

CS062

DATA STRUCTURES AND ADVANCED PROGRAMMING

12: Insertion Sort & Mergesort



Alexandra Papoutsaki
she/her/hers



Tom Yeh
he/him/his

Lecture 12: Insertion Sort & Mergesort

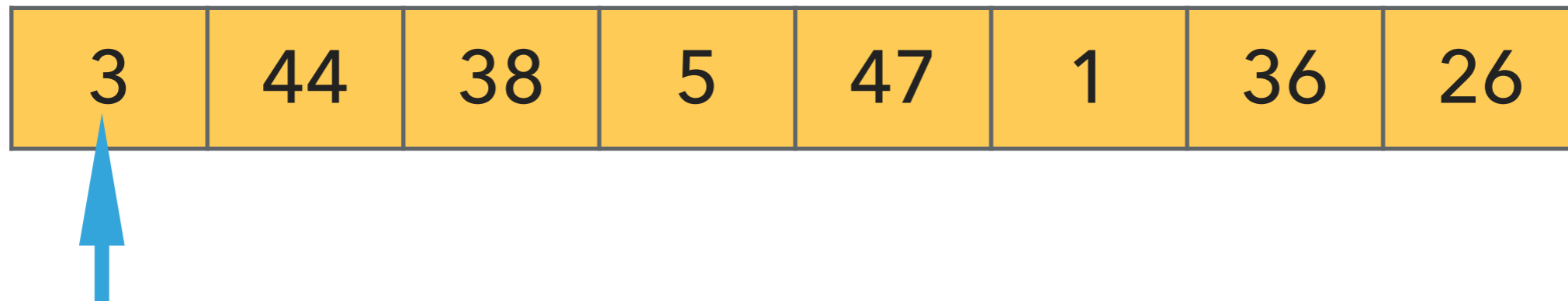
- ▶ Insertion Sort
- ▶ Mergesort

Insertion sort

3	44	38	5	47	1	36	26
---	----	----	---	----	---	----	----

- ▶ Keep a *partially sorted subarray* on the left and an *unsorted subarray* on the right
- ▶ Repeat:
 - ▶ Examine the next element in the unsorted subarray.
 - ▶ Insert this element by exchanging with every entry to the left that is greater.
 - ▶ Move subarray boundaries one element to the right.

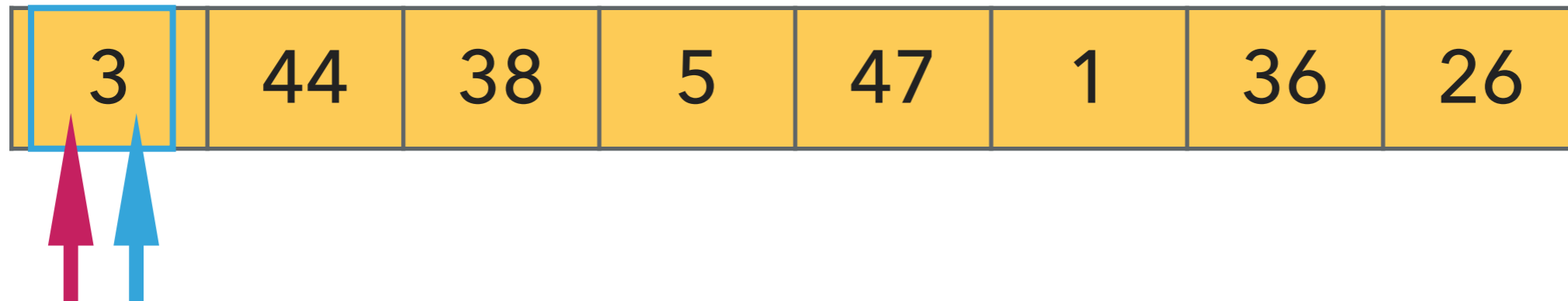
Insertion sort



▶ Repeat:

- ▶ Examine the next element in the unsorted subarray.
- ▶ Insert this element by exchanging with every entry to the left that is greater.
- ▶ Move subarray boundaries one element to the right.

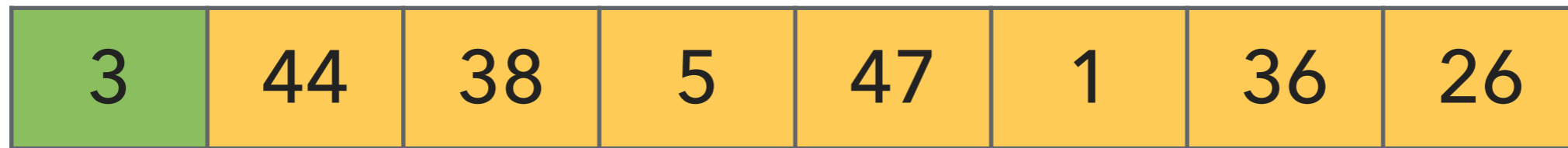
Insertion sort



▶ Repeat:

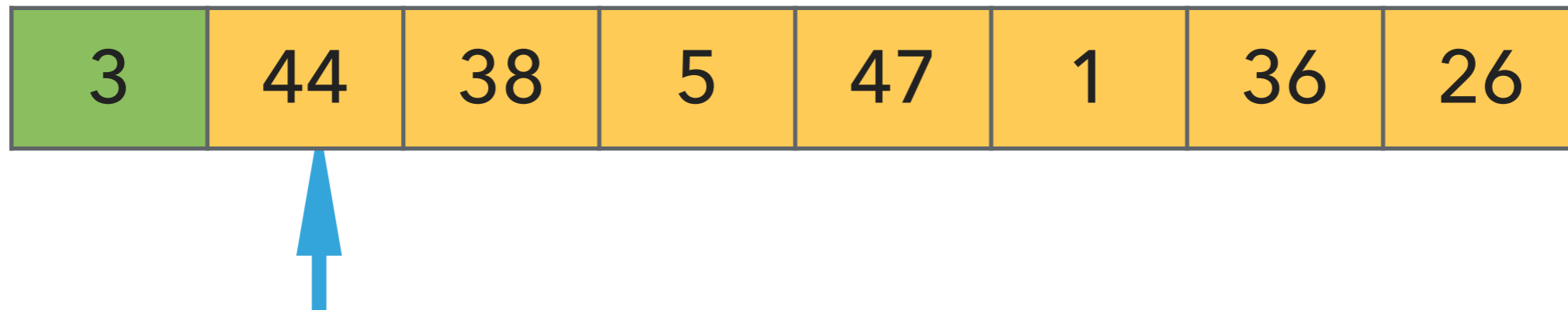
- ▶ Examine the next element in the unsorted subarray.
- ▶ Exchange this element with every entry to the left that is greater.
- ▶ Move subarray boundaries one element to the right.

Insertion sort



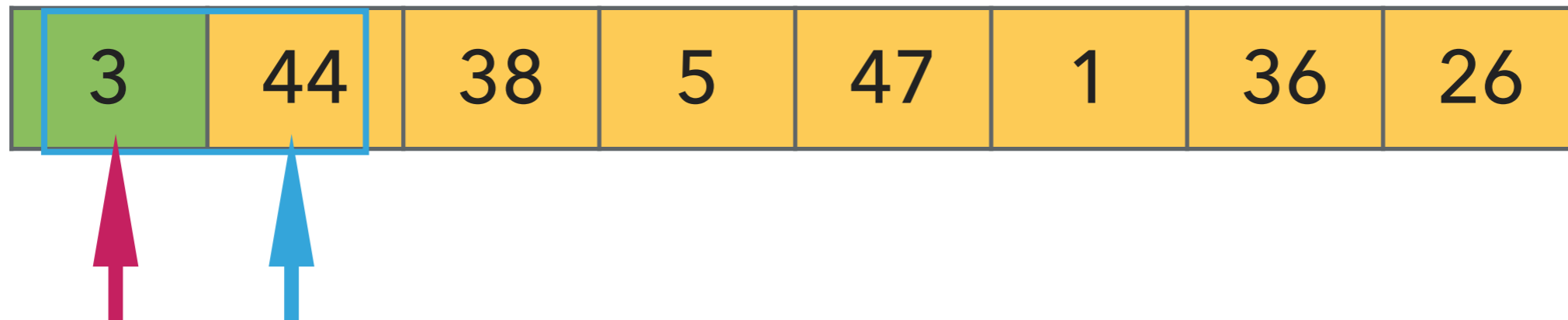
- ▶ Repeat:
 - ▶ Examine the next element in the unsorted subarray.
 - ▶ Exchange this element with every entry to the left that is greater.
 - ▶ Move subarray boundaries one element to the right.

Insertion sort



- ▶ Repeat:
 - ▶ Examine the next element in the unsorted subarray.
 - ▶ Exchange this element with every entry to the left that is greater.
 - ▶ Move subarray boundaries one element to the right.

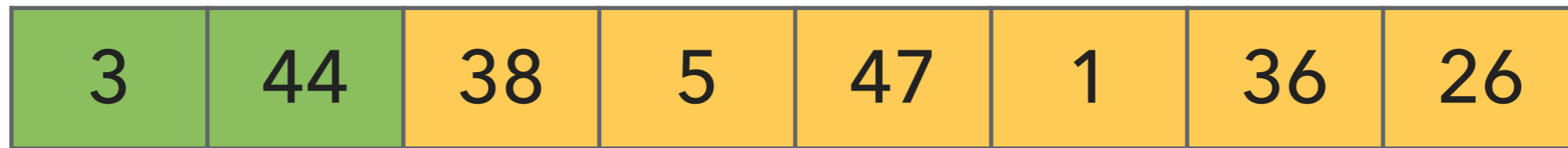
Insertion sort



▶ Repeat:

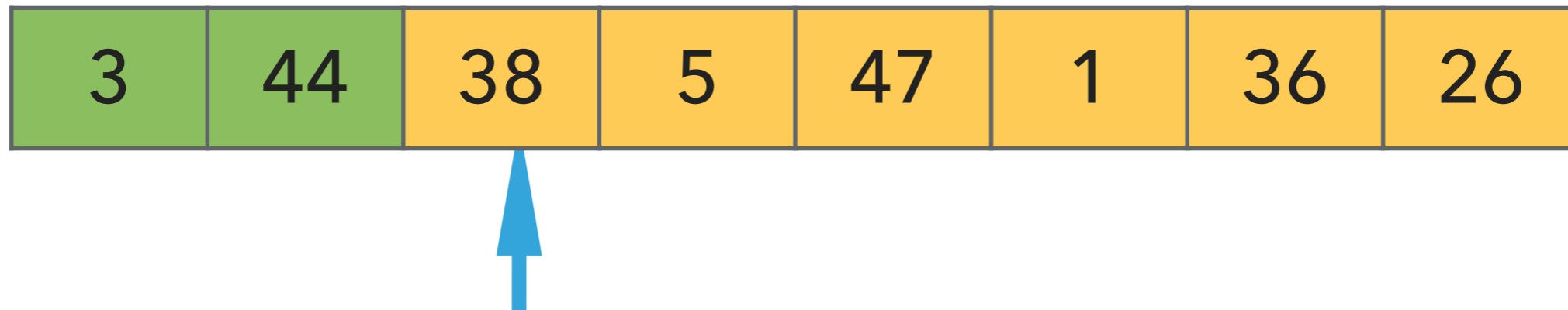
- ▶ Examine the next element in the unsorted subarray.
- ▶ Exchange this element with every entry to the left that is greater.
- ▶ Move subarray boundaries one element to the right.

Insertion sort



- ▶ Repeat:
 - ▶ Examine the next element in the unsorted subarray.
 - ▶ Exchange this element with every entry to the left that is greater.
 - ▶ Move subarray boundaries one element to the right.

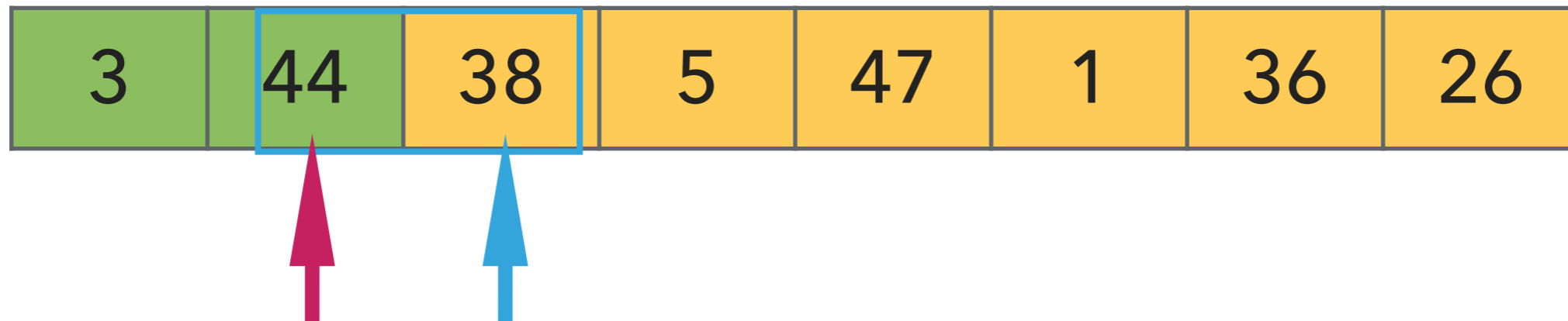
Insertion sort



▶ Repeat:

- ▶ Examine the next element in the unsorted subarray.
- ▶ Exchange this element with every entry to the left that is greater.
- ▶ Move subarray boundaries one element to the right.

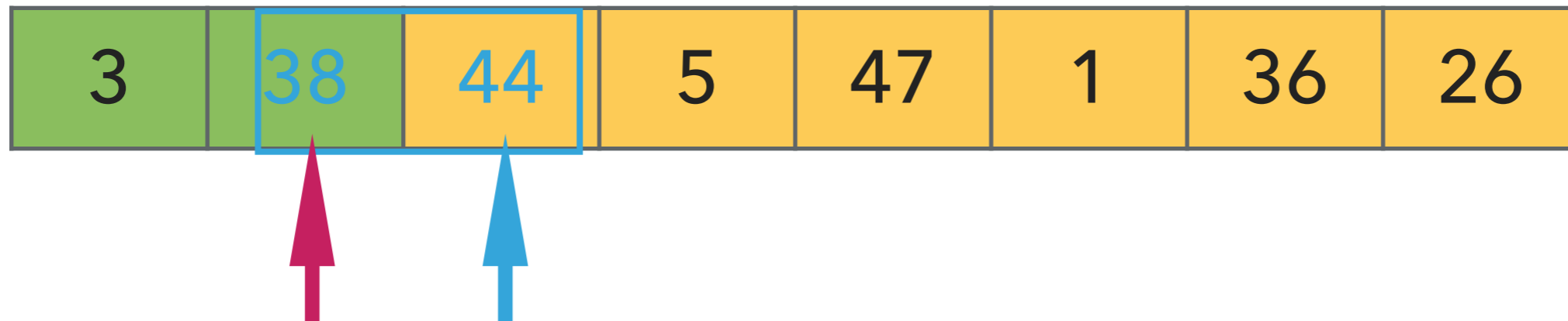
Insertion sort



▶ Repeat:

- ▶ Examine the next element in the unsorted subarray.
- ▶ Exchange this element with every entry to the left that is greater.
- ▶ Move subarray boundaries one element to the right.

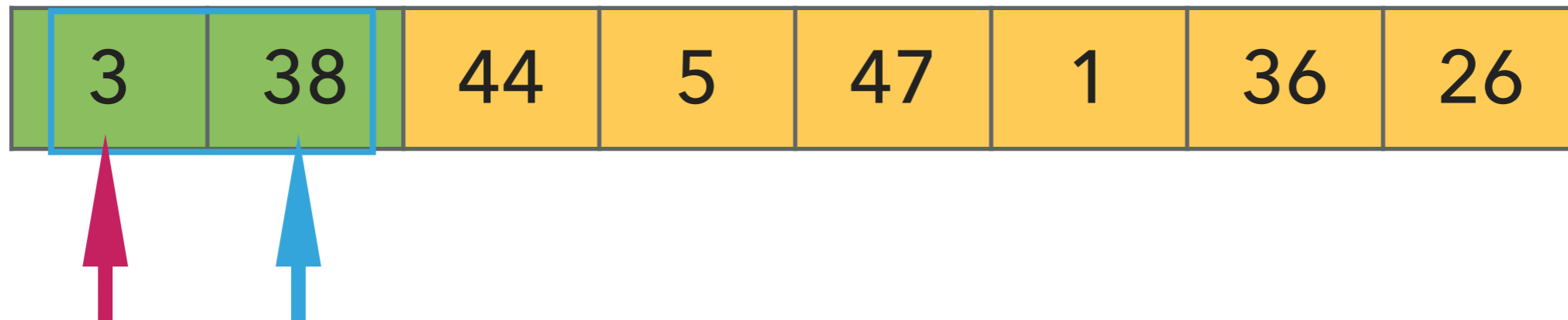
Insertion sort



▶ Repeat:

- ▶ Examine the next element in the unsorted subarray.
- ▶ Exchange this element with every entry to the left that is greater.
- ▶ Move subarray boundaries one element to the right.

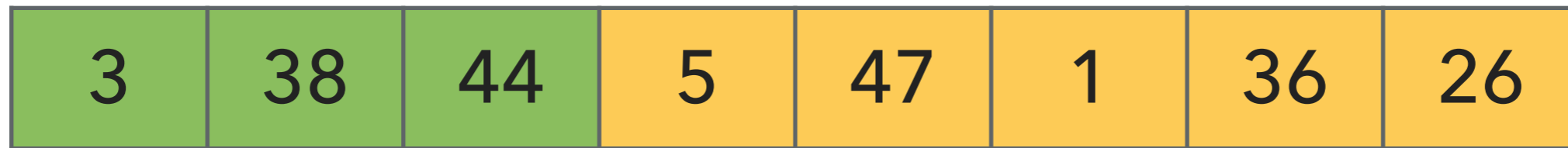
Insertion sort



▶ Repeat:

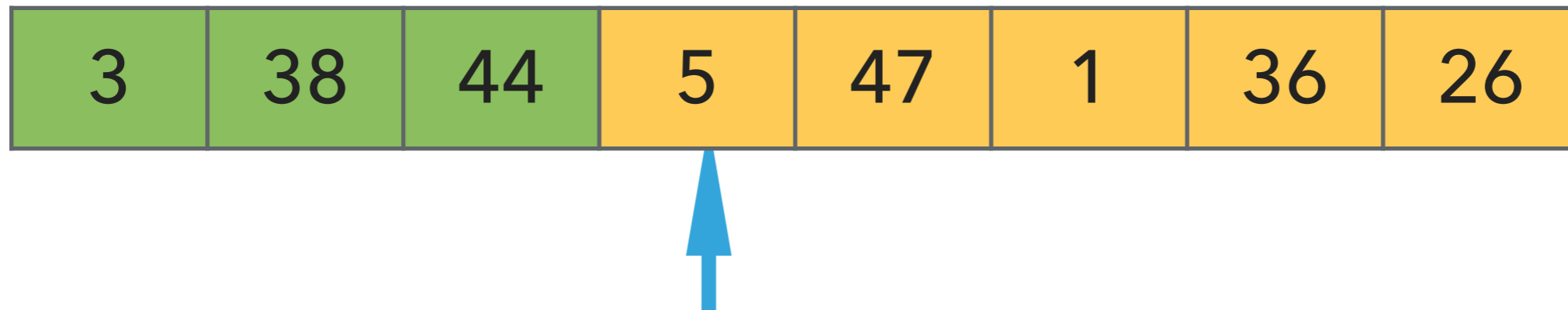
- ▶ Examine the next element in the unsorted subarray.
- ▶ Exchange this element with every entry to the left that is greater.
- ▶ Move subarray boundaries one element to the right.

Insertion sort



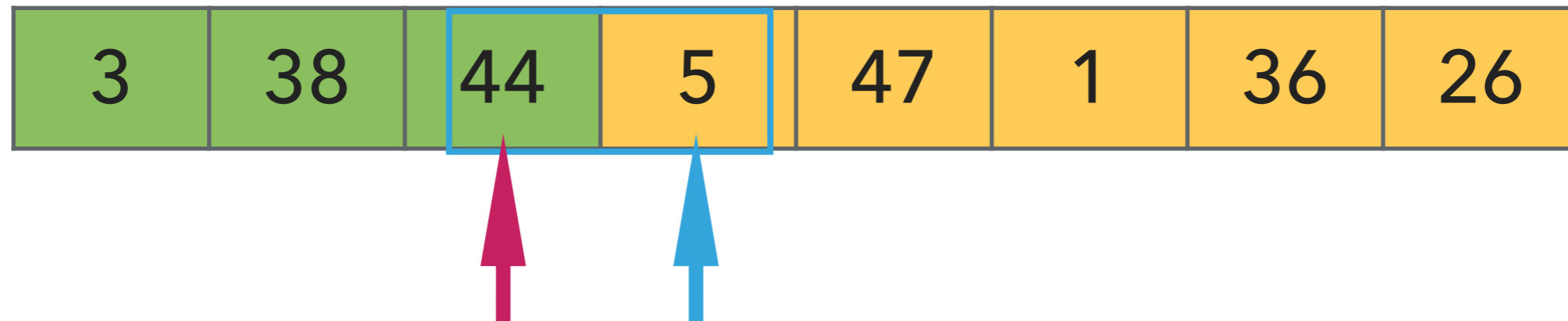
- ▶ Repeat:
 - ▶ Examine the next element in the unsorted subarray.
 - ▶ Exchange this element with every entry to the left that is greater.
 - ▶ Move subarray boundaries one element to the right.

Insertion sort



- ▶ Repeat:
 - ▶ Examine the next element in the unsorted subarray.
 - ▶ Exchange this element with every entry to the left that is greater.
 - ▶ Move subarray boundaries one element to the right.

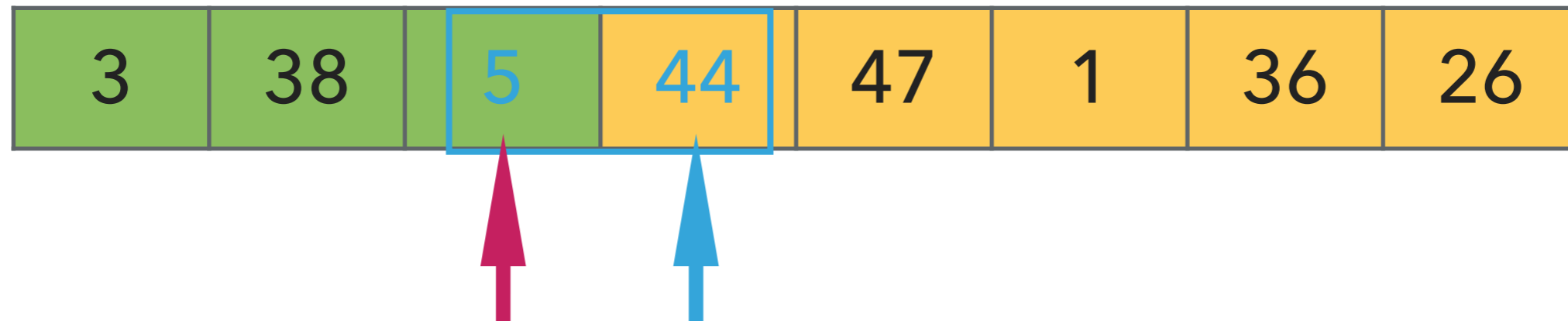
Insertion sort



▶ Repeat:

- ▶ Examine the next element in the unsorted subarray.
- ▶ Exchange this element with every entry to the left that is greater.
- ▶ Move subarray boundaries one element to the right.

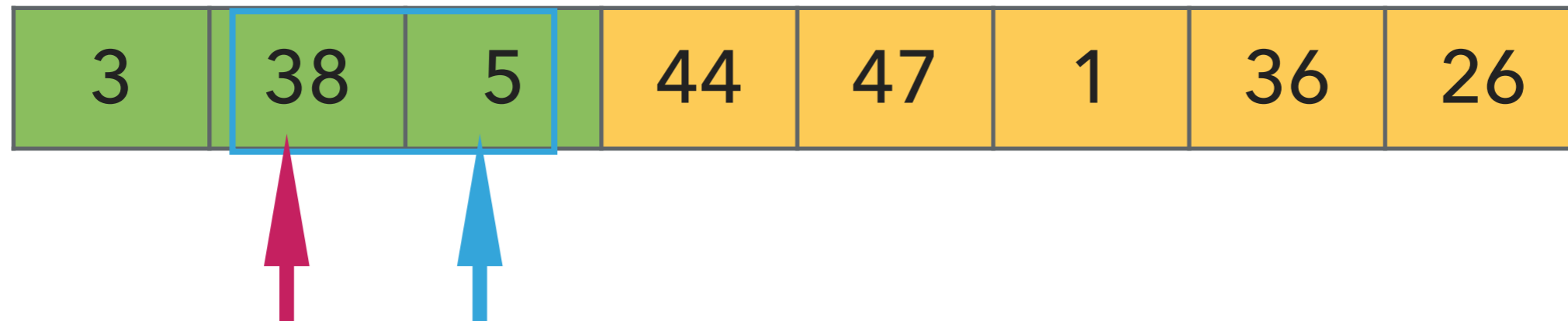
Insertion sort



▶ Repeat:

- ▶ Examine the next element in the unsorted subarray.
- ▶ Exchange this element with every entry to the left that is greater.
- ▶ Move subarray boundaries one element to the right.

