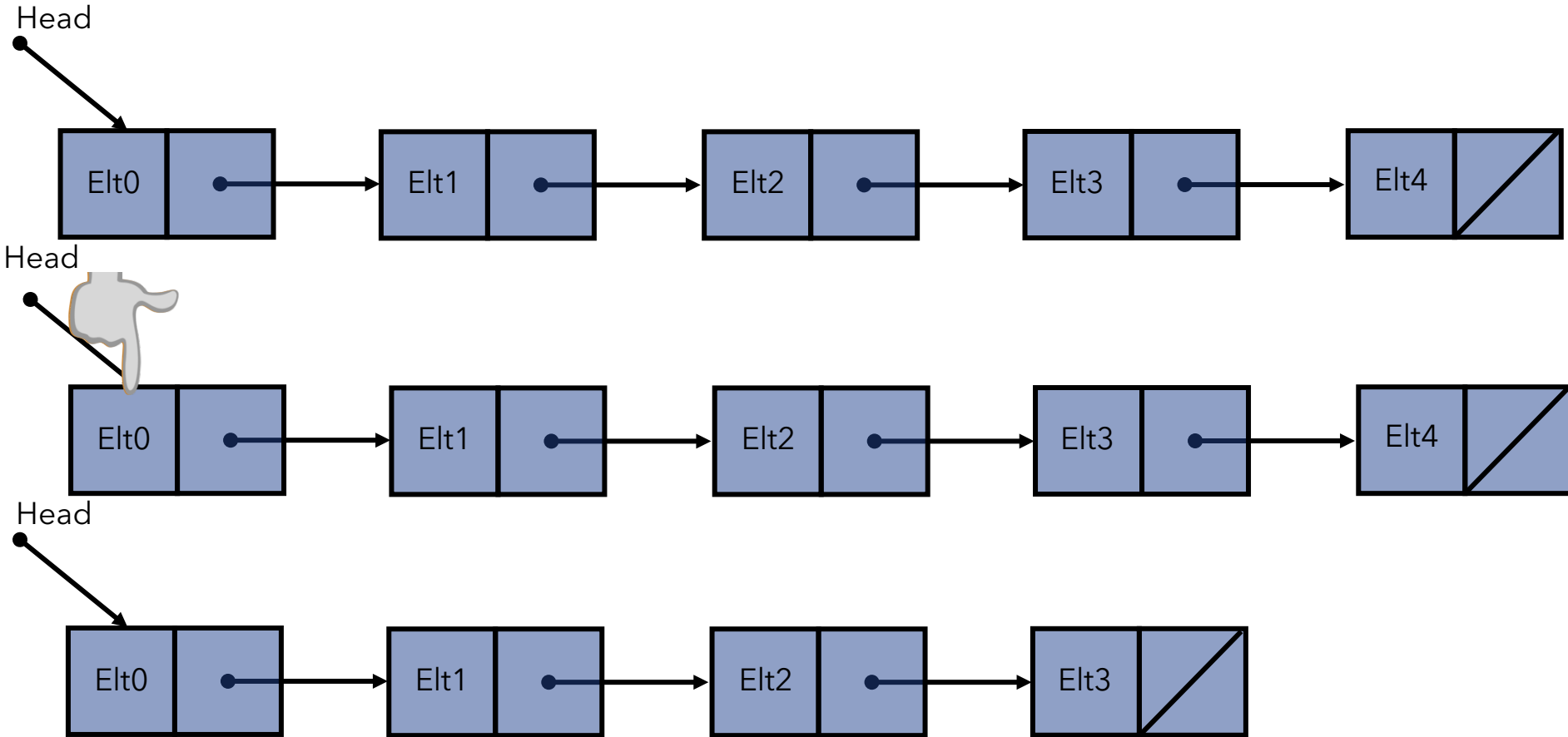
Adding at the end



Adding at the end is $O(n)$

```
public void addLast(E value) {  
    // location for new value  
    Node<E> temp = new Node<E>(value,null);  
    if (head != null) {  
        // pointer to possible tail  
        Node<E> finger = head;  
        while (finger.next() != null) {  
            finger = finger.next();  
        }  
        finger.setNext(temp);  
    }  
    else head = temp; //empty list  
    count++;  
}
```

