

# CS062

## DATA STRUCTURES AND ADVANCED PROGRAMMING

### 27: Shortest Paths

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she/her/hers



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he/him/his

## Lecture 27: Shortest Paths

- ▶ Introduction to Shortest Paths
- ▶ API
- ▶ Properties
- ▶ Dijkstra's Algorithm
- ▶ Belman-Ford Algorithm

## Edge-weighted digraph

- ▶ **Edge-weighted digraph**: a digraph where we associate weights or costs with each edge.

### edge-weighted digraph

4→5	0.35
5→4	0.35
4→7	0.37
5→7	0.28
7→5	0.28
5→1	0.32
0→4	0.38
0→2	0.26
7→3	0.39
1→3	0.29
2→7	0.34
6→2	0.40
3→6	0.52
6→0	0.58
6→4	0.93



## Shortest Paths

- ▶ **Shortest path from vertex  $s$  to vertex  $t$ :** a directed path from  $s$  to  $t$  with the property that no other such path has a lower weight (total weight sum of edges it consists of).

### edge-weighted digraph

4→5	0.35
5→4	0.35
4→7	0.37
5→7	0.28
7→5	0.28
5→1	0.32
0→4	0.38
0→2	0.26
7→3	0.39
1→3	0.29
2→7	0.34
6→2	0.40
3→6	0.52
6→0	0.58
6→4	0.93



### shortest path from 0 to 6

0→2	0.26
2→7	0.34
7→3	0.39
3→6	0.52

An edge-weighted digraph and a shortest path

## Shortest Path variants

- ▶ **Single source**: from one vertex  $s$  to every other vertex.
- ▶ **Single sink**: from every vertex to one vertex  $t$ .
- ▶ **Source-sink**: from one vertex  $s$  to another vertex  $t$ .
- ▶ **All pairs**: from every vertex to every other vertex.
- ▶ What version is there in your navigation app?

## Shortest Paths Assumptions

- ▶ Not all vertices need to be reachable.
  - ▶ We will assume so in this lecture.
- ▶ Weights are non-negative.
  - ▶ There are algorithms that can handle negative weights.
- ▶ Shortest paths are not necessarily unique but they are simple.

## Lecture 27: Shortest Paths

- ▶ Introduction to Shortest Paths
- ▶ API
- ▶ Properties
- ▶ Dijkstra's Algorithm
- ▶ Belman-Ford Algorithm

## Weighted directed edge API

- ▶ `public class DirectedEdge`
  - ▶ `DirectedEdge(int v, int w, double weight)`
    - ▶ Constructs a weighted edge from  $v$  to  $w$  ( $v \rightarrow w$ ) with the provided `weight`.
  - ▶ `int from()`
    - ▶ Returns vertex source of this edge.
  - ▶ `int to()`
    - ▶ Returns vertex destination of this edge.
  - ▶ `double weight()`
    - ▶ Returns weight of this edge.
  - ▶ `String toString()`
    - ▶ Returns the string representation of this edge.



## Weighted directed edge in Java

```
public class DirectedEdge {
    private final int v;
    private final int w;
    private final double weight;

    public DirectedEdge(int v, int w, double weight) {
        this.v = v;
        this.w = w;
        this.weight = weight;
    }

    public int from() {
        return v;
    }

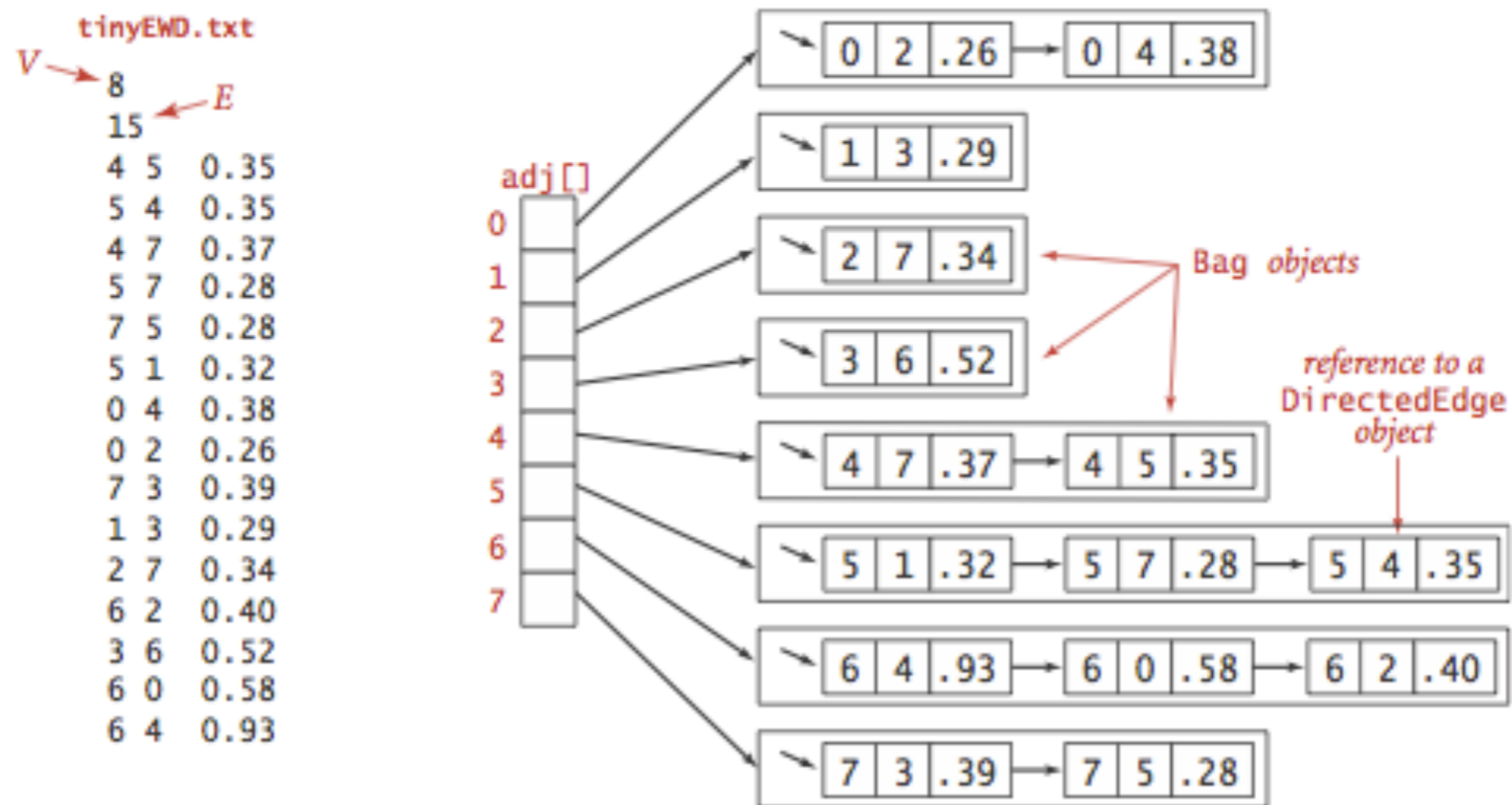
    public int to() {
        return w;
    }

    public double weight() {
        return weight;
    }
}
```

# Edge-weighted digraph API

- ▶ `public class` EdgeWeightedDigraph
  - ▶ `EdgeWeightedDigraph(int v)`
    - ▶ Constructs an edge-weighted digraph with `v` vertices.
  - ▶ `void` `addEdge(DirectedEdge e)`
    - ▶ Add weighted directed edge `e`.
  - ▶ `Iterable<DirectedEdge>` `adj(int v)`
    - ▶ Returns edges adjacent from `v`.
  - ▶ `int` `V()`
    - ▶ Returns number of vertices.
  - ▶ `int` `E()`
    - ▶ Returns number of edges.
  - ▶ `Iterable<DirectedEdge>` `edges()`
    - ▶ Returns all edges.

# Edge-weighted digraph adjacency list representation



Edge-weighted digraph representation

## Edge-weighted digraph in Java

```
public class EdgeWeightedDigraph {
    private final int V;           // number of vertices in this digraph
    private int E;                 // number of edges in this digraph
    private Bag<DirectedEdge>[] adj; // adj[v] = adjacency list for vertex v

    public EdgeWeightedDigraph(int V) {
        this.V = V;
        this.E = 0;
        adj = (Bag<DirectedEdge>[]) new Bag[V];
        for (int v = 0; v < V; v++)
            adj[v] = new Bag<DirectedEdge>();
    }

    public void addEdge(DirectedEdge e) {
        int v = e.from();
        int w = e.to();
        adj[v].add(e);
        E++;
    }

    public Iterable<DirectedEdge> adj(int v) {
        return adj[v];
    }
}
```

## Single-source shortest path API

- ▶ **Goal**: find shortest path from  $s$  to every other vertex in the digraph.
- ▶ **public class** SP
  - ▶ `SP(EdgeWeightedDigraph G, int s)`
    - ▶ Shortest paths from  $s$  in digraph  $G$ .
  - ▶ **double** `distTo(int v)`
    - ▶ Length of shortest path from  $s$  to  $v$ .
  - ▶ `Iterable<DirectedEdge> pathTo(int v)`
    - ▶ Returns edges along the shortest path from  $s$  to  $v$ .
  - ▶ **boolean** `hasPathTo(int v)`
    - ▶ Returns whether there is a path from  $s$  to  $v$ .

## Lecture 27: Shortest Paths

- ▶ Introduction to Shortest Paths
- ▶ API
- ▶ **Properties**
- ▶ Dijkstra's Algorithm
- ▶ Belman-Ford Algorithm

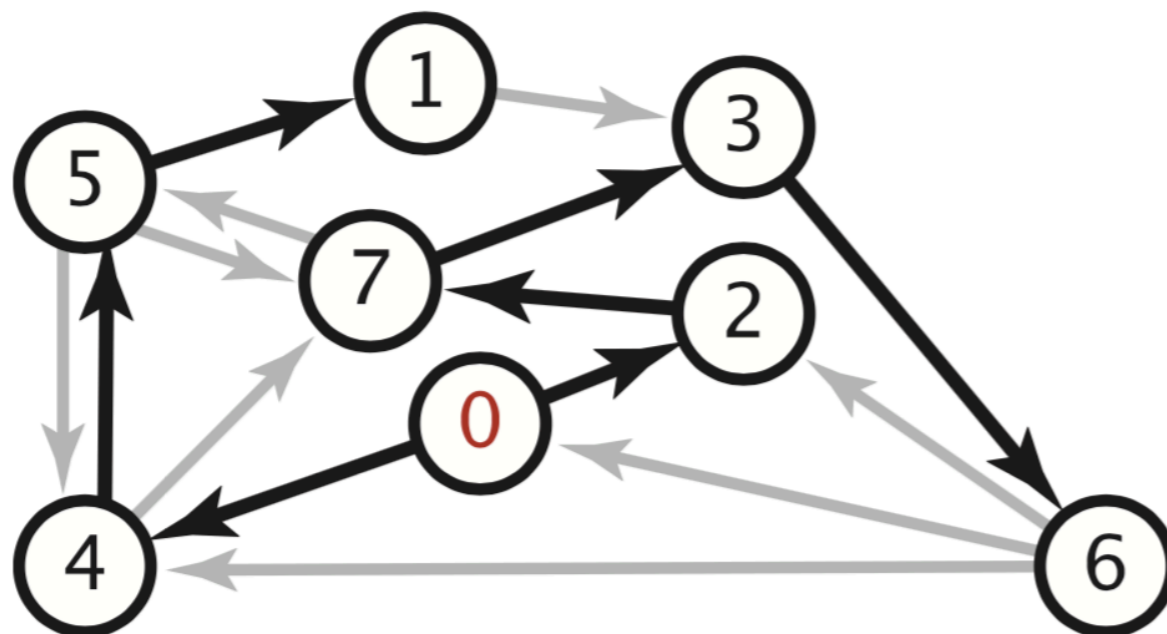
## Data structures for single-source shortest paths

- ▶ **Goal**: find shortest path from  $s$  to every other vertex in the digraph.
- ▶ **Shortest-paths tree (SPT)**: a subgraph containing  $s$  and all the vertices reachable from  $s$  that forms a directed tree rooted at  $s$  such that every tree path in the SPT is a shortest path in the digraph.
- ▶ Representation of shortest paths with two vertex-indexed arrays.
  - ▶ **Edges on the shortest-paths tree**:  $\text{edgeTo}[v]$  is the last edge on a shortest path from  $s$  to  $v$ .
  - ▶ **Distance to the source**:  $\text{distTo}[v]$  is the length of the shortest path from  $s$  to  $v$ .

# PROPERTIES

```
public Iterable<DirectedEdge> pathTo(int v) {  
    Stack<DirectedEdge> path = new Stack<DirectedEdge>();  
    for (DirectedEdge e = edgeTo[v]; e != null; e = edgeTo[e.from()]) {  
        path.push(e);  
    }  
    return path;  
}
```

4->5	0.35
5->4	0.35
4->7	0.37
5->7	0.28
7->5	0.28
5->1	0.32
0->4	0.38
0->2	0.26
7->3	0.39
1->3	0.29
2->7	0.34
6->2	0.40
3->6	0.52
6->0	0.58
6->4	0.93



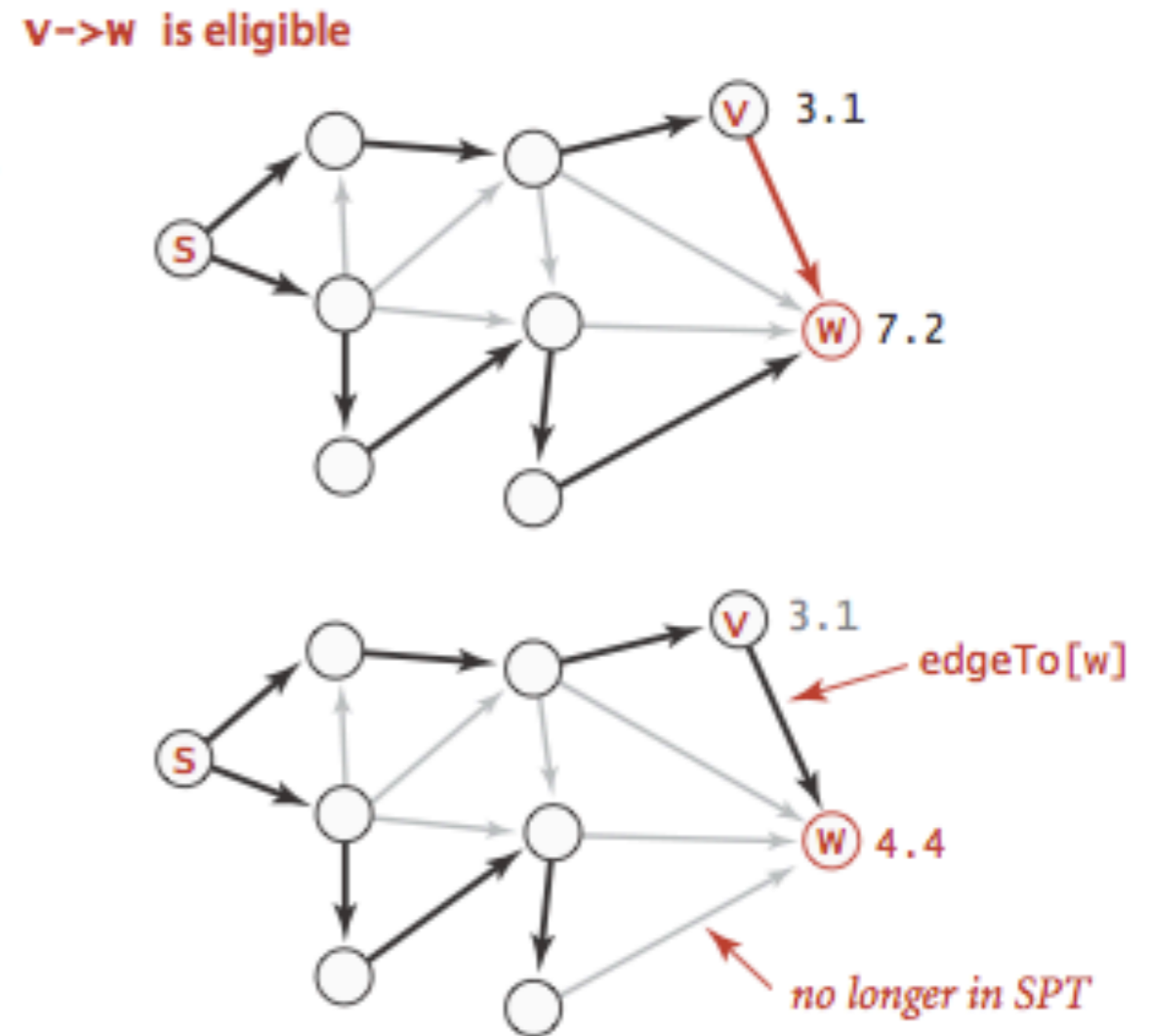
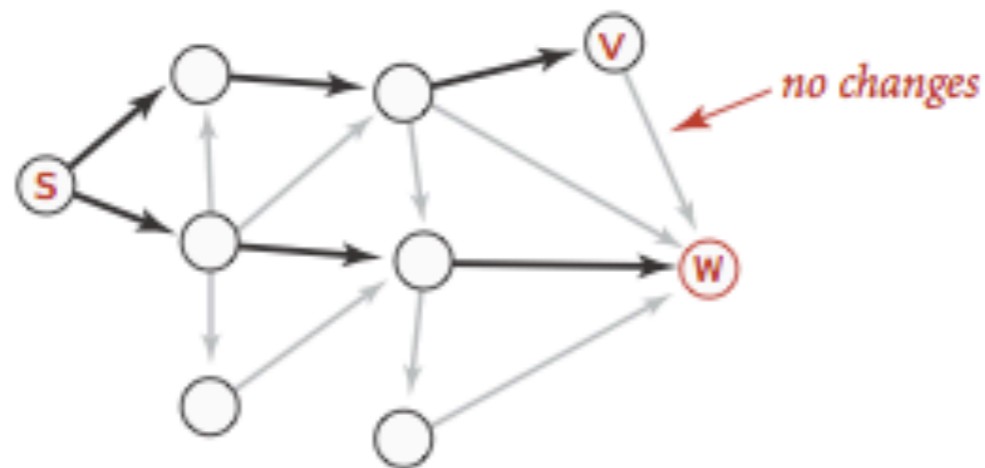
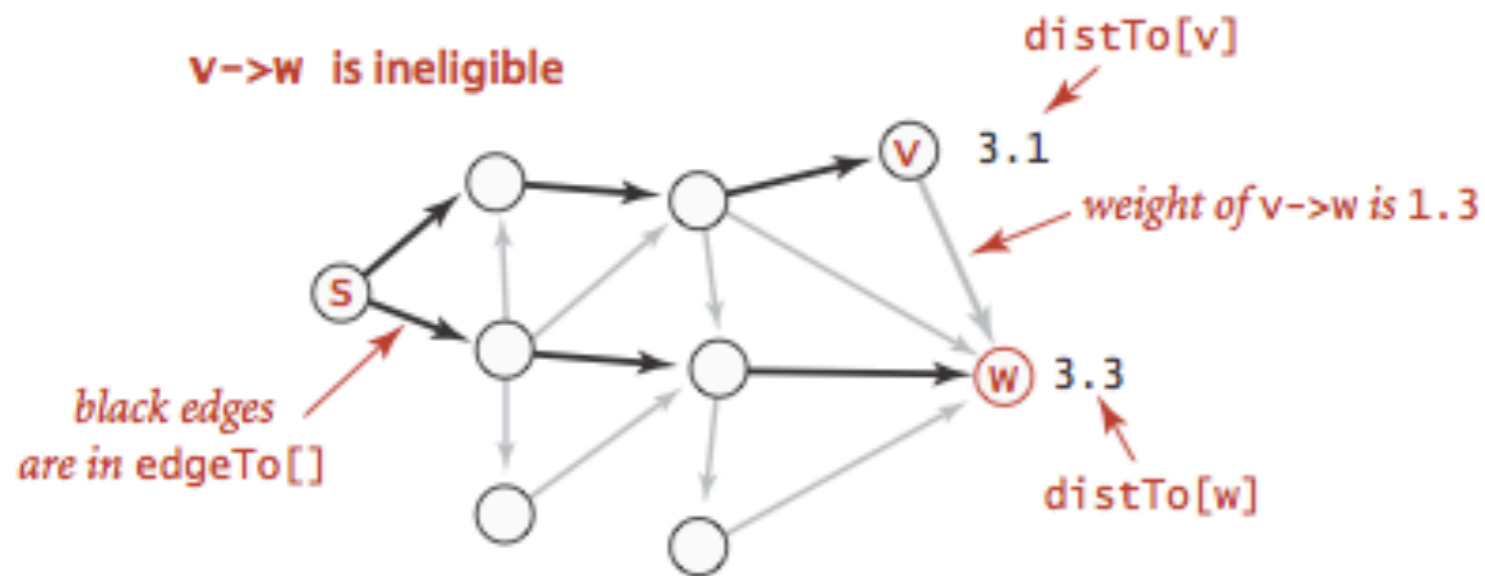
	edgeTo[]	distTo[]
0	null	0
1	5->1 0.32	1.05
2	0->2 0.26	0.26
3	7->3 0.39	0.99
4	0->4 0.38	0.38
5	4->5 0.35	0.73
6	3->6 0.52	1.51
7	2->7 0.34	0.60



## Edge relaxation

- ▶ Relax edge  $e = v \rightarrow w$ 
  - ▶  $\text{distTo}[v]$  is the length of the shortest **known** path from  $S$  to  $v$ .
  - ▶  $\text{distTo}[w]$  is the length of the shortest **known** path from  $S$  to  $w$ .
  - ▶  $\text{edgeTo}[w]$  is the last edge on shortest **known** path from  $S$  to  $w$ .
  - ▶ If  $e = v \rightarrow w$  yields shorter path to  $w$ , update  $\text{distTo}[w]$  and  $\text{edgeTo}[w]$ .

# Edge relaxation



## Edge relaxation implementation

```
private void relax(DirectedEdge e) {  
    int v = e.from(), w = e.to();  
    if (distTo[w] > distTo[v] + e.weight()) {  
        distTo[w] = distTo[v] + e.weight();  
        edgeTo[w] = e;  
    }  
}
```

## Framework for shortest-paths algorithm

- ▶ Generic algorithm to compute a SPT from  $s$ 
  - ▶  $\text{distTo}[v] = \infty$  for each vertex  $v$ .
  - ▶  $\text{edgeTo}[v] = \text{null}$  for each vertex  $v$ .
  - ▶  $\text{distTo}[s] = 0$ .
  - ▶ Repeat until done:
    - ▶ Relax any edge.
- ▶  $\text{distTo}[v]$  is the length of a simple path from  $s$  to  $v$ .
- ▶  $\text{distTo}[v]$  does not increase.

## Lecture 27: Shortest Paths

- ▶ Introduction to Shortest Paths
- ▶ API
- ▶ Properties
- ▶ **Dijkstra's Algorithm**
- ▶ Belman-Ford Algorithm



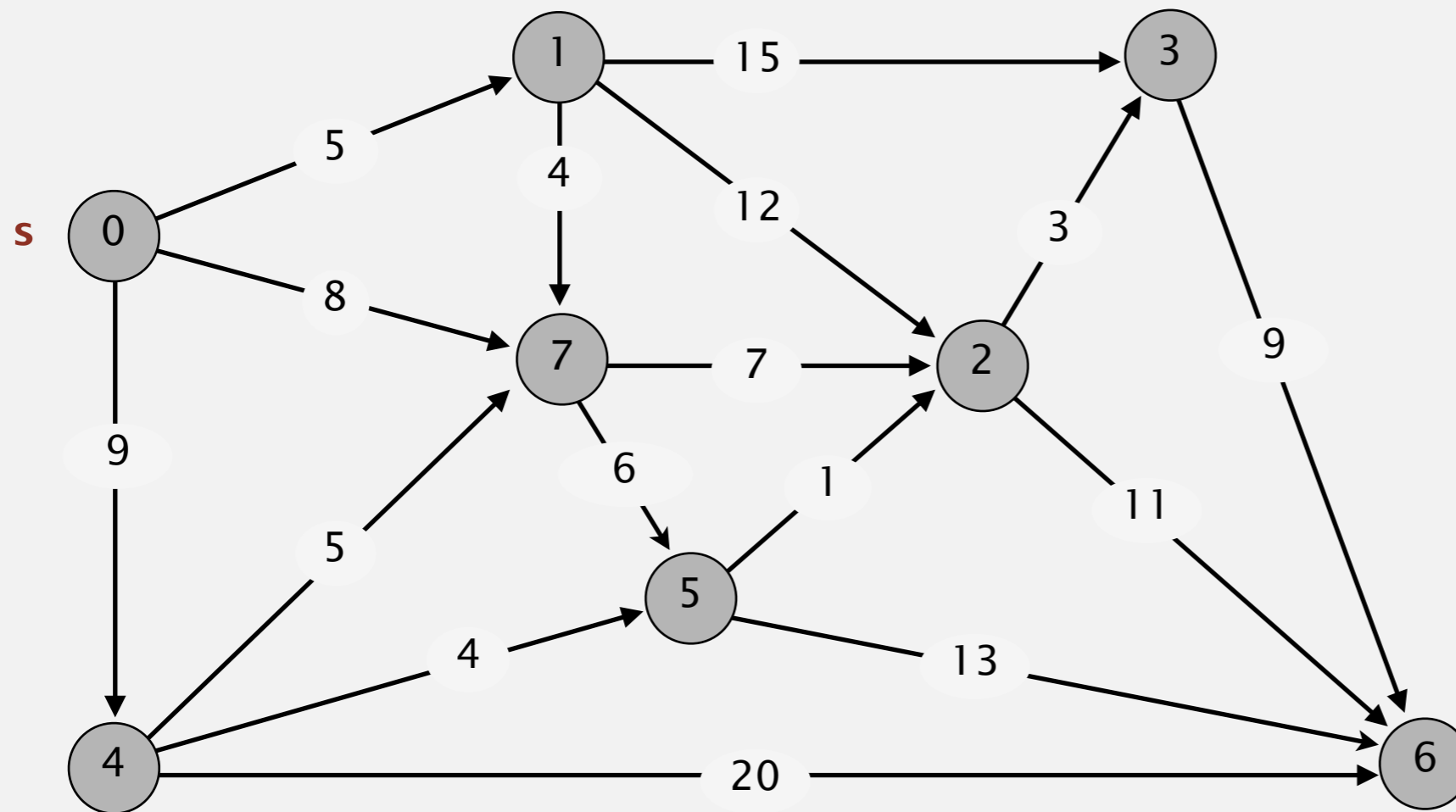
<http://algs4.cs.princeton.edu>

## DIJKSTRA'S ALGORITHM DEMO

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# Dijkstra's algorithm demo

- Consider vertices in increasing order of distance from  $s$  (non-tree vertex with the lowest  $\text{distTo}[]$  value).
- Add vertex to tree and relax all edges adjacent from that vertex.



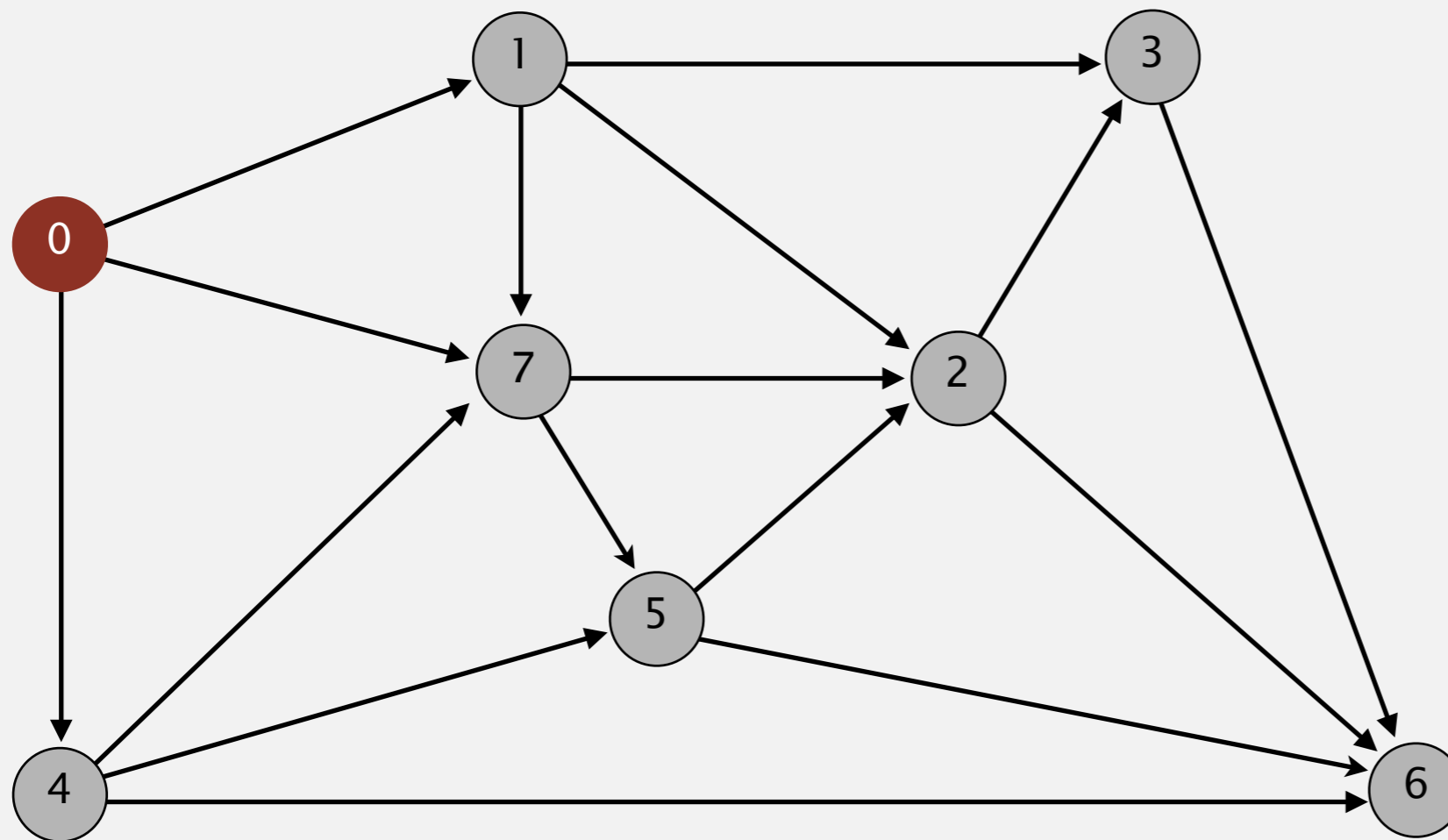
0→1	5.0
0→4	9.0
0→7	8.0
1→2	12.0
1→3	15.0
1→7	4.0
2→3	3.0
2→6	11.0
3→6	9.0
4→5	4.0
4→6	20.0
4→7	5.0
5→2	1.0
5→6	13.0
7→5	6.0
7→2	7.0

**an edge-weighted digraph**

# Dijkstra's algorithm demo

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- Consider vertices in increasing order of distance from  $s$  (non-tree vertex with the lowest  $\text{distTo}[]$  value).
- Add vertex to tree and relax all edges adjacent from that vertex.



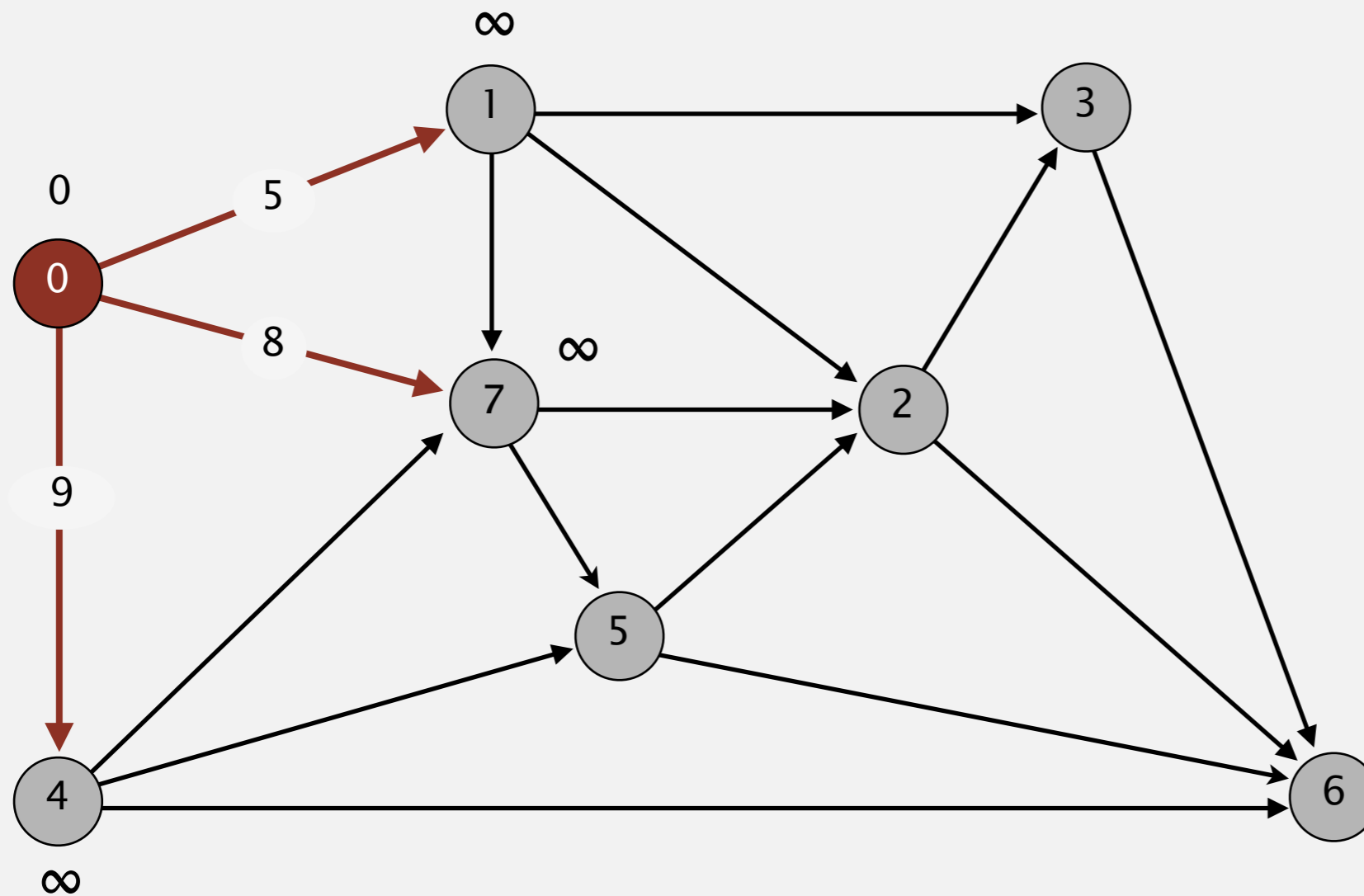
$v$	$\text{distTo}[]$	$\text{edgeTo}[]$
→ 0	0.0	-
1		
2		
3		
4		
5		
6		
7		

choose source vertex 0



# Dijkstra's algorithm demo

- Consider vertices in increasing order of distance from  $s$  (non-tree vertex with the lowest  $\text{distTo}[]$  value).
- Add vertex to tree and relax all edges adjacent from that vertex.

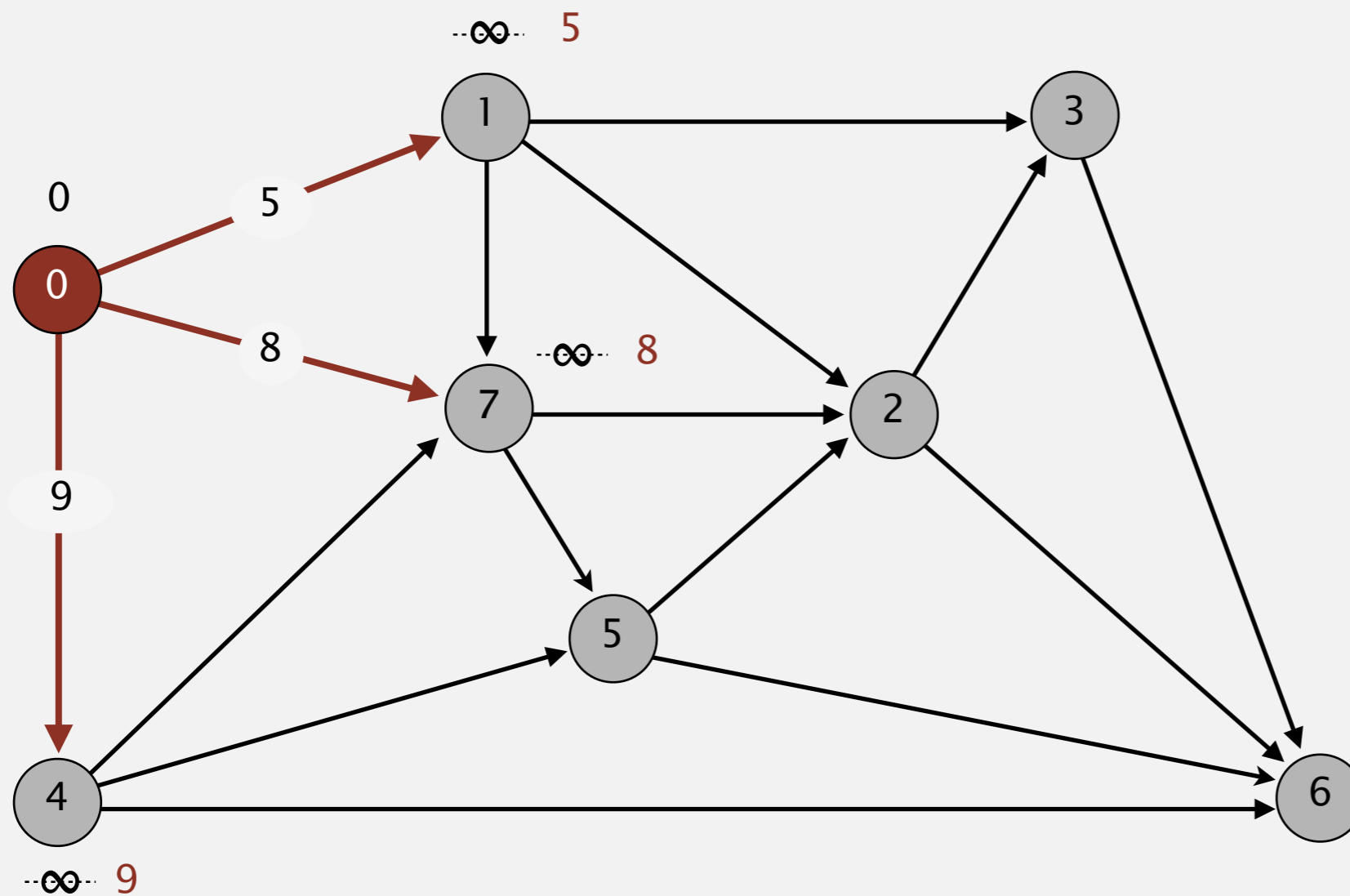


$v$	$\text{distTo}[]$	$\text{edgeTo}[]$
→ 0	0.0	-
1		
2		
3		
4		
5		
6		
7		

relax all edges adjacent from 0

# Dijkstra's algorithm demo

- Consider vertices in increasing order of distance from  $s$  (non-tree vertex with the lowest  $\text{distTo}[]$  value).
- Add vertex to tree and relax all edges adjacent from that vertex.

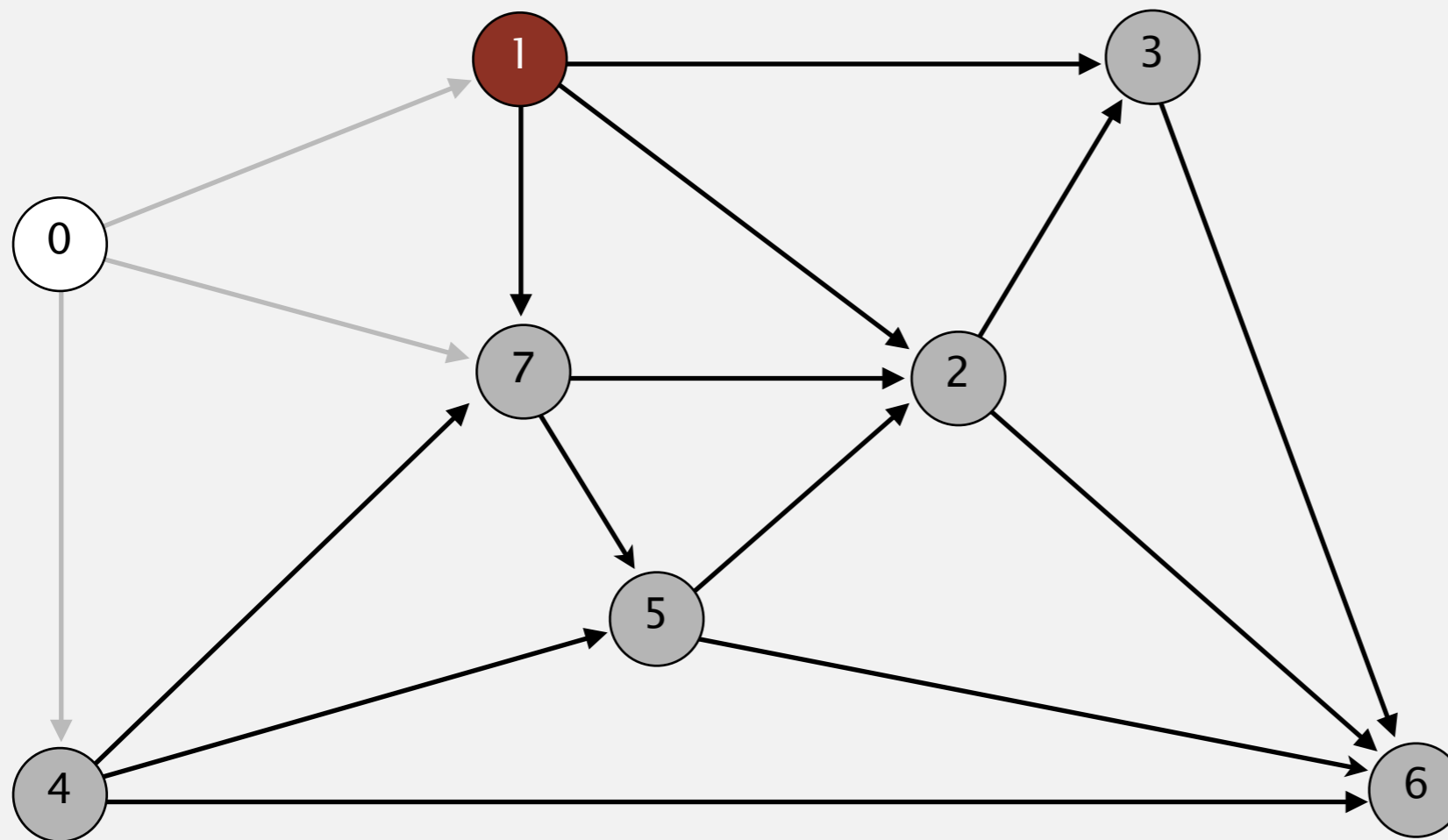


$v$	$\text{distTo}[]$	$\text{edgeTo}[]$
→ 0	0.0	-
1	5.0	0→1
2		
3		
4	9.0	0→4
5		
6		
7	8.0	0→7

relax all edges adjacent from 0

# Dijkstra's algorithm demo

- Consider vertices in increasing order of distance from  $s$  (non-tree vertex with the lowest  $\text{distTo}[]$  value).
- Add vertex to tree and relax all edges adjacent from that vertex.



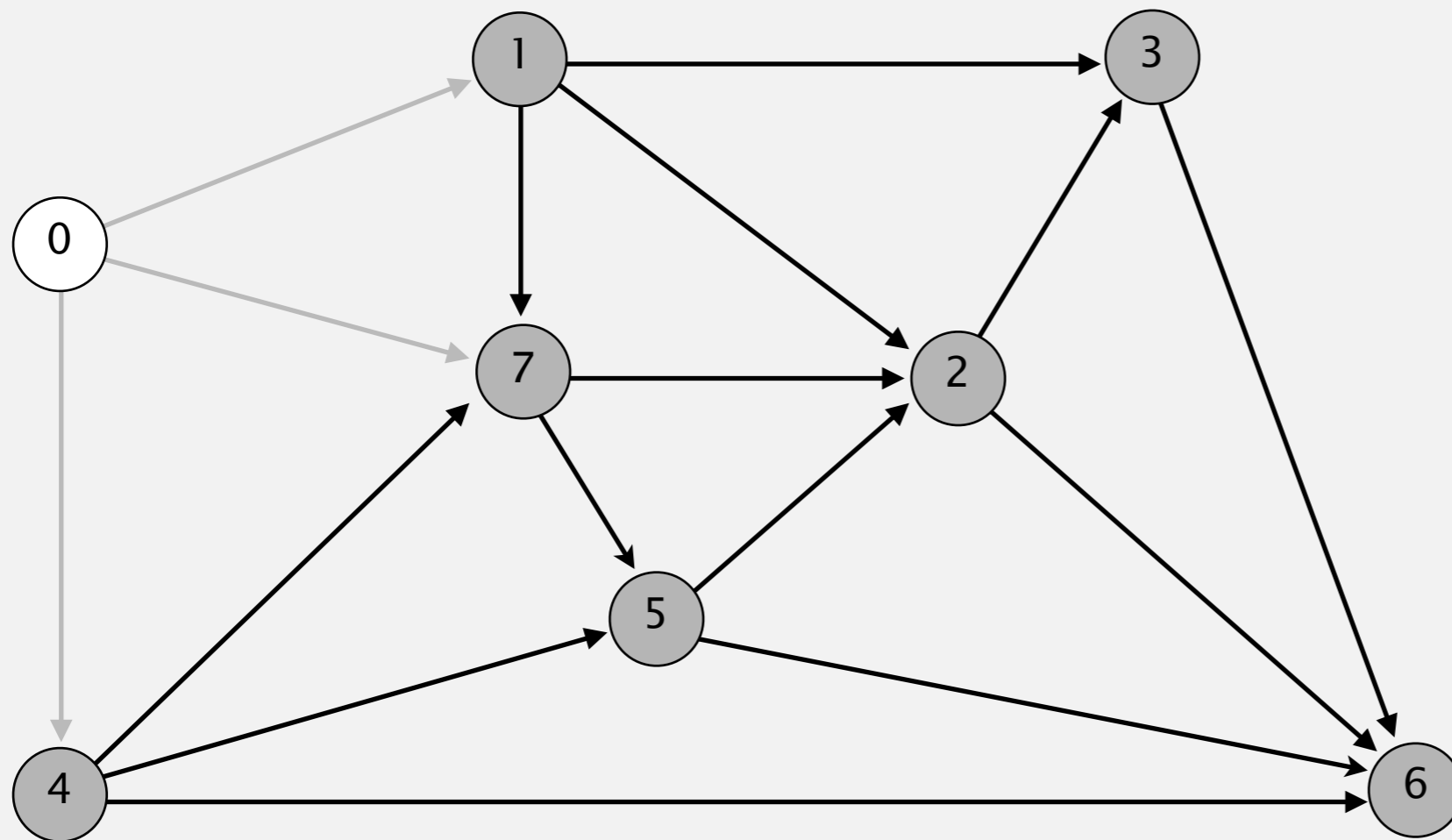
$v$	$\text{distTo}[]$	$\text{edgeTo}[]$
0	0.0	-
→ 1	5.0	0→1
2		
3		
4	9.0	0→4
5		
6		
7	8.0	0→7

choose vertex 1

# Dijkstra's algorithm demo

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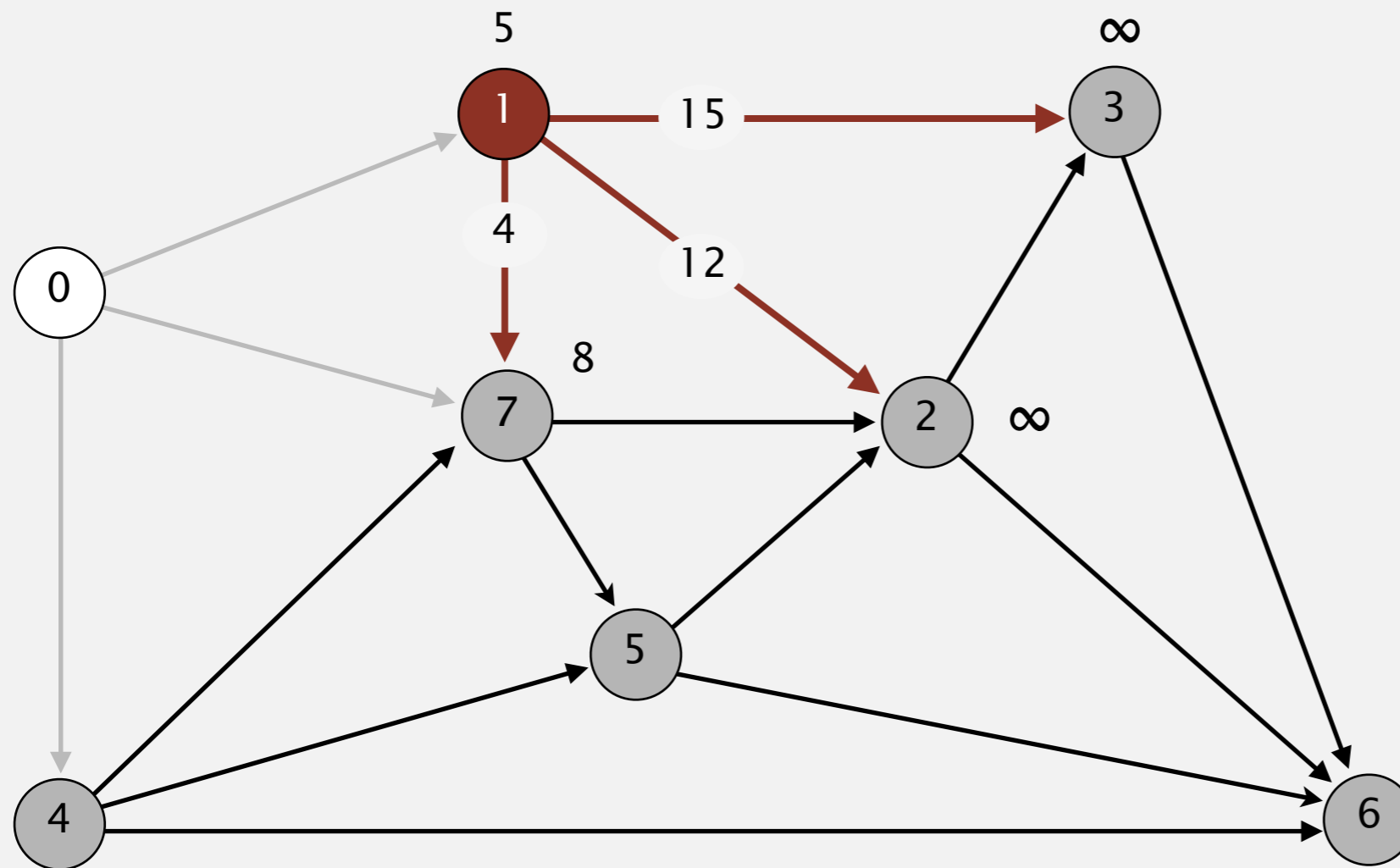
- Consider vertices in increasing order of distance from  $s$  (non-tree vertex with the lowest  $\text{distTo}[]$  value).
- Add vertex to tree and relax all edges adjacent from that vertex.



$v$	$\text{distTo}[]$	$\text{edgeTo}[]$
0	0.0	-
1	5.0	0→1
2		
3		
4	9.0	0→4
5		
6		
7	8.0	0→7

# Dijkstra's algorithm demo

- Consider vertices in increasing order of distance from  $s$  (non-tree vertex with the lowest  $\text{distTo}[]$  value).
- Add vertex to tree and relax all edges adjacent from that vertex.

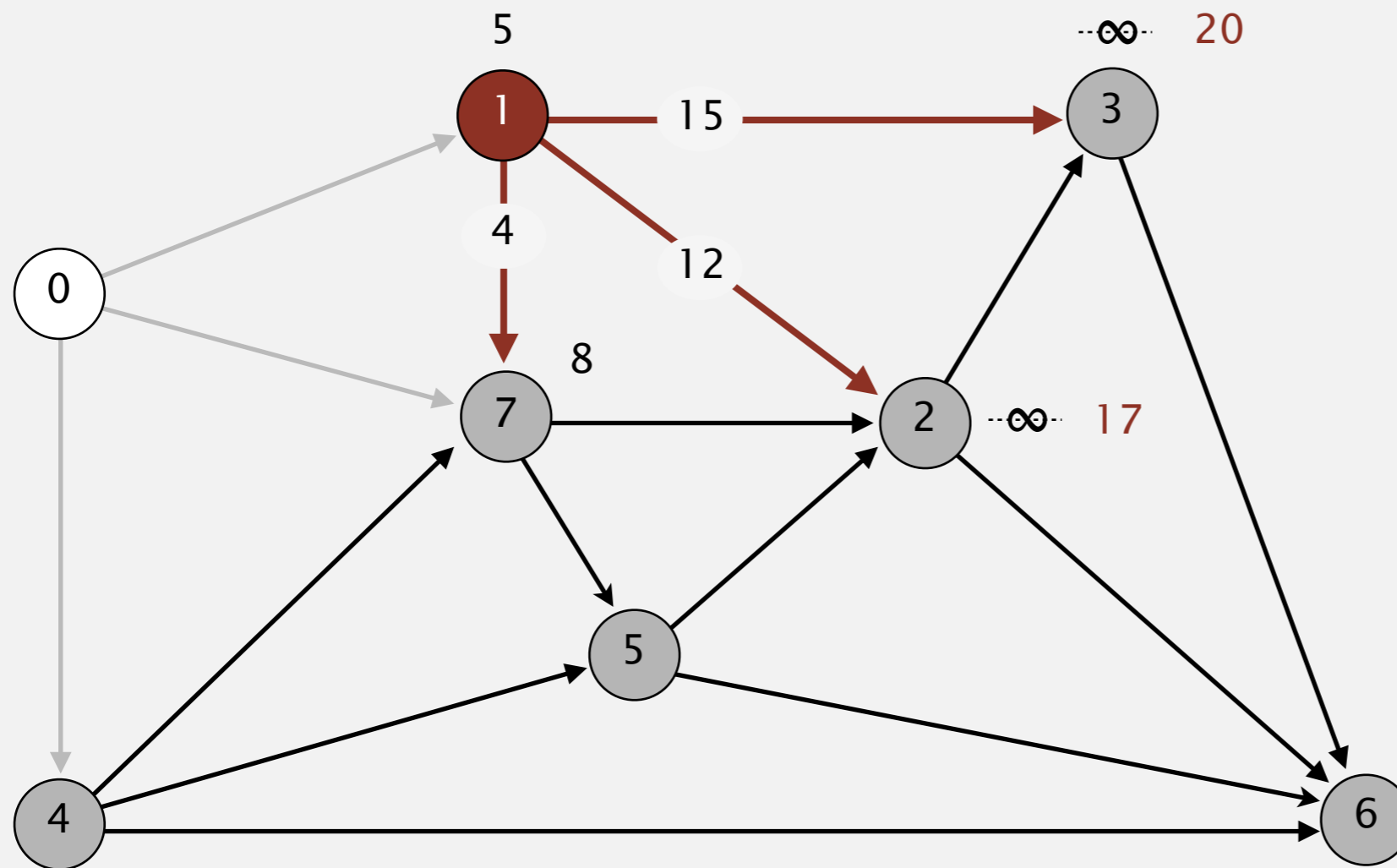


$v$	$\text{distTo}[]$	$\text{edgeTo}[]$
0	0.0	-
→ 1	5.0	0→1
2		
3		
4	9.0	0→4
5		
6		
7	8.0	0→7

relax all edges adjacent from 1

# Dijkstra's algorithm demo

- Consider vertices in increasing order of distance from  $s$  (non-tree vertex with the lowest  $\text{distTo}[]$  value).
- Add vertex to tree and relax all edges adjacent from that vertex.



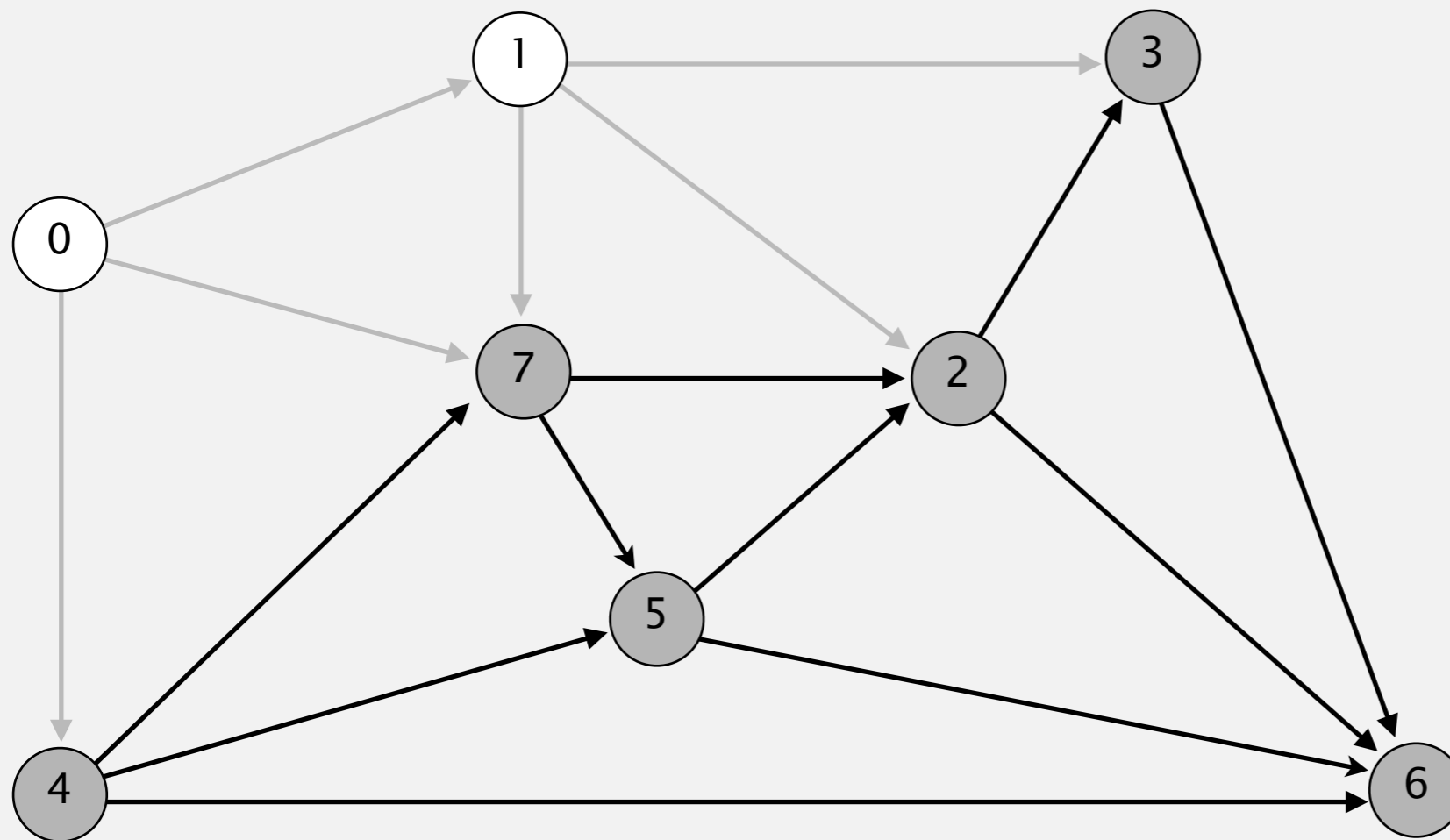
$v$	$\text{distTo}[]$	$\text{edgeTo}[]$
0	0.0	-
→ 1	5.0	0→1
2	17.0	1→2
3	20.0	1→3
4	9.0	0→4
5		
6		
7	8.0 ✓	0→7

relax all edges adjacent from 1

# Dijkstra's algorithm demo

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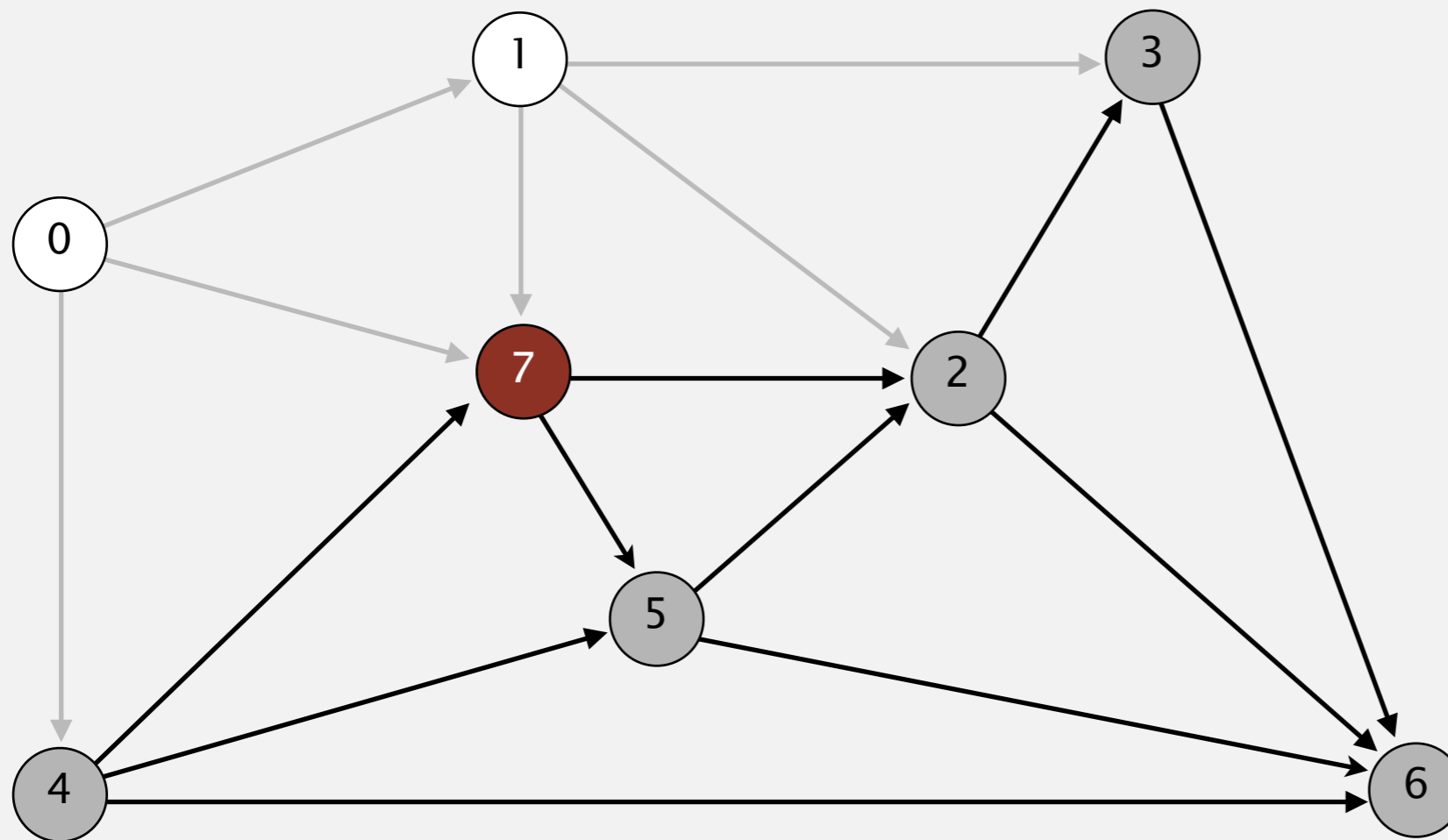
- Consider vertices in increasing order of distance from  $s$  (non-tree vertex with the lowest  $\text{distTo}[]$  value).
- Add vertex to tree and relax all edges adjacent from that vertex.



$v$	$\text{distTo}[]$	$\text{edgeTo}[]$
0	0.0	-
1	5.0	0→1
2	17.0	1→2
3	20.0	1→3
4	9.0	0→4
5		
6		
7	8.0	0→7

# Dijkstra's algorithm demo

- Consider vertices in increasing order of distance from  $s$  (non-tree vertex with the lowest  $\text{distTo}[]$  value).
- Add vertex to tree and relax all edges adjacent from that vertex.



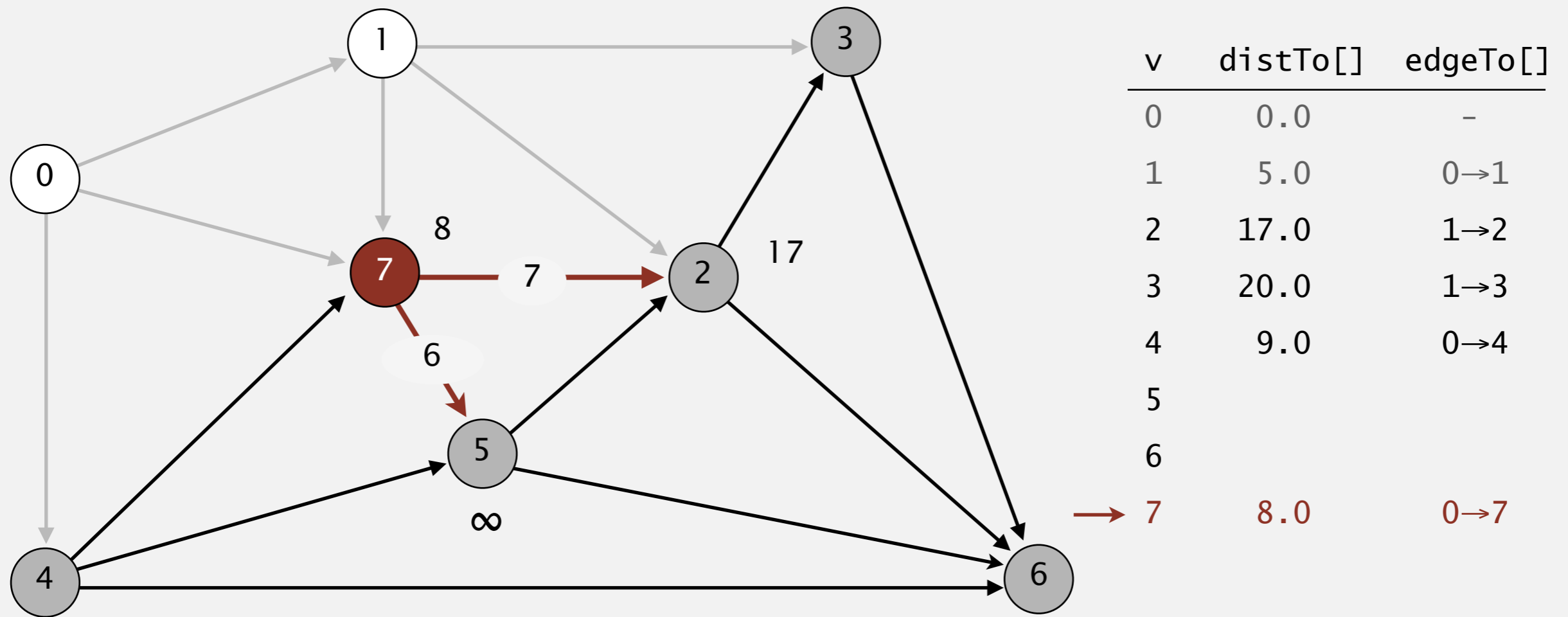
$v$	$\text{distTo}[]$	$\text{edgeTo}[]$
0	0.0	-
1	5.0	0→1
2	17.0	1→2
3	20.0	1→3
4	9.0	0→4
5		
6		
→ 7	8.0	0→7

choose vertex 7



# Dijkstra's algorithm demo

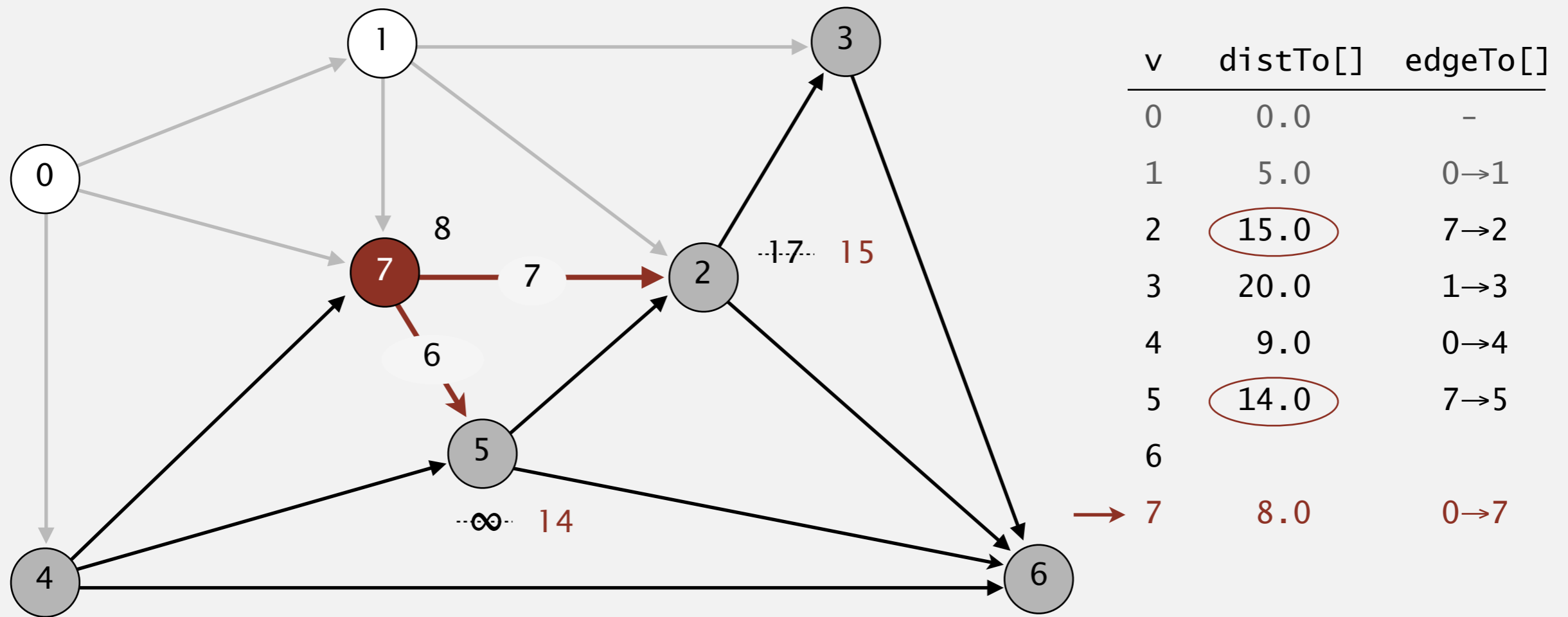
- Consider vertices in increasing order of distance from  $s$  (non-tree vertex with the lowest  $\text{distTo}[]$  value).
- Add vertex to tree and relax all edges adjacent from that vertex.



relax all edges adjacent from 7

# Dijkstra's algorithm demo

- Consider vertices in increasing order of distance from  $s$  (non-tree vertex with the lowest  $\text{distTo}[]$  value).
- Add vertex to tree and relax all edges adjacent from that vertex.