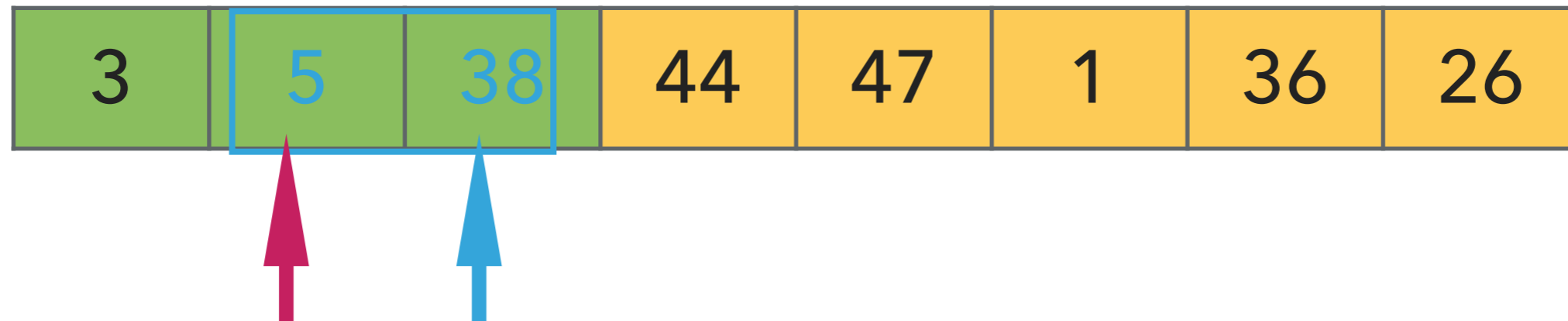
Insertion sort



▶ Repeat:

- ▶ Examine the next element in the unsorted subarray.
- ▶ Exchange this element with every entry to the left that is greater.
- ▶ Move subarray boundaries one element to the right.

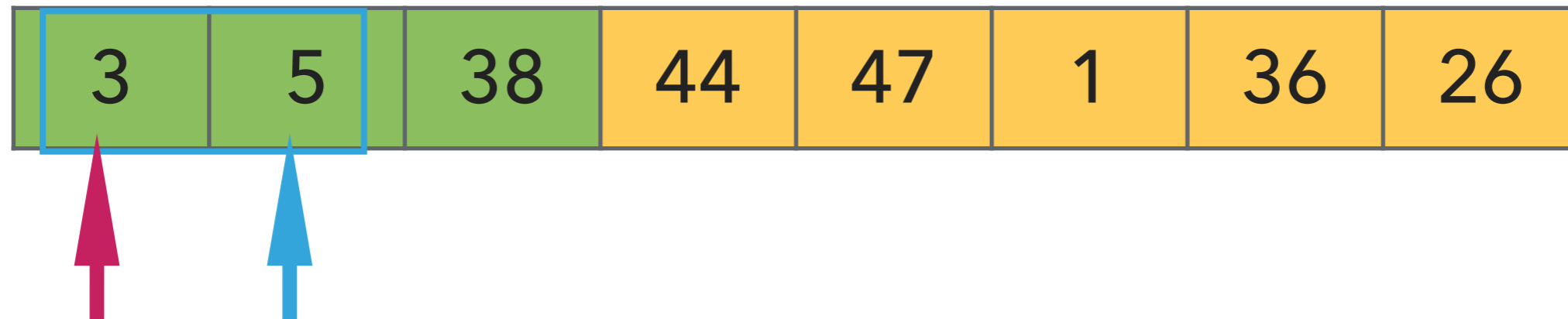
Insertion sort



▶ Repeat:

- ▶ Examine the next element in the unsorted subarray.
- ▶ Exchange this element with every entry to the left that is greater.
- ▶ Move subarray boundaries one element to the right.

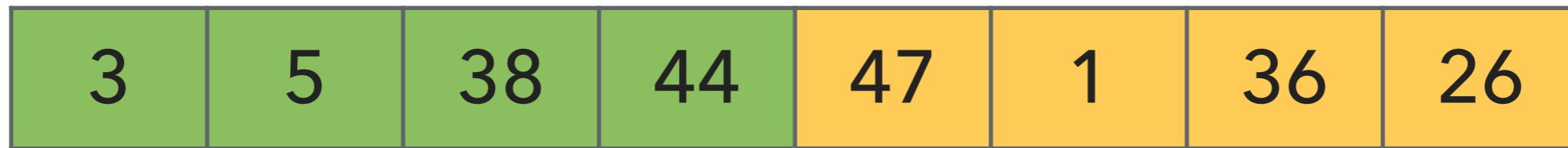
Insertion sort



▶ Repeat:

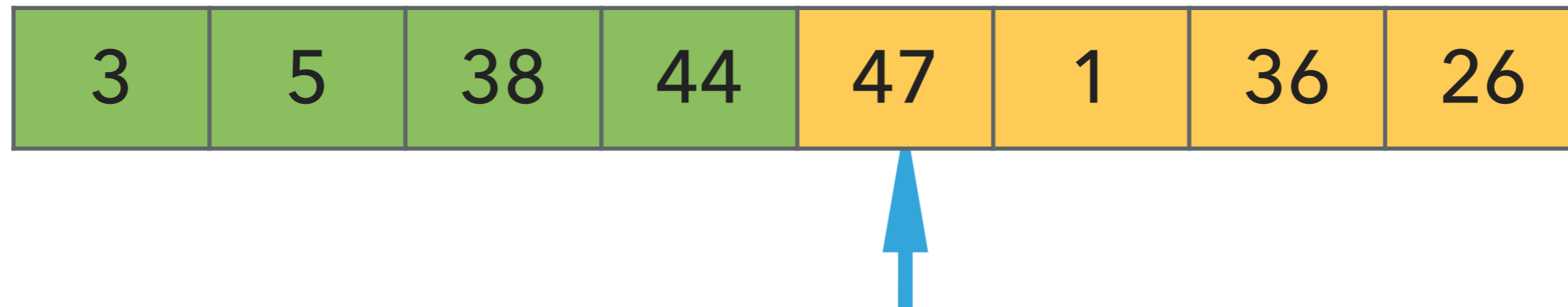
- ▶ Examine the next element in the unsorted subarray.
- ▶ Exchange this element with every entry to the left that is greater.
- ▶ Move subarray boundaries one element to the right.

Insertion sort



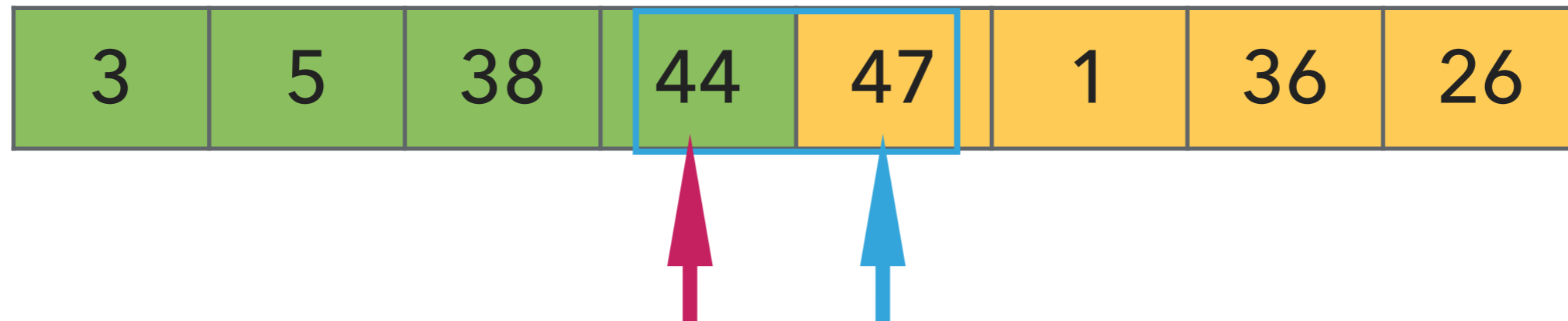
- ▶ Repeat:
 - ▶ Examine the next element in the unsorted subarray.
 - ▶ Exchange this element with every entry to the left that is greater.
 - ▶ Move subarray boundaries one element to the right.

Insertion sort



- ▶ Repeat:
 - ▶ Examine the next element in the unsorted subarray.
 - ▶ Exchange this element with every entry to the left that is greater.
 - ▶ Move subarray boundaries one element to the right.

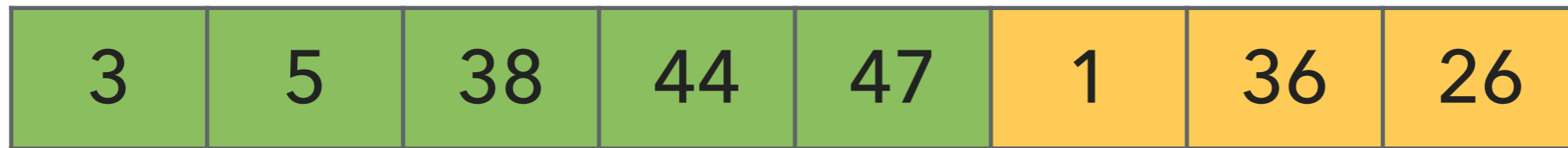
Insertion sort



▶ Repeat:

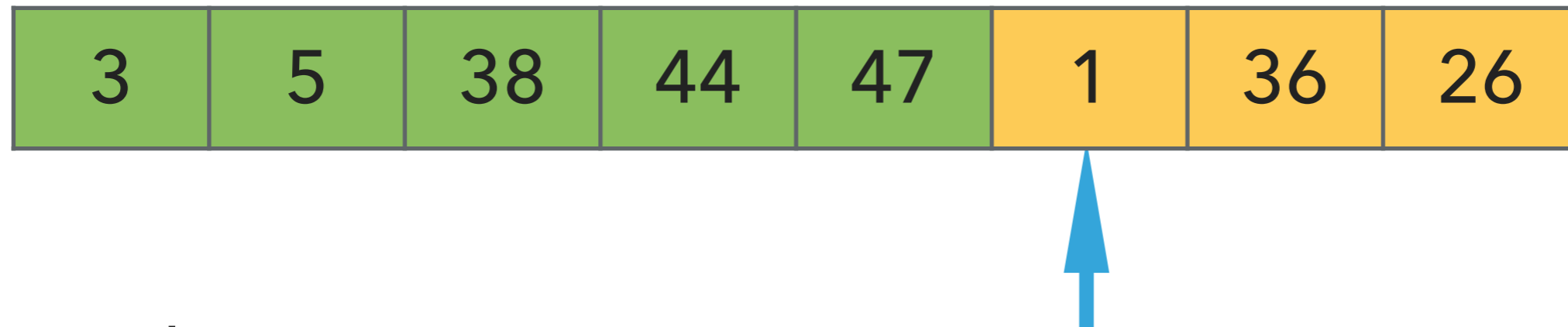
- ▶ Examine the next element in the unsorted subarray.
- ▶ Exchange this element with every entry to the left that is greater.
- ▶ Move subarray boundaries one element to the right.

Insertion sort



- ▶ Repeat:
 - ▶ Examine the next element in the unsorted subarray.
 - ▶ Exchange this element with every entry to the left that is greater.
 - ▶ Move subarray boundaries one element to the right.

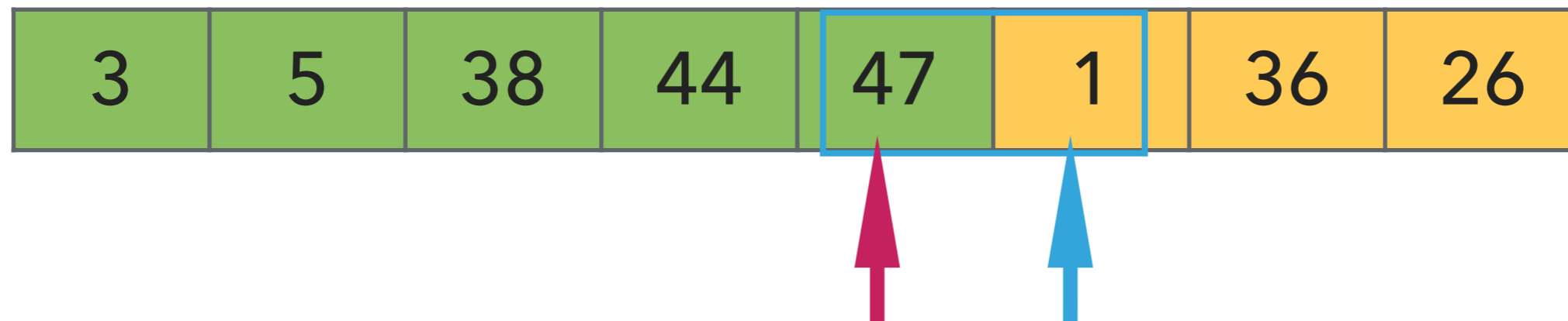
Insertion sort



▶ Repeat:

- ▶ Examine the next element in the unsorted subarray.
- ▶ Exchange this element with every entry to the left that is greater.
- ▶ Move subarray boundaries one element to the right.

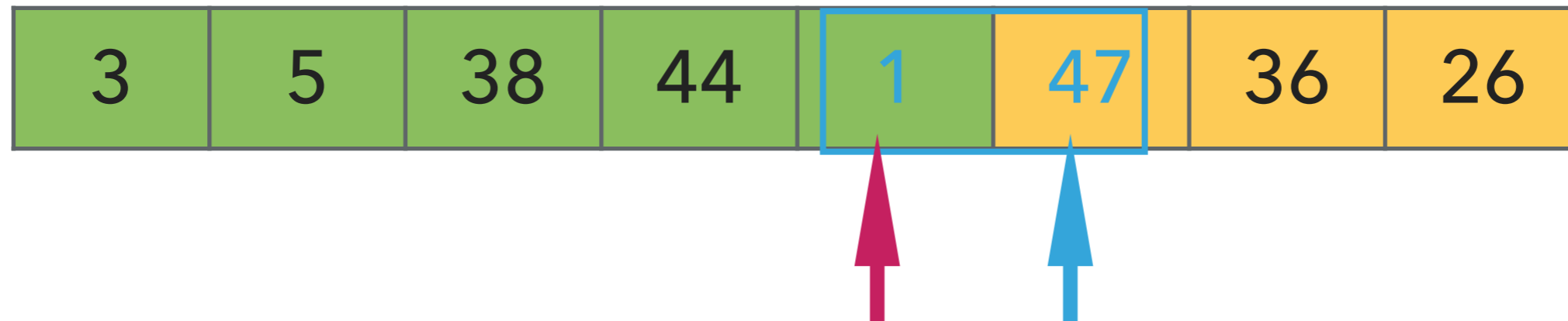
Insertion sort



▶ Repeat:

- ▶ Examine the next element in the unsorted subarray.
- ▶ Exchange this element with every entry to the left that is greater.
- ▶ Move subarray boundaries one element to the right.

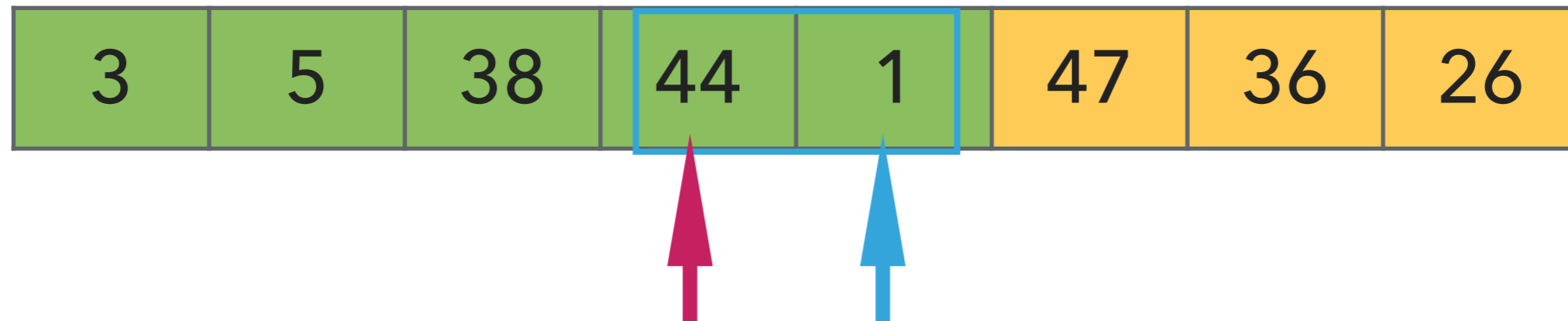
Insertion sort



▶ Repeat:

- ▶ Examine the next element in the unsorted subarray.
- ▶ Exchange this element with every entry to the left that is greater.
- ▶ Move subarray boundaries one element to the right.

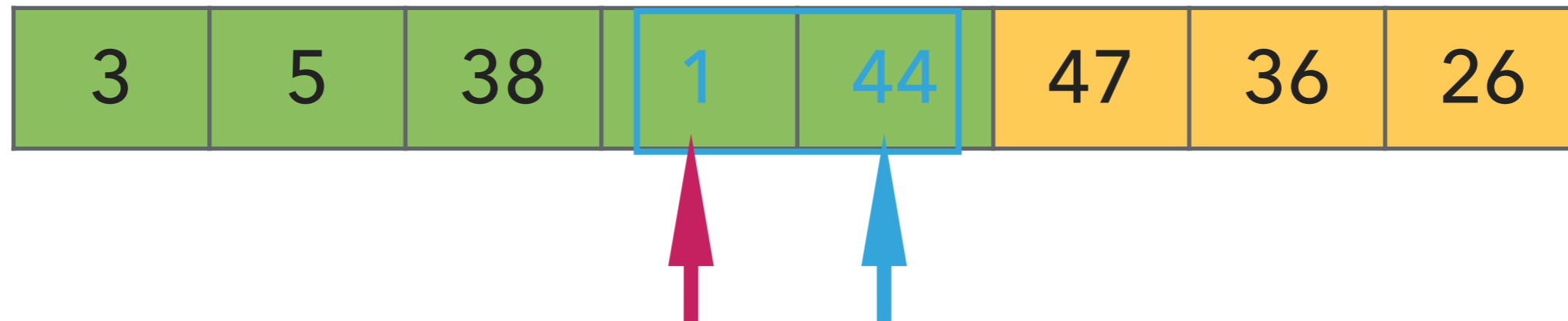
Insertion sort



▶ Repeat:

- ▶ Examine the next element in the unsorted subarray.
- ▶ Exchange this element with every entry to the left that is greater.
- ▶ Move subarray boundaries one element to the right.

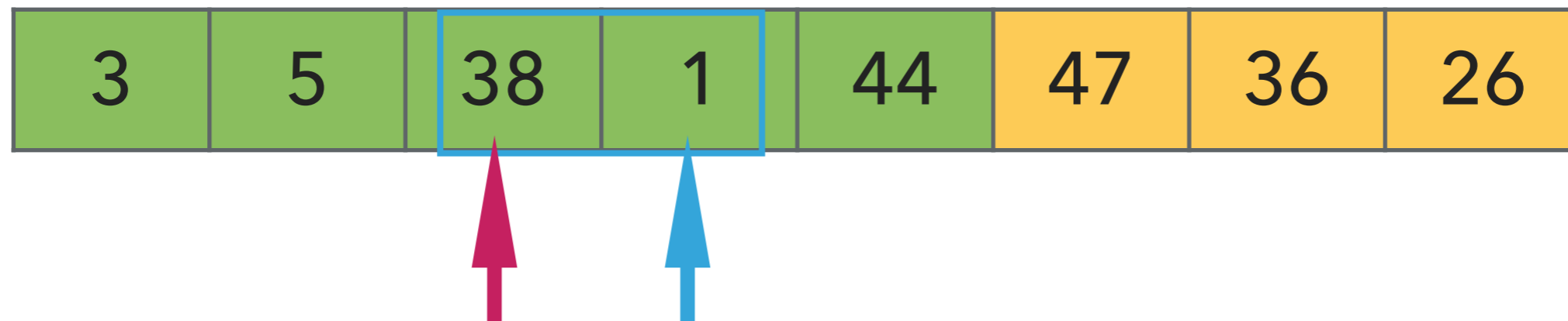
Insertion sort



▶ Repeat:

- ▶ Examine the next element in the unsorted subarray.
- ▶ Exchange this element with every entry to the left that is greater.
- ▶ Move subarray boundaries one element to the right.

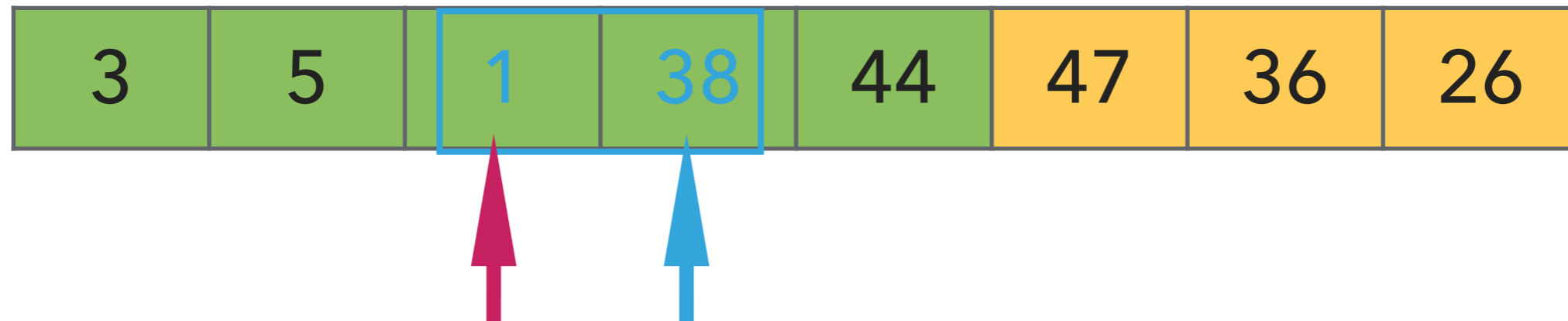
Insertion sort



▶ Repeat:

- ▶ Examine the next element in the unsorted subarray.
- ▶ Exchange this element with every entry to the left that is greater.
- ▶ Move subarray boundaries one element to the right.

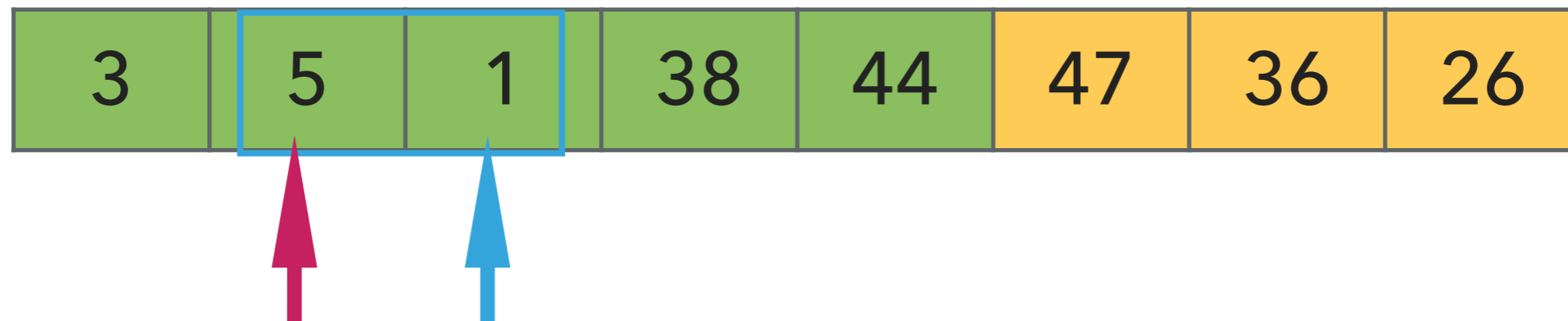
Insertion sort



▶ Repeat:

- ▶ Examine the next element in the unsorted subarray.
- ▶ Exchange this element with every entry to the left that is greater.
- ▶ Move subarray boundaries one element to the right.

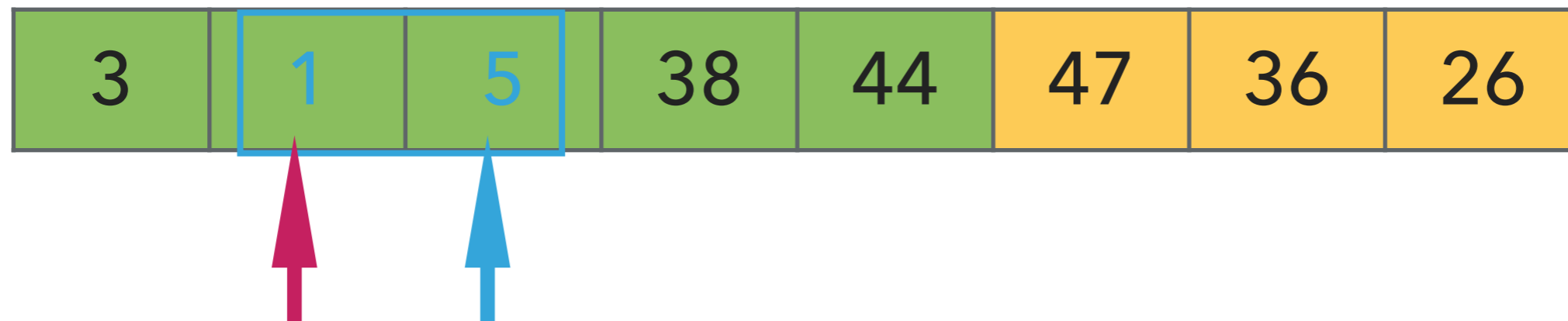
Insertion sort



▶ Repeat:

- ▶ Examine the next element in the unsorted subarray.
- ▶ Exchange this element with every entry to the left that is greater.
- ▶ Move subarray boundaries one element to the right.

