# **CS062** DATA STRUCTURES AND ADVANCED PROGRAMMING

# 23: Shortest Paths



Alexandra Papoutsaki she/her/hers Lecture 23: Shortest Paths

- Introduction to Shortest Paths
- API
- Properties
- Dijkstra's Algorithm

### Edge-weighted digraph

Edge-weighted digraph: a digraph where we associate weights or costs with each edge.

#### edge-weighted digraph

4->5	0.35	-
5->4	0.35	$(1) \rightarrow (3)$
4->7	0.37	(5)
5->7	0.28	$1  \bigcirc \frown \frown$
7->5	0.28	
5->1	0.32	
0->4	0.38	
0->2	0.26	
7->3	0.39	
1->3	0.29	
2->7	0.34	
6->2	0.40	
3->6	0.52	
6->0	0.58	
6->4	0.93	

**Shortest Paths** 

Shortest path from vertex S to vertex t: a directed path from S to t with the property that no other such path has a lower weight (total weight sum of edges it consists of).

edge-weighted digraph



An edge-weighted digraph and a shortest path

Shortest Path variants

- Single source: from one vertex S to every other vertex.
- Single sink: from every vertex to one vertex t.
- Source-sink: from one vertex S to another vertex t.
- All pairs: from every vertex to every other vertex.

What version is there in your navigation app?

**Shortest Paths Assumptions** 

- Not all vertices need to be reachable.
  - We will assume so in this lecture.
- Weights are non-negative.
  - > There are algorithms that can handle negative weights.
- Shortest paths are not necessarily unique but they are simple.

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#### Weighted directed edge API

- > public class DirectedEdge
  - DirectedEdge(int v, int w, double weight)
    - Constructs a weighted edge from v to w (v->w) with the provided weight.
  - int from()
    - Returns vertex source of this edge.
  - int to()
    - Returns vertex destination of this edge.
  - > double weight()
    - Returns weight of this edge.
  - > String toString()
    - Returns the string representation of this edge.

#### Weighted directed edge in Java

```
public class DirectedEdge {
    private final int v;
    private final int w;
    private final double weight;
   public DirectedEdge(int v, int w, double weight) {
        this.v = v;
        this.w = w;
        this.weight = weight;
    }
   public int from() {
        return v;
    }
    public int to() {
        return w;
    }
    public double weight() {
        return weight;
    }
```

### Edge-weighted digraph API

#### > public class EdgeWeightedDigraph

- > EdgeWeightedDigraph(int v)
  - Constructs an edge-weighted digraph with v vertices.
- void addEdge(DirectedEdge e)
  - Add weighted directed edge e.
- > Iterable<DirectedEdge> adj(int v)
  - Returns edges adjacent from V.
- int V()
  - Returns number of vertices.
- int E()
  - Returns number of edges.
- > Iterable<DirectedEdge> edges()
  - Returns all edges.

Edge-weighted digraph adjacency list representation



Edge-weighted digraph representation

#### Edge-weighted digraph in Java

```
public class EdgeWeightedDigraph {
    private final int V;
                                         // number of vertices in this digraph
    private int E;
                                         // number of edges in this digraph
    private Bag<DirectedEdge>[] adj;
                                         // adj[v] = adjacency list for vertex v
    public EdgeWeightedDigraph(int V) {
       this.V = V;
       this.E = 0;
       adj = (Bag<DirectedEdge>[]) new Bag[V];
        for (int v = 0; v < V; v++)
            adj[v] = new Bag<DirectedEdge>();
    }
    public void addEdge(DirectedEdge e) {
        int v = e.from();
        int w = e.to();
        adj[v].add(e);
        E++;
    }
   public Iterable<DirectedEdge> adj(int v) {
       return adj[v];
    }
```

#### Single-source shortest path API

- Goal: find shortest path from S to every other vertex in the digraph.
- > public class SP
  - > SP(EdgeWeightedDigraph G, int s)
    - Shortest paths from s in digraph G.
  - > double distTo(int v)
    - Length of shortest path from S to V.
  - Iterable<DirectedEdge> pathTo(int v)
    - Returns edges along the shortest path from S to V.
  - boolean hasPathTo(int v)
    - Returns whether there is a path from S to V.

