# CSO62 DATA STRUCTURES AND ADVANCED PROGRAMMING 

23: Shortest Paths

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## Lecture 23: Shortest Paths

- Introduction to Shortest Paths
- API
- Properties
- Dijkstra's Algorithm

Edge-weighted digraph

- Edge-weighted digraph: a digraph where we associate weights or costs with each edge.

```
edge-weighted digraph
    4->5 0.35
    5->4 0.35
    4->7 0.37
    5->7 0.28
    7->5 0.28
    5->1 0.32
    0->4 0.38
    0->2 0.26
    7->3 0.39
    1->3 0.29
    2->7 0.34
    6->2 0.40
    3->6 0.52
    6->0 0.58
    6->4 0.93
```



## Shortest Paths

- Shortest path from vertex S to vertex t : a directed path from $s$ to $t$ with the property that no other such path has a lower weight (total weight sum of edges it consists of).


An edge-weighted digraph and a shortest path

## Shortest Path variants

- Single source: from one vertex s to every other vertex.
- Single sink: from every vertex to one vertex $t$.
- Source-sink: from one vertex $s$ to another vertex $t$.
- All pairs: from every vertex to every other vertex.
- What version is there in your navigation app?


## Shortest Paths Assumptions

- Not all vertices need to be reachable.
- We will assume so in this lecture.
, Weights are non-negative.
, There are algorithms that can handle negative weights.
- Shortest paths are not necessarily unique but they are simple.


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## Weighted directed edge API

- public class DirectedEdge
- DirectedEdge(int v, int w, double weight)
- Constructs a weighted edge from $v$ to $w(v->w)$ with the provided weight.
- int from()
- Returns vertex source of this edge.
- int to()
- Returns vertex destination of this edge.
- double weight()
- Returns weight of this edge.
- String toString()
- Returns the string representation of this edge.


## Weighted directed edge in Java

```
public class DirectedEdge {
    private final int v;
    private final int w;
    private final double weight;
    public DirectedEdge(int v, int w, double weight) {
            this.v = v;
            this.w = w;
            this.weight = weight;
    }
    public int from() {
        return v;
    }
    public int to() {
        return w;
    }
    public double weight() {
        return weight;
    }
```


## Edge-weighted digraph API

p public class EdgeWeightedDigraph

- EdgeWeightedDigraph(int v)
- Constructs an edge-weighted digraph with $v$ vertices.
- void addEdge(DirectedEdge e)
- Add weighted directed edge e.
- Iterable<DirectedEdge> adj(int v)
- Returns edges adjacent from v .
- int V()
- Returns number of vertices.
- int E()
- Returns number of edges.
- Iterable<DirectedEdge> edges()
- Returns all edges.


## Edge-weighted digraph adjacency list representation

| tinyEWD.txt |  |  |
| :---: | :---: | :---: |
|  |  |  |
| 15 |  |  |
| 4 | 5 | 0.35 |
| 5 | 4 | 0.35 |
| 4 | 7 | 0.37 |
| 5 | 7 | 0.28 |
| 7 | 5 | 0.28 |
| 5 | 1 | 0.32 |
| 0 | 4 | 0.38 |
| 0 | 2 | 0.26 |
| 7 | 3 | 0.39 |
| 1 | 3 | 0.29 |
| 2 | 7 | 0.34 |
| 6 | 2 | 0.40 |
| 3 | 6 | 0.52 |
| 6 | 0 | 0.58 |
| 6 | 4 | 0.93 |



Edge-weighted digraph representation

## Edge-weighted digraph in Java

```
public class EdgeWeightedDigraph {
    private final int V; // number of vertices in this digraph
    private int E; // number of edges in this digraph
    private Bag<DirectedEdge>[] adj; // adj[v] = adjacency list for vertex v
    public EdgeWeightedDigraph(int V) {
        this.V = V;
        this.E = 0;
        adj = (Bag<DirectedEdge>[]) new Bag[V];
            for (int v = 0; v < V; v++)
                adj[v] = new Bag<DirectedEdge>();
    }
    public void addEdge(DirectedEdge e) {
        int v = e.from();
        int w = e.to();
        adj[v].add(e);
        E++;
    }
    public Iterable<DirectedEdge> adj(int v) {
        return adj[v];
    }
```