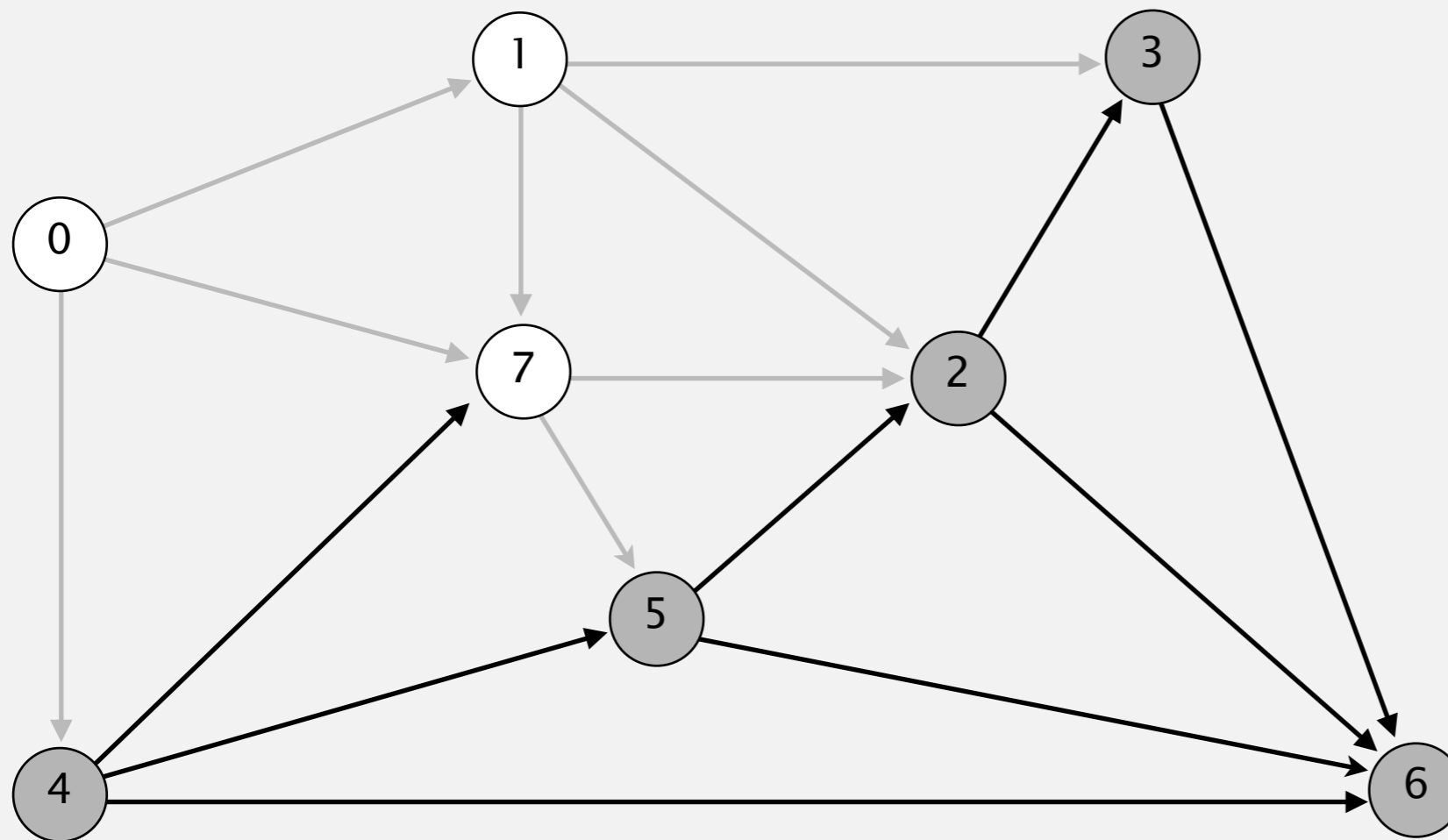


relax all edges adjacent from 7

# Dijkstra's algorithm demo

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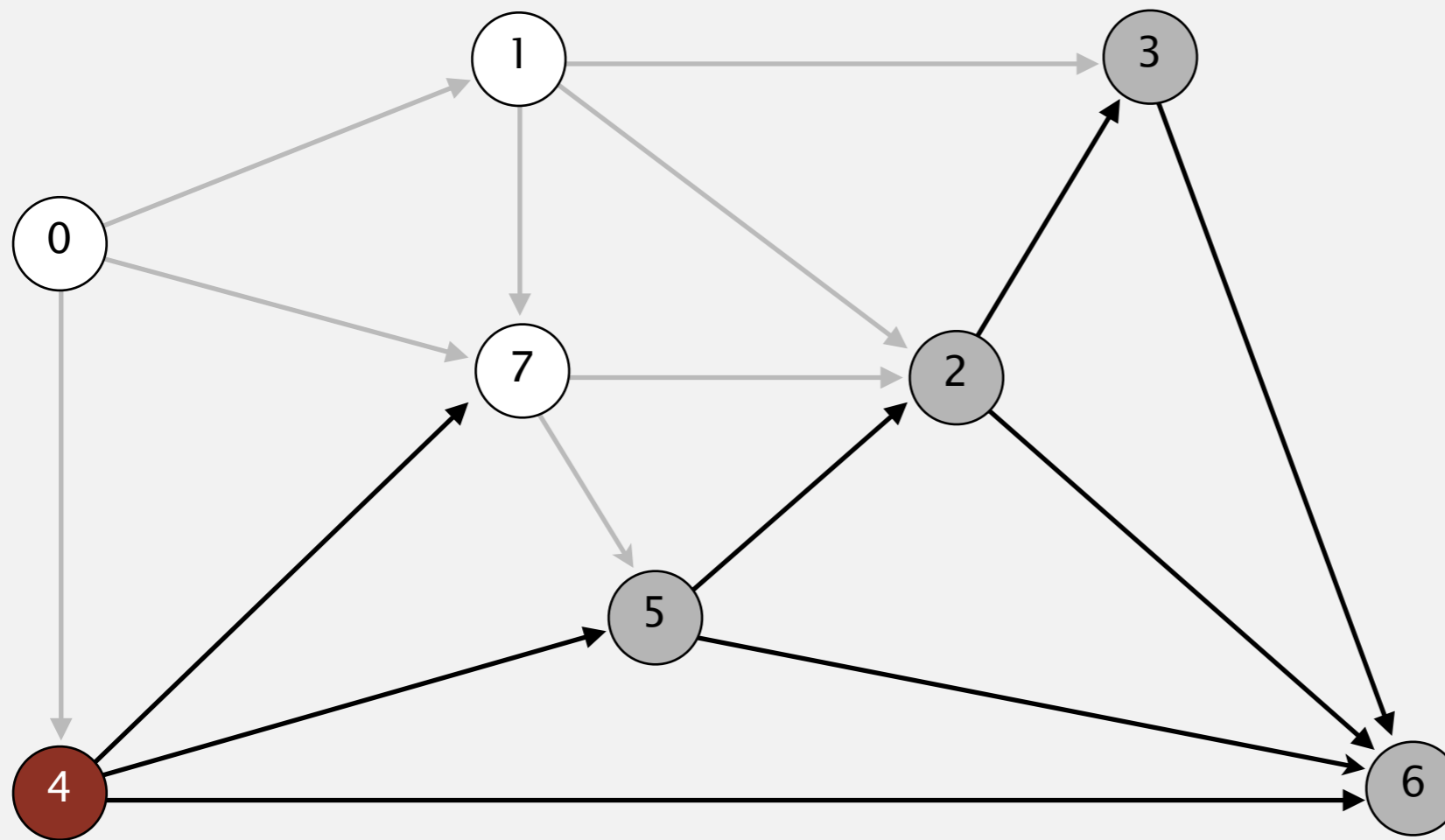
- Consider vertices in increasing order of distance from  $s$  (non-tree vertex with the lowest  $\text{distTo}[]$  value).
- Add vertex to tree and relax all edges adjacent from that vertex.



$v$	$\text{distTo}[]$	$\text{edgeTo}[]$
0	0.0	-
1	5.0	0→1
2	15.0	7→2
3	20.0	1→3
4	9.0	0→4
5	14.0	7→5
6		
7	8.0	0→7

# Dijkstra's algorithm demo

- Consider vertices in increasing order of distance from  $s$  (non-tree vertex with the lowest  $\text{distTo}[]$  value).
- Add vertex to tree and relax all edges adjacent from that vertex.

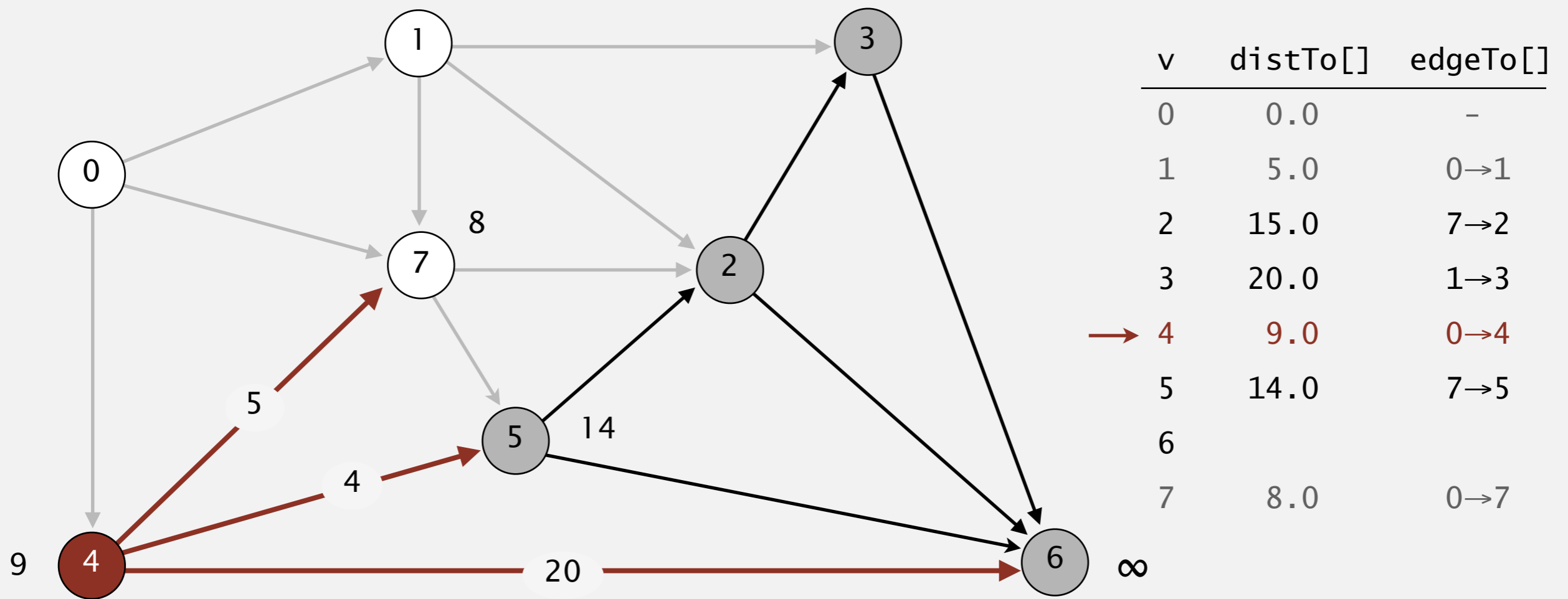


$v$	$\text{distTo}[]$	$\text{edgeTo}[]$
0	0.0	-
1	5.0	0→1
2	15.0	7→2
3	20.0	1→3
→ 4	9.0	0→4
5	14.0	7→5
6		
7	8.0	0→7

**select vertex 4**

# Dijkstra's algorithm demo

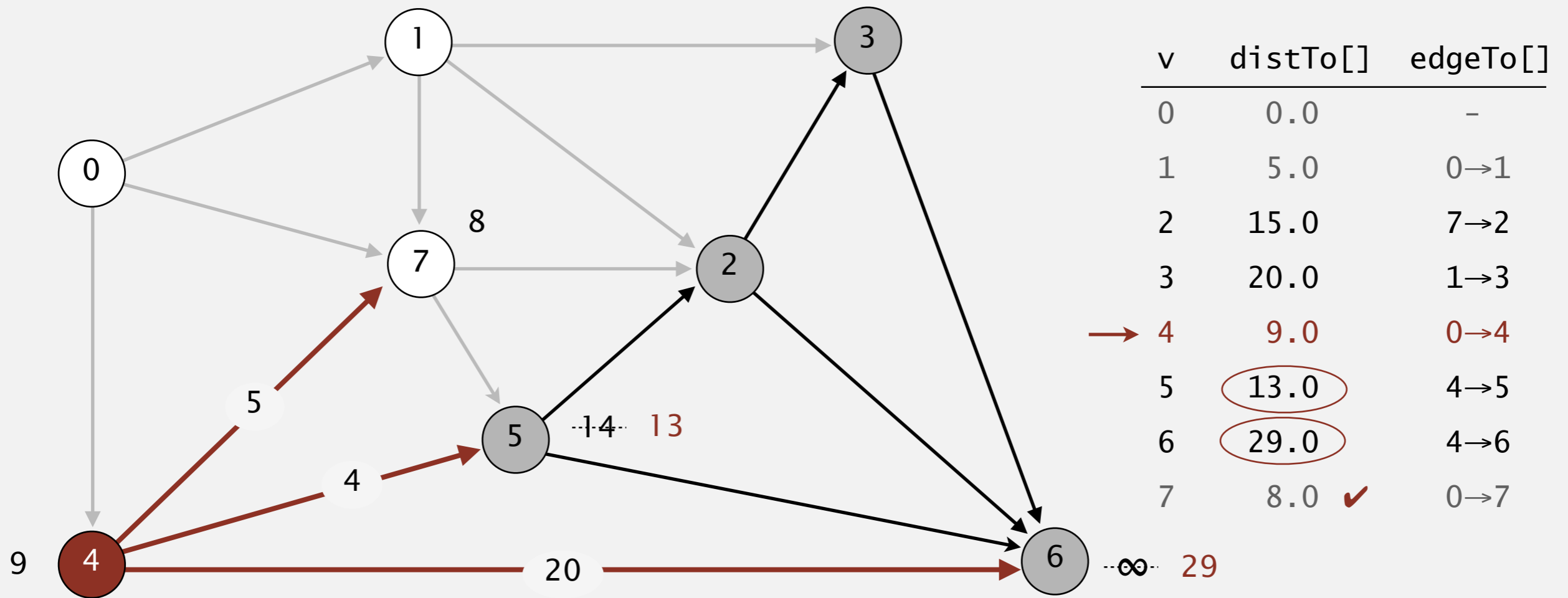
- Consider vertices in increasing order of distance from  $s$  (non-tree vertex with the lowest  $\text{distTo}[]$  value).
- Add vertex to tree and relax all edges adjacent from that vertex.



relax all edges adjacent from 4

# Dijkstra's algorithm demo

- Consider vertices in increasing order of distance from  $s$  (non-tree vertex with the lowest  $\text{distTo}[]$  value).
- Add vertex to tree and relax all edges adjacent from that vertex.

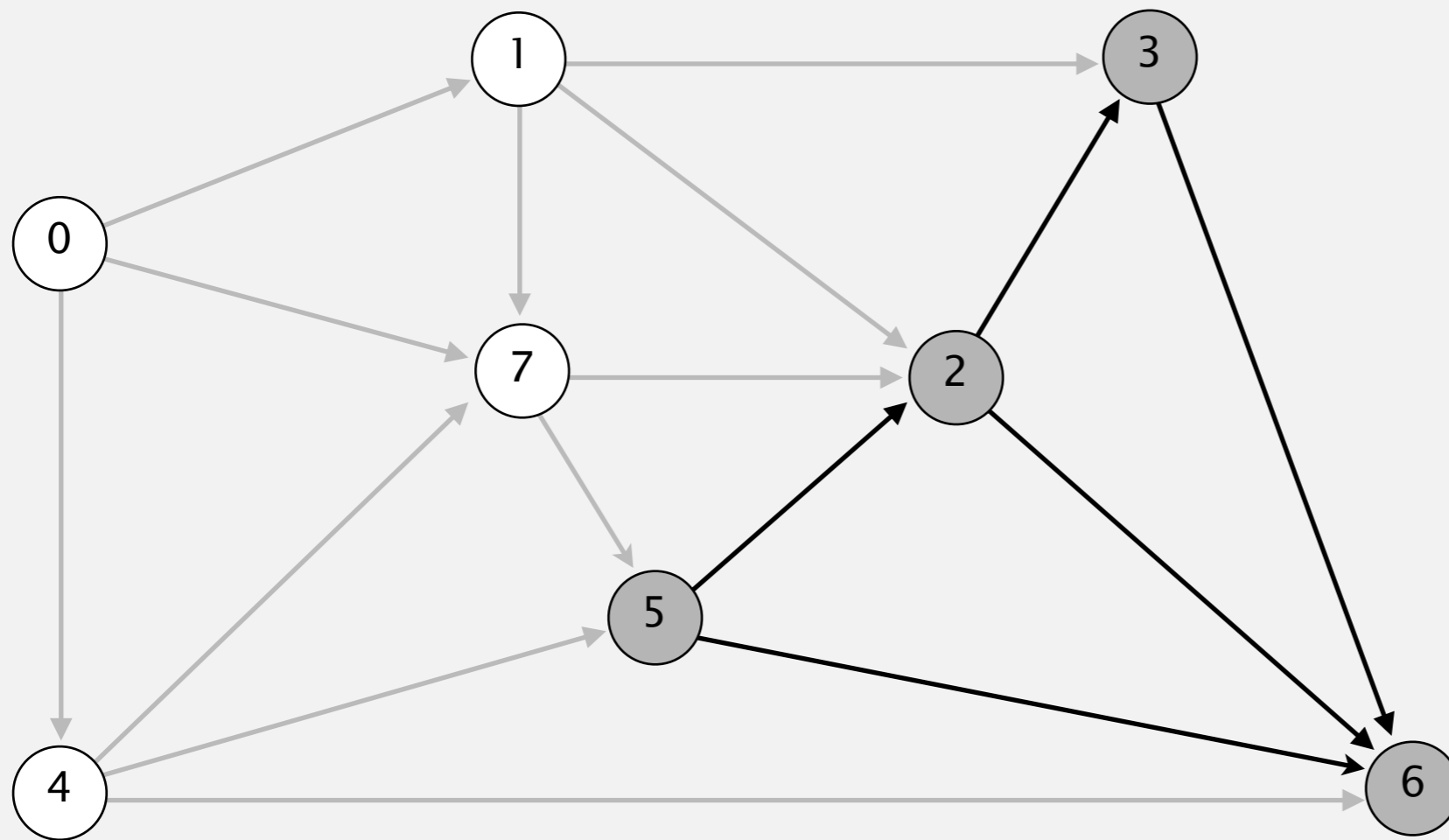


relax all edges adjacent from 4

# Dijkstra's algorithm demo

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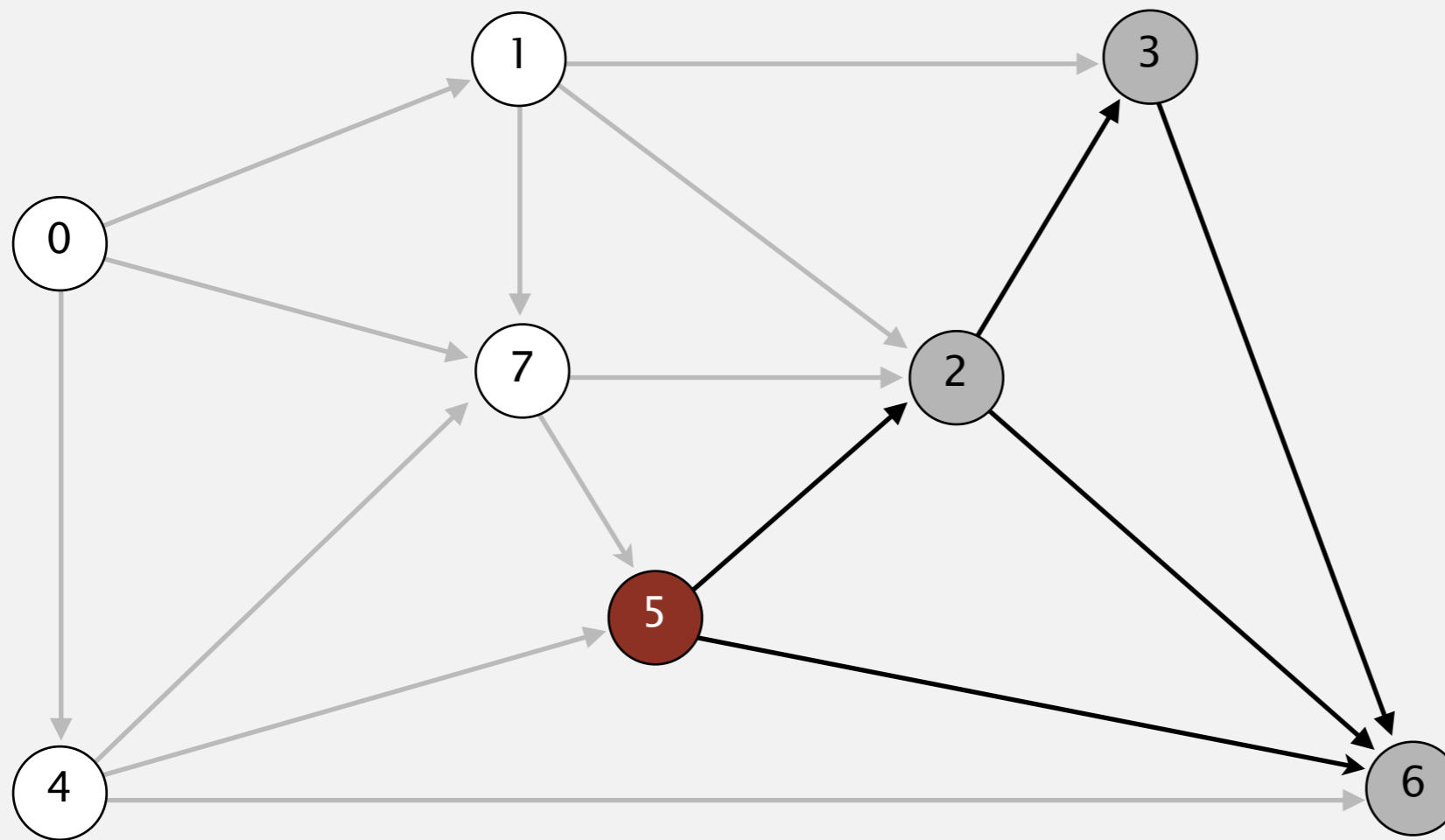
- Consider vertices in increasing order of distance from  $s$  (non-tree vertex with the lowest `distTo[]` value).
- Add vertex to tree and relax all edges adjacent from that vertex.



<code>v</code>	<code>distTo[]</code>	<code>edgeTo[]</code>
0	0.0	-
1	5.0	0→1
2	15.0	7→2
3	20.0	1→3
4	9.0	0→4
5	13.0	4→5
6	29.0	4→6
7	8.0	0→7

# Dijkstra's algorithm demo

- Consider vertices in increasing order of distance from  $s$  (non-tree vertex with the lowest  $\text{distTo}[]$  value).
- Add vertex to tree and relax all edges adjacent from that vertex.



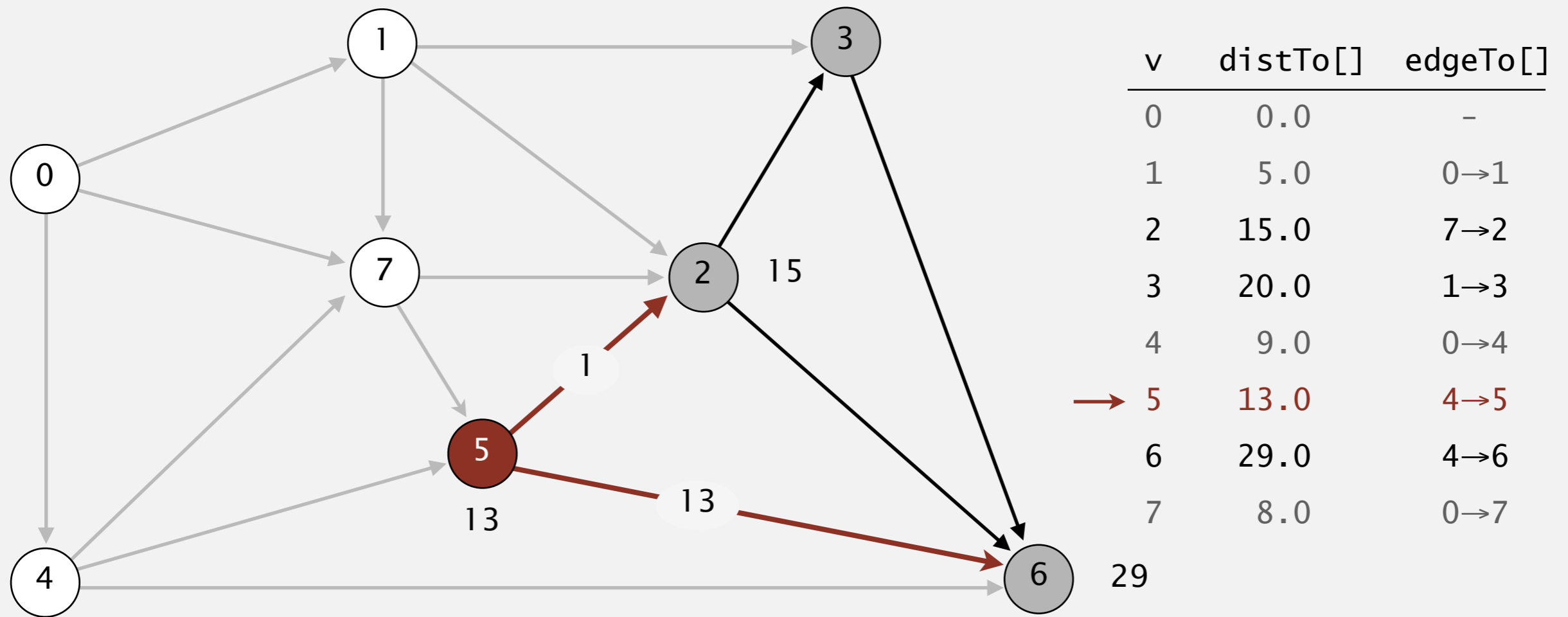
<u>v</u>	<u>distTo[]</u>	<u>edgeTo[]</u>
0	0.0	-
1	5.0	0→1
2	15.0	7→2
3	20.0	1→3
4	9.0	0→4
→ 5	13.0	4→5
6	29.0	4→6
7	8.0	0→7

**select vertex 5**



# Dijkstra's algorithm demo

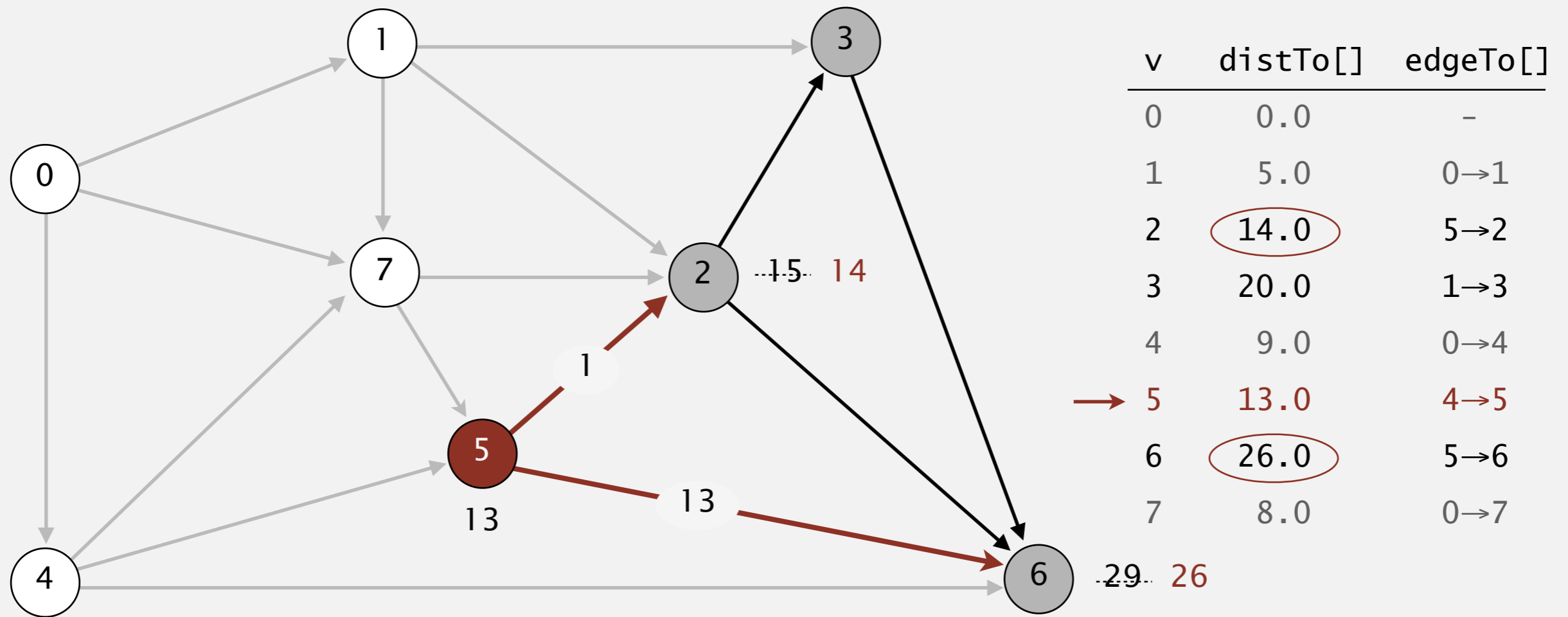
- Consider vertices in increasing order of distance from  $s$  (non-tree vertex with the lowest  $\text{distTo}[]$  value).
- Add vertex to tree and relax all edges adjacent from that vertex.



relax all edges adjacent from 5

# Dijkstra's algorithm demo

- Consider vertices in increasing order of distance from  $s$  (non-tree vertex with the lowest `distTo[]` value).
- Add vertex to tree and relax all edges adjacent from that vertex.

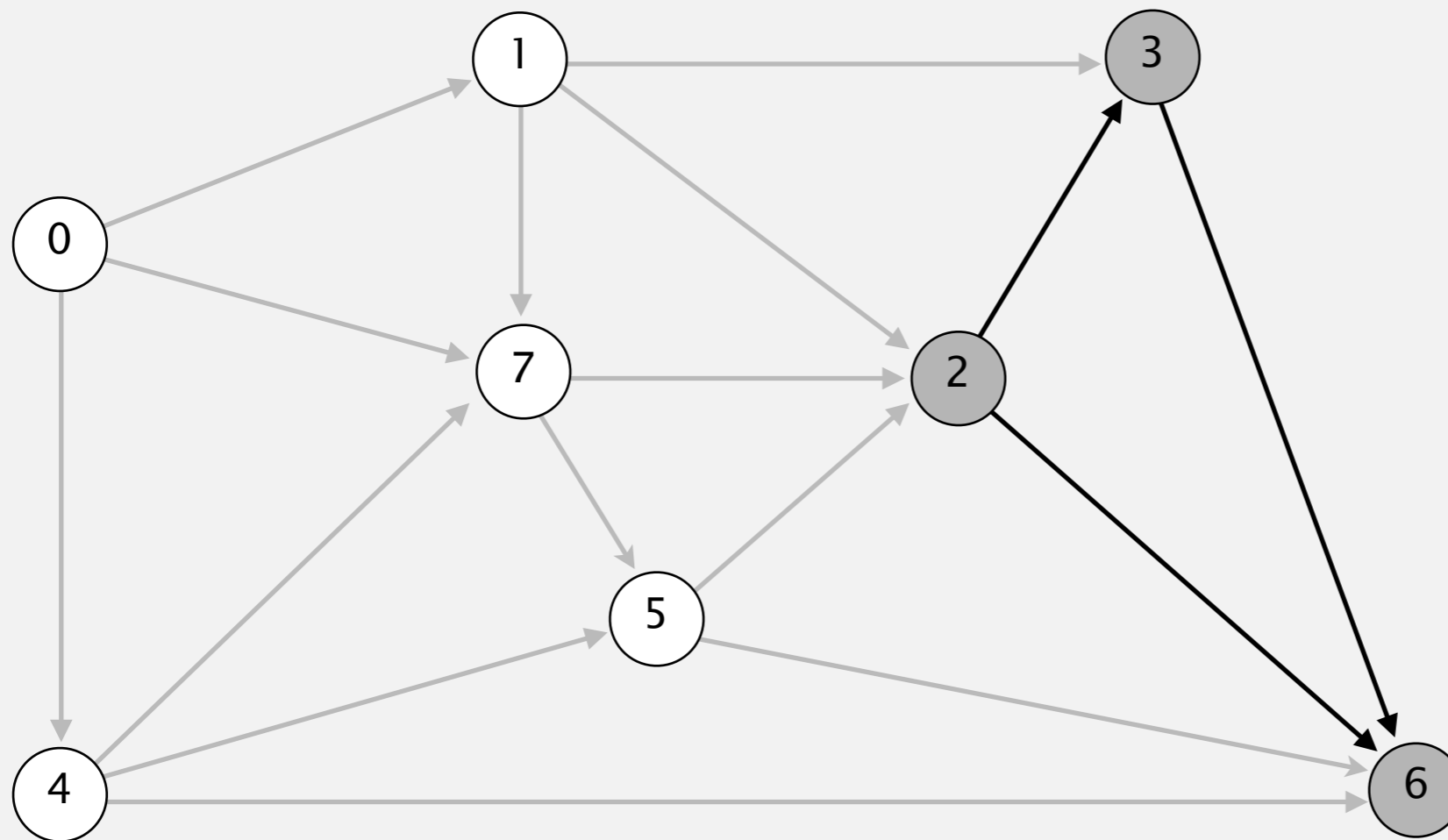


relax all edges adjacent from 5

# Dijkstra's algorithm demo

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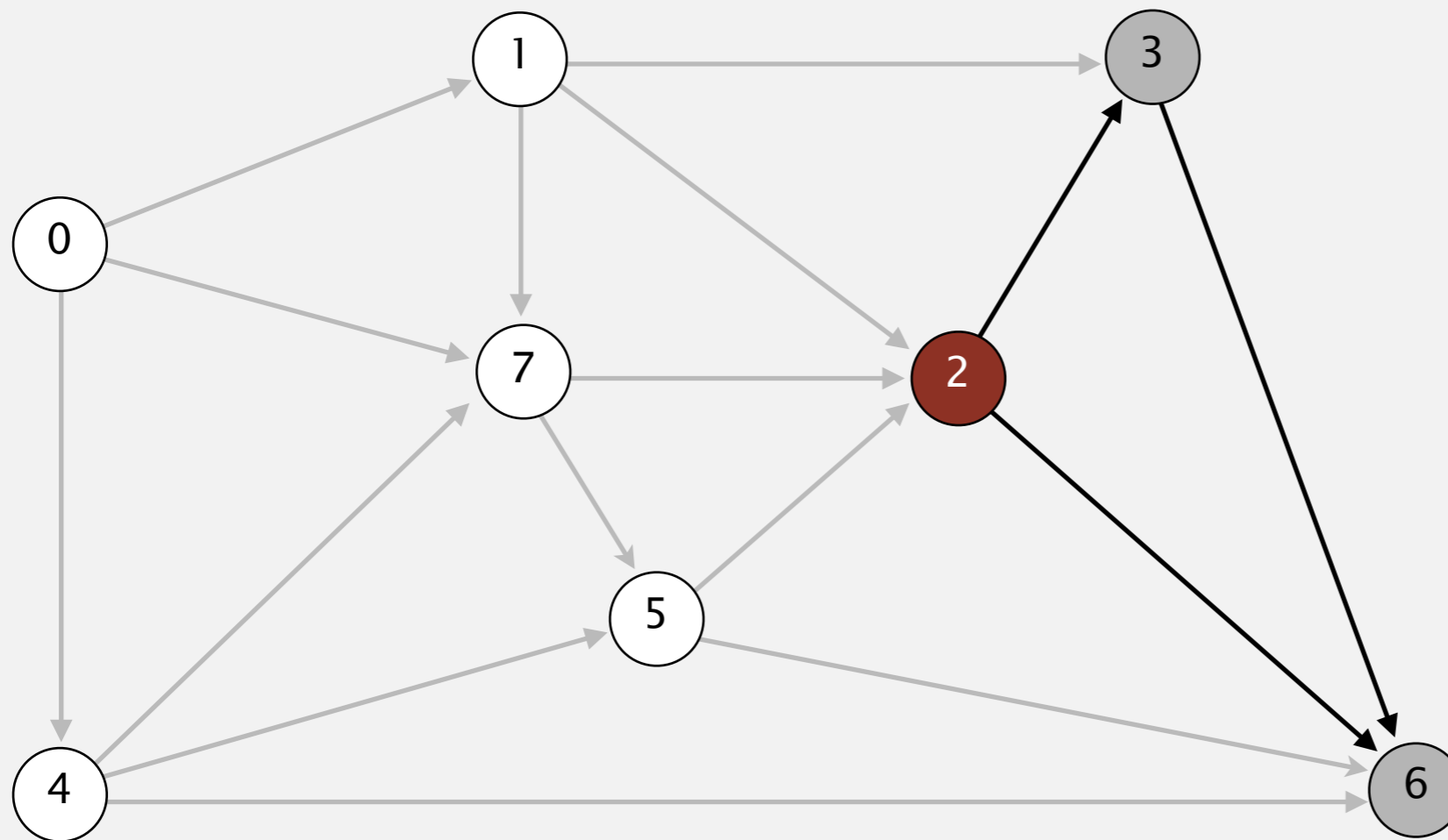
- Consider vertices in increasing order of distance from  $s$  (non-tree vertex with the lowest `distTo[]` value).
- Add vertex to tree and relax all edges adjacent from that vertex.



<code>v</code>	<code>distTo[]</code>	<code>edgeTo[]</code>
0	0.0	-
1	5.0	0→1
2	14.0	5→2
3	20.0	1→3
4	9.0	0→4
5	13.0	4→5
6	26.0	5→6
7	8.0	0→7

# Dijkstra's algorithm demo

- Consider vertices in increasing order of distance from  $s$  (non-tree vertex with the lowest  $\text{distTo}[]$  value).
- Add vertex to tree and relax all edges adjacent from that vertex.

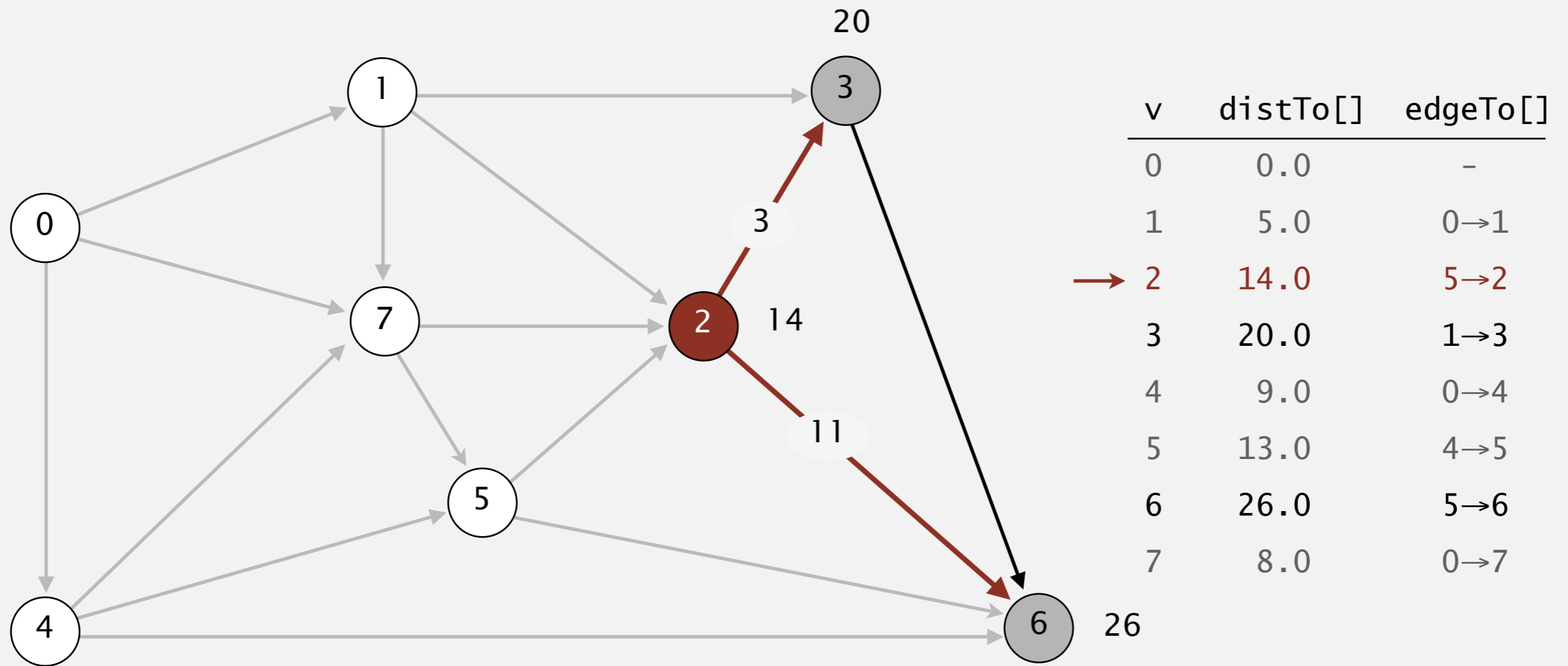


<u>v</u>	<u>distTo[]</u>	<u>edgeTo[]</u>
0	0.0	-
1	5.0	0→1
→ 2	14.0	5→2
3	20.0	1→3
4	9.0	0→4
5	13.0	4→5
6	26.0	5→6
7	8.0	0→7

**select vertex 2**

# Dijkstra's algorithm demo

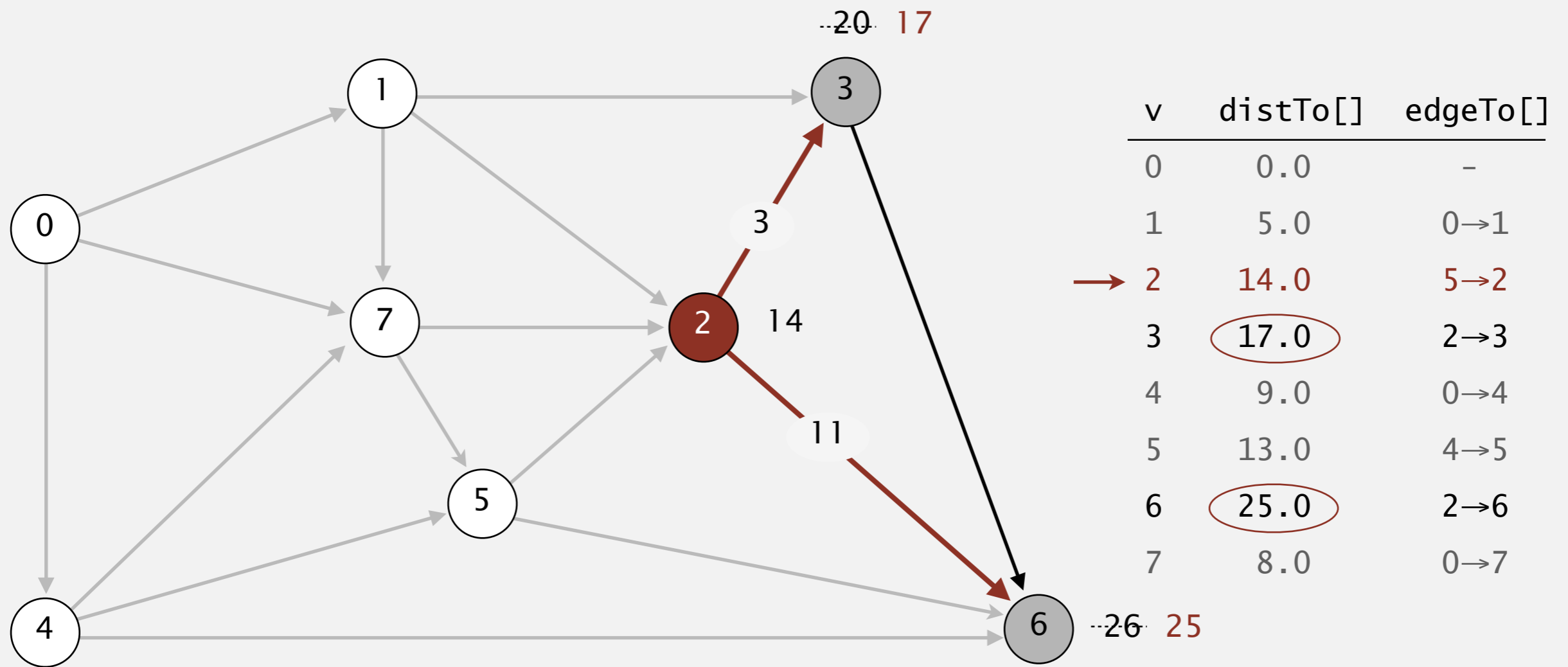
- Consider vertices in increasing order of distance from  $s$  (non-tree vertex with the lowest  $\text{distTo}[]$  value).
- Add vertex to tree and relax all edges adjacent from that vertex.



relax all edges adjacent from 2

# Dijkstra's algorithm demo

- Consider vertices in increasing order of distance from  $s$  (non-tree vertex with the lowest  $\text{distTo}[]$  value).
- Add vertex to tree and relax all edges adjacent from that vertex.

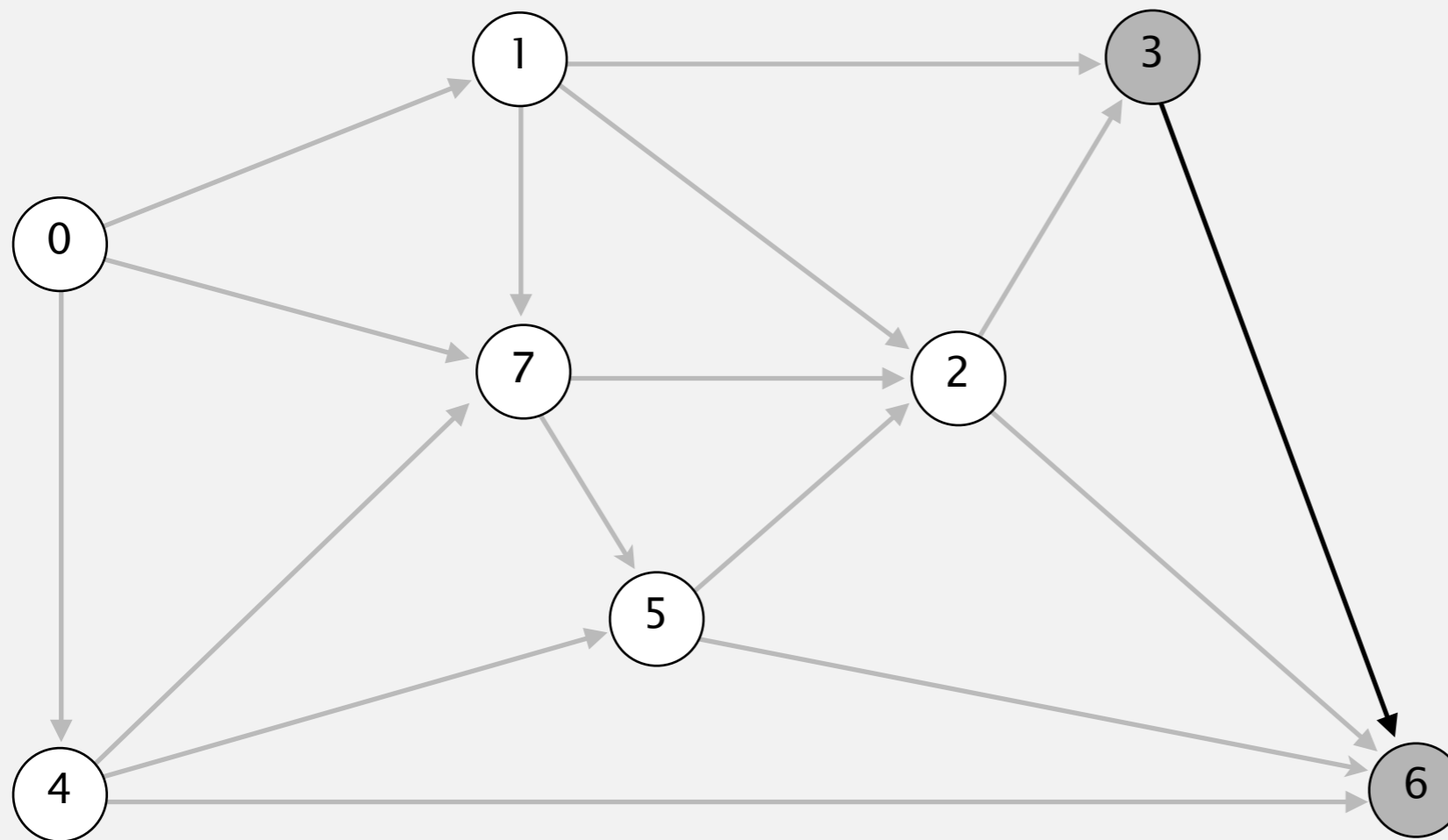


relax all edges adjacent from 2

# Dijkstra's algorithm demo

---

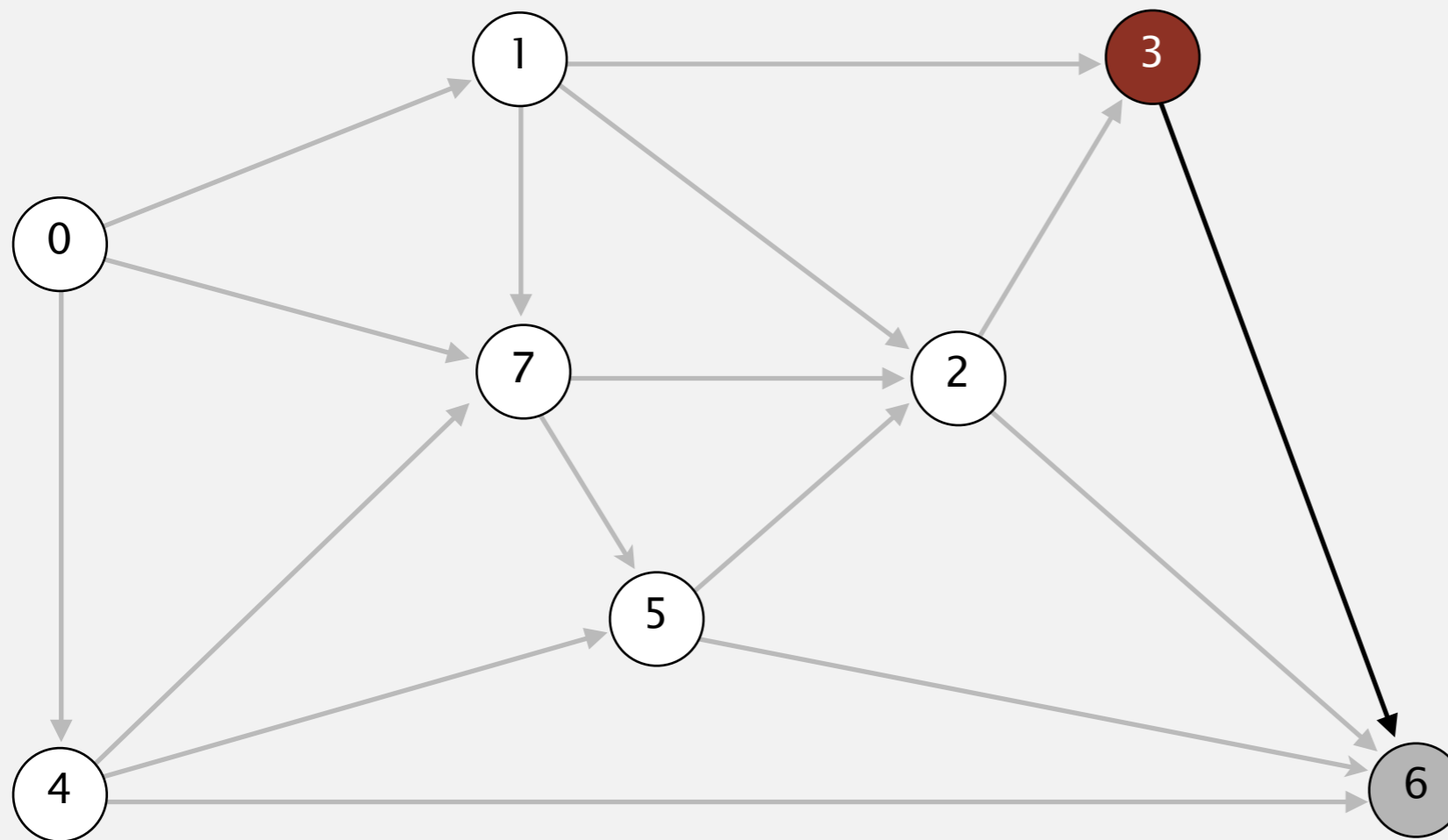
- Consider vertices in increasing order of distance from  $s$  (non-tree vertex with the lowest  $\text{distTo}[]$  value).
- Add vertex to tree and relax all edges adjacent from that vertex.



$v$	$\text{distTo}[]$	$\text{edgeTo}[]$
0	0.0	-
1	5.0	0→1
2	14.0	5→2
3	17.0	2→3
4	9.0	0→4
5	13.0	4→5
6	25.0	2→6
7	8.0	0→7

# Dijkstra's algorithm demo

- Consider vertices in increasing order of distance from  $s$  (non-tree vertex with the lowest  $\text{distTo}[]$  value).
- Add vertex to tree and relax all edges adjacent from that vertex.



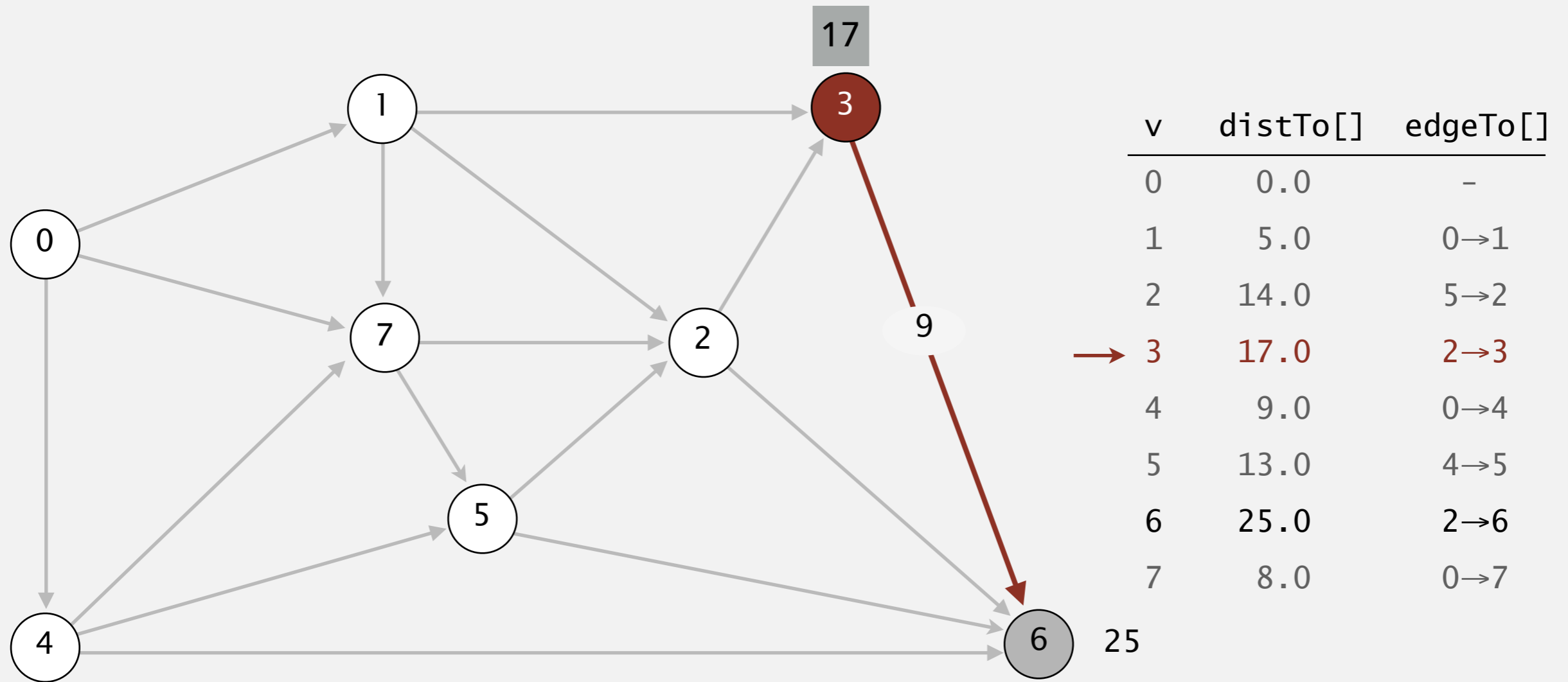
<u>v</u>	<u>distTo[]</u>	<u>edgeTo[]</u>
0	0.0	-
1	5.0	0→1
2	14.0	5→2
→ 3	17.0	2→3
4	9.0	0→4
5	13.0	4→5
6	25.0	2→6
7	8.0	0→7

**select vertex 3**



# Dijkstra's algorithm demo

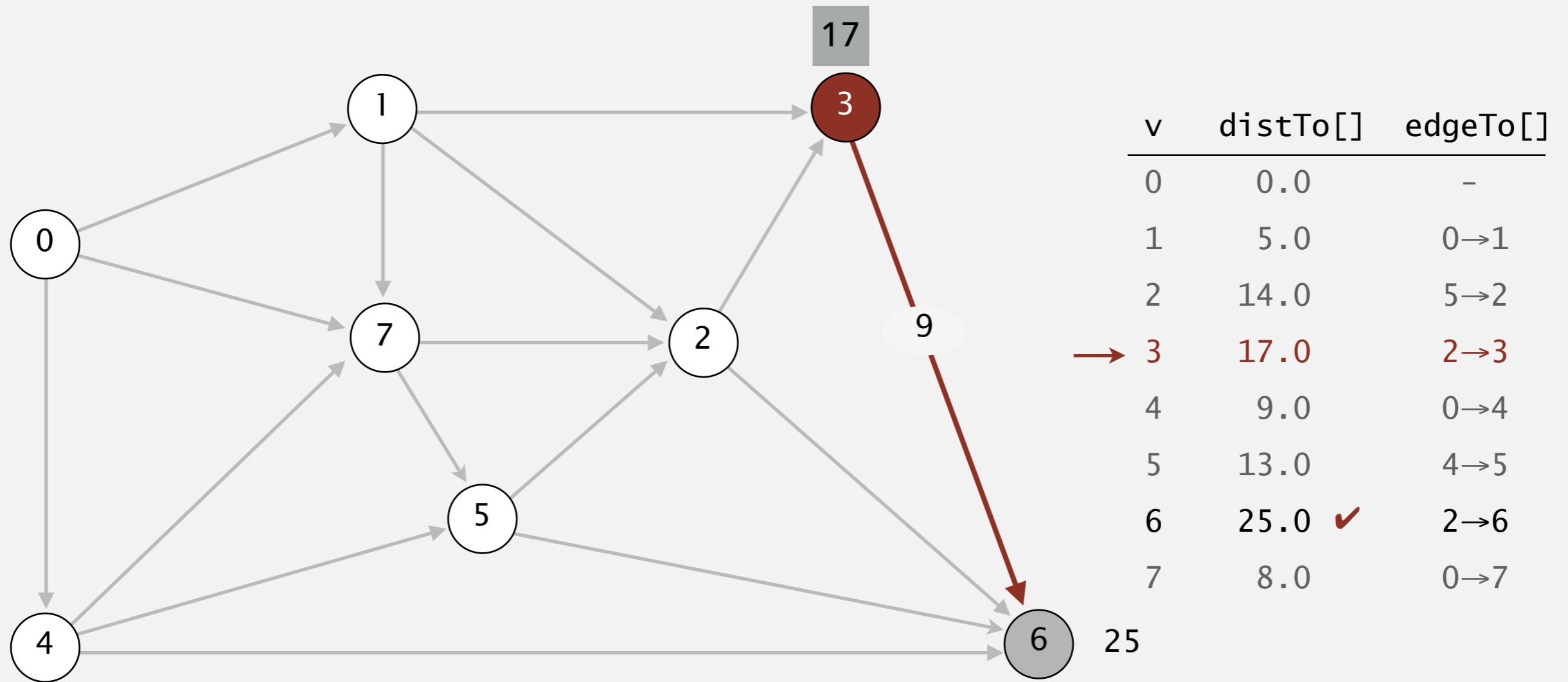
- Consider vertices in increasing order of distance from  $s$  (non-tree vertex with the lowest  $\text{distTo}[]$  value).
- Add vertex to tree and relax all edges adjacent from that vertex.



relax all edges adjacent from 3

# Dijkstra's algorithm demo

- Consider vertices in increasing order of distance from  $s$  (non-tree vertex with the lowest  $\text{distTo}[]$  value).
- Add vertex to tree and relax all edges adjacent from that vertex.

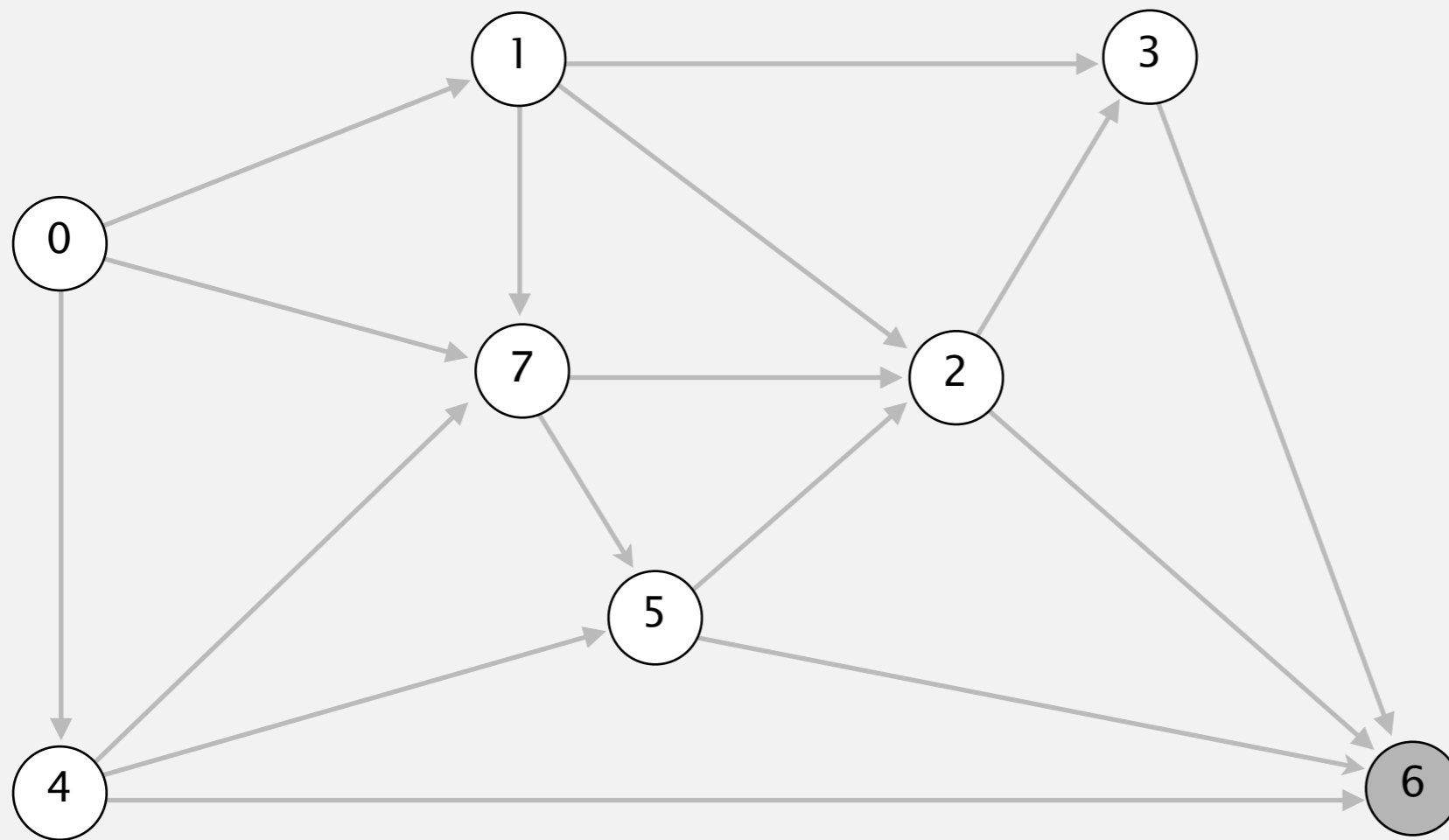


relax all edges adjacent from 3

# Dijkstra's algorithm demo

---

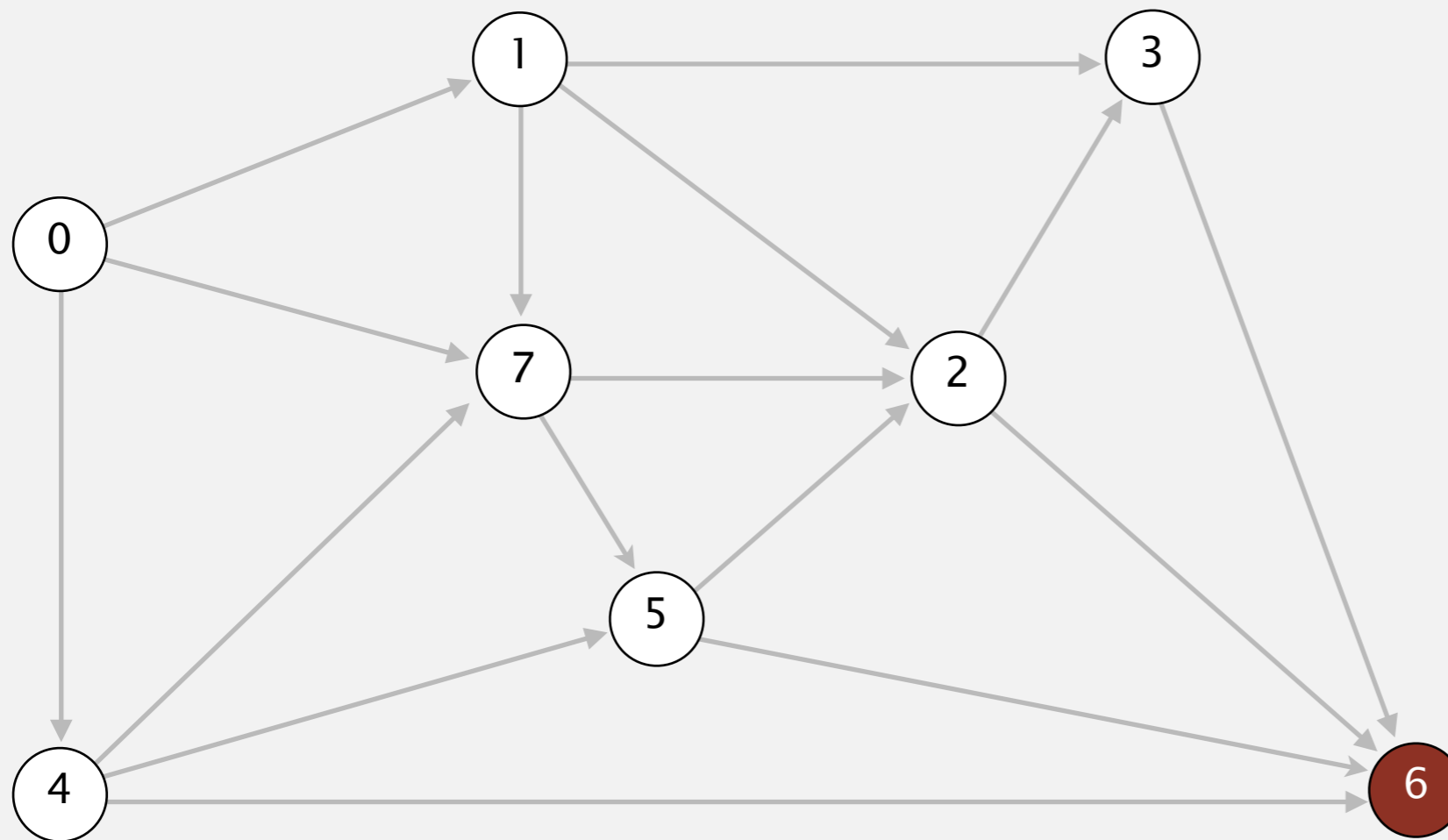
- Consider vertices in increasing order of distance from  $s$  (non-tree vertex with the lowest  $\text{distTo}[]$  value).
- Add vertex to tree and relax all edges adjacent from that vertex.