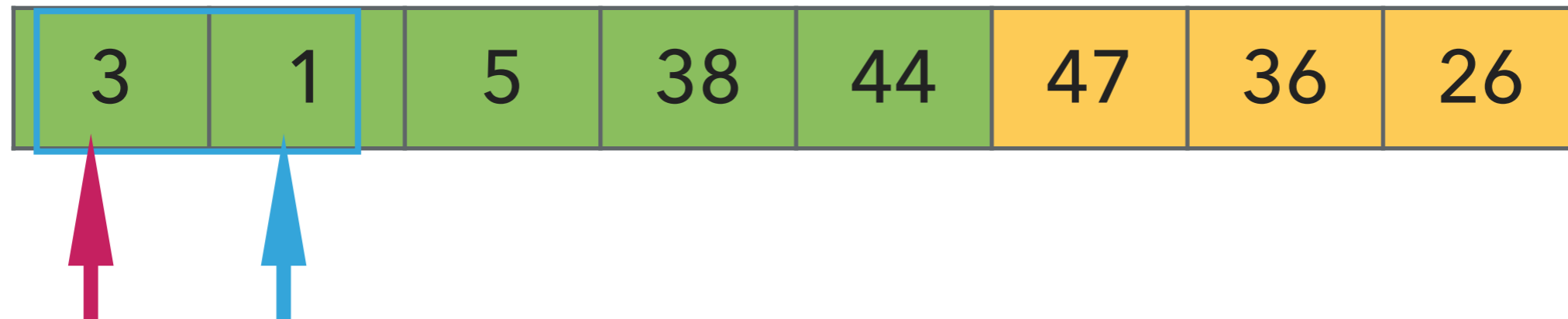
Insertion sort



▶ Repeat:

- ▶ Examine the next element in the unsorted subarray.
- ▶ Exchange this element with every entry to the left that is greater.
- ▶ Move subarray boundaries one element to the right.

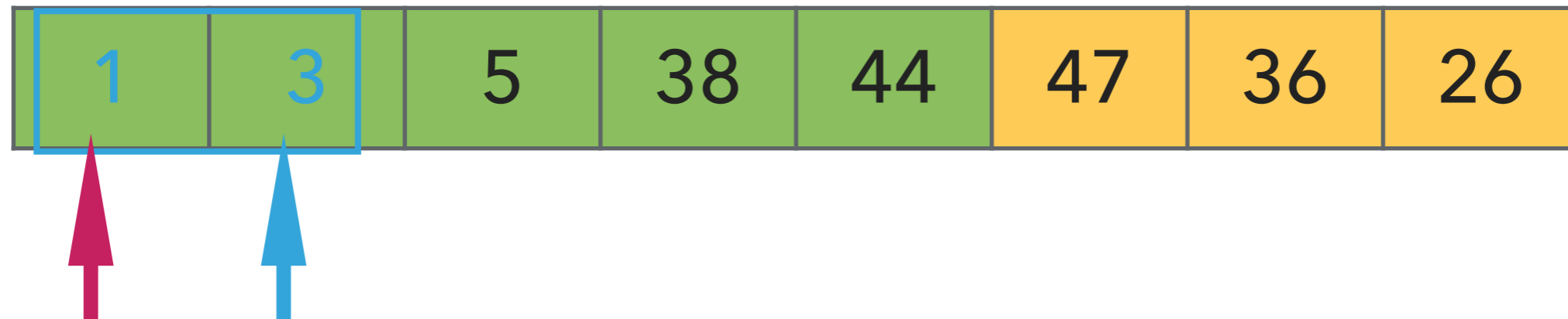
Insertion sort



▶ Repeat:

- ▶ Examine the next element in the unsorted subarray.
- ▶ Exchange this element with every entry to the left that is greater.
- ▶ Move subarray boundaries one element to the right.

Insertion sort



▶ Repeat:

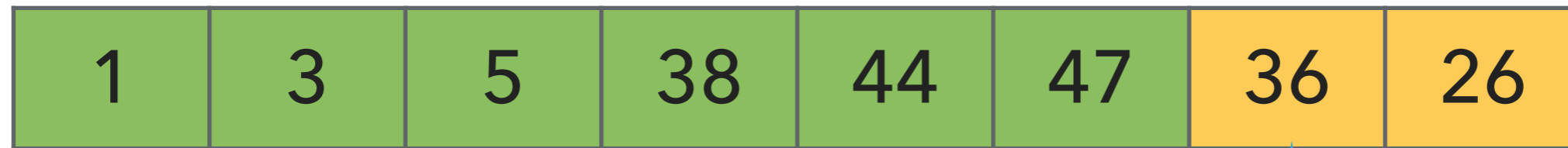
- ▶ Examine the next element in the unsorted subarray.
- ▶ Exchange this element with every entry to the left that is greater.
- ▶ Move subarray boundaries one element to the right.

Insertion sort



- ▶ Repeat:
 - ▶ Examine the next element in the unsorted subarray.
 - ▶ Exchange this element with every entry to the left that is greater.
 - ▶ Move subarray boundaries one element to the right.

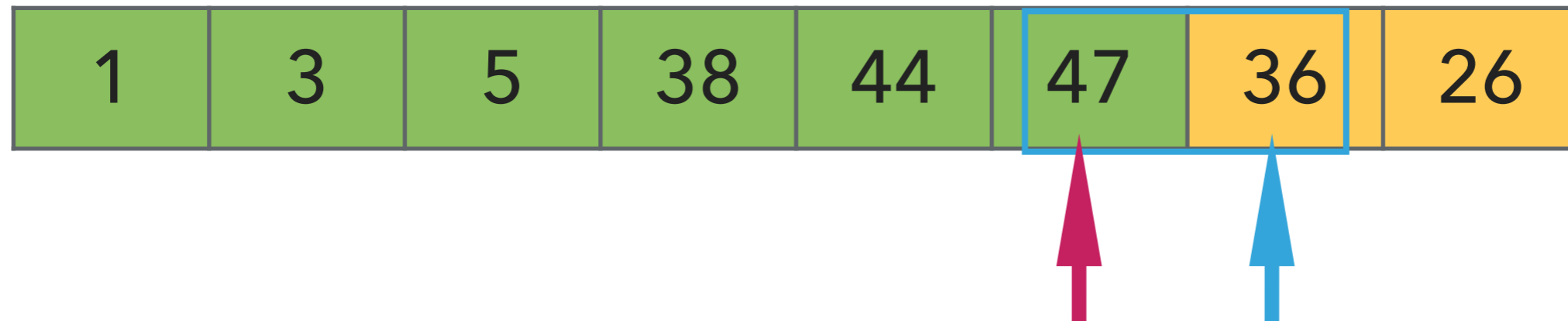
Insertion sort



▶ Repeat:

- ▶ Examine the next element in the unsorted subarray.
- ▶ Exchange this element with every entry to the left that is greater.
- ▶ Move subarray boundaries one element to the right.

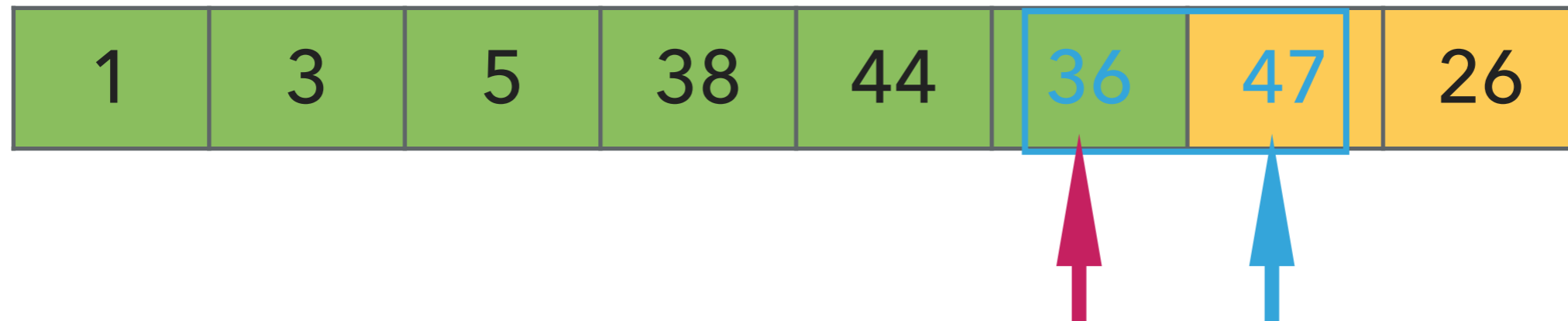
Insertion sort



▶ Repeat:

- ▶ Examine the next element in the unsorted subarray.
- ▶ Exchange this element with every entry to the left that is greater.
- ▶ Move subarray boundaries one element to the right.

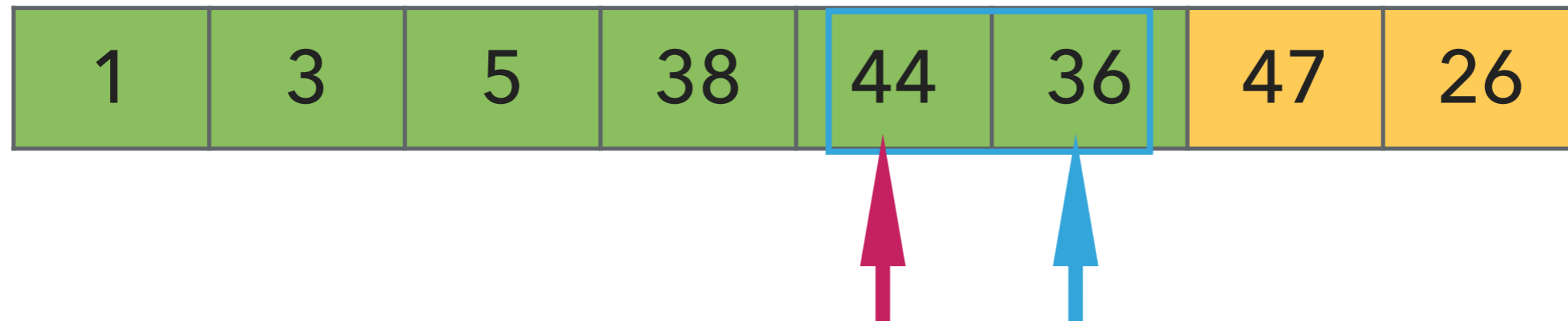
Insertion sort



▶ Repeat:

- ▶ Examine the next element in the unsorted subarray.
- ▶ Exchange this element with every entry to the left that is greater.
- ▶ Move subarray boundaries one element to the right.

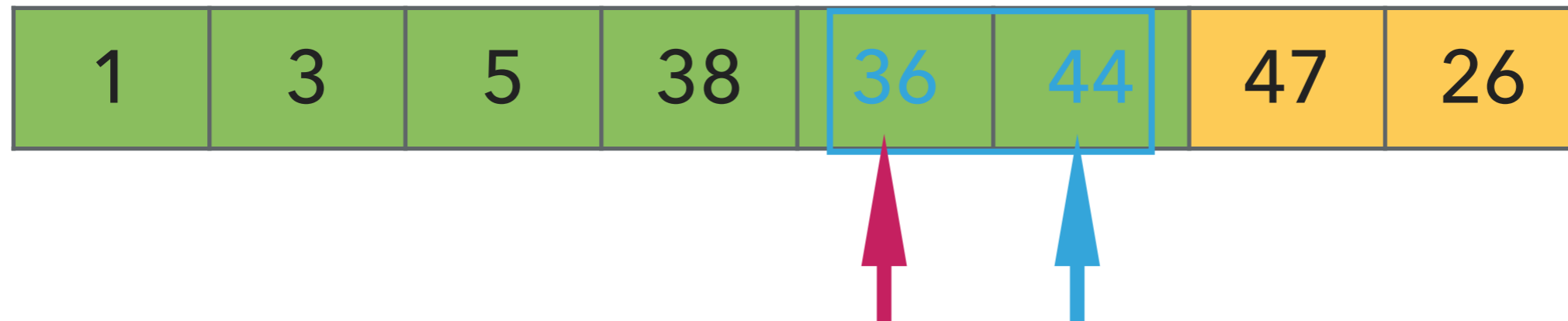
Insertion sort



▶ Repeat:

- ▶ Examine the next element in the unsorted subarray.
- ▶ Exchange this element with every entry to the left that is greater.
- ▶ Move subarray boundaries one element to the right.

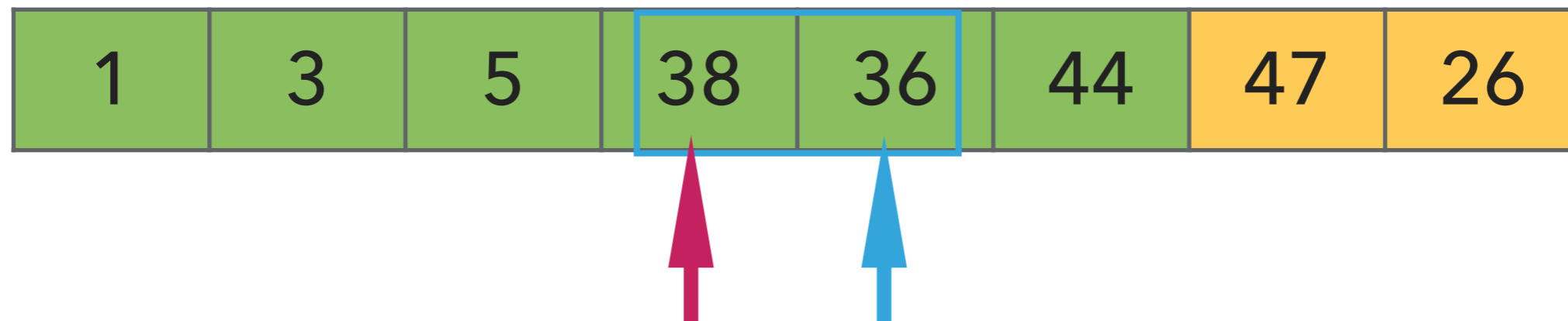
Insertion sort



▶ Repeat:

- ▶ Examine the next element in the unsorted subarray.
- ▶ Exchange this element with every entry to the left that is greater.
- ▶ Move subarray boundaries one element to the right.

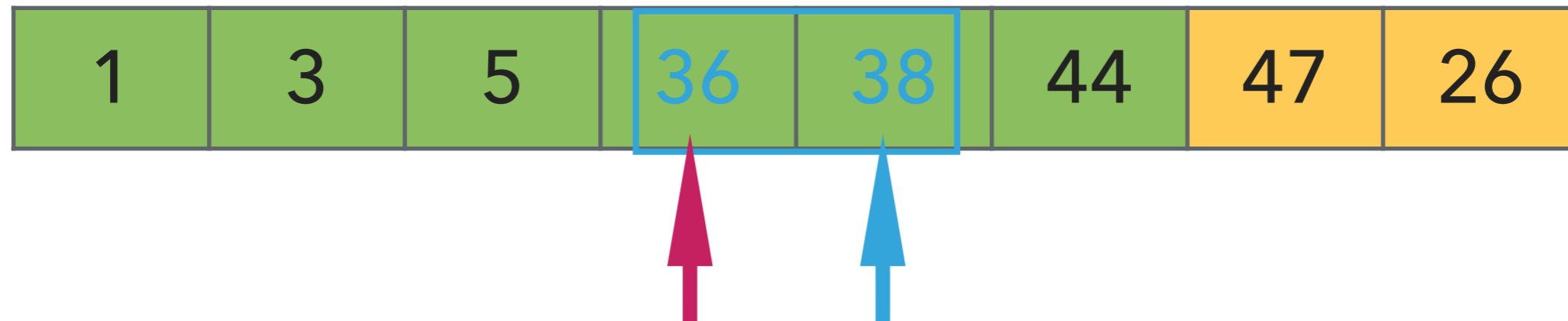
Insertion sort



▶ Repeat:

- ▶ Examine the next element in the unsorted subarray.
- ▶ Exchange this element with every entry to the left that is greater.
- ▶ Move subarray boundaries one element to the right.

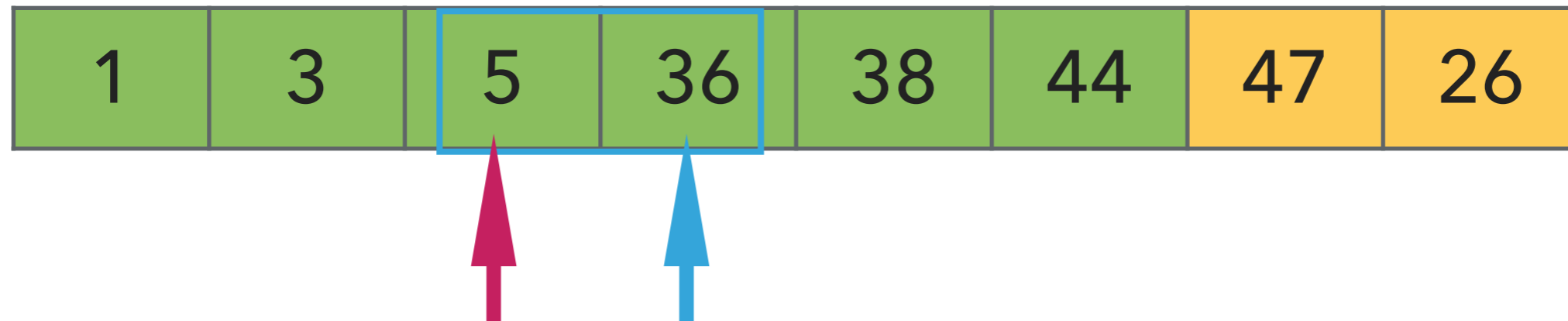
Insertion sort



▶ Repeat:

- ▶ Examine the next element in the unsorted subarray.
- ▶ Exchange this element with every entry to the left that is greater.
- ▶ Move subarray boundaries one element to the right.

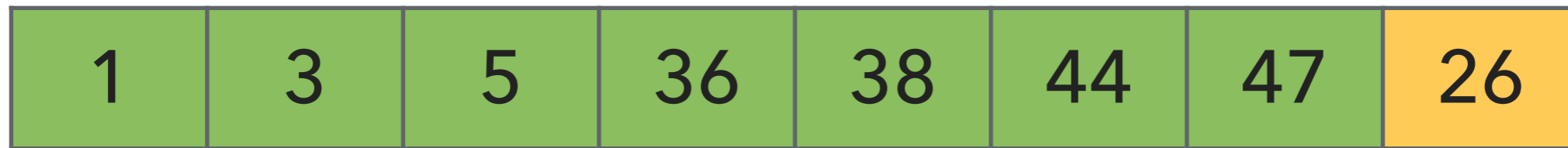
Insertion sort



▶ Repeat:

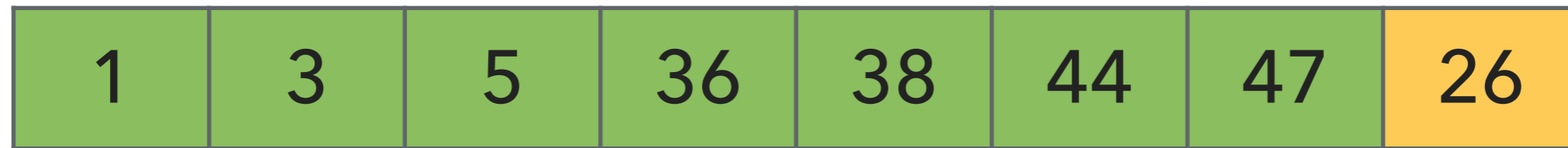
- ▶ Examine the next element in the unsorted subarray.
- ▶ Exchange this element with every entry to the left that is greater.
- ▶ Move subarray boundaries one element to the right.

Insertion sort



- ▶ Repeat:
 - ▶ Examine the next element in the unsorted subarray.
 - ▶ Exchange this element with every entry to the left that is greater.
 - ▶ Move subarray boundaries one element to the right.

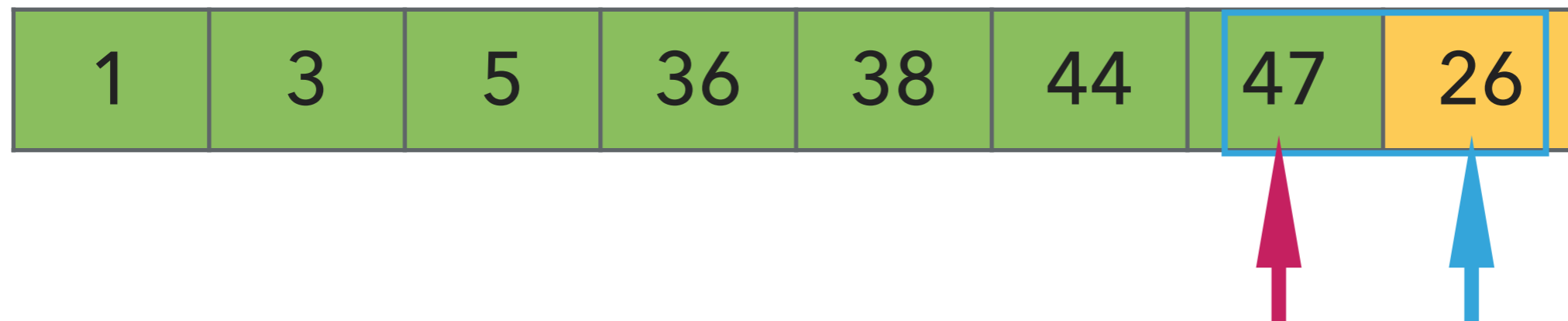
Insertion sort



▶ Repeat:

- ▶ Examine the next element in the unsorted subarray.
- ▶ Exchange this element with every entry to the left that is greater.
- ▶ Move subarray boundaries one element to the right.

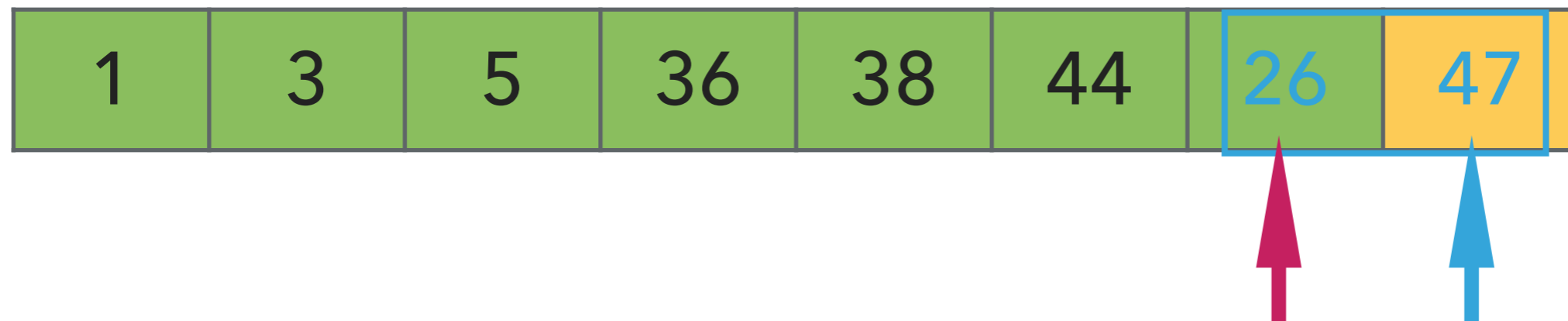
Insertion sort



▶ Repeat:

- ▶ Examine the next element in the unsorted subarray.
- ▶ Exchange this element with every entry to the left that is greater.
- ▶ Move subarray boundaries one element to the right.

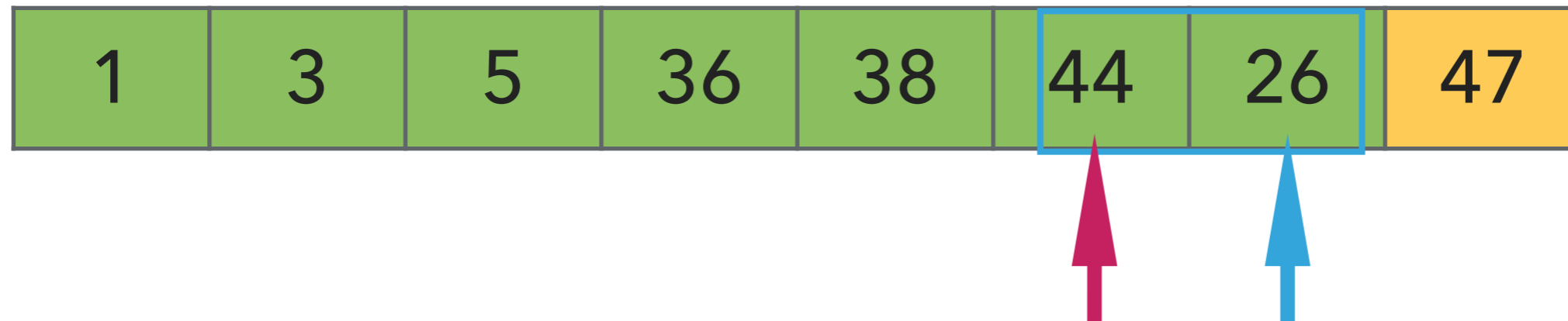
Insertion sort



▶ Repeat:

- ▶ Examine the next element in the unsorted subarray.
- ▶ Exchange this element with every entry to the left that is greater.
- ▶ Move subarray boundaries one element to the right.