## Single-source shortest path API

- Goal: find shortest path from S to every other vertex in the digraph.
- public class SP
- SP(EdgeWeightedDigraph G, int s)
- Shortest paths from s in digraph G .
- double distTo(int v)
- Length of shortest path from S to V .
- Iterable<DirectedEdge> pathTo(int v)
- Returns edges along the shortest path from s to V .
- boolean hasPathTo(int v)
- Returns whether there is a path from s to V .


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Data structures for single-source shortest paths

- Goal: find shortest path from S to every other vertex in the digraph.
- Shortest-paths tree (SPT): a subgraph containing $s$ and all the vertices reachable from s that forms a directed tree rooted at s such that every tree path in the SPT is a shortest path in the digraph.
- Representation of shortest paths with two vertex-indexed arrays.
- Edges on the shortest-paths tree: edgeTo[v] is the last edge on a shortest path from s to V .
- Distance to the source: distTo[v] is the length of the shortest path from $s$ to $v$.



## Edge relaxation

- Relax edge $\mathrm{e}=\mathrm{v}->\mathrm{W}$
- distTo[v] is the length of the shortest known path from $s$ to V.
- distTo[w] is the length of the shortest known path from s to W.
, edgeTo[w] is the last edge on shortest known path from s to w.
- If $\mathrm{e}=\mathrm{v}->\mathrm{w}$ yields shorter path to w , update distTo[w] and edgeTo[w].


## Edge relaxation



## Edge relaxation implementation

```
private void relax(DirectedEdge e) {
    int v = e.from(), w = e.to();
    if (distTo[w] > distTo[v] + e.weight()) {
        distTo[w] = distTo[v] + e.weight();
        edgeTo[w] = e;
    }
}
```

Framework for shortest-paths algorithm

- Generic algorithm to compute a SPT from s
- distTo[ V$]=\infty$ for each vertex V .
- edgeTo[v]=null for each vertex v .
- distTo[s]=0.
- Repeat until done:
- Relax any edge.
- distTo[v] is the length of a simple path from $s$ to $v$.
, distTo[v] does not increase.


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Algorithms

Dijkstra's Algorithm Demo

Robert Sedgewick I Kevin Wayne
http://algs4.cs.princeton.edu

## Dijkstra's algorithm demo

- Consider vertices in increasing order of distance from s (non-tree vertex with the lowest distTo[] value).
- Add vertex to tree and relax all edges adjacent from that vertex.



## Dijkstra's algorithm demo

- Consider vertices in increasing order of distance from s (non-tree vertex with the lowest distTo[] value).
- Add vertex to tree and relax all edges adjacent from that vertex.


$\longrightarrow 0$| $v$ | distTo[] | edgeTo[] |
| :---: | :---: | :---: |
| 0 | 0.0 | - |

choose source vertex 0

## Dijkstra's algorithm demo

- Consider vertices in increasing order of distance from s (non-tree vertex with the lowest distTo[] value).
- Add vertex to tree and relax all edges adjacent from that vertex.

relax all edges adjacent from 0


## Dijkstra's algorithm demo

- Consider vertices in increasing order of distance from s (non-tree vertex with the lowest distTo[] value).
- Add vertex to tree and relax all edges adjacent from that vertex.

relax all edges adjacent from 0


## Dijkstra's algorithm demo

- Consider vertices in increasing order of distance from s (non-tree vertex with the lowest distTo[] value).
- Add vertex to tree and relax all edges adjacent from that vertex.