$v$	$\text{distTo}[]$	$\text{edgeTo}[]$
0	0.0	-
1	5.0	0→1
2	14.0	5→2
3	17.0	2→3
4	9.0	0→4
5	13.0	4→5
6	25.0	2→6
7	8.0	0→7

# Dijkstra's algorithm demo

- Consider vertices in increasing order of distance from  $s$  (non-tree vertex with the lowest  $\text{distTo}[]$  value).
- Add vertex to tree and relax all edges adjacent from that vertex.

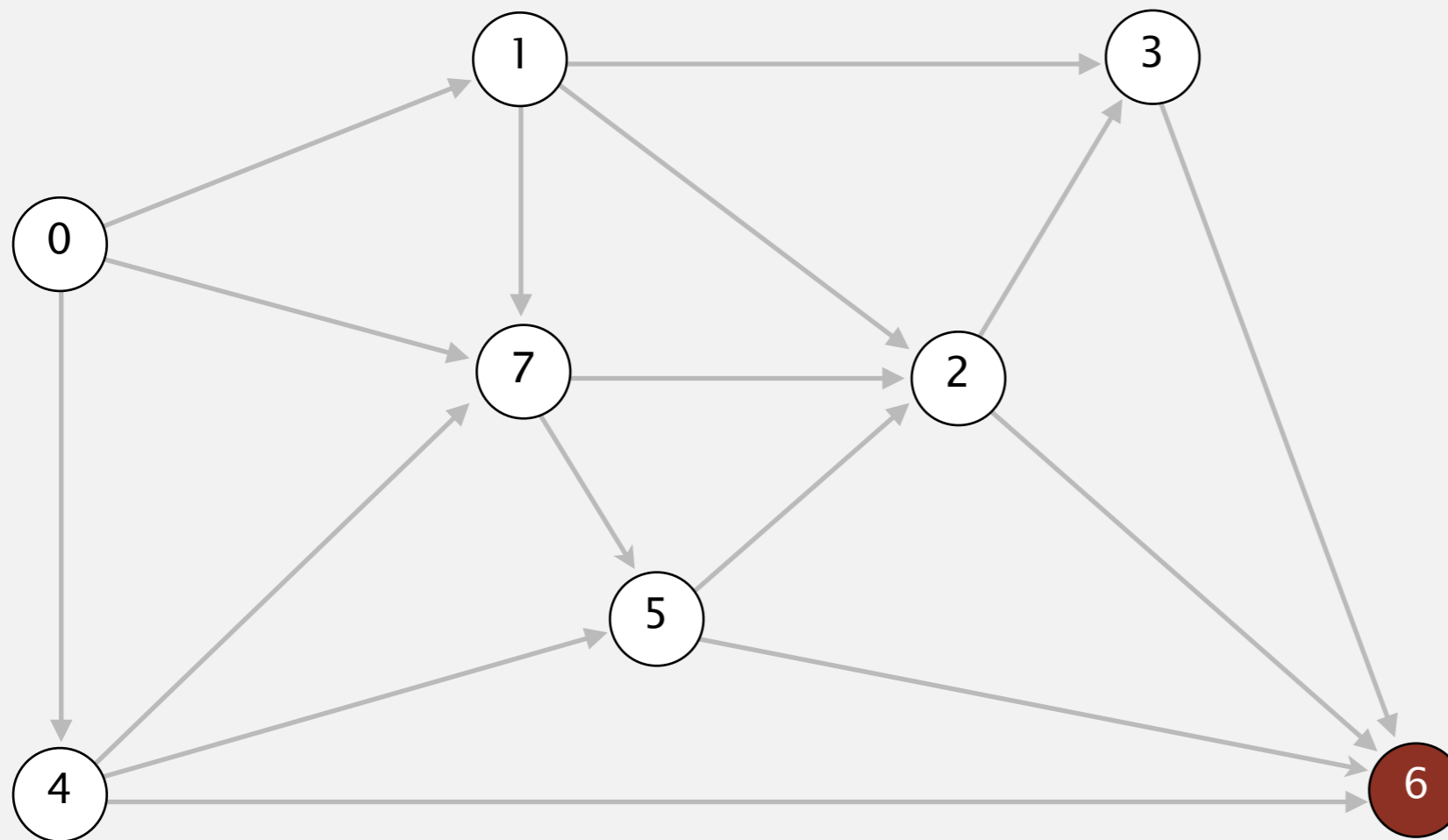


<u>v</u>	<u>distTo[]</u>	<u>edgeTo[]</u>
0	0.0	-
1	5.0	0→1
2	14.0	5→2
3	17.0	2→3
4	9.0	0→4
5	13.0	4→5
→ 6	25.0	2→6
7	8.0	0→7

**select vertex 6**

# Dijkstra's algorithm demo

- Consider vertices in increasing order of distance from  $s$  (non-tree vertex with the lowest  $\text{distTo}[]$  value).
- Add vertex to tree and relax all edges adjacent from that vertex.



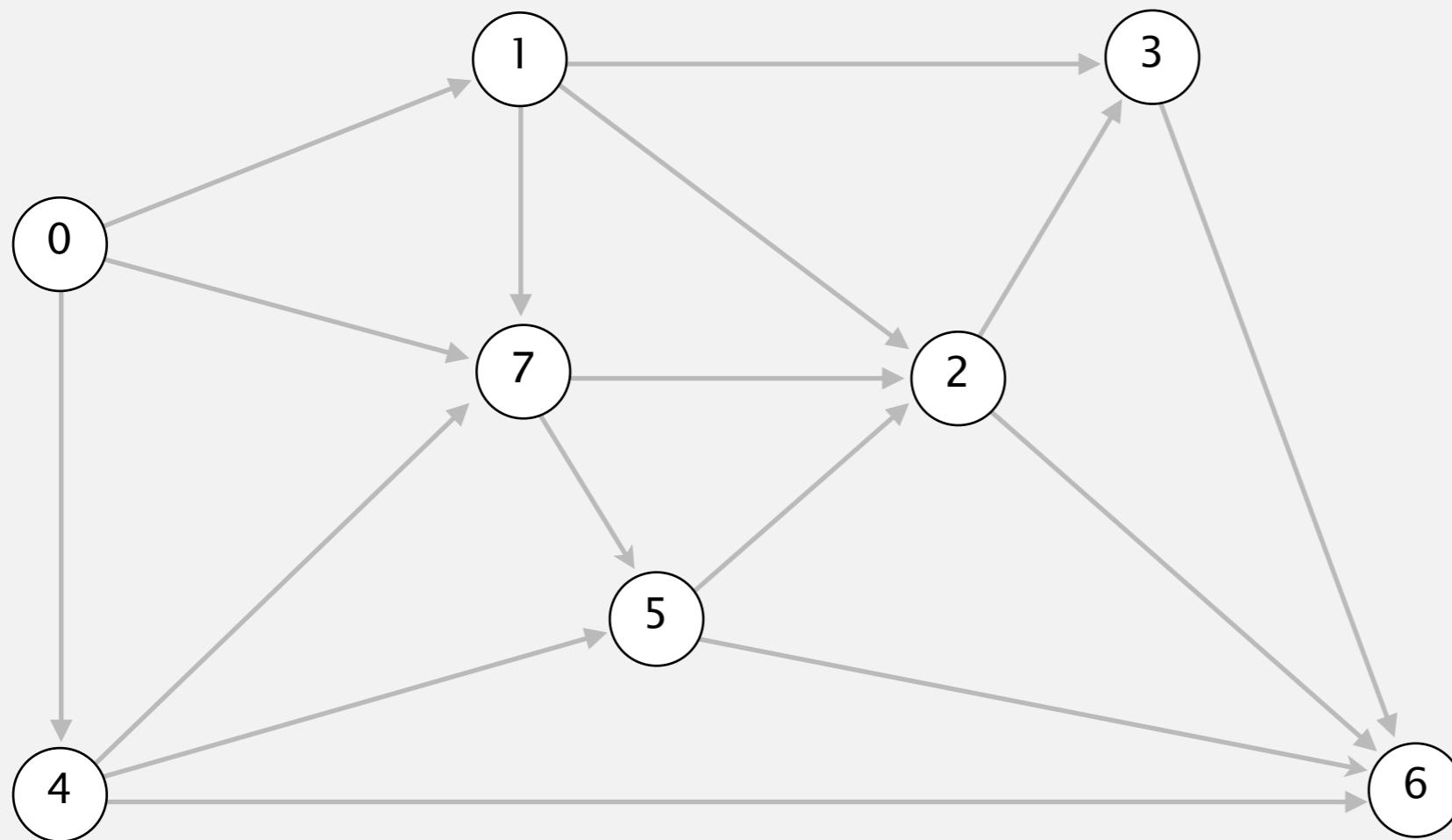
$v$	$\text{distTo}[]$	$\text{edgeTo}[]$
0	0.0	-
1	5.0	0→1
2	14.0	5→2
3	17.0	2→3
4	9.0	0→4
5	13.0	4→5
→ 6	25.0	2→6
7	8.0	0→7

relax all edges adjacent from 6

# Dijkstra's algorithm demo

---

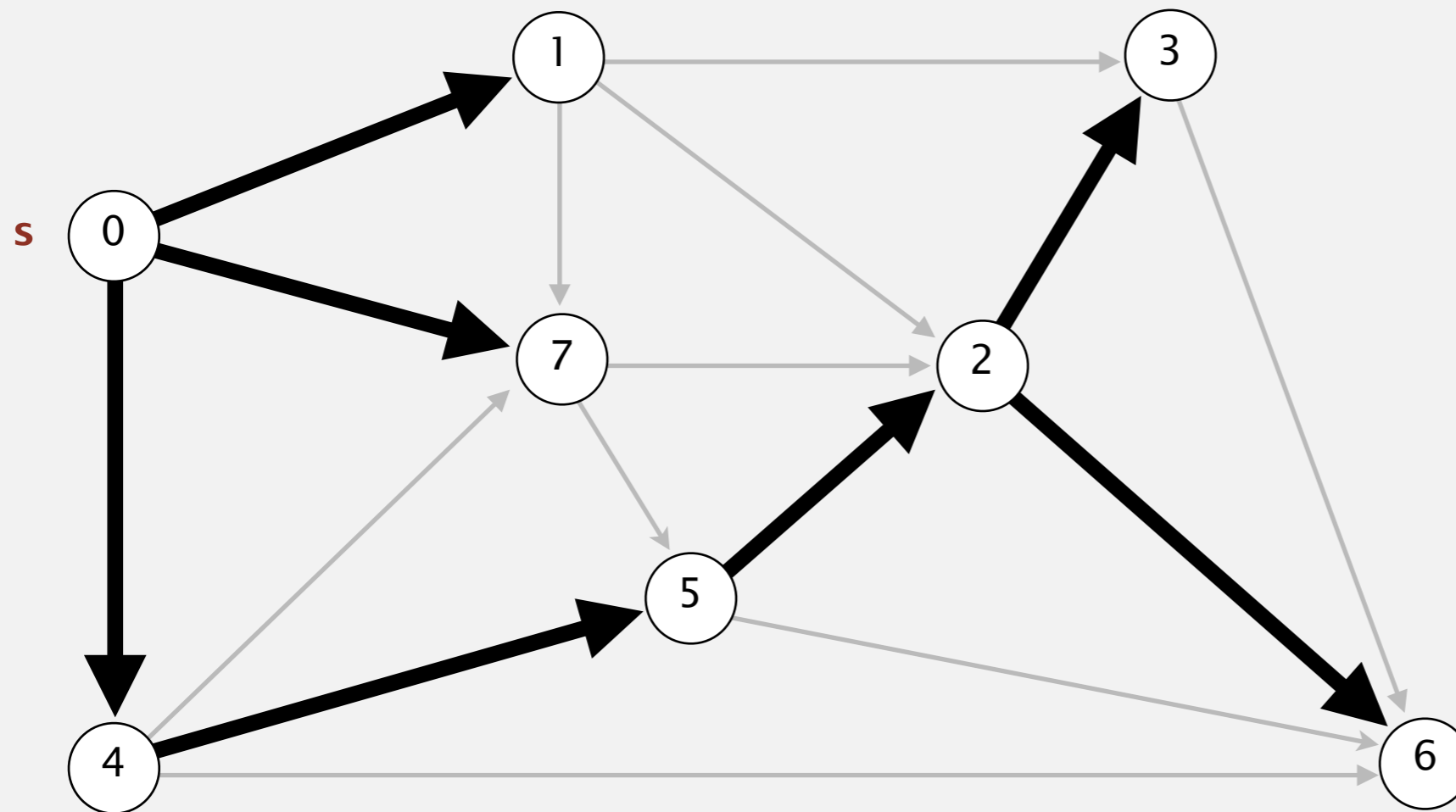
- Consider vertices in increasing order of distance from  $s$  (non-tree vertex with the lowest  $\text{distTo}[]$  value).
- Add vertex to tree and relax all edges adjacent from that vertex.



$v$	$\text{distTo}[]$	$\text{edgeTo}[]$
0	0.0	-
1	5.0	0→1
2	14.0	5→2
3	17.0	2→3
4	9.0	0→4
5	13.0	4→5
6	25.0	2→6
7	8.0	0→7

# Dijkstra's algorithm demo

- Consider vertices in increasing order of distance from  $s$  (non-tree vertex with the lowest  $\text{distTo}[]$  value).
- Add vertex to tree and relax all edges adjacent from that vertex.



$v$	$\text{distTo}[]$	$\text{edgeTo}[]$
0	0.0	-
1	5.0	0→1
2	14.0	5→2
3	17.0	2→3
4	9.0	0→4
5	13.0	4→5
6	25.0	2→6
7	8.0	0→7

shortest-paths tree from vertex  $s$

## Indexed min-priority queue (Section 2.4 in textbook)

- ▶ Associate an index between 0 and  $n-1$  with each key in a priority queue.
  - ▶ Insert a key associated with a given index.
  - ▶ Delete a minimum key and return associated index.
  - ▶ Decrease the key associated with a given index.
- ▶ `public class` IndexMinPQ<Key extends Comparable<Key>>
  - ▶ IndexMinPQ(`int` n)
    - ▶ Create indexed PQ with indices 0,1,...n-1
  - ▶ `void` insert(`int` i, Key key)
    - ▶ Associate key with index i.
  - ▶ `int` delMin()
    - ▶ Remove a minimal key and return its associated index.
  - ▶ `void` decreaseKey(`int` i, Key key)
    - ▶ Decrease the key with index i to the specified value.



```
public class DijkstraSP {
    private double[] distTo;           // distTo[v] = distance of shortest s->v path
    private DirectedEdge[] edgeTo;    // edgeTo[v] = last edge on shortest s->v path
    private IndexMinPQ<Double> pq;    // priority queue of vertices

    public DijkstraSP(EdgeWeightedDigraph G, int s) {
        distTo = new double[G.V()];
        edgeTo = new DirectedEdge[G.V()];

        for (int v = 0; v < G.V(); v++)
            distTo[v] = Double.POSITIVE_INFINITY;
        distTo[s] = 0.0;

        // relax vertices in order of distance from s
        pq = new IndexMinPQ<Double>(G.V());
        pq.insert(s, distTo[s]);
        while (!pq.isEmpty()) {
            int v = pq.delMin();
            for (DirectedEdge e : G.adj(v))
                relax(e);
        }
    }

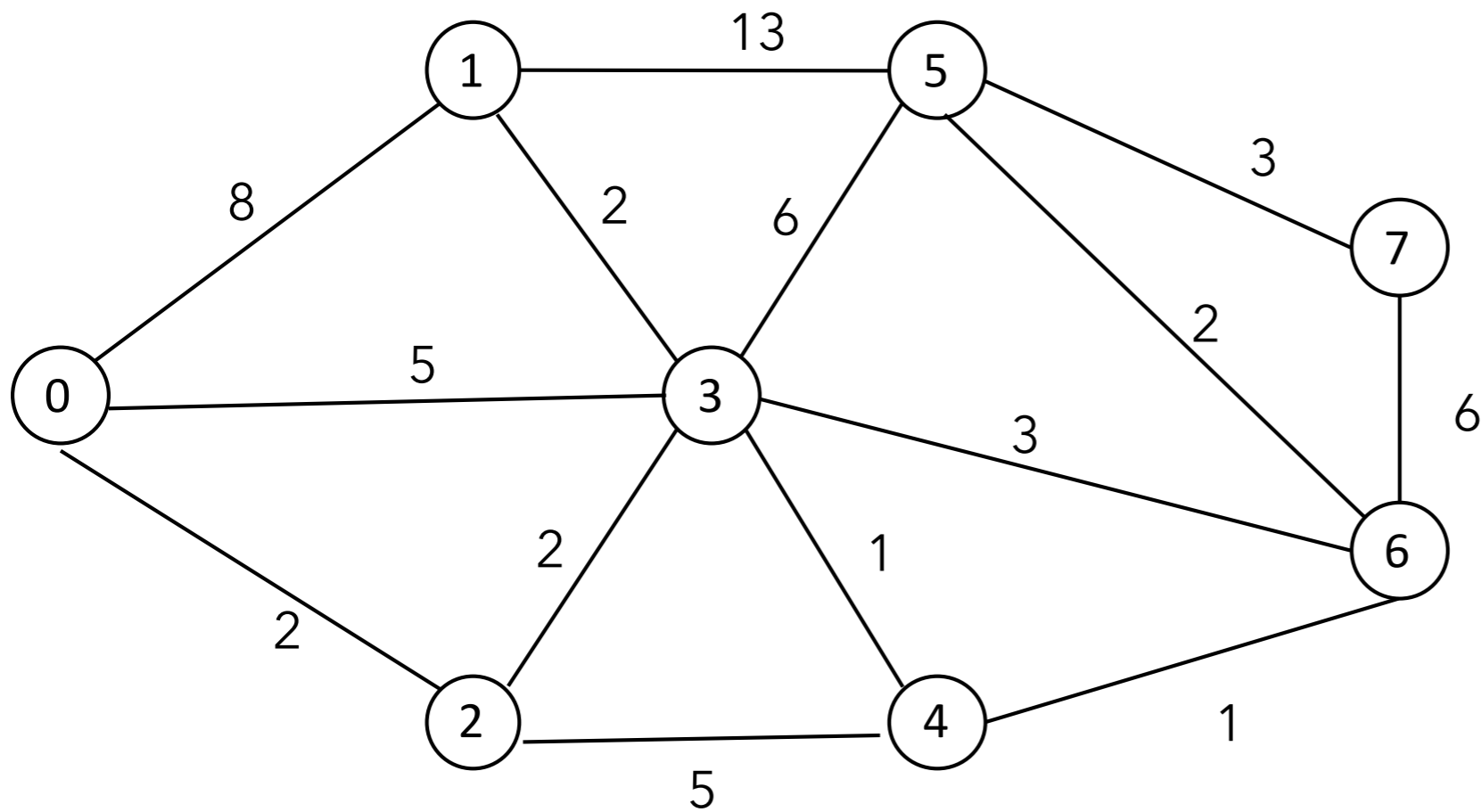
    // relax edge e and update pq if changed
    private void relax(DirectedEdge e) {
        int v = e.from(), w = e.to();
        if (distTo[w] > distTo[v] + e.weight()) {
            distTo[w] = distTo[v] + e.weight();
            edgeTo[w] = e;
            if (pq.contains(w)) pq.decreaseKey(w, distTo[w]);
            else pq.insert(w, distTo[w]);
        }
    }
}
```

Running time depends on PQ implementation

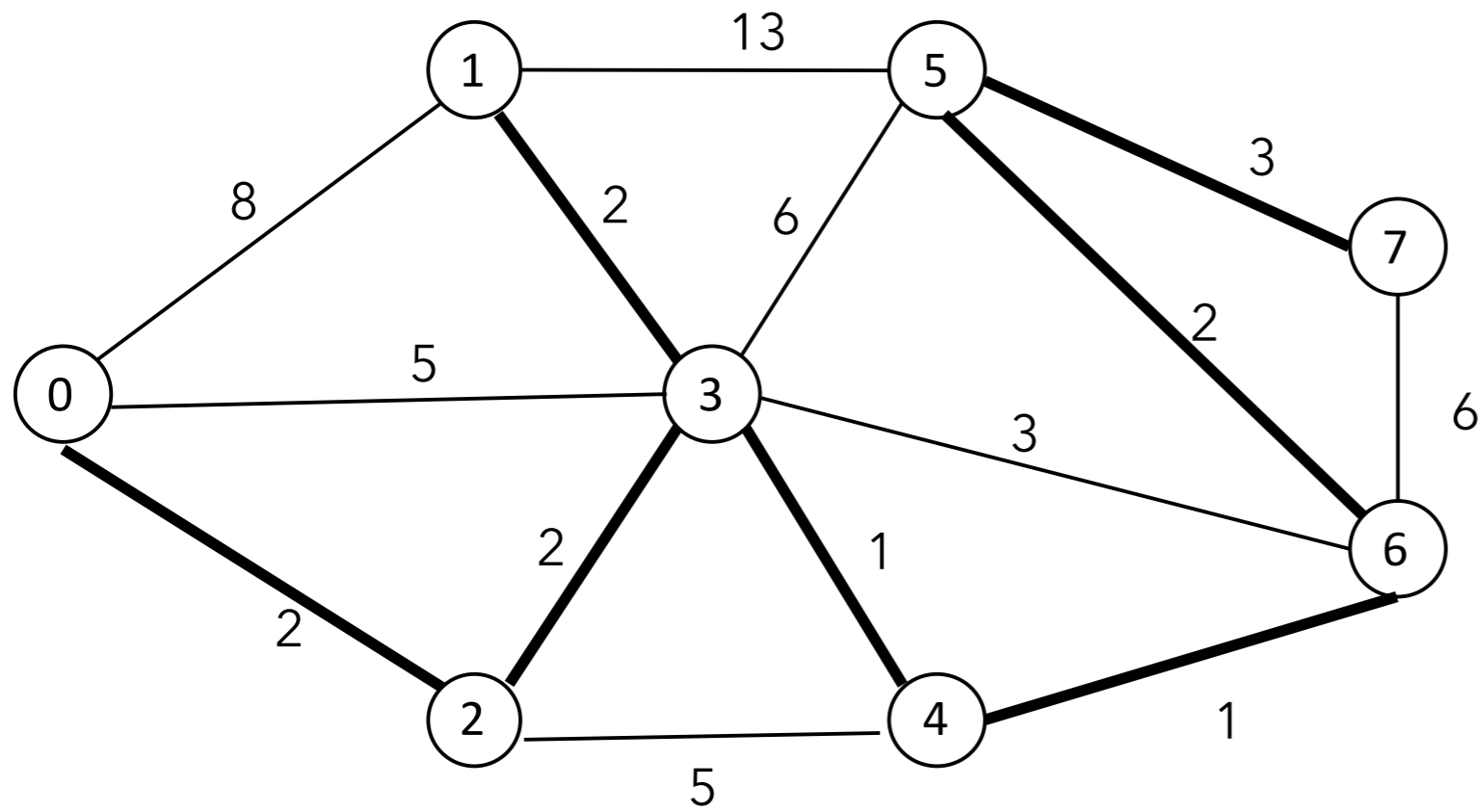
- ▶ Many variations. Assuming binary heap, running time is proportional to  $|E| \log |V|$  and  $|V|$  extra space.
  - ▶ Cost of insert, delete-min, decrease-key are all  $\log V$ .
- ▶ More complicated version with a Fibonacci heap (CS140...) takes  $O(|E| + |V| \log |V|)$  time but in practice it's not worth implementing.

## Practice Time

- ▶ Run Dijkstra's algorithm on the following graph with 0 being the starting vertex.



Answer



v	distTo[]	edgeTo[]
0	0	-
1	6	3->1
2	2	0->2
3	4	2->3
4	5	3->4
5	8	6->5
6	6	4->6
7	11	5->7

## Lecture 27: Shortest Paths

- ▶ Introduction to Shortest Paths
- ▶ API
- ▶ Properties
- ▶ Dijkstra's Algorithm
- ▶ **Belman-Ford Algorithm**

## Framework for shortest-paths algorithm

- ▶ Generic algorithm to compute a SPT from  $s$ 
  - ▶  $\text{distTo}[v] = \infty$  for each vertex  $v$ .
  - ▶  $\text{edgeTo}[v] = \text{null}$  for each vertex  $v$ .
  - ▶  $\text{distTo}[s] = 0$ .
  - ▶ Repeat until done:
    - ▶ Relax any edge.
- ▶  $\text{distTo}[v]$  is the length of a simple path from  $s$  to  $v$ .
- ▶  $\text{distTo}[v]$  does not increase.

## Bellman-Ford algorithm

- ▶  $\text{distTo}[v] = \infty$  for each vertex  $v$ .
- ▶  $\text{edgeTo}[v] = \text{null}$  for each vertex  $v$ .
- ▶  $\text{distTo}[s] = 0$ .
- ▶ Repeat  $|V|-1$  times:
  - ▶ Relax all edges.



<http://algs4.cs.princeton.edu>

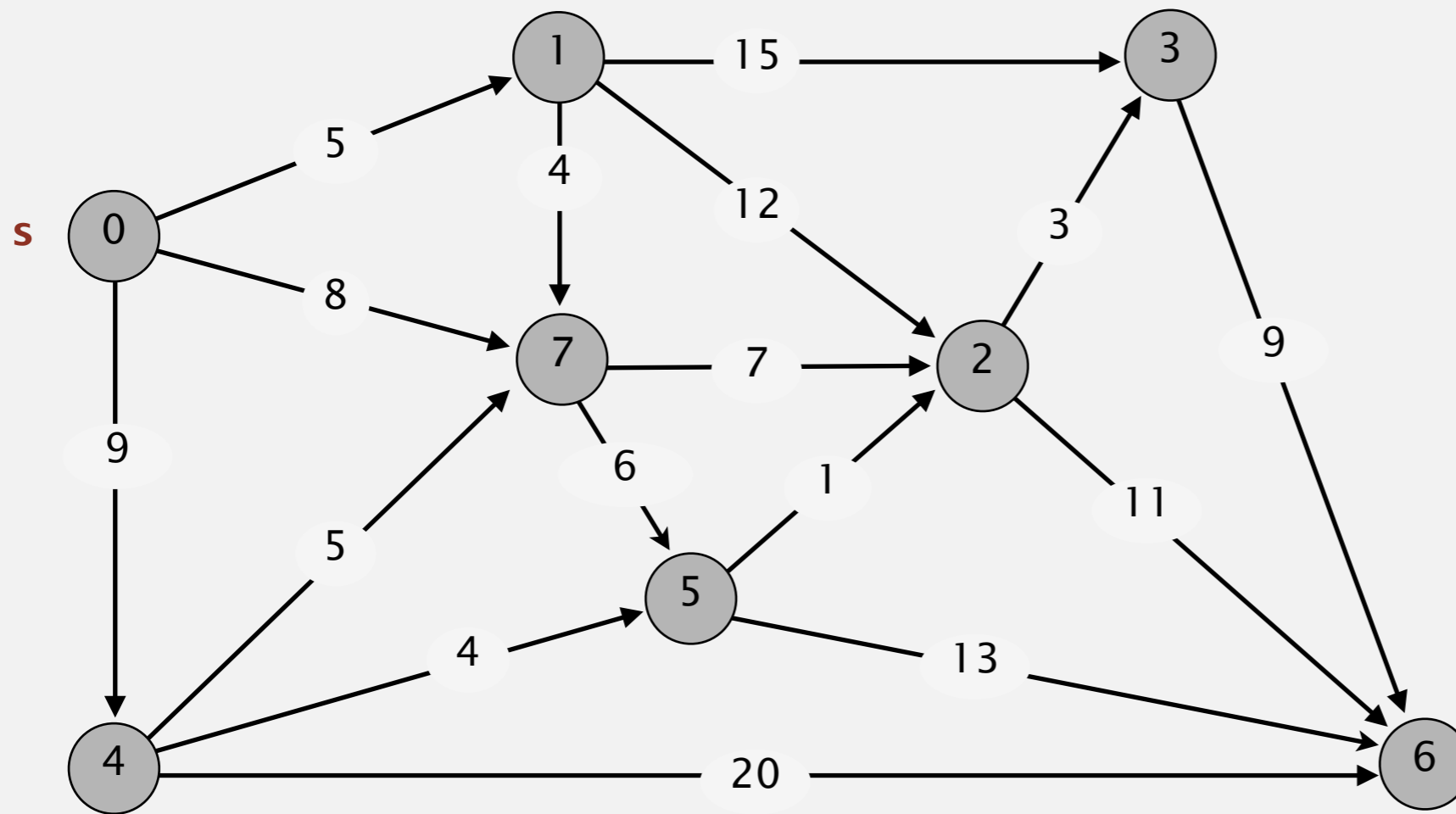
## BELLMAN-FORD DEMO

---



# Bellman-Ford algorithm demo

Repeat  $V$  times: relax all  $E$  edges.



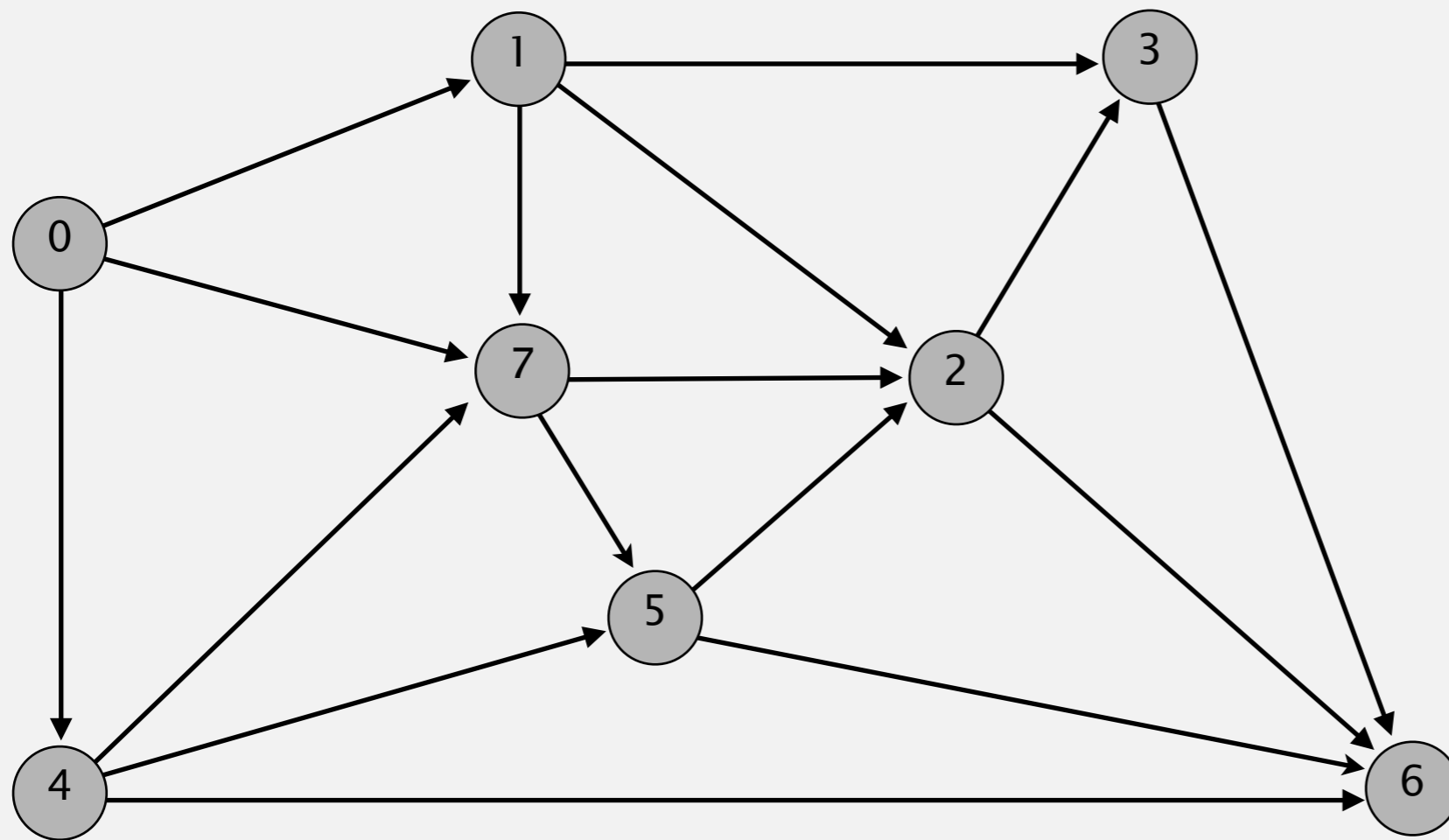
0→1	5.0
0→4	9.0
0→7	8.0
1→2	12.0
1→3	15.0
1→7	4.0
2→3	3.0
2→6	11.0
3→6	9.0
4→5	4.0
4→6	20.0
4→7	5.0
5→2	1.0
5→6	13.0
7→5	6.0
7→2	7.0

**an edge-weighted digraph**

# Bellman-Ford algorithm demo

---

Repeat  $V$  times: relax all  $E$  edges.

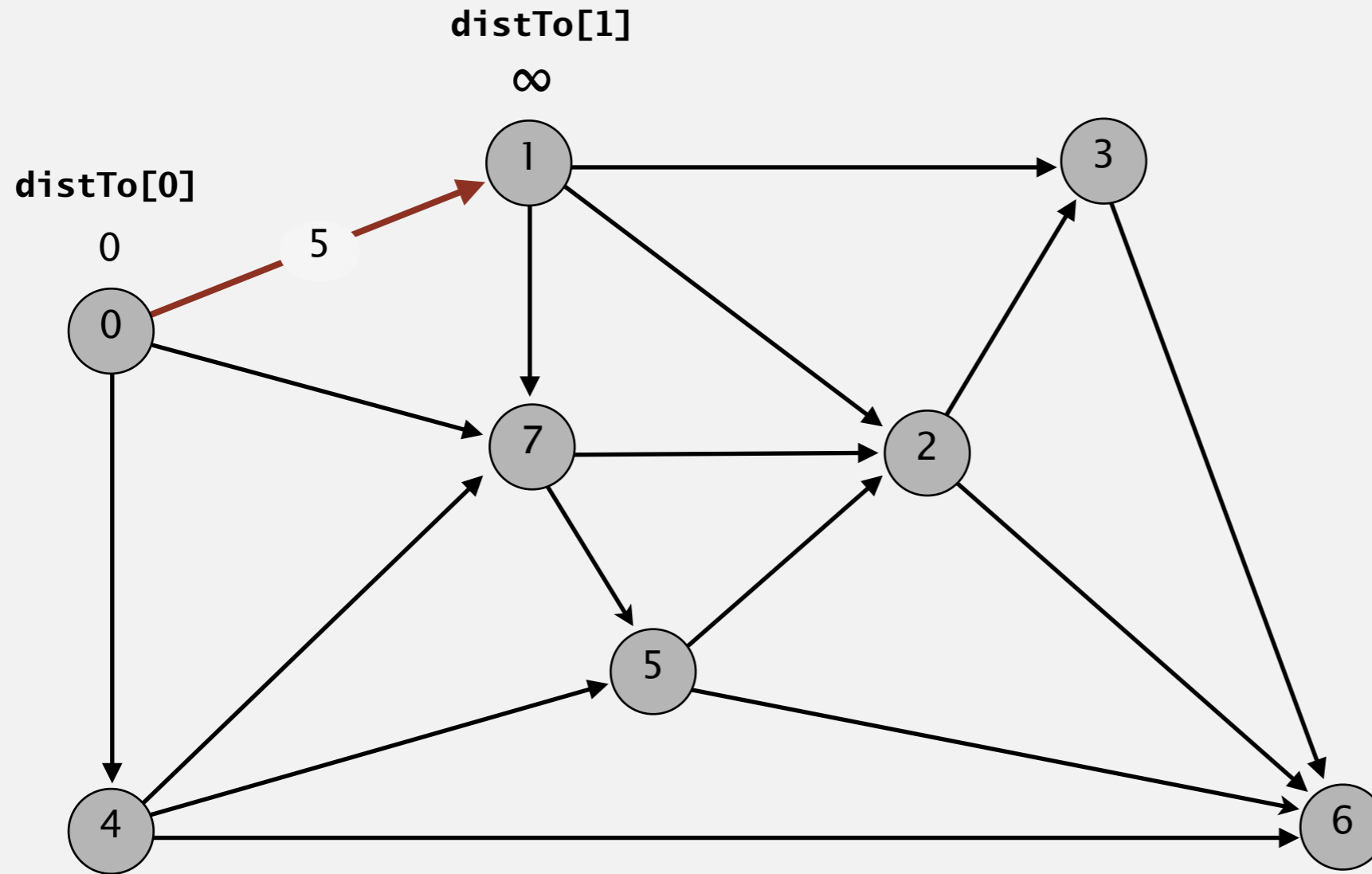


$v$	distTo[]	edgeTo[]
0	0.0	-
1		
2		
3		
4		
5		
6		
7		

**initialize**

# Bellman-Ford algorithm demo

Repeat  $V$  times: relax all  $E$  edges.



$v$	$\text{distTo}[]$	$\text{edgeTo}[]$
0	0.0	-
1		
2		
3		
4		
5		
6		
7		

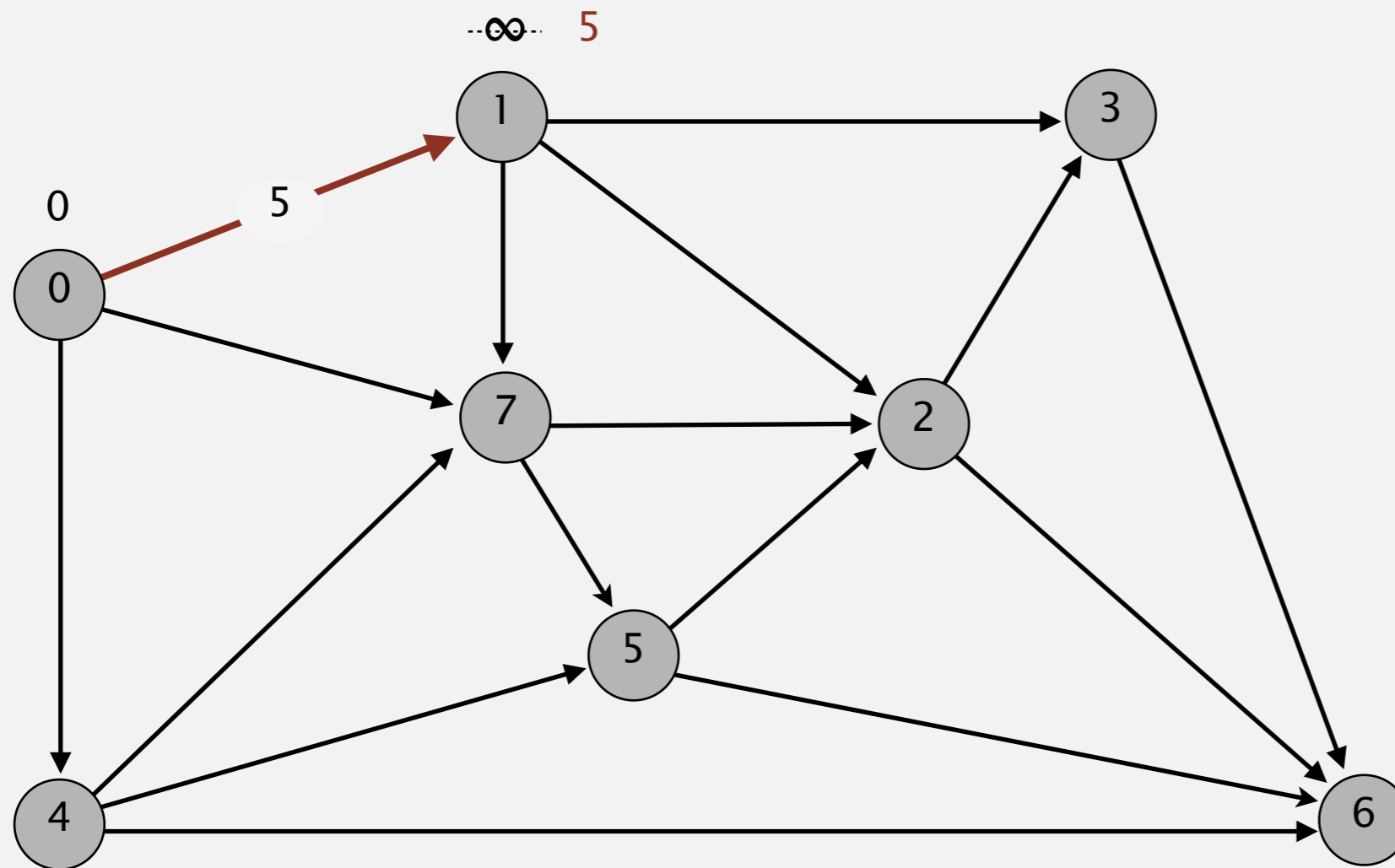
pass 0

0→1 0→4 0→7 1→2 1→3 1→7 2→3 2→6 3→6 4→5 4→6 4→7 5→2 5→6 7→5 7→2



# Bellman-Ford algorithm demo

Repeat  $V$  times: relax all  $E$  edges.



v	distTo[]	edgeTo[]
0	0.0	-
1	5.0	0 → 1
2	$\infty$	-
3	$\infty$	-
4	$\infty$	-
5	$\infty$	-
6	$\infty$	-
7	5.0	0 → 7

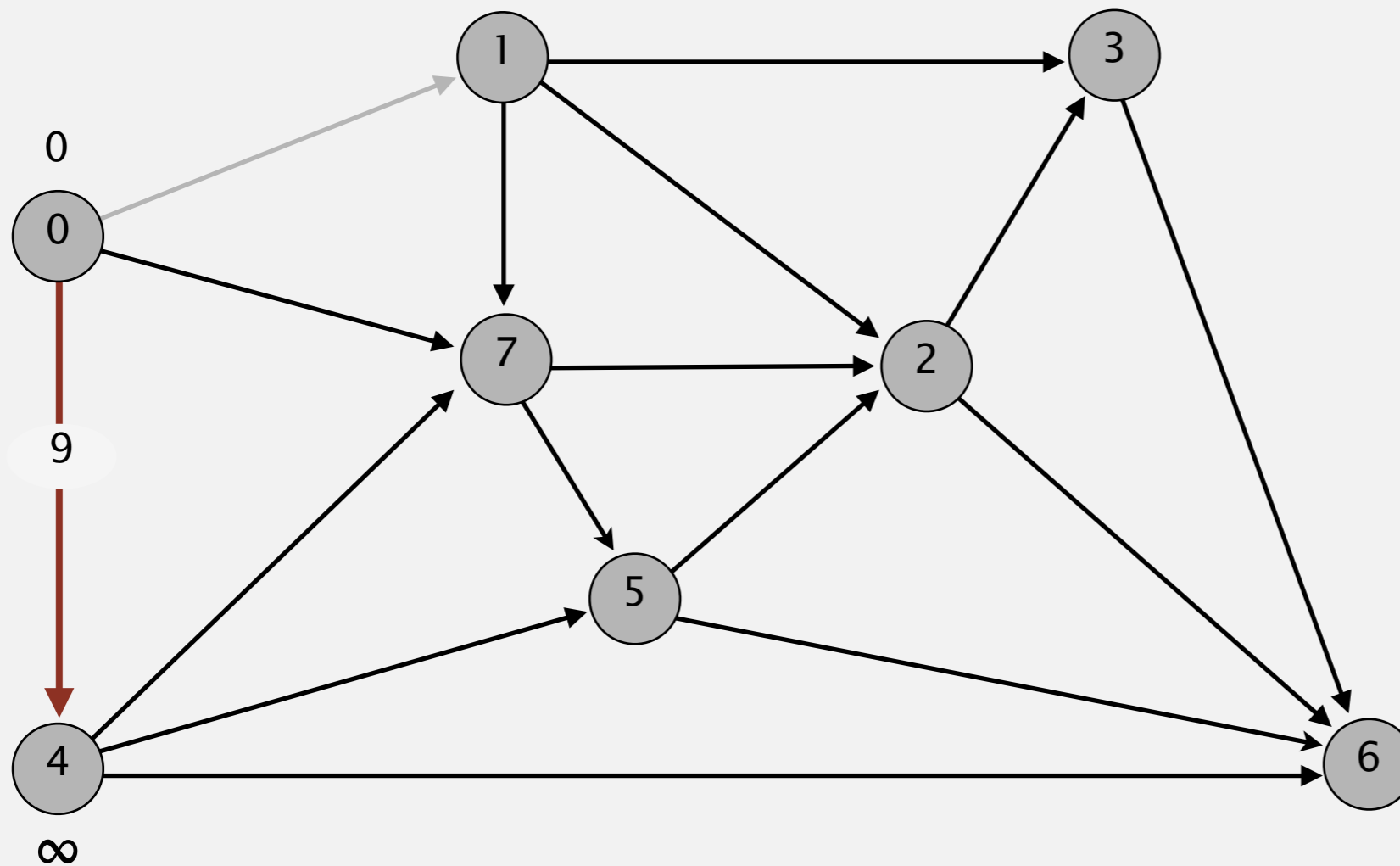
pass 0

0 → 1 0 → 4 0 → 7 1 → 2 1 → 3 1 → 7 2 → 3 2 → 6 3 → 6 4 → 5 4 → 6 4 → 7 5 → 2 5 → 6 7 → 5 7 → 2



# Bellman-Ford algorithm demo

Repeat  $V$  times: relax all  $E$  edges.



v	distTo[]	edgeTo[]
0	0.0	-
1	5.0	0→1
2		
3		
4		
5		
6		
7		

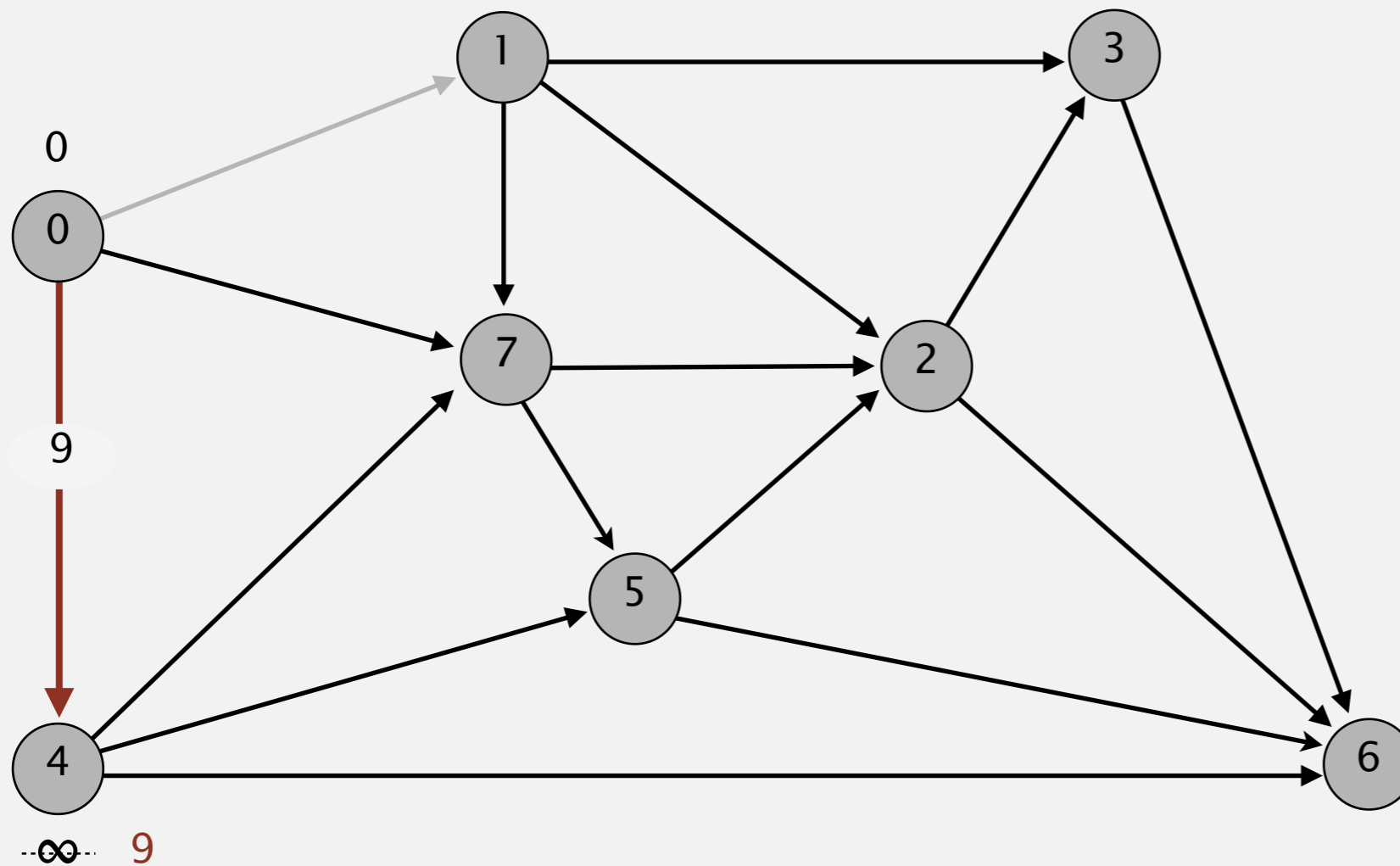
pass 0

0→1 0→4 0→7 1→2 1→3 1→7 2→3 2→6 3→6 4→5 4→6 4→7 5→2 5→6 7→5 7→2



# Bellman-Ford algorithm demo

Repeat  $V$  times: relax all  $E$  edges.



v	distTo[]	edgeTo[]
0	0.0	-
1	5.0	0→1
2		
3		
4	9.0	0→4
5		
6		
7		

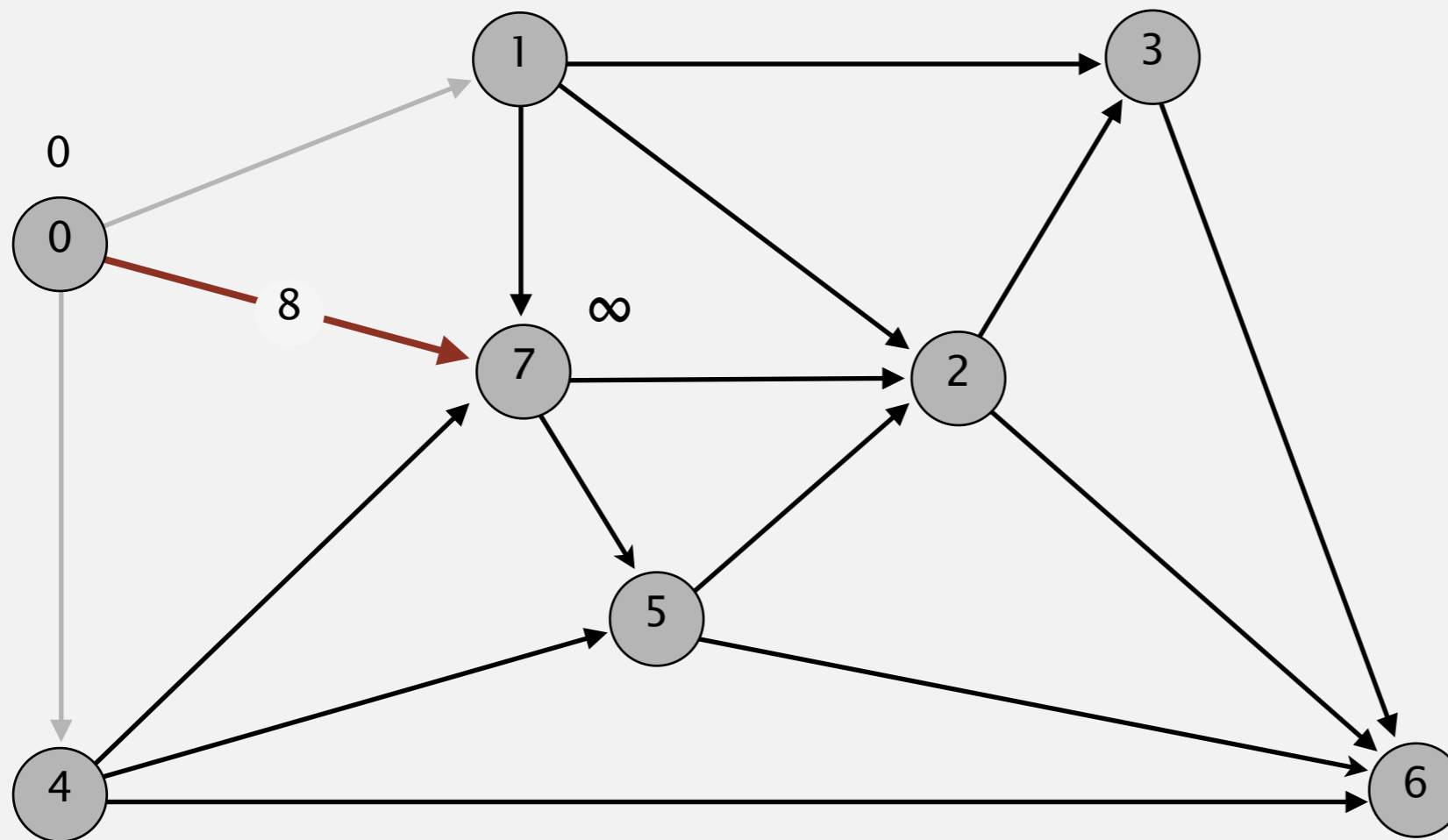
pass 0

0→1 0→4 0→7 1→2 1→3 1→7 2→3 2→6 3→6 4→5 4→6 4→7 5→2 5→6 7→5 7→2



# Bellman-Ford algorithm demo

Repeat  $V$  times: relax all  $E$  edges.



$v$	distTo[]	edgeTo[]
0	0.0	-
1	5.0	0→1
2		
3		
4	9.0	0→4
5		
6		
7		

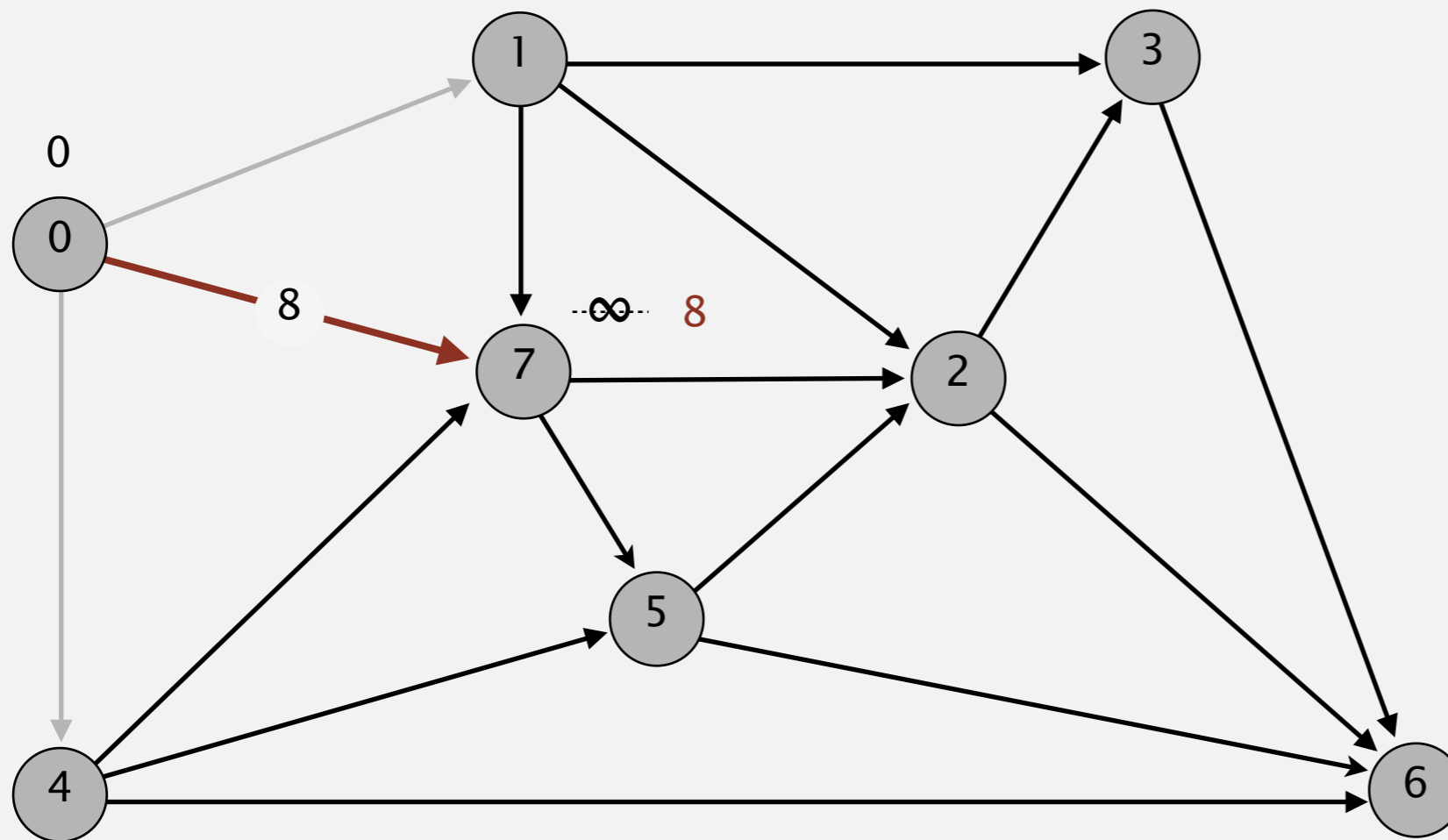
pass 0

0→1 0→4 0→7 1→2 1→3 1→7 2→3 2→6 3→6 4→5 4→6 4→7 5→2 5→6 7→5 7→2



# Bellman-Ford algorithm demo

Repeat  $V$  times: relax all  $E$  edges.



$v$	distTo[]	edgeTo[]
0	0.0	-
1	5.0	0→1
2		
3		
4	9.0	0→4
5		
6		
7	8.0	0→7

pass 0

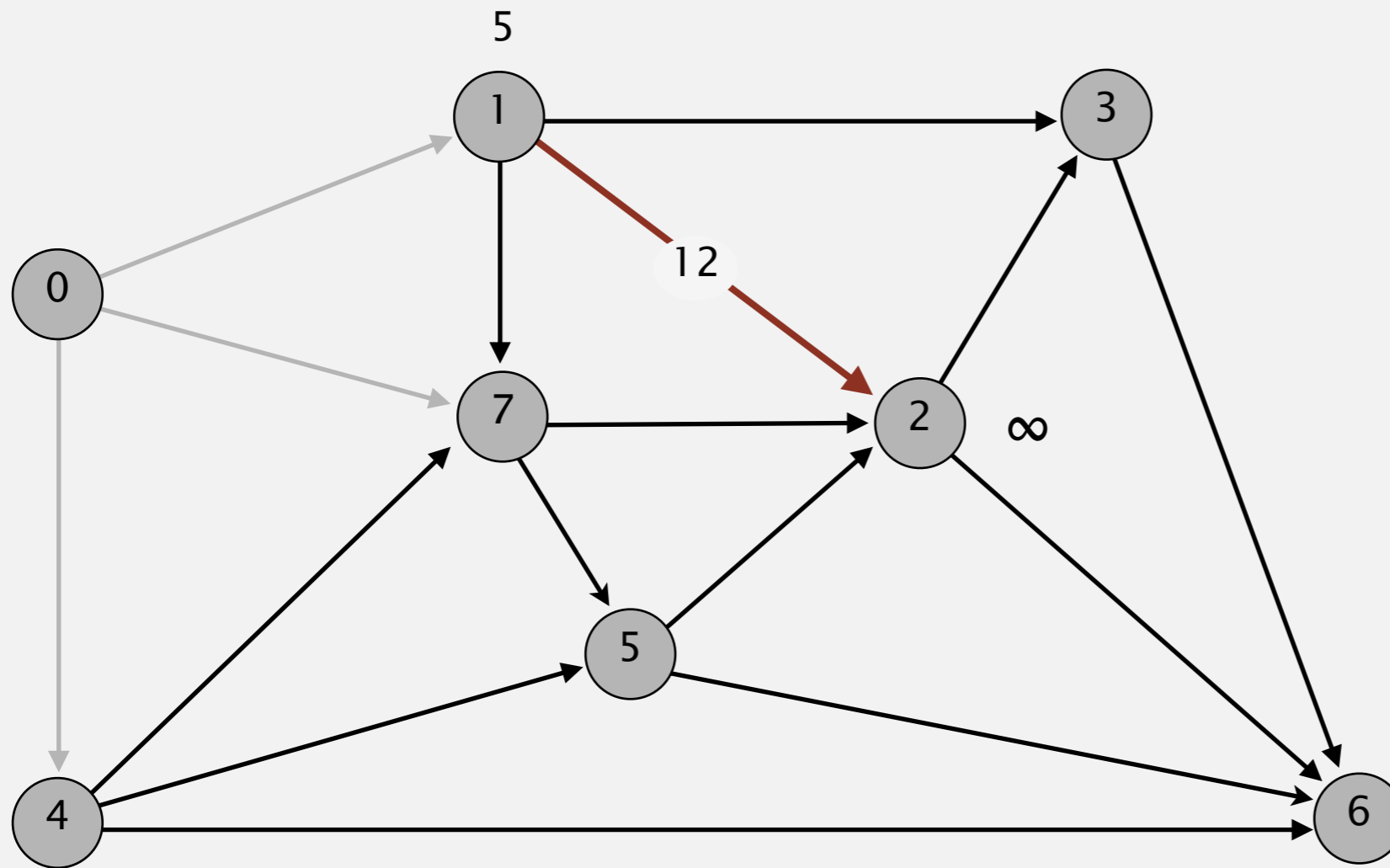
0→1 0→4 0→7 1→2 1→3 1→7 2→3 2→6 3→6 4→5 4→6 4→7 5→2 5→6 7→5 7→2





# Bellman-Ford algorithm demo

Repeat  $V$  times: relax all  $E$  edges.