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Data structures for single-source shortest paths

- Goal: find shortest path from S to every other vertex in the digraph.
- Shortest-paths tree (SPT): a subgraph containing S and all the vertices reachable from S that forms a directed tree rooted at s such that every tree path in the SPT is a shortest path in the digraph.
- Representation of shortest paths with two vertex-indexed arrays.
  - Edges on the shortest-paths tree: edgeTo[v] is the last edge on a shortest path from S to V.
  - Distance to the source: distTo[v] is the length of the shortest path from S to V.

#### PROPERTIES

```
public Iterable<DirectedEdge> pathTo(int v) {
    Stack<DirectedEdge> path = new Stack<DirectedEdge>();
    for (DirectedEdge e = edgeTo[v]; e != null; e = edgeTo[e.from()]) {
        path.push(e);
    }
    return path;
}
```

	edgeTo[]	distTo[]
0	null	0
1	5->1 0.3	32 1.05
2	0->2 0.2	0.26
3	7->3 0.3	0.99
4	0->4 0.3	0.38
5	4->5 0.3	35 0.73
6	3->6 0.5	1.51
7	2->7 0.3	84 0.60



16

4->5 0.35

5->4 0.35

4->7 0.37

5->7 0.28

7->5 0.28

5->1 0.32

0->4 0.38

0->2 0.26

7->3 0.39

1->3 0.29

2->7 0.34

6->2 0.40

3->6 0.52

6->0 0.58

6->4 0.93

### Edge relaxation

- Relax edge e = v->w
  - distTo[v] is the length of the shortest known path from S to v.
  - distTo[w] is the length of the shortest known path from S to w.
  - edgeTo[w] is the last edge on shortest known path from S to W.
  - If e = v->w yields shorter path to w, update distTo[w] and edgeTo[w].

### Edge relaxation



#### Edge relaxation implementation

```
private void relax(DirectedEdge e) {
    int v = e.from(), w = e.to();
    if (distTo[w] > distTo[v] + e.weight()) {
        distTo[w] = distTo[v] + e.weight();
        edgeTo[w] = e;
    }
}
```

Framework for shortest-paths algorithm

- Generic algorithm to compute a SPT from s
  - distTo $[v] = \infty$  for each vertex v.
  - edgeTo[v]=null for each vertex v.
  - distTo[s]=0.
  - Repeat until done:
    - Relax any edge.
- distTo[v] is the length of a simple path from s to v.
- distTo[v] does not increase.

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## Algorithms

 $\checkmark$ 

#### ROBERT SEDGEWICK | KEVIN WAYNE

### DIJKSTRA'S ALGORITHM DEMO



Robert Sedgewick | Kevin Wayne

http://algs4.cs.princeton.edu

- Consider vertices in increasing order of distance from s (non-tree vertex with the lowest distTo[] value).
- Add vertex to tree and relax all edges adjacent from that vertex.







#### choose source vertex 0

5.0

9.0

8.0

15.0

4.0

3.0

9.0

4.0 20.0

5.0

1.0 13.0

11.0

1→2 12.0

0→1

0**→**4

1→3

1→7

2→3

2→6

3→6

4→6

4→7

5→6



relax all edges adjacent from 0



• Add vertex to tree and relax all edges adjacent from that vertex.



relax all edges adjacent from 0

0→1

0→4

0→7

1→3 1→7

2→3

2→6

3→6

4→5

4→6

4→7

1→2 12.0

5.0 9.0

8.0

15.0

11.0

9.0

4.0

5.0

20.0



• Add vertex to tree and relax all edges adjacent from that vertex.



#### choose vertex 1

0→1

 $0 \rightarrow 4$ 

0**→**7

1→3 1→7

2→3

2→6

3→6

4→5

4→6

4→7

5→2

1→2 12.0

5.0 9.0

8.0

15.0

11.0

9.0

4.0

5.0

1.0

20.0

- Consider vertices in increasing order of distance from s (non-tree vertex with the lowest distTo[] value).
- Add vertex to tree and relax all edges adjacent from that vertex.



0→1

 $0 \rightarrow 4$ 

0**→**7

1→3 1→7

2→3

2→6

3→6

4→5

4→6

4→7

5→2

1→2 12.0

5.0 9.0

8.0

15.0

11.0

9.0

4.0

5.0

1.0

20.0



relax all edges adjacent from 1



relax all edges adjacent from 1

 Consider vertices in increasing order of distance from s (non-tree vertex with the lowest distTo[] value).