| $v$ | distTo[] | edgeTo[ |
| :---: | :---: | :---: |
| 0 | 0.0 | - |
| 1 | 5.0 | $0 \rightarrow 1$ |
| 2 |  |  |
| 3 |  |  |
| 4 | 9.0 | $0 \rightarrow 4$ |
| 5 |  |  |
| 6 |  | $0 \rightarrow 7$ |

choose vertex 1

## Dijkstra's algorithm demo

- Consider vertices in increasing order of distance from s (non-tree vertex with the lowest distTo[] value).
- Add vertex to tree and relax all edges adjacent from that vertex.


| $v$ | distTo[] | edgeTo[] |
| :---: | :---: | :---: |
| 0 | 0.0 | - |
| 1 | 5.0 | $0 \rightarrow 1$ |
| 2 |  |  |
| 3 |  |  |
| 4 | 9.0 | $0 \rightarrow 4$ |
| 5 |  |  |
| 6 |  |  |
| 7 | 8.0 | $0 \rightarrow 7$ |

## Dijkstra's algorithm demo

- Consider vertices in increasing order of distance from s (non-tree vertex with the lowest distTo[] value).
- Add vertex to tree and relax all edges adjacent from that vertex.


| $v$ | distTo[] | edgeTo[] |
| :---: | :---: | :---: |
| 0 | 0.0 | - |
| 1 | 5.0 | $0 \rightarrow 1$ |
| 2 |  |  |
| 3 |  |  |
| 4 | 9.0 | $0 \rightarrow 4$ |
| 5 |  |  |
| 6 |  |  |
| 7 | 8.0 | $0 \rightarrow 7$ |

relax all edges adjacent from 1

## Dijkstra's algorithm demo

- Consider vertices in increasing order of distance from s (non-tree vertex with the lowest distTo[] value).
- Add vertex to tree and relax all edges adjacent from that vertex.


| $v$ | distTo[] | edgeTo[ |
| :---: | :---: | :---: |
| 0 | 0.0 | - |
| $\rightarrow 1$ | 5.0 | $0 \rightarrow 1$ |
| 2 | 17.0 | $1 \rightarrow 2$ |
| 3 | 20.0 | $1 \rightarrow 3$ |
| 4 | 9.0 | $0 \rightarrow 4$ |
| 5 |  |  |
| 6 |  |  |
| 7 | 8.0 | $0 \rightarrow 7$ |

relax all edges adjacent from 1

## Dijkstra's algorithm demo

- Consider vertices in increasing order of distance from s (non-tree vertex with the lowest distTo[] value).
- Add vertex to tree and relax all edges adjacent from that vertex.


| $v$ | distTo[] | edgeTo[] |
| :---: | :---: | :---: |
| 0 | 0.0 | - |
| 1 | 5.0 | $0 \rightarrow 1$ |
| 2 | 17.0 | $1 \rightarrow 2$ |
| 3 | 20.0 | $1 \rightarrow 3$ |
| 4 | 9.0 | $0 \rightarrow 4$ |
| 5 |  |  |
| 6 |  |  |
| 7 | 8.0 | $0 \rightarrow 7$ |

## Dijkstra's algorithm demo

- Consider vertices in increasing order of distance from s (non-tree vertex with the lowest distTo[] value).
- Add vertex to tree and relax all edges adjacent from that vertex.

choose vertex 7


## Dijkstra's algorithm demo

- Consider vertices in increasing order of distance from s (non-tree vertex with the lowest distTo[] value).
- Add vertex to tree and relax all edges adjacent from that vertex.

relax all edges adjacent from 7


## Dijkstra's algorithm demo

- Consider vertices in increasing order of distance from s (non-tree vertex with the lowest distTo[] value).
- Add vertex to tree and relax all edges adjacent from that vertex.

relax all edges adjacent from 7


## Dijkstra's algorithm demo

- Consider vertices in increasing order of distance from s (non-tree vertex with the lowest distTo[] value).
- Add vertex to tree and relax all edges adjacent from that vertex.