$v$	distTo[]	edgeTo[]
0	0.0	-
1	5.0	0→1
2		
3		
4	9.0	0→4
5		
6		
7	8.0	0→7

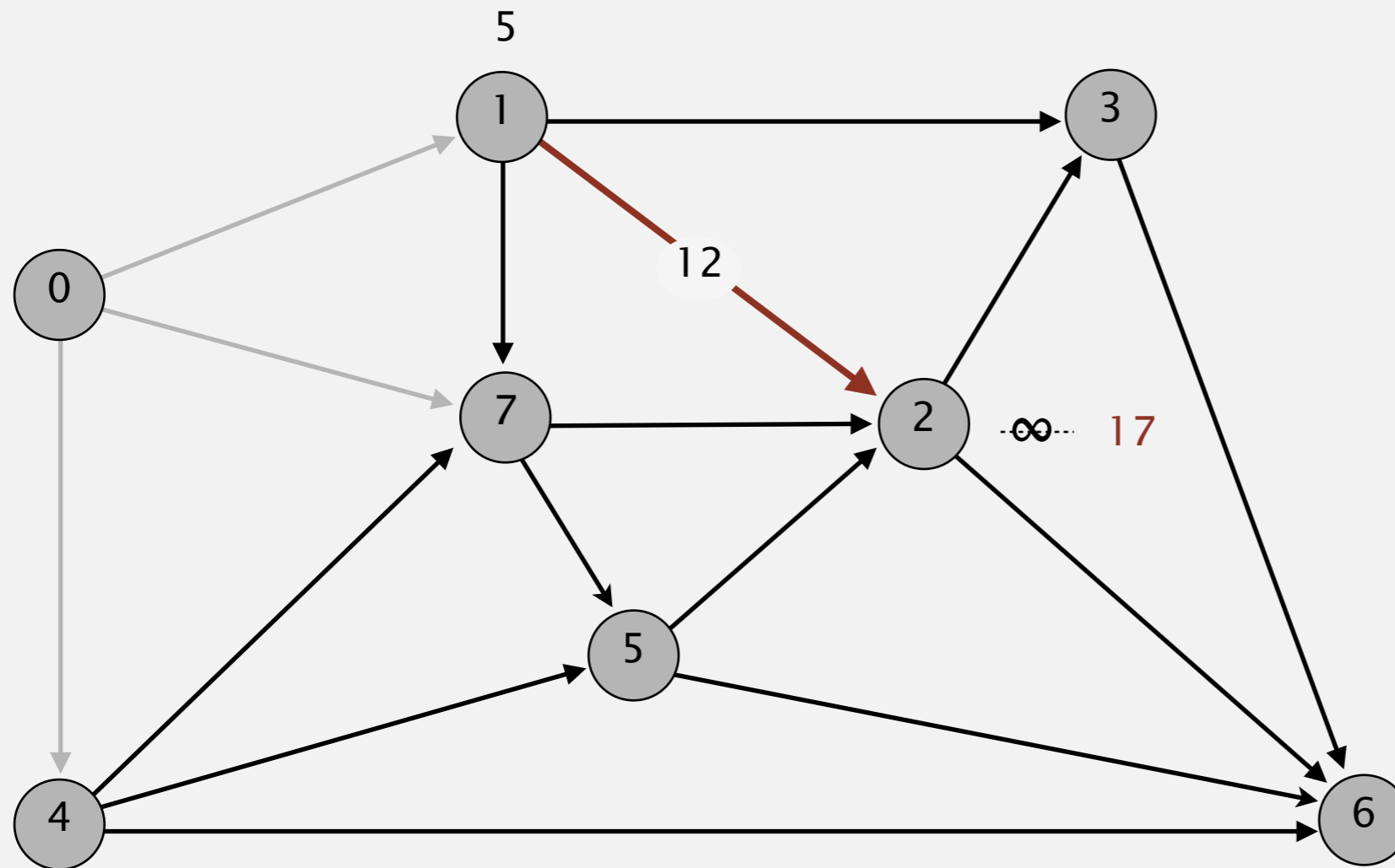
pass 0

0→1 0→4 0→7 1→2 1→3 1→7 2→3 2→6 3→6 4→5 4→6 4→7 5→2 5→6 7→5 7→2



# Bellman-Ford algorithm demo

Repeat  $V$  times: relax all  $E$  edges.



v	distTo[]	edgeTo[]
0	0.0	-
1	5.0	0→1
2	17.0	1→2
3		
4	9.0	0→4
5		
6		
7	8.0	0→7

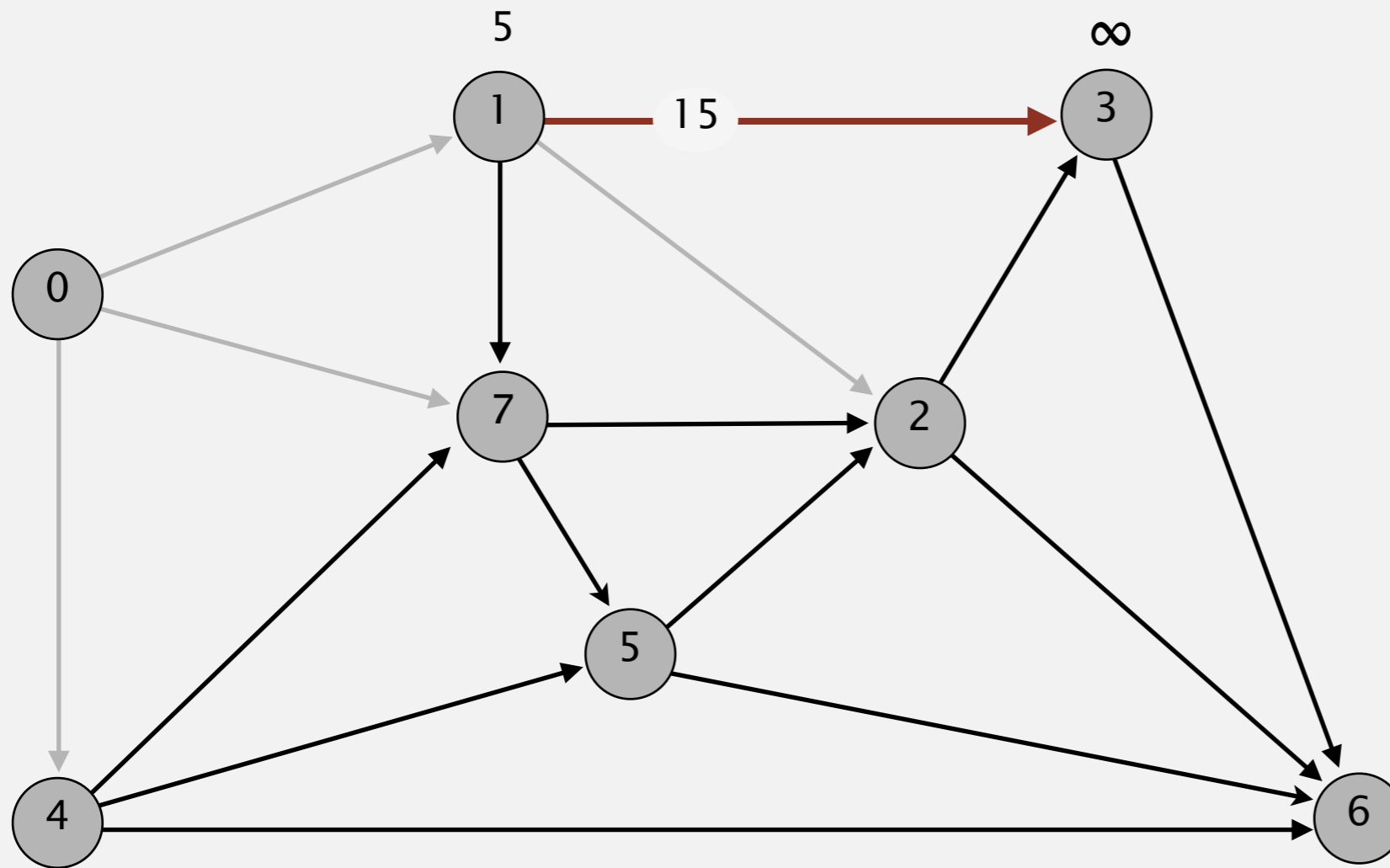
pass 0

0→1 0→4 0→7 1→2 1→3 1→7 2→3 2→6 3→6 4→5 4→6 4→7 5→2 5→6 7→5 7→2



# Bellman-Ford algorithm demo

Repeat  $V$  times: relax all  $E$  edges.



$v$	distTo[]	edgeTo[]
0	0.0	-
1	5.0	0→1
2	17.0	1→2
3	$\infty$	
4	9.0	0→4
5		
6		
7	8.0	0→7

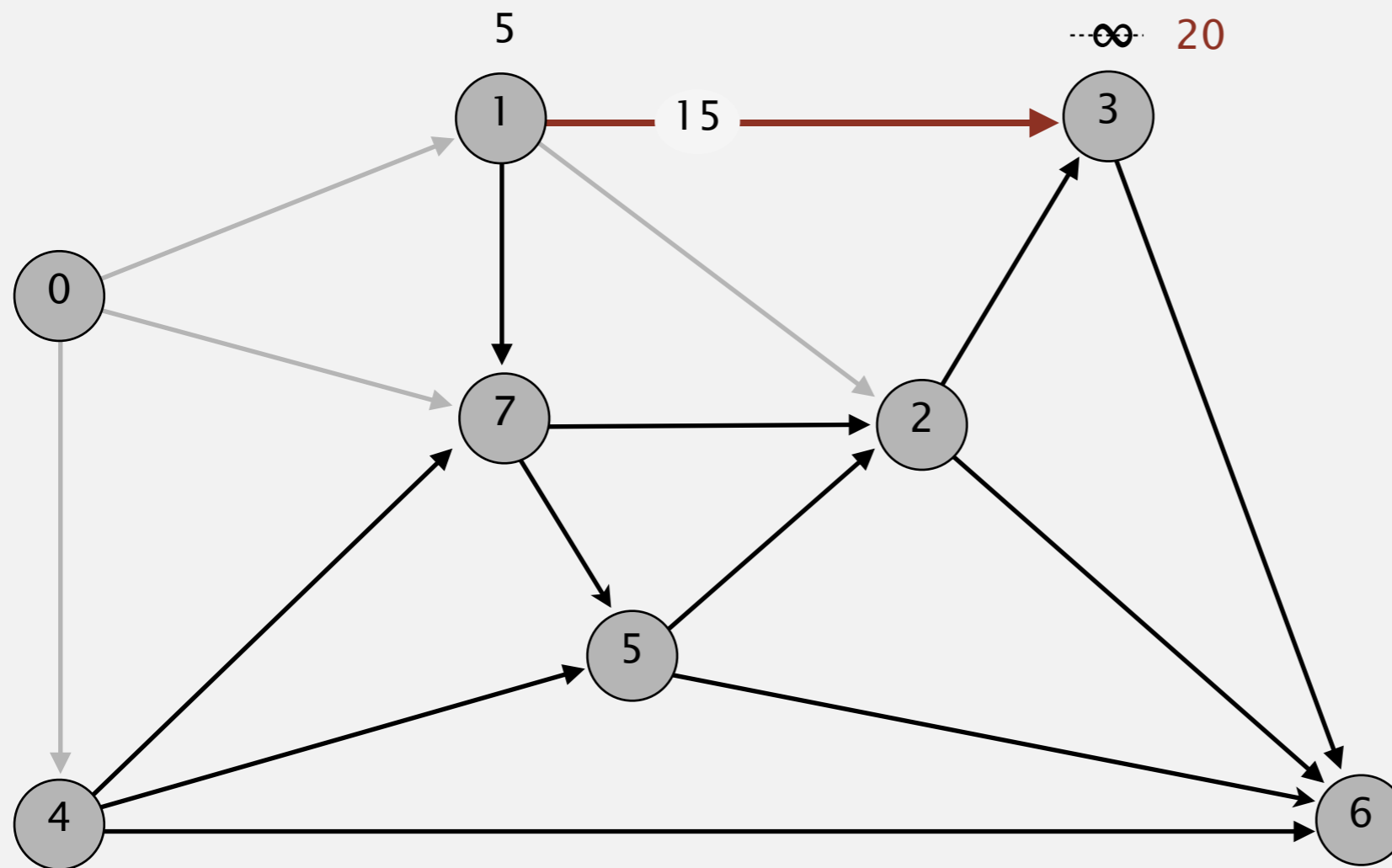
pass 0

0→1 0→4 0→7 1→2 1→3 1→7 2→3 2→6 3→6 4→5 4→6 4→7 5→2 5→6 7→5 7→2



# Bellman-Ford algorithm demo

Repeat  $V$  times: relax all  $E$  edges.



$v$	distTo[]	edgeTo[]
0	0.0	-
1	5.0	0→1
2	17.0	1→2
3	20.0	1→3
4	9.0	0→4
5		
6		
7	8.0	0→7

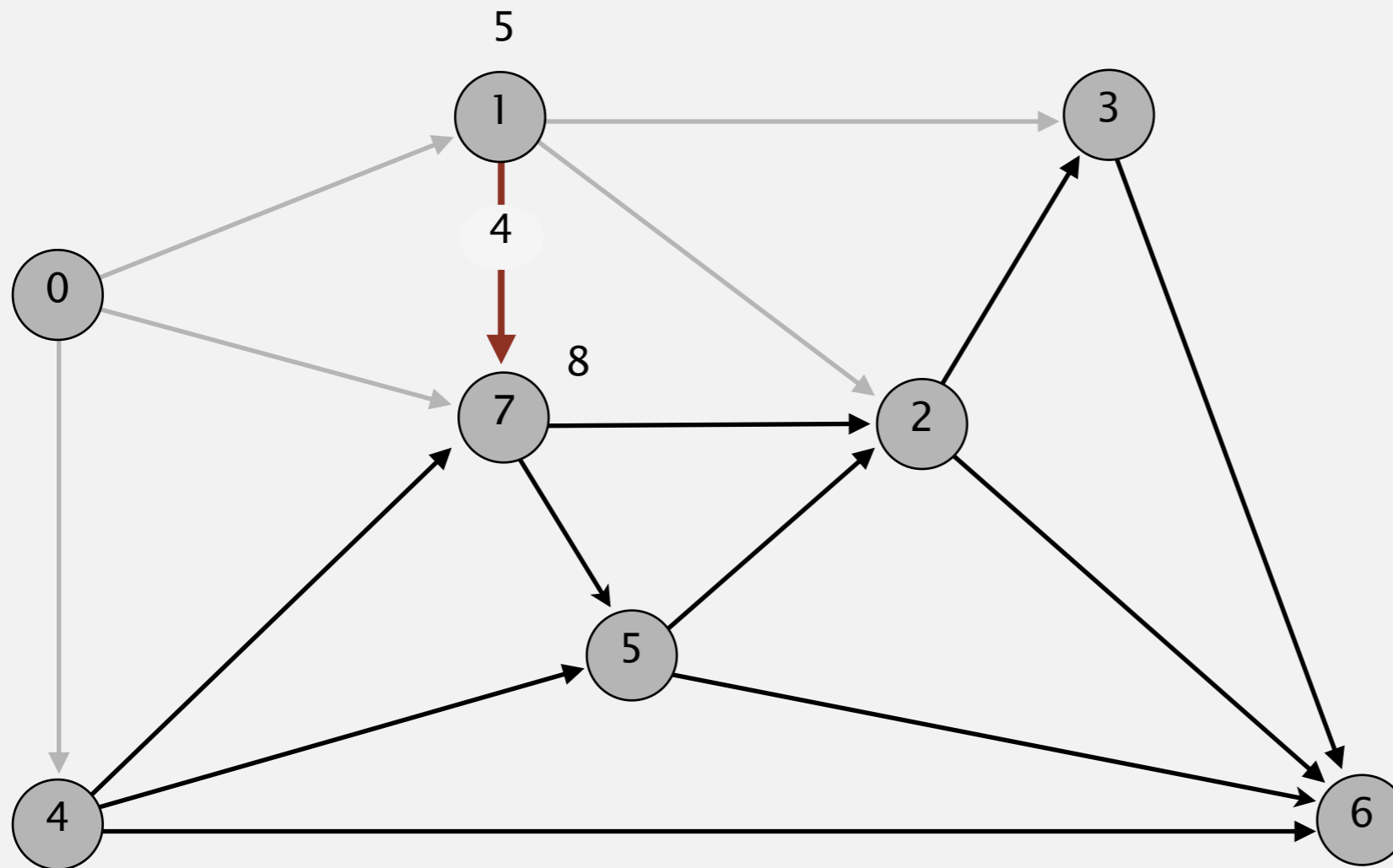
pass 0

0→1 0→4 0→7 1→2 1→3 1→7 2→3 2→6 3→6 4→5 4→6 4→7 5→2 5→6 7→5 7→2



# Bellman-Ford algorithm demo

Repeat  $V$  times: relax all  $E$  edges.



$v$	distTo[]	edgeTo[]
0	0.0	-
1	5.0	0→1
2	17.0	1→2
3	20.0	1→3
4	9.0	0→4
5		
6		
7	8.0	0→7

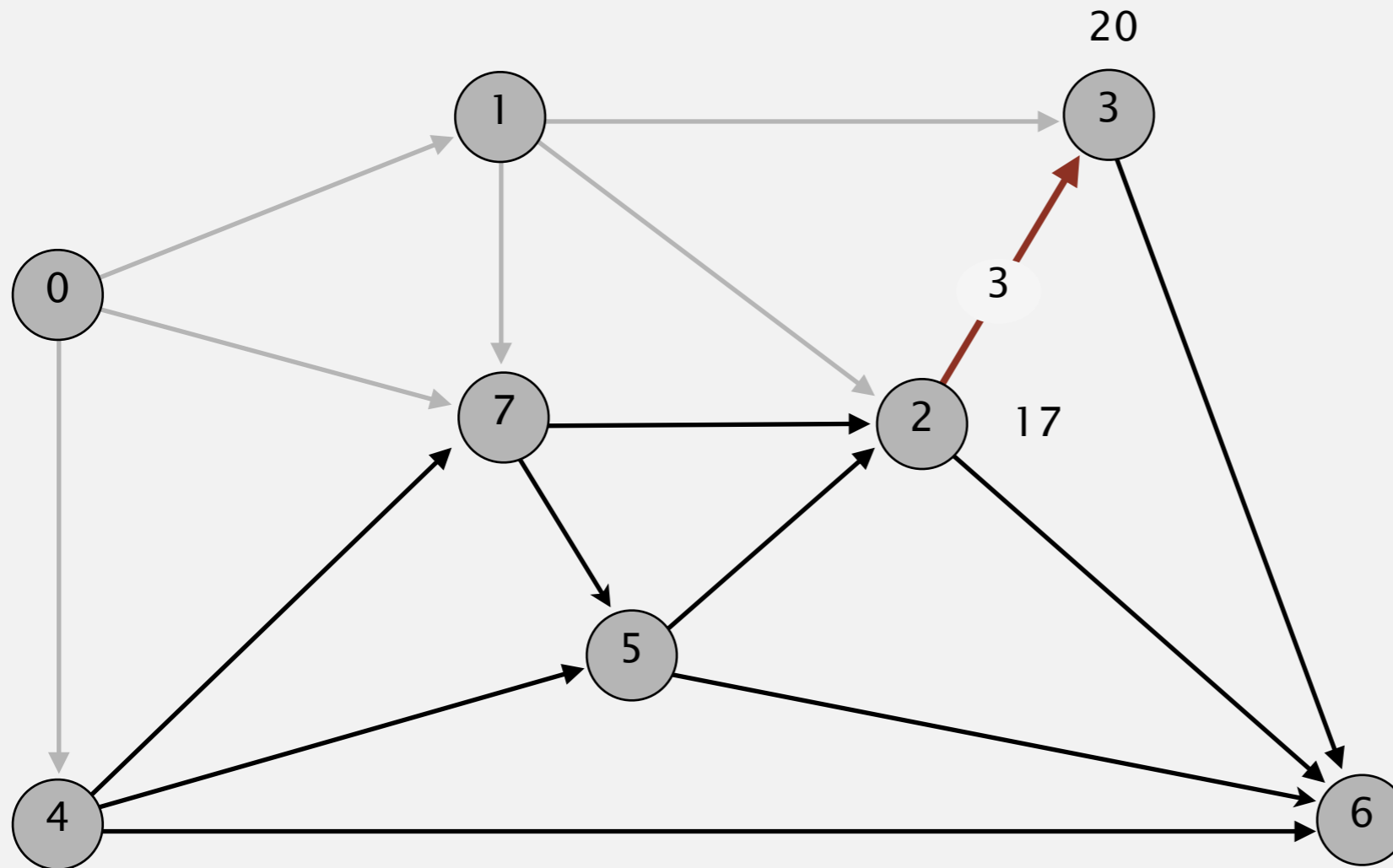
**pass 0**

0→1 0→4 0→7 1→2 1→3 1→7 2→3 2→6 3→6 4→5 4→6 4→7 5→2 5→6 7→5 7→2



# Bellman-Ford algorithm demo

Repeat  $V$  times: relax all  $E$  edges.



$v$	distTo[]	edgeTo[]
0	0.0	-
1	5.0	0→1
2	17.0	1→2
3	20.0	1→3
4	9.0	0→4
5		
6		
7	8.0	0→7

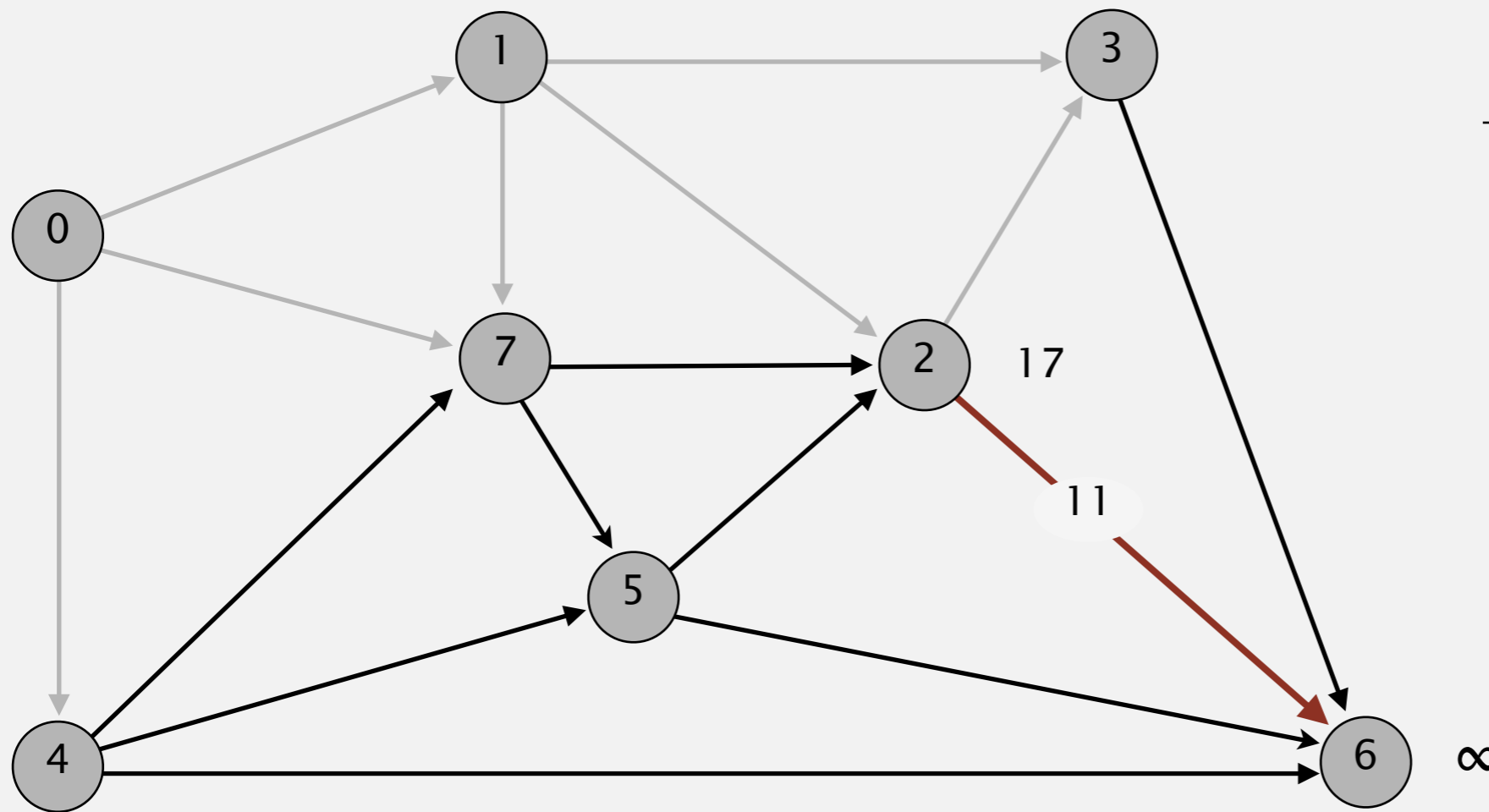
**pass 0**

0→1 0→4 0→7 1→2 1→3 1→7 2→3 2→6 3→6 4→5 4→6 4→7 5→2 5→6 7→5 7→2



# Bellman-Ford algorithm demo

Repeat  $V$  times: relax all  $E$  edges.



v	distTo[]	edgeTo[]
0	0.0	-
1	5.0	0→1
2	17.0	1→2
3	20.0	1→3
4	9.0	0→4
5		
6		
7	8.0	0→7

∞

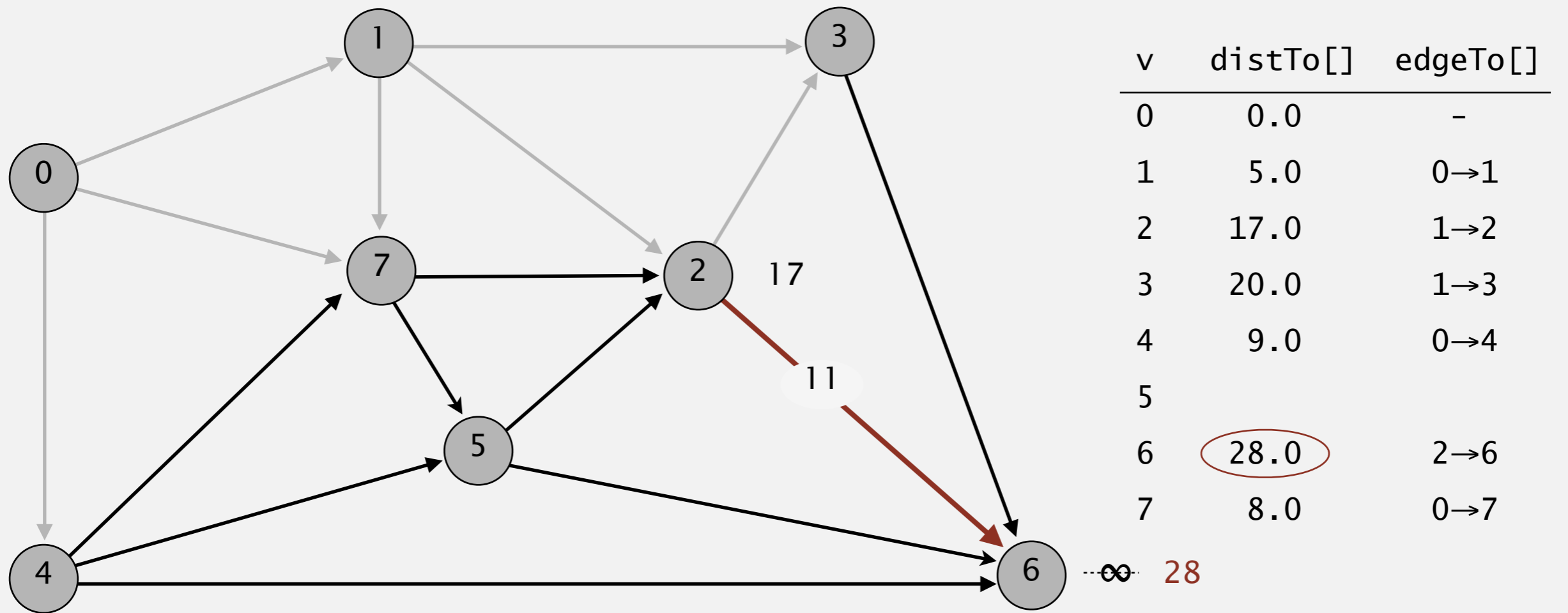
pass 0

0→1 0→4 0→7 1→2 1→3 1→7 2→3 2→6 3→6 4→5 4→6 4→7 5→2 5→6 7→5 7→2



# Bellman-Ford algorithm demo

Repeat  $V$  times: relax all  $E$  edges.



pass 0

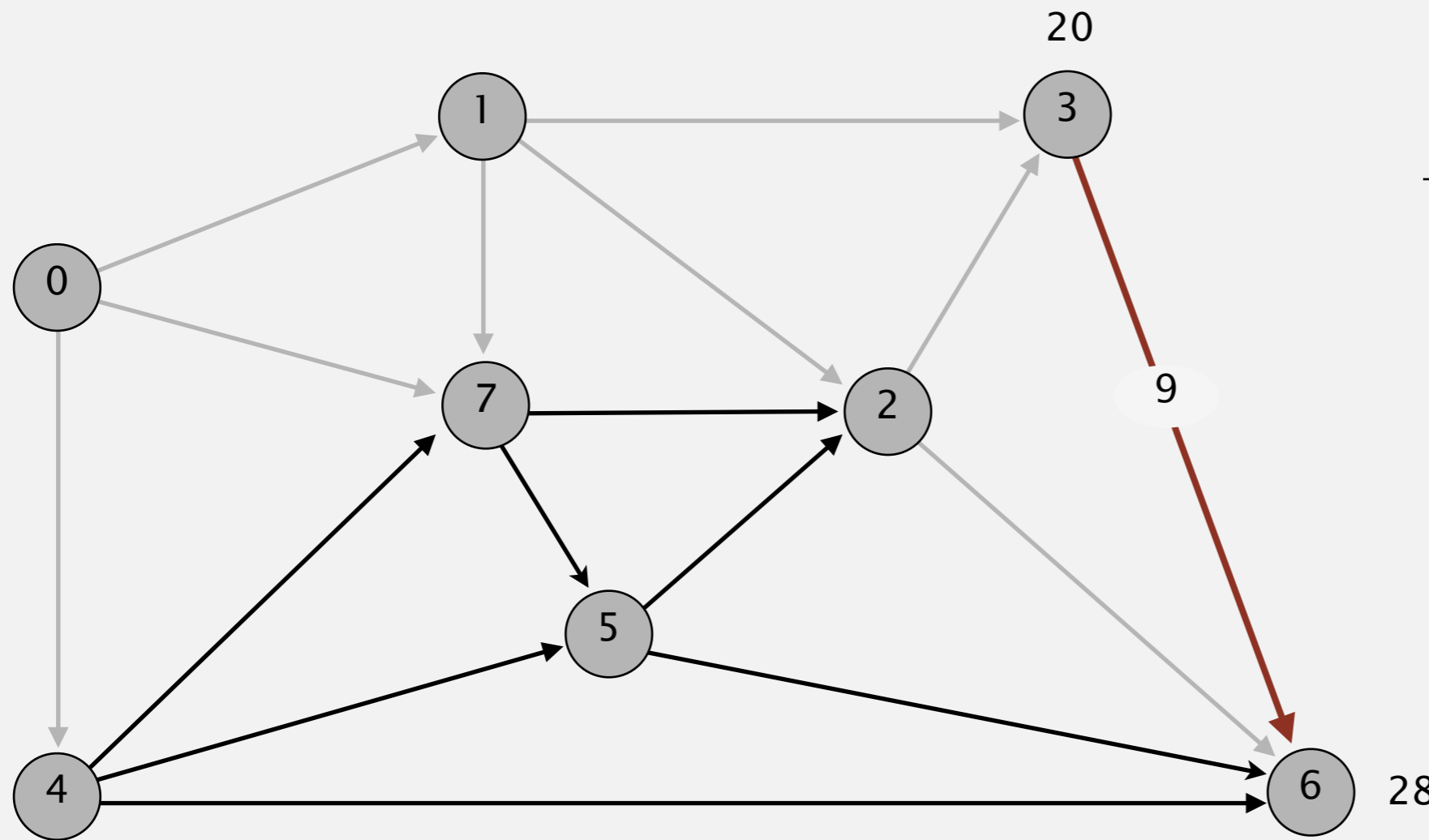
0→1 0→4 0→7 1→2 1→3 1→7 2→3 2→6 3→6 4→5 4→6 4→7 5→2 5→6 7→5 7→2





# Bellman-Ford algorithm demo

Repeat  $V$  times: relax all  $E$  edges.



v	distTo[]	edgeTo[]
0	0.0	-
1	5.0	0→1
2	17.0	1→2
3	20.0	1→3
4	9.0	0→4
5		
6	28.0	2→6
7	8.0	0→7

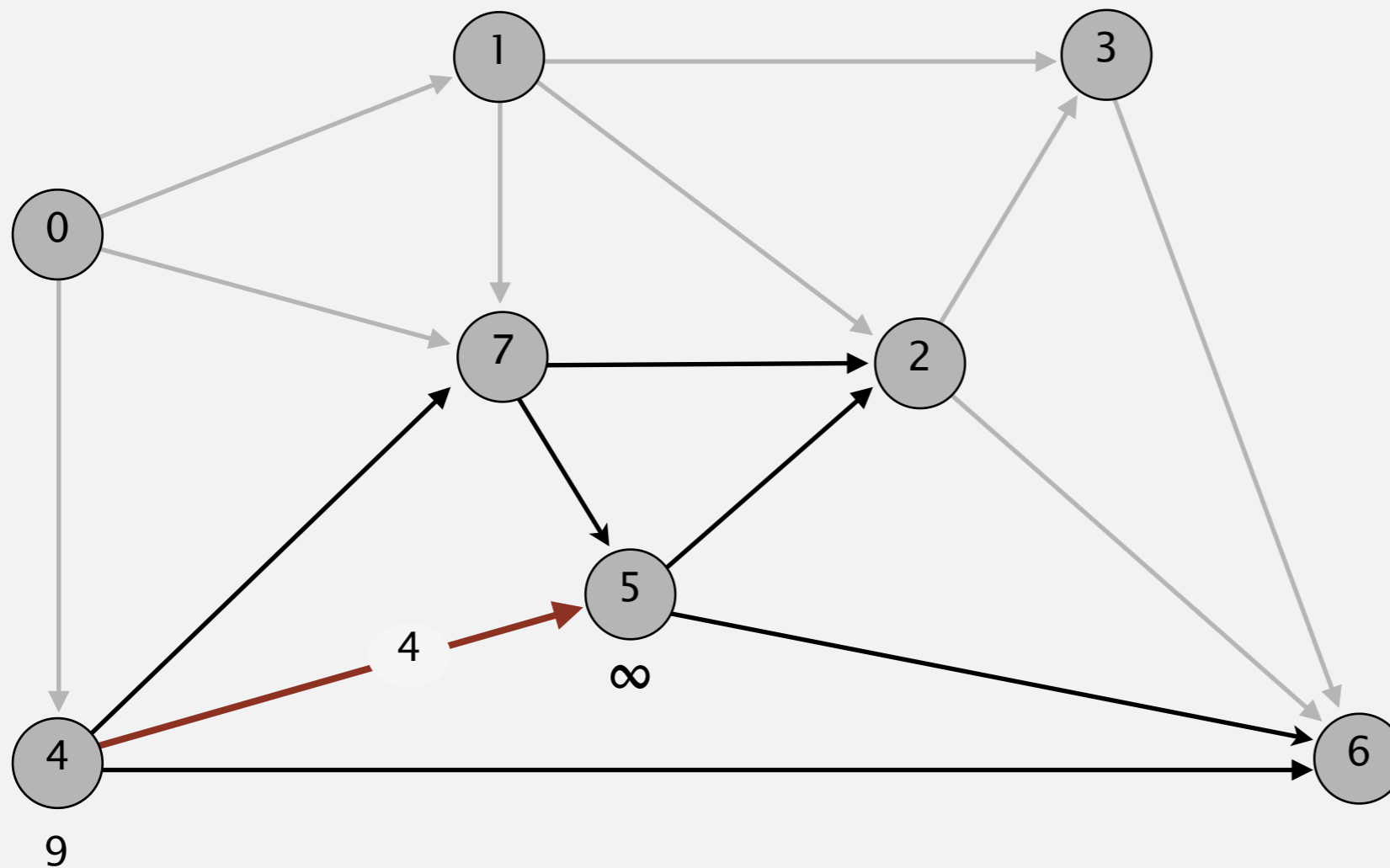
**pass 0**

0→1 0→4 0→7 1→2 1→3 1→7 2→3 2→6 3→6 4→5 4→6 4→7 5→2 5→6 7→5 7→2



# Bellman-Ford algorithm demo

Repeat  $V$  times: relax all  $E$  edges.



$v$	distTo[]	edgeTo[]
0	0.0	-
1	5.0	0→1
2	17.0	1→2
3	20.0	1→3
4	9.0	0→4
5		
6	28.0	2→6
7	8.0	0→7

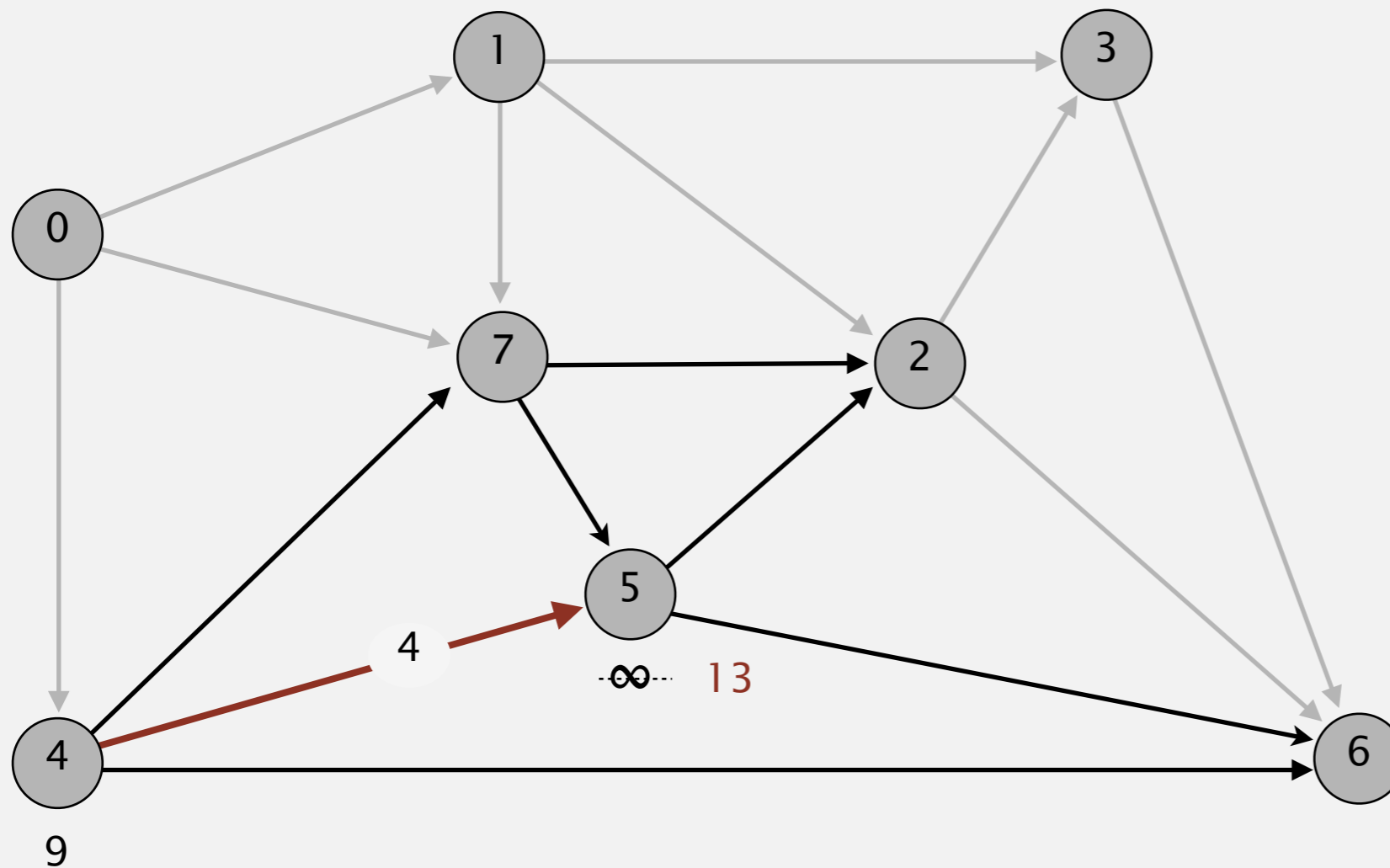
pass 0

0→1 0→4 0→7 1→2 1→3 1→7 2→3 2→6 3→6 4→5 4→6 4→7 5→2 5→6 7→5 7→2



# Bellman-Ford algorithm demo

Repeat  $V$  times: relax all  $E$  edges.



v	distTo[]	edgeTo[]
0	0.0	-
1	5.0	0→1
2	17.0	1→2
3	20.0	1→3
4	9.0	0→4
5	13.0	4→5
6	28.0	2→6
7	8.0	0→7

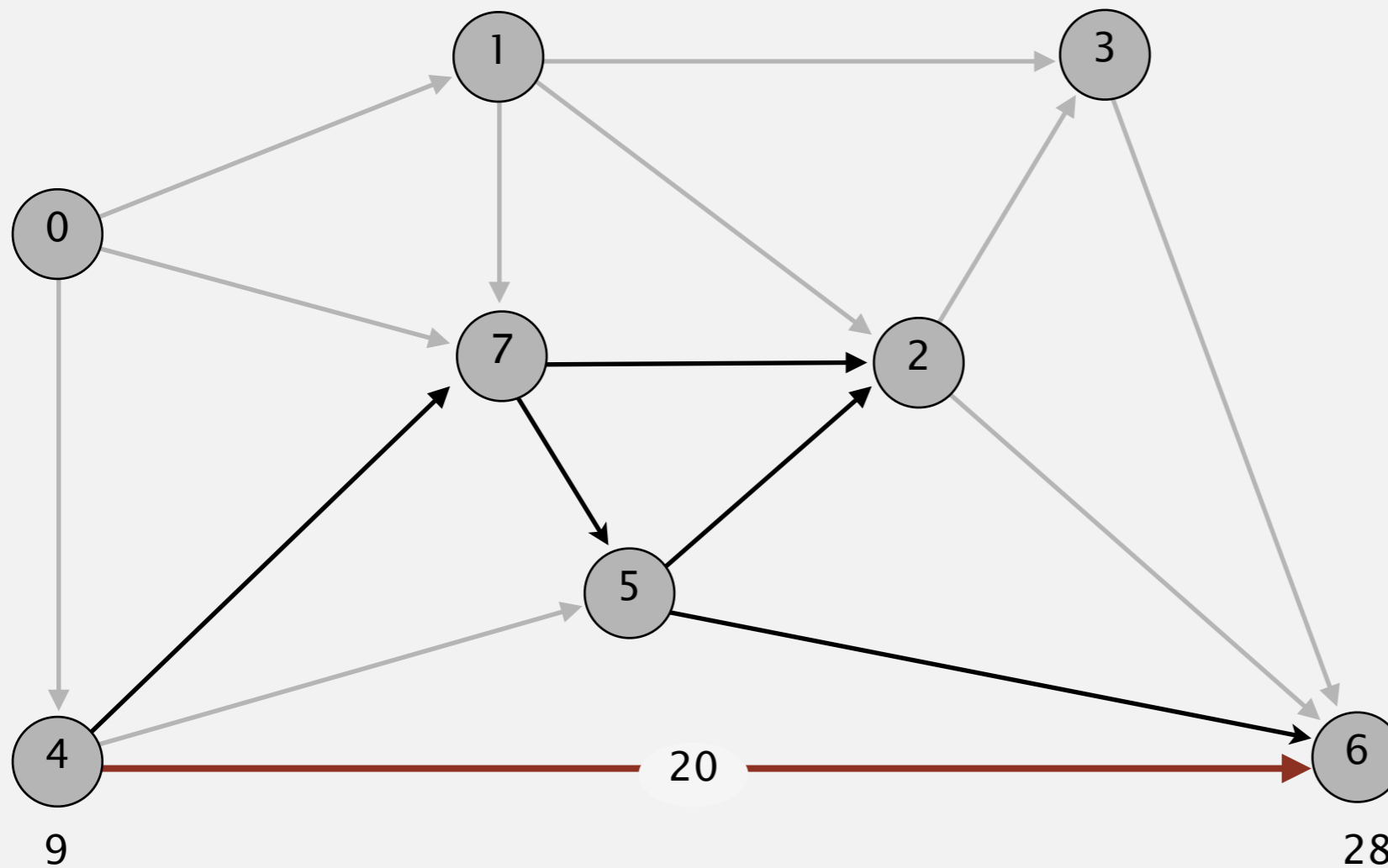
pass 0

0→1 0→4 0→7 1→2 1→3 1→7 2→3 2→6 3→6 4→5 4→6 4→7 5→2 5→6 7→5 7→2



# Bellman-Ford algorithm demo

Repeat  $V$  times: relax all  $E$  edges.



$v$	distTo[]	edgeTo[]
0	0.0	-
1	5.0	0→1
2	17.0	1→2
3	20.0	1→3
4	9.0	0→4
5	13.0	4→5
6	28.0	2→6
7	8.0	0→7

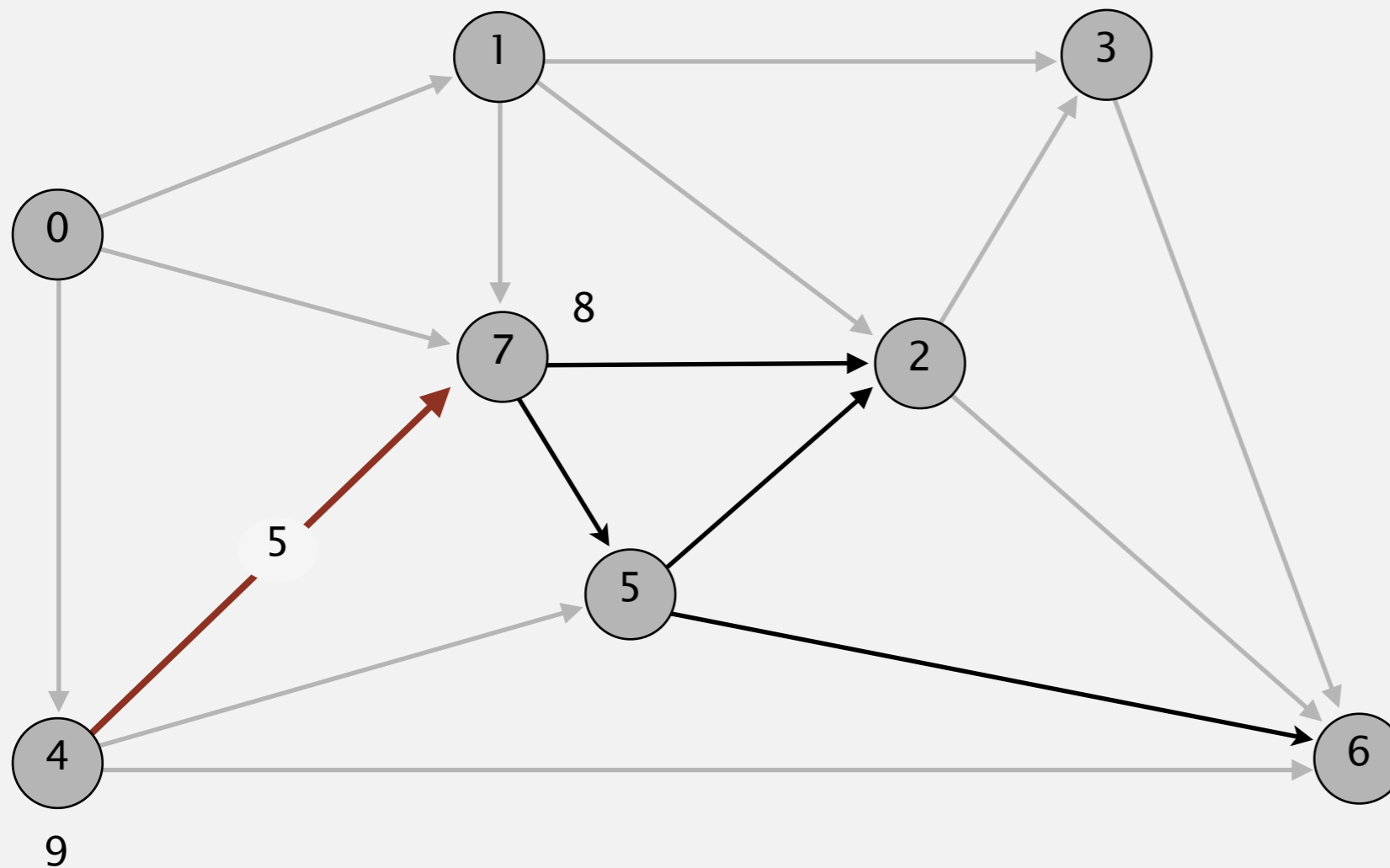
pass 0

0→1 0→4 0→7 1→2 1→3 1→7 2→3 2→6 3→6 4→5 4→6 4→7 5→2 5→6 7→5 7→2



# Bellman-Ford algorithm demo

Repeat  $V$  times: relax all  $E$  edges.



v	distTo[]	edgeTo[]
0	0.0	-
1	5.0	0→1
2	17.0	1→2
3	20.0	1→3
4	9.0	0→4
5	13.0	4→5
6	28.0	2→6
7	8.0	0→7

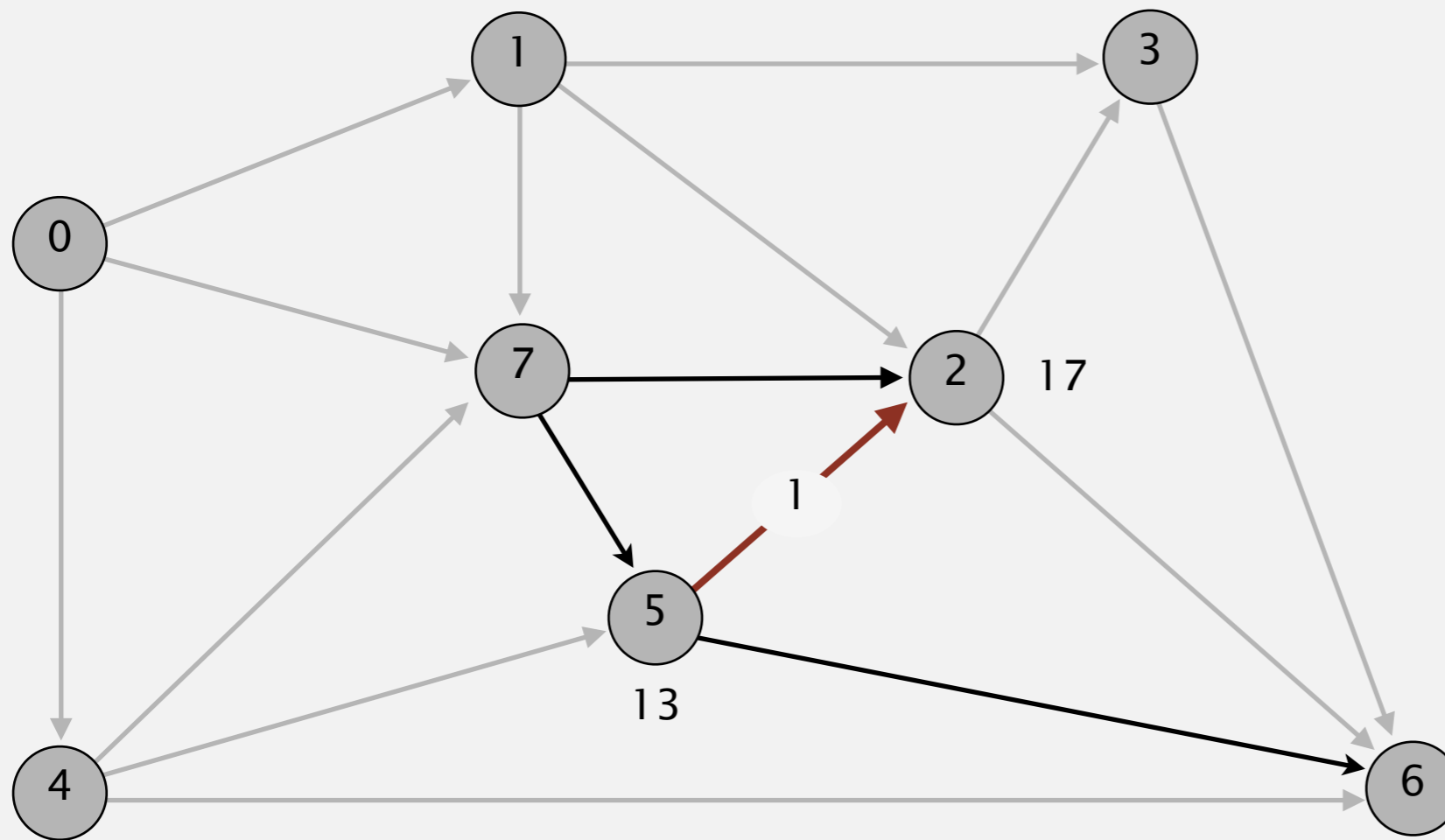
**pass 0**

0→1 0→4 0→7 1→2 1→3 1→7 2→3 2→6 3→6 4→5 4→6 4→7 5→2 5→6 7→5 7→2



# Bellman-Ford algorithm demo

Repeat  $V$  times: relax all  $E$  edges.



v	distTo[]	edgeTo[]
0	0.0	-
1	5.0	0→1
2	17.0	1→2
3	20.0	1→3
4	9.0	0→4
5	13.0	4→5
6	28.0	2→6
7	8.0	0→7

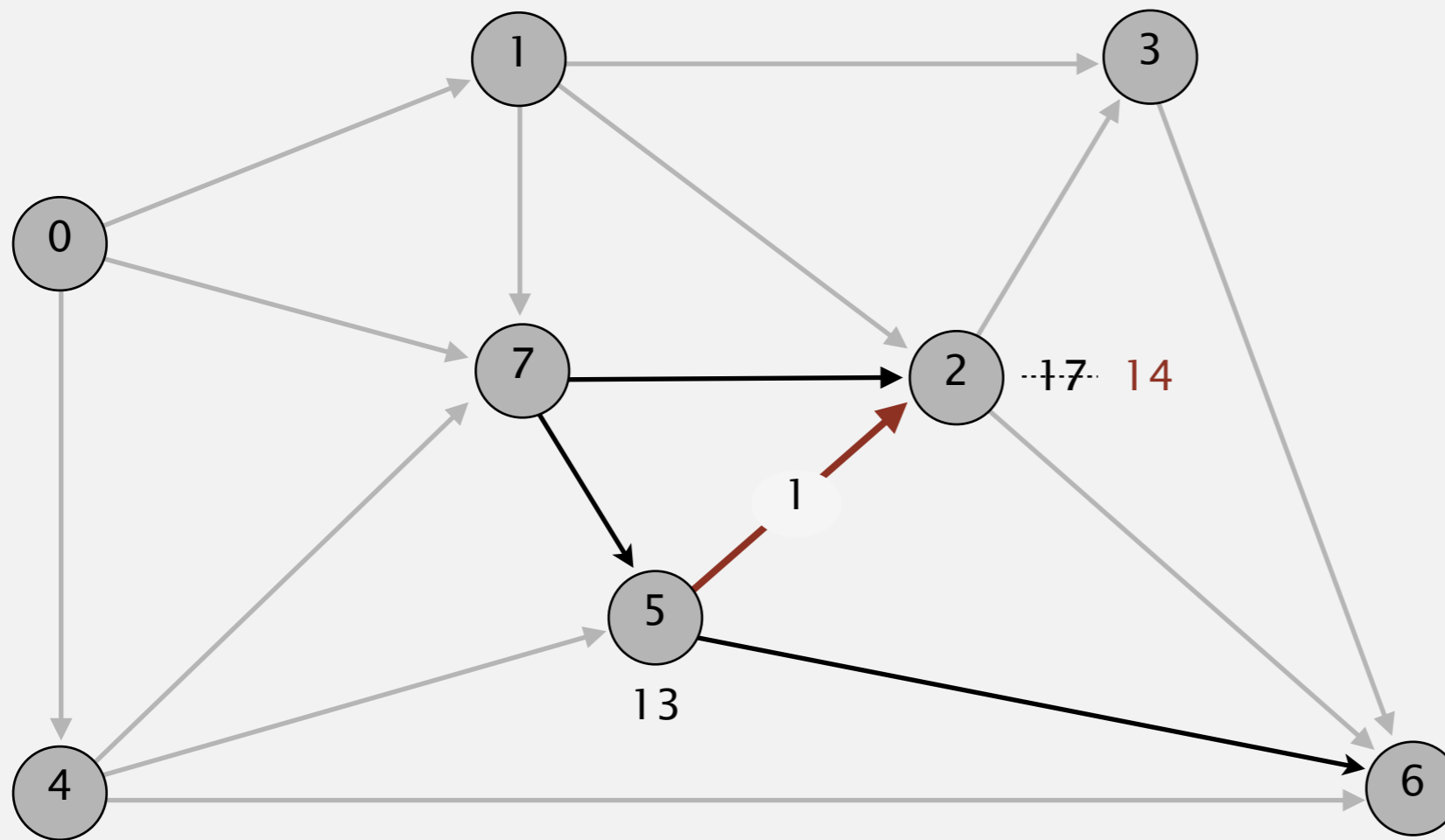
**pass 0**

0→1 0→4 0→7 1→2 1→3 1→7 2→3 2→6 3→6 4→5 4→6 4→7 5→2 5→6 7→5 7→2



# Bellman-Ford algorithm demo

Repeat  $V$  times: relax all  $E$  edges.



v	distTo[]	edgeTo[]
0	0.0	-
1	5.0	0→1
2	14.0	5→2
3	20.0	1→3
4	9.0	0→4
5	13.0	4→5
6	28.0	2→6
7	8.0	0→7

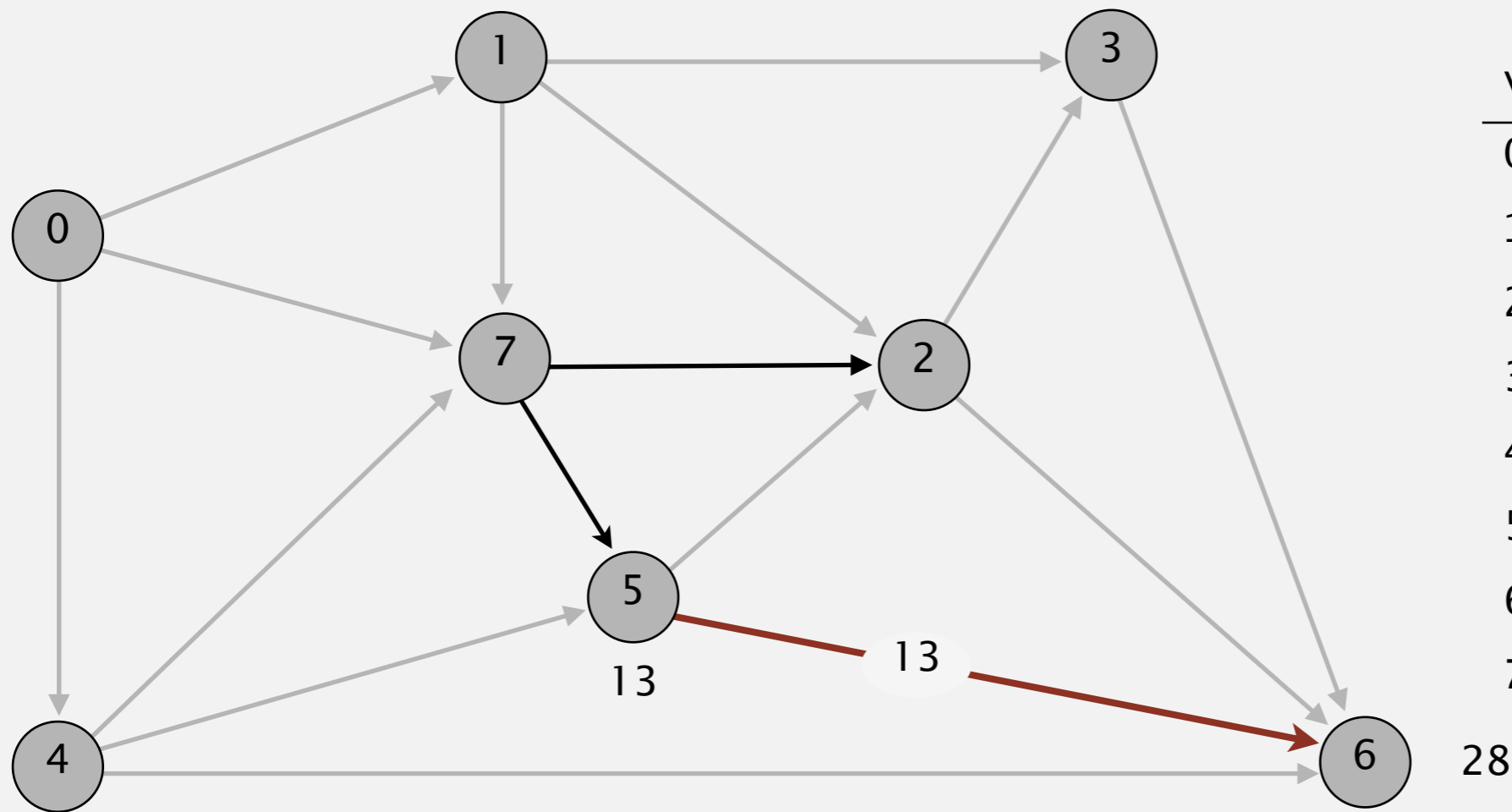
pass 0

0→1 0→4 0→7 1→2 1→3 1→7 2→3 2→6 3→6 4→5 4→6 4→7 5→2 5→6 7→5 7→2



# Bellman-Ford algorithm demo

Repeat  $V$  times: relax all  $E$  edges.



v	distTo[]	edgeTo[]
0	0.0	-
1	5.0	0→1
2	14.0	5→2
3	20.0	1→3
4	9.0	0→4
5	13.0	4→5
6	28.0	2→6
7	8.0	0→7

pass 0

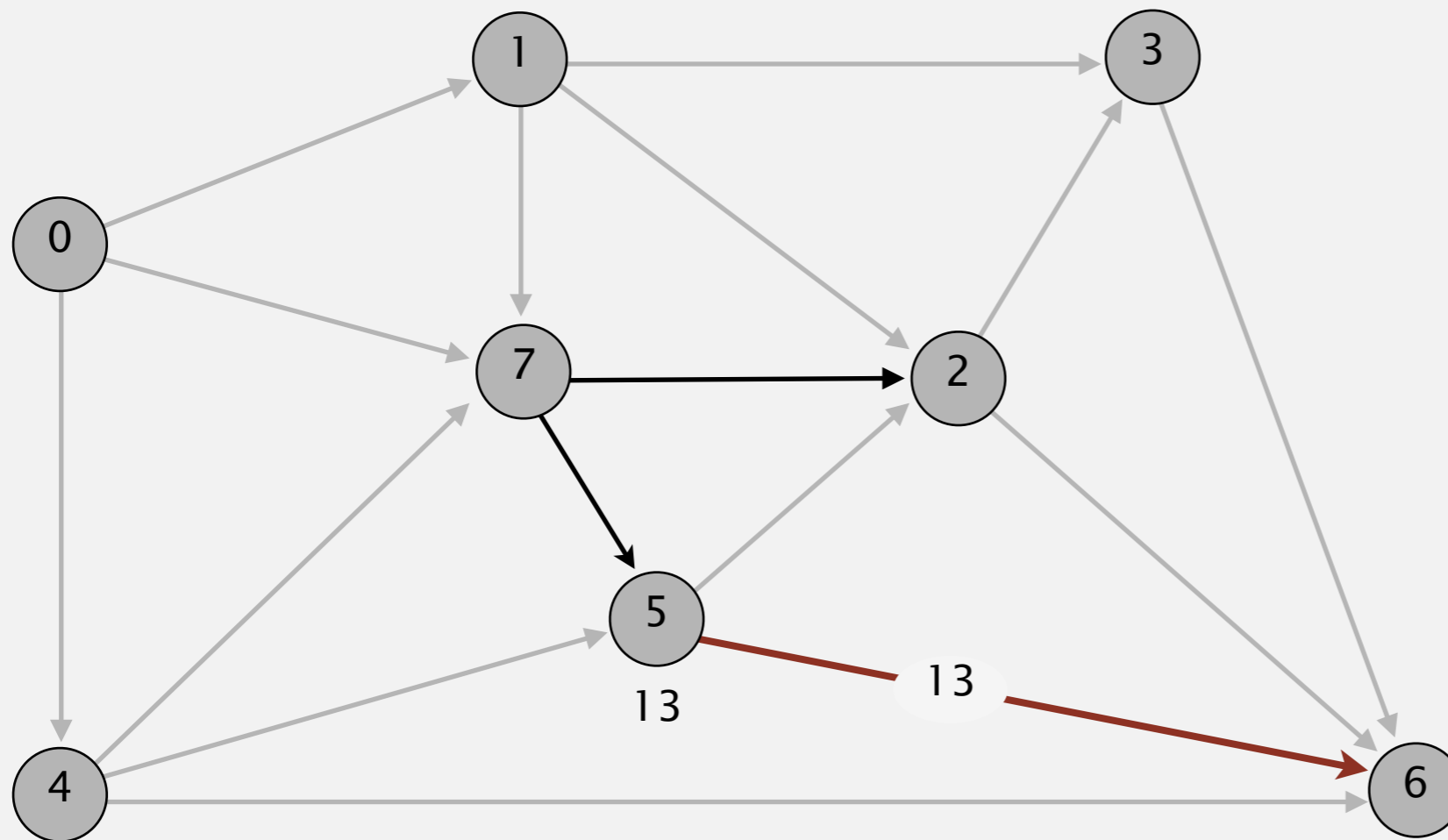
0→1 0→4 0→7 1→2 1→3 1→7 2→3 2→6 3→6 4→5 4→6 4→7 5→2 5→6 7→5 7→2





# Bellman-Ford algorithm demo

Repeat  $V$  times: relax all  $E$  edges.



v	distTo[]	edgeTo[]
0	0.0	-
1	5.0	0→1
2	14.0	5→2
3	20.0	1→3
4	9.0	0→4
5	13.0	4→5
6	26.0	5→6
7	8.0	0→7

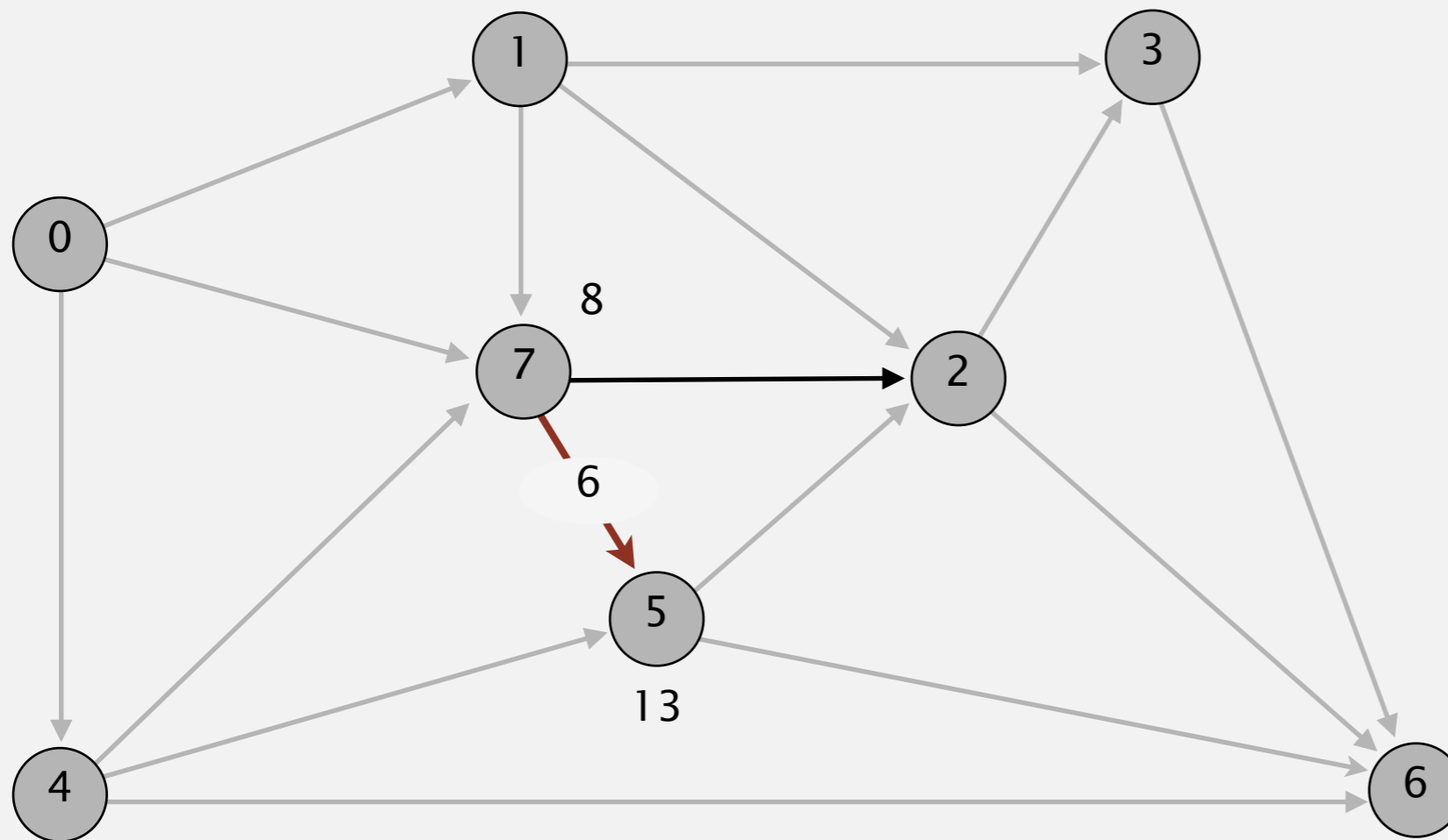
pass 0

0→1 0→4 0→7 1→2 1→3 1→7 2→3 2→6 3→6 4→5 4→6 4→7 5→2 5→6 7→5 7→2



# Bellman-Ford algorithm demo

Repeat  $V$  times: relax all  $E$  edges.



v	distTo[]	edgeTo[]
0	0.0	-
1	5.0	0→1
2	14.0	5→2
3	20.0	1→3
4	9.0	0→4
5	13.0	4→5
6	26.0	5→6
7	8.0	0→7

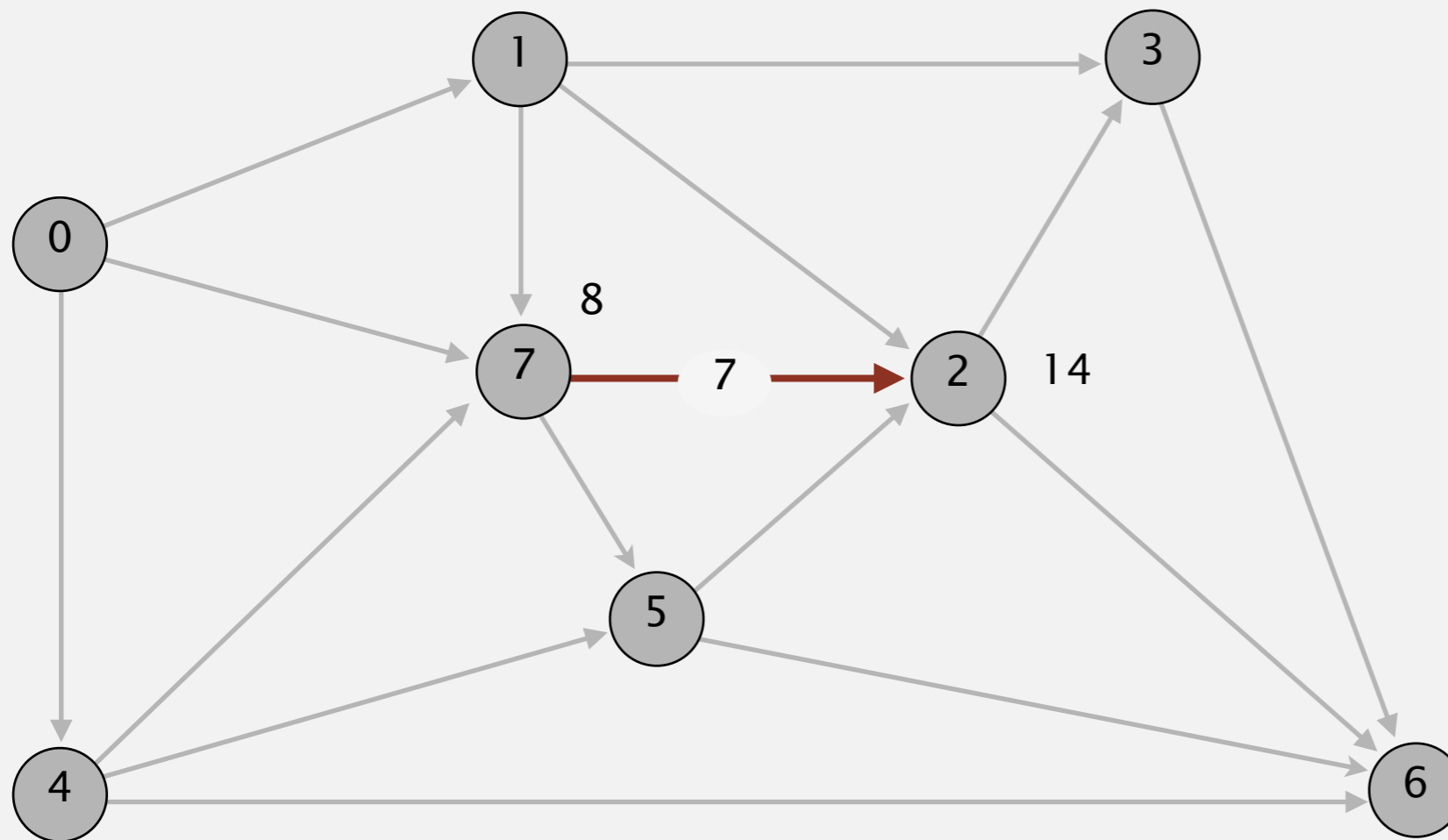
**pass 0**

0→1 0→4 0→7 1→2 1→3 1→7 2→3 2→6 3→6 4→5 4→6 4→7 5→2 5→6 7→5 7→2



# Bellman-Ford algorithm demo

Repeat  $V$  times: relax all  $E$  edges.



v	distTo[]	edgeTo[]
0	0.0	-
1	5.0	0→1
2	14.0	5→2
3	20.0	1→3
4	9.0	0→4
5	13.0	4→5
6	26.0	5→6
7	8.0	0→7

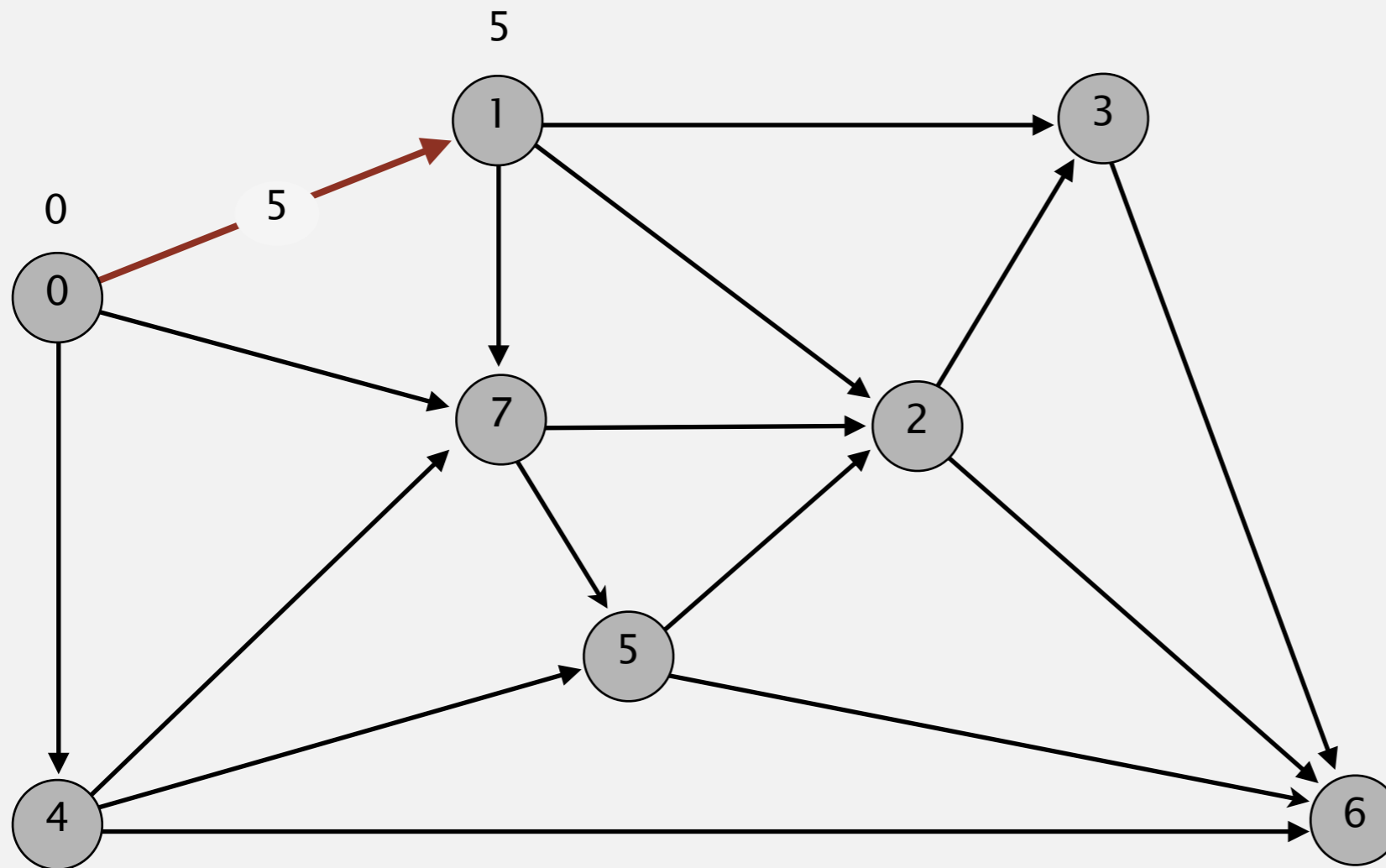
**pass 0**

0→1 0→4 0→7 1→2 1→3 1→7 2→3 2→6 3→6 4→5 4→6 4→7 5→2 5→6 7→5 7→2



# Bellman-Ford algorithm demo

Repeat  $V$  times: relax all  $E$  edges.