• Add vertex to tree and relax all edges adjacent from that vertex.



0→1

 $0 \rightarrow 4$ 

0**→**7

1→3 1→7

2→3

2→6

3→6

4→5

4→6

4→7

5→2

1→2 12.0

5.0 9.0

8.0

15.0

11.0

9.0

4.0

5.0

1.0

20.0



• Add vertex to tree and relax all edges adjacent from that vertex.



#### choose vertex 7

0→1

 $0 \rightarrow 4$ 

0**→**7

1→3 1→7

2→3

2→6

3→6

4→5

4→6

4→7

5→2

1→2 12.0

5.0 9.0

8.0

15.0

11.0

9.0

4.0

5.0

1.0

20.0



0→1

 $0 \rightarrow 4$ 

0**→**7

1→3 1→7

2→3

2→6

3→6

4→5

4→6

4→7

5→2

5→6

**7**→5 7→2

0→1

1→2

1→3

0→4

0→7

1→2 12.0

5.0 9.0

8.0

15.0

11.0

9.0

4.0

5.0

1.0

6.0

7.0

13.0

20.0



• Add vertex to tree and relax all edges adjacent from that vertex.



relax all edges adjacent from 7

0→1

 $0 \rightarrow 4$ 

0**→**7

1→3 1→7

2→3

2→6

3→6

4→5

4→6

4→7

5→2

1→2 12.0

5.0 9.0

8.0

15.0

11.0

9.0

4.0

5.0

1.0

20.0

 Consider vertices in increasing order of distance from s (non-tree vertex with the lowest distTo[] value).

• Add vertex to tree and relax all edges adjacent from that vertex.



0→1

0→4

0→7

1→3 1→7

2→3

2→6

3→6

4→5

4→6

4→7

5→2

1→2 12.0

5.0 9.0

8.0

15.0

11.0

9.0

4.0

5.0

1.0

20.0



• Add vertex to tree and relax all edges adjacent from that vertex.



#### select vertex 4

0→1

0**→**4

0**→**7

1→3 1→7

2→3

2→6

3→6

4→5

4→6

4→7

5→2

1→2 12.0

5.0 9.0

8.0

15.0

11.0

9.0

4.0

5.0

1.0

20.0



14

20

9

relax all edges adjacent from 4

0→1

0**→**4

0**→**7

1→3 1→7

2→3

2→6

3→6

4→6

4→7

5→2

5→6

**7**→5

7→2

edgeTo[]

0→1

7→2

1→3

0→4

**7**→5

0→7

6

7

 $\infty$ 

6

8.0

1→2 12.0

5.0 9.0

8.0

15.0

11.0

9.0

4.0 20.0

5.0

1.0

6.0

7.0

13.0



• Add vertex to tree and relax all edges adjacent from that vertex.



relax all edges adjacent from 4

0→1

0→4

0**→**7

1→3 1→7

2→3

2→6

3→6

4→5

4→6

4→7

5→2

1→2 12.0

5.0 9.0

8.0

15.0

11.0

9.0

4.0

5.0

1.0

20.0

 Consider vertices in increasing order of distance from s (non-tree vertex with the lowest distTo[] value).

• Add vertex to tree and relax all edges adjacent from that vertex.



0→1

0→4

0→7

1→3 1→7

2→3

2→6

3→6

4→5

4→6

4→7

5→2

1→2 12.0

5.0 9.0

8.0

15.0

11.0

9.0

4.0

5.0

1.0

20.0



• Add vertex to tree and relax all edges adjacent from that vertex.



#### select vertex 5

0→1

0**→**4

0→7

1→3 1→7

2→3

2→6

3→6

4→5

4→6

4→7

5→2

1→2 12.0

5.0 9.0

8.0

15.0

11.0

9.0

4.0

5.0

1.0

20.0



• Add vertex to tree and relax all edges adjacent from that vertex.



0→1

0**→**4

0**→**7

1→3 1→7

2→3

2→6

3→6

4→5

4→6

4→7

5→2

1→2 12.0

5.0 9.0

8.0

15.0

11.0

9.0

4.0

5.0

1.0

20.0



• Add vertex to tree and relax all edges adjacent from that vertex.



