CS062

DATA STRUCTURES AND ADVANCED PROGRAMMING

19: Binary Search Trees

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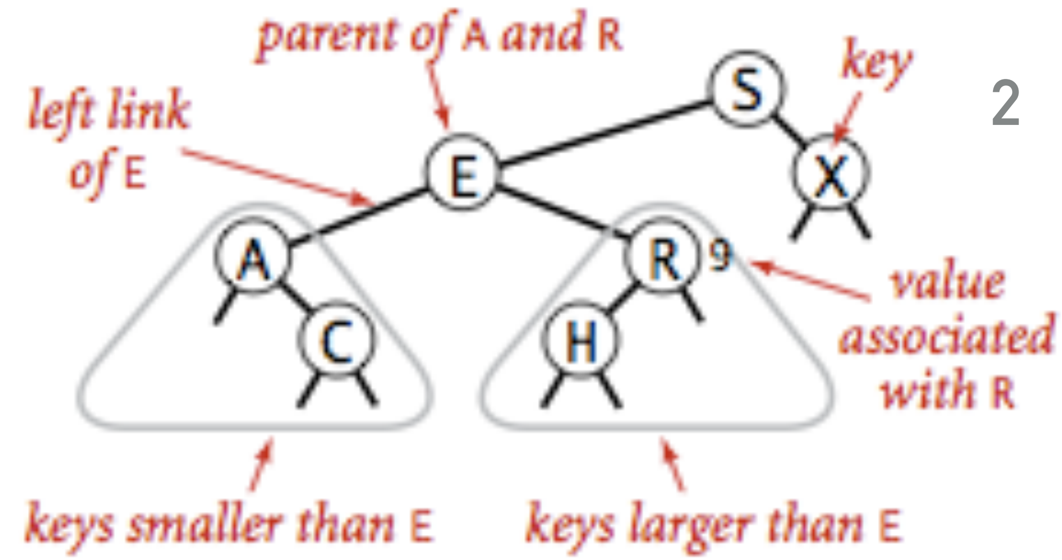
she/her/hers

Definitions



- Binary Search Tree: A binary tree in symmetric order.
- Symmetric order: Each node has a key, and every node's key is:
  - Larger than all keys in its left subtree.
  - Smaller than all keys in its right subtree.
- Our textbook uses BSTs to implement dictionaries, therefore each node holds a key-value pair. Other implementations hold only a key.

Differences between heaps and BSTs

	Heap	BST
Used to implement	Priority queues	Dictionaries
Supported operations	Insert, delete max	insert, search, delete, ordered operations
What is inserted	Keys	Key-value pairs
Underlying data structure	(Resizing) array	Linked nodes
Tree shape	Complete binary tree	Depends on data
Ordering of keys	Heap-ordered	Symmetrically-ordered
Duplicate keys allowed?	Yes	No*

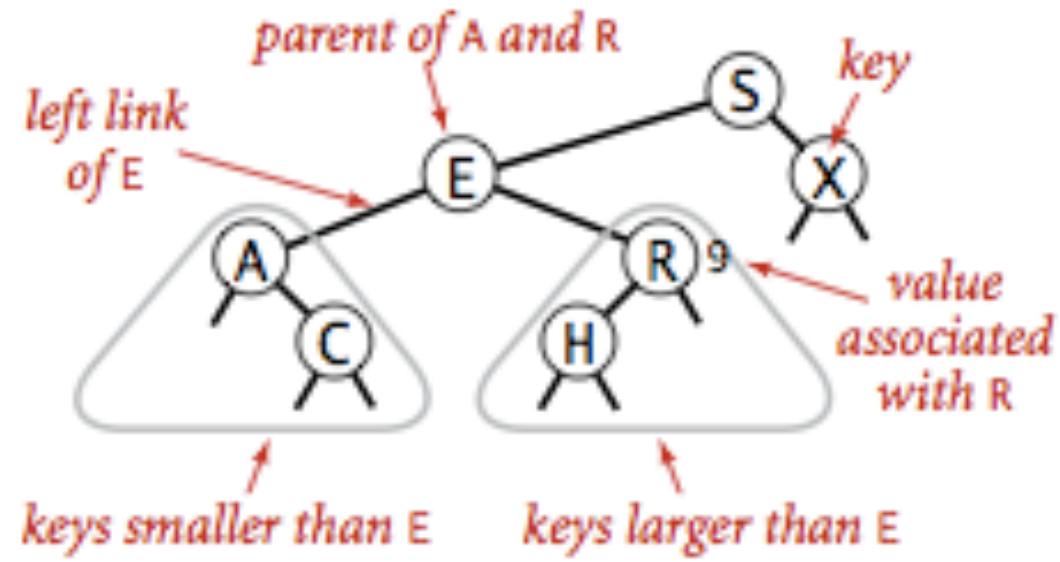
*: when are BSTs used to implement dictionaries.

BST representation of dictionaries

- We will use an inner class Node that is composed by:
  - A Key that is comparable and a `Value`
  - A reference to the root nodes of the left (smaller keys) and right (larger keys) subtrees.
  - Potentially, the total number of nodes in the subtree that has root this node.
- A BST has a reference to a Node `root`.