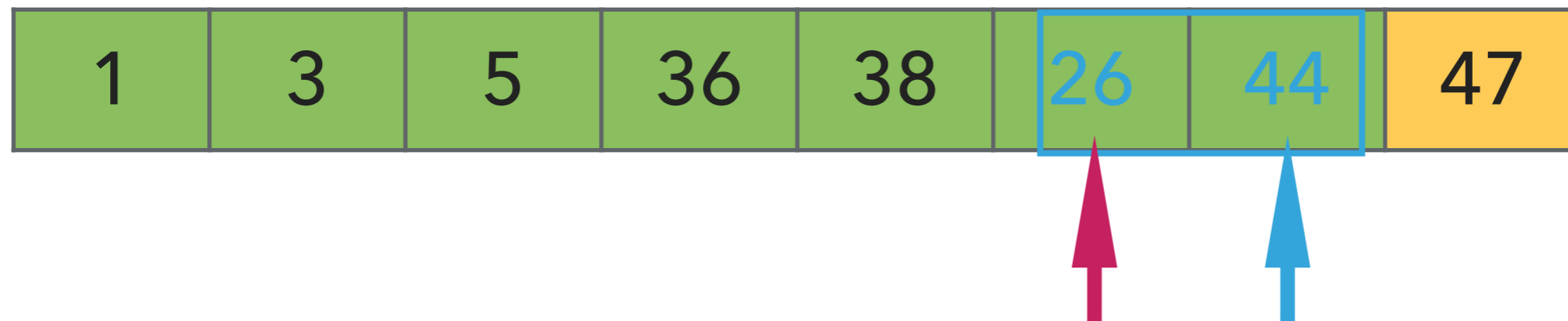
Insertion sort



▶ Repeat:

- ▶ Examine the next element in the unsorted subarray.
- ▶ Exchange this element with every entry to the left that is greater.
- ▶ Move subarray boundaries one element to the right.

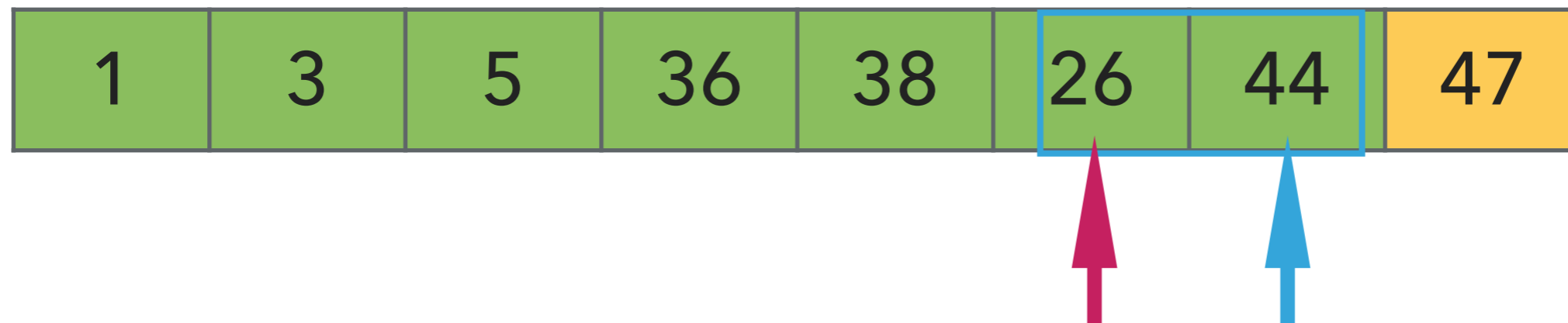
Insertion sort



▶ Repeat:

- ▶ Examine the next element in the unsorted subarray.
- ▶ Exchange this element with every entry to the left that is greater.
- ▶ Move subarray boundaries one element to the right.

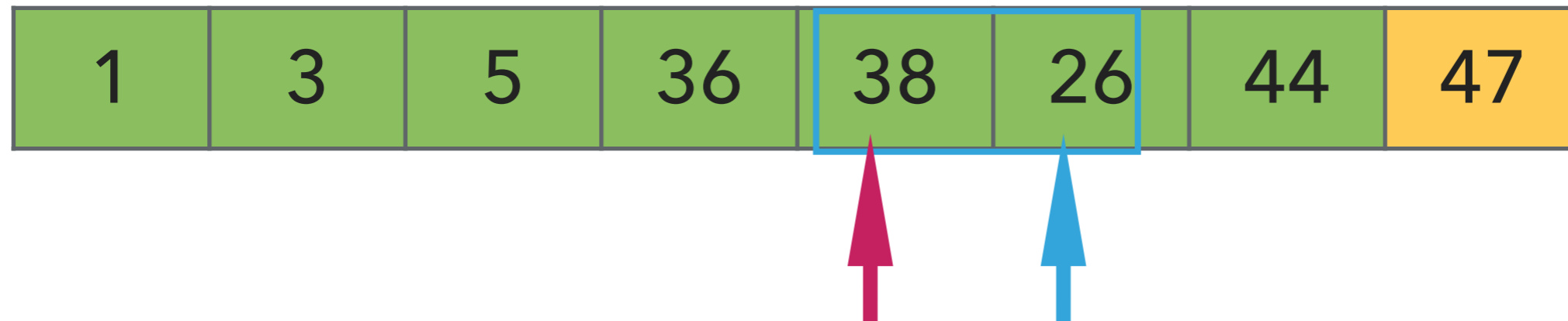
Insertion sort



▶ Repeat:

- ▶ Examine the next element in the unsorted subarray.
- ▶ Exchange this element with every entry to the left that is greater.
- ▶ Move subarray boundaries one element to the right.

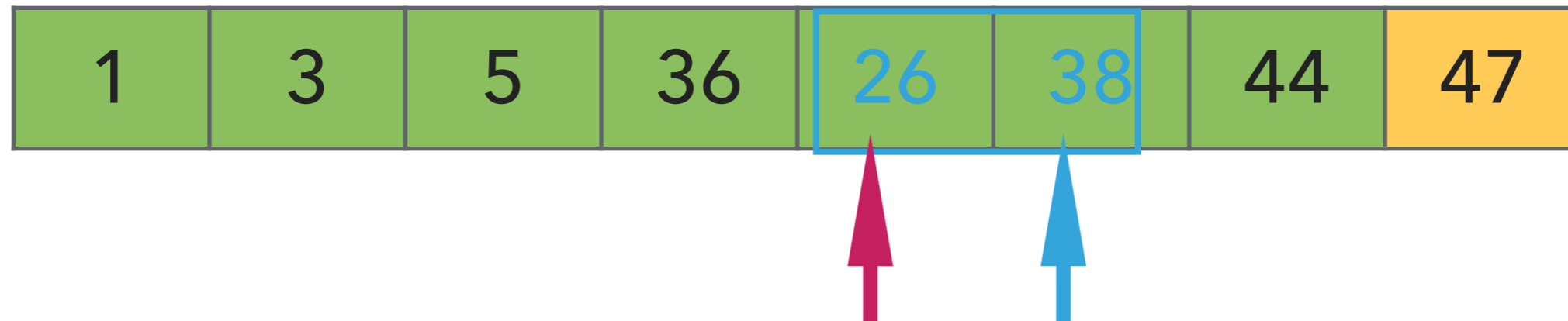
Insertion sort



▶ Repeat:

- ▶ Examine the next element in the unsorted subarray.
- ▶ Exchange this element with every entry to the left that is greater.
- ▶ Move subarray boundaries one element to the right.

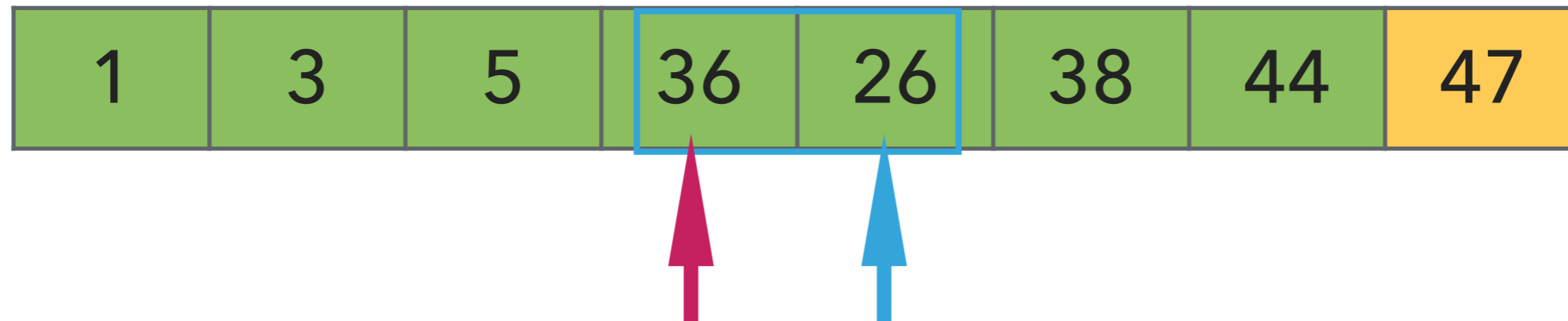
Insertion sort



▶ Repeat:

- ▶ Examine the next element in the unsorted subarray.
- ▶ Exchange this element with every entry to the left that is greater.
- ▶ Move subarray boundaries one element to the right.

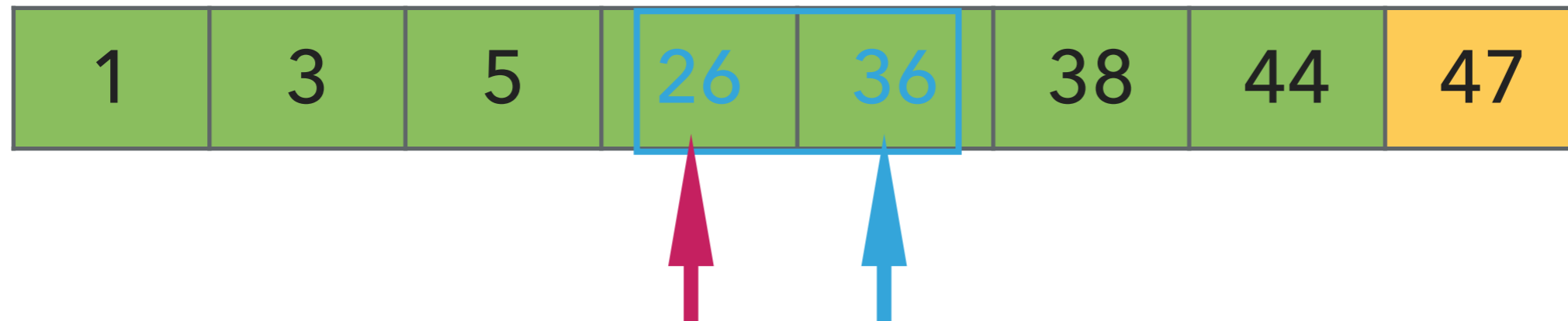
Insertion sort



▶ Repeat:

- ▶ Examine the next element in the unsorted subarray.
- ▶ Exchange this element with every entry to the left that is greater.
- ▶ Move subarray boundaries one element to the right.

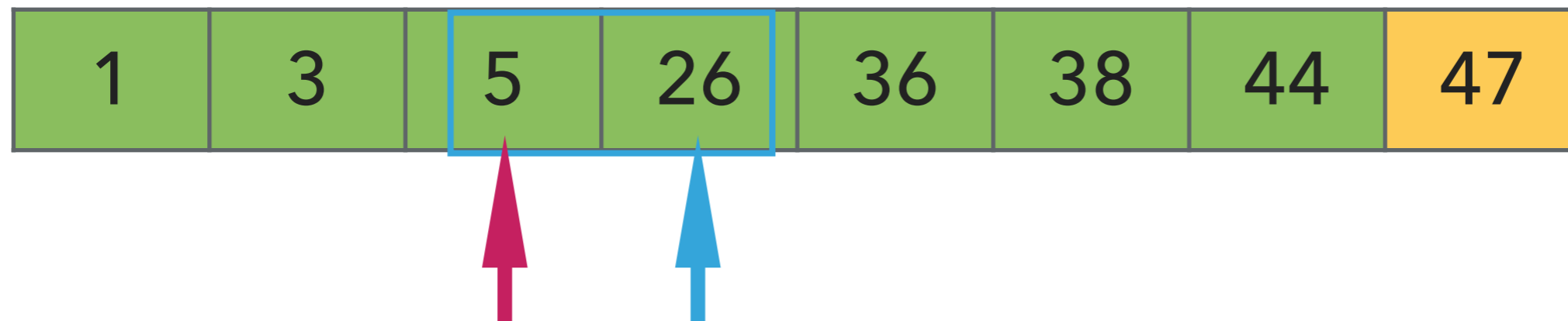
Insertion sort



▶ Repeat:

- ▶ Examine the next element in the unsorted subarray.
- ▶ Exchange this element with every entry to the left that is greater.
- ▶ Move subarray boundaries one element to the right.

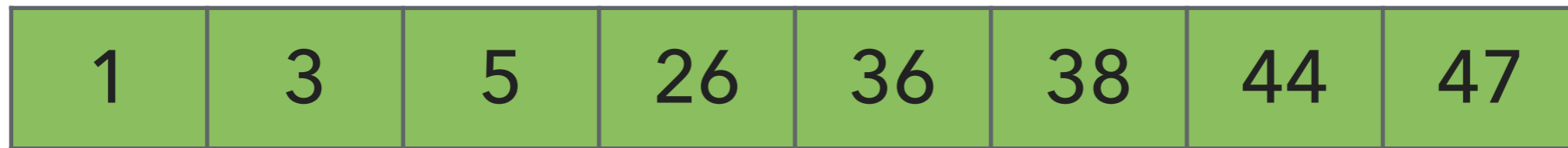
Insertion sort



▶ Repeat:

- ▶ Examine the next element in the unsorted subarray.
- ▶ Exchange this element with every entry to the left that is greater.
- ▶ Move subarray boundaries one element to the right.

Insertion sort



- ▶ Repeat:
 - ▶ Examine the next element in the unsorted subarray.
 - ▶ Exchange this element with every entry to the left that is greater.
 - ▶ Move subarray boundaries one element to the right.

2.1 INSERTION SORT DEMO

Demo with Cards



<http://algs4.cs.princeton.edu>

INSERTION SORT

In case you didn't get this...

- ▶ <https://www.youtube.com/watch?v=ROalU379l3U>

INSERTION SORT

Insertion sort

```
public static void sort(Comparable[] a) {  
    // for loop to iterate through each element of the array  
  
    // Moving right to left, exchange a[i] with every larger  
    // entry to its left  
  
}
```

INSERTION SORT

Insertion sort

```
public static void sort(Comparable[] a) {  
    int n = a.length;  
    for (int i = 0; i < n; i++) {  
        for (int j = i; j > 0; j--) {  
            if (less(a[j], a[j-1]))  
                exch(a, j, j-1);  
            else  
                break;  
        }  
    }  
}
```

← In iteration i

← Move from right to left,
exchange $a[i]$ with entry to
the left, if it's larger

▶ **Invariants:** At the end of each iteration i :

▶ the array a is sorted in ascending order for the first $i+1$ elements $a[0..i]$

Insertion sort: mathematical analysis for **worst-case**

```
public static void sort(Comparable[] a) {  
    int n = a.length;  
    for (int i = 0; i < n; i++) {  
        for (int j = i; j > 0; j--) {  
            if (less(a[j], a[j-1]))  
                exch(a, j, j-1);  
            else  
                break;  
        }  
    }  
}
```

▶ Comparisons:

▶ Exchanges: ?

▶ In-place?

▶ Stable?

Insertion sort: mathematical analysis for **worst-case**

```
public static void sort(Comparable[] a) {  
    int n = a.length;  
    for (int i = 0; i < n; i++) {  
        for (int j = i; j > 0; j--) {  
            if (less(a[j], a[j-1]))  
                exch(a, j, j-1);  
            else  
                break;  
        }  
    }  
}
```

▶ **Comparisons:** $0 + 1 + 2 + \dots + (n - 2) + (n - 1) \sim n^2/2$, that is $O(n^2)$.

▶ **Exchanges:** ?

▶ **In-place?**

▶ **Stable?**

Insertion sort: mathematical analysis for **worst-case**

```
public static void sort(Comparable[] a) {
    int n = a.length;
    for (int i = 0; i < n; i++) {
        for (int j = i; j > 0; j--) {
            if (less(a[j], a[j-1]))
                exch(a, j, j-1);
            else
                break;
        }
    }
}
```

- ▶ **Comparisons:** $0 + 1 + 2 + \dots + (n - 2) + (n - 1) \sim n^2/2$, that is $O(n^2)$.
- ▶ **Exchanges:** $0 + 1 + 2 + \dots + (n - 2) + (n - 1) \sim n^2/2$, that is $O(n^2)$.
- ▶ Worst-case running time is **quadratic**. Worst case = array sorted in reverse order.
- ▶ Every element moves all the way to the left.
- ▶ **In-place**, requires almost no additional memory.
- ▶ **Stable**

Insertion sort: average and best case

```
public static void sort(Comparable[] a) {
    int n = a.length;
    for (int i = 0; i < n; i++) {
        for (int j = i; j > 0; j--) {
            if (less(a[j], a[j-1]))
                exch(a, j, j-1);
            else
                break;
        }
    }
}
```

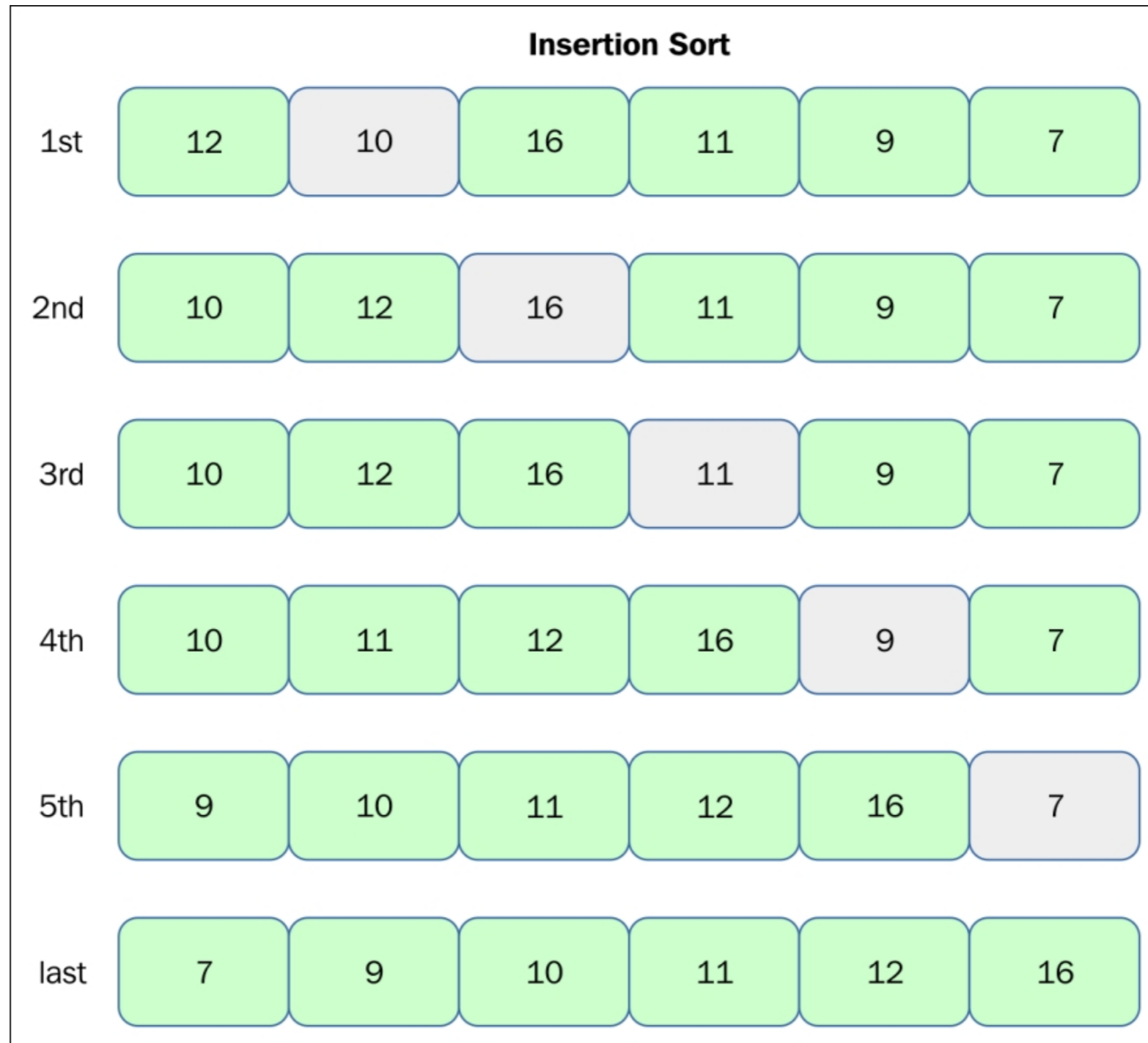
- ▶ **Average case:** quadratic for both comparisons and exchanges $\sim n^2/4$ when sorting a randomly ordered array. (2X faster than selection sort on average)
 - ▶ Expect each entry to move halfway back: $0 + 0.5 + 1 + \dots + (n-1)/2 \sim (n/2) * (n/2) \sim n^2/4$
- ▶ **Best case:** $n - 1$ comparisons (validate) and 0 exchanges for an already sorted array.

Practice Time (Use your cards)

- ▶ Using insertion sort, sort the array with elements [12,10,16,11,9,7].
- ▶ Visualize your work for every iteration of the algorithm.

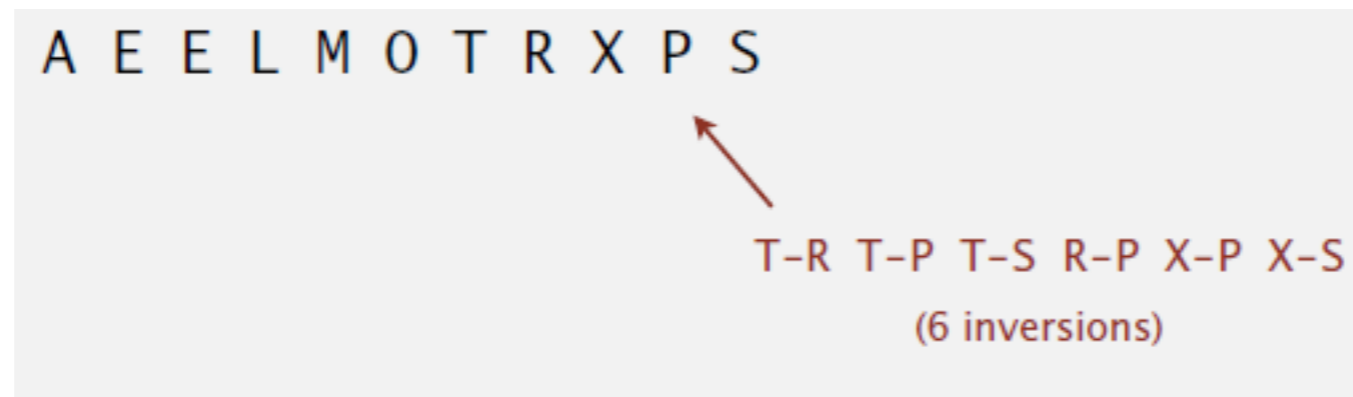
INSERTION SORT

Answer



Insertion Sort

- ▶ For partially-sorted arrays, insertion sort runs in **linear time**
- ▶ Number of exchanges equals number of **inversions**
- ▶ **Inversion** = pair of keys that are out of order



- ▶ Partially sorted examples
 - ▶ 1) Appending a subarray of size 10 to a sorted subarray of size N
 - ▶ 2) An array of size N with only 10 entries out of place

Lecture 12: Insertion Sort & Mergesort

- ▶ Insertion sort
- ▶ Comparators
- ▶ Mergesort

Comparable

- ▶ Interface with a single method that we need to implement:
`public int compareTo(T that)`
- ▶ Implement it so that `v.compareTo(w)`:
 - ▶ Returns >0 if `v` is greater than `w`.
 - ▶ Returns <0 if `v` is smaller than `w`.
 - ▶ Returns 0 if `v` is equal to `w`.
- ▶ Corresponds to [natural ordering](#).

How to make your class T comparable?

1. Implement Comparable<T> interface.
2. Implement compareTo(T that) method to compare this T object to that based on natural ordering.