0→1

0**→**4

0**→**7

1→3 1→7

2→3

2→6

3→6

4→5

4→6

4→7

5→2

1→2 12.0

5.0 9.0

8.0

15.0

11.0

9.0

4.0

5.0

1.0

20.0

 Consider vertices in increasing order of distance from s (non-tree vertex with the lowest distTo[] value).

• Add vertex to tree and relax all edges adjacent from that vertex.



0→1

0→4

0→7

1→3 1→7

2→3

3→6

4→5

4→6

4→7

5→2

1→2 12.0

2→6 11.0

5.0 9.0

8.0

15.0

4.0 3.0

9.0

4.0

5.0

1.0

 Consider vertices in increasing order of distance from s (non-tree vertex with the lowest distTo[] value).

• Add vertex to tree and relax all edges adjacent from that vertex.



#### select vertex 2

0→1

0**→**4

0→7

1→3 1→7

2→3

3→6

4→5

4→6

4→7

5→2

1→2 12.0

2→6 11.0

5.0 9.0

8.0

15.0

4.0 3.0

9.0

4.0

5.0

1.0





relax all edges adjacent from 2

 Consider vertices in increasing order of distance from s (non-tree vertex with the lowest distTo[] value).

• Add vertex to tree and relax all edges adjacent from that vertex.



0→1

0→4

0→7

1→3 1→7

2→3

3→6

4→5

4→6

4→7

5→2

1→2 12.0

2→6 11.0

5.0 9.0

8.0

15.0

4.0 3.0

9.0

4.0

5.0

1.0

 Consider vertices in increasing order of distance from s (non-tree vertex with the lowest distTo[] value).

• Add vertex to tree and relax all edges adjacent from that vertex.



#### select vertex 3

0→1

0→4

0→7

1→3 1→7

2→3

3→6

4→5

4→6

4→7

5→2

1→2 12.0

2→6 11.0

5.0 9.0

8.0

15.0

4.0 3.0

9.0

4.0

5.0

1.0





 Consider vertices in increasing order of distance from s (non-tree vertex with the lowest distTo[] value).

Add vertex to tree and relax all edges adjacent from that vertex



	0→4	9.0
	0→7	8.0
	1→2	12.0
	1→3	15.0
	1→7	4.0
	2→3	3.0
	2→6	11.0
	3→0 1→5	9.0
	4 <i>→</i> 5	20.0
at vartav	4→7	5.0
	5→2	1.0
	5→6	13.0
	7→5	6.0
	7→2	7.0
distTo[]	edgeTo[]	
0.0	-	
5.0	0→1	
14.0	5→2	
14.0 17.0	5→2 2→3	
14.0 17.0 9.0	$5 \rightarrow 2$ $2 \rightarrow 3$ $0 \rightarrow 4$	
14.0 17.0 9.0 13.0	$5 \rightarrow 2$ $2 \rightarrow 3$ $0 \rightarrow 4$ $4 \rightarrow 5$	
14.0 17.0 9.0 13.0 25.0	$5 \rightarrow 2$ $2 \rightarrow 3$ $0 \rightarrow 4$ $4 \rightarrow 5$ $2 \rightarrow 6$	
14.0 17.0 9.0 13.0 25.0 8.0	$5 \rightarrow 2$ $2 \rightarrow 3$ $0 \rightarrow 4$ $4 \rightarrow 5$ $2 \rightarrow 6$ $0 \rightarrow 7$	

V

0

1

2

3

4

5

6

7

0→1



• Add vertex to tree and relax all edges adjacent from that vertex.



#### select vertex 6

0→1

0→4

0→7

1→3 1→7

2→3

3→6

4→5

4→6

4→7

5→2

1→2 12.0

2→6 11.0

5.0 9.0

8.0

15.0

4.0 3.0

9.0

4.0

5.0

1.0



(non-tree vertex with the lowest distTo[] value).

• Add vertex to tree and relax all edges adjacent from that vertex.



0→1

0→4

0→7

1→3 1→7

2→3

3→6

4→5

4→6

4→7

5→2

1→2 12.0

2→6 11.0

5.0 9.0

8.0

15.0

4.0 3.0

9.0

4.0

5.0

1.0

 Consider vertices in increasing order of distance from s (non-tree vertex with the lowest distTo[] value).