BST and Node implementation

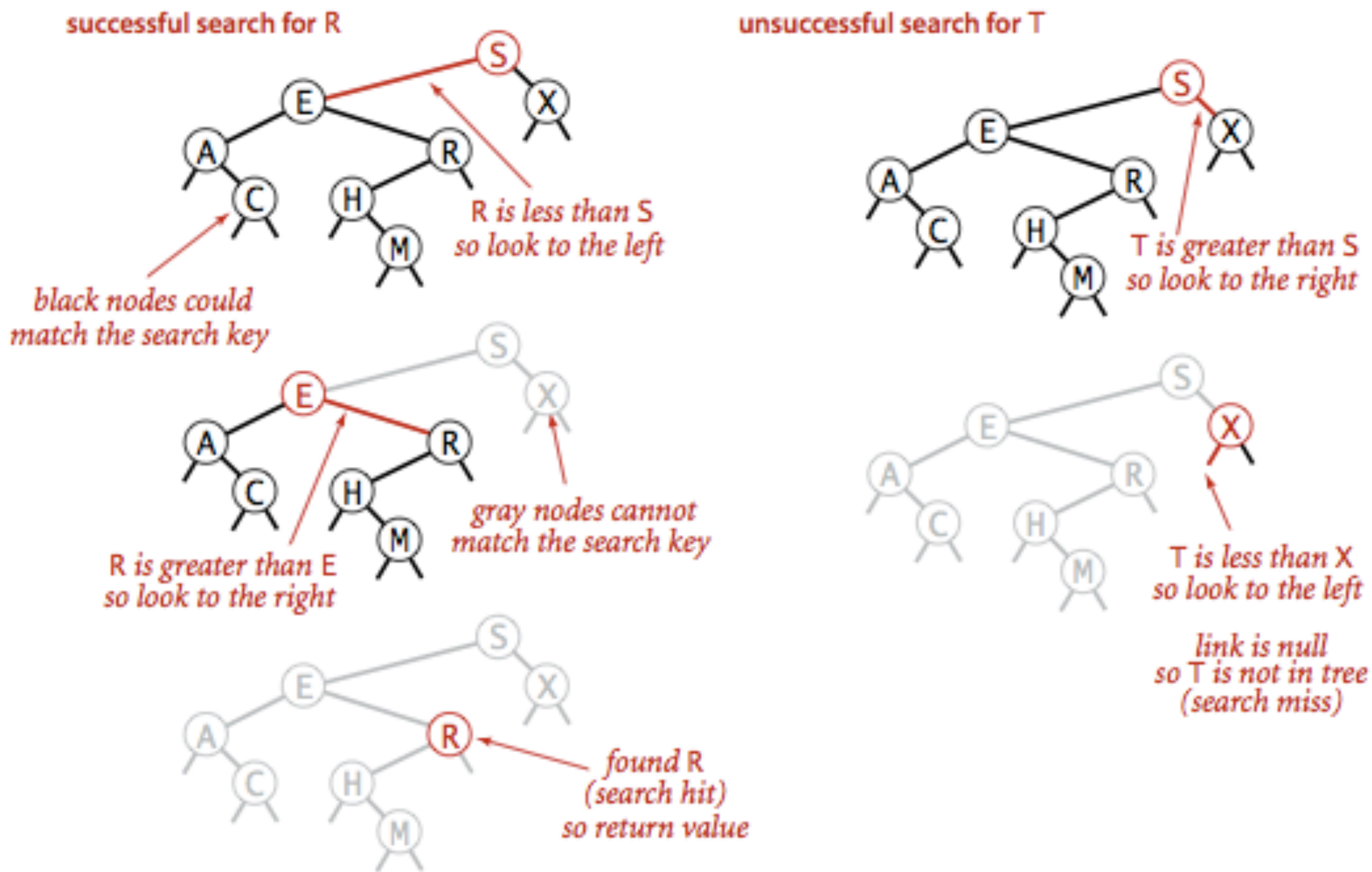
```
public class BST<Key extends Comparable<Key>, Value> {
   private Node root;                // root of BST

   private class Node {
        private Key key;             // sorted by key
        private Value val;           // associated value
        private Node left, right;    // roots of left and right subtrees
        private int size;            // #nodes in subtree rooted at this

        public Node(Key key, Value val, int size) {
            this.key = key;
            this.val = val;
            this.size = size;
        }
    }
```



Search for a key

- If less than key in node go to left subtree.
- If greater than key in node go to right subtree.
- If given key and key at examined node are equal, search hit.
- Return value corresponding to given key, or `null` if no such key.
    - In other implementations, you return the last node you reached.
- Number of compares is equal to the depth of the node + 1.