Comparator

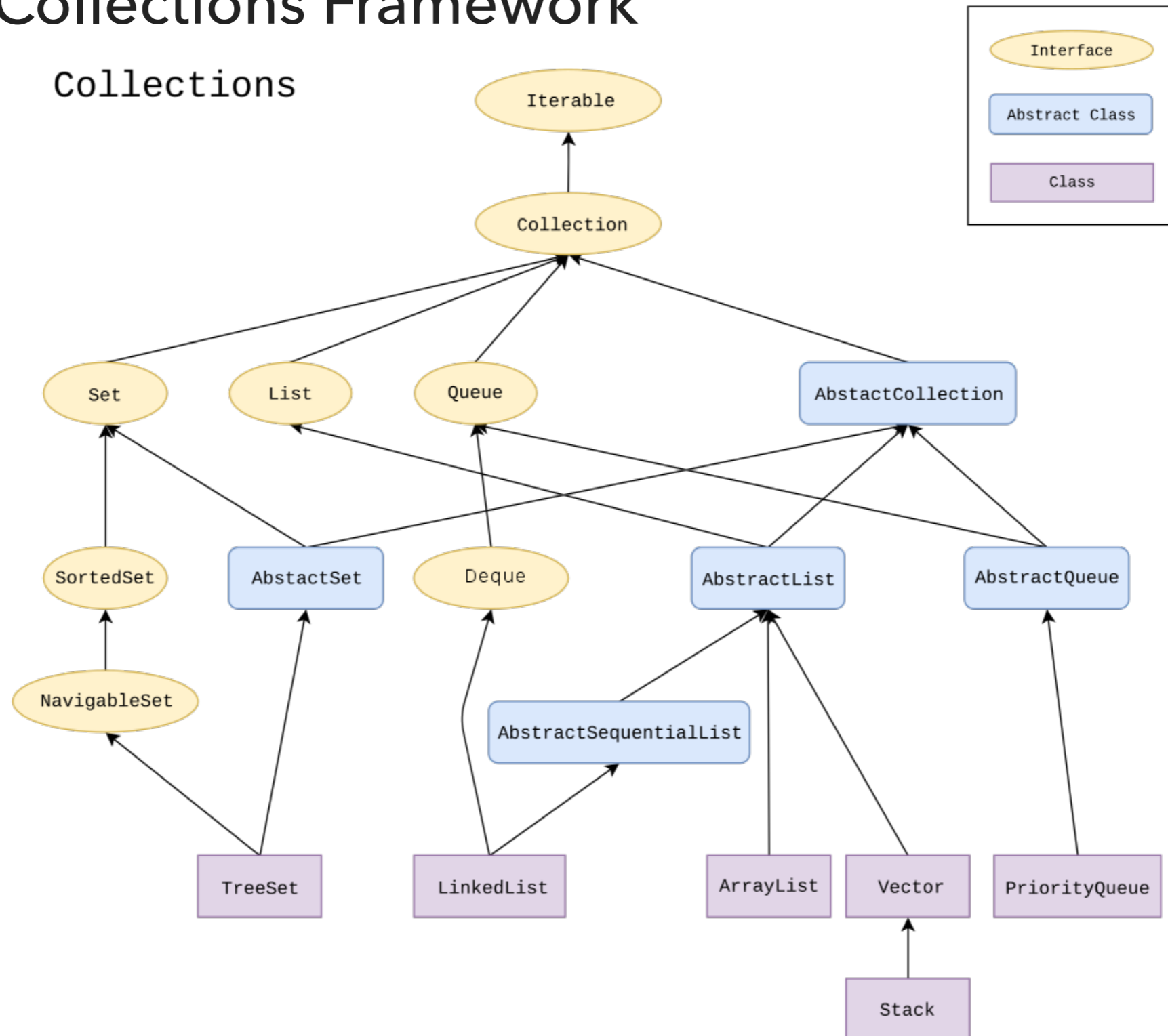
- ▶ Sometimes the natural ordering is **not** the type of ordering we want.
- ▶ Comparator is an interface which allows us to **dictate what kind of ordering** we want by implementing the method:
`public int compare(T this, T that)`
- ▶ Implement it so that `compare(v, w)`:
 - ▶ Returns >0 if v is greater than w .
 - ▶ Returns <0 if v is smaller than w .
 - ▶ Returns 0 if v is equal to w .

How to define an alternative ordering for your class T?

1. Make a new class that implements `Comparator<T>` interface.
2. Implement `compare(T t1, T t2)` method to compare `t1` object to `t2` based on an alternative ordering.
3. Alternatively, implement an anonymous inner class:

```
public static Comparator<T> nameOfComparator = new Comparator<T>()
{
    @Override // indicates method overriding the superclass' method
    public int compare(T t1, T t2) {
        {
            //return something;
        }
    }
};
```

The Java Collections Framework



Alternative sorting of Collections

- ▶ Collections class contains:
 - ▶ `static <T> void sort(List<T> list, Comparator<? super T> c)`
- ▶ `Collections.sort(list, someComparator);`
 - ▶ `Collections.sort(list, new ExternalComparatorClass());` or:
 - ▶ `Collections.sort(list, T.InnerAnonymousClass);`
 - ▶ If list's elements do not implement `Comparable` or cannot be compared with `Comparator`, throw `ClassCastException`.

Example: Natural and alternative sorting for Employees

<https://github.com/pomonacs622021fa/LectureCode/blob/main/Lecture11/Employee.java>

Lecture 12: Insertion Sort & Mergesort

- ▶ Insertion sort
- ▶ Comparators
- ▶ Mergesort

Lecture 12: Insertion Sort & Mergesort

- ▶ Mergesort

MERGESORT

input	M	E	R	G	E	S	O	R	T	E	X	A	M	P	L	E	
sort left half	E	E	G	M	O	R	R	S		T	E	X	A	M	P	L	E
sort right half	E	E	G	M	O	R	R	S		A	E	E	L	M	P	T	X
merge results	A	E	E	E	E	G	L	M	M	O	P	R	R	S	T	X	

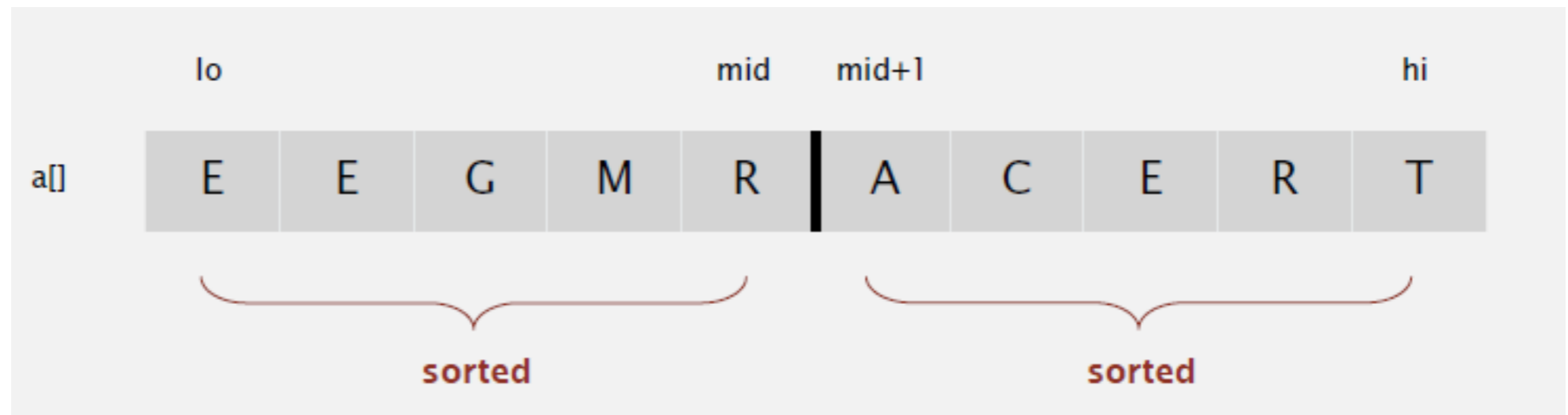
Basics

Mergesort overview

- ▶ Invented by John von Neumann in 1945
- ▶ Algorithm sketch:
 - ▶ Divide array into two halves.
 - ▶ Recursively sort each half.
 - ▶ Merge the two halves



Merging two already sorted halves into one sorted array



Copy to auxiliary array

Merging Example - copying to auxiliary array

Array a

A	G	L	O	R	H	I	M	S	T
0	1	2	3	4	5	6	7	8	9

Array aux

0	1	2	3	4	5	6	7	8	9

Merging Example - copying to auxiliary array

Array a

A	G	L	O	R	H	I	M	S	T
0	1	2	3	4	5	6	7	8	9

Array aux

A	G	L	O	R	H	I	M	S	T
0	1	2	3	4	5	6	7	8	9

Merging Example - copy elements back to original array in order

Maintain 3 indices: i , j , k

Array a

A	G	L	O	R	H	I	M	S	T
0	1	2	3	4	5	6	7	8	9

k

Array aux

A	G	L	O	R	H	I	M	S	T
0	1	2	3	4	5	6	7	8	9

i

j

Merging Example - copy elements back to original array in order

Compare minimum in each subarray

Array a (sorted result)

A	G	L	O	R	H	I	M	S	T
0	1	2	3	4	5	6	7	8	9

k

Array aux

A	G	L	O	R	H	I	M	S	T
0	1	2	3	4	5	6	7	8	9

i

j

Merging Example - copy elements back to original array in order

Copy smaller element back to a, increment indices i and j

Array a (sorted result)

A	G	L	O	R	H	I	M	S	T
0	1	2	3	4	5	6	7	8	9

k

Array aux

A	G	L	O	R	H	I	M	S	T
0	1	2	3	4	5	6	7	8	9

i

j

Merging Example - copy elements back to original array in order

Compare

Array a (sorted result)

A	G	L	O	R	H	I	M	S	T
0	1	2	3	4	5	6	7	8	9

k

Array aux

A	G	L	O	R	H	I	M	S	T
0	1	2	3	4	5	6	7	8	9

i

j

Merging Example - copy elements back to original array in order

Compare

Array a (sorted result)

A	G	L	O	R	H	I	M	S	T
0	1	2	3	4	5	6	7	8	9

k

Array aux

A	G	L	O	R	H	I	M	S	T
0	1	2	3	4	5	6	7	8	9

i

j

Merging Example - copy elements back to original array in order

Compare

Array a (sorted result)

A	G	L	O	R	H	I	M	S	T
0	1	2	3	4	5	6	7	8	9

k

Array aux

A	G	L	O	R	H	I	M	S	T
0	1	2	3	4	5	6	7	8	9

i

j

Merging Example - copy elements back to original array in order

Compare

Array a (sorted result)

A	G	H	O	R	H	I	M	S	T
0	1	2	3	4	5	6	7	8	9

k

Array aux

A	G	L	O	R	H	I	M	S	T
0	1	2	3	4	5	6	7	8	9

i

j

Merging Example - copy elements back to original array in order

Compare

Array a (sorted result)

A	G	H	O	R	H	I	M	S	T
0	1	2	3	4	5	6	7	8	9

k

Array aux

A	G	L	O	R	H	I	M	S	T
0	1	2	3	4	5	6	7	8	9

i

j

Merging Example - copy elements back to original array in order

Compare

Array a (sorted result)

A	G	H	I	R	H	I	M	S	T
0	1	2	3	4	5	6	7	8	9

k

Array aux

A	G	L	O	R	H	I	M	S	T
0	1	2	3	4	5	6	7	8	9

i

j

Merging Example - copy elements back to original array in order

Compare

Array a (sorted result)

A	G	H	I	L	H	I	M	S	T
0	1	2	3	4	5	6	7	8	9

k

Array aux

A	G	L	O	R	H	I	M	S	T
0	1	2	3	4	5	6	7	8	9

i

j

Merging Example - copy elements back to original array in order

Compare

Array a (sorted result)

A	G	H	I	L	M	I	M	S	T
0	1	2	3	4	5	6	7	8	9

k

Array aux

A	G	L	O	R	H	I	M	S	T
0	1	2	3	4	5	6	7	8	9

i

j

Merging Example - copy elements back to original array in order

Compare

Array a (sorted result)

A	G	H	I	L	M	O	M	S	T
0	1	2	3	4	5	6	7	8	9

k

Array aux

A	G	L	O	R	H	I	M	S	T
0	1	2	3	4	5	6	7	8	9

i

j

Merging Example - copy elements back to original array in order

Compare

Array a (sorted result)

A	G	H	I	L	M	O	R	S	T
0	1	2	3	4	5	6	7	8	9

k

Array aux

A	G	L	O	R	H	I	M	S	T
0	1	2	3	4	5	6	7	8	9

i

j

Merging Example - copy elements back to original array in order

Compare

Array a (sorted result)

A	G	H	I	L	M	O	R	S	T
0	1	2	3	4	5	6	7	8	9

k

Array aux

A	G	L	O	R	H	I	M	S	T
0	1	2	3	4	5	6	7	8	9

i

j

Merging Example - copy elements back to original array in order

Done

Array a (sorted result)

A	G	H	I	L	M	O	R	S	T
0	1	2	3	4	5	6	7	8	9

k

Array aux

A	G	L	O	R	H	I	M	S	T
0	1	2	3	4	5	6	7	8	9

i

j

Merging two already sorted halves into one sorted array

```
private static void merge(Comparable[] a, Comparable[] aux, int lo, int mid, int hi) {
    for (int k = lo; k <= hi; k++) // copy to aux array
        aux[k] = a[k];

    int i = lo, j = mid+1; // lo and mid+1 are the start of the 2 sorted halves

    for (int k = lo; k <= hi; k++) {
        if (i > mid) //ran out of elements in the left subarray
            a[k] = aux[j++];
        else if (j > hi) //ran out of elements in the right subarray
            a[k] = aux[i++];
        else if (less(aux[j], aux[i])) // Compares left and right subarray
            a[k] = aux[j++];
        else
            a[k] = aux[i++];
    }
}
```

2.2 MERGING DEMO



<http://algs4.cs.princeton.edu>

Practice time

How many calls does `merge()` make to `less()` in order to merge two already sorted subarrays, each of length $n/2$ into a sorted array of length n ?

A. $\sim 1/4n$ to $\sim 1/2n$

B. $\sim 1/2n$

C. $\sim 1/2n$ to n

D. $\sim n$

```
private static void merge(Comparable[] a, Comparable[] aux, int lo, int hi) {
    for (int k = lo; k <= hi; k++)
        aux[k] = a[k];

    int i = lo, j = mid+1;
    for (int k = lo; k <= hi; k++) {
        if (i > mid) //ran out of elements in the left subarray
            a[k] = aux[j++];
        else if (j > hi) //ran out of elements in the right subarray
            a[k] = aux[i++];
        else if (less(aux[j], aux[i]))
            a[k] = aux[j++];
        else
            a[k] = aux[i++];
    }
}
```

Answer

How many calls does `merge()` make to `less()` in order to merge two already sorted subarrays, each of length $n/2$ into a sorted array of length n ?

C. $\sim 1/2n$ to n , that is at most $n - 1$ or $O(n)$

Best case example

Merging `[1,2,3]` and `[4,5,6]` requires 3 calls to `less()`
(1 with 4, 2 with 4, 3 with 4).

Worst case example

Merging `[1,3,5]` and `[2, 4, 6]` requires 5 calls to `less()`
(1 with 2, 2 with 3, 3 with 4, 4 with 5, 5 with 6)

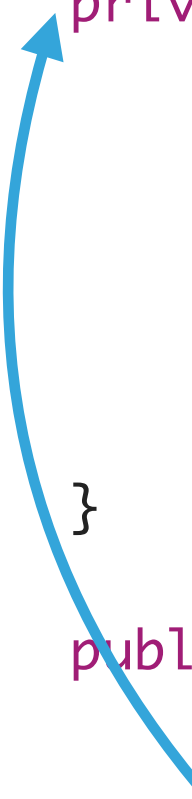
Mergesort - the quintessential example of divide-and-conquer

```
private static void sort(Comparable[] a, Comparable[] aux, int lo, int hi) {
    if (hi <= lo)
        return;
    int mid = lo + (hi - lo) / 2;    // Computes midpoint
    sort(a, aux, lo, mid);          // Sort the first half
    sort(a, aux, mid+1, hi);        // Sort the second half
    merge(a, aux, lo, mid, hi);     // Merge the 2 halves
}

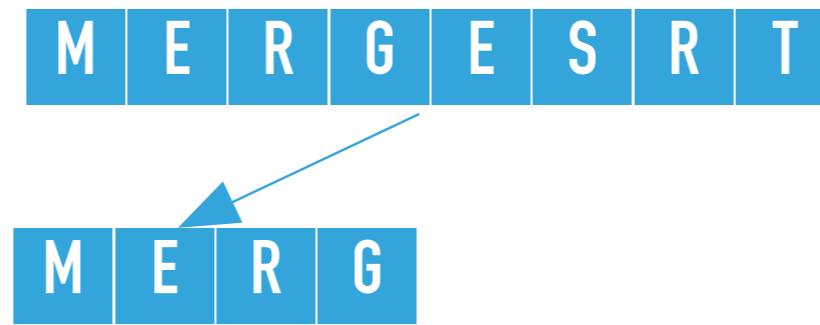
public static void sort(Comparable[] a) {
    Comparable[] aux = new Comparable[a.length];    // Create aux array
    sort(a, aux, 0, a.length - 1);                // Recursively call sort
}
```

```
private static void sort(Comparable[] a, Comparable[] aux, int lo, int hi) {
    if (hi <= lo)
        return;
    int mid = lo + (hi - lo) / 2;
    sort(a, aux, lo, mid);
    sort(a, aux, mid+1, hi);
    merge(a, aux, lo, mid, hi);
}

public static void sort(Comparable[] a) {
    Comparable[] aux = new Comparable[a.length];
    sort(a, aux, 0, a.length - 1);
}
```

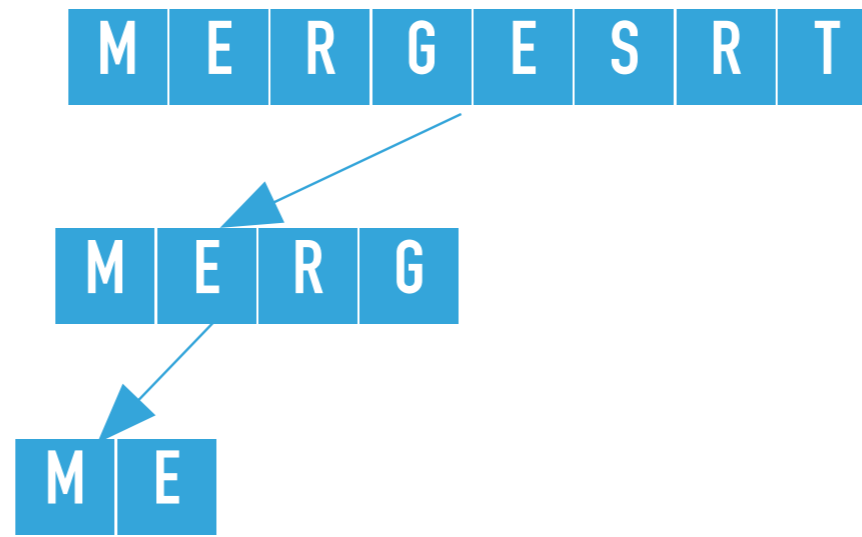


`sort([M, E, R, G, E, S, R, T])` calls
`sort([M, E, R, G, E, S, R, T], [null, null, null, null, null, null, null, null], 0, 7)` where the array of `nulls` is the auxiliary array, `lo = 0` and `hi = 7`.



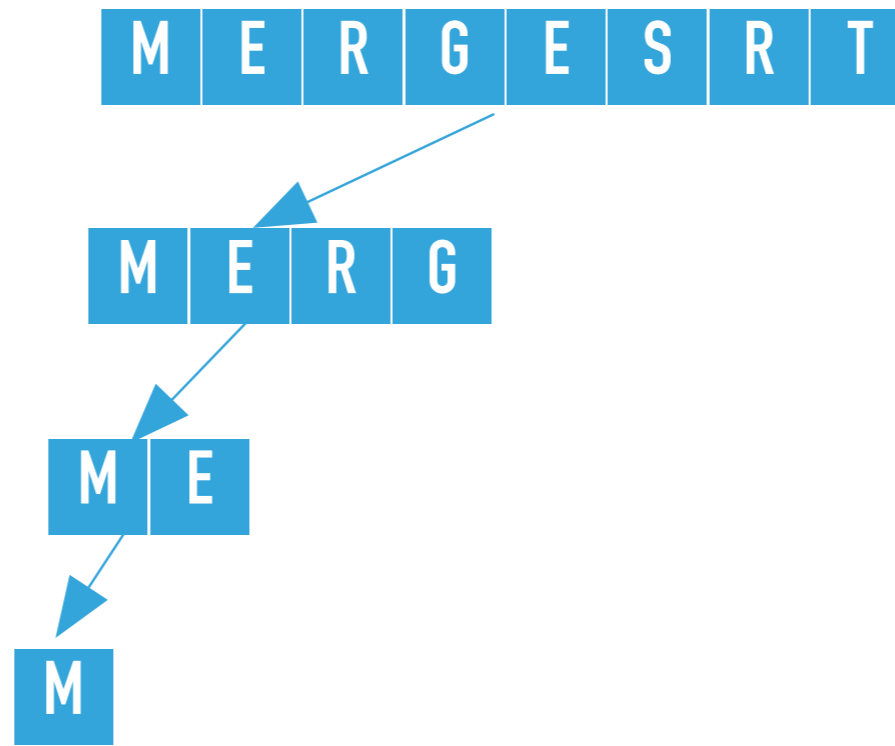
```
private static void sort(Comparable[] a, Comparable[] aux, int lo, int hi) {  
    if (hi <= lo)  
        return;  
    int mid = lo + (hi - lo) / 2;  
    sort(a, aux, lo, mid);  
    sort(a, aux, mid+1, hi);  
    merge(a, aux, lo, mid, hi);  
}
```

`sort([M, E, R, G, E, S, R, T], [null, null, null, null, null, null, null, null], 0, 7)` calculates the `mid = 3` and calls recursively `sort` on the left subarray, that is `sort([M, E, R, G, E, S, R, T], [null, null, null, null, null, null, null, null], 0, 3)`, where `lo = 0, hi = 3`



```
private static void sort(Comparable[] a, Comparable[] aux, int lo, int hi) {
    if (hi <= lo)
        return;
    int mid = lo + (hi - lo) / 2;
    sort(a, aux, lo, mid);
    sort(a, aux, mid+1, hi);
    merge(a, aux, lo, mid, hi);
}
```

`sort([M, E, R, G, E, S, R, T], [null, null, null, null, null, null, null, null], 0, 3)` calculates the `mid = 1` and calls recursively `sort` on the left subarray, that is `sort([M, E, R, G, E, S, R, T], [null, null, null, null, null, null, null, null], 0, 1)`, where `lo = 0, hi = 1`

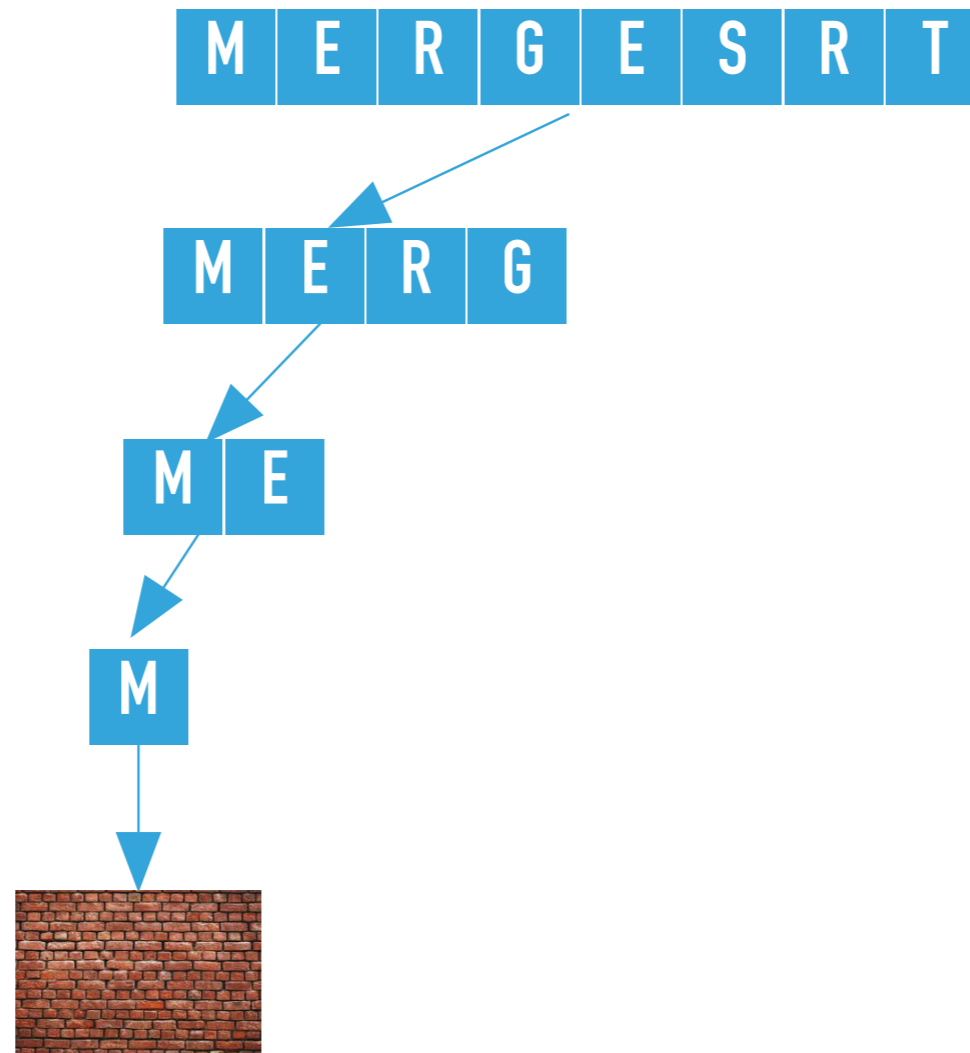


```

private static void sort(Comparable[] a, Comparable[] aux, int lo, int hi) {
    if (hi <= lo)
        return;
    int mid = lo + (hi - lo) / 2;
    sort(a, aux, lo, mid);
    sort(a, aux, mid+1, hi);
    merge(a, aux, lo, mid, hi);
}

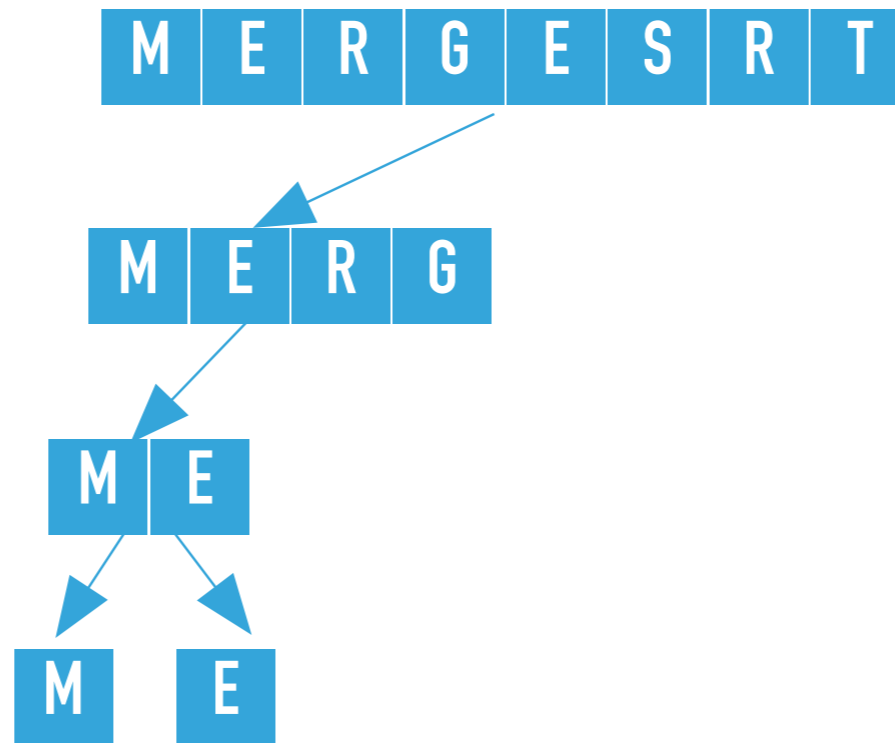
```

`sort([M, E, R, G, E, S, R, T], [null, null, null, null, null, null, null, null], 0, 1)` calculates the `mid = 0` and calls recursively `sort` on the left subarray, that is `sort([M, E, R, G, E, S, R, T], [null, null, null, null, null, null, null, null], 0, 0)`, where `lo = 0, hi = 0`



```
private static void sort(Comparable[] a, Comparable[] aux, int lo, int hi) {
    if (hi <= lo)
        return;
    int mid = lo + (hi - lo) / 2;
    sort(a, aux, lo, mid);
    sort(a, aux, mid+1, hi);
    merge(a, aux, lo, mid, hi);
}
```

`sort([M, E, R, G, E, S, R, T], [null, null, null, null, null, null, null, null], 0, 0)` finds `hi <= lo` and returns.