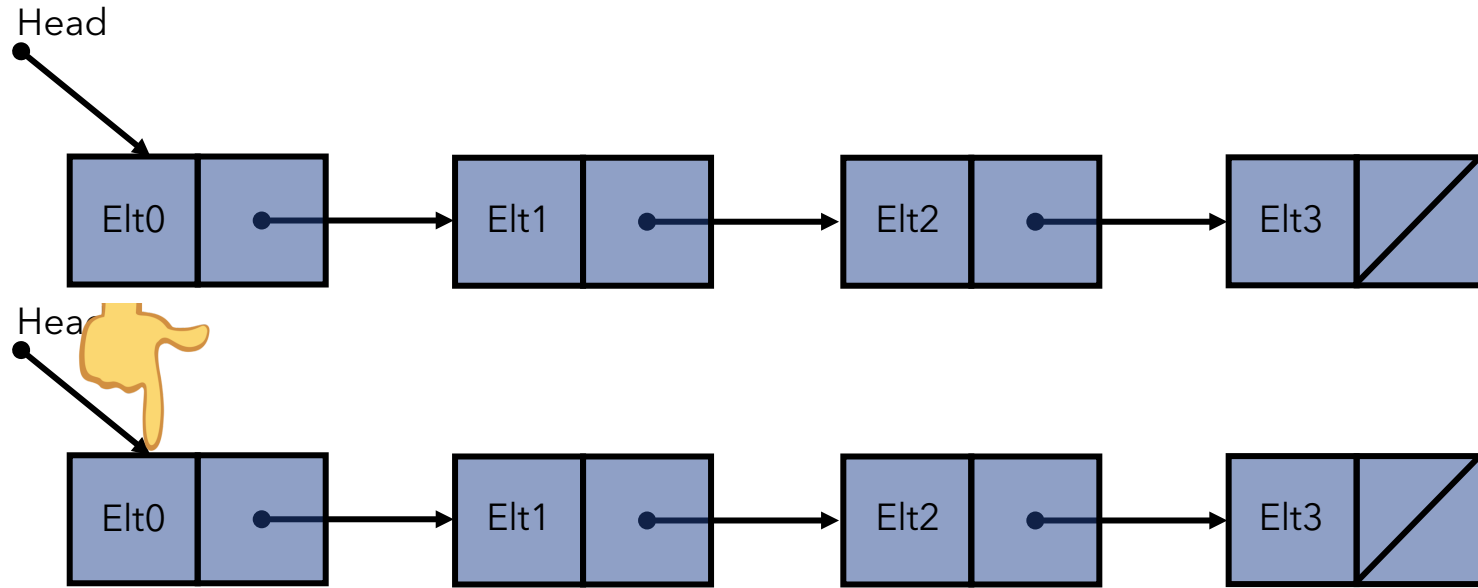
Removing from the end



Removing from the end is $O(n)$

```
public E removeLast() {
    Node<E> finger = head;
    Node<E> previous = null;
    //throw an exception if list is empty
    while (finger.next() != null)
    {
        previous = finger;    // find end of list
        finger = finger.next();
    }
    // finger is null, or points to end of list
    if (previous == null) {
        head = null;    // has exactly one element
    } else{
        previous.setNext(null); // pointer to last element is reset
    }
    count--;
    return finger.value();
}
```