Search example



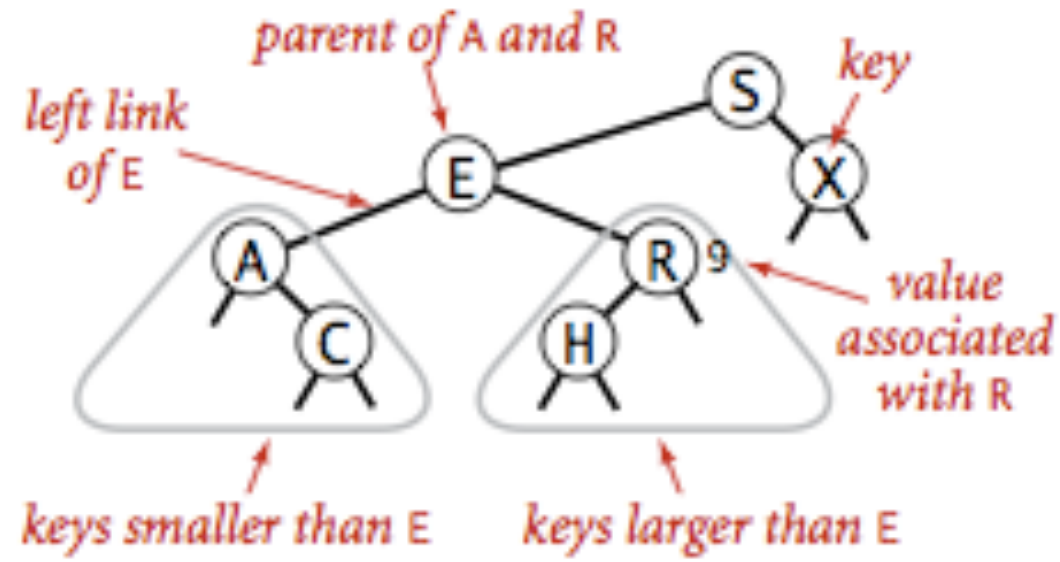
Successful (left) and unsuccessful (right) search in a BST

Search - iterative implementation

```
public Value get(Key key) {
        Node x = root;
        while (x != null) {
                int cmp = key.compareTo(x.key);
                if (cmp < 0)
                        x = x.left;
                else if (cmp > 0)
                        x = x.right;
                else if (cmp == 0)
                        return x.val;
         }
         return null;
}
```

Search - recursive implementation

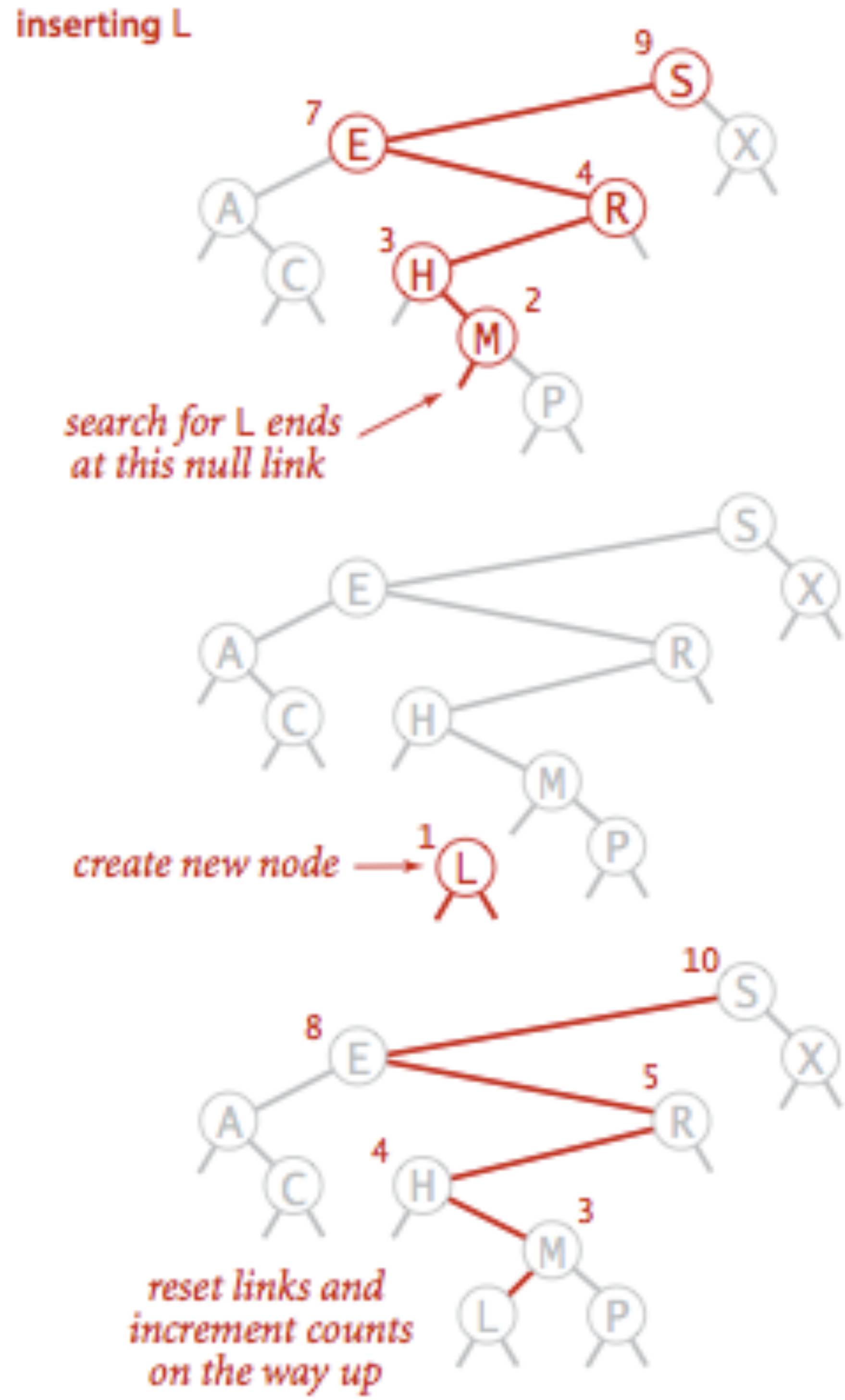
- ```
  public Value get(Key key) {
          return get(root, key);
  }
  ```
- ```
  private Value get(Node x, Key key) {
          if (x == null)
                  return null;
          int cmp = key.compareTo(x.key);
          if (cmp < 0)
              return get(x.left, key);
          else if (cmp > 0)
              return get(x.right, key);
          else
              return x.val;
  }
  ```