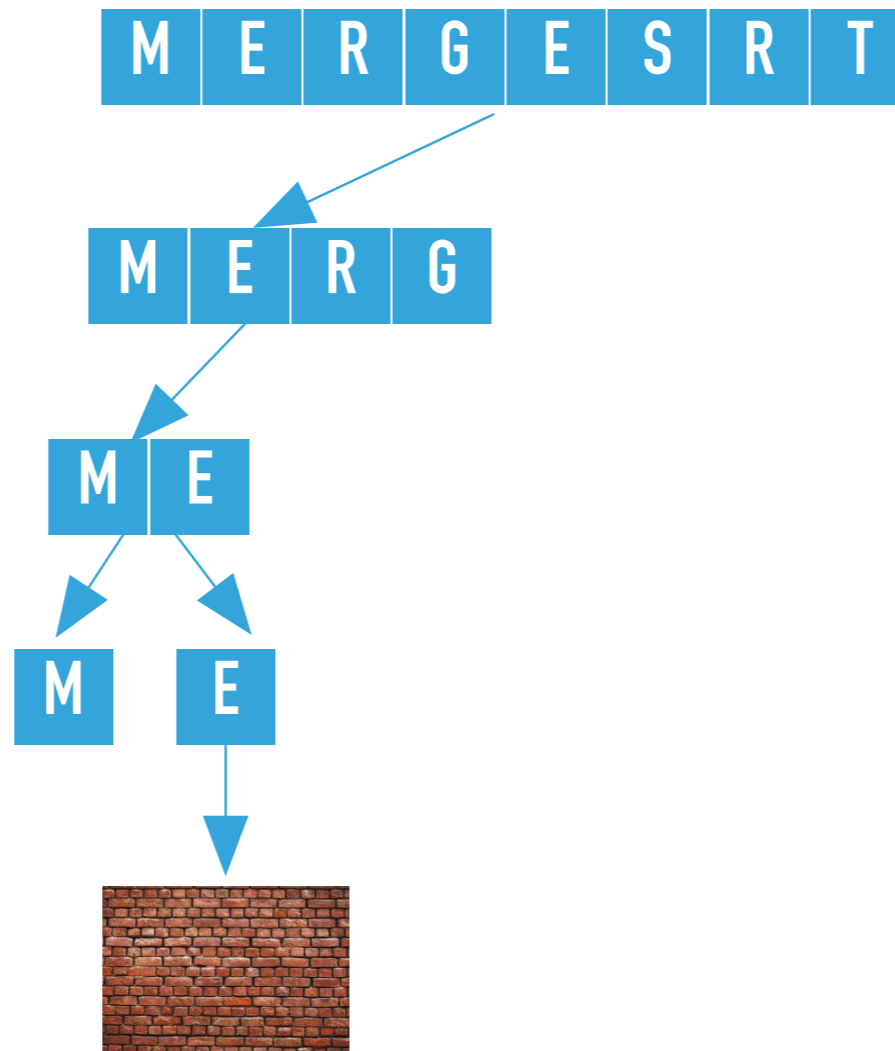


```

private static void sort(Comparable[] a, Comparable[] aux, int lo, int hi) {
    if (hi <= lo)
        return;
    int mid = lo + (hi - lo) / 2;
    sort(a, aux, lo, mid);
    sort(a, aux, mid+1, hi);
    merge(a, aux, lo, mid, hi);
}

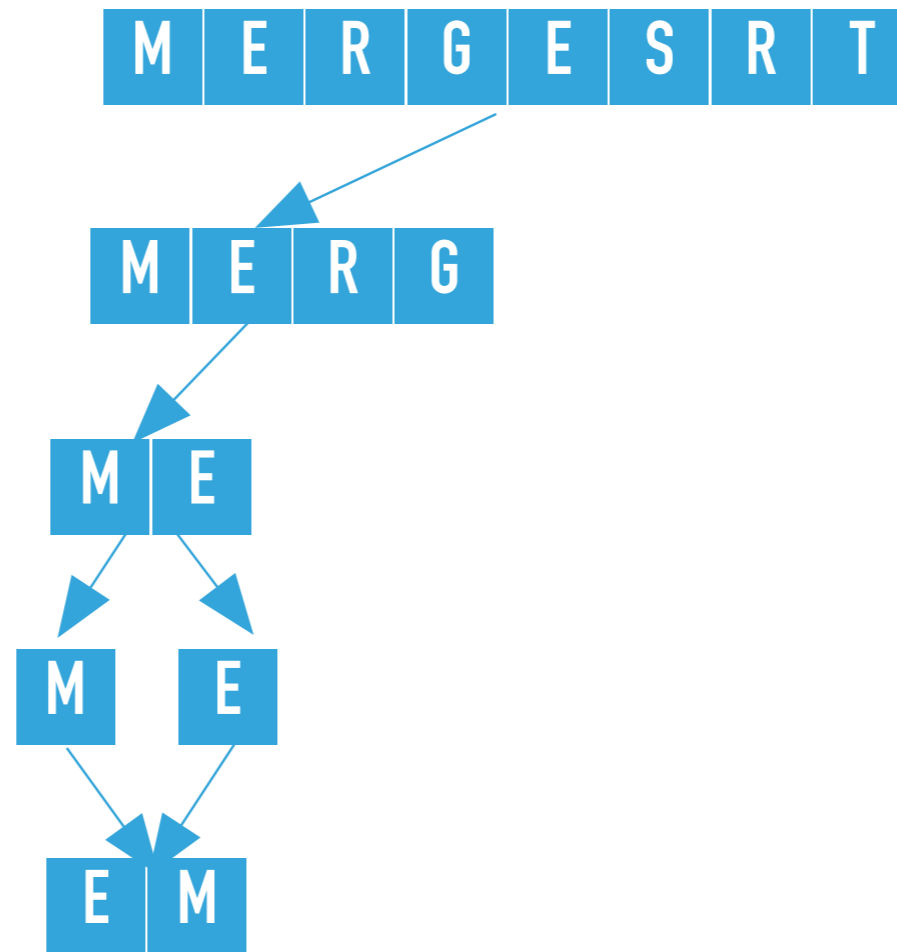
```

sort([M, E, R, G, E, S, R, T], [null, null, null, null, null, null, null, null], 0, 1) calls recursively sort on the right subarray, that is sort([M, E, R, G, E, S, R, T], [null, null, null, null, null, null, null, null], 1, 1), where lo = 1, hi = 1



```
private static void sort(Comparable[] a, Comparable[] aux, int lo, int hi) {
    if (hi <= lo)
        return;
    int mid = lo + (hi - lo) / 2;
    sort(a, aux, lo, mid);
    sort(a, aux, mid+1, hi);
    merge(a, aux, lo, mid, hi);
}
```

`sort([M, E, R, G, E, S, R, T], [null, null, null, null, null, null, null, null], 1, 1)` finds `hi <= lo` and returns.

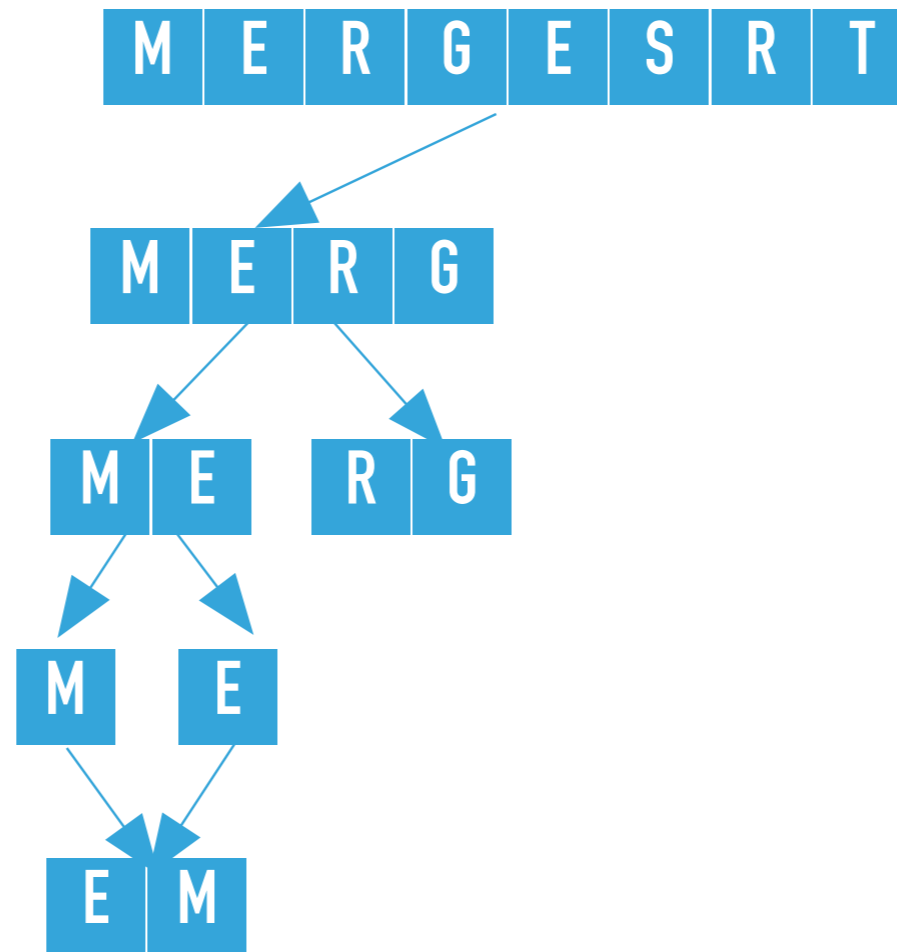


```

private static void sort(Comparable[] a, Comparable[] aux, int lo, int hi) {
    if (hi <= lo)
        return;
    int mid = lo + (hi - lo) / 2;
    sort(a, aux, lo, mid);
    sort(a, aux, mid+1, hi);
    merge(a, aux, lo, mid, hi);
}

```

`sort([M, E, R, G, E, S, R, T], [null, null, null, null, null, null, null, null], 0, 1)` merges the two subarrays that is calls `merge([M, E, R, G, E, S, R, T], [null, null, null, null, null, null, null, null], 0, 0, 1)`, where `lo = 0, mid = 0, and hi = 1`. The resulting partially sorted array is `[E, M, R, G, E, S, R T]`.

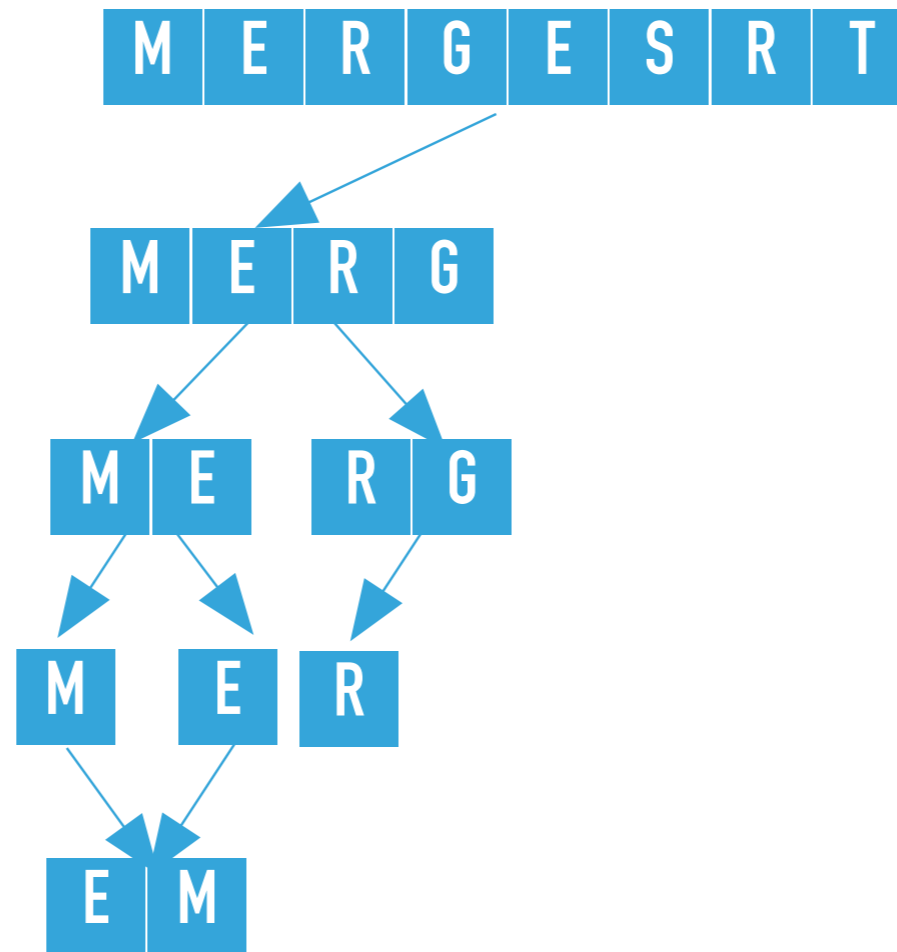


```

private static void sort(Comparable[] a, Comparable[] aux, int lo, int hi) {
    if (hi <= lo)
        return;
    int mid = lo + (hi - lo) / 2;
    sort(a, aux, lo, mid);
    sort(a, aux, mid+1, hi);
    merge(a, aux, lo, mid, hi);
}

```

`sort([E, M, R, G, E, S, R, T], [M, E, null, null, null, null, null, null], 0, 3)` calls recursively `sort` on the right subarray, that is `sort([E, M, R, G, E, S, R, T], [M, E, null, null, null, null, null, null], 2, 3)`, where `lo = 2, hi = 3`

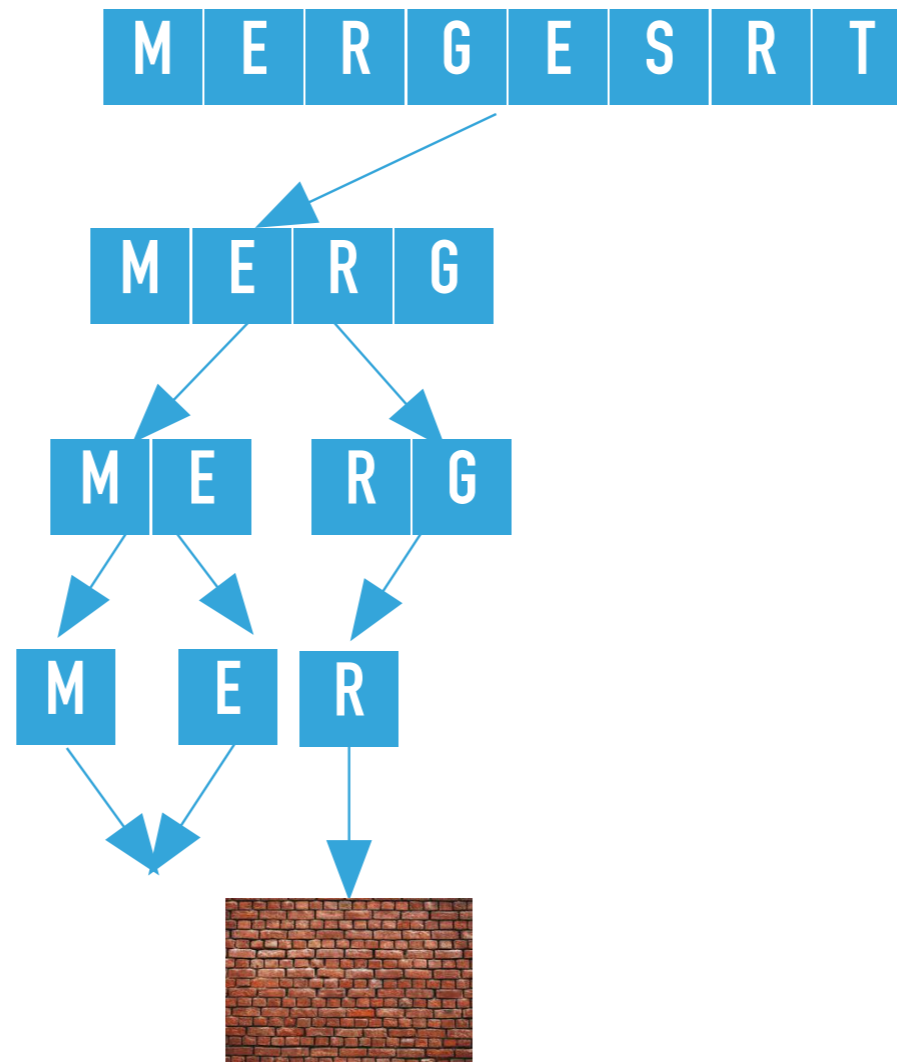


```

private static void sort(Comparable[] a, Comparable[] aux, int lo, int hi) {
    if (hi <= lo)
        return;
    int mid = lo + (hi - lo) / 2;
    sort(a, aux, lo, mid);
    sort(a, aux, mid+1, hi);
    merge(a, aux, lo, mid, hi);
}

```

`sort([E, M, R, G, E, S, R, T], [M, E, null, null, null, null, null, null], 2, 3)`
 calculates the `mid = 2` and calls recursively `sort` on the left subarray, that is `sort([E, M, R, G, E, S, R, T], [M, E, null, null, null, null, null, null], 2, 2)`, where `lo = 2, hi = 2`

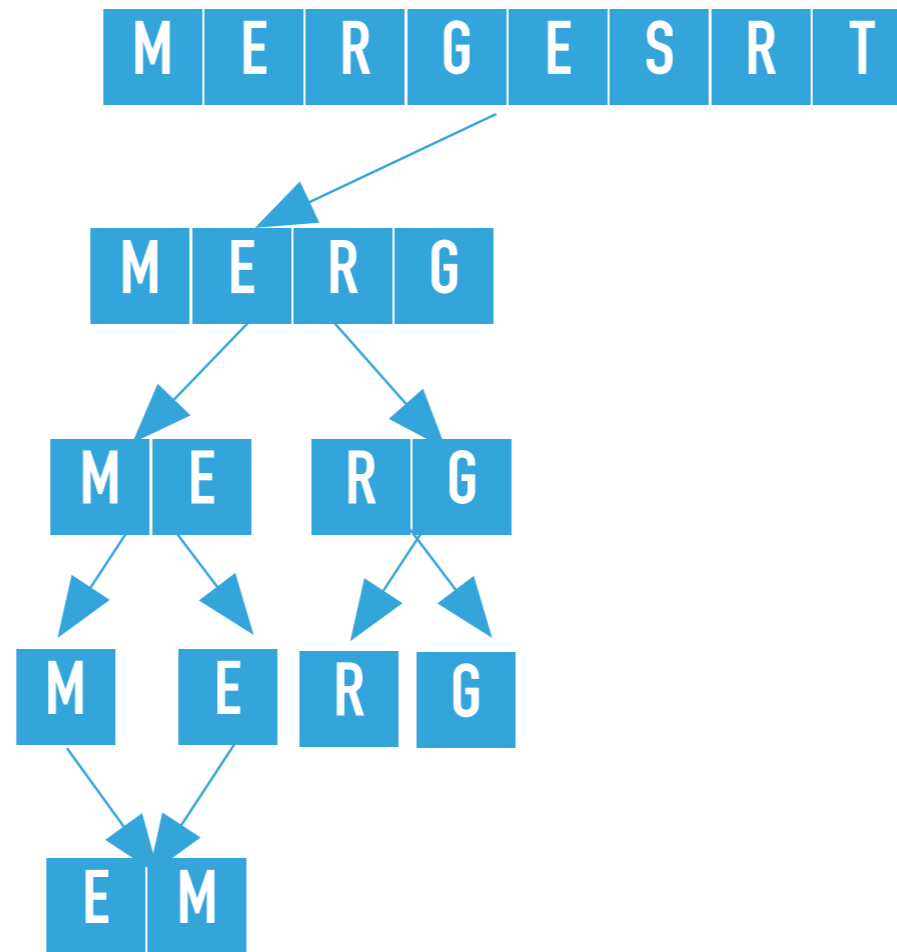


```

private static void sort(Comparable[] a, Comparable[] aux, int lo, int hi) {
    if (hi <= lo)
        return;
    int mid = lo + (hi - lo) / 2;
    sort(a, aux, lo, mid);
    sort(a, aux, mid+1, hi);
    merge(a, aux, lo, mid, hi);
}

```

sort([E, M, R, G, E, S, R, T], [M, E, null, null, null, null, null, null], 2, 2) finds $hi \leq lo$ and returns.