$v$	distTo[]	edgeTo[]
0	0.0	-
1	5.0	0→1
2	14.0	5→2
3	20.0	1→3
4	9.0	0→4
5	13.0	4→5
6	26.0	5→6
7	8.0	0→7

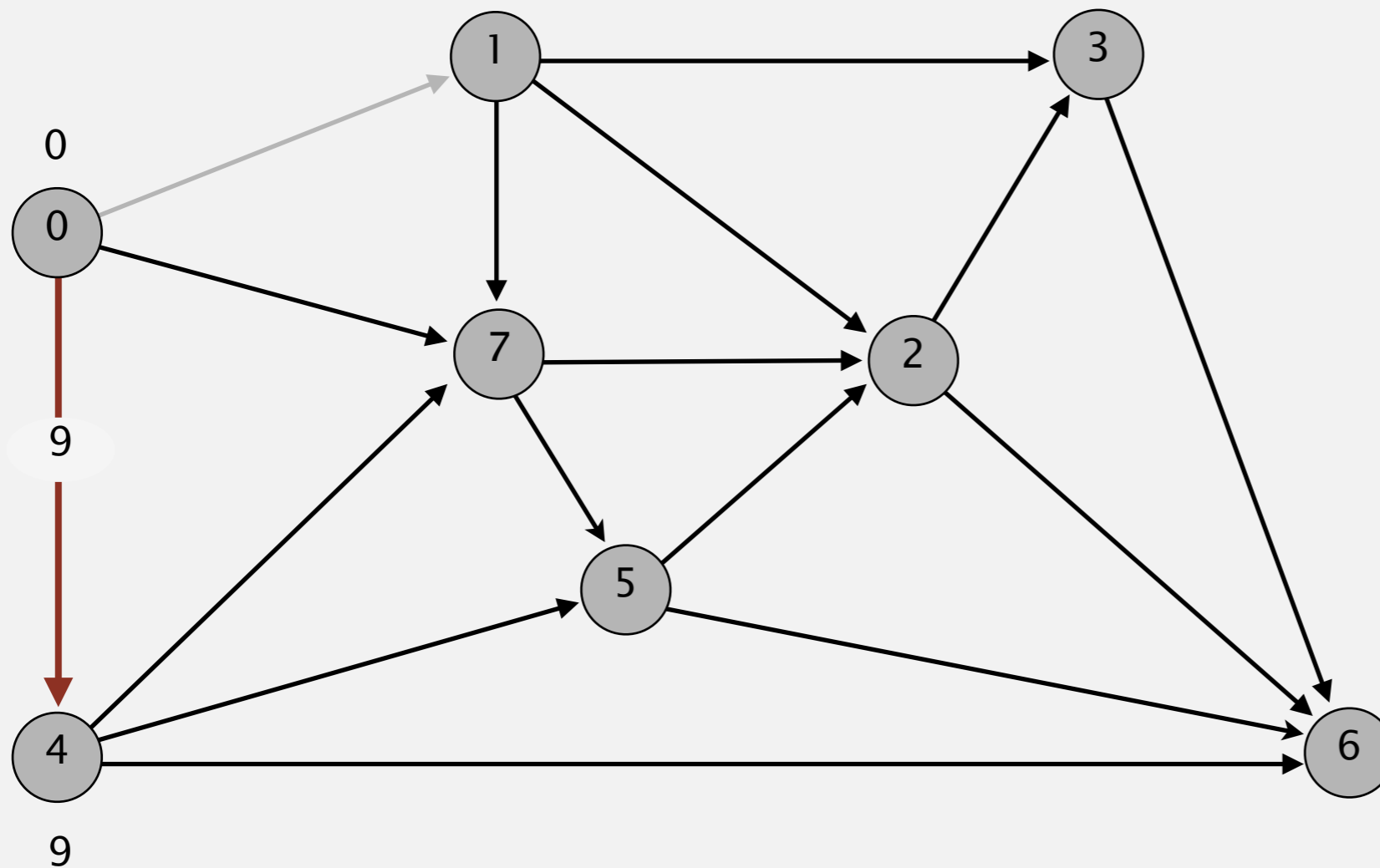
**pass 1**

0→1 0→4 0→7 1→2 1→3 1→7 2→3 2→6 3→6 4→5 4→6 4→7 5→2 5→6 7→5 7→2



# Bellman-Ford algorithm demo

Repeat  $V$  times: relax all  $E$  edges.



v	distTo[]	edgeTo[]
0	0.0	-
1	5.0	0→1
2	14.0	5→2
3	20.0	1→3
4	9.0	0→4
5	13.0	4→5
6	26.0	5→6
7	8.0	0→7

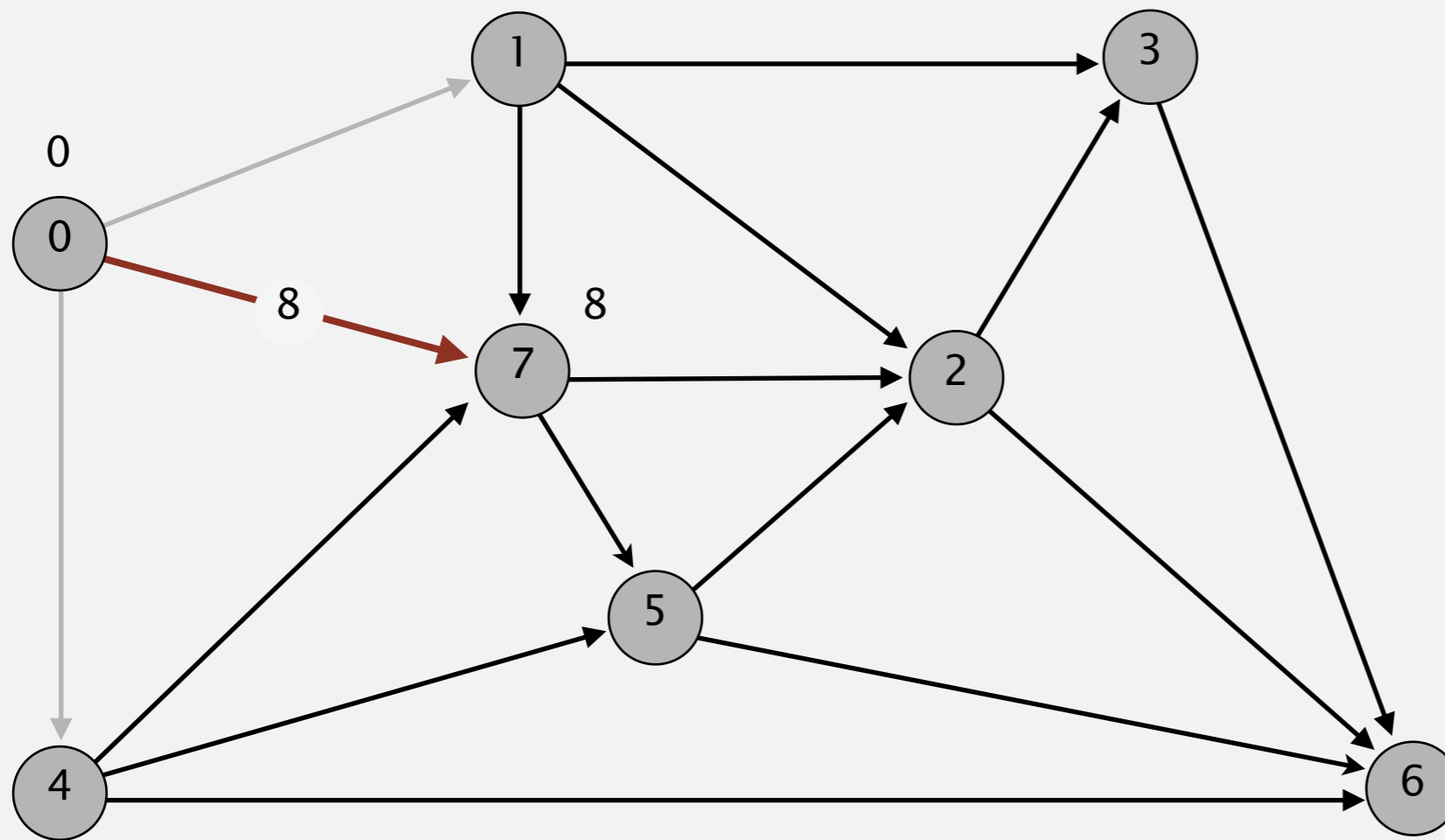
**pass 1**

0→1 0→4 0→7 1→2 1→3 1→7 2→3 2→6 3→6 4→5 4→6 4→7 5→2 5→6 7→5 7→2



# Bellman-Ford algorithm demo

Repeat  $V$  times: relax all  $E$  edges.



$v$	distTo[]	edgeTo[]
0	0.0	-
1	5.0	0→1
2	14.0	5→2
3	20.0	1→3
4	9.0	0→4
5	13.0	4→5
6	26.0	5→6
7	8.0	0→7

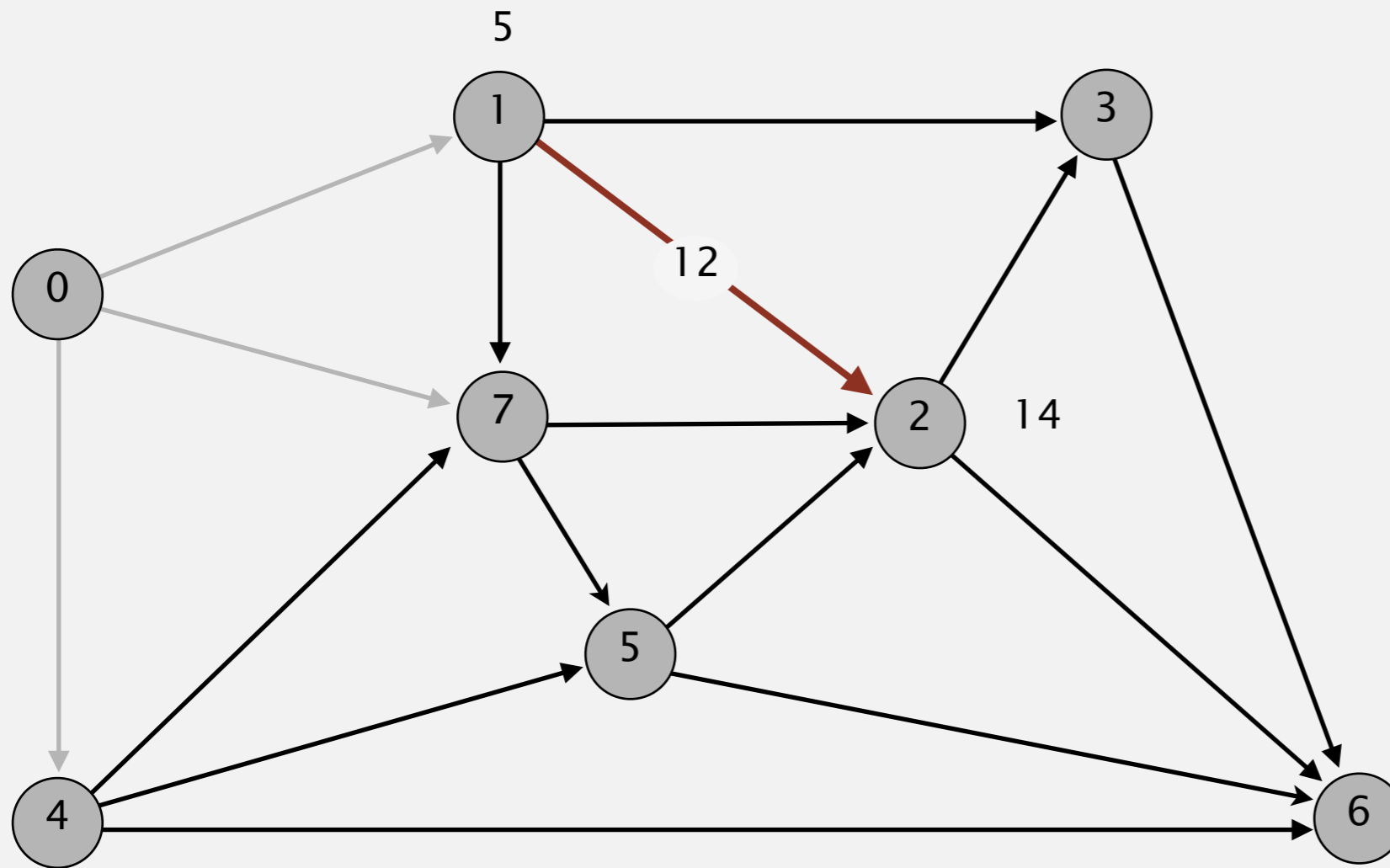
**pass 1**

0→1 0→4 0→7 1→2 1→3 1→7 2→3 2→6 3→6 4→5 4→6 4→7 5→2 5→6 7→5 7→2



# Bellman-Ford algorithm demo

Repeat  $V$  times: relax all  $E$  edges.



$v$	distTo[]	edgeTo[]
0	0.0	-
1	5.0	0→1
2	14.0	5→2
3	20.0	1→3
4	9.0	0→4
5	13.0	4→5
6	26.0	5→6
7	8.0	0→7

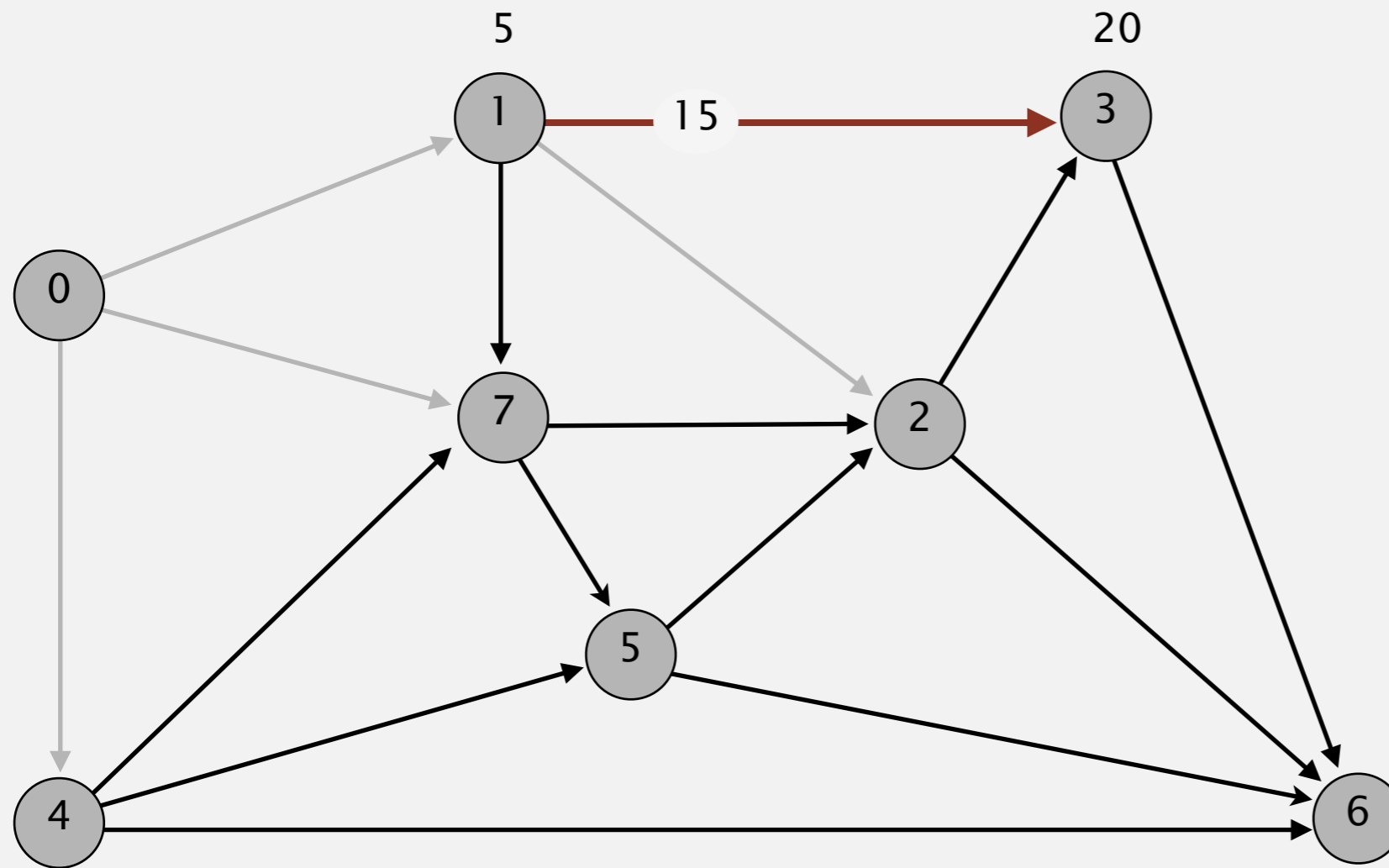
**pass 1**

0→1 0→4 0→7 1→2 1→3 1→7 2→3 2→6 3→6 4→5 4→6 4→7 5→2 5→6 7→5 7→2



# Bellman-Ford algorithm demo

Repeat  $V$  times: relax all  $E$  edges.



$v$	distTo[]	edgeTo[]
0	0.0	-
1	5.0	0→1
2	14.0	5→2
3	20.0	1→3
4	9.0	0→4
5	13.0	4→5
6	26.0	5→6
7	8.0	0→7

**pass 1**

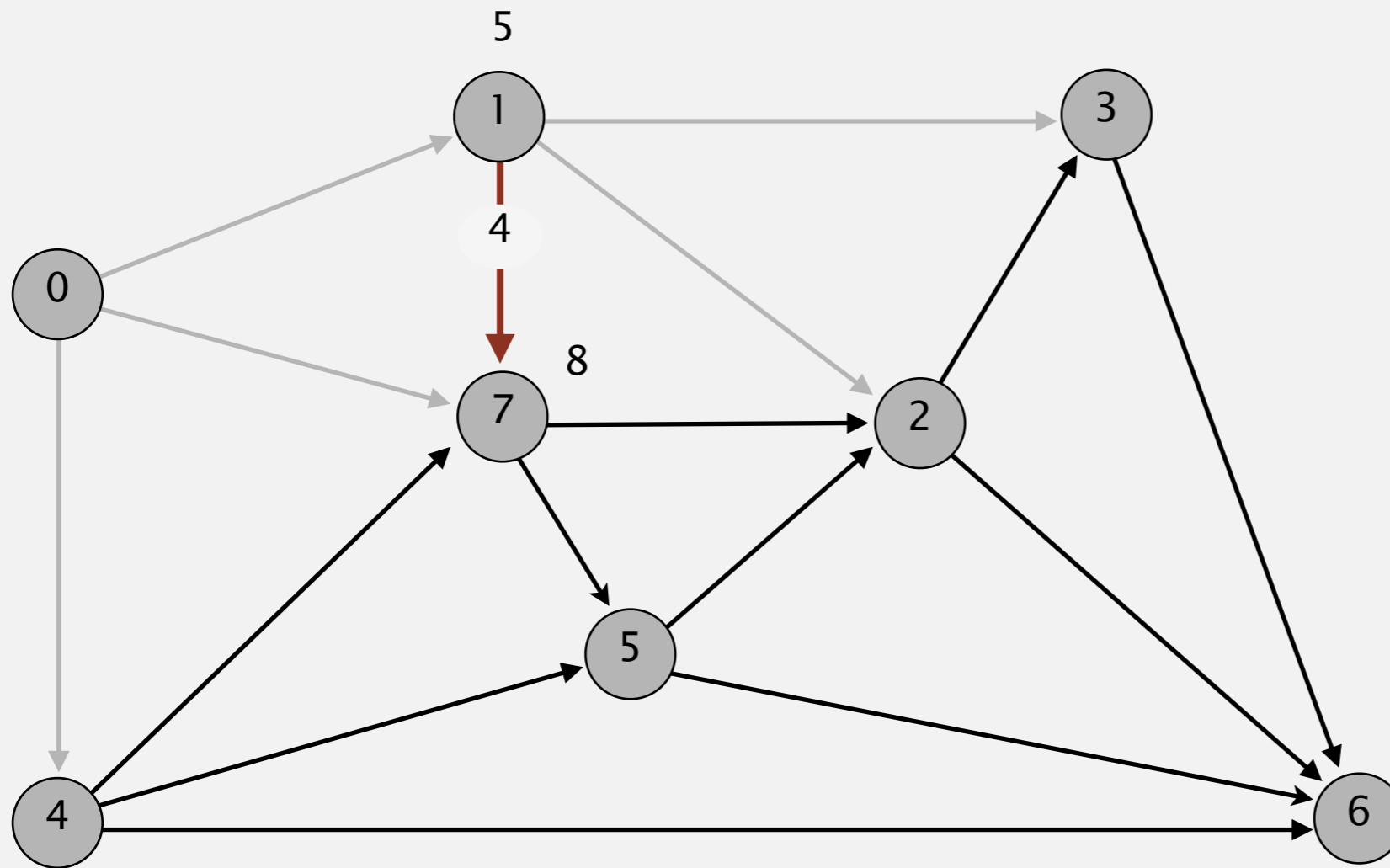
0→1 0→4 0→7 1→2 1→3 1→7 2→3 2→6 3→6 4→5 4→6 4→7 5→2 5→6 7→5 7→2





# Bellman-Ford algorithm demo

Repeat  $V$  times: relax all  $E$  edges.



$v$	distTo[]	edgeTo[]
0	0.0	-
1	5.0	0→1
2	14.0	5→2
3	20.0	1→3
4	9.0	0→4
5	13.0	4→5
6	26.0	5→6
7	8.0	0→7

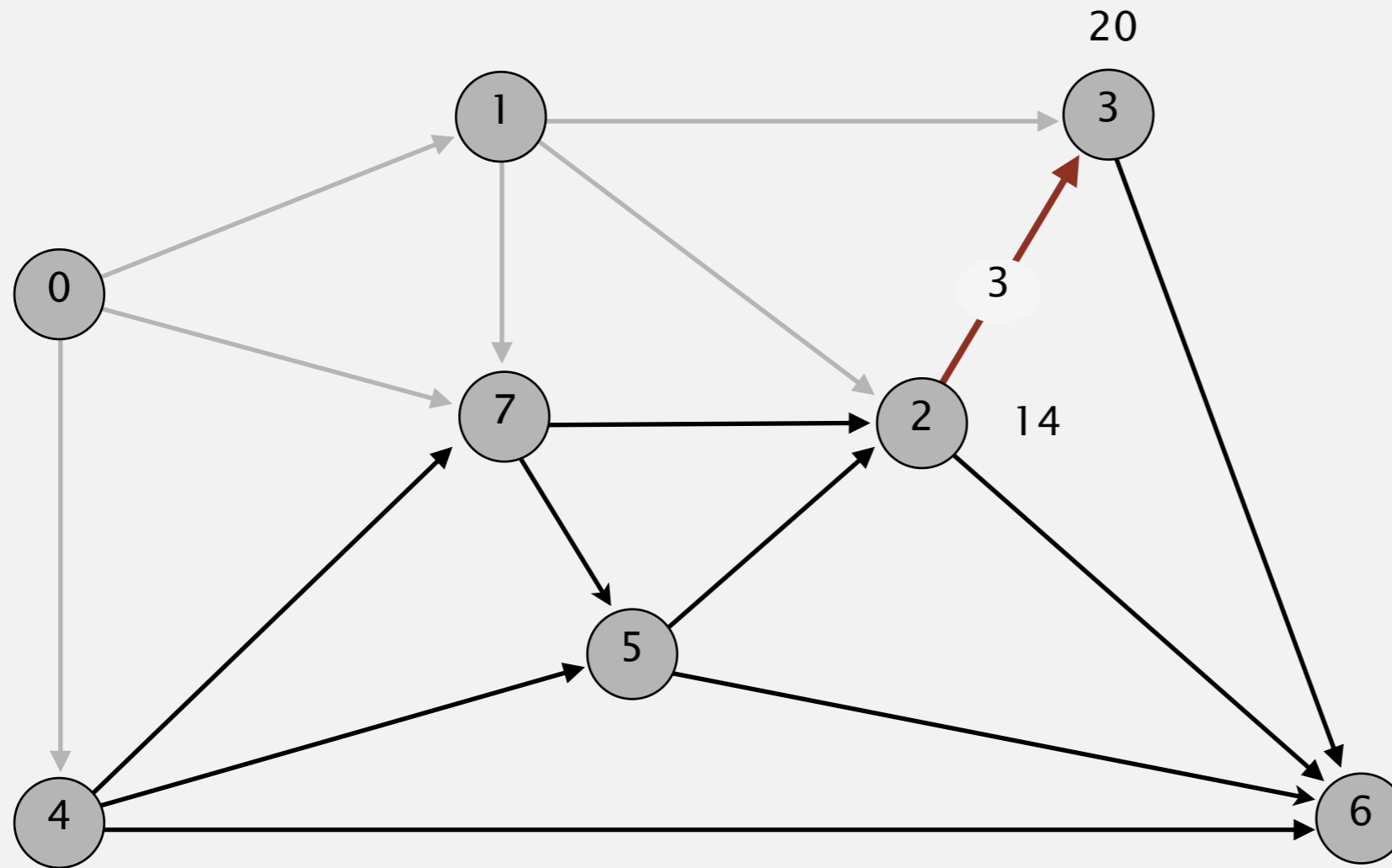
**pass 1**

0→1 0→4 0→7 1→2 1→3 1→7 2→3 2→6 3→6 4→5 4→6 4→7 5→2 5→6 7→5 7→2



# Bellman-Ford algorithm demo

Repeat  $V$  times: relax all  $E$  edges.



$v$	distTo[]	edgeTo[]
0	0.0	-
1	5.0	0→1
2	14.0	5→2
3	20.0	1→3
4	9.0	0→4
5	13.0	4→5
6	26.0	5→6
7	8.0	0→7

**pass 1**

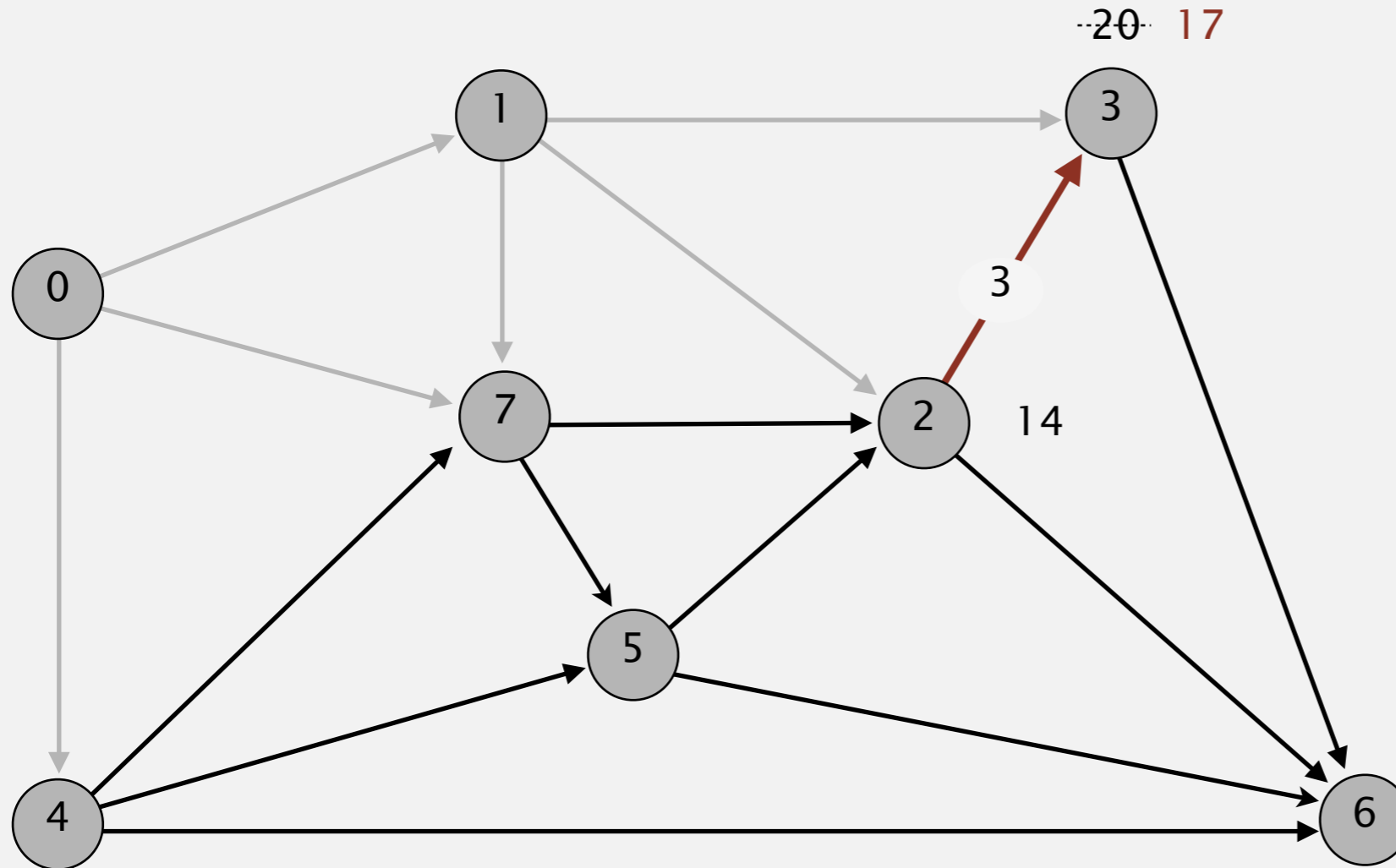
0→1 0→4 0→7 1→2 1→3 1→7 2→3 2→6 3→6 4→5 4→6 4→7 5→2 5→6 7→5 7→2



# Bellman-Ford algorithm demo

Repeat  $V$  times: relax all  $E$  edges.

2-3 successfully relaxed in pass 1, but not pass 0



v	distTo[]	edgeTo[]
0	0.0	-
1	5.0	0→1
2	14.0	5→2
3	17.0	2→3
4	9.0	0→4
5	13.0	4→5
6	26.0	5→6
7	8.0	0→7

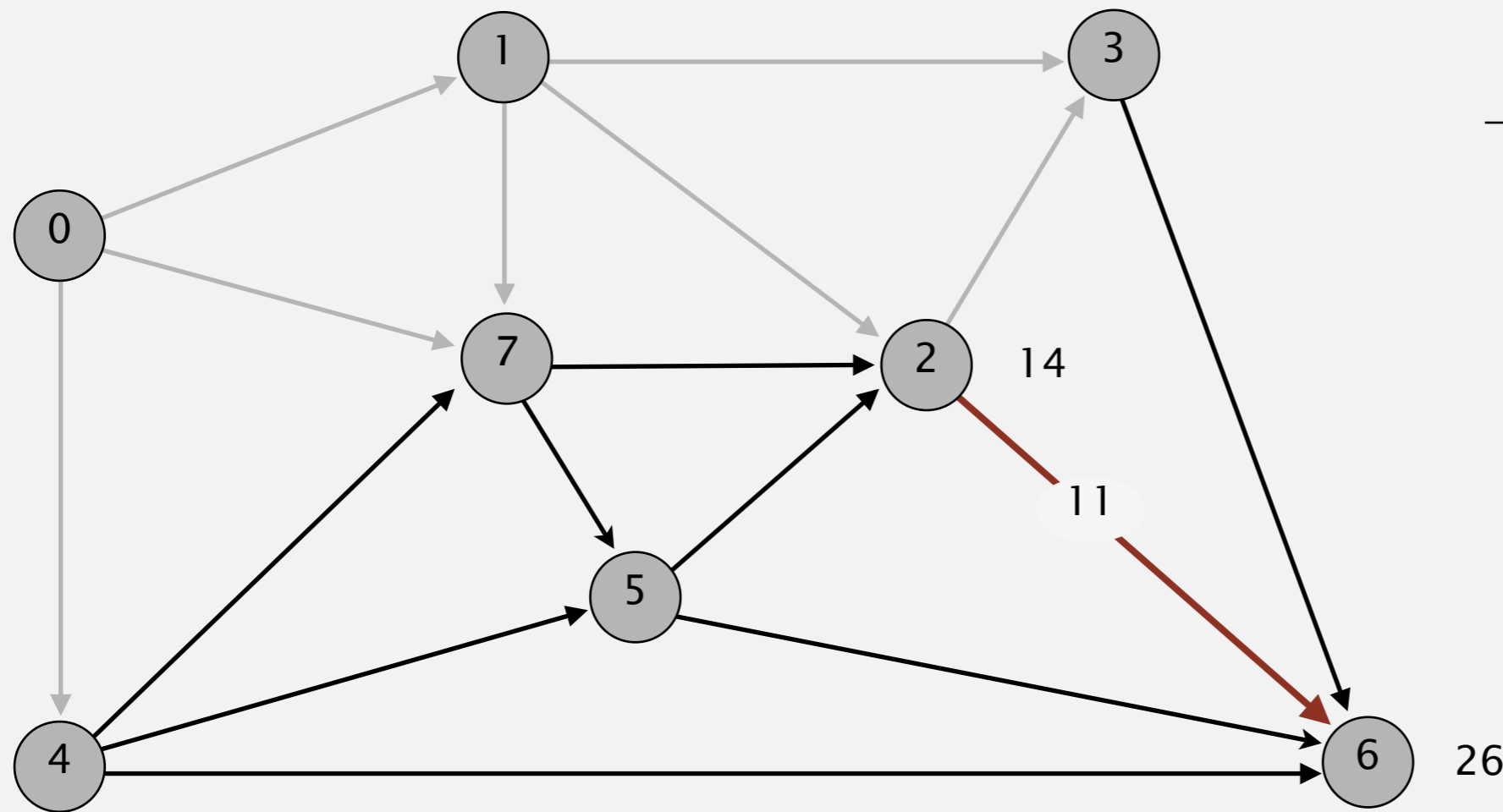
pass 1

0→1 0→4 0→7 1→2 1→3 1→7 2→3 2→6 3→6 4→5 4→6 4→7 5→2 5→6 7→5 7→2



# Bellman-Ford algorithm demo

Repeat  $V$  times: relax all  $E$  edges.



v	distTo[]	edgeTo[]
0	0.0	-
1	5.0	0→1
2	14.0	5→2
3	17.0	2→3
4	9.0	0→4
5	13.0	4→5
6	26.0	5→6
7	8.0	0→7

**pass 1**

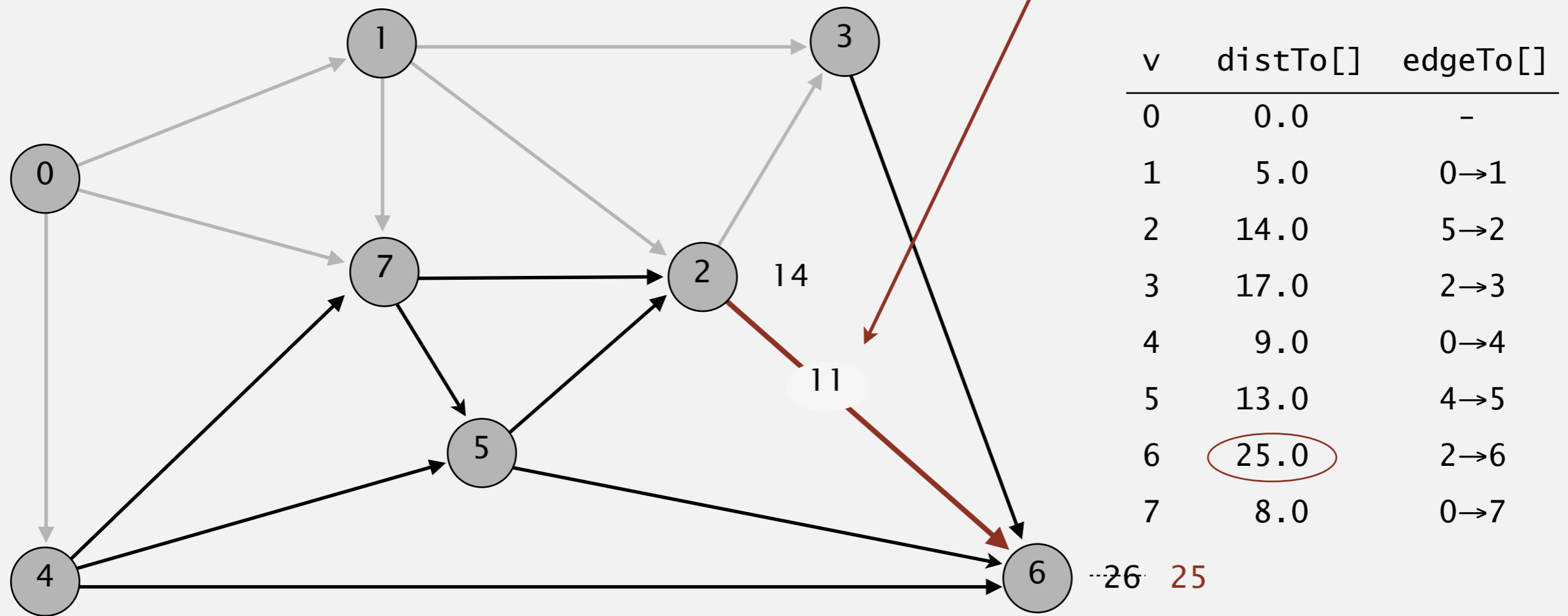
0→1 0→4 0→7 1→2 1→3 1→7 2→3 2→6 3→6 4→5 4→6 4→7 5→2 5→6 7→5 7→2



# Bellman-Ford algorithm demo

Repeat  $V$  times: relax all  $E$  edges.

2-6 successfully relaxed  
in pass 0 and pass 1



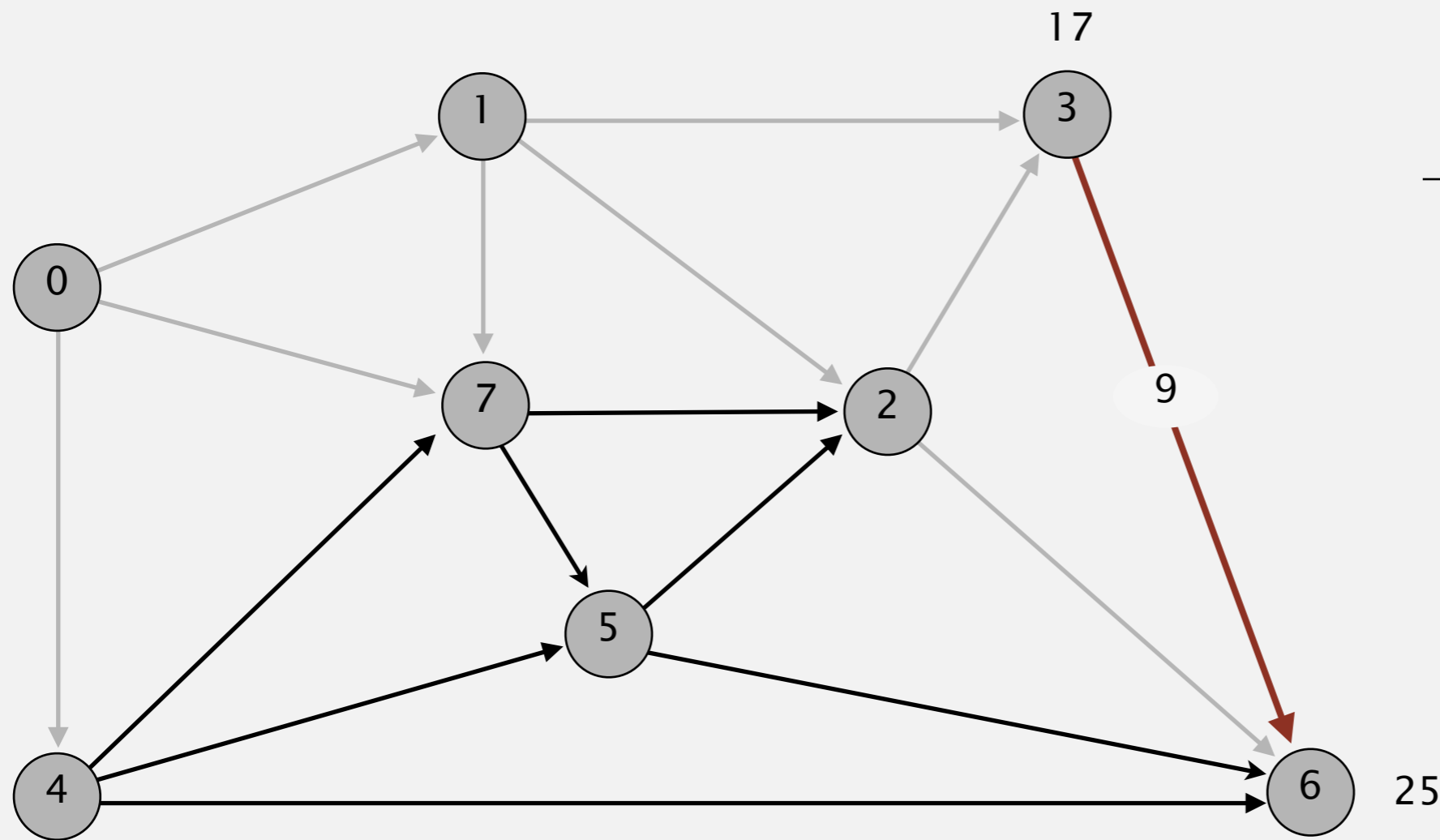
pass 1

0→1 0→4 0→7 1→2 1→3 1→7 2→3 2→6 3→6 4→5 4→6 4→7 5→2 5→6 7→5 7→2



# Bellman-Ford algorithm demo

Repeat  $V$  times: relax all  $E$  edges.



v	distTo[]	edgeTo[]
0	0.0	-
1	5.0	0→1
2	14.0	5→2
3	17.0	2→3
4	9.0	0→4
5	13.0	4→5
6	25.0	2→6
7	8.0	0→7

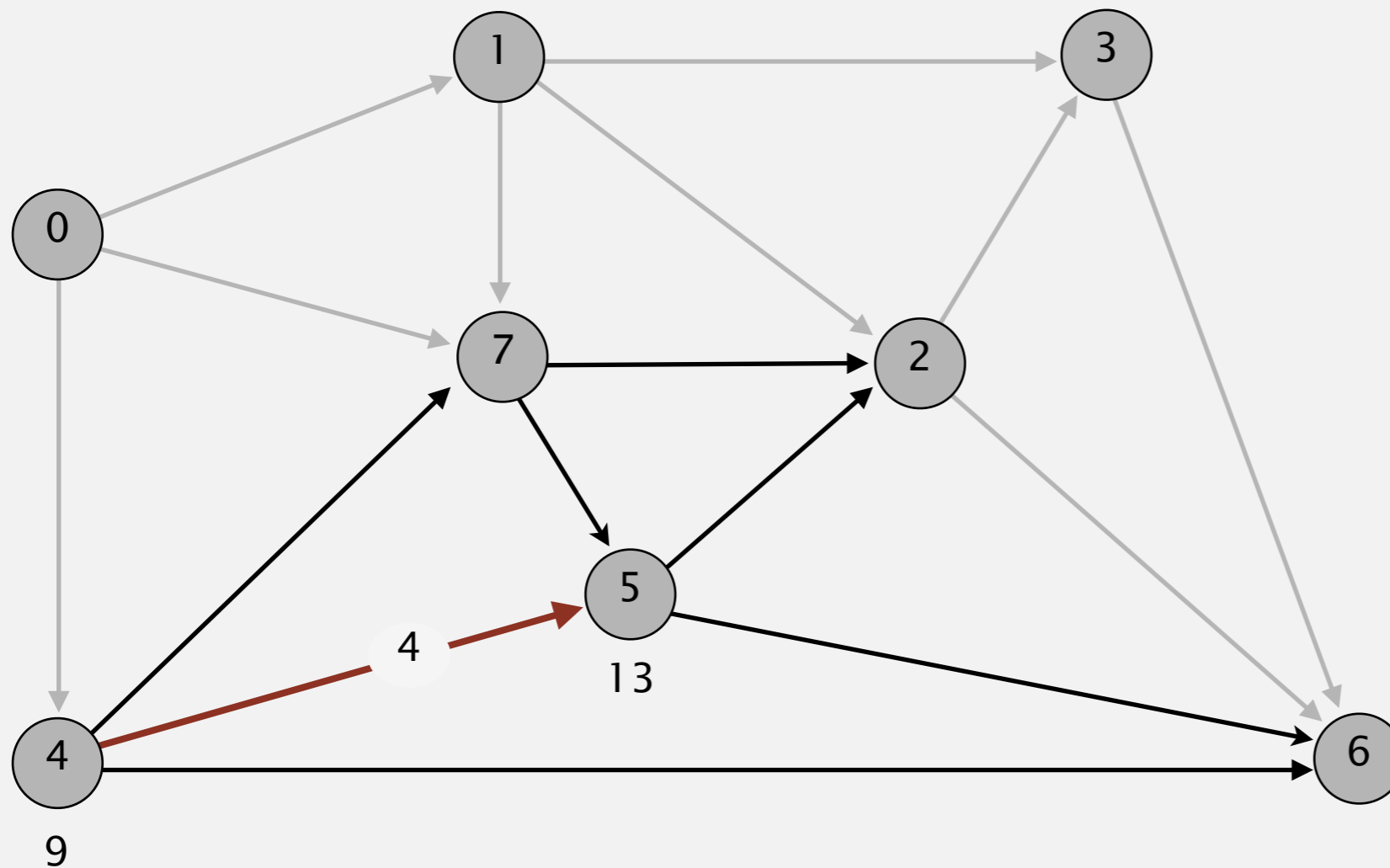
**pass 1**

0→1 0→4 0→7 1→2 1→3 1→7 2→3 2→6 3→6 4→5 4→6 4→7 5→2 5→6 7→5 7→2



# Bellman-Ford algorithm demo

Repeat  $V$  times: relax all  $E$  edges.



v	distTo[]	edgeTo[]
0	0.0	-
1	5.0	0→1
2	14.0	5→2
3	17.0	2→3
4	9.0	0→4
5	13.0	4→5
6	25.0	2→6
7	8.0	0→7

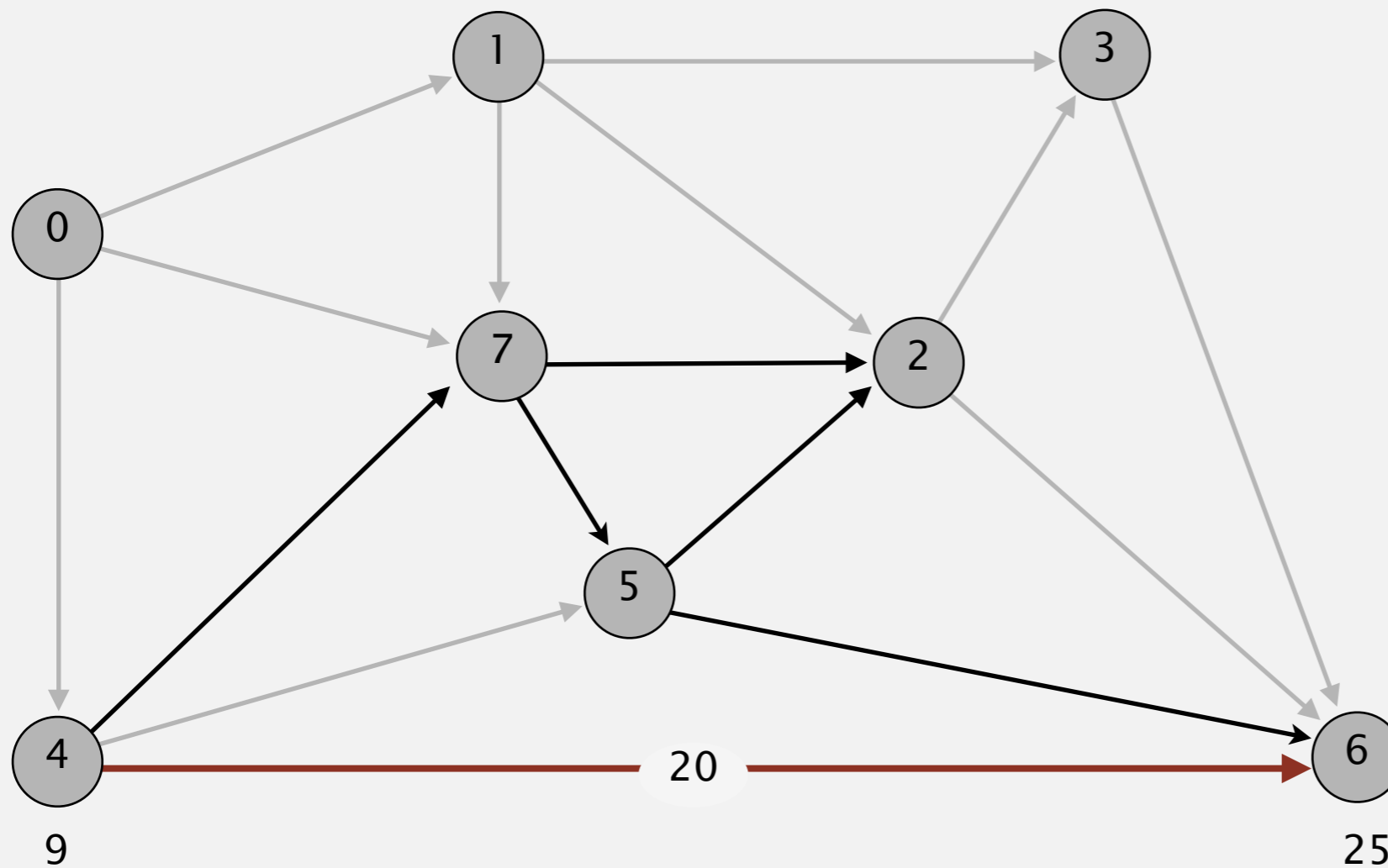
**pass 1**

0→1 0→4 0→7 1→2 1→3 1→7 2→3 2→6 3→6 4→5 4→6 4→7 5→2 5→6 7→5 7→2



# Bellman-Ford algorithm demo

Repeat  $V$  times: relax all  $E$  edges.



v	distTo[]	edgeTo[]
0	0.0	-
1	5.0	0→1
2	14.0	5→2
3	17.0	2→3
4	9.0	0→4
5	13.0	4→5
6	25.0	2→6
7	8.0	0→7

**pass 1**

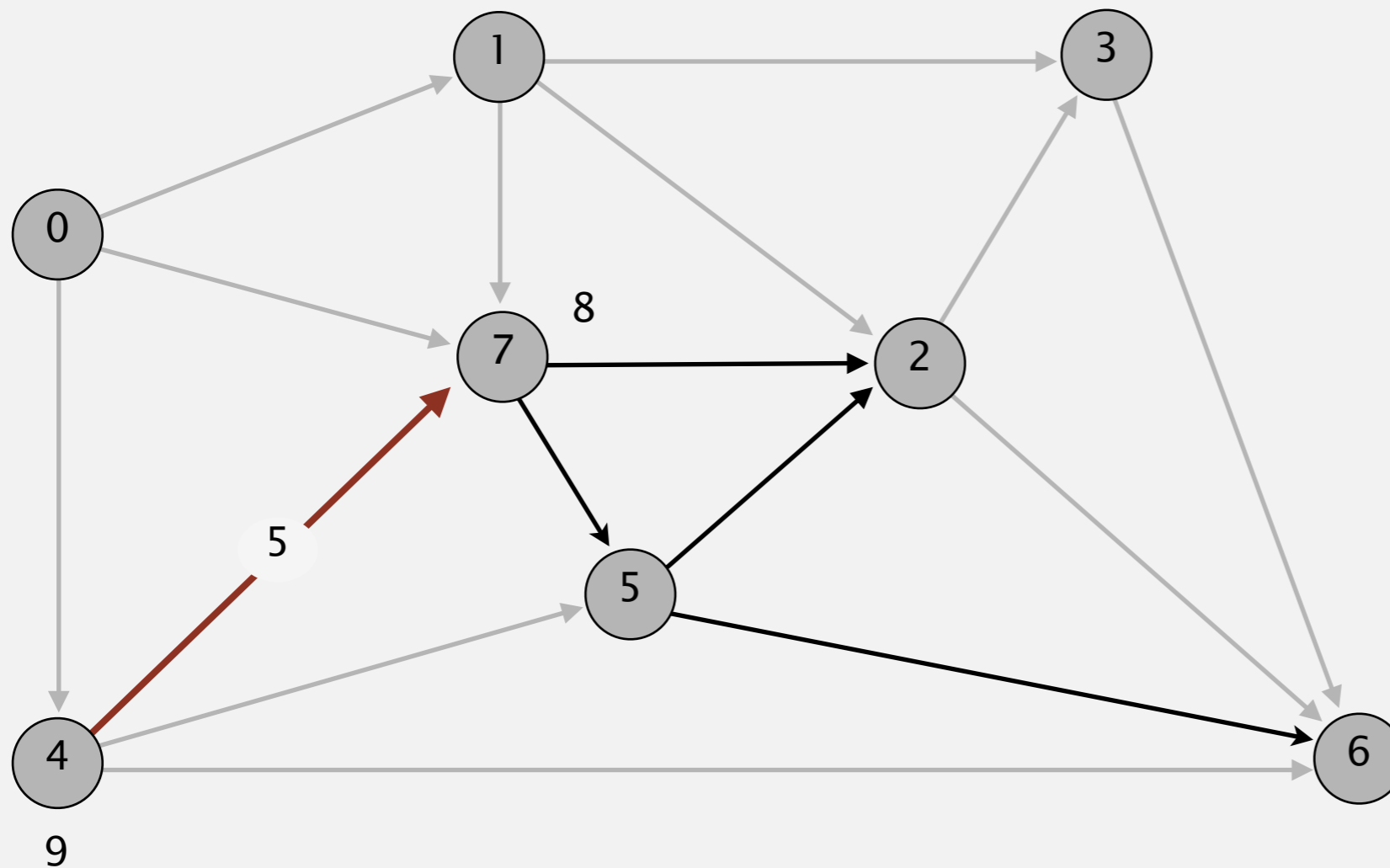
0→1 0→4 0→7 1→2 1→3 1→7 2→3 2→6 3→6 4→5 4→6 4→7 5→2 5→6 7→5 7→2





# Bellman-Ford algorithm demo

Repeat  $V$  times: relax all  $E$  edges.



v	distTo[]	edgeTo[]
0	0.0	-
1	5.0	0→1
2	14.0	5→2
3	17.0	2→3
4	9.0	0→4
5	13.0	4→5
6	25.0	2→6
7	8.0	0→7

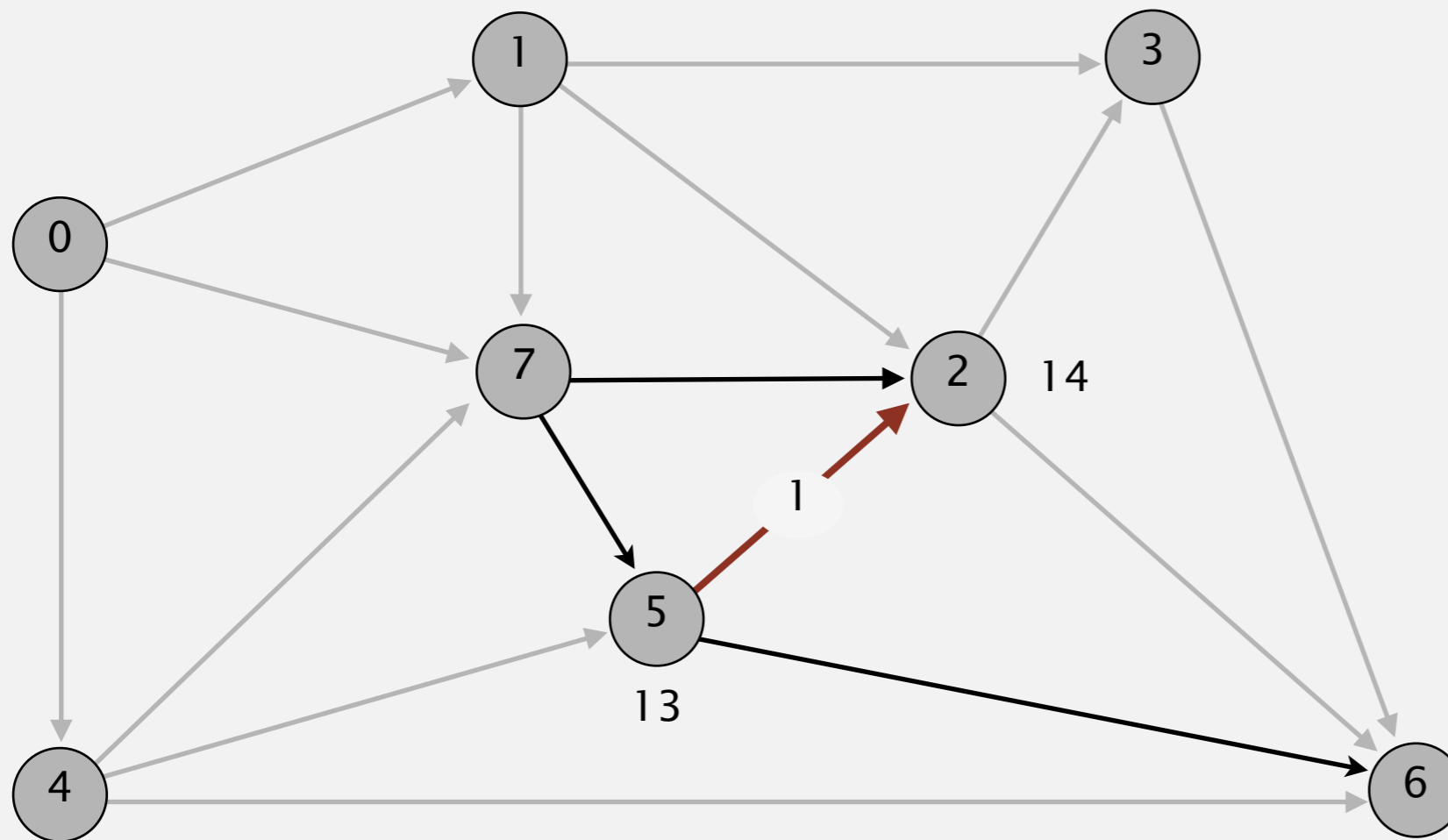
**pass 1**

0→1 0→4 0→7 1→2 1→3 1→7 2→3 2→6 3→6 4→5 4→6 4→7 5→2 5→6 7→5 7→2



# Bellman-Ford algorithm demo

Repeat  $V$  times: relax all  $E$  edges.



v	distTo[]	edgeTo[]
0	0.0	-
1	5.0	0→1
2	14.0	5→2
3	17.0	2→3
4	9.0	0→4
5	13.0	4→5
6	25.0	2→6
7	8.0	0→7

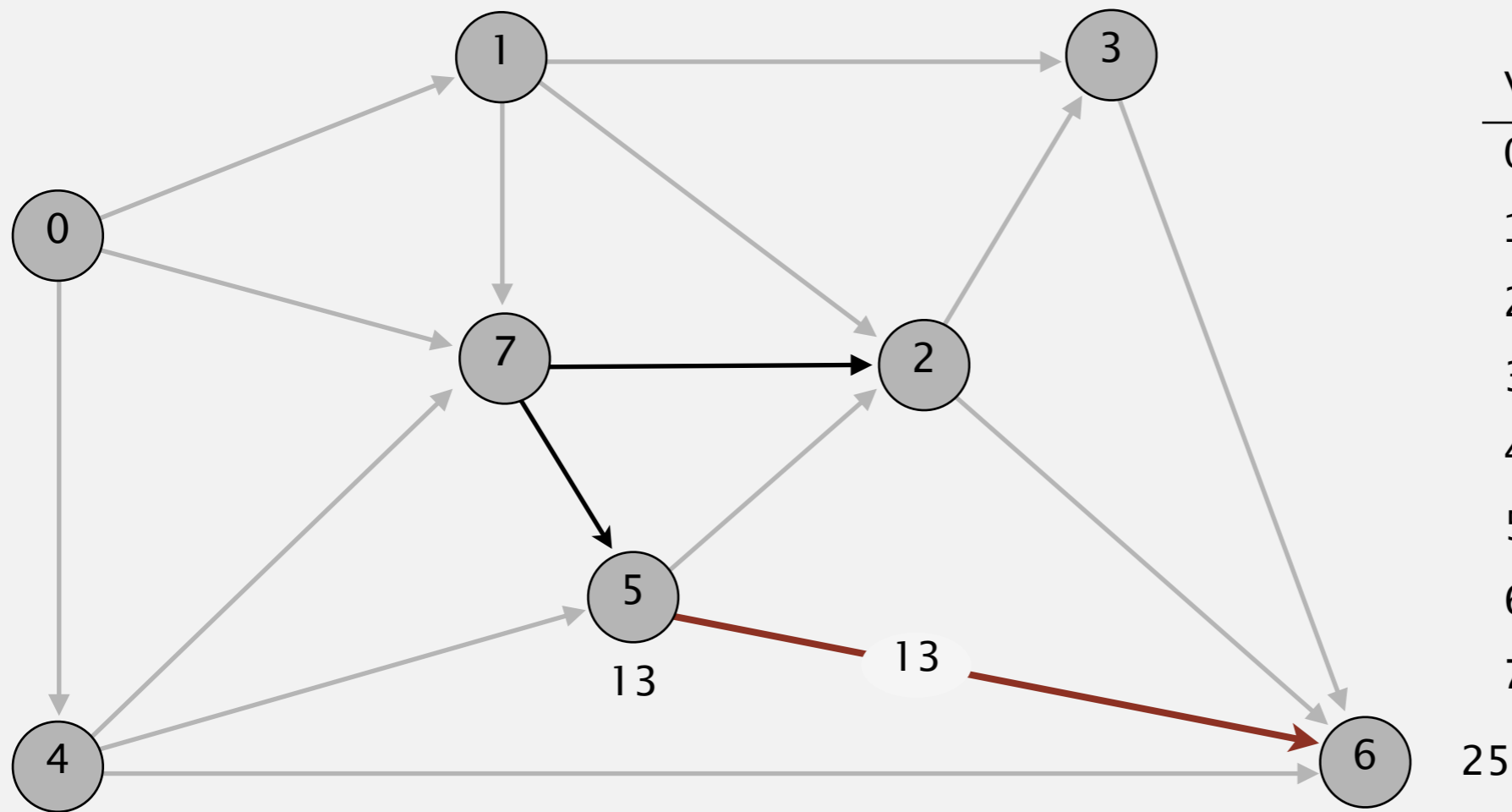
**pass 1**

0→1 0→4 0→7 1→2 1→3 1→7 2→3 2→6 3→6 4→5 4→6 4→7 5→2 5→6 7→5 7→2



# Bellman-Ford algorithm demo

Repeat  $V$  times: relax all  $E$  edges.