Insert

- If less than key in node go left.
- If greater than key in node go right.
- If null, insert.
- If already exists, update value.
- Number of compares is equal to the depth of the node + 1.

Insert example



Insertion into a BST

Insert

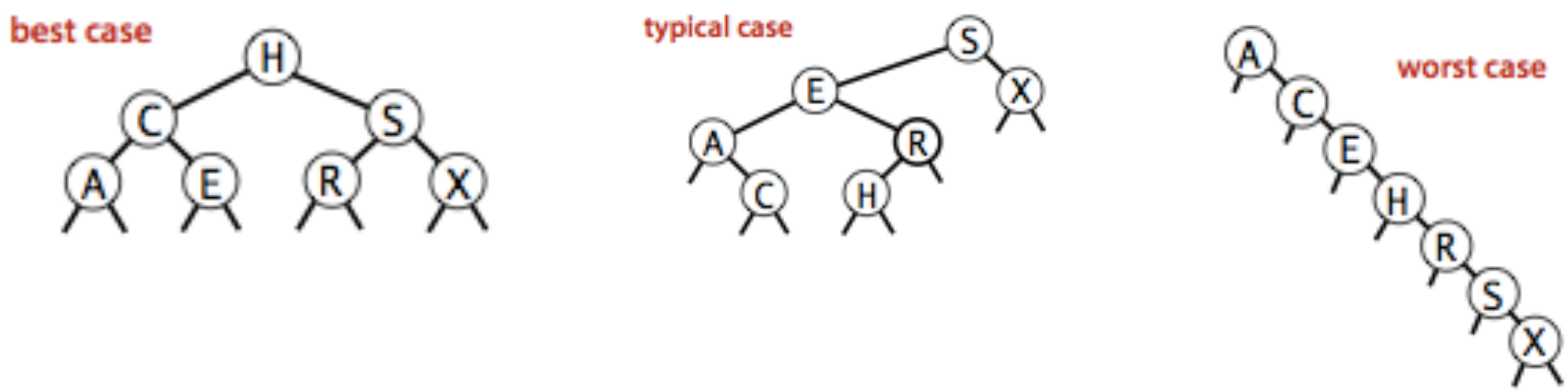
```
public void put(Key key, Value val) {
     root = put(root, key, val);
}
private Node put(Node x, Key key, Value val) {
     if (x == null)
           return new Node(key, val, 1);
     int cmp = key.compareTo(x.key);
     if (cmp < 0)
         x.left = put(x.left, key, val);
     else if (cmp > 0)
         x.right = put(x.right, key, val);
     else
         x.val = val;
     x.size = 1 + size(x.left) + size(x.right);
     return x;
}
```

http://algs4.cs.princeton.edu

3.2 BINARY SEARCH TREE DEMO

Tree shape

- The same set of keys can result to different BSTs based on their order of insertion.
- Number of compares for search/insert is equal to depth of node +1.