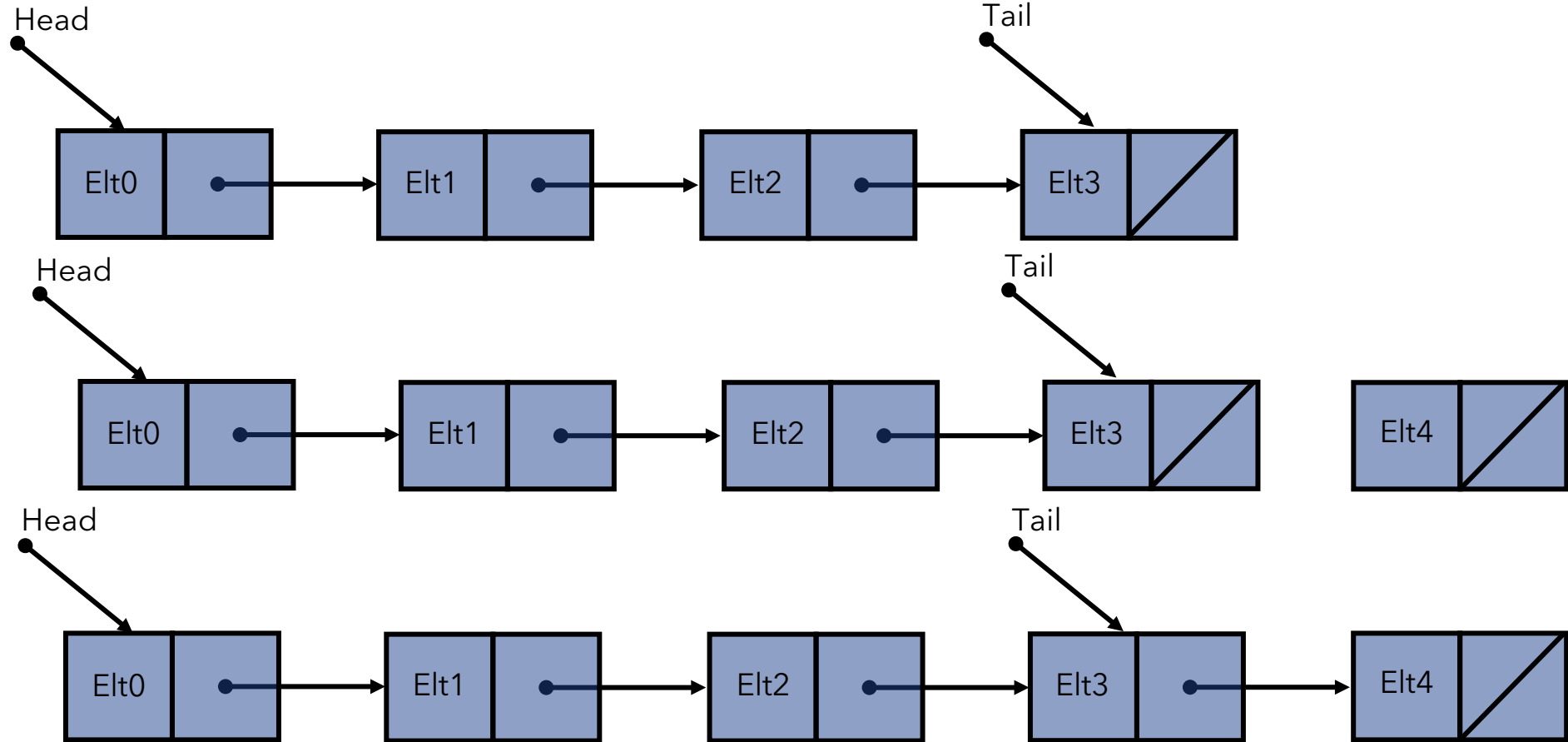

Getting the last element



Getting the last element is $O(n)$

```
public E getLast() {  
    Node<E> finger = head;  
    //throw exception if list is empty  
    while (finger != null && finger.next() != null)  
    {  
        finger = finger.next();  
    }  
    return finger.value();  
}
```

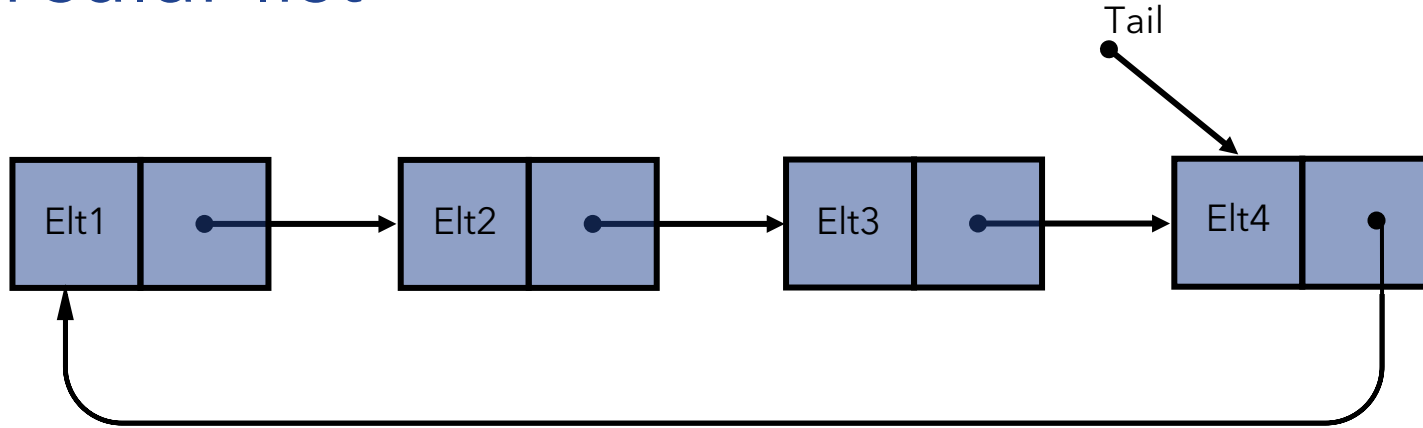
Adding when having a tail pointer is $O(1)$



Removing at the tail is $O(n)$

- We still have to traverse the whole linked list - not efficient!

Circular list



- Accessing/modifying the head or the tail is $O(1)$.
- Circular lists are as space-efficient as singly linked lists but tail-related operations are less costly.