| $v$ | distTo[] | edgeTo[] |
| :---: | :---: | :---: |
| 0 | 0.0 | - |
| 1 | 5.0 | $0 \rightarrow 1$ |
| 2 | 15.0 | $7 \rightarrow 2$ |
| 3 | 20.0 | $1 \rightarrow 3$ |
| 4 | 9.0 | $0 \rightarrow 4$ |
| 5 | 14.0 | $7 \rightarrow 5$ |
| 6 |  |  |
| 7 | 8.0 | $0 \rightarrow 7$ |

## Dijkstra's algorithm demo

- Consider vertices in increasing order of distance from s (non-tree vertex with the lowest distTo[] value).
- Add vertex to tree and relax all edges adjacent from that vertex.

select vertex 4


## Dijkstra's algorithm demo

- Consider vertices in increasing order of distance from s (non-tree vertex with the lowest distTo[] value).
- Add vertex to tree and relax all edges adjacent from that vertex.

relax all edges adjacent from 4


## Dijkstra's algorithm demo

- Consider vertices in increasing order of distance from s (non-tree vertex with the lowest distTo[] value).
- Add vertex to tree and relax all edges adjacent from that vertex.

relax all edges adjacent from 4


## Dijkstra's algorithm demo

- Consider vertices in increasing order of distance from s (non-tree vertex with the lowest distTo[] value).
- Add vertex to tree and relax all edges adjacent from that vertex.


| $v$ | distTo[] | edgeTo[] |
| :---: | :---: | :---: |
| 0 | 0.0 | - |
| 1 | 5.0 | $0 \rightarrow 1$ |
| 2 | 15.0 | $7 \rightarrow 2$ |
| 3 | 20.0 | $1 \rightarrow 3$ |
| 4 | 9.0 | $0 \rightarrow 4$ |
| 5 | 13.0 | $4 \rightarrow 5$ |
| 6 | 29.0 | $4 \rightarrow 6$ |
| 7 | 8.0 | $0 \rightarrow 7$ |

## Dijkstra's algorithm demo

- Consider vertices in increasing order of distance from s (non-tree vertex with the lowest distTo[] value).
- Add vertex to tree and relax all edges adjacent from that vertex.

select vertex 5


## Dijkstra's algorithm demo

- Consider vertices in increasing order of distance from s (non-tree vertex with the lowest distTo[] value).
- Add vertex to tree and relax all edges adjacent from that vertex.

relax all edges adjacent from 5


## Dijkstra's algorithm demo

- Consider vertices in increasing order of distance from s (non-tree vertex with the lowest distTo[] value).
- Add vertex to tree and relax all edges adjacent from that vertex.

relax all edges adjacent from 5


## Dijkstra's algorithm demo

- Consider vertices in increasing order of distance from s (non-tree vertex with the lowest distTo[] value).
- Add vertex to tree and relax all edges adjacent from that vertex.


| $v$ | distTo[] | edgeTo[] |
| :---: | :---: | :---: |
| 0 | 0.0 | - |
| 1 | 5.0 | $0 \rightarrow 1$ |
| 2 | 14.0 | $5 \rightarrow 2$ |
| 3 | 20.0 | $1 \rightarrow 3$ |
| 4 | 9.0 | $0 \rightarrow 4$ |
| 5 | 13.0 | $4 \rightarrow 5$ |
| 6 | 26.0 | $5 \rightarrow 6$ |
| 7 | 8.0 | $0 \rightarrow 7$ |

## Dijkstra's algorithm demo

- Consider vertices in increasing order of distance from s (non-tree vertex with the lowest distTo[] value).
- Add vertex to tree and relax all edges adjacent from that vertex.

select vertex 2


## Dijkstra's algorithm demo

- Consider vertices in increasing order of distance from s (non-tree vertex with the lowest distTo[] value).
- Add vertex to tree and relax all edges adjacent from that vertex.