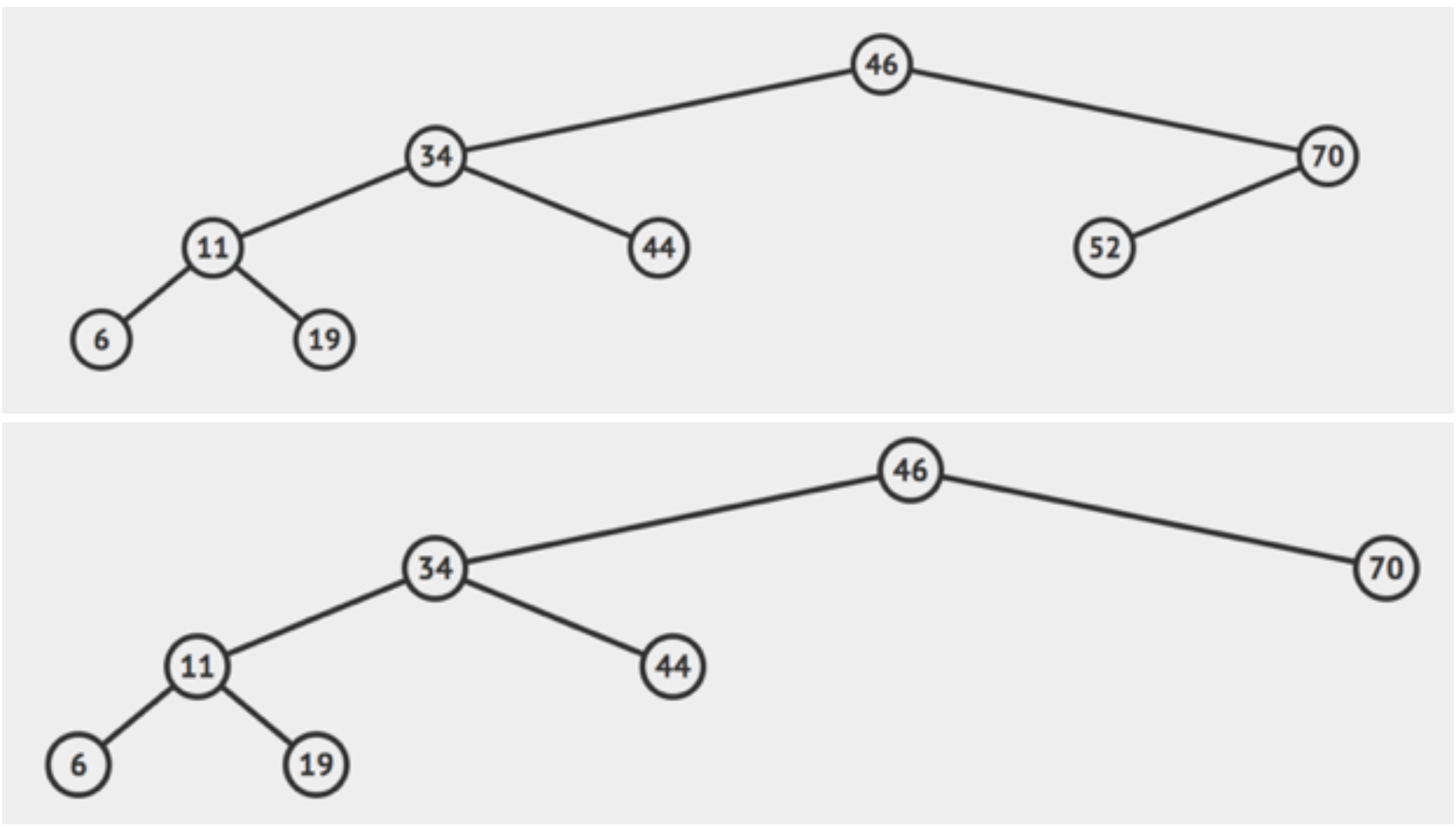


BSTs mathematical analysis

- If n distinct keys are inserted into a BST in random order, the expected number of compares of search/insert is $O(\log n)$.
  - If n distinct keys are inserted into a BST in random order, the expected height of tree is $O(\log n)$. [Reed, 2003].
- Worst case height is n but highly unlikely.
  - Keys would have to come (reversely) sorted!
- All ordered operations in a dictionary implemented with a BST depend on the height of the BST.