v	distTo[]	edgeTo[]
0	0.0	-
1	5.0	0→1
2	14.0	5→2
3	17.0	2→3
4	9.0	0→4
5	13.0	4→5
6	25.0	2→6
7	8.0	0→7

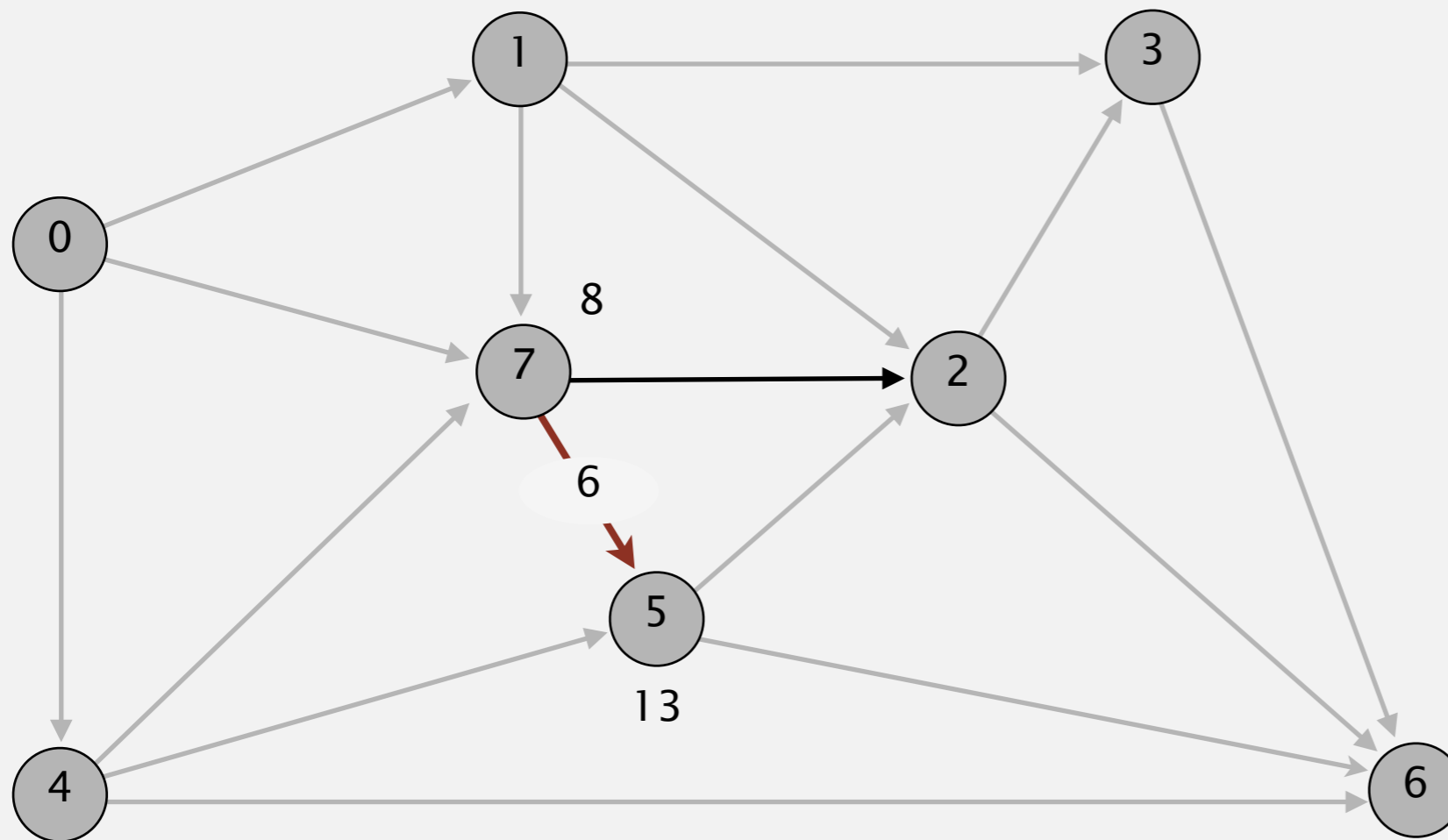
**pass 1**

0→1 0→4 0→7 1→2 1→3 1→7 2→3 2→6 3→6 4→5 4→6 4→7 5→2 5→6 7→5 7→2



# Bellman-Ford algorithm demo

Repeat  $V$  times: relax all  $E$  edges.



v	distTo[]	edgeTo[]
0	0.0	-
1	5.0	0→1
2	14.0	5→2
3	17.0	2→3
4	9.0	0→4
5	13.0	4→5
6	25.0	2→6
7	8.0	0→7

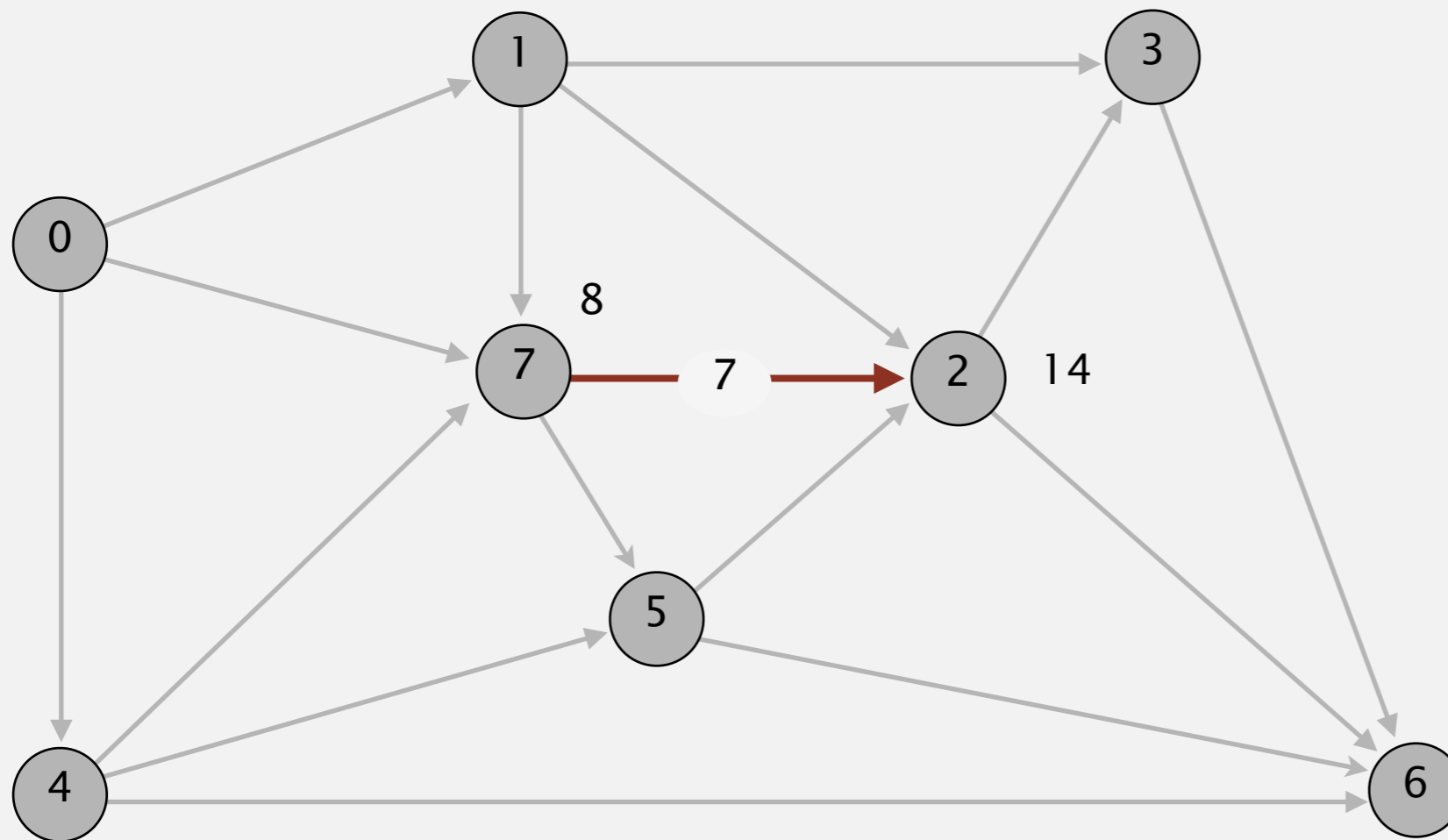
**pass 1**

0→1 0→4 0→7 1→2 1→3 1→7 2→3 2→6 3→6 4→5 4→6 4→7 5→2 5→6 7→5 7→2



# Bellman-Ford algorithm demo

Repeat  $V$  times: relax all  $E$  edges.



v	distTo[]	edgeTo[]
0	0.0	-
1	5.0	0→1
2	14.0	5→2
3	17.0	2→3
4	9.0	0→4
5	13.0	4→5
6	25.0	2→6
7	8.0	0→7

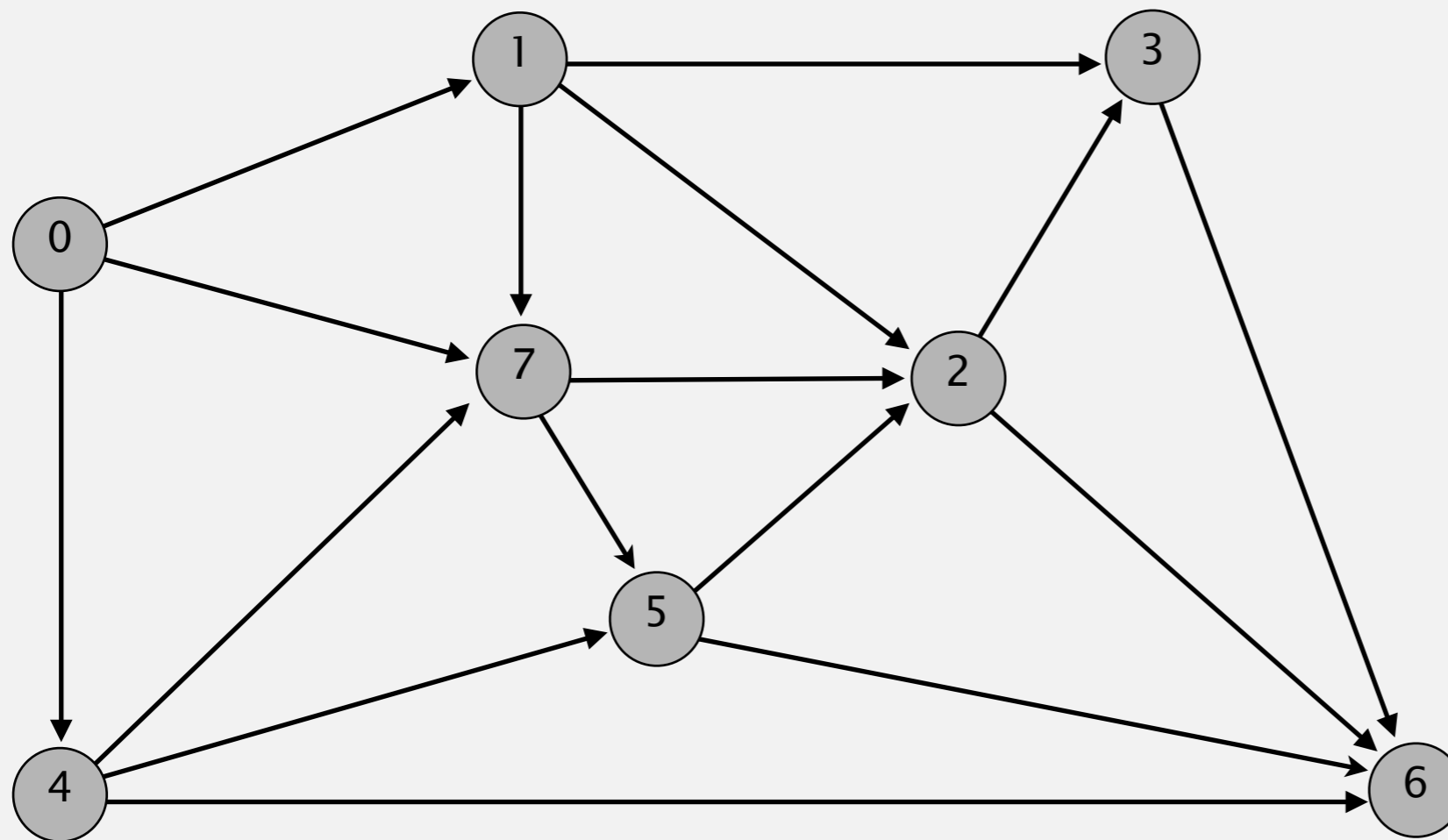
**pass 1**

0→1 0→4 0→7 1→2 1→3 1→7 2→3 2→6 3→6 4→5 4→6 4→7 5→2 5→6 7→5 7→2



# Bellman-Ford algorithm demo

Repeat  $V$  times: relax all  $E$  edges.



$v$	distTo[]	edgeTo[]
0	0.0	-
1	5.0	0→1
2	14.0	5→2
3	17.0	2→3
4	9.0	0→4
5	13.0	4→5
6	25.0	2→6
7	8.0	0→7

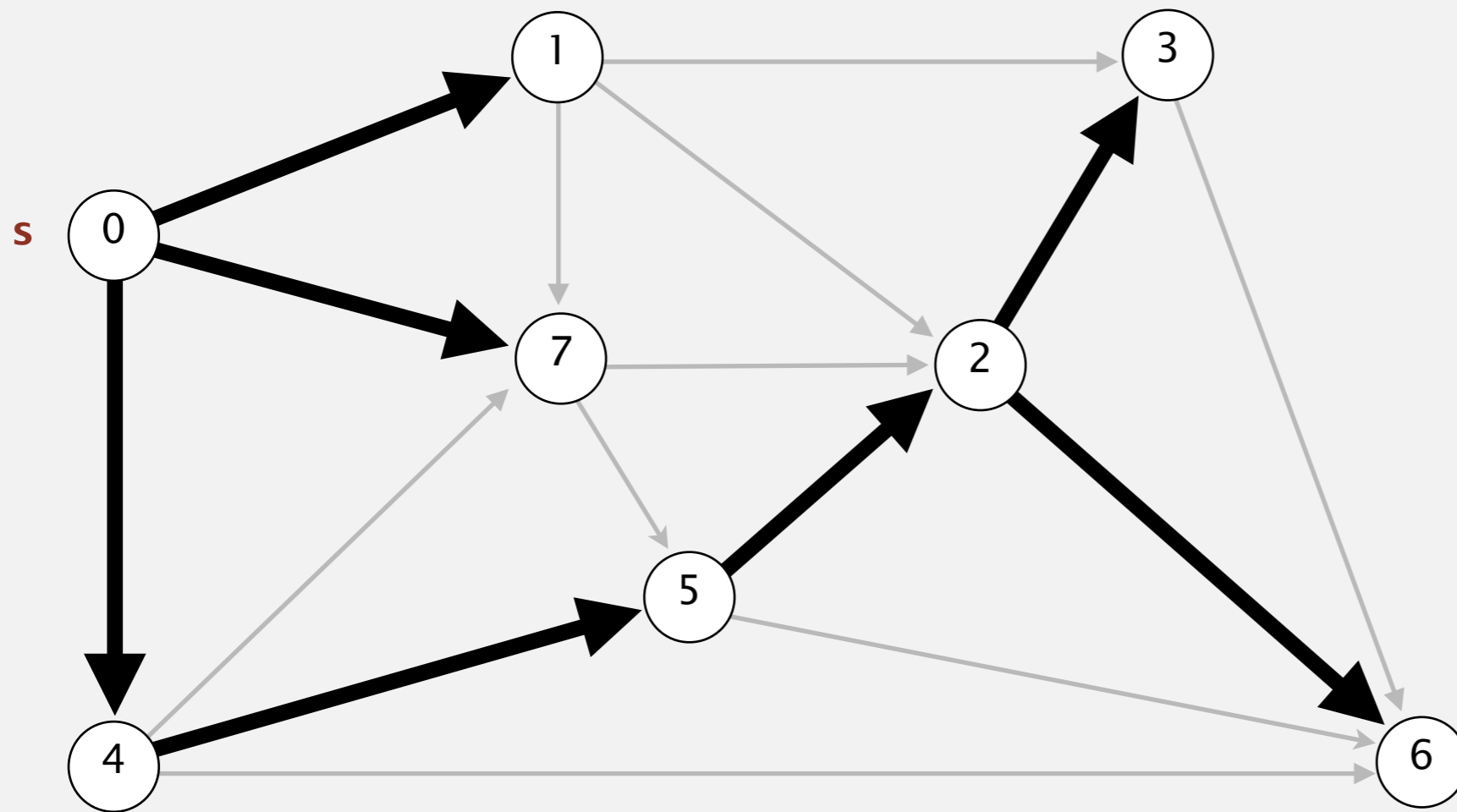
pass 2, 3, 4, 5, 6, 7 (no further changes)

0→1 0→4 0→7 1→2 1→3 1→7 2→3 2→6 3→6 4→5 4→6 4→7 5→2 5→6 7→5 7→2



# Bellman-Ford algorithm demo

Repeat  $V$  times: relax all  $E$  edges.



$v$	distTo[]	edgeTo[]
0	0.0	-
1	5.0	0→1
2	14.0	5→2
3	17.0	2→3
4	9.0	0→4
5	13.0	4→5
6	25.0	2→6
7	8.0	0→7

**shortest-paths tree from vertex s**

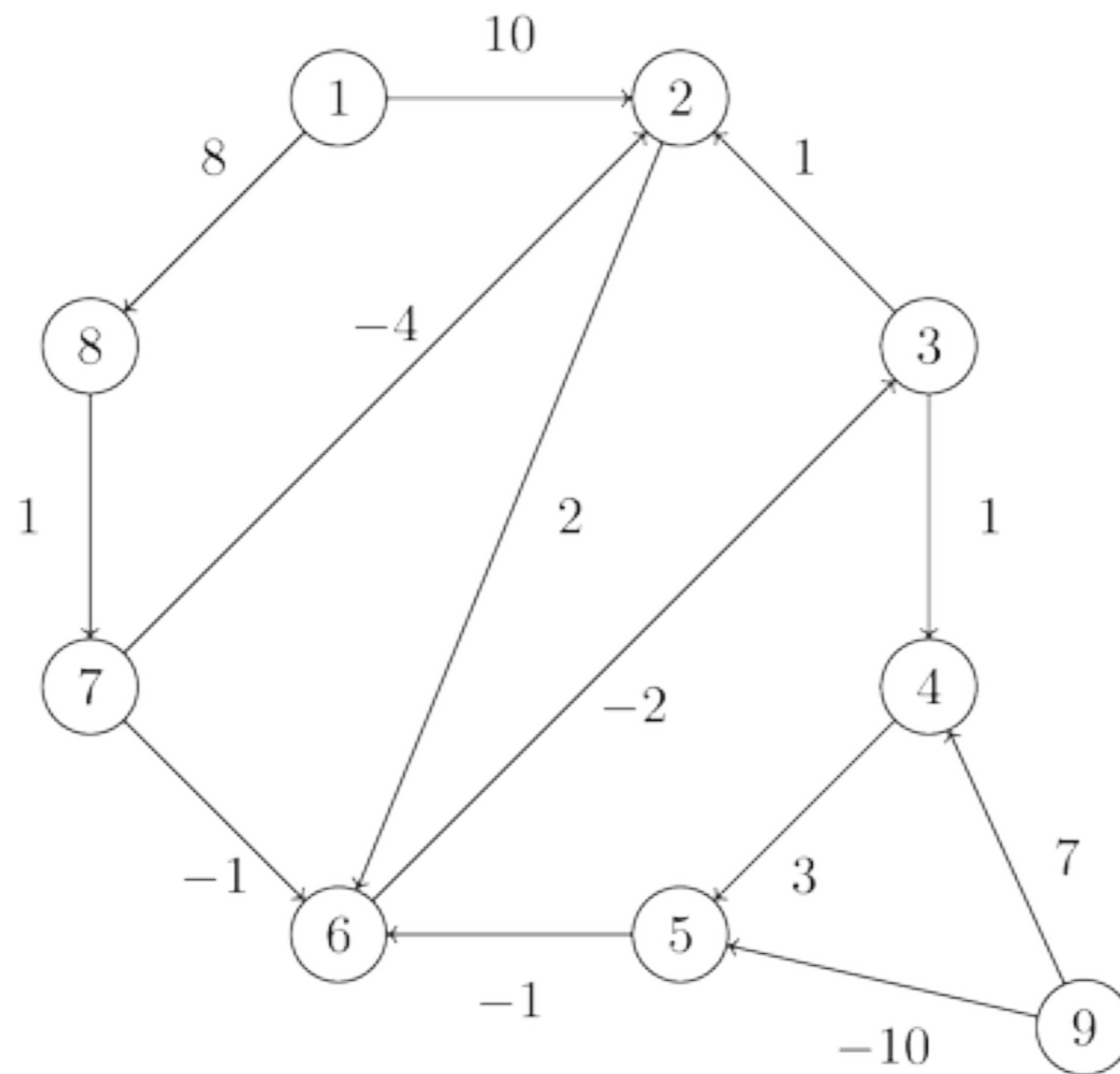
## Bellman-Ford vs Dijkstra

- ▶ Bellman-Ford's worst-case running time is  $|E| |V|$  vs Dijkstra's  $|E| \log |V|$ .
  - ▶ Bellman-Ford's algorithm is queue-based.
- ▶ Both require  $|V|$  extra space and can handle graphs with cycles.
- ▶ Only Bellman-Ford can handle negative weights, as long as there are no cycles that sum to a negative weight.



# Practice Time

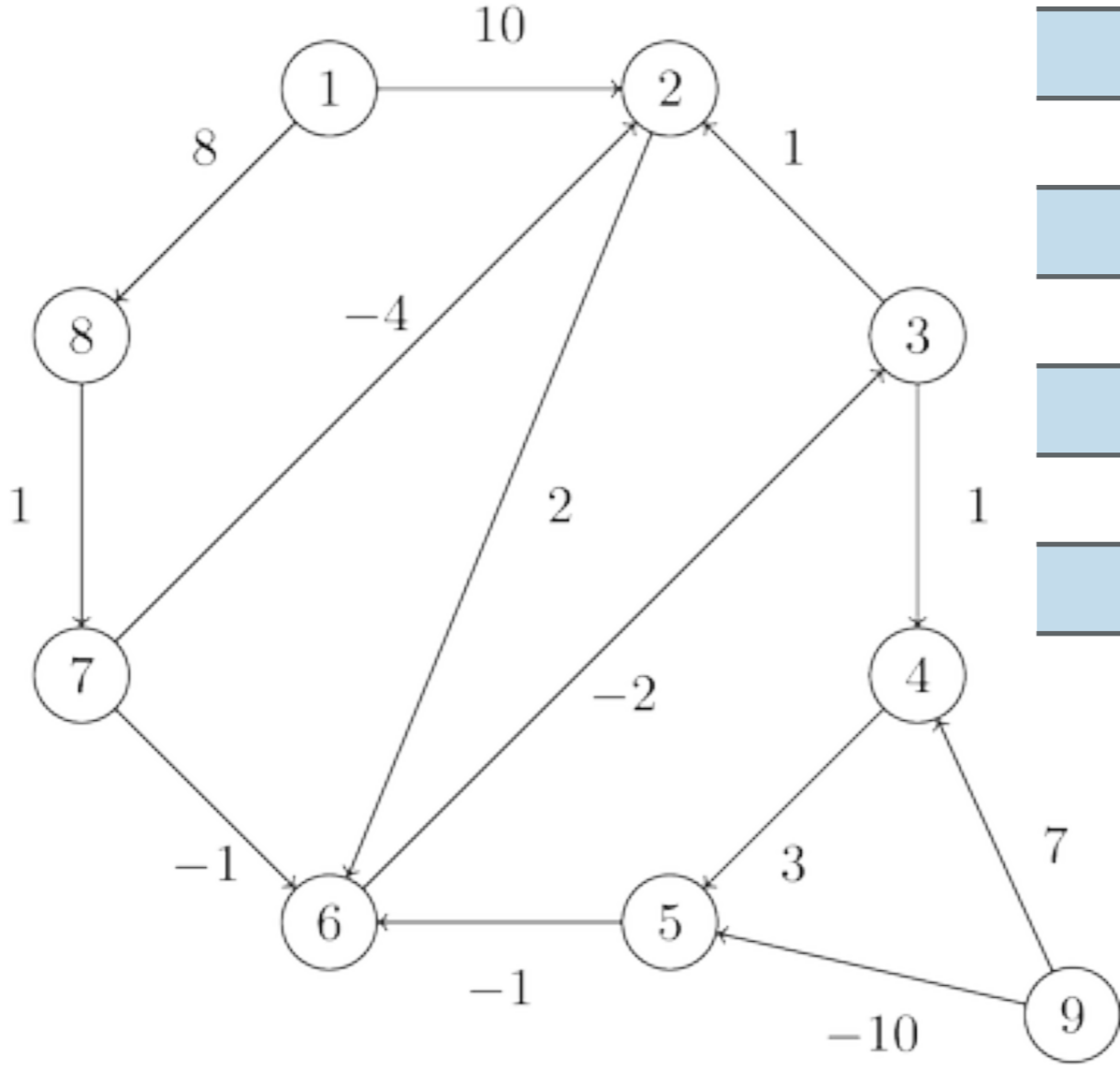
1 2 10  
 3 2 1  
 3 4 1  
 4 5 3  
 5 6 -1  
 7 6 -1  
 8 7 1  
 1 8 8  
 7 2 -4  
 2 6 2  
 6 3 -2  
 9 5 -10  
 9 4 7



<http://rosalind.info/problems/bf/>

# Practice Time

1 2 10  
 3 2 1  
 3 4 1  
 4 5 3  
 5 6 -1  
 7 6 -1  
 8 7 1  
 1 8 8  
 7 2 -4  
 2 6 2  
 6 3 -2  
 9 5 -10  
 9 4 7

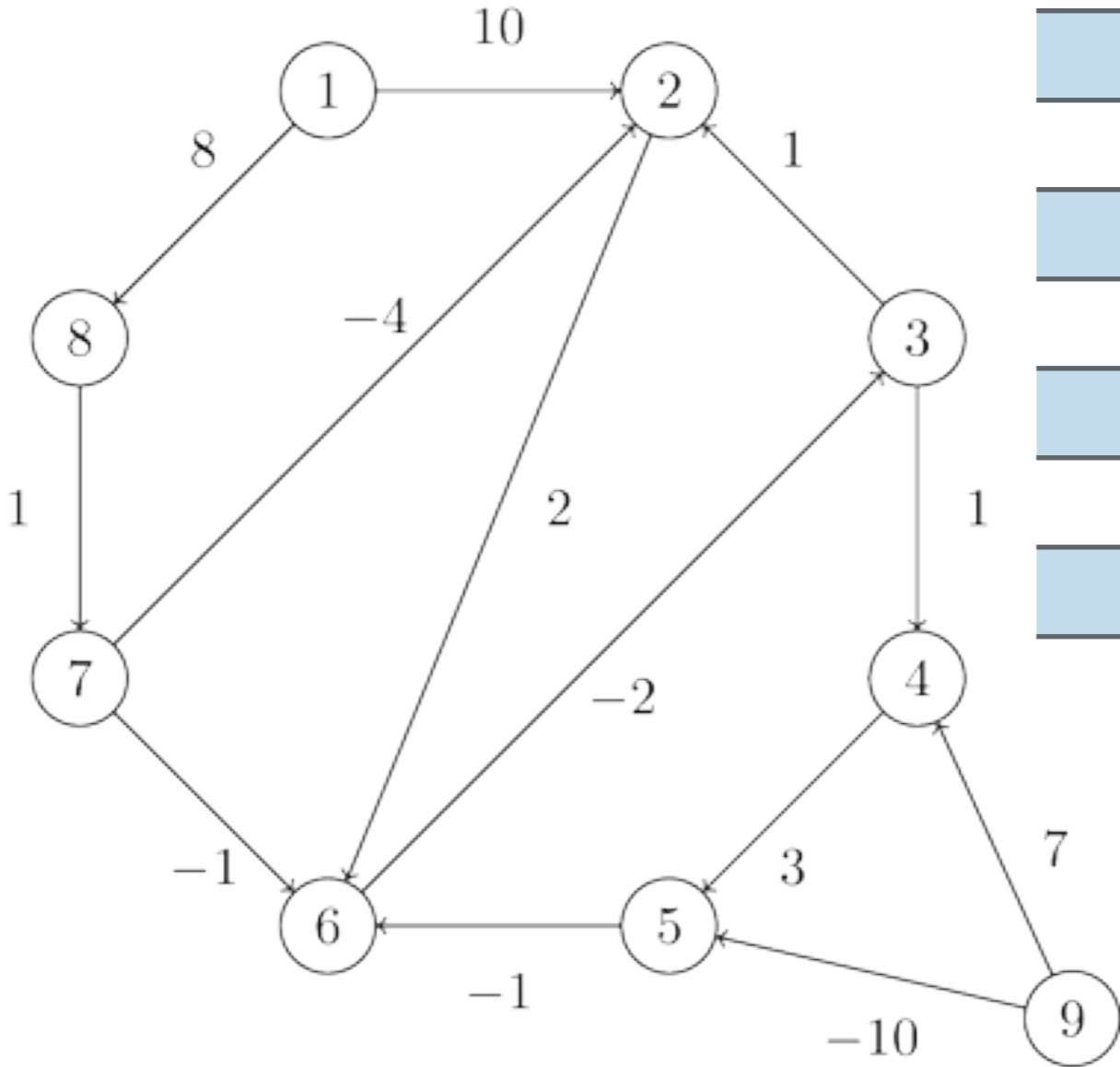


v	distTo[]	edgeTo[]
1	0	-
2	Inf	null
3	Inf	null
4	Inf	null
5	Inf	null
6	Inf	null
7	Inf	null
8	Inf	null
9	Inf	null

<http://rosalind.info/problems/bf/>

# Practice Time - Pass 0

1 2 10  
 3 2 1  
 3 4 1  
 4 5 3  
 5 6 -1  
 7 6 -1  
 8 7 1  
 1 8 8  
 7 2 -4  
 2 6 2  
 6 3 -2  
 9 5 -10  
 9 4 7

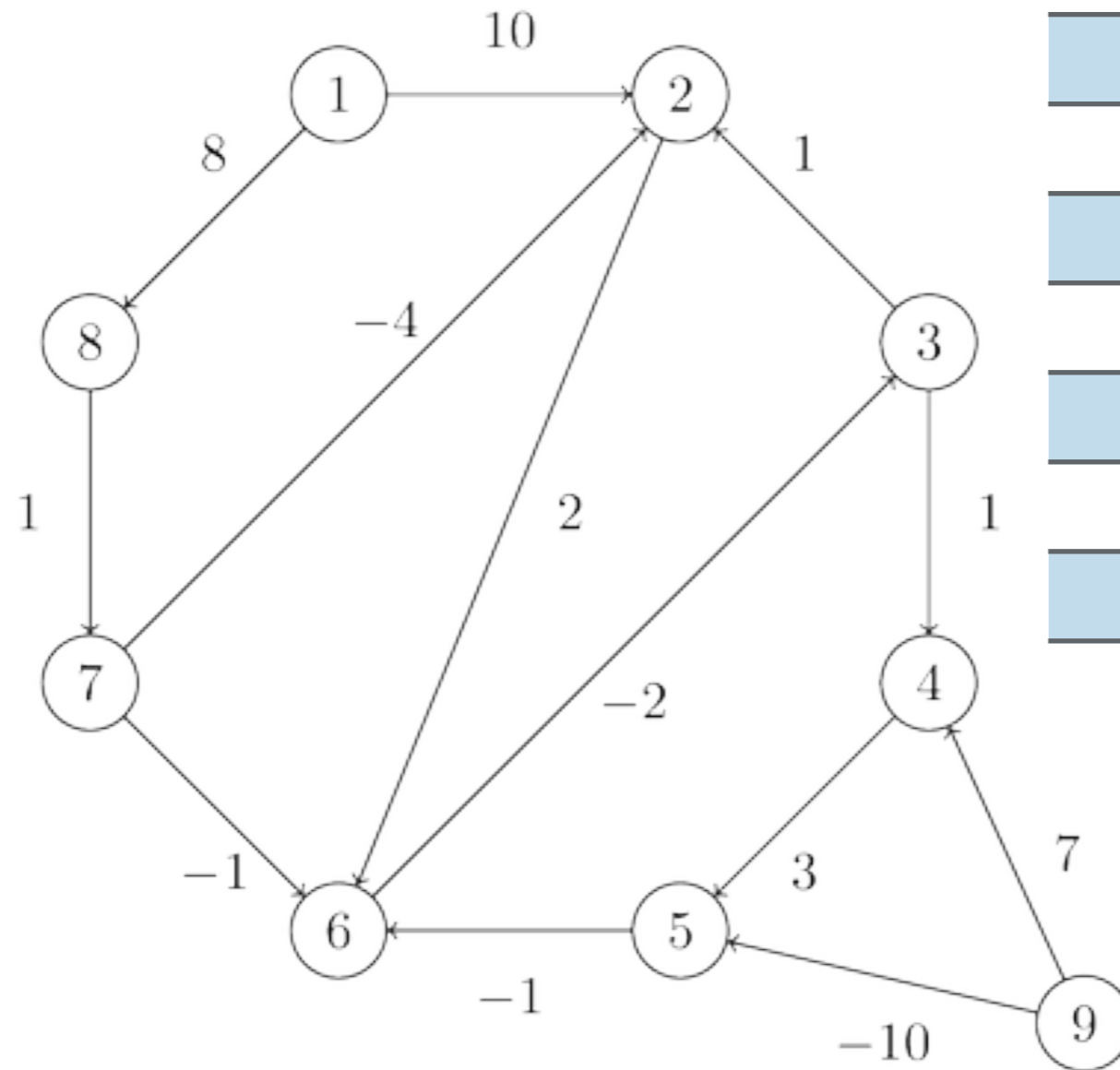


v	distTo[]	edgeTo[]
1	0	-
2	10	1->2
3	Inf	null
4	Inf	null
5	Inf	null
6	Inf	null
7	Inf	null
8	Inf	null
9	Inf	null

<http://rosalind.info/problems/bf/>

# Practice Time - Pass 0

1 2 10  
 3 2 1  
 3 4 1  
 4 5 3  
 5 6 -1  
 7 6 -1  
 8 7 1  
 1 8 8  
 7 2 -4  
 2 6 2  
 6 3 -2  
 9 5 -10  
 9 4 7



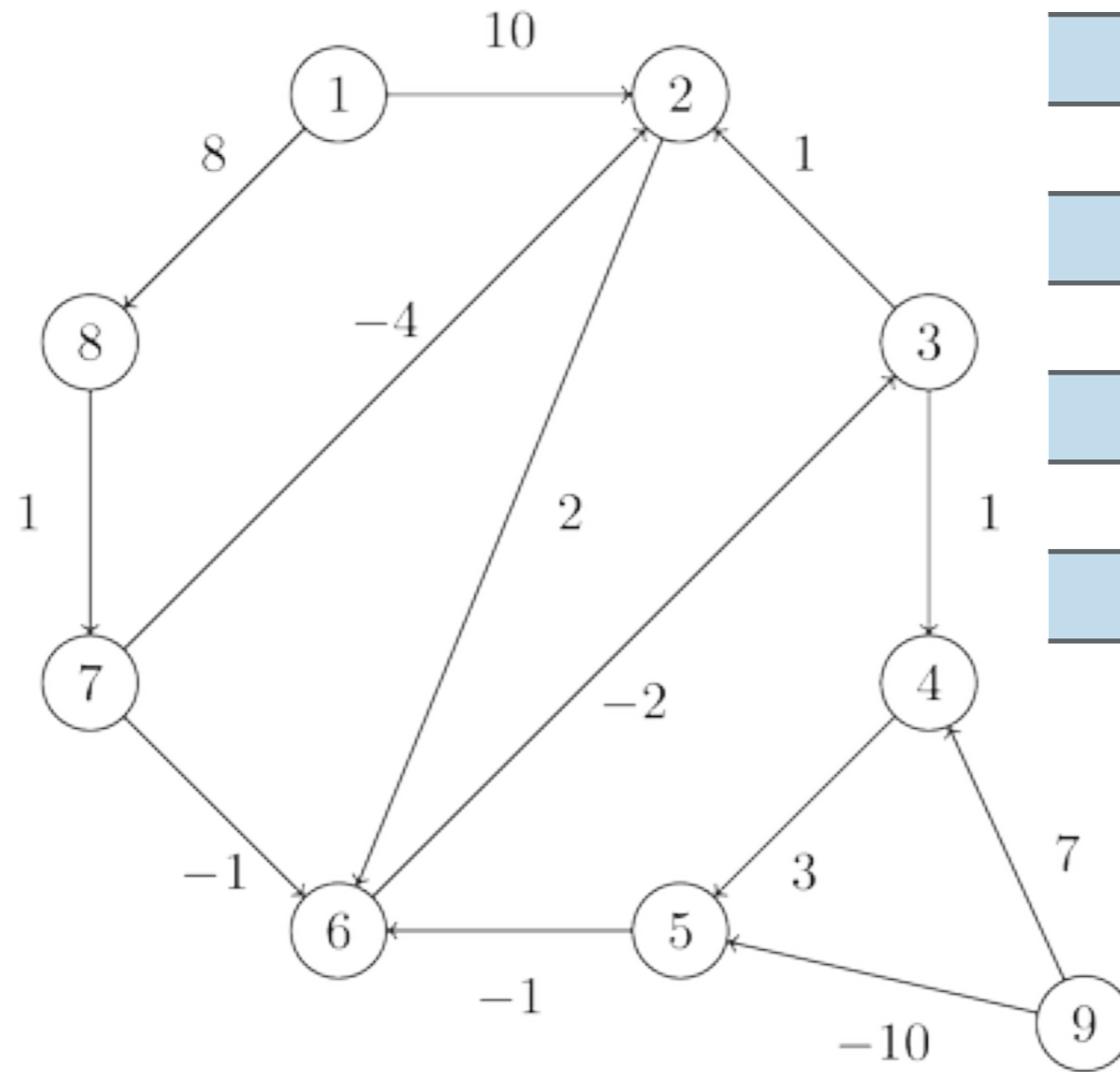
v	distTo[]	edgeTo[]
1	0	-
2	10	1-> 2
3	Inf	null
4	Inf	null
5	Inf	null
6	Inf	null
7	Inf	null
8	Inf	null
9	Inf	null

<http://rosalind.info/problems/bf/>

# Practice Time - Pass 0

```

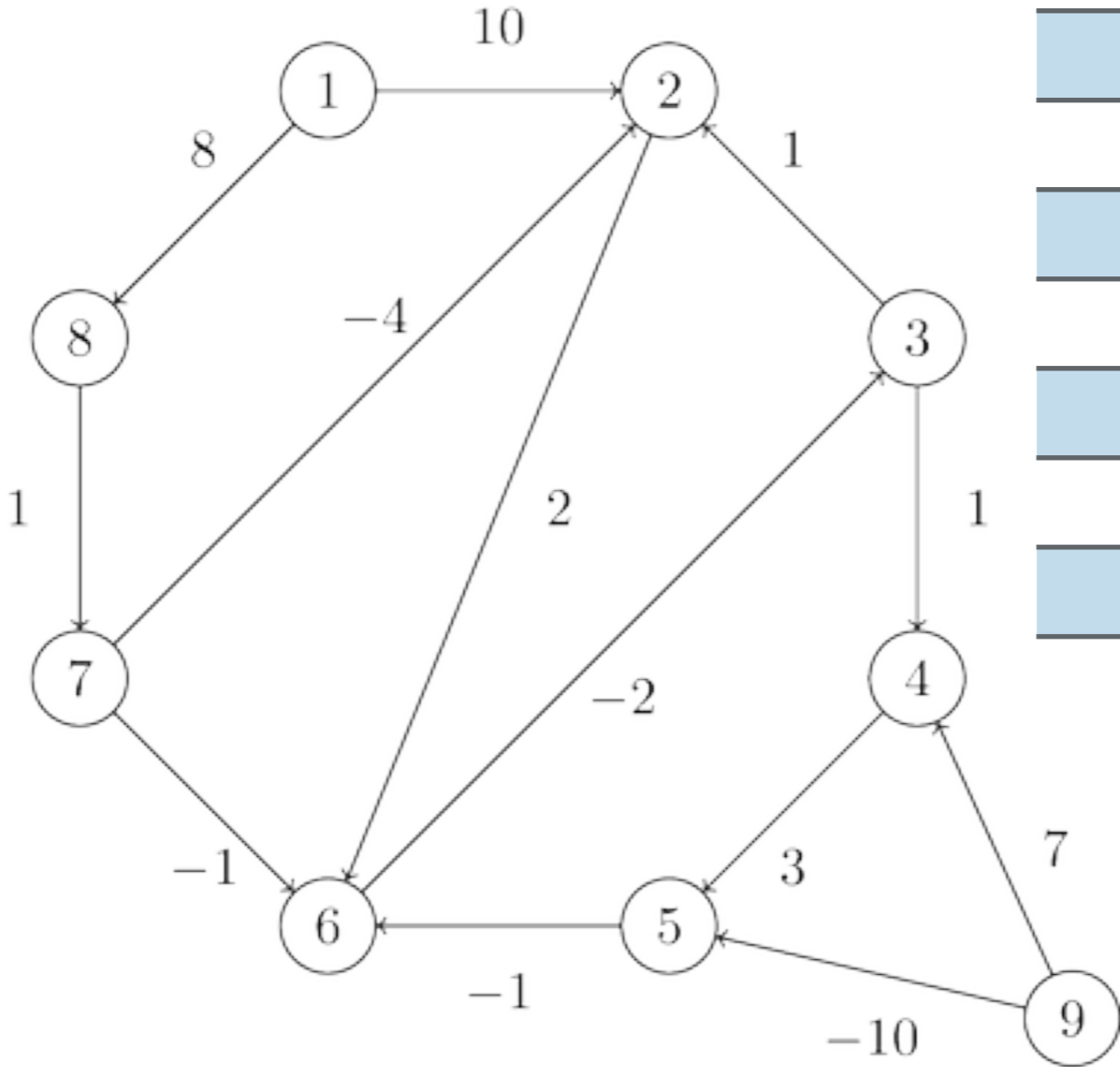
1 2 10
3 2 1
3 4 1
4 5 3
5 6 -1
7 6 -1
8 7 1
1 8 8
7 2 -4
2 6 2
6 3 -2
9 5 -10
9 4 7
    
```



v	distTo[]	edgeTo[]
1	0	-
2	10	1->2
3	Inf	null
4	Inf	null
5	Inf	null
6	Inf	null
7	Inf	null
8	Inf	null
9	Inf	null

# Practice Time - Pass 0

1 2 10  
 3 2 1  
 3 4 1  
 4 5 3  
 5 6 -1  
 7 6 -1  
 8 7 1  
 1 8 8  
 7 2 -4  
 2 6 2  
 6 3 -2  
 9 5 -10  
 9 4 7

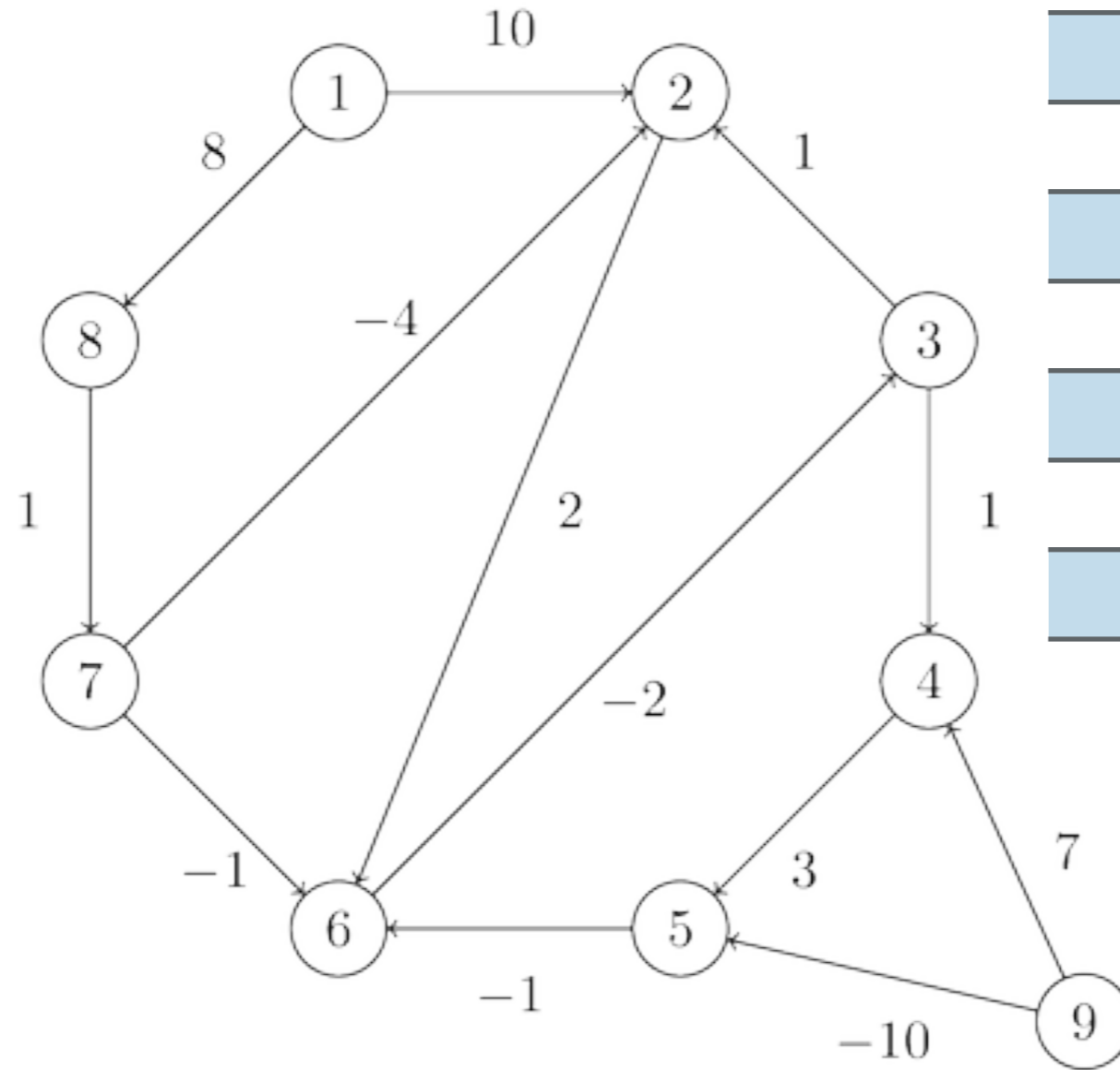


v	distTo[]	edgeTo[]
1	0	-
2	10	1-> 2
3	Inf	null
4	Inf	null
5	Inf	null
6	Inf	null
7	Inf	null
8	Inf	null
9	Inf	null

<http://rosalind.info/problems/bf/>

# Practice Time - Pass 0

1 2 10  
 3 2 1  
 3 4 1  
 4 5 3  
 5 6 -1  
 7 6 -1  
 8 7 1  
 1 8 8  
 7 2 -4  
 2 6 2  
 6 3 -2  
 9 5 -10  
 9 4 7

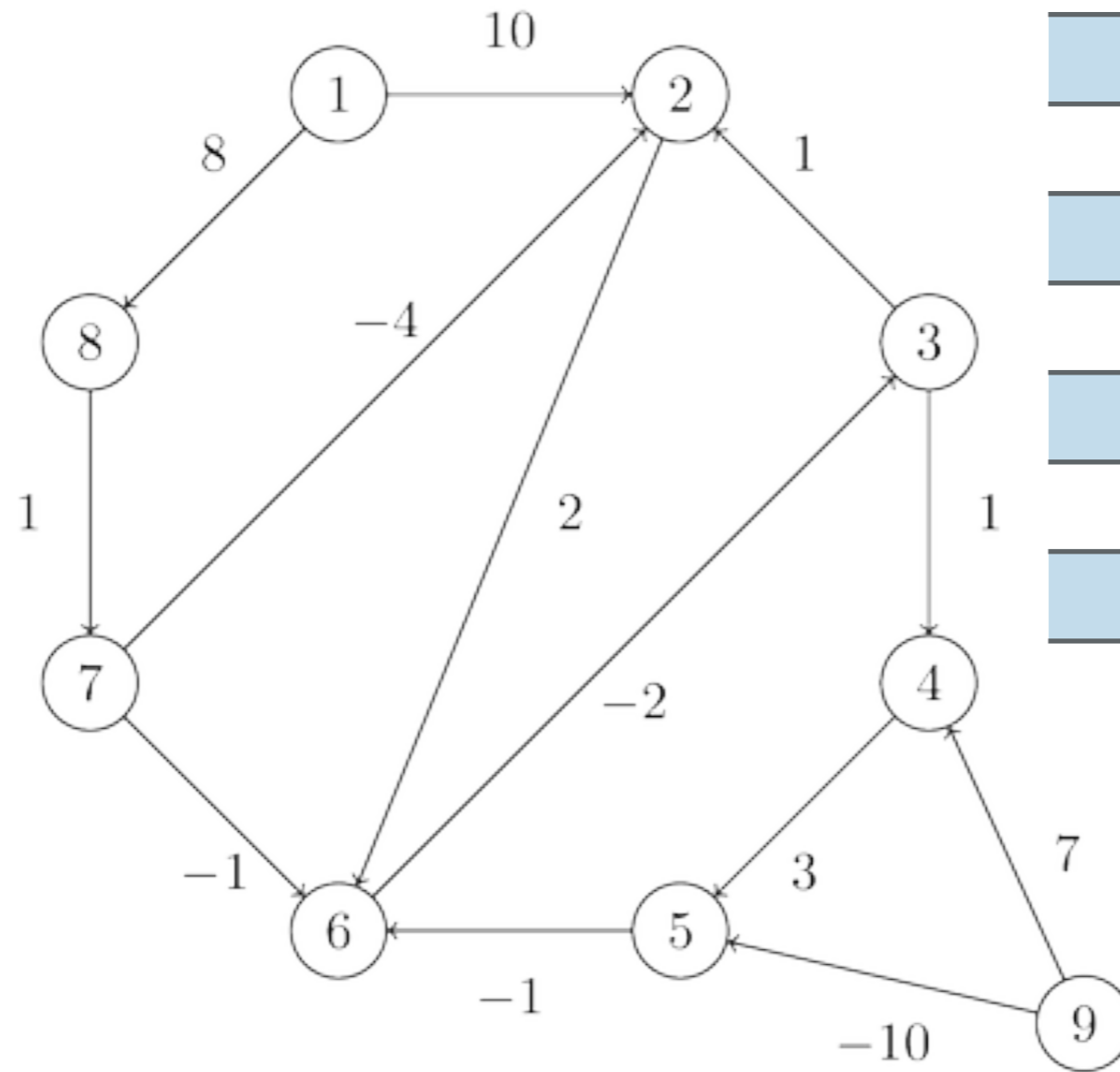


v	distTo[]	edgeTo[]
1	0	-
2	10	1-> 2
3	Inf	null
4	Inf	null
5	Inf	null
6	Inf	null
7	Inf	null
8	Inf	null
9	Inf	null

# Practice Time - Pass 0

```

1 2 10
3 2 1
3 4 1
4 5 3
5 6 -1
7 6 -1
8 7 1
1 8 8
7 2 -4
2 6 2
6 3 -2
9 5 -10
9 4 7
    
```

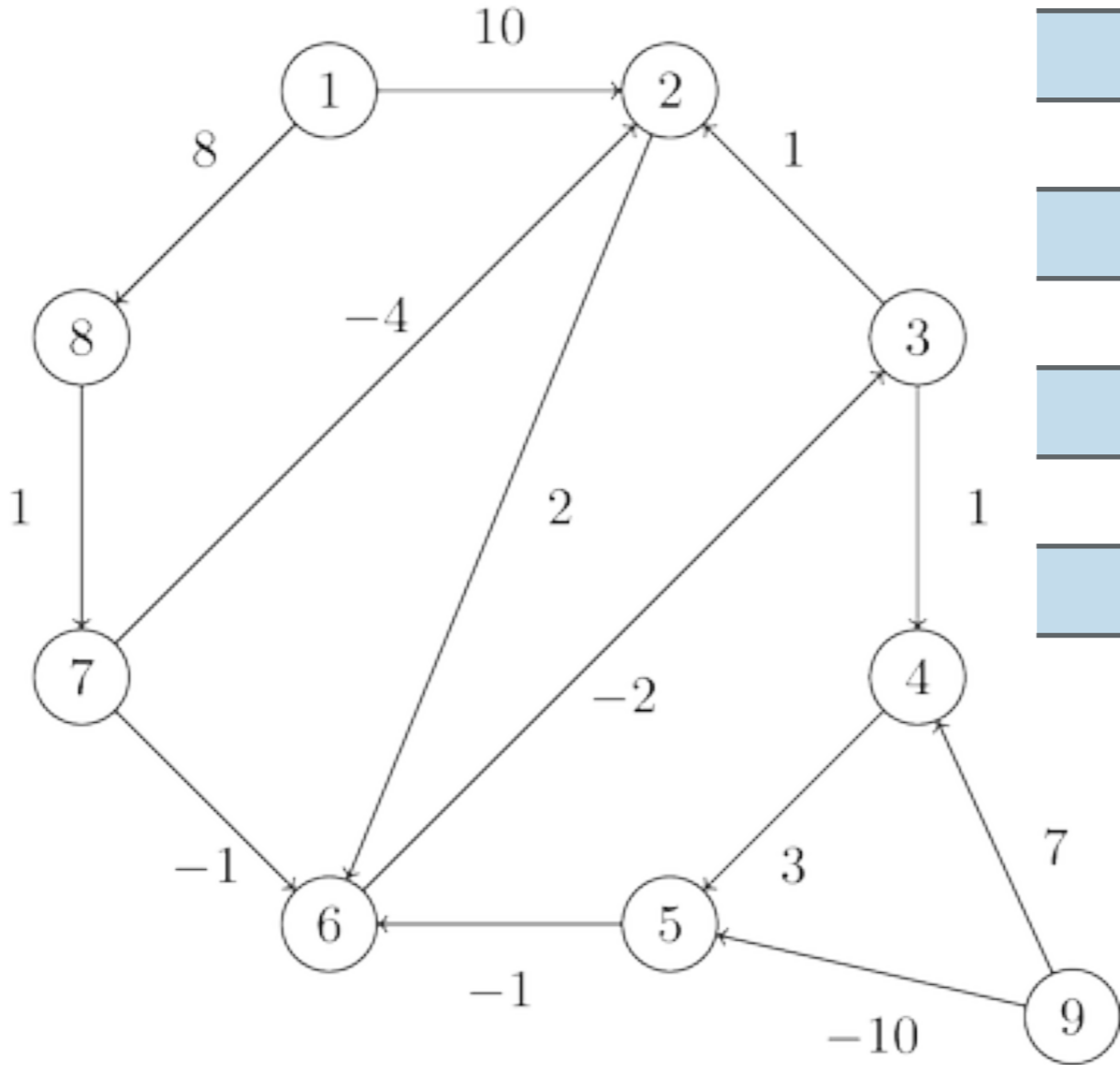


v	distTo[]	edgeTo[]
1	0	-
2	10	1->2
3	Inf	null
4	Inf	null
5	Inf	null
6	Inf	null
7	Inf	null
8	Inf	null
9	Inf	null



# Practice Time - Pass 0

1 2 10  
 3 2 1  
 3 4 1  
 4 5 3  
 5 6 -1  
 7 6 -1  
 8 7 1  
 1 8 8  
 7 2 -4  
 2 6 2  
 6 3 -2  
 9 5 -10  
 9 4 7

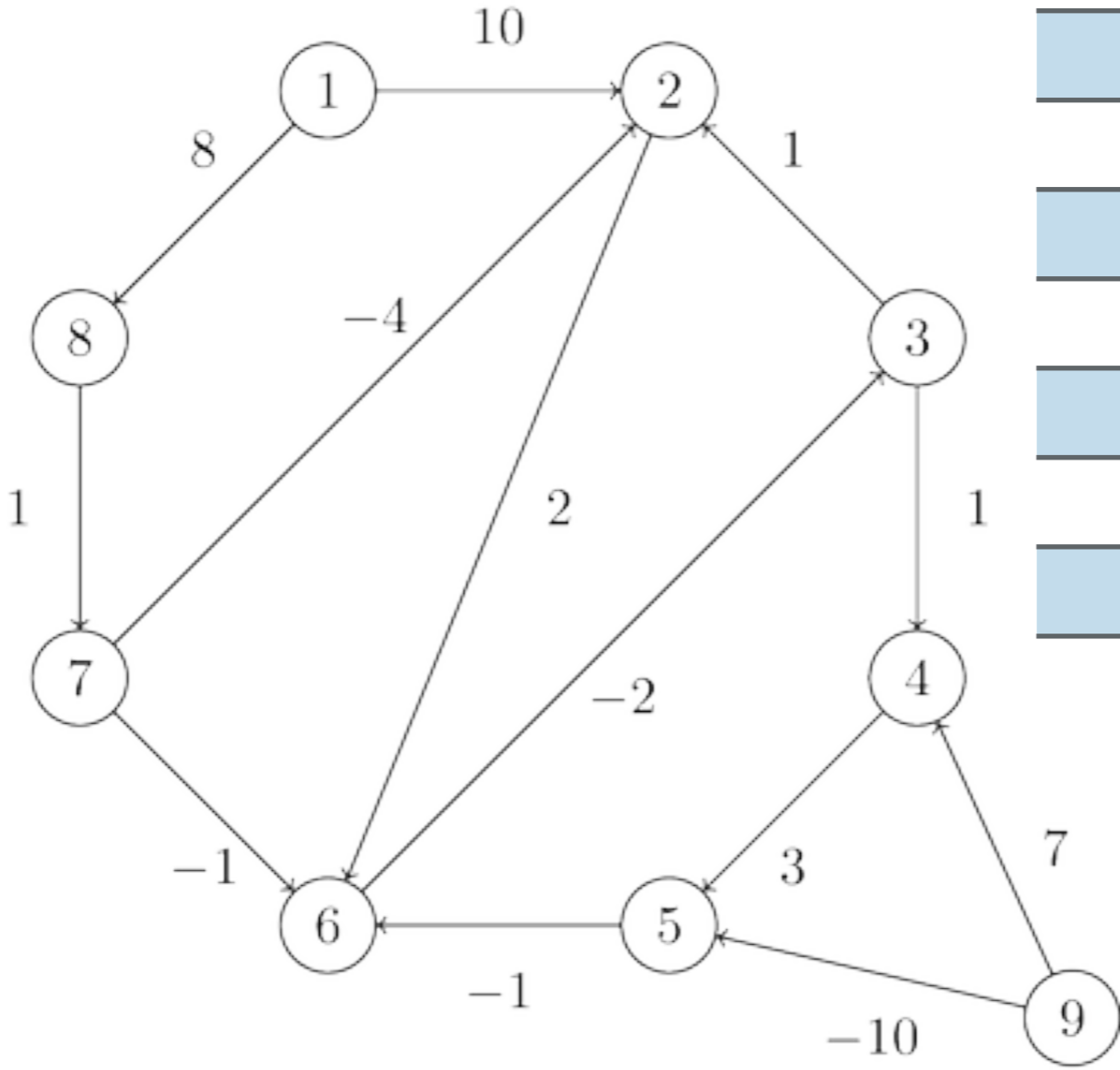


v	distTo[]	edgeTo[]
1	0	-
2	10	1-> 2
3	Inf	null
4	Inf	null
5	Inf	null
6	Inf	null
7	Inf	null
8	Inf	null
9	Inf	null

<http://rosalind.info/problems/bf/>

# Practice Time - Pass 0

1 2 10  
 3 2 1  
 3 4 1  
 4 5 3  
 5 6 -1  
 7 6 -1  
 8 7 1  
 1 8 8  
 7 2 -4  
 2 6 2  
 6 3 -2  
 9 5 -10  
 9 4 7

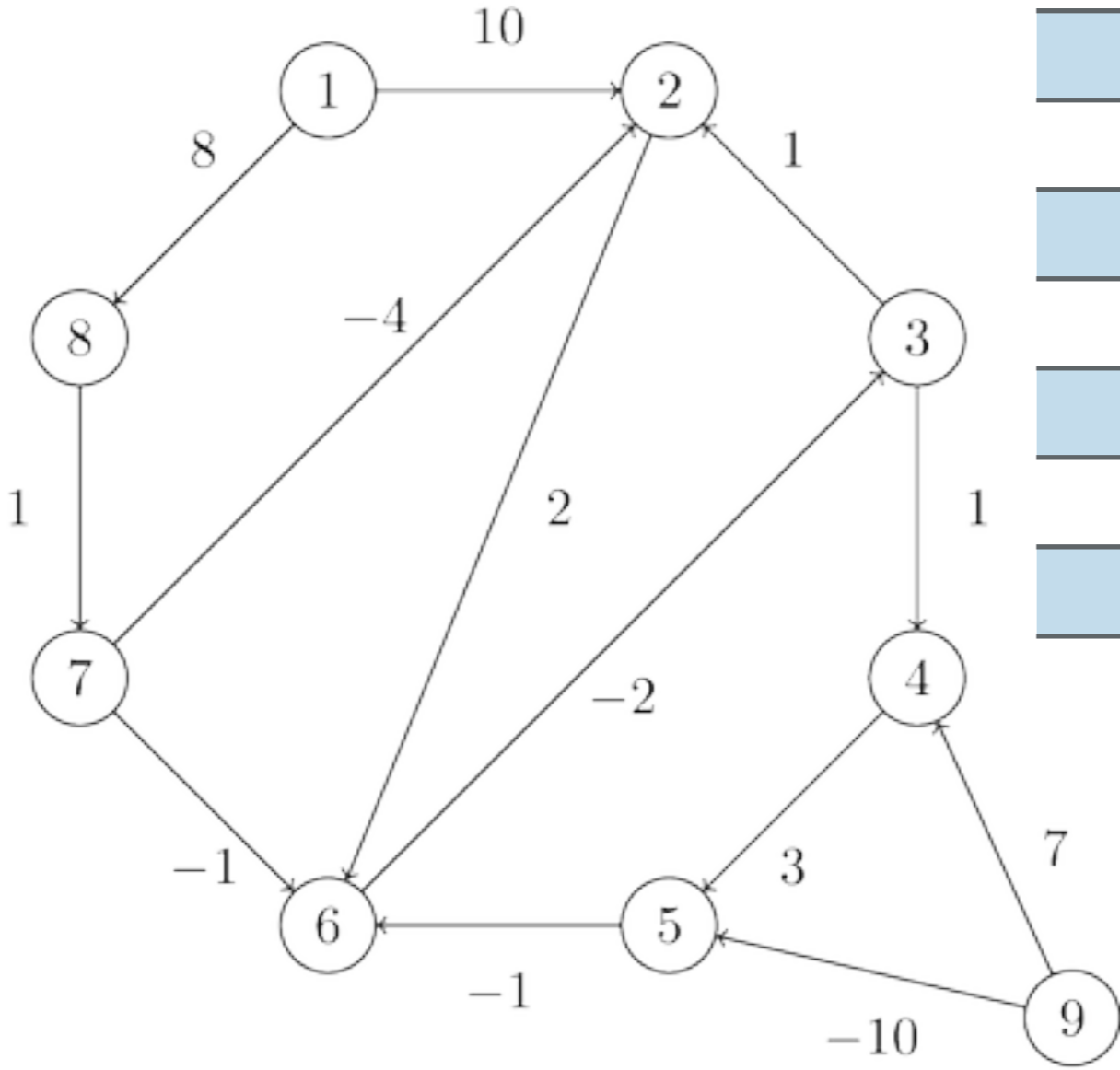


v	distTo[]	edgeTo[]
1	0	-
2	10	1->2
3	Inf	null
4	Inf	null
5	Inf	null
6	Inf	null
7	Inf	null
8	8	1->8
9	Inf	null

<http://rosalind.info/problems/bf/>

# Practice Time - Pass 0

1 2 10  
 3 2 1  
 3 4 1  
 4 5 3  
 5 6 -1  
 7 6 -1  
 8 7 1  
 1 8 8  
**7 2 -4**  
 2 6 2  
 6 3 -2  
 9 5 -10  
 9 4 7

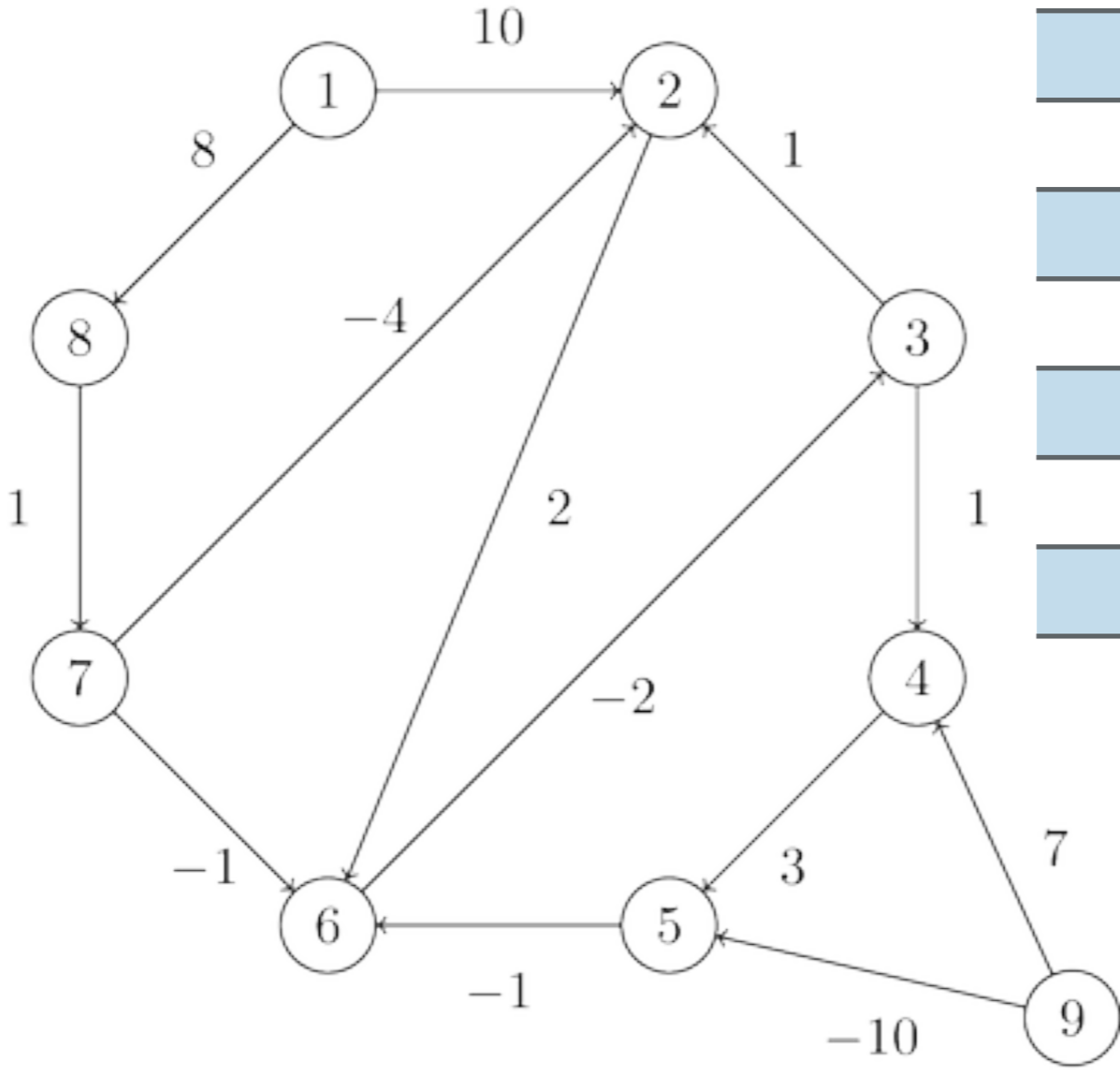


v	distTo[]	edgeTo[]
1	0	-
2	10	1->2
3	Inf	null
4	Inf	null
5	Inf	null
6	Inf	null
7	Inf	null
8	8	1->8
9	Inf	null