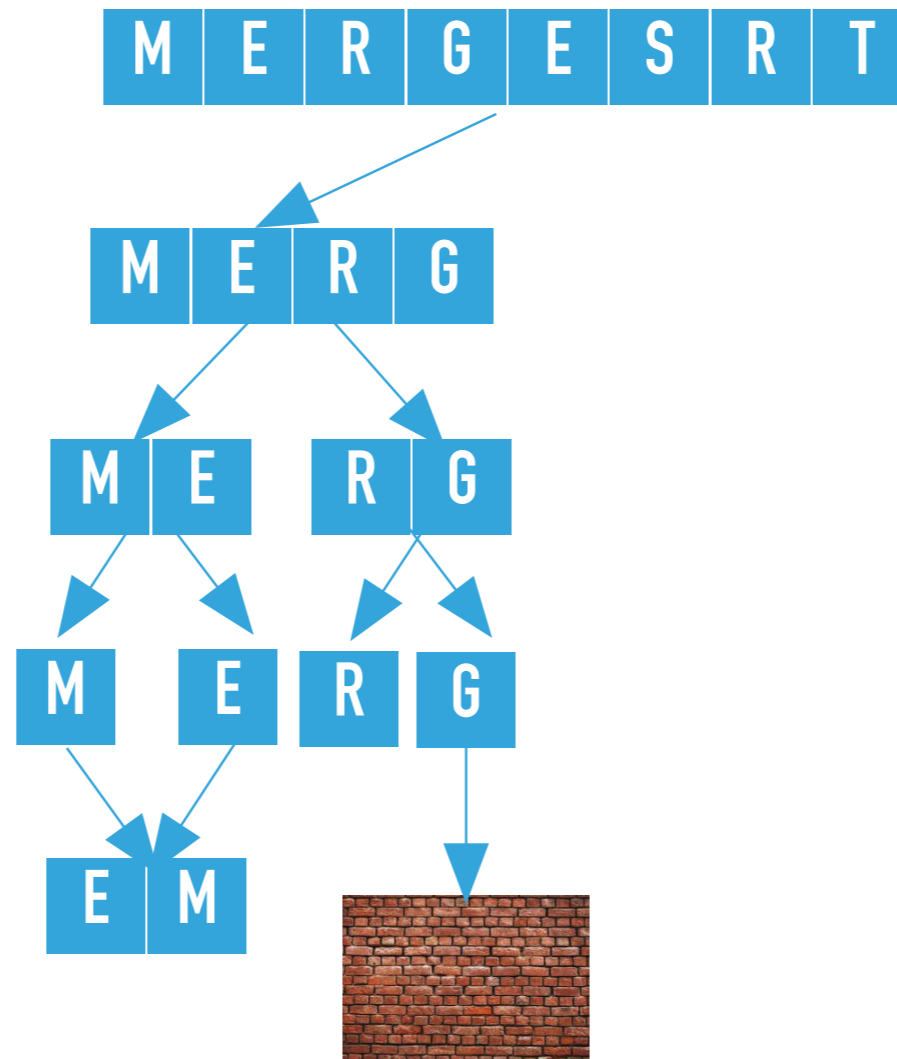


```

private static void sort(Comparable[] a, Comparable[] aux, int lo, int hi) {
    if (hi <= lo)
        return;
    int mid = lo + (hi - lo) / 2;
    sort(a, aux, lo, mid);
    sort(a, aux, mid+1, hi);
    merge(a, aux, lo, mid, hi);
}

```

sort([E, M, R, G, E, S, R, T], [M, E, null, null, null, null, null, null], 2, 3) calls recursively sort on the right subarray, that is sort([E, M, R, G, E, S, R, T], [M, E, null, null, null, null, null, null], 3, 3), where lo = 3, hi = 3

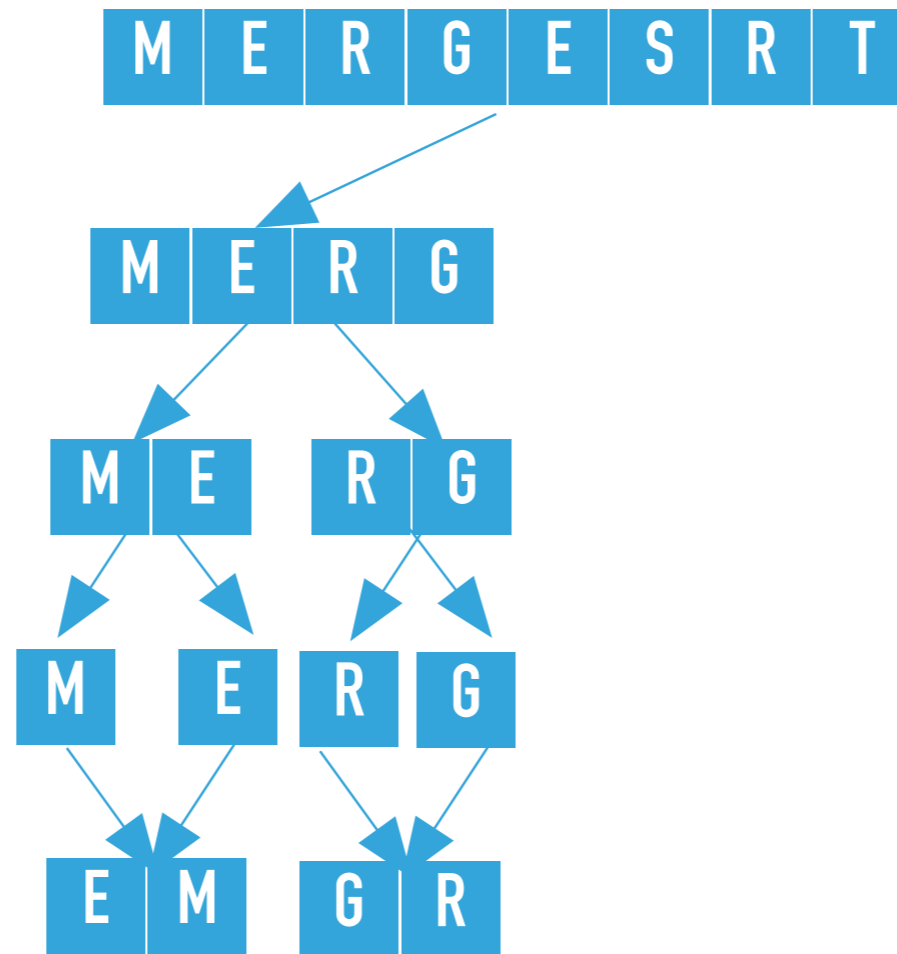


```

private static void sort(Comparable[] a, Comparable[] aux, int lo, int hi) {
    if (hi <= lo)
        return;
    int mid = lo + (hi - lo) / 2;
    sort(a, aux, lo, mid);
    sort(a, aux, mid+1, hi);
    merge(a, aux, lo, mid, hi);
}

```

`sort([E, M, R, G, E, S, R, T], [M, E, null, null, null, null, null, null], 3, 3)` finds `hi <= lo` and returns.

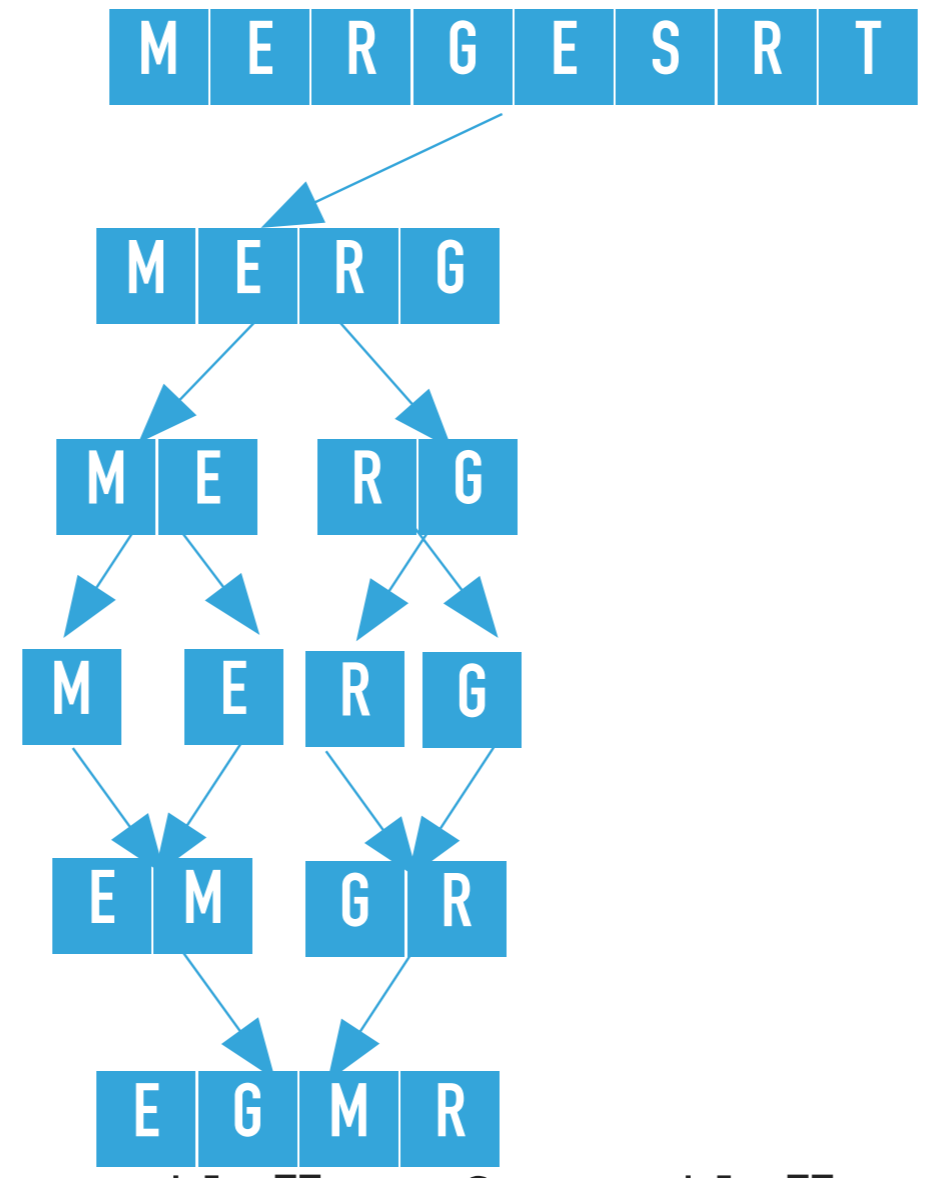


```

private static void sort(Comparable[] a, Comparable[] aux, int lo, int hi) {
    if (hi <= lo)
        return;
    int mid = lo + (hi - lo) / 2;
    sort(a, aux, lo, mid);
    sort(a, aux, mid+1, hi);
    merge(a, aux, lo, mid, hi);
}

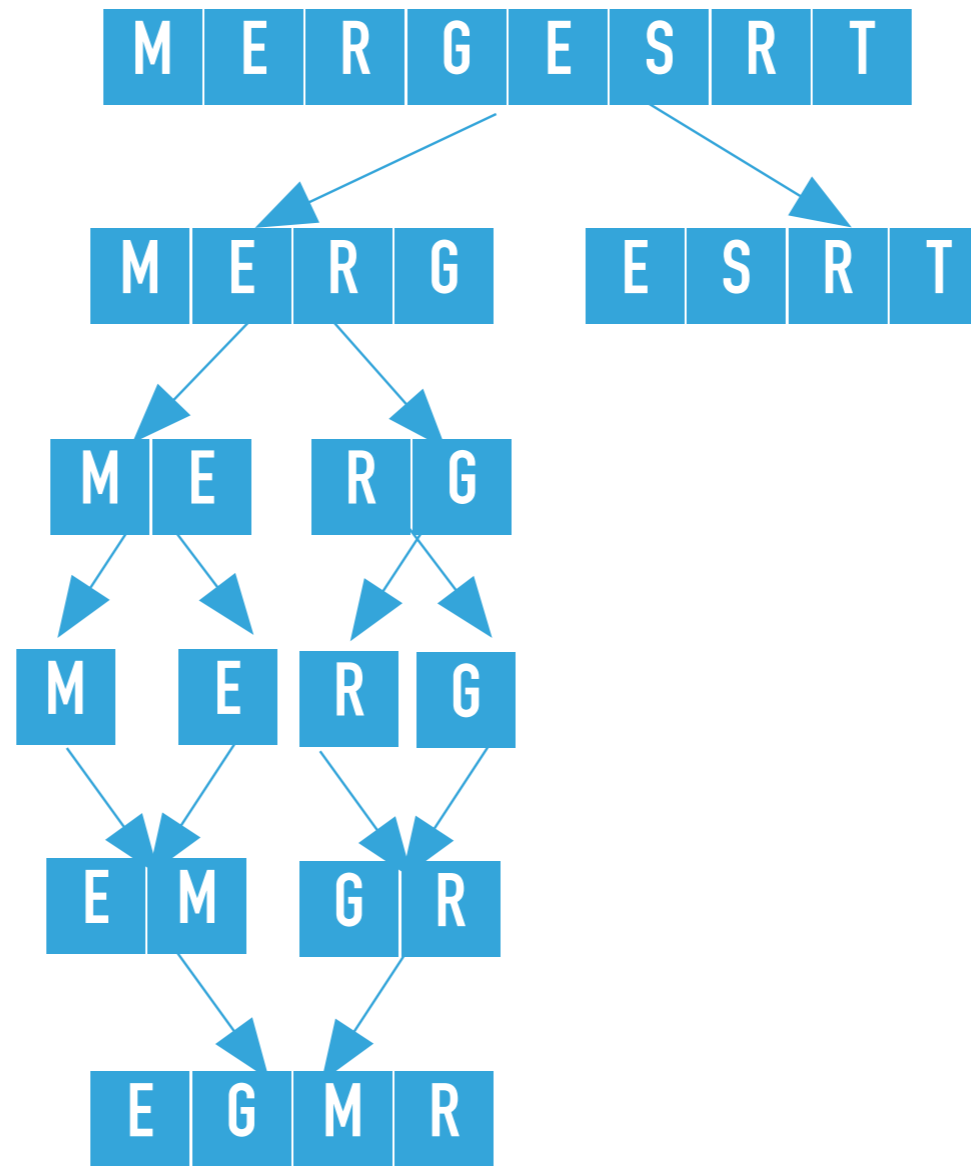
```

`sort([E, M, R, G, E, S, R, T], [M, E, null, null, null, null, null, null], 2, 3)`
 merges the two subarrays that is calls `merge([E, M, R, G, E, S, R, T], [M, E, null, null, null, null, null, null], 2, 2, 3)`, where `lo = 2, mid = 2, and hi = 3`. The resulting partially sorted array is `[E, M, G, R, E, S, R T]`.



```
private static void sort(Comparable[] a, Comparable[] aux, int lo, int hi) {
    if (hi <= lo)
        return;
    int mid = lo + (hi - lo) / 2;
    sort(a, aux, lo, mid);
    sort(a, aux, mid+1, hi);
    merge(a, aux, lo, mid, hi);
}
```

sort([E, M, G, R, E, S, R, T], [M, E, R, G, null, null, null, null], 0, 3) merges the two subarrays that is calls merge([E, M, G, R, E, S, R, T], [M, E, R, G, null, null, null, null], 0, 1, 3), where lo = 0, mid = 1, and hi = 3. The resulting partially sorted array is [E, G, M, R, E, S, R T].

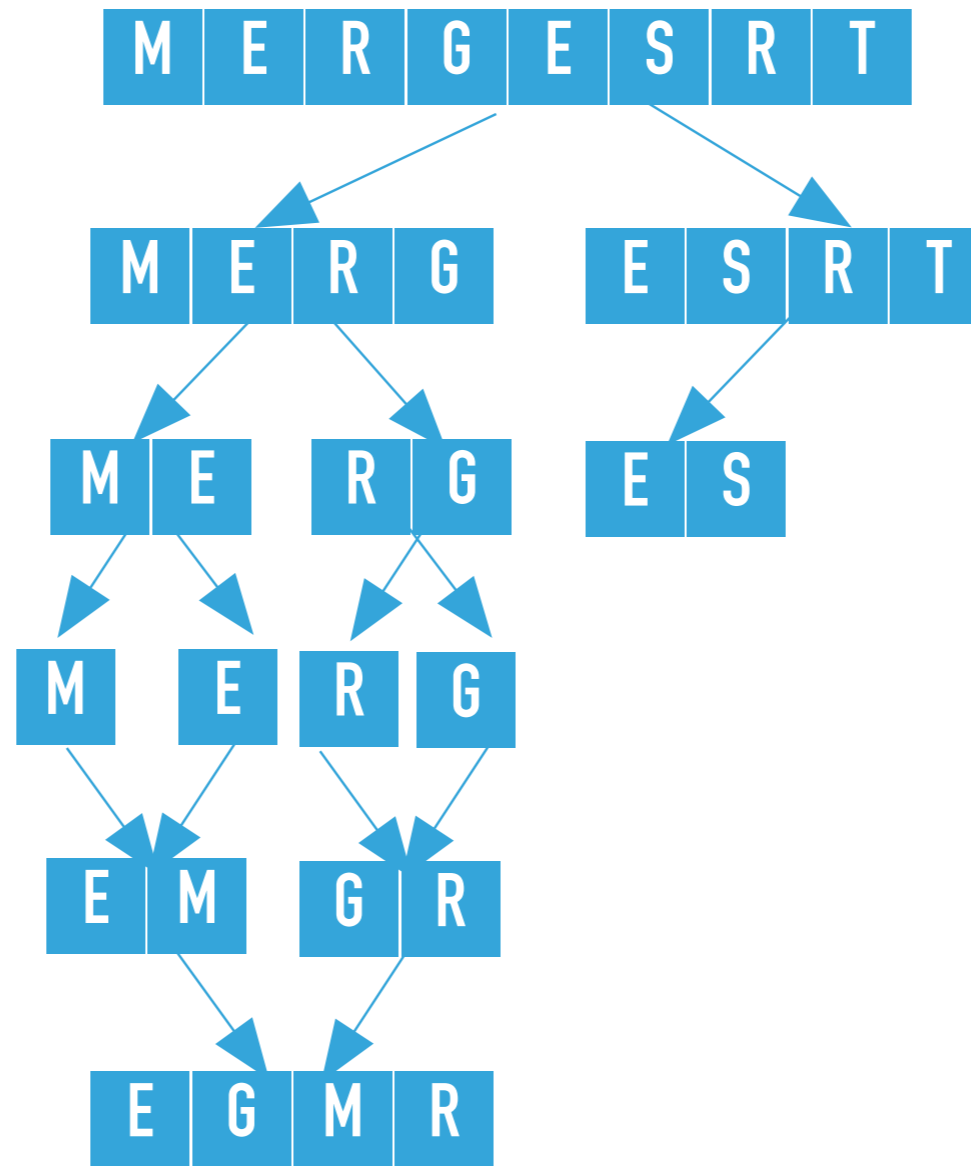


```

private static void sort(Comparable[] a, Comparable[] aux, int lo, int hi) {
    if (hi <= lo)
        return;
    int mid = lo + (hi - lo) / 2;
    sort(a, aux, lo, mid);
    sort(a, aux, mid+1, hi);
    merge(a, aux, lo, mid, hi);
}

```

sort([E, G, M, R, E, S, R, T], [E, M, G, R, null, null, null, null], 0, 7) calls recursively sort on the right subarray, that is sort([E, G, M, R, E, S, R, T], [E, M, G, R, null, null, null, null], 4, 7), where lo = 4, hi = 7

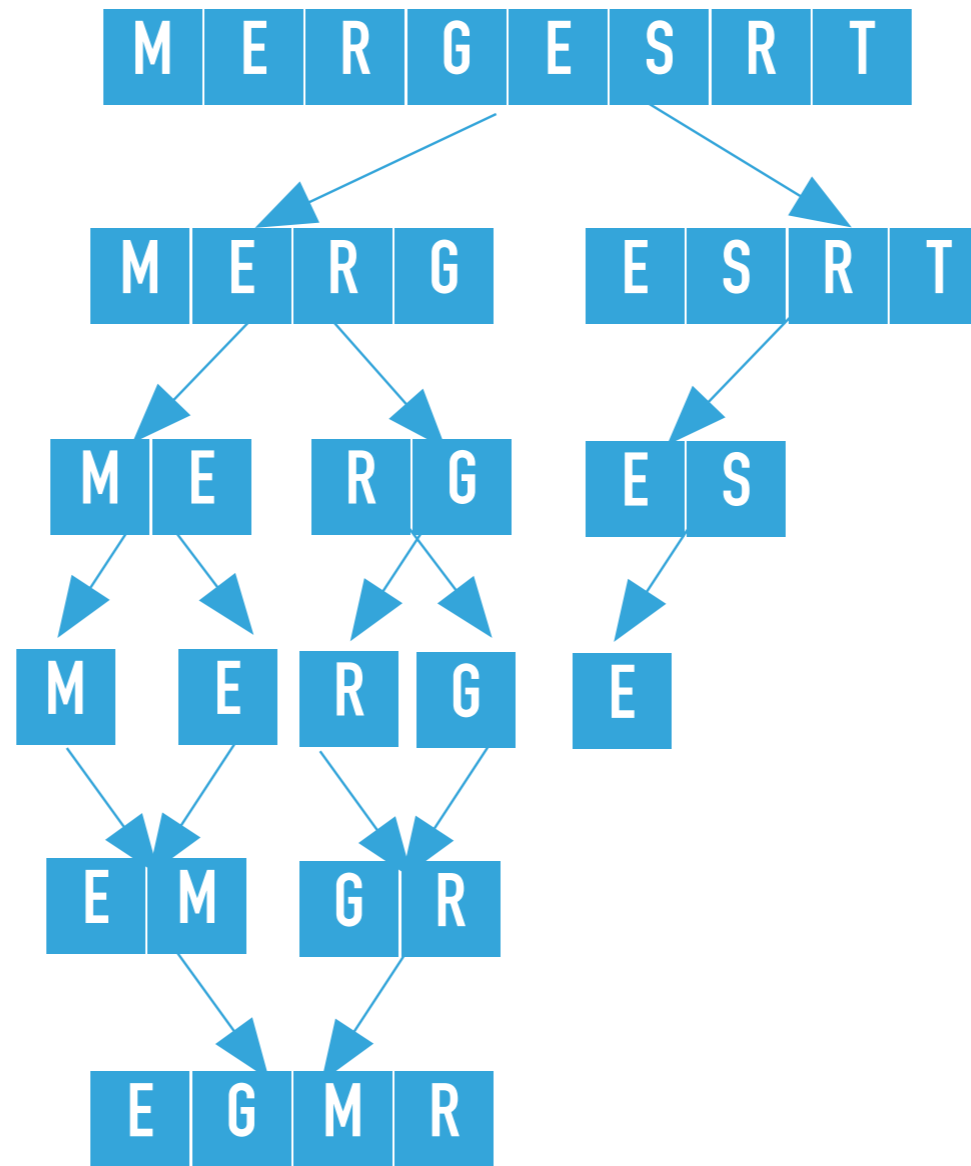


```

private static void sort(Comparable[] a, Comparable[] aux, int lo, int hi) {
    if (hi <= lo)
        return;
    int mid = lo + (hi - lo) / 2;
    sort(a, aux, lo, mid);
    sort(a, aux, mid+1, hi);
    merge(a, aux, lo, mid, hi);
}

```

`sort([E, G, M, R, E, S, R, T], [E, M, G, R, null, null, null, null], 4, 7)` calculates the `mid = 5` and calls recursively `sort` on the left subarray, that is `sort([E, G, M, R, E, S, R, T], [E, M, G, R, null, null, null, null], 4, 5)`, where `lo = 4, hi = 5`.

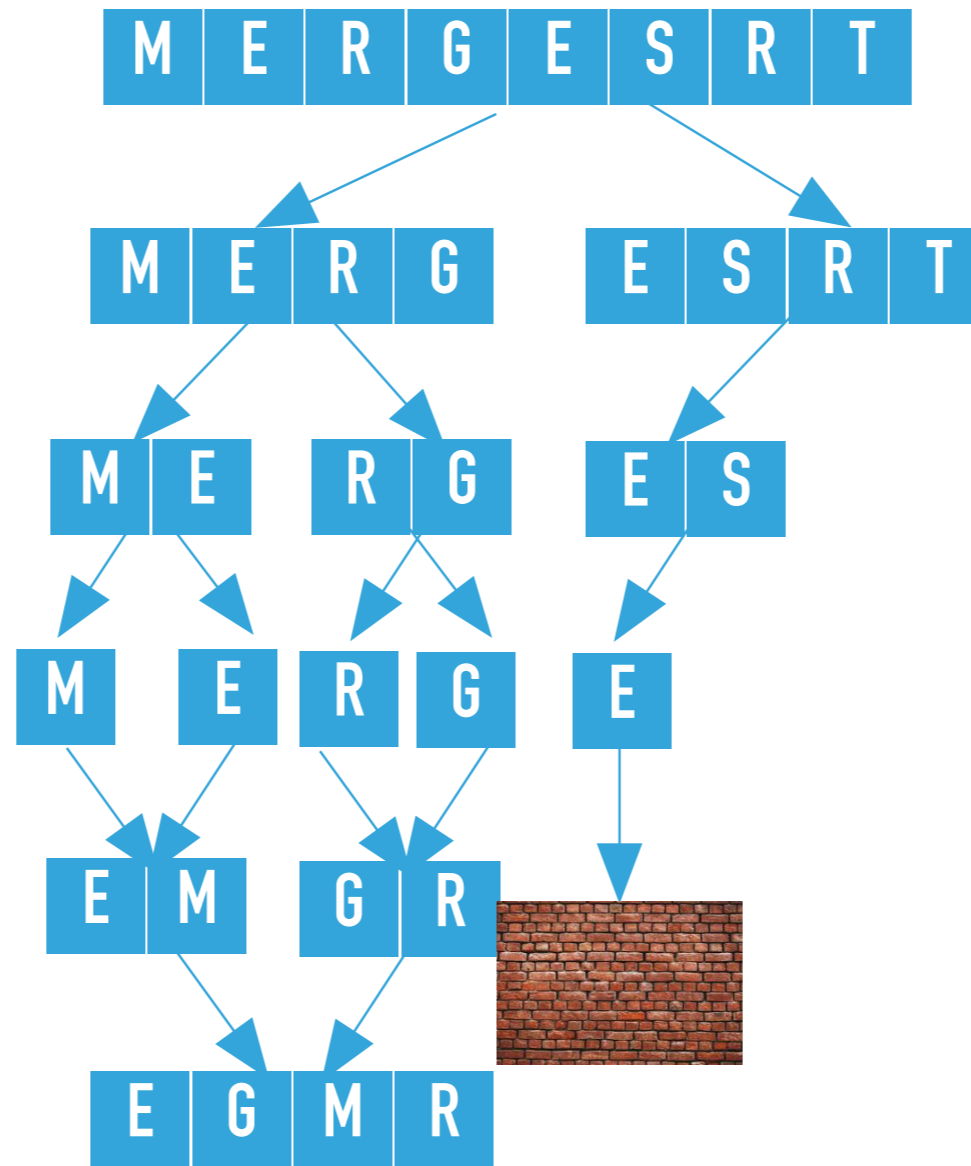


```

private static void sort(Comparable[] a, Comparable[] aux, int lo, int hi) {
    if (hi <= lo)
        return;
    int mid = lo + (hi - lo) / 2;
    sort(a, aux, lo, mid);
    sort(a, aux, mid+1, hi);
    merge(a, aux, lo, mid, hi);
}

```

`sort([E, G, M, R, E, S, R, T], [E, M, G, R, null, null, null, null], 4, 5)` calculates the `mid = 4` and calls recursively `sort` on the left subarray, that is `sort([E, G, M, R, E, S, R, T], [E, M, G, R, null, null, null, null], 4, 4)`, where `lo = 4, hi = 4`.