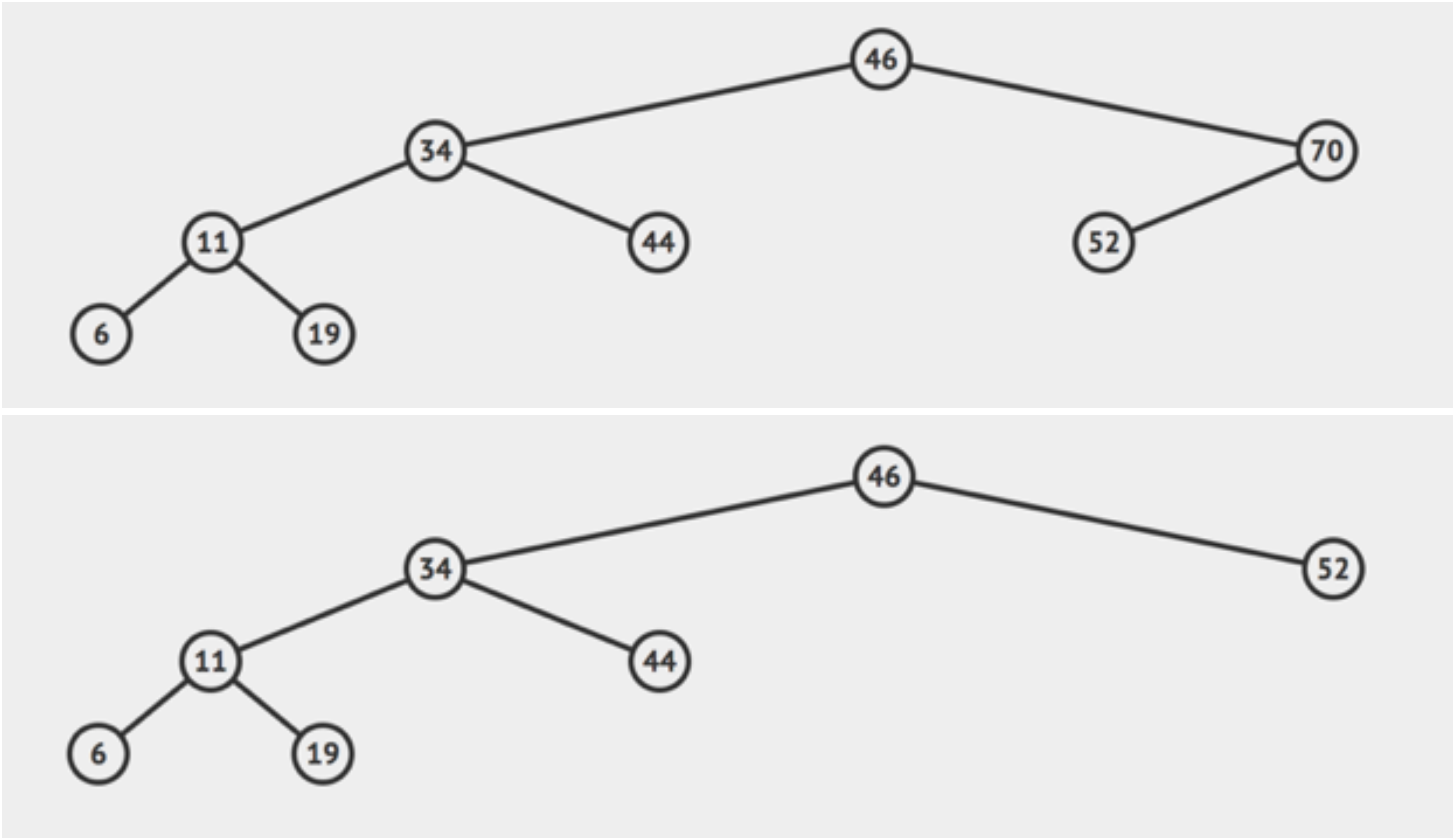
Hibbard deletion: Delete node which is a leaf

- Simply delete node.
- Example: delete 52 locates a node which is a leaf and removes it.



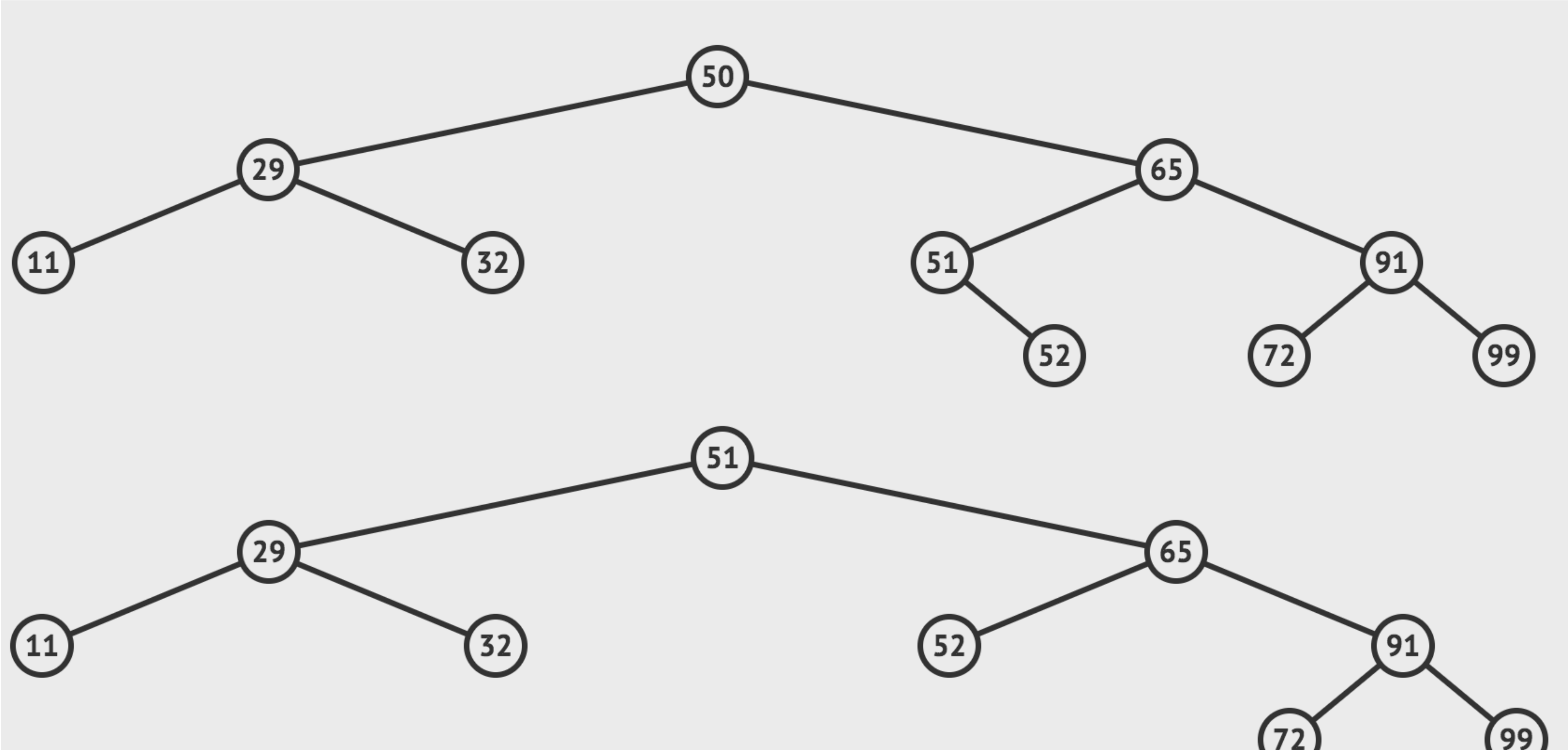
Hibbard deletion: Delete node with one child