<http://rosalind.info/problems/bf/>

# Practice Time - Pass 0

1 2 10  
 3 2 1  
 3 4 1  
 4 5 3  
 5 6 -1  
 7 6 -1  
 8 7 1  
 1 8 8  
 7 2 -4  
 2 6 2  
 6 3 -2  
 9 5 -10  
 9 4 7

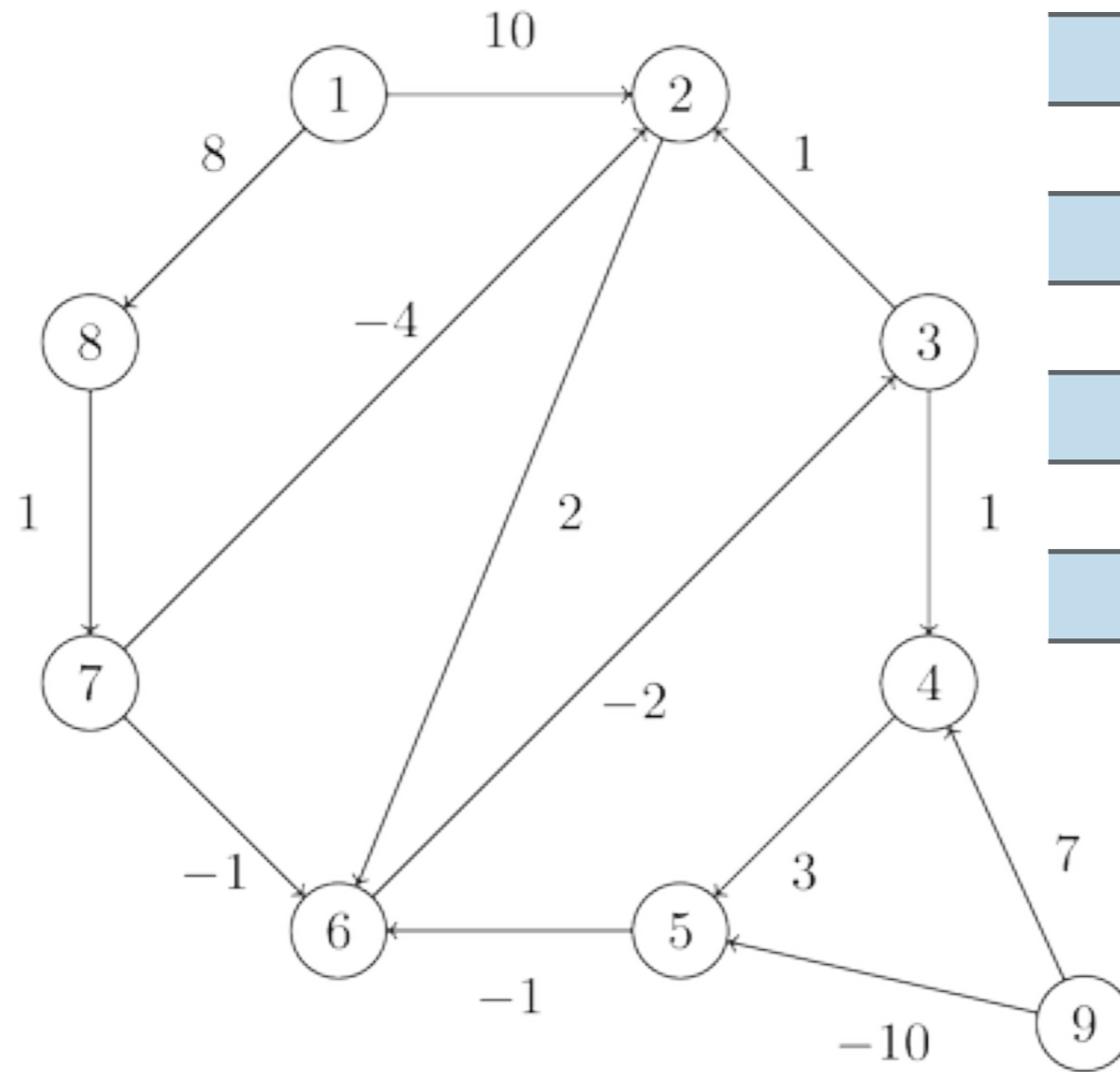


v	distTo[]	edgeTo[]
1	0	-
2	10	1->2
3	Inf	null
4	Inf	null
5	Inf	null
6	12	2->6
7	Inf	null
8	8	1->8
9	Inf	null

<http://rosalind.info/problems/bf/>

# Practice Time - Pass 0

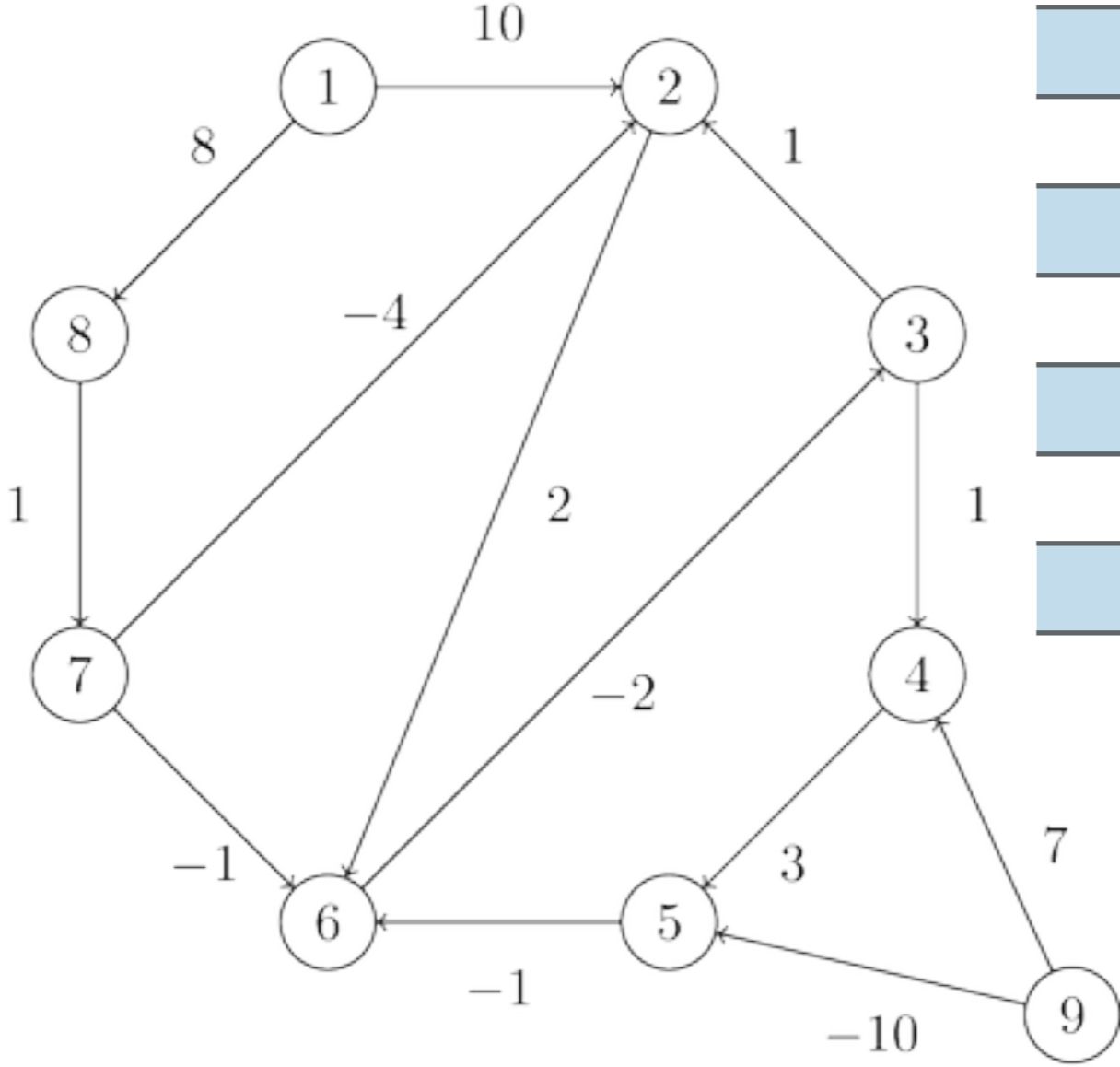
1 2 10  
 3 2 1  
 3 4 1  
 4 5 3  
 5 6 -1  
 7 6 -1  
 8 7 1  
 1 8 8  
 7 2 -4  
 2 6 2  
 6 3 -2  
 9 5 -10  
 9 4 7



v	distTo[]	edgeTo[]
1	0	-
2	10	1->2
3	10	6->3
4	Inf	null
5	Inf	null
6	12	2->6
7	Inf	null
8	8	1->8
9	Inf	null

# Practice Time - Pass 0

1 2 10  
 3 2 1  
 3 4 1  
 4 5 3  
 5 6 -1  
 7 6 -1  
 8 7 1  
 1 8 8  
 7 2 -4  
 2 6 2  
 6 3 -2  
 9 5 -10  
 9 4 7

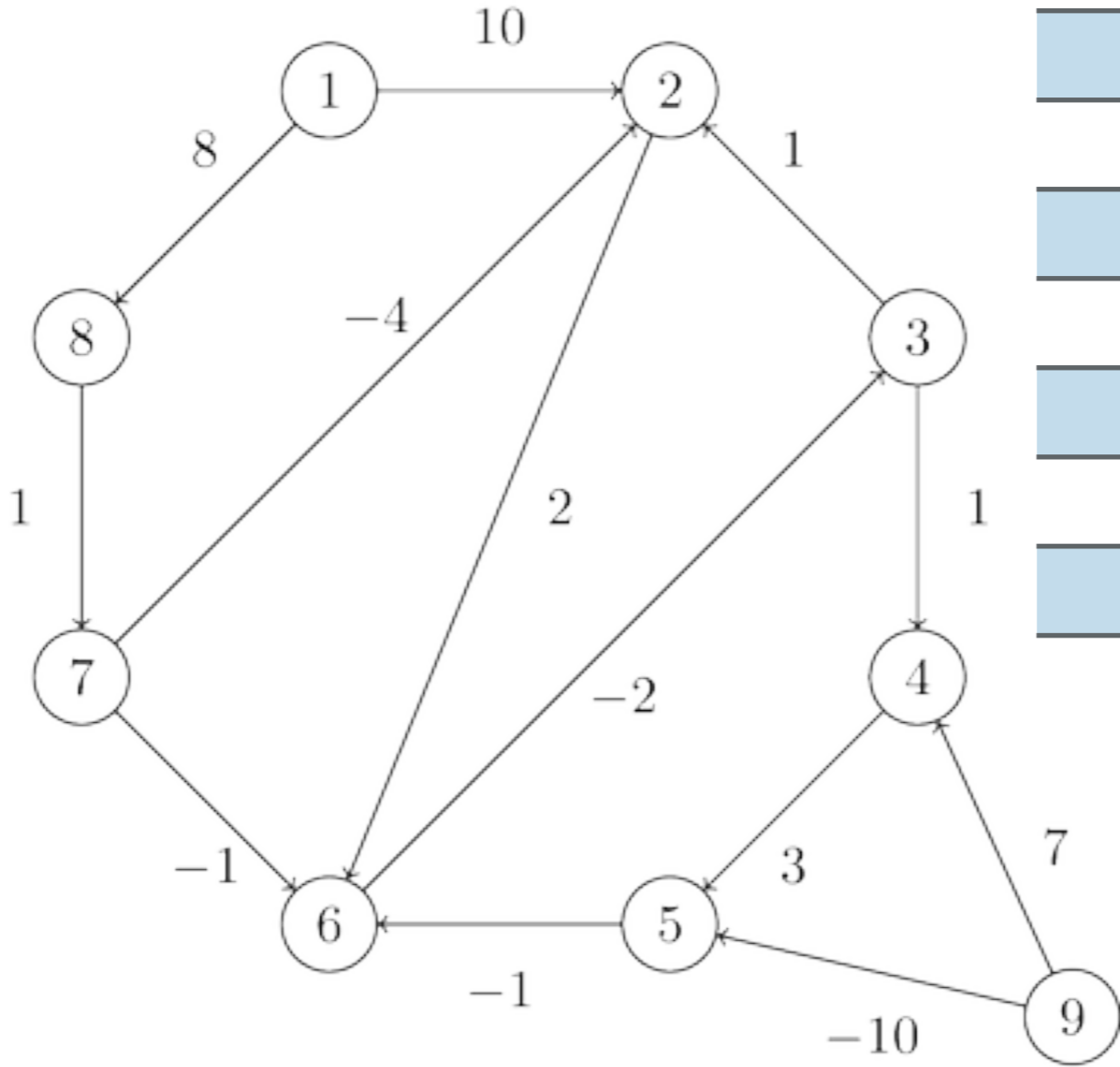


v	distTo[]	edgeTo[]
1	0	-
2	10	1->2
3	10	6->3
4	Inf	null
5	Inf	null
6	12	2->6
7	Inf	null
8	8	1->8
9	Inf	null

<http://rosalind.info/problems/bf/>

# Practice Time - Pass 0

1 2 10  
 3 2 1  
 3 4 1  
 4 5 3  
 5 6 -1  
 7 6 -1  
 8 7 1  
 1 8 8  
 7 2 -4  
 2 6 2  
 6 3 -2  
 9 5 -10  
 9 4 7

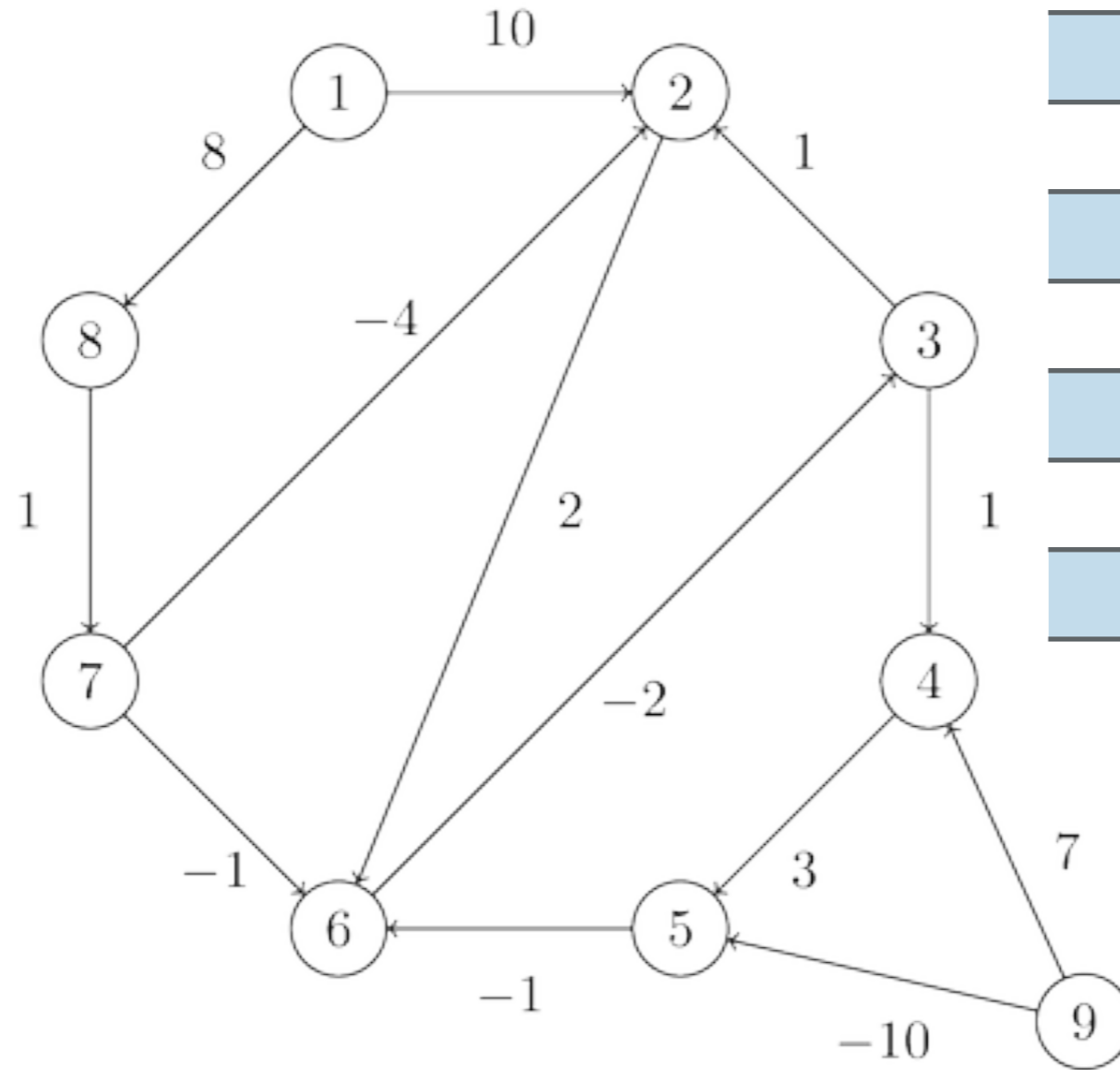


v	distTo[]	edgeTo[]
1	0	-
2	10	1->2
3	10	6->3
4	Inf	null
5	Inf	null
6	12	2->6
7	Inf	null
8	8	1->8
9	Inf	null

<http://rosalind.info/problems/bf/>

# Practice Time - Pass 1

1 2 10  
 3 2 1  
 3 4 1  
 4 5 3  
 5 6 -1  
 7 6 -1  
 8 7 1  
 1 8 8  
 7 2 -4  
 2 6 2  
 6 3 -2  
 9 5 -10  
 9 4 7

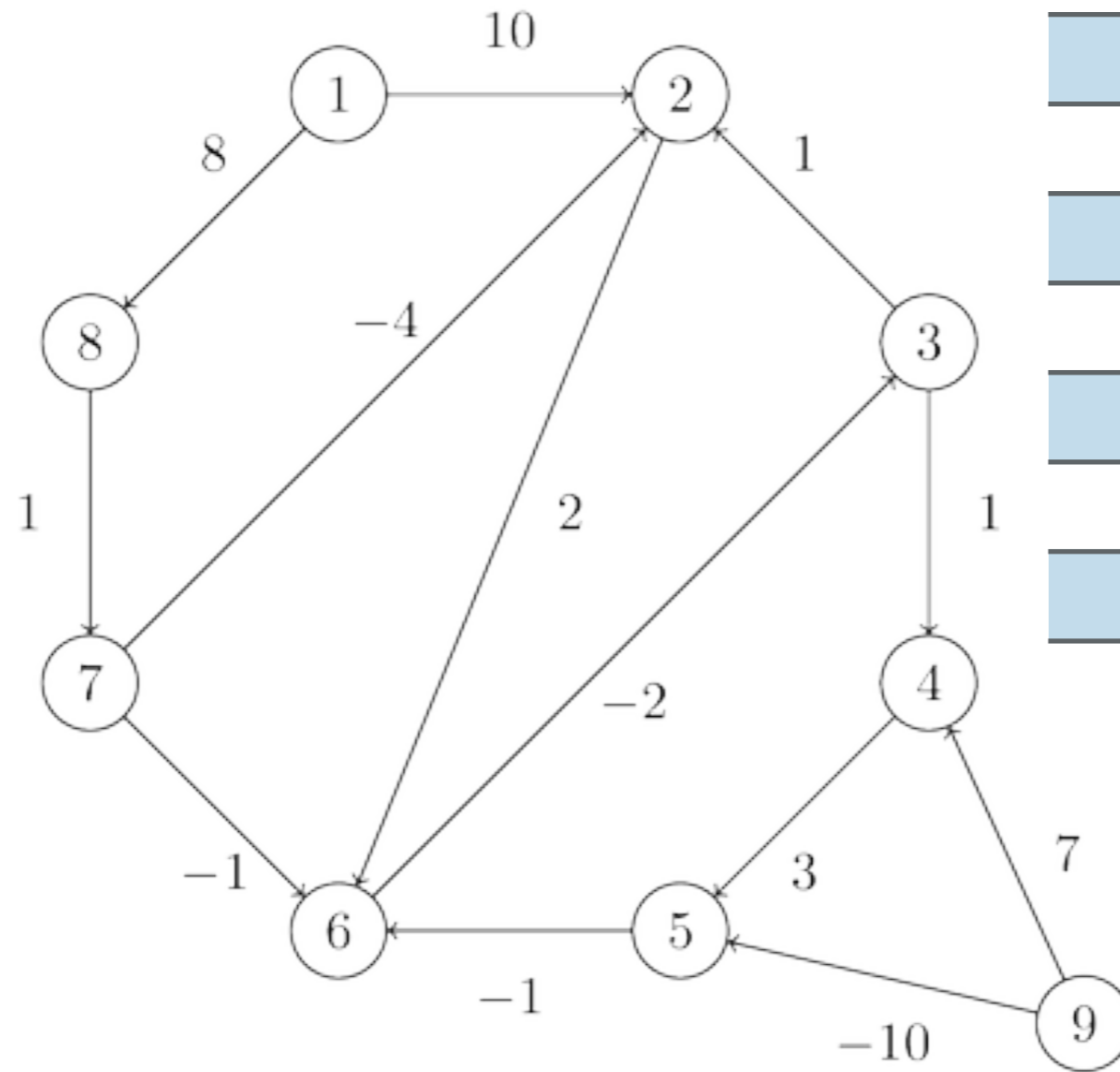


v	distTo[]	edgeTo[]
1	0	-
2	10	1->2
3	10	6->3
4	Inf	null
5	Inf	null
6	12	2->6
7	Inf	null
8	8	1->8
9	Inf	null



# Practice Time - Pass 1

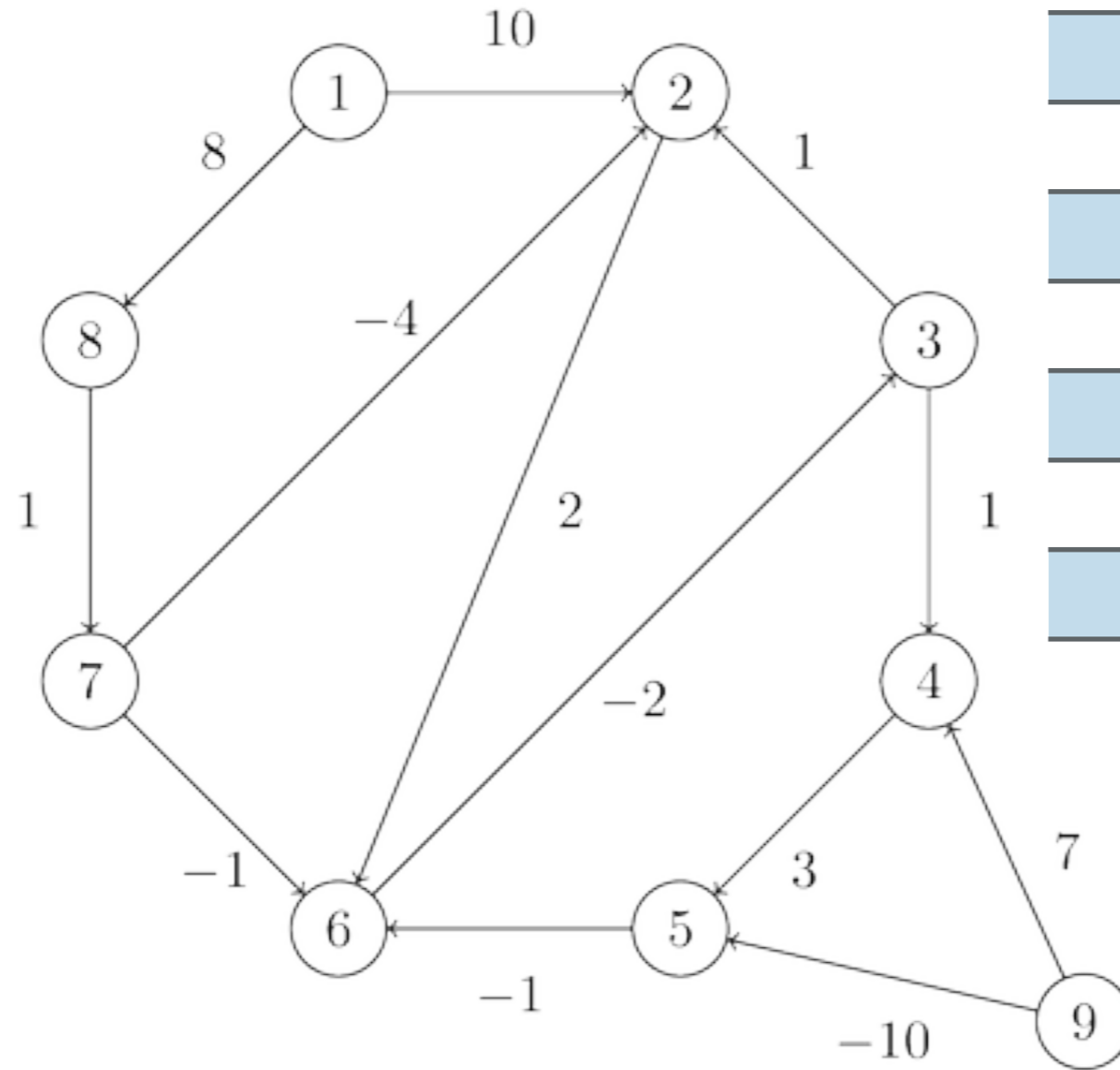
1 2 10  
 3 2 1  
 3 4 1  
 4 5 3  
 5 6 -1  
 7 6 -1  
 8 7 1  
 1 8 8  
 7 2 -4  
 2 6 2  
 6 3 -2  
 9 5 -10  
 9 4 7



v	distTo[]	edgeTo[]
1	0	-
2	10	1->2
3	10	6->3
4	Inf	null
5	Inf	null
6	12	2->6
7	Inf	null
8	8	1->8
9	Inf	null

# Practice Time - Pass 1

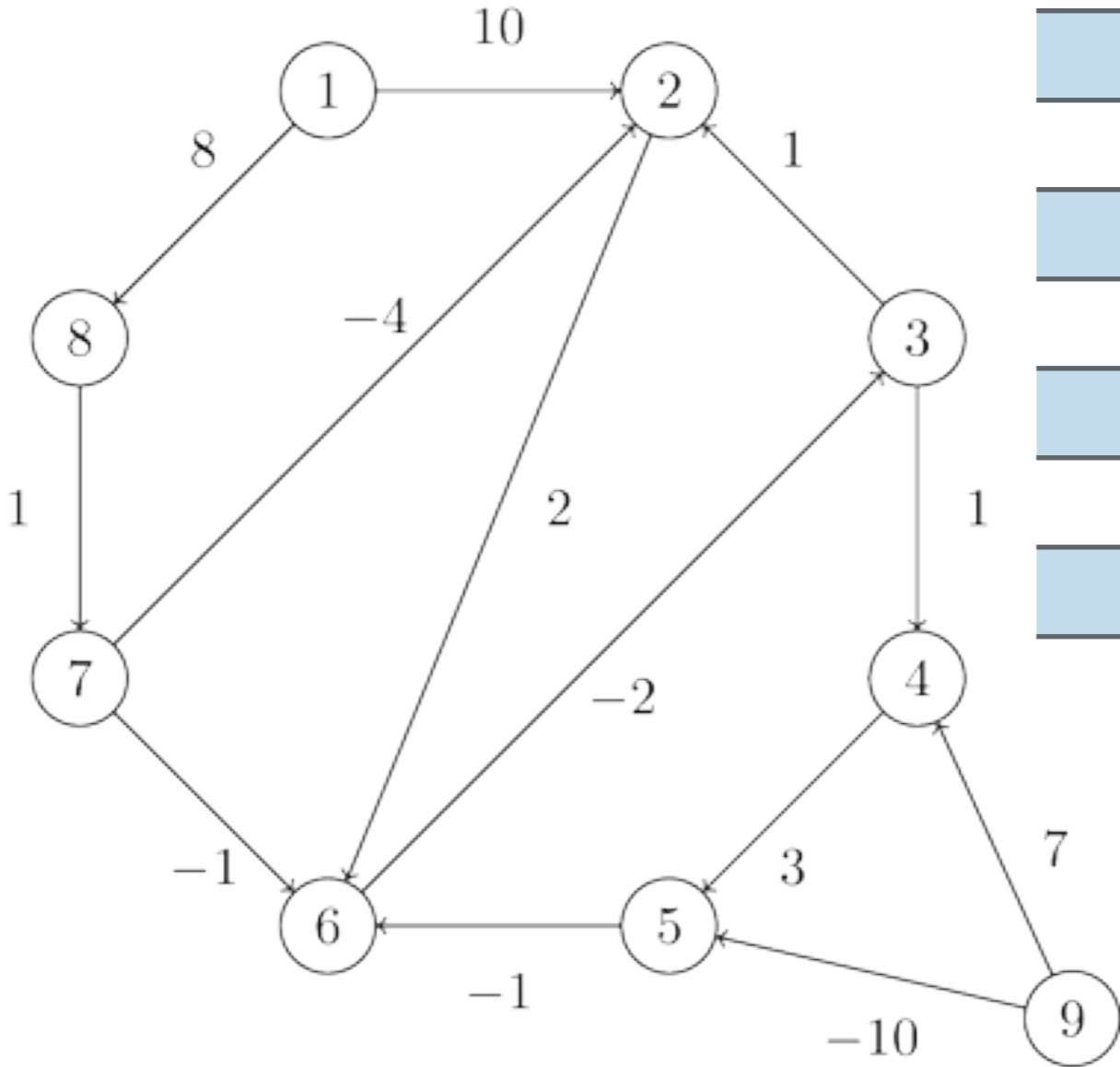
1 2 10  
 3 2 1  
 3 4 1  
 4 5 3  
 5 6 -1  
 7 6 -1  
 8 7 1  
 1 8 8  
 7 2 -4  
 2 6 2  
 6 3 -2  
 9 5 -10  
 9 4 7



v	distTo[]	edgeTo[]
1	0	-
2	10	1->2
3	10	6->3
4	11	3->4
5	Inf	null
6	12	2->6
7	Inf	null
8	8	1->8
9	Inf	null

# Practice Time - Pass 1

1 2 10  
 3 2 1  
 3 4 1  
 4 5 3  
 5 6 -1  
 7 6 -1  
 8 7 1  
 1 8 8  
 7 2 -4  
 2 6 2  
 6 3 -2  
 9 5 -10  
 9 4 7

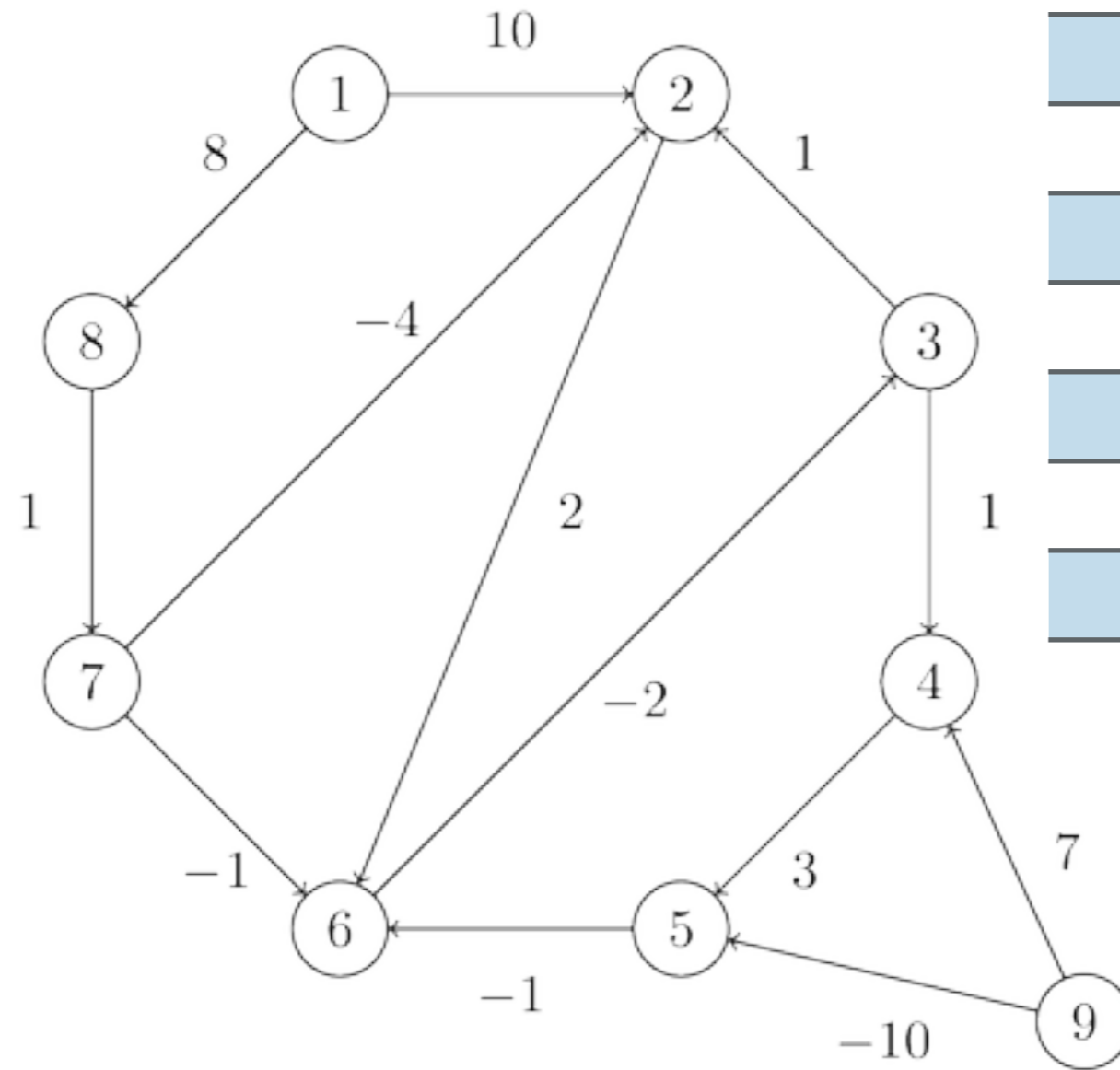


v	distTo[]	edgeTo[]
1	0	-
2	10	1->2
3	10	6->3
4	11	3->4
5	14	4->5
6	12	2->6
7	Inf	null
8	8	1->8
9	Inf	null

<http://rosalind.info/problems/bf/>

# Practice Time - Pass 1

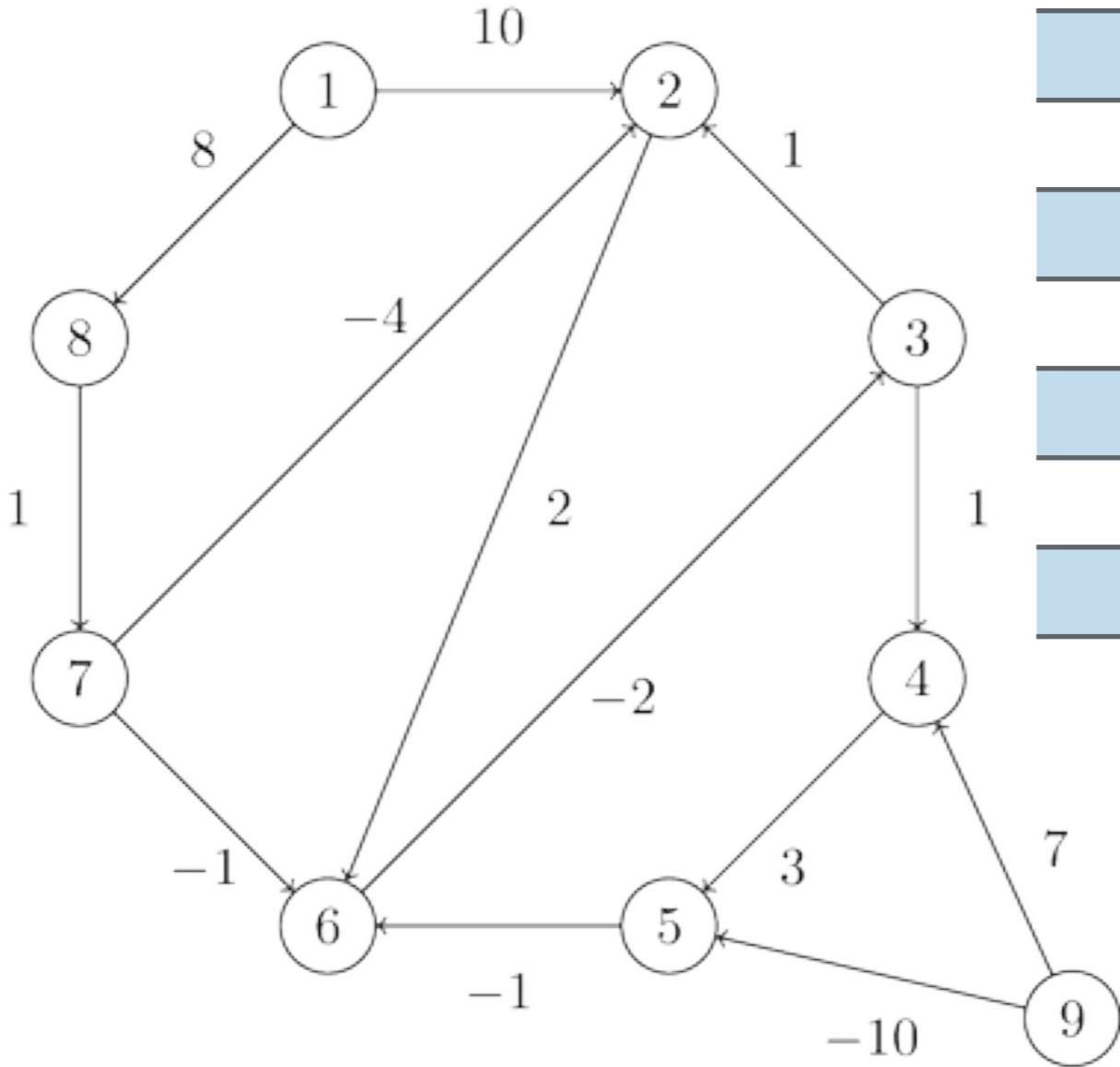
1 2 10  
 3 2 1  
 3 4 1  
 4 5 3  
 5 6 -1  
 7 6 -1  
 8 7 1  
 1 8 8  
 7 2 -4  
 2 6 2  
 6 3 -2  
 9 5 -10  
 9 4 7



v	distTo[]	edgeTo[]
1	0	-
2	10	1->2
3	10	6->3
4	11	3->4
5	14	4->5
6	12	2->6
7	Inf	null
8	8	1->8
9	Inf	null

# Practice Time - Pass 1

1 2 10  
 3 2 1  
 3 4 1  
 4 5 3  
 5 6 -1  
 7 6 -1  
 8 7 1  
 1 8 8  
 7 2 -4  
 2 6 2  
 6 3 -2  
 9 5 -10  
 9 4 7

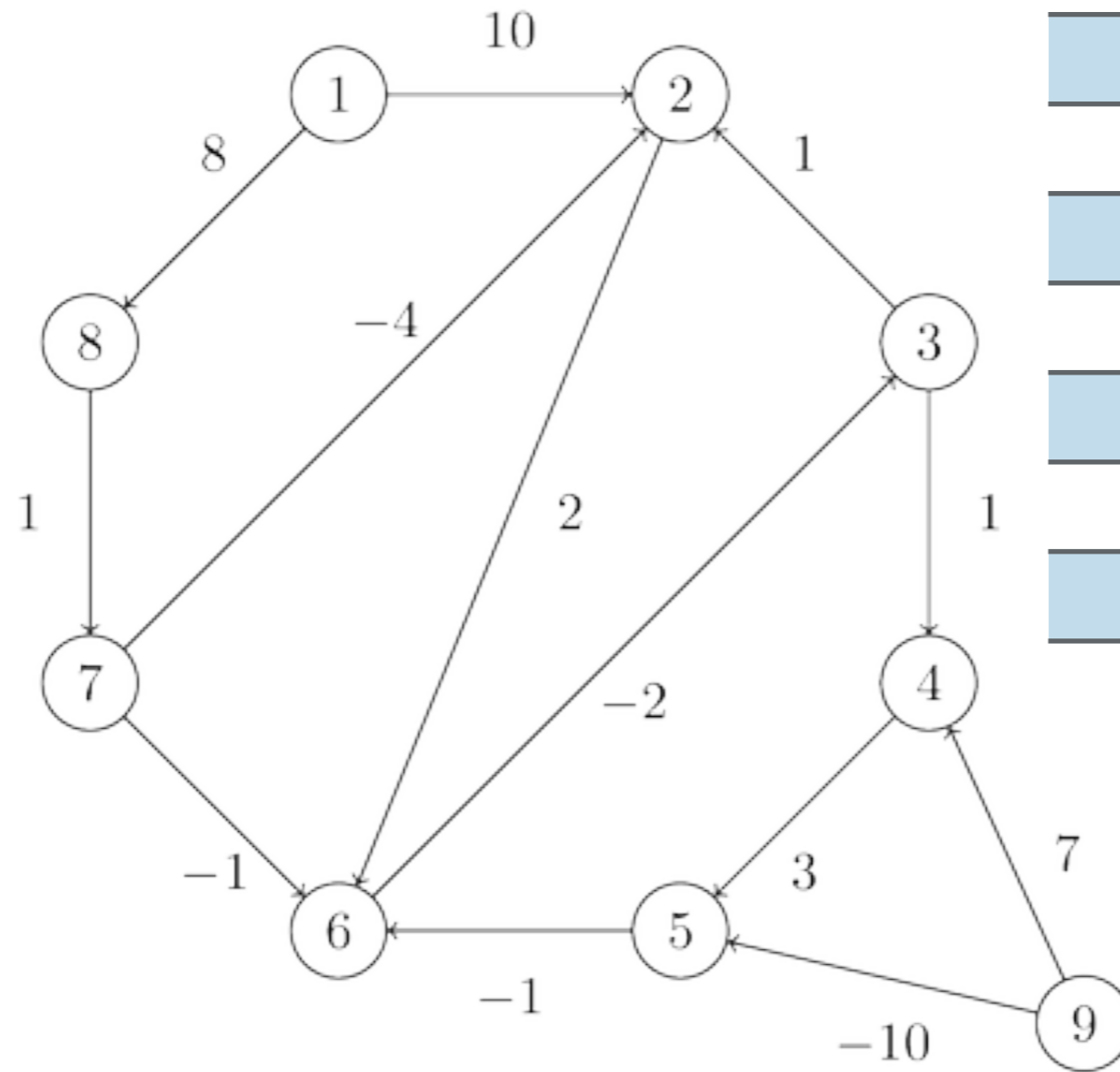


v	distTo[]	edgeTo[]
1	0	-
2	10	1->2
3	10	6->3
4	11	3->4
5	14	4->5
6	12	2->6
7	Inf	null
8	8	1->8
9	Inf	null

<http://rosalind.info/problems/bf/>

# Practice Time - Pass 1

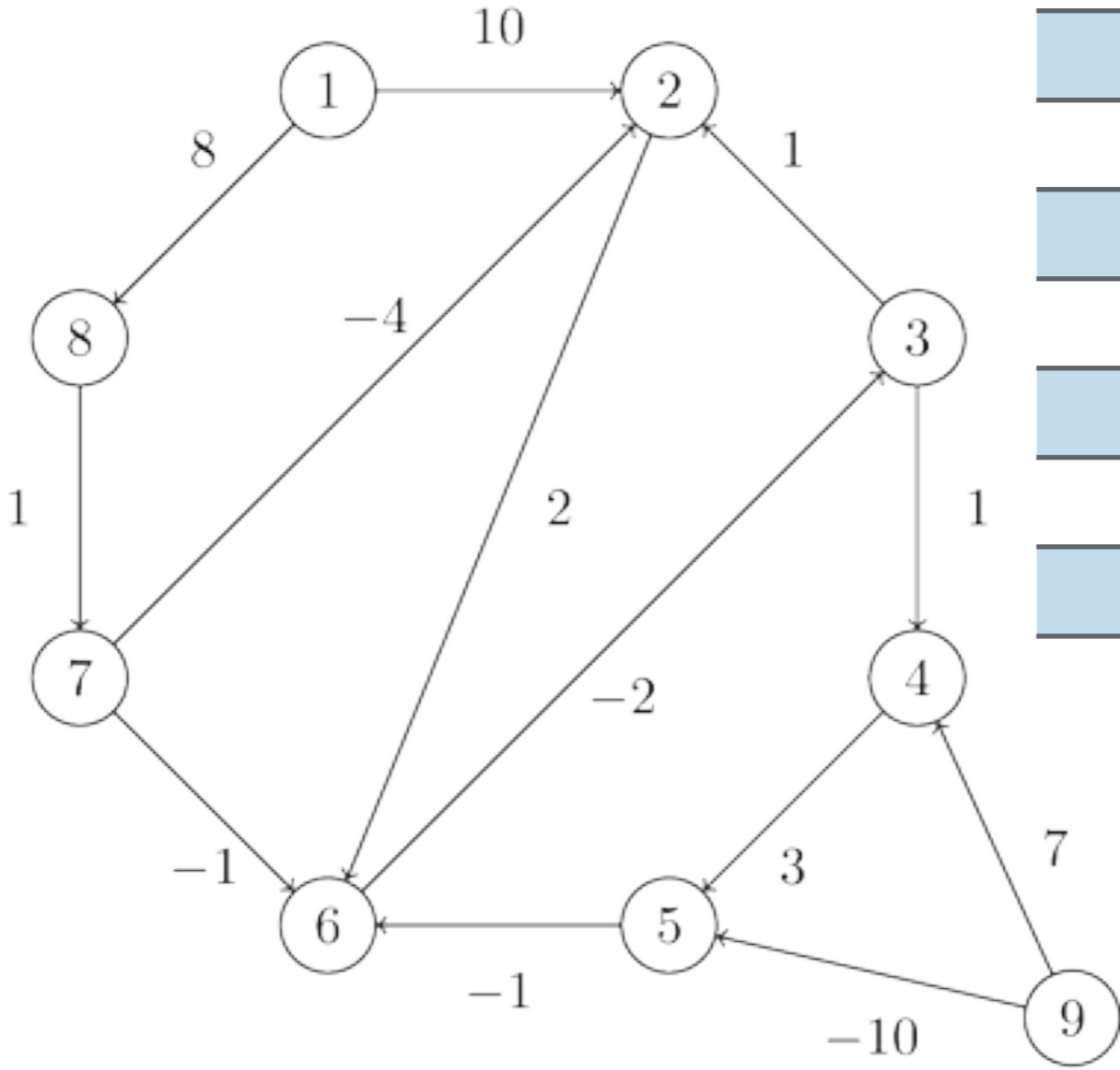
1 2 10  
 3 2 1  
 3 4 1  
 4 5 3  
 5 6 -1  
 7 6 -1  
 8 7 1  
 1 8 8  
 7 2 -4  
 2 6 2  
 6 3 -2  
 9 5 -10  
 9 4 7



v	distTo[]	edgeTo[]
1	0	-
2	10	1->2
3	10	6->3
4	11	3->4
5	14	4->5
6	12	2->6
7	9	8->7
8	8	1->8
9	Inf	null

# Practice Time - Pass 1

1 2 10  
 3 2 1  
 3 4 1  
 4 5 3  
 5 6 -1  
 7 6 -1  
 8 7 1  
 1 8 8  
 7 2 -4  
 2 6 2  
 6 3 -2  
 9 5 -10  
 9 4 7

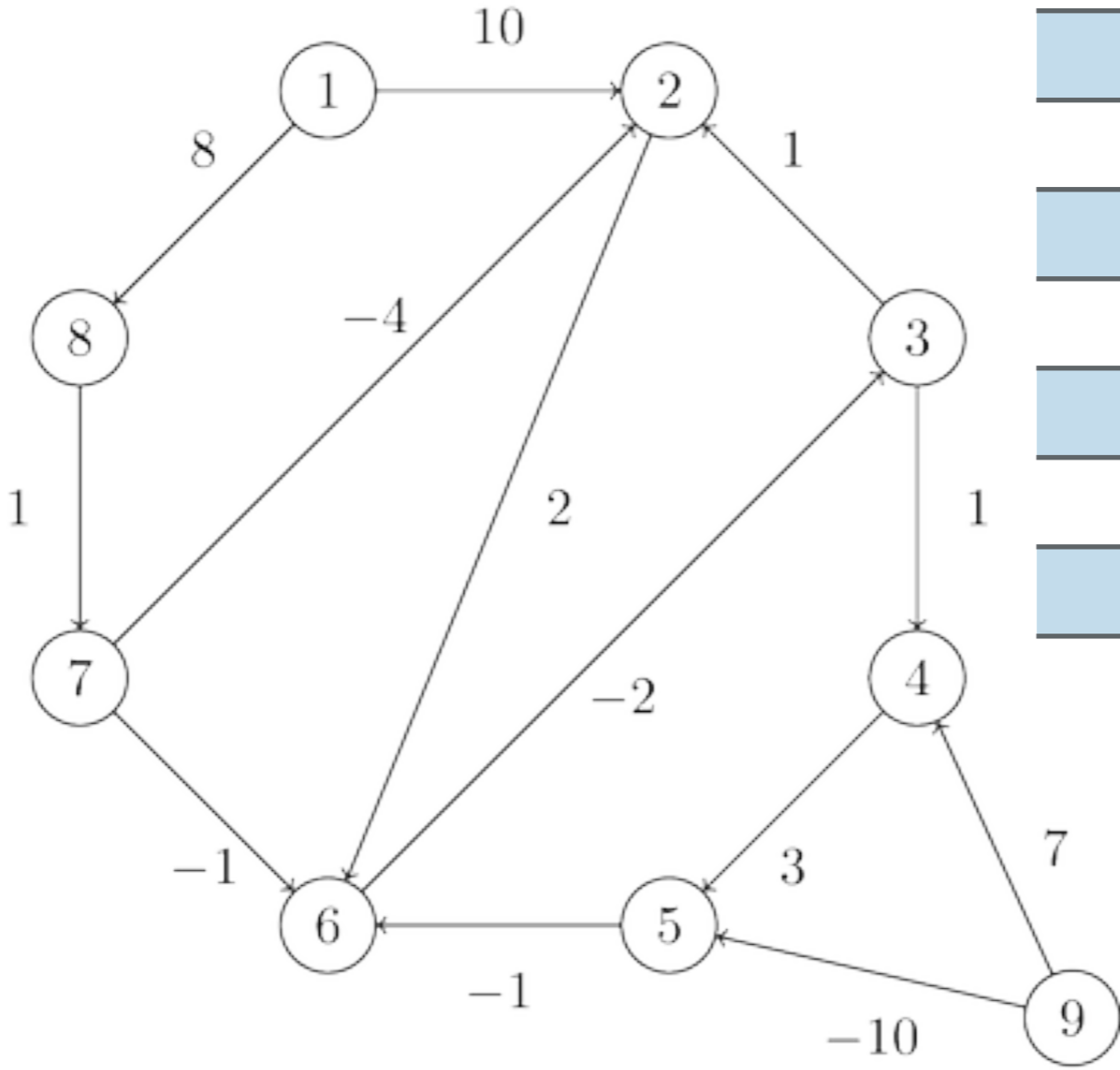


v	distTo[]	edgeTo[]
1	0	-
2	10	1->2
3	10	6->3
4	11	3->4
5	14	4->5
6	12	2->6
7	9	8->7
8	8	1->8
9	Inf	null

<http://rosalind.info/problems/bf/>

# Practice Time - Pass 1

1 2 10  
 3 2 1  
 3 4 1  
 4 5 3  
 5 6 -1  
 7 6 -1  
 8 7 1  
 1 8 8  
**7 2 -4**  
 2 6 2  
 6 3 -2  
 9 5 -10  
 9 4 7



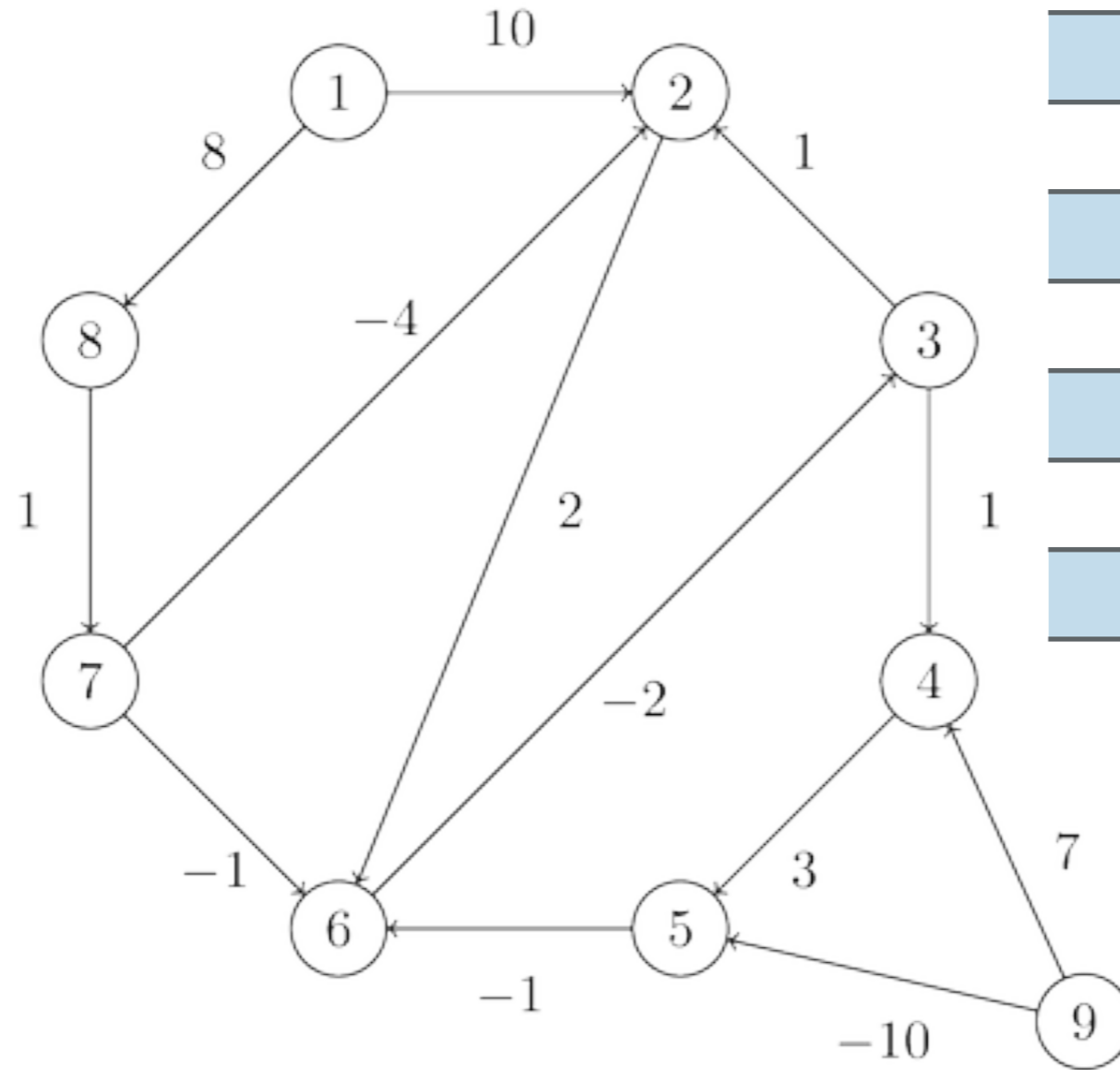
v	distTo[]	edgeTo[]
1	0	-
2	5	7->2
3	10	6->3
4	11	3->4
5	14	4->5
6	12	2->6
7	9	8->7
8	8	1->8
9	Inf	null

<http://rosalind.info/problems/bf/>



# Practice Time - Pass 1

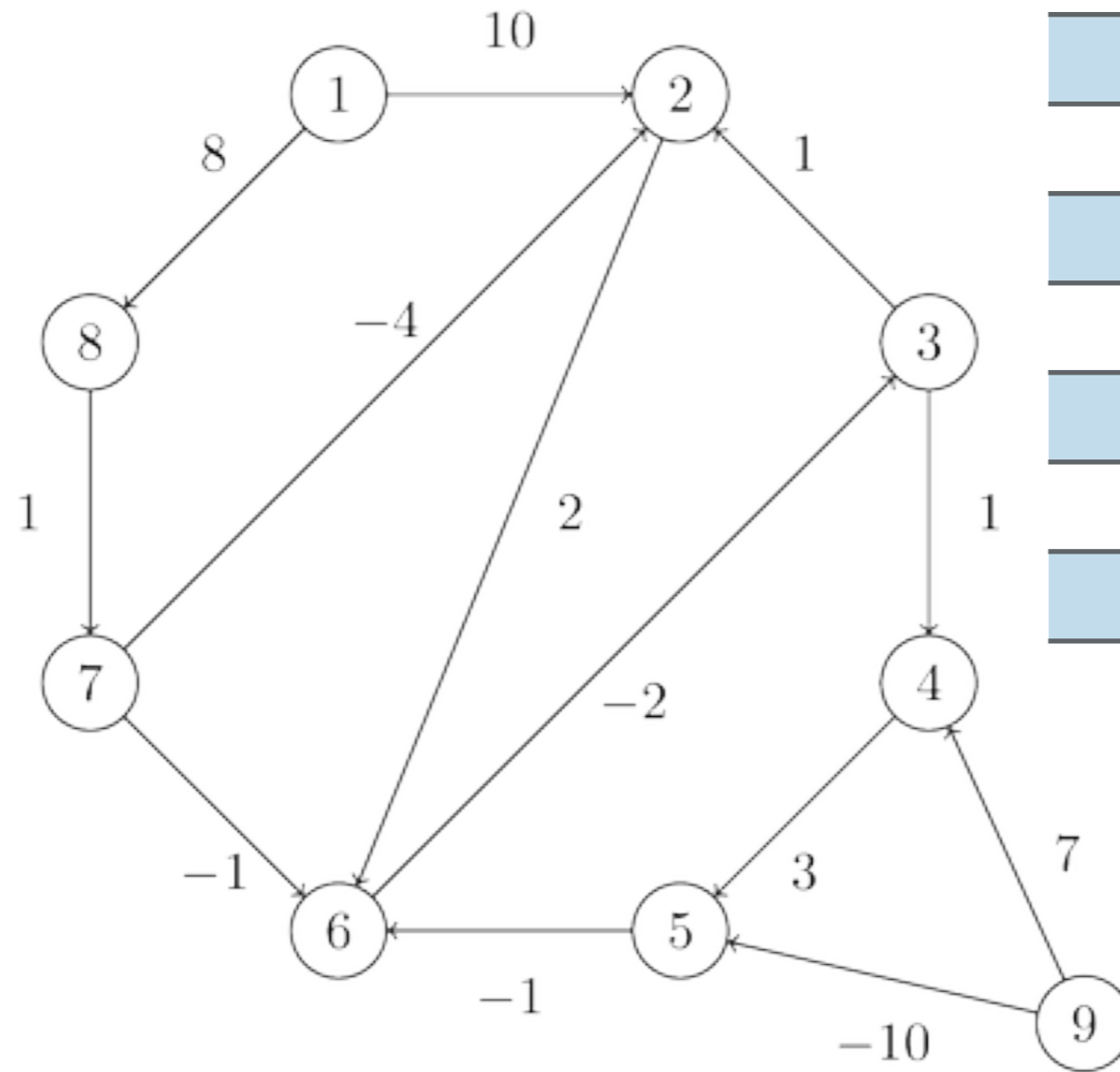
1 2 10  
 3 2 1  
 3 4 1  
 4 5 3  
 5 6 -1  
 7 6 -1  
 8 7 1  
 1 8 8  
 7 2 -4  
 2 6 2  
 6 3 -2  
 9 5 -10  
 9 4 7



v	distTo[]	edgeTo[]
1	0	-
2	5	7->2
3	10	6->3
4	11	3->4
5	14	4->5
6	7	2->6
7	9	8->7
8	8	1->8
9	Inf	null

# Practice Time - Pass 1

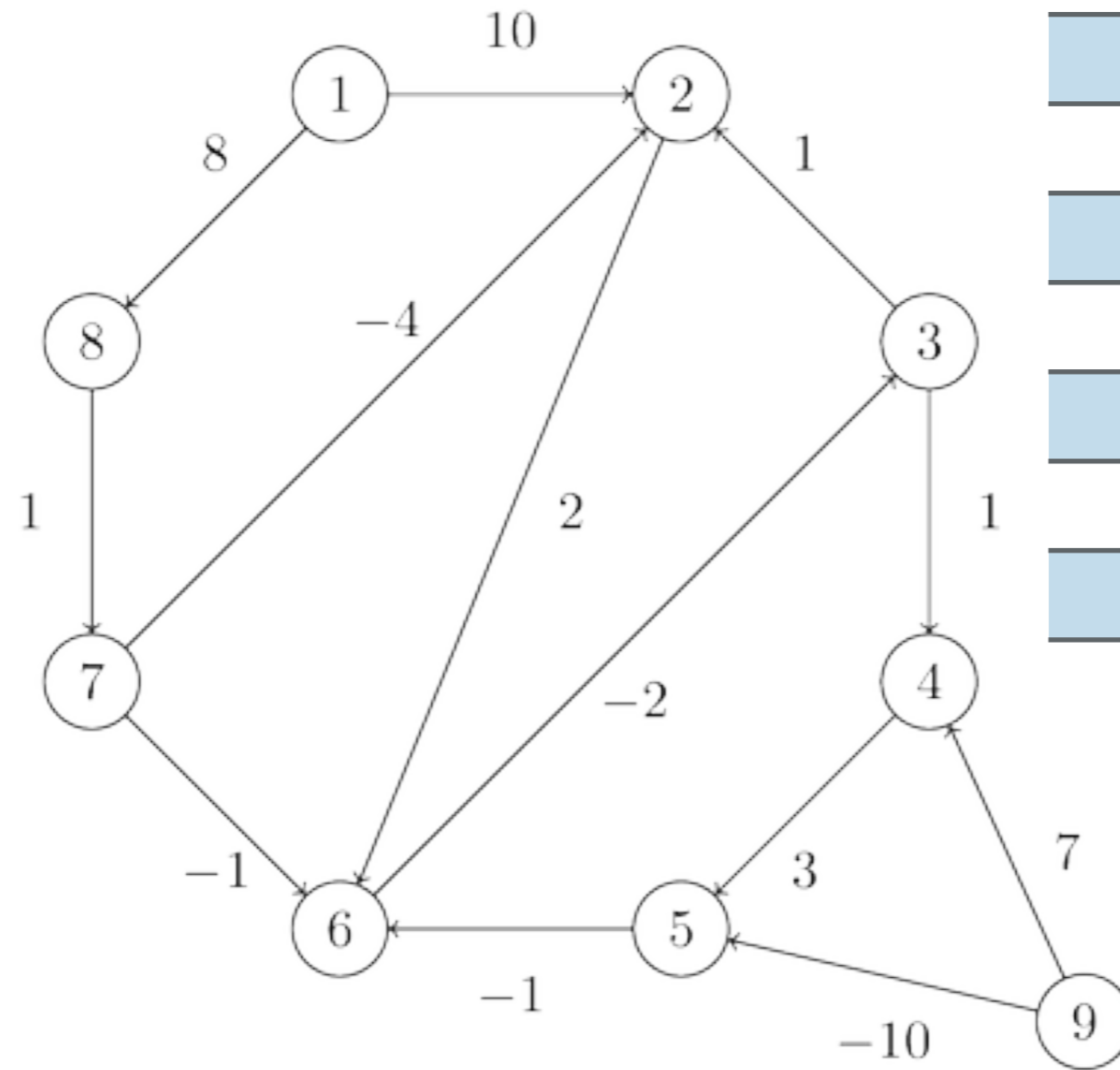
1 2 10  
 3 2 1  
 3 4 1  
 4 5 3  
 5 6 -1  
 7 6 -1  
 8 7 1  
 1 8 8  
 7 2 -4  
 2 6 2  
 6 3 -2  
 9 5 -10  
 9 4 7



v	distTo[]	edgeTo[]
1	0	-
2	5	7->2
3	5	6->3
4	11	3->4
5	14	4->5
6	7	2->6
7	9	8->7
8	8	1->8
9	Inf	null

# Practice Time - Pass 1

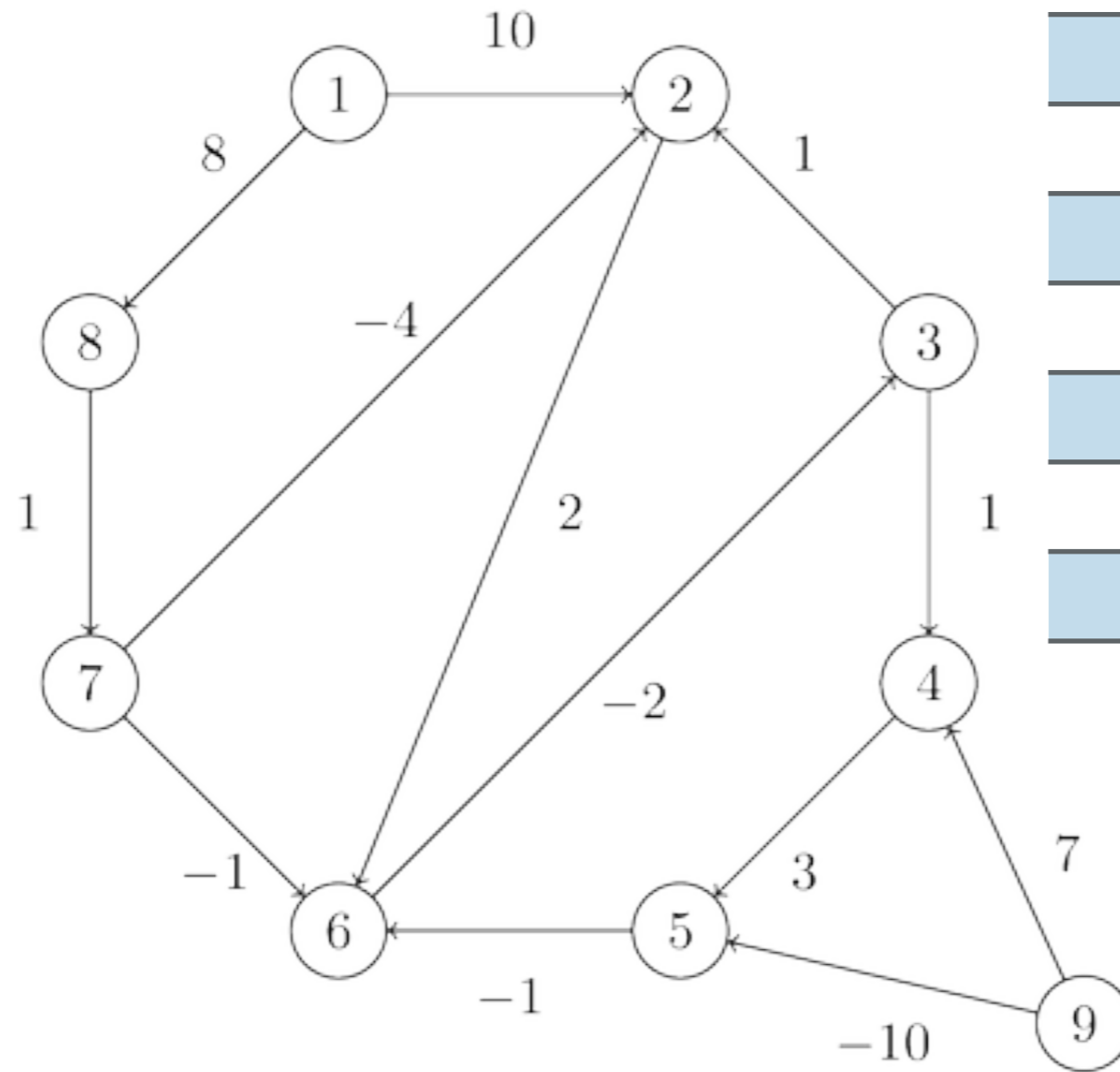
1 2 10  
 3 2 1  
 3 4 1  
 4 5 3  
 5 6 -1  
 7 6 -1  
 8 7 1  
 1 8 8  
 7 2 -4  
 2 6 2  
 6 3 -2  
 9 5 -10  
 9 4 7



v	distTo[]	edgeTo[]
1	0	-
2	5	7->2
3	5	6->3
4	11	3->4
5	14	4->5
6	7	2->6
7	9	8->7
8	8	1->8
9	Inf	null