Add vertex to tree and relax all edges adjacent from that vertex



	0→1	5.0
	0→4	9.0
	0→7	8.0
	1→2	12.0
	1→3	15.0
	1→7	4.0
	2→3 2→6	3.0
	2->0 3→6	9.0
	4→5	4.0
	4→6	20.0
at vertex.	4→7	5.0
	5→2	1.0
	5→6	13.0
	7→5	6.0
11 53	7→2	7.0
distic[]	edgelo[]	
0.0	-	
5.0	0→1	
14.0	5→2	
17.0	2→3	
9.0	0→4	
13.0		
	4→5	
25.0	4→5 2→6	
25.0 8.0	4→5 2→6 0→7	

V

 Consider vertices in increasing order of distance from s (non-tree vertex with the lowest distTo[] value).

• Add vertex to tree and relax all edges adjacent from that vertex.



#### shortest-paths tree from vertex s

5.0

9.0

8.0

15.0

11.0

9.0

4.0 20.0

5.0

1.0

4.0 3.0

0→1

0**→**4

0**→**7

1→3 1→7

2→3

2→6

3→6

4→6

4→7

5→2

1→2 12.0

### Indexed min-priority queue (Section 2.4 in textbook)

- Associate an index between 0 and n-1 with each key in a priority queue.
  - Insert a key associated with a given index.
  - > Delete a minimum key and return associated index.
  - Decrease the key associated with a given index.
- > public class IndexMinPQ<Key extends Comparable<Key>>
  - IndexMinPQ(int n)
    - Create indexed PQ with indices 0,1,...n-1
  - void insert(int i, Key key)
    - Associate key with index i.
  - int delMin()
    - Remove a minimal key and return its associated index.
  - void decreaseKey(int i, Key key)
    - > Decrease the key with index i to the specified valye.

```
public class DijkstraSP {
                                    // distTo[v] = distance of shortest s->v path
   private double[] distTo;
   private DirectedEdge[] edgeTo; // edgeTo[v] = last edge on shortest s->v path
   private IndexMinPQ<Double> pq;
                                      // priority queue of vertices
   public DijkstraSP(EdgeWeightedDigraph G, int s) {
        distTo = new double[G.V()];
        edgeTo = new DirectedEdge[G.V()];
        for (int v = 0; v < G.V(); v++)
            distTo[v] = Double.POSITIVE INFINITY;
        distTo[s] = 0.0;
        // relax vertices in order of distance from s
        pq = new IndexMinPQ<Double>(G.V());
        pq.insert(s, distTo[s]);
       while (!pq.isEmpty()) {
            int v = pq.delMin();
           for (DirectedEdge e : G.adj(v))
                relax(e);
        }
   }
   // relax edge e and update pq if changed
   private void relax(DirectedEdge e) {
        int v = e.from(), w = e.to();
        if (distTo[w] > distTo[v] + e.weight()) {
            distTo[w] = distTo[v] + e.weight();
            edgeTo[w] = e;
           if (pq.contains(w)) pq.decreaseKey(w, distTo[w]);
                                pq.insert(w, distTo[w]);
            else
        }
   }
```

Running time depends on PQ implementation

- Many variations. Assuming binary heap, running time is proportional to  $|E|\log|V|$  and |V| extra space.
  - Cost of insert, delete-min, decrease-key are all  $\log |V|$ .
- More complicated version with a Fibonacci heap (CS140...) takes  $O(|E| + |V|\log|V|)$  time but in practice it's not worth implementing.

#### **Practice Time**

Run Dijkstra's algorithm on the following graph with 0 being the starting vertex.



#### Answer



V	distTo[]	edgeTo[]
0	0	-
1	6	3->1
2	2	0->2
3	4	2->3
4	5	3->4
5	8	6->5
6	6	4->6
7	11	5->7

Lecture 23: Shortest Paths

- Introduction to Shortest Paths
- API
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- Dijkstra's Algorithm

### **Readings:**

- Recommended Textbook: Chapter 4.4 (Pages 638-676)
- Website:
  - https://algs4.cs.princeton.edu/44sp/

### **Practice Problems:**

Run Dijkstra's algorithm on the graph on the right with A being the starting vertex.