relax all edges adjacent from 2


## Dijkstra's algorithm demo

- Consider vertices in increasing order of distance from s (non-tree vertex with the lowest distTo[] value).
- Add vertex to tree and relax all edges adjacent from that vertex.

relax all edges adjacent from 2


## Dijkstra's algorithm demo

- Consider vertices in increasing order of distance from s (non-tree vertex with the lowest distTo[] value).
- Add vertex to tree and relax all edges adjacent from that vertex.


| $v$ | distTo[] | edgeTo[] |
| :---: | :---: | :---: |
| 0 | 0.0 | - |
| 1 | 5.0 | $0 \rightarrow 1$ |
| 2 | 14.0 | $5 \rightarrow 2$ |
| 3 | 17.0 | $2 \rightarrow 3$ |
| 4 | 9.0 | $0 \rightarrow 4$ |
| 5 | 13.0 | $4 \rightarrow 5$ |
| 6 | 25.0 | $2 \rightarrow 6$ |
| 7 | 8.0 | $0 \rightarrow 7$ |

## Dijkstra's algorithm demo

- Consider vertices in increasing order of distance from s (non-tree vertex with the lowest distTo[] value).
- Add vertex to tree and relax all edges adjacent from that vertex.

select vertex 3


## Dijkstra's algorithm demo

- Consider vertices in increasing order of distance from s (non-tree vertex with the lowest distTo[] value).
- Add vertex to tree and relax all edges adjacent from that vertex.

relax all edges adjacent from 3


## Dijkstra's algorithm demo

- Consider vertices in increasing order of distance from s (non-tree vertex with the lowest distTo[] value).
- Add vertex to tree and relax all edges adjacent from that vertex.

relax all edges adjacent from 3


## Dijkstra's algorithm demo

- Consider vertices in increasing order of distance from s (non-tree vertex with the lowest distTo[] value).
- Add vertex to tree and relax all edges adjacent from that vertex.


| $v$ | distTo[] | edgeTo[] |
| :---: | :---: | :---: |
| 0 | 0.0 | - |
| 1 | 5.0 | $0 \rightarrow 1$ |
| 2 | 14.0 | $5 \rightarrow 2$ |
| 3 | 17.0 | $2 \rightarrow 3$ |
| 4 | 9.0 | $0 \rightarrow 4$ |
| 5 | 13.0 | $4 \rightarrow 5$ |
| 6 | 25.0 | $2 \rightarrow 6$ |
| 7 | 8.0 | $0 \rightarrow 7$ |

## Dijkstra's algorithm demo

- Consider vertices in increasing order of distance from s (non-tree vertex with the lowest distTo[] value).
- Add vertex to tree and relax all edges adjacent from that vertex.

select vertex 6


## Dijkstra's algorithm demo

- Consider vertices in increasing order of distance from s (non-tree vertex with the lowest distTo[] value).
- Add vertex to tree and relax all edges adjacent from that vertex.

relax all edges adjacent from 6


## Dijkstra's algorithm demo

- Consider vertices in increasing order of distance from s (non-tree vertex with the lowest distTo[] value).
- Add vertex to tree and relax all edges adjacent from that vertex.


| $v$ | distTo[] | edgeTo[] |
| :---: | :---: | :---: |
| 0 | 0.0 | - |
| 1 | 5.0 | $0 \rightarrow 1$ |
| 2 | 14.0 | $5 \rightarrow 2$ |
| 3 | 17.0 | $2 \rightarrow 3$ |
| 4 | 9.0 | $0 \rightarrow 4$ |
| 5 | 13.0 | $4 \rightarrow 5$ |
| 6 | 25.0 | $2 \rightarrow 6$ |
| 7 | 8.0 | $0 \rightarrow 7$ |