# Practice Time - Pass 1

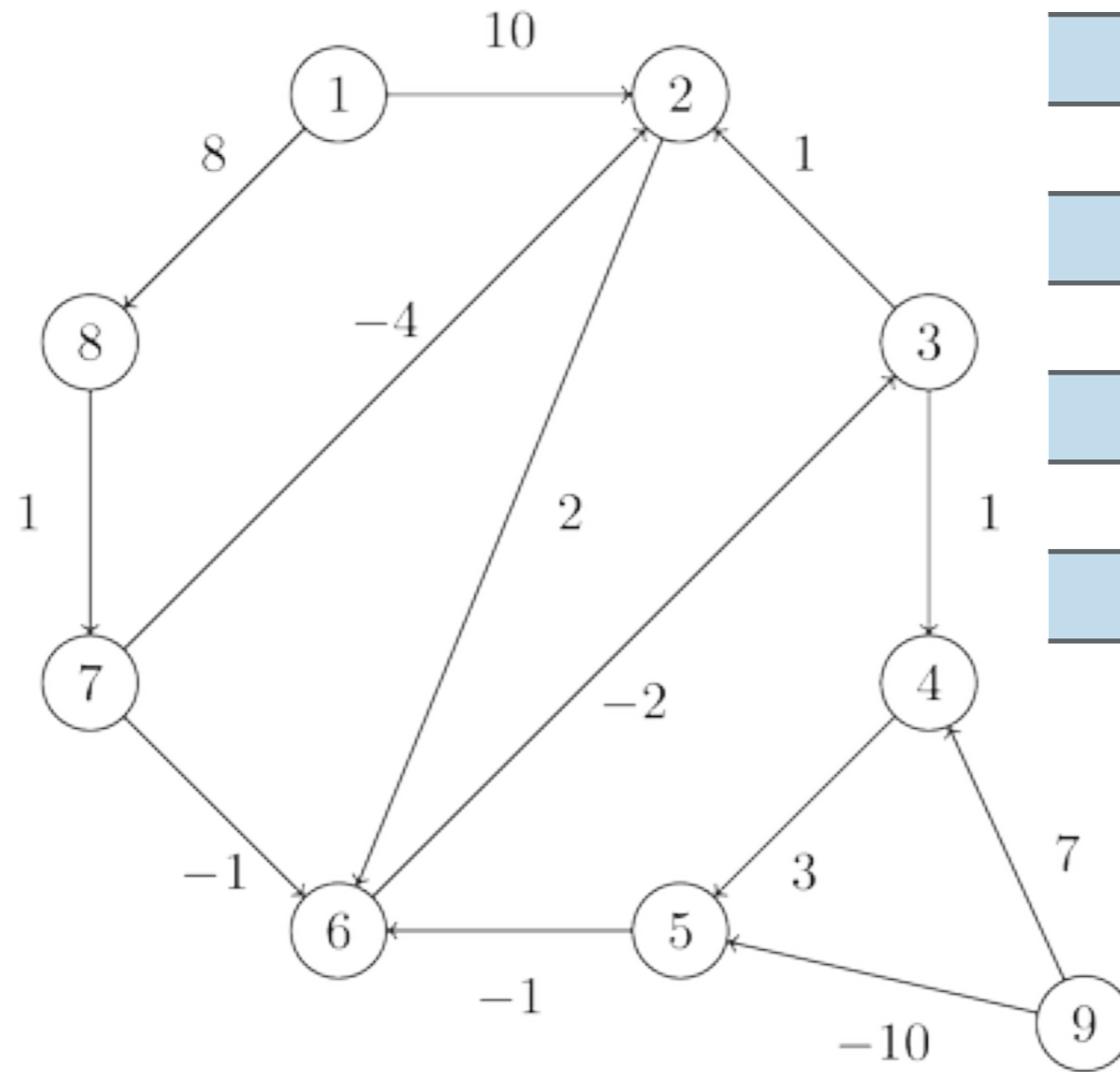
1 2 10  
 3 2 1  
 3 4 1  
 4 5 3  
 5 6 -1  
 7 6 -1  
 8 7 1  
 1 8 8  
 7 2 -4  
 2 6 2  
 6 3 -2  
 9 5 -10  
 9 4 7



v	distTo[]	edgeTo[]
1	0	-
2	5	7->2
3	5	6->3
4	11	3->4
5	14	4->5
6	7	2->6
7	9	8->7
8	8	1->8
9	Inf	null

# Practice Time - Pass 2

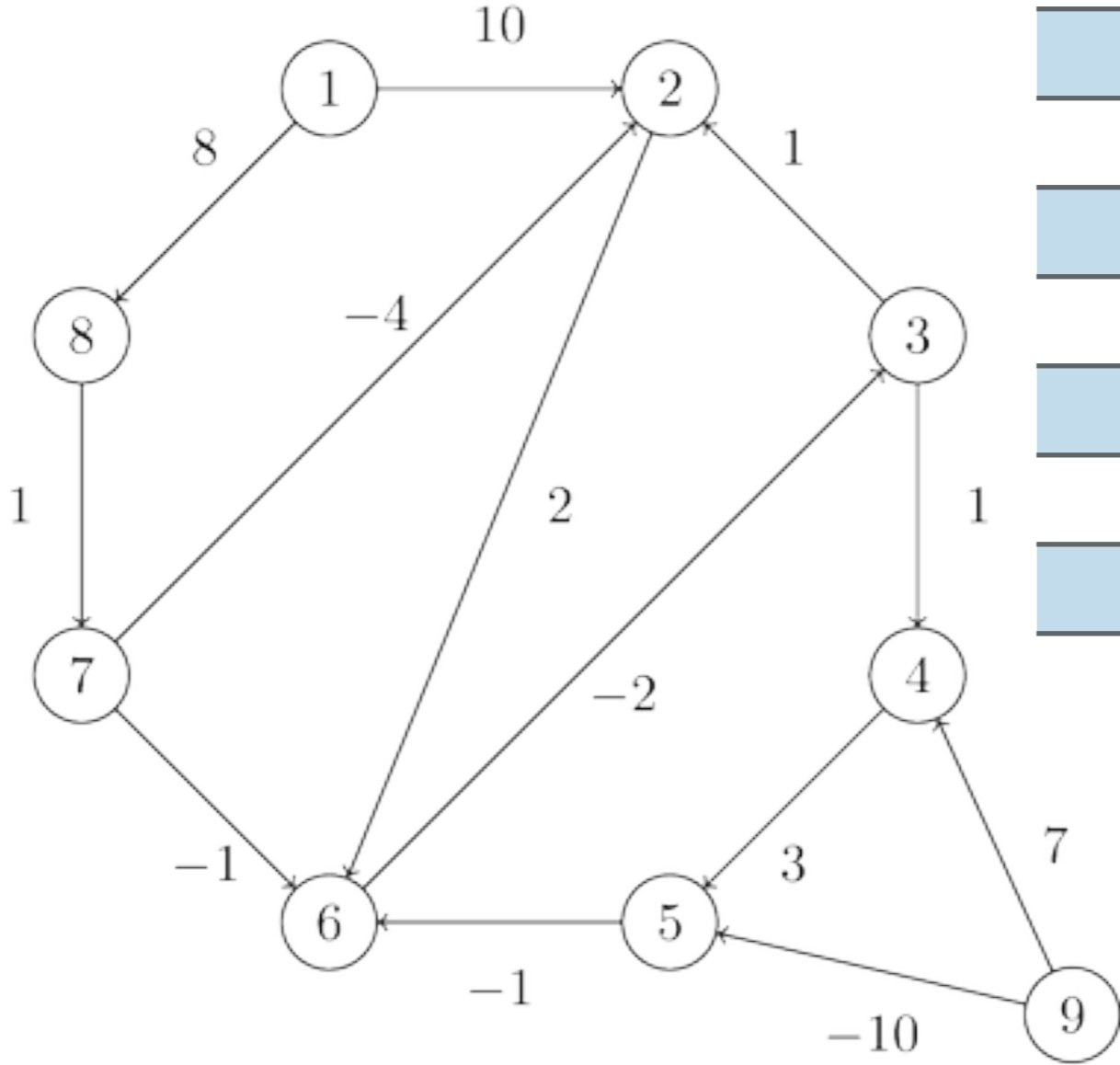
1 2 10  
 3 2 1  
 3 4 1  
 4 5 3  
 5 6 -1  
 7 6 -1  
 8 7 1  
 1 8 8  
 7 2 -4  
 2 6 2  
 6 3 -2  
 9 5 -10  
 9 4 7



v	distTo[]	edgeTo[]
1	0	-
2	5	7->2
3	5	6->3
4	11	3->4
5	14	4->5
6	7	2->6
7	9	8->7
8	8	1->8
9	Inf	null

# Practice Time - Pass 2

1 2 10  
 3 2 1  
 3 4 1  
 4 5 3  
 5 6 -1  
 7 6 -1  
 8 7 1  
 1 8 8  
 7 2 -4  
 2 6 2  
 6 3 -2  
 9 5 -10  
 9 4 7

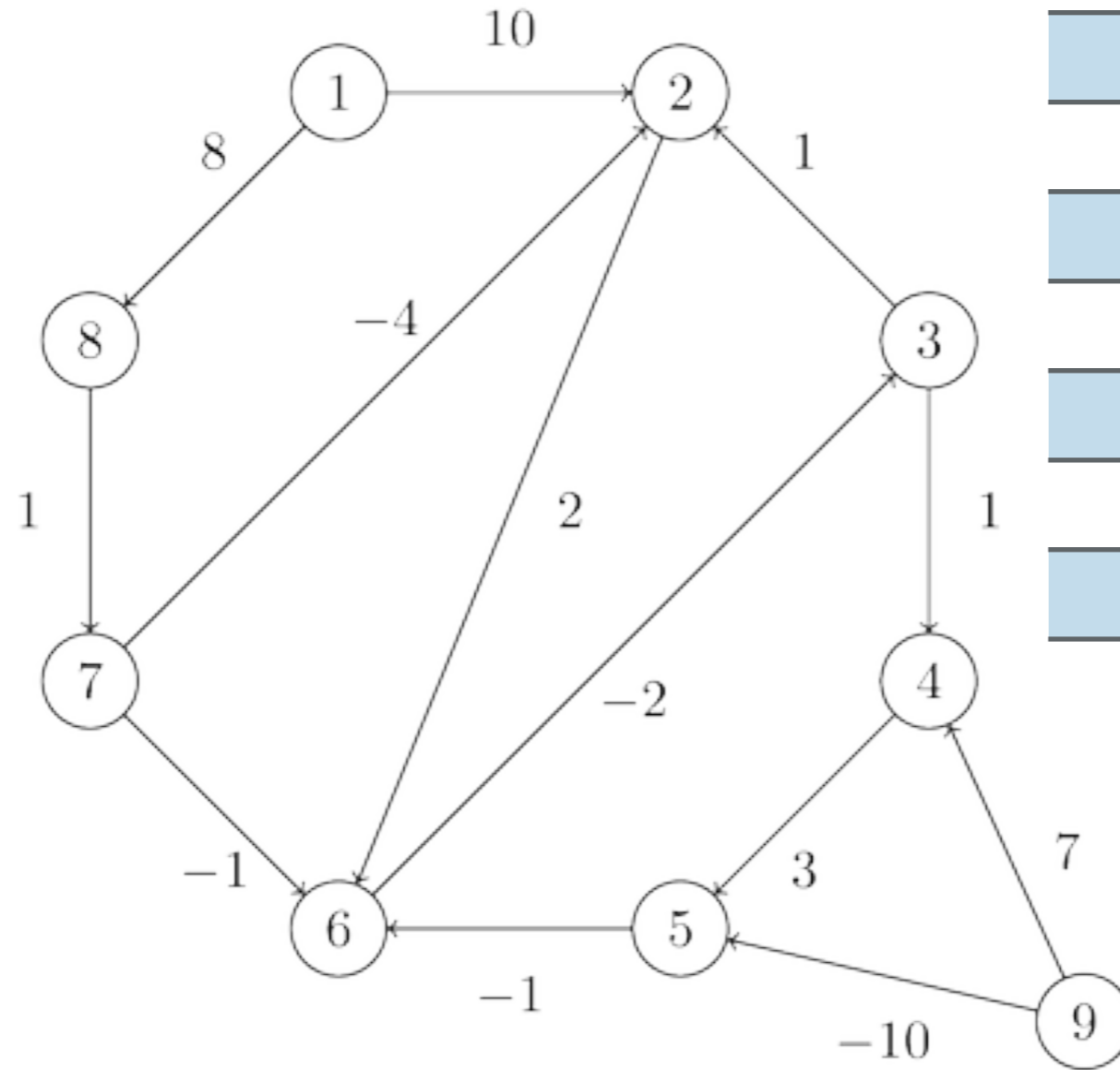


v	distTo[]	edgeTo[]
1	0	-
2	5	7->2
3	5	6->3
4	11	3->4
5	14	4->5
6	7	2->6
7	9	8->7
8	8	1->8
9	Inf	null

<http://rosalind.info/problems/bf/>

# Practice Time - Pass 2

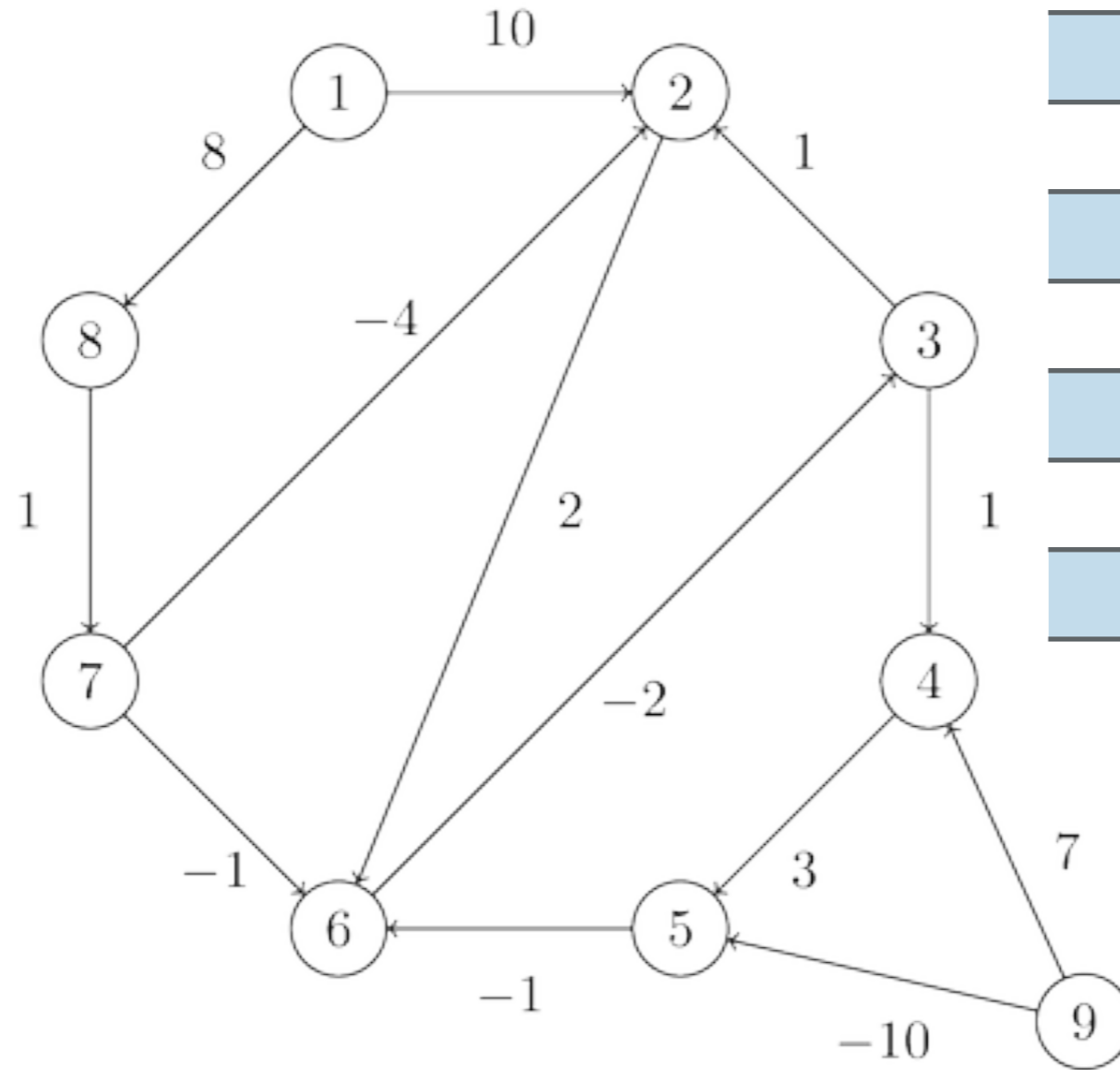
1 2 10  
 3 2 1  
 3 4 1  
 4 5 3  
 5 6 -1  
 7 6 -1  
 8 7 1  
 1 8 8  
 7 2 -4  
 2 6 2  
 6 3 -2  
 9 5 -10  
 9 4 7



v	distTo[]	edgeTo[]
1	0	-
2	5	7->2
3	5	6->3
4	6	3->4
5	14	4->5
6	7	2->6
7	9	8->7
8	8	1->8
9	Inf	null

# Practice Time - Pass 2

1 2 10  
 3 2 1  
 3 4 1  
 4 5 3  
 5 6 -1  
 7 6 -1  
 8 7 1  
 1 8 8  
 7 2 -4  
 2 6 2  
 6 3 -2  
 9 5 -10  
 9 4 7

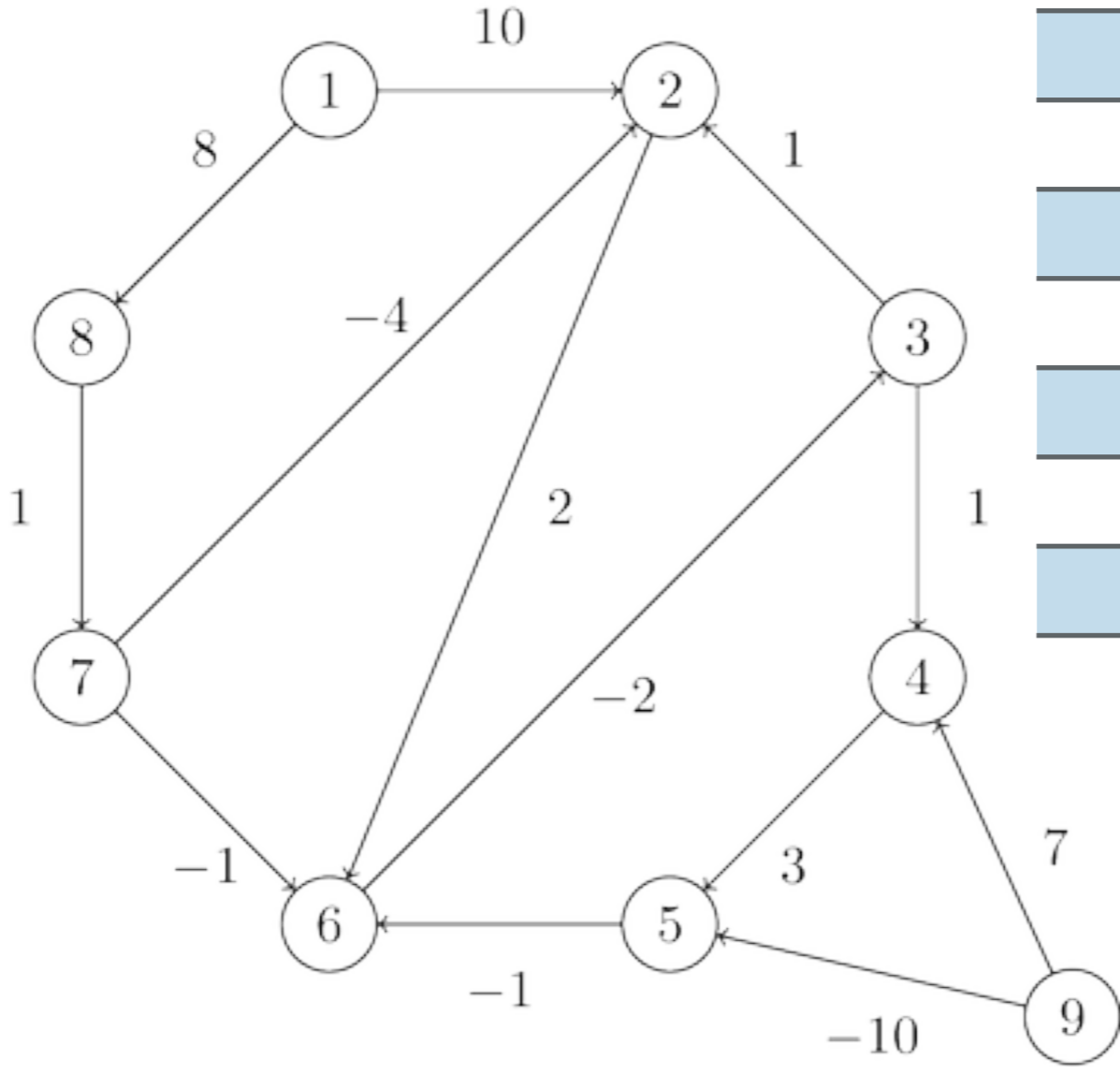


v	distTo[]	edgeTo[]
1	0	-
2	5	7->2
3	5	6->3
4	6	3->4
5	9	4->5
6	7	2->6
7	9	8->7
8	8	1->8
9	Inf	null



# Practice Time - Pass 2

1 2 10  
 3 2 1  
 3 4 1  
 4 5 3  
 5 6 -1  
 7 6 -1  
 8 7 1  
 1 8 8  
 7 2 -4  
 2 6 2  
 6 3 -2  
 9 5 -10  
 9 4 7

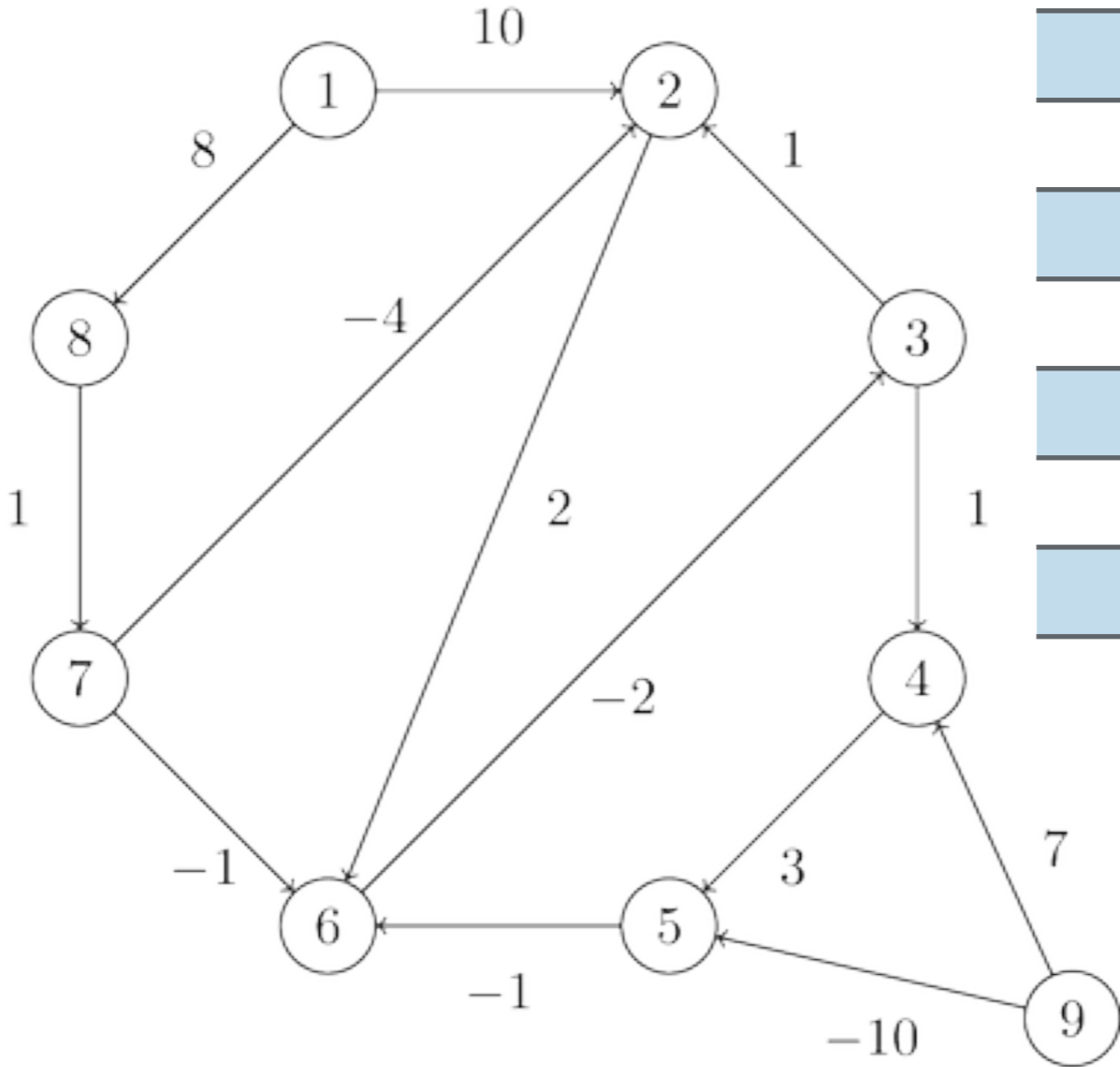


v	distTo[]	edgeTo[]
1	0	-
2	5	7->2
3	5	6->3
4	6	3->4
5	9	4->5
6	7	2->6
7	9	8->7
8	8	1->8
9	Inf	null

<http://rosalind.info/problems/bf/>

# Practice Time - Pass 2

1 2 10  
 3 2 1  
 3 4 1  
 4 5 3  
 5 6 -1  
 7 6 -1  
 8 7 1  
 1 8 8  
 7 2 -4  
 2 6 2  
 6 3 -2  
 9 5 -10  
 9 4 7

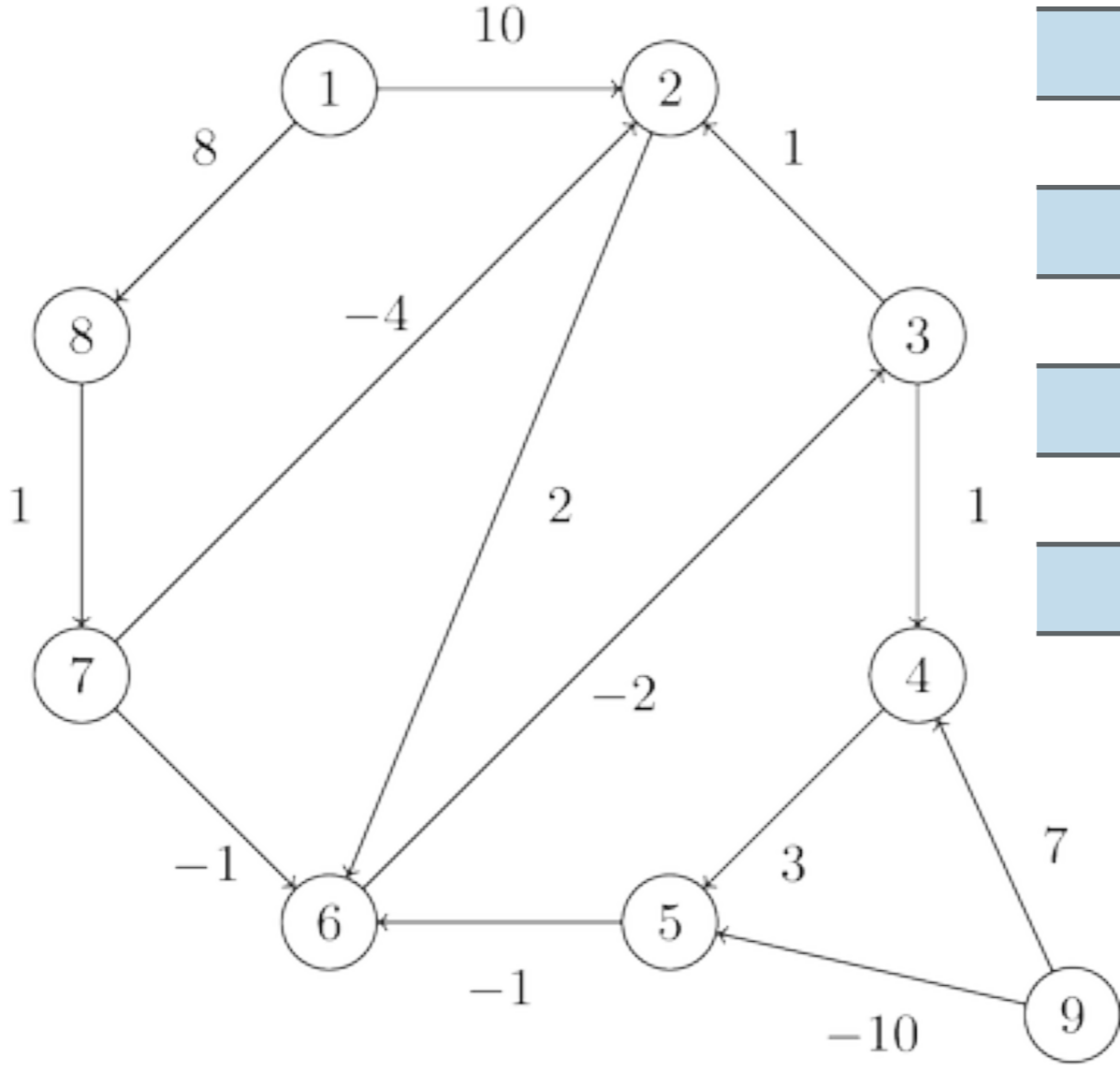


v	distTo[]	edgeTo[]
1	0	-
2	5	7->2
3	5	6->3
4	6	3->4
5	9	4->5
6	7	2->6
7	9	8->7
8	8	1->8
9	Inf	null

<http://rosalind.info/problems/bf/>

# Practice Time - Pass 2

1 2 10  
 3 2 1  
 3 4 1  
 4 5 3  
 5 6 -1  
 7 6 -1  
 8 7 1  
 1 8 8  
 7 2 -4  
 2 6 2  
 6 3 -2  
 9 5 -10  
 9 4 7

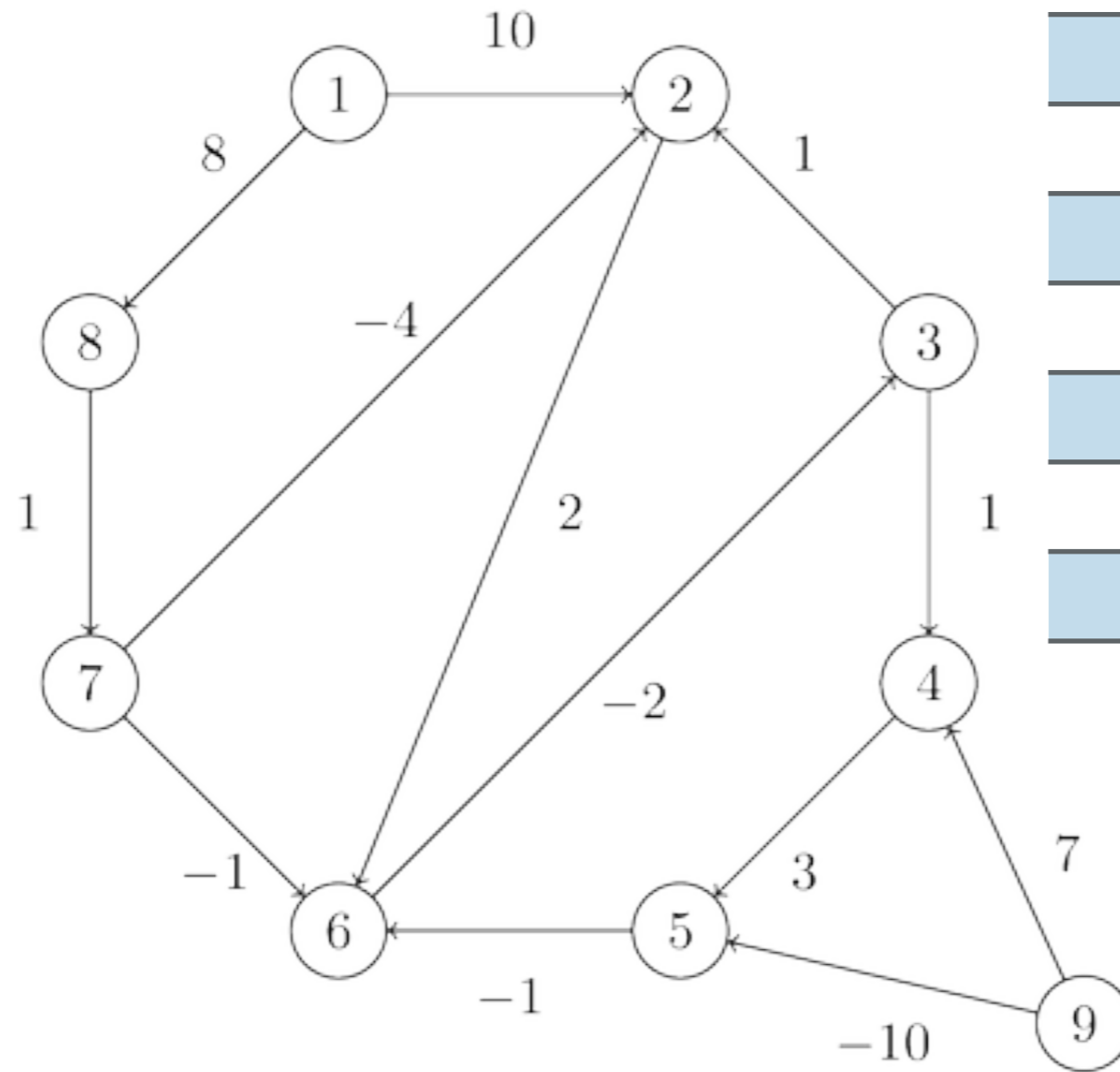


v	distTo[]	edgeTo[]
1	0	-
2	5	7->2
3	5	6->3
4	6	3->4
5	9	4->5
6	7	2->6
7	9	8->7
8	8	1->8
9	Inf	null

<http://rosalind.info/problems/bf/>

# Practice Time - Pass 2

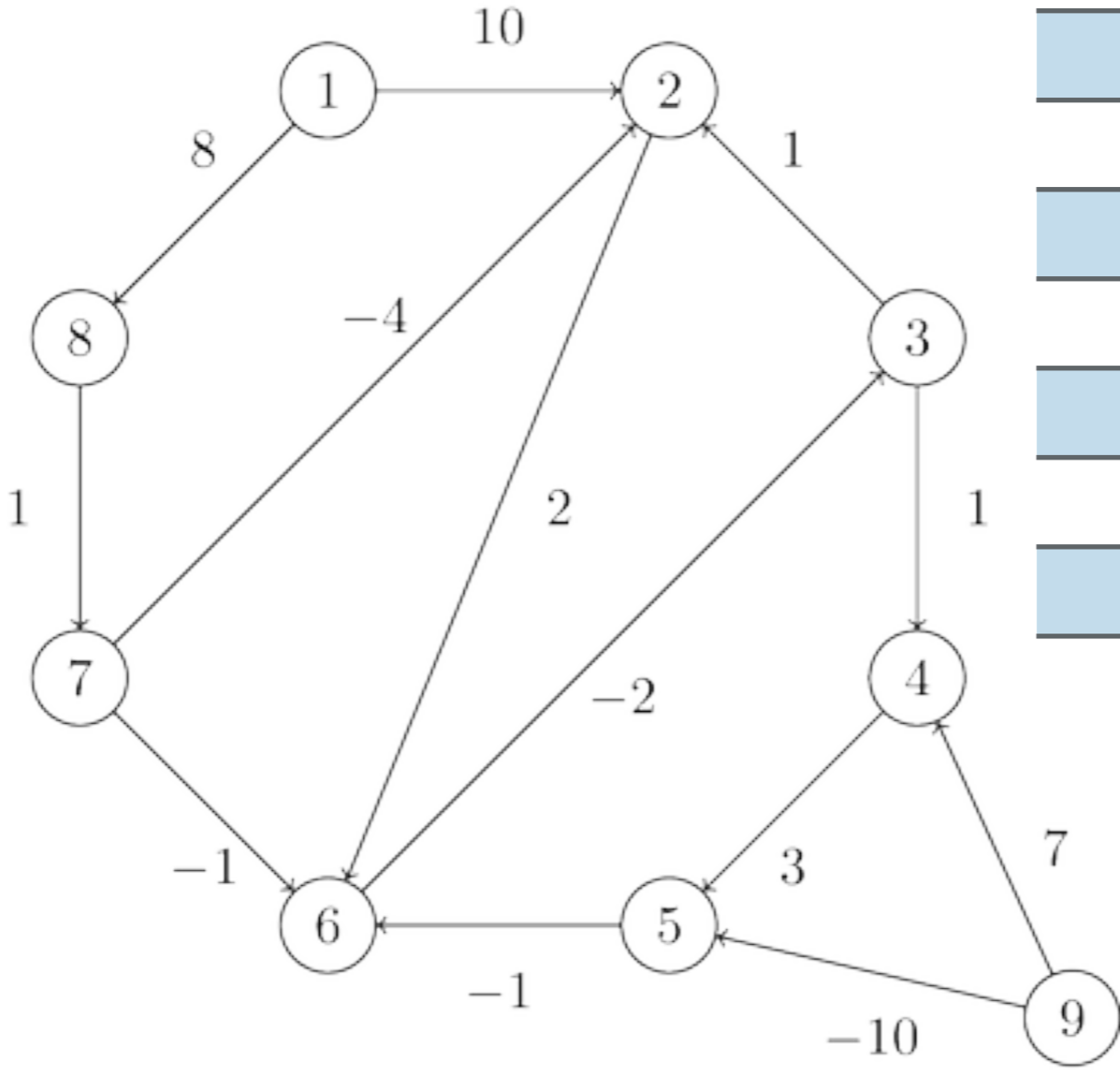
1 2 10  
 3 2 1  
 3 4 1  
 4 5 3  
 5 6 -1  
 7 6 -1  
 8 7 1  
 1 8 8  
 7 2 -4  
 2 6 2  
 6 3 -2  
 9 5 -10  
 9 4 7



v	distTo[]	edgeTo[]
1	0	-
2	5	7->2
3	5	6->3
4	6	3->4
5	9	4->5
6	7	2->6
7	9	8->7
8	8	1->8
9	Inf	null

# Practice Time - Pass 2

1 2 10  
 3 2 1  
 3 4 1  
 4 5 3  
 5 6 -1  
 7 6 -1  
 8 7 1  
 1 8 8  
**7 2 -4**  
 2 6 2  
 6 3 -2  
 9 5 -10  
 9 4 7

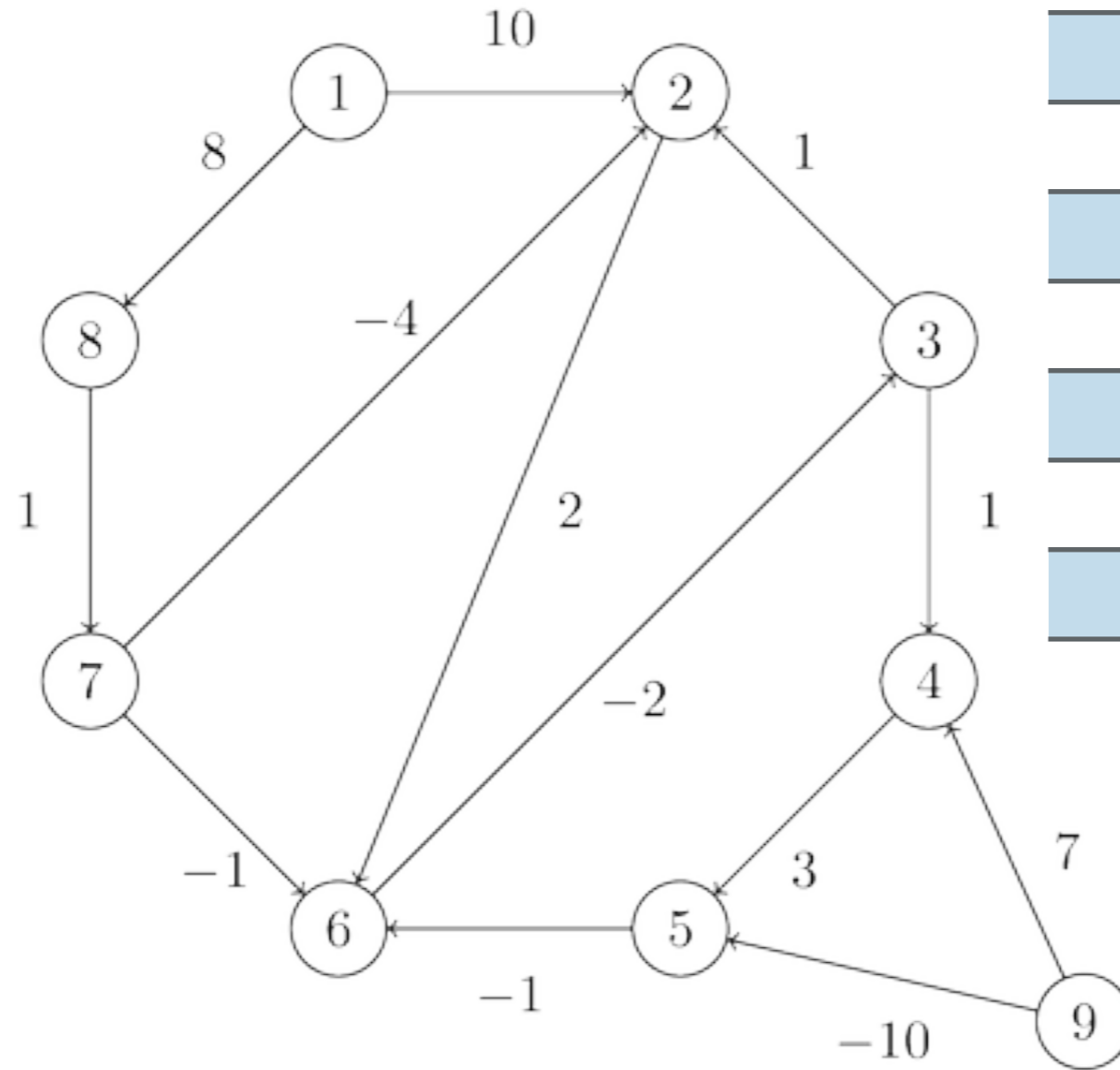


v	distTo[]	edgeTo[]
1	0	-
2	5	7->2
3	5	6->3
4	6	3->4
5	9	4->5
6	7	2->6
7	9	8->7
8	8	1->8
9	Inf	null

<http://rosalind.info/problems/bf/>

# Practice Time - Pass 2

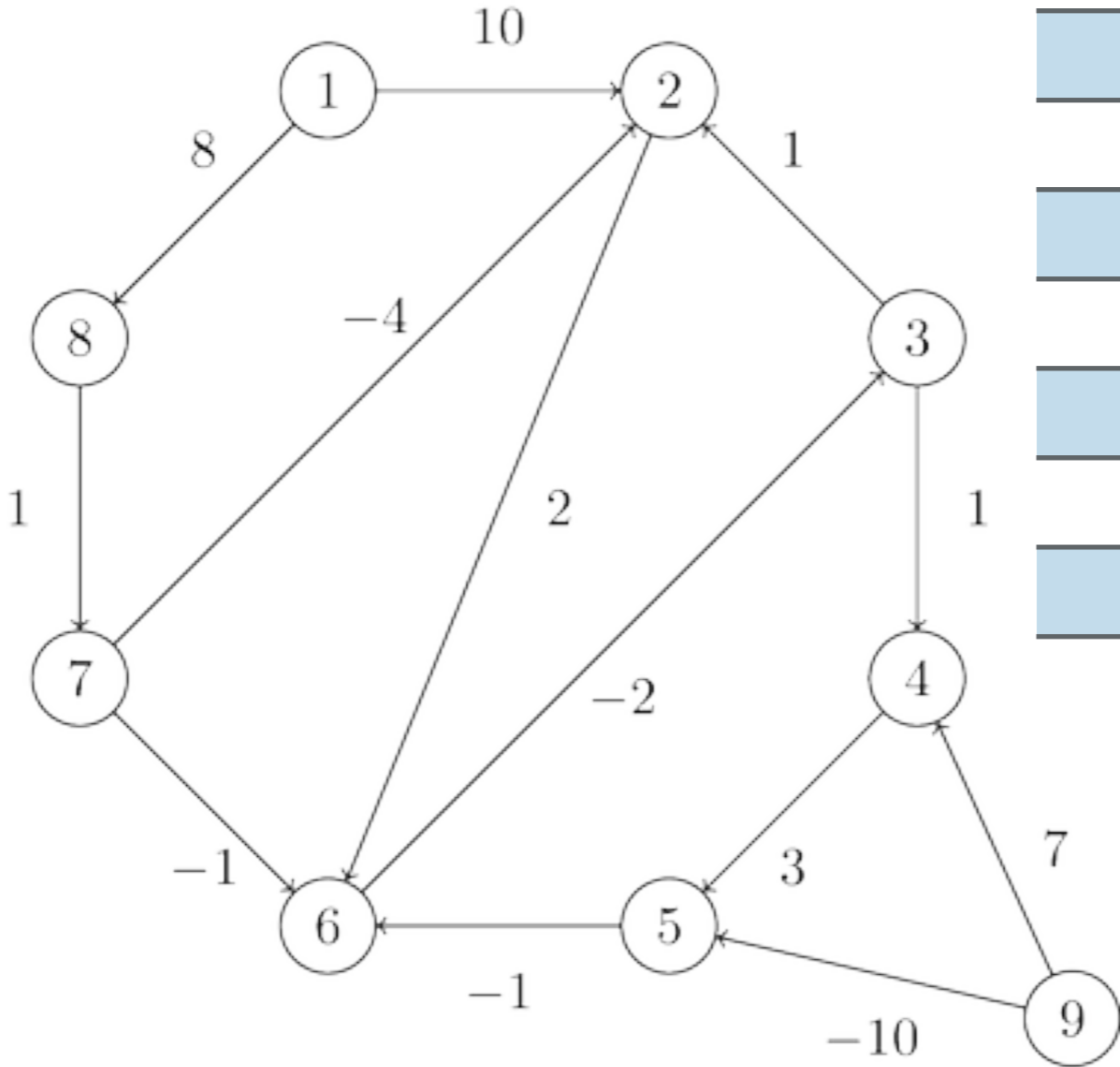
1 2 10  
 3 2 1  
 3 4 1  
 4 5 3  
 5 6 -1  
 7 6 -1  
 8 7 1  
 1 8 8  
 7 2 -4  
 2 6 2  
 6 3 -2  
 9 5 -10  
 9 4 7



v	distTo[]	edgeTo[]
1	0	-
2	5	7->2
3	5	6->3
4	6	3->4
5	9	4->5
6	7	2->6
7	9	8->7
8	8	1->8
9	Inf	null

# Practice Time - Pass 2

1 2 10  
 3 2 1  
 3 4 1  
 4 5 3  
 5 6 -1  
 7 6 -1  
 8 7 1  
 1 8 8  
 7 2 -4  
 2 6 2  
 6 3 -2  
 9 5 -10  
 9 4 7

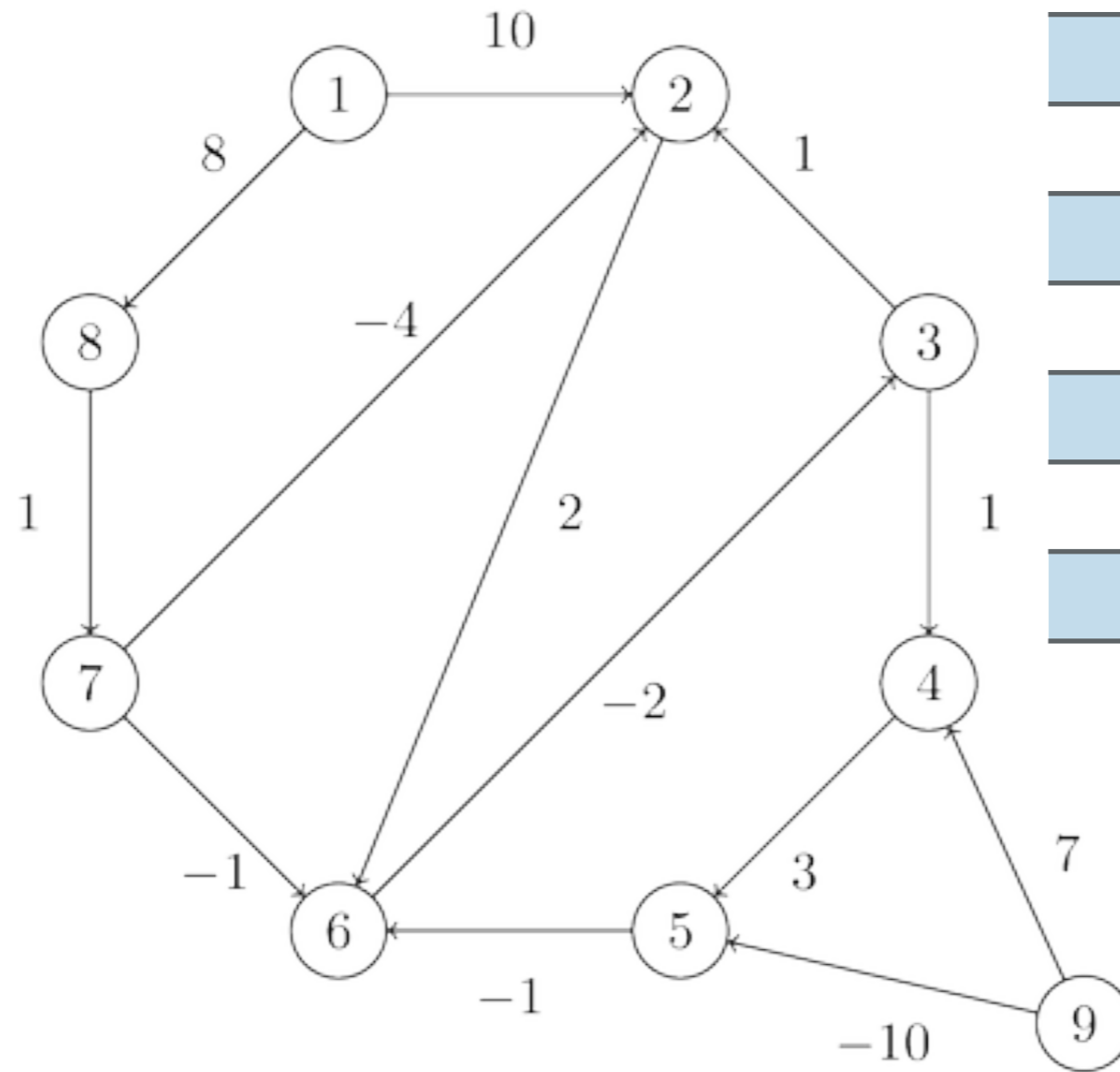


v	distTo[]	edgeTo[]
1	0	-
2	5	7->2
3	5	6->3
4	6	3->4
5	9	4->5
6	7	2->6
7	9	8->7
8	8	1->8
9	Inf	null

<http://rosalind.info/problems/bf/>

# Practice Time - Pass 2

1 2 10  
 3 2 1  
 3 4 1  
 4 5 3  
 5 6 -1  
 7 6 -1  
 8 7 1  
 1 8 8  
 7 2 -4  
 2 6 2  
 6 3 -2  
 9 5 -10  
 9 4 7

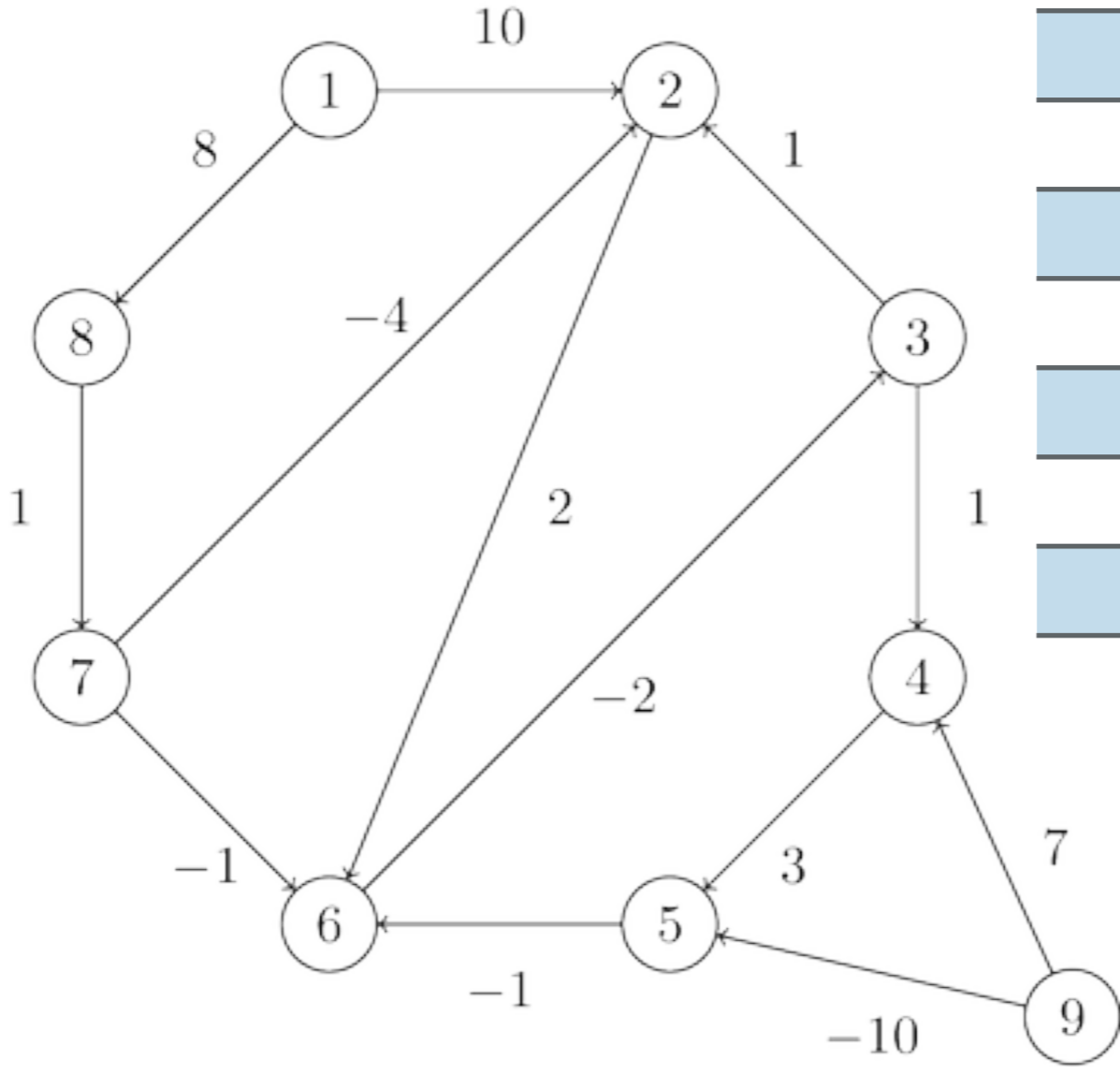


v	distTo[]	edgeTo[]
1	0	-
2	5	7->2
3	5	6->3
4	6	3->4
5	9	4->5
6	7	2->6
7	9	8->7
8	8	1->8
9	Inf	null



# Practice Time - Pass 2

1 2 10  
 3 2 1  
 3 4 1  
 4 5 3  
 5 6 -1  
 7 6 -1  
 8 7 1  
 1 8 8  
 7 2 -4  
 2 6 2  
 6 3 -2  
 9 5 -10  
 9 4 7

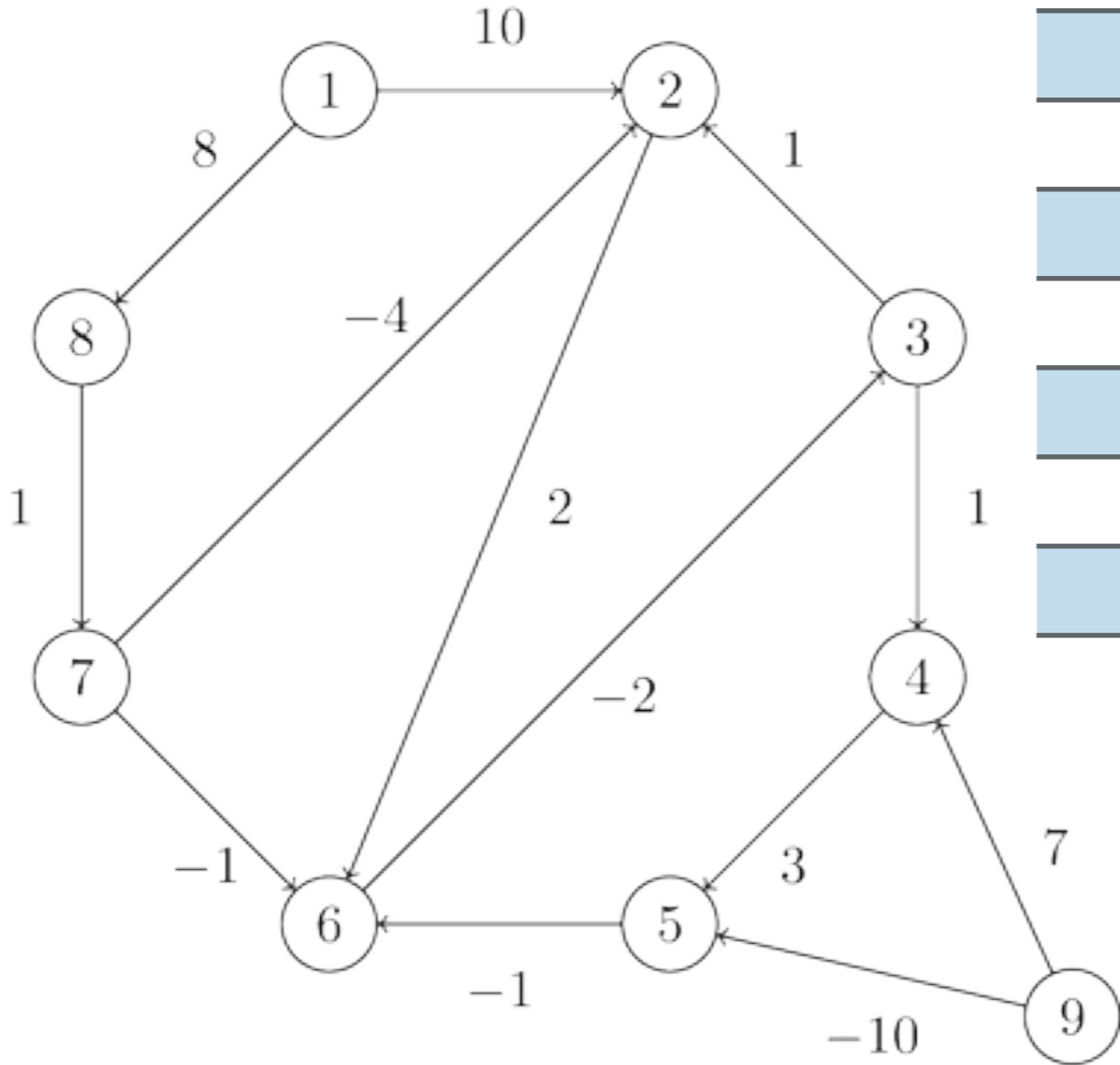


v	distTo[]	edgeTo[]
1	0	-
2	5	7->2
3	5	6->3
4	6	3->4
5	9	4->5
6	7	2->6
7	9	8->7
8	8	1->8
9	Inf	null

<http://rosalind.info/problems/bf/>

# Practice Time - No changes in passes 3-8

1 2 10  
 3 2 1  
 3 4 1  
 4 5 3  
 5 6 -1  
 7 6 -1  
 8 7 1  
 1 8 8  
 7 2 -4  
 2 6 2  
 6 3 -2  
 9 5 -10  
 9 4 7



v	distTo[]	edgeTo[]
1	0	-
2	5	7->2
3	5	6->3
4	6	3->4
5	9	4->5
6	7	2->6
7	9	8->7
8	8	1->8
9	Inf	null

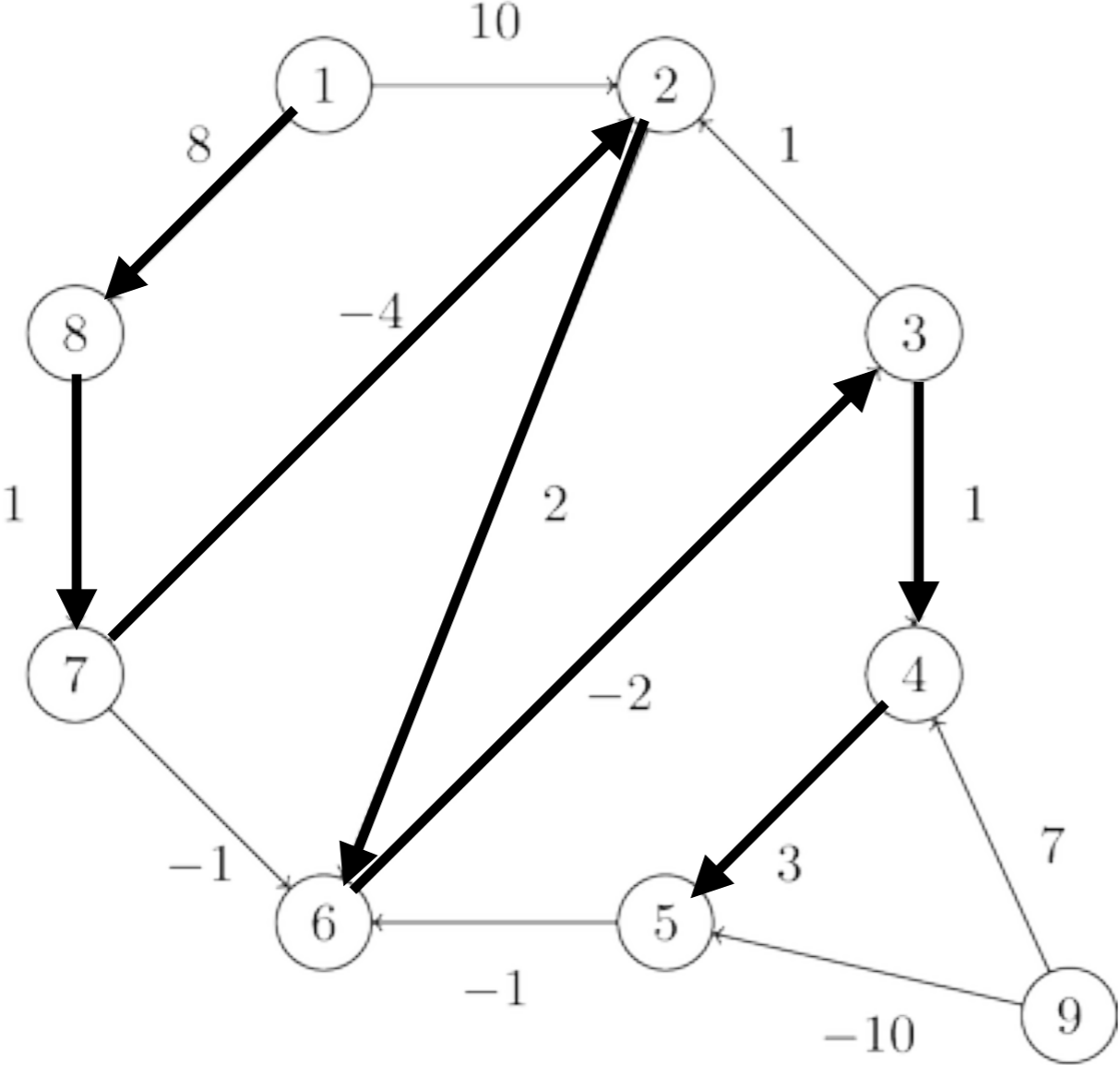
<http://rosalind.info/problems/bf/>

# Answer

```

1 2 10
3 2 1
3 4 1
4 5 3
5 6 -1
7 6 -1
8 7 1
1 8 8
7 2 -4
2 6 2
6 3 -2
9 5 -10
9 4 7

```



v	distTo[]	edgeTo[]
1	0	-
2	5	7->2
3	5	6->3
4	6	3->4
5	9	4->5
6	7	2->6
7	9	8->7
8	8	1->8
9	Inf	null

<http://rosalind.info/problems/bf/>

## Lecture 27: Shortest Paths

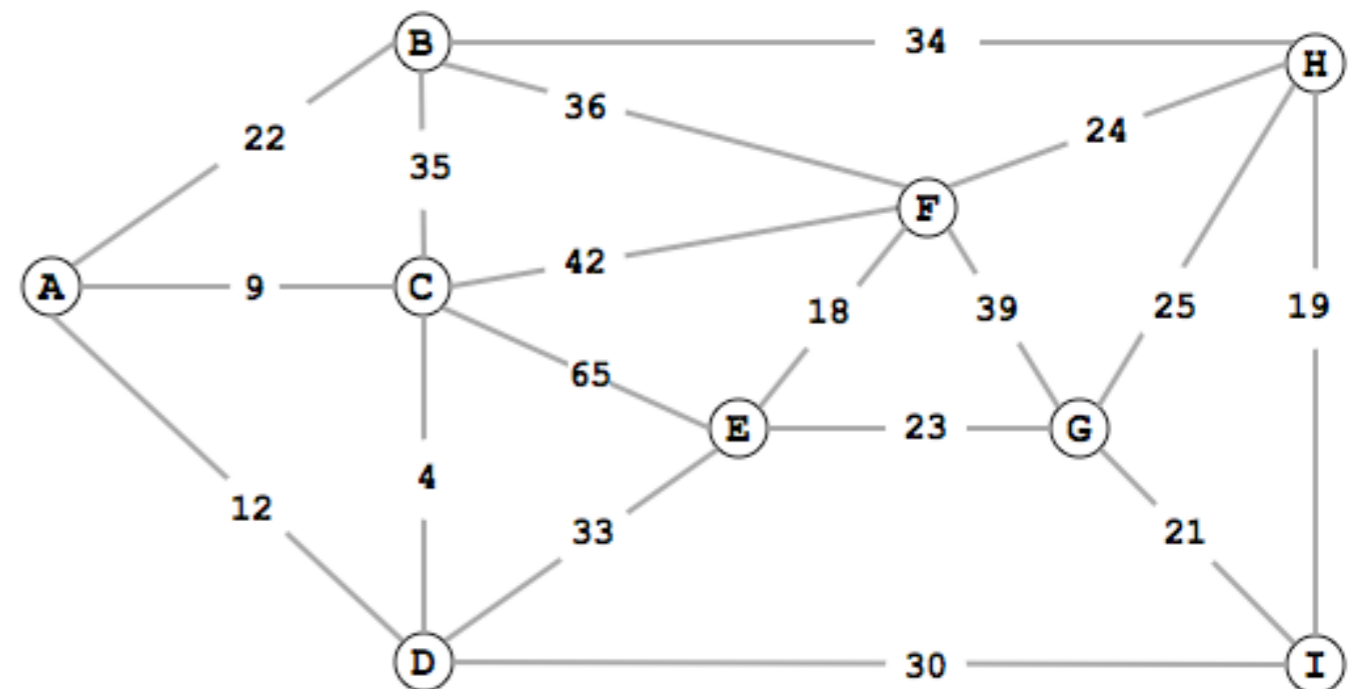
- ▶ Introduction to Shortest Paths
- ▶ API
- ▶ Properties
- ▶ Dijkstra's Algorithm
- ▶ Belman-Ford Algorithm

## Readings:

- ▶ Textbook: Chapter 4.4 (Pages 638-676)
- ▶ Website:
  - ▶ <https://algs4.cs.princeton.edu/44sp/>

## Practice Problems:

Run Dijkstra's algorithm on the graph on the right with A being the starting vertex.



## Practice Problems:

Run the Bellman-Ford algorithm on this directed graph using vertex  $z$  as the source. In each pass show the values  $d$  and  $\pi$ . In the graph,  $V = \{s, t, v, x, z\}$  and the weighted, directed edges are  $E = \{(s, t, 6), (s, v, 7), (t, v, 8), (t, z, -4), (t, x, 5), (v, x, -3), (v, z, 9), (x, t, -2), (z, s, 2), (z, x, 4)\}$ .

<https://www.chegg.com/homework-help/questions-and-answers/run-bellman-ford-algorithm-following-directed-graph-using-vertex-z-source-pass-show-values-q17182493>