- Delete node and replace it with its child.
- Example: delete 70 locates a node which has one child and replaces it with the child.



Hibbard deletion: Delete node with two children

- Delete node and replace it with successor (node with smallest of the larger keys). Move successor's child (if any) where successor was.
- Example: delete 50 locates a node which has two children. Successor is 51.



https://visualgo.net/en/bst

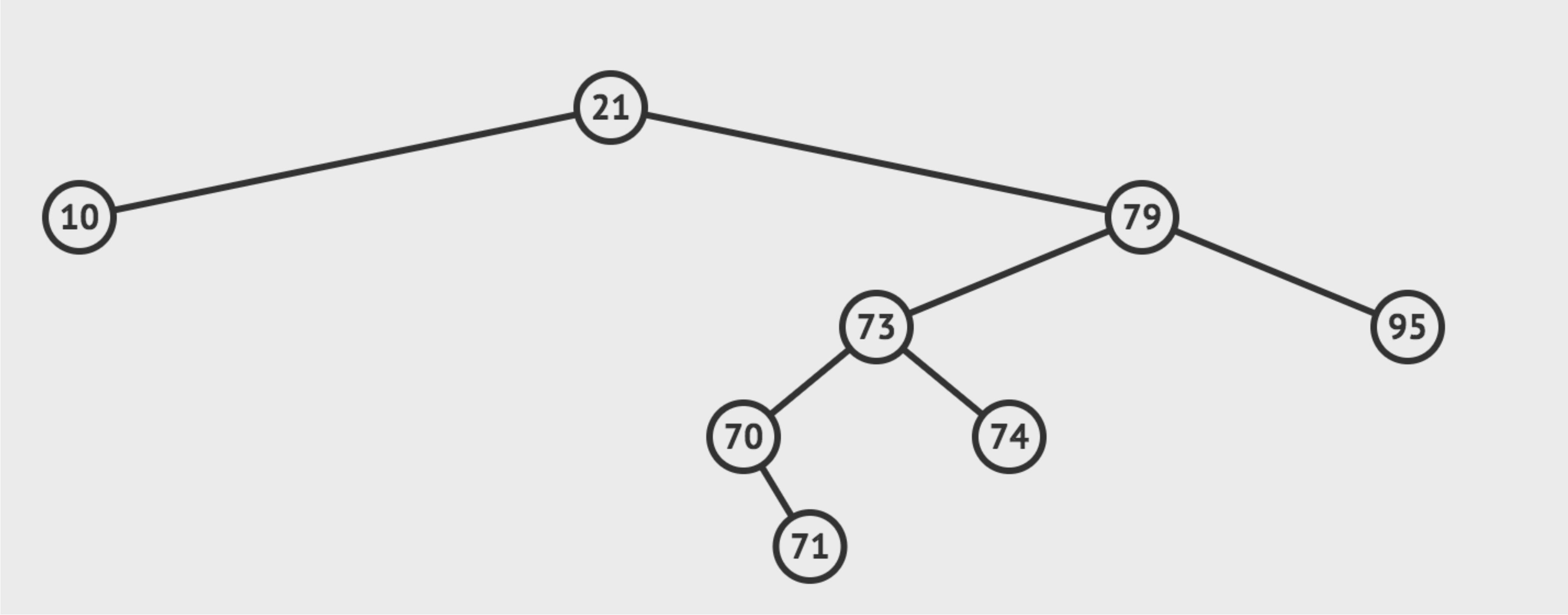
```
public void delete(Key key) {
    root = delete(root, key);
}

 private Node delete(Node x, Key key) {
     if (x == null) return null;

     int cmp = key.compareTo(x.key);
     if (cmp < 0)
         x.left  = delete(x.left,  key);
     else if (cmp > 0)
         x.right = delete(x.right, key);
     else {
         if (x.right == null)
             return x.left;
         if (x.left  == null)
             return x.right;
         Node t = x; //replace with successor
         x = min(t.right);
         x.right = deleteMin(t.right);
         x.left = t.left;
     }
     x.size = size(x.left) + size(x.right) + 1;
     return x;
 }
```

```
/**
 * Removes the smallest key and associated value from the symbo
 *
 * @throws NoSuchElementException if the symbol table is empty
 */
public void deleteMin() {
    if (isEmpty()) throw new NoSuchElementException();
    root = deleteMin(root);
}

private Node deleteMin(Node x) {
    if (x.left == null) return x.right;
    x.left = deleteMin(x.left);
    x.size = size(x.left) + size(x.right) + 1;
    return x;
}
```

```
private Node min(Node x) {
    if (x.left == null) return x;
    else                return min(x.left);
}
```