## Dijkstra's algorithm demo

- Consider vertices in increasing order of distance from s (non-tree vertex with the lowest distTo[] value).
- Add vertex to tree and relax all edges adjacent from that vertex.


| $v$ | distTo[] | edgeTo[] |
| :---: | :---: | :---: |
| 0 | 0.0 | - |
| 1 | 5.0 | $0 \rightarrow 1$ |
| 2 | 14.0 | $5 \rightarrow 2$ |
| 3 | 17.0 | $2 \rightarrow 3$ |
| 4 | 9.0 | $0 \rightarrow 4$ |
| 5 | 13.0 | $4 \rightarrow 5$ |
| 6 | 25.0 | $2 \rightarrow 6$ |
| 7 | 8.0 | $0 \rightarrow 7$ |

shortest-paths tree from vertex s

## Indexed min-priority queue (Section 2.4 in textbook)

- Associate an index between 0 and $\mathrm{n}-1$ with each key in a priority queue.
- Insert a key associated with a given index.
- Delete a minimum key and return associated index.
- Decrease the key associated with a given index.
p public class IndexMinPQ<Key extends Comparable<Key>>
- IndexMinPQ(int $n$ )
- Create indexed PQ with indices $0,1, \ldots \mathrm{n}-1$
- void insert(int i, Key key)
- Associate key with index i.
- int delMin()
- Remove a minimal key and return its associated index.
- void decreaseKey(int i, Key key)
- Decrease the key with index ito the specified valye.

```
public class DijkstraSP {
    private double[] distTo; // distTo[v] = distance of shortest s->V path
    private DirectedEdge[] edgeTo; // edgeTo[v] = last edge on shortest s->V path
    private IndexMinPQ<Double> pq; // priority queue of vertices
    public DijkstraSP(EdgeWeightedDigraph G, int s) {
        distTo = new double[G.V()];
        edgeTo = new DirectedEdge[G.V()];
        for (int v = 0; v < G.V(); v++)
            distTo[v] = Double.POSITIVE_INFINITY;
        distTo[s] = 0.0;
        // relax vertices in order of distance from s
        pq = new IndexMinPQ<Double>(G.V());
        pq.insert(s, distTo[s]);
        while (!pq.isEmpty()) {
            int v = pq.delMin();
            for (DirectedEdge e : G.adj(v))
                relax(e);
        }
    }
    // relax edge e and update pq if changed
    private void relax(DirectedEdge e) {
        int v = e.from(), w = e.to();
        if (distTo[w] > distTo[v] + e.weight()) {
            distTo[w] = distTo[v] + e.weight();
            edgeTo[w] = e;
            if (pq.contains(w)) pq.decreaseKey(w, distTo[w]);
            else pq.insert(w, distTo[w]);
        }
    }
```

Running time depends on PQ implementation

- Many variations. Assuming binary heap, running time is proportional to $|E| \log |V|$ and $|V|$ extra space.
- Cost of insert, delete-min, decrease-key are all $\log |V|$.
- More complicated version with a Fibonacci heap (CS140...) takes $O(|E|+|V| \log |V|)$ time but in practice it's not worth implementing.


## Practice Time

- Run Dijkstra's algorithm on the following graph with 0 being the starting vertex.



## Answer



| $v$ | distTo[] | edgeTo[] |
| :---: | :---: | :---: |
| 0 | 0 | - |
| 1 | 6 | $3->1$ |
| 2 | 2 | $0->2$ |
| 3 | 4 | $2->3$ |
| 4 | 5 | $3->4$ |
| 5 | 8 | $6->5$ |
| 6 | 6 | $4->6$ |
| 7 | 11 | $5->7$ |

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## Readings:

- Recommended Textbook: Chapter 4.4 (Pages 638-676)
, Website:
- https://algs4.cs.princeton.edu/44sp/


## Practice Problems:

Run Dijkstra's algorithm on the graph on the right with A being the starting vertex.