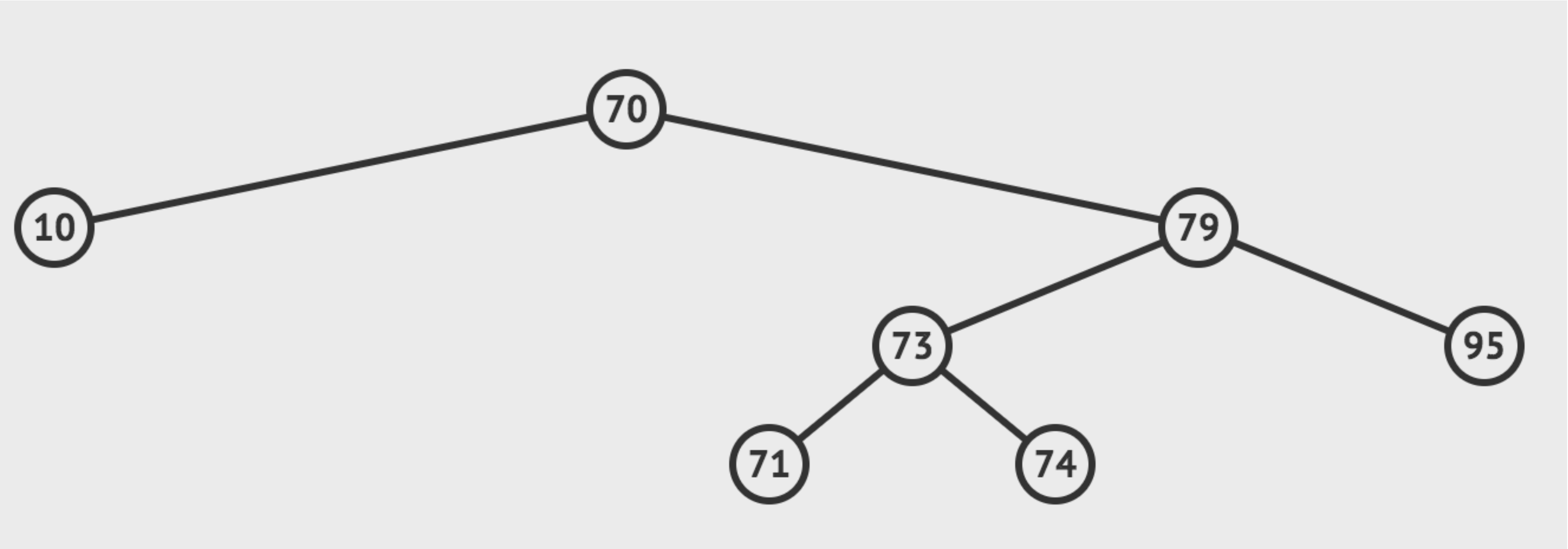
Practice Time

- Delete the node 21 following Hibbard's deletion



Answer

- Delete the node 21 following Hibbard's deletion



Hibbard's deletion

- Unsatisfactory solution. If we were to perform many insertions and deletions the BST ends up being not symmetric and skewed to the left.
  - Extremely complicated analysis, but average cost of deletion ends up being $\sqrt{n}$. Let's simplify things by saying it stays $O(\log n)$.
  - No one has proven that alternating between the predecessor and successor will fix this.
- Hibbard devised the algorithm in 1962. Still no algorithm for efficient deletion in Binary Search Trees!
- Overall, BSTs can have $O(n)$ worst-case for search, insert, and delete. We want to do better (see future lectures).

Lecture 19: Binary Search Trees

- Binary Search Trees

The story so far

- The symbol table/dictionary is a fundamental data type.
- Naive implementations (arrays/linked lists sorted or unsorted) are way too slow.
- Binary search trees work well in the average case, but can grow too tall and imbalanced in the worst case.
- Question of the day: How to balance search trees?

Order of growth for symbol table/dictionary operations

	Worst case			Average case		
	Search	**Insert**	**Delete**	**Search**	**Insert**	**Delete**
BST	n	n	n	$\log n$	$\log n$	$\sqrt{n}$
Goal	$\log n$	$\log n$	$\log n$	$\log n$	$\log n$	$\log n$

Readings:

- Recommended Textbook: Chapters 3.2 (Pages 396–414)
- Website:
  - https://algs4.cs.princeton.edu/32bst/
- Visualization:
  - https://visualgo.net/en/bst

Practice Problems:

- In-class worksheet