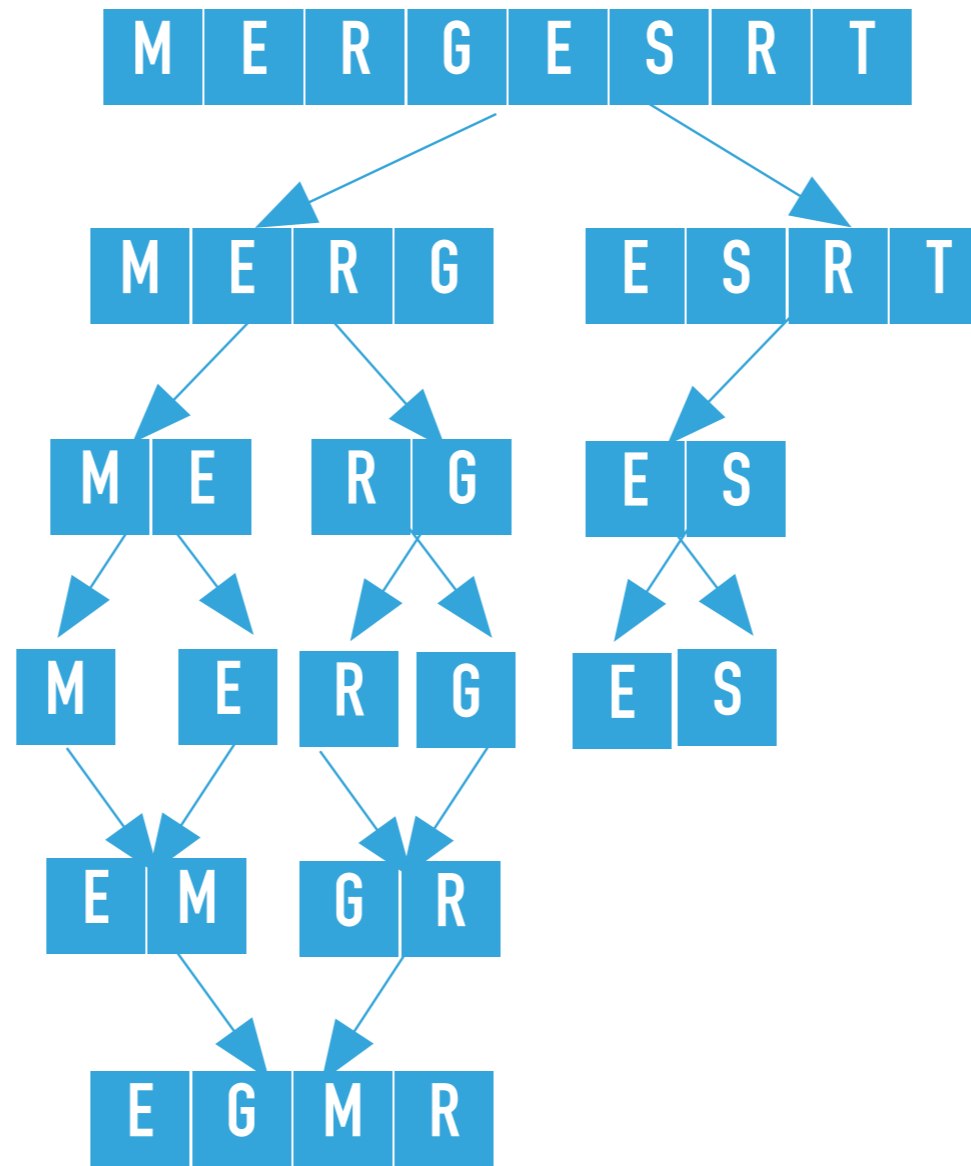


```

private static void sort(Comparable[] a, Comparable[] aux, int lo, int hi) {
    if (hi <= lo)
        return;
    int mid = lo + (hi - lo) / 2;
    sort(a, aux, lo, mid);
    sort(a, aux, mid+1, hi);
    merge(a, aux, lo, mid, hi);
}

```

`sort([E, G, M, R, E, S, R, T], [E, M, G, R, null, null, null, null], 4, 4)` finds `hi <= lo` and returns.

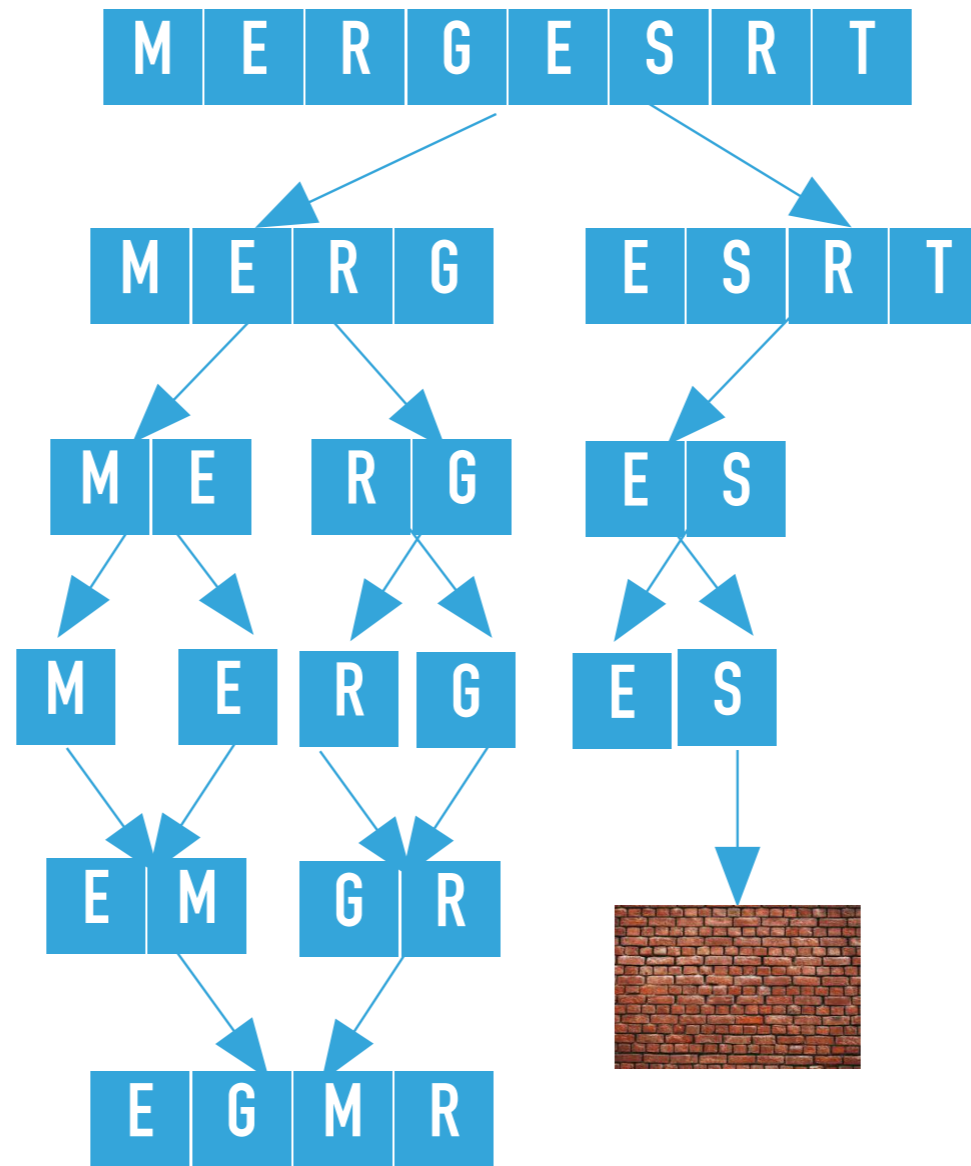


```

private static void sort(Comparable[] a, Comparable[] aux, int lo, int hi) {
    if (hi <= lo)
        return;
    int mid = lo + (hi - lo) / 2;
    sort(a, aux, lo, mid);
    sort(a, aux, mid+1, hi);
    merge(a, aux, lo, mid, hi);
}

```

sort([E, G, M, R, E, S, R, T], [E, M, G, R, null, null, null, null], 4, 5) calls recursively sort on the right subarray, that is sort([E, G, M, R, E, S, R, T], [E, M, G, R, null, null, null, null], 5, 5), where lo = 5, hi = 5

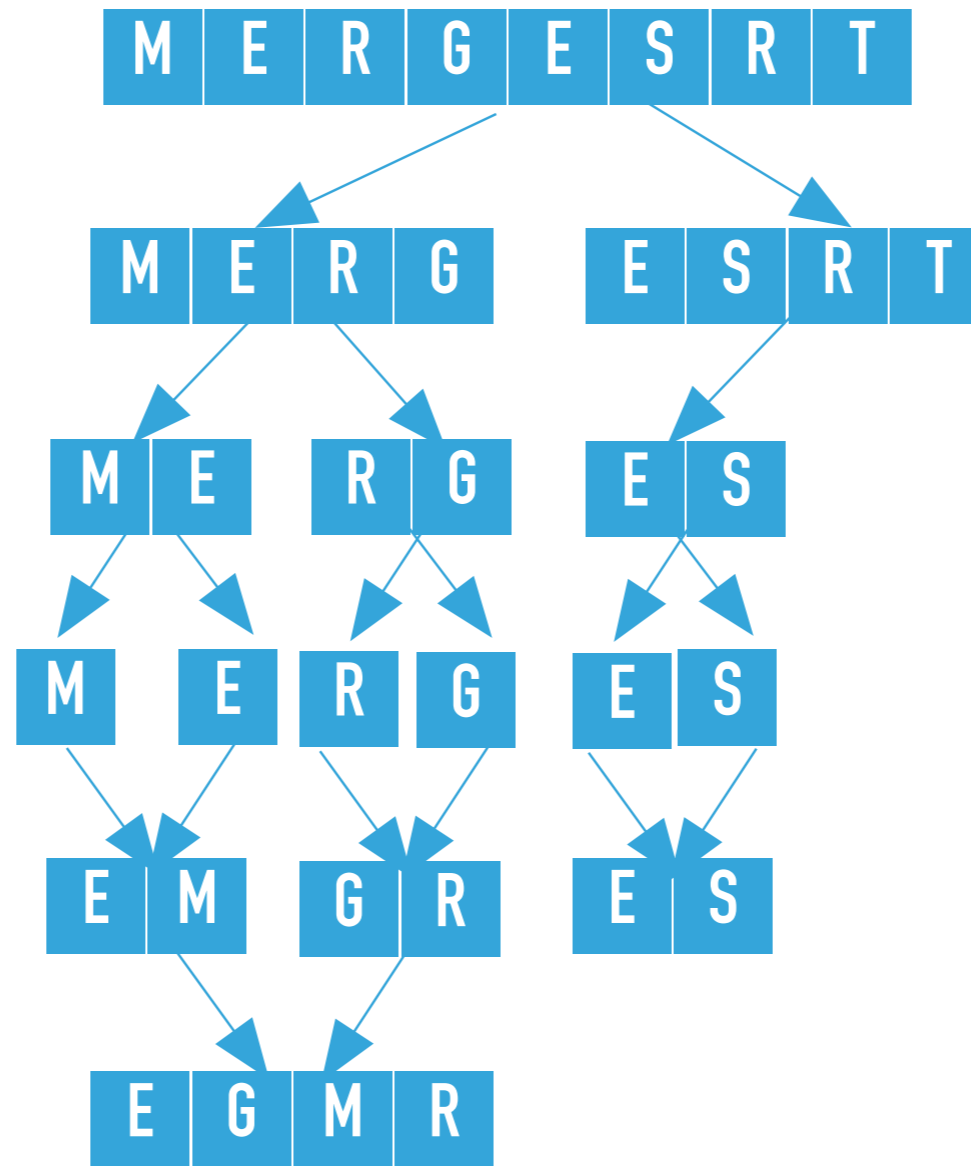


```

private static void sort(Comparable[] a, Comparable[] aux, int lo, int hi) {
    if (hi <= lo)
        return;
    int mid = lo + (hi - lo) / 2;
    sort(a, aux, lo, mid);
    sort(a, aux, mid+1, hi);
    merge(a, aux, lo, mid, hi);
}

```

`sort([E, G, M, R, E, S, R, T], [E, M, G, R, null, null, null, null], 5, 5)` finds `hi <= lo` and returns.

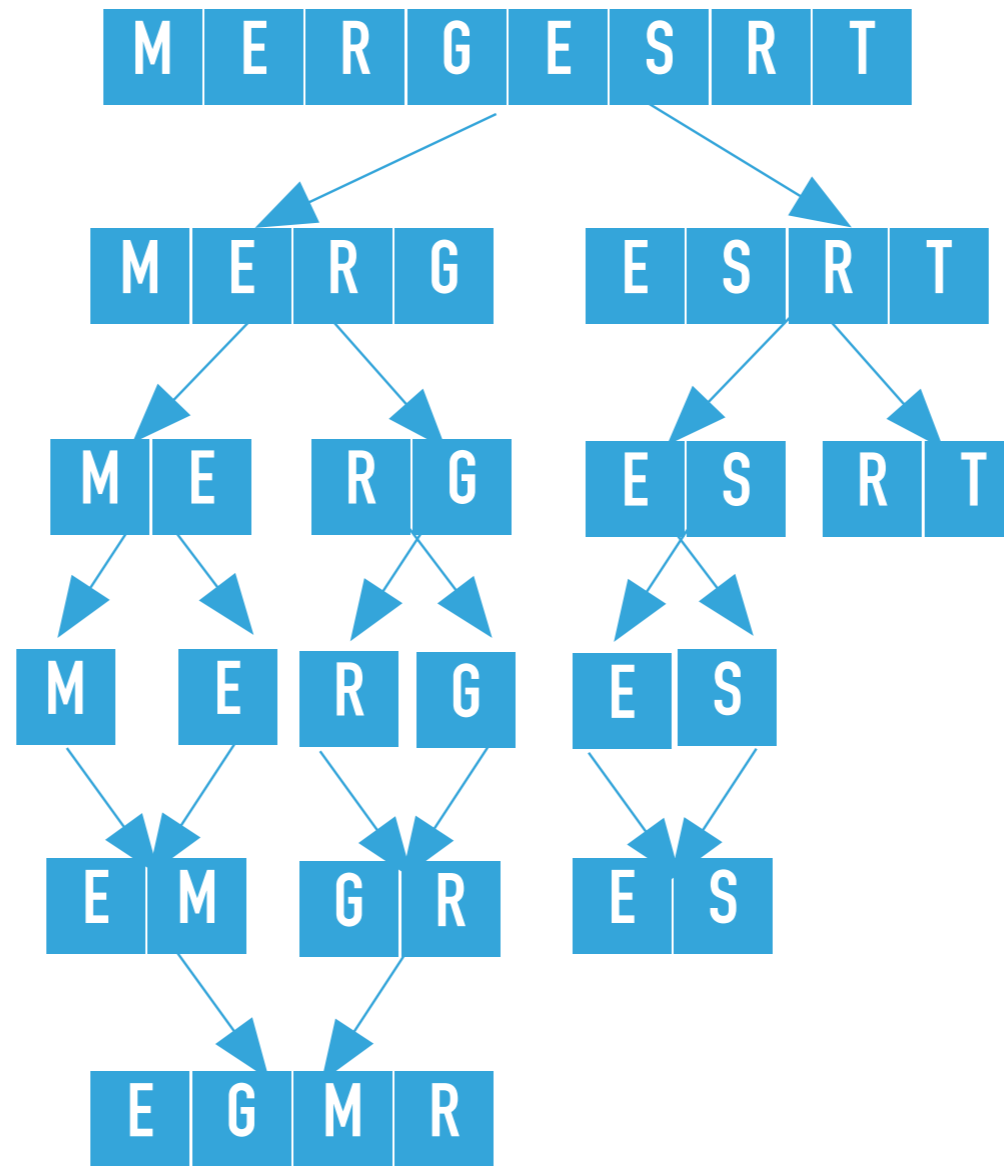


```

private static void sort(Comparable[] a, Comparable[] aux, int lo, int hi) {
    if (hi <= lo)
        return;
    int mid = lo + (hi - lo) / 2;
    sort(a, aux, lo, mid);
    sort(a, aux, mid+1, hi);
    merge(a, aux, lo, mid, hi);
}

```

`sort([E, G, M, R, E, S, R, T], [E, M, G, R, null, null, null, null], 4, 5)` merges the two subarrays that is calls `merge([E, G, M, R, E, S, R, T], [E, M, G, R, null, null, null, null], 4, 4, 5)`, where `lo = 4, mid = 4, and hi = 5`. The resulting partially sorted array is `[E, G, M, R, E, S, R, T]`.

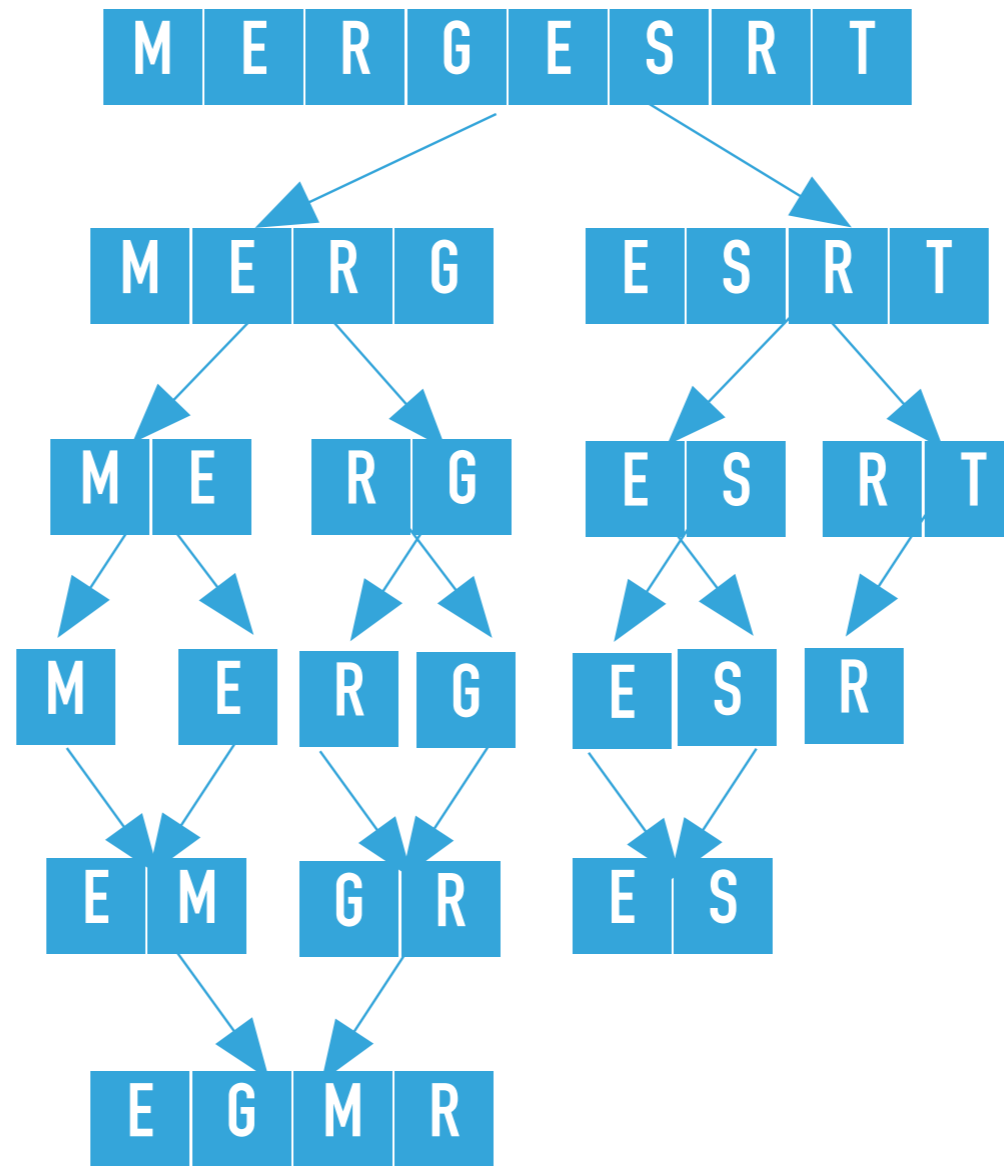


```

private static void sort(Comparable[] a, Comparable[] aux, int lo, int hi) {
    if (hi <= lo)
        return;
    int mid = lo + (hi - lo) / 2;
    sort(a, aux, lo, mid);
    sort(a, aux, mid+1, hi);
    merge(a, aux, lo, mid, hi);
}

```

sort([E, G, M, R, E, S, R, T], [E, M, G, R, E, S, null, null], 4, 7) calls recursively sort on the right subarray, that is sort([E, G, M, R, E, S, R, T], [E, M, G, R, E, S, null, null], 6, 7), where lo = 6, hi = 7

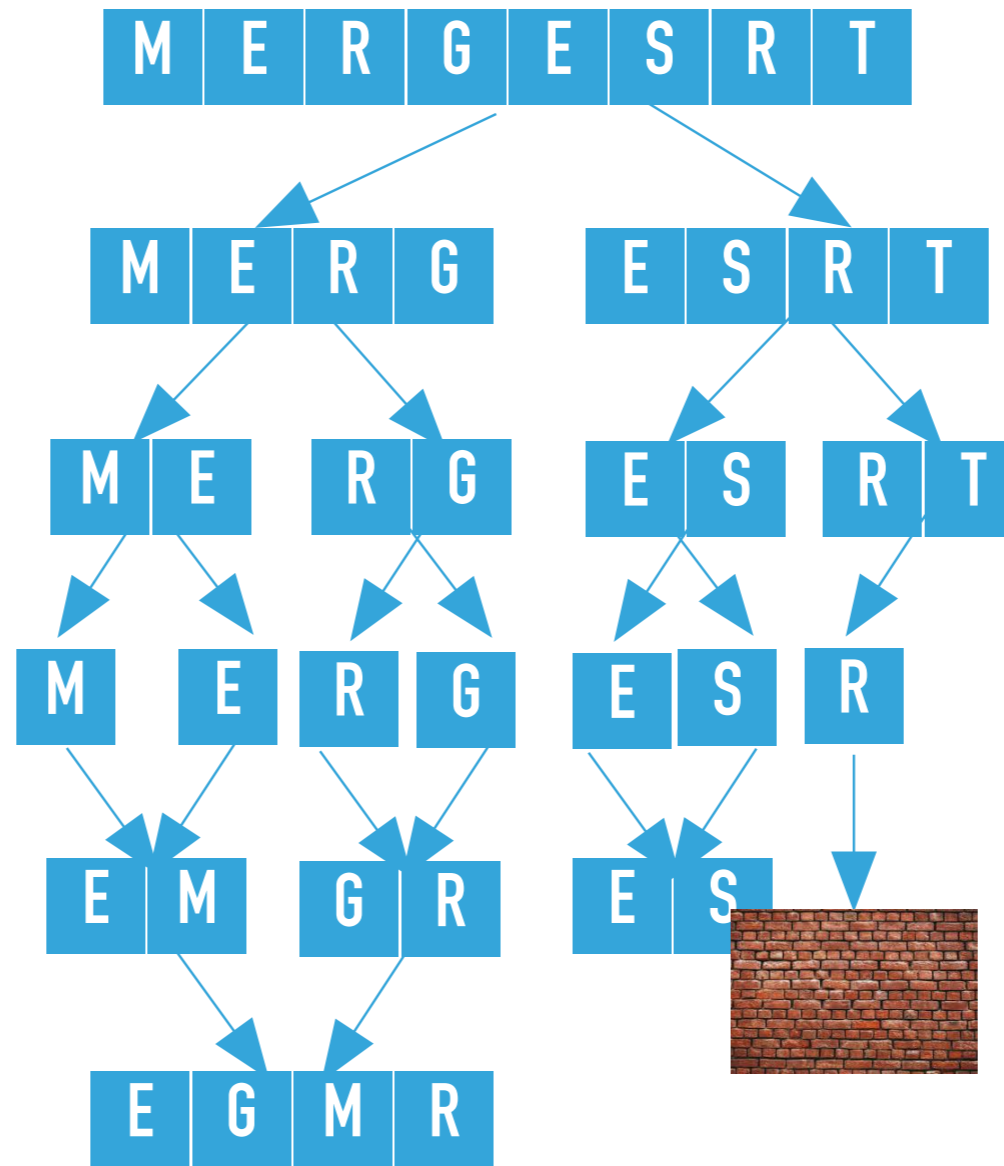


```

private static void sort(Comparable[] a, Comparable[] aux, int lo, int hi) {
    if (hi <= lo)
        return;
    int mid = lo + (hi - lo) / 2;
    sort(a, aux, lo, mid);
    sort(a, aux, mid+1, hi);
    merge(a, aux, lo, mid, hi);
}

```

sort([E, G, M, R, E, S, R, T], [E, M, G, R, E, S, null, null], 6, 7) calculates the mid = 6 and calls recursively sort on the left subarray, that is sort([E, G, M, R, E, S, R, T], [E, M, G, R, E, S, null, null], 6, 6), where lo = 6, hi = 6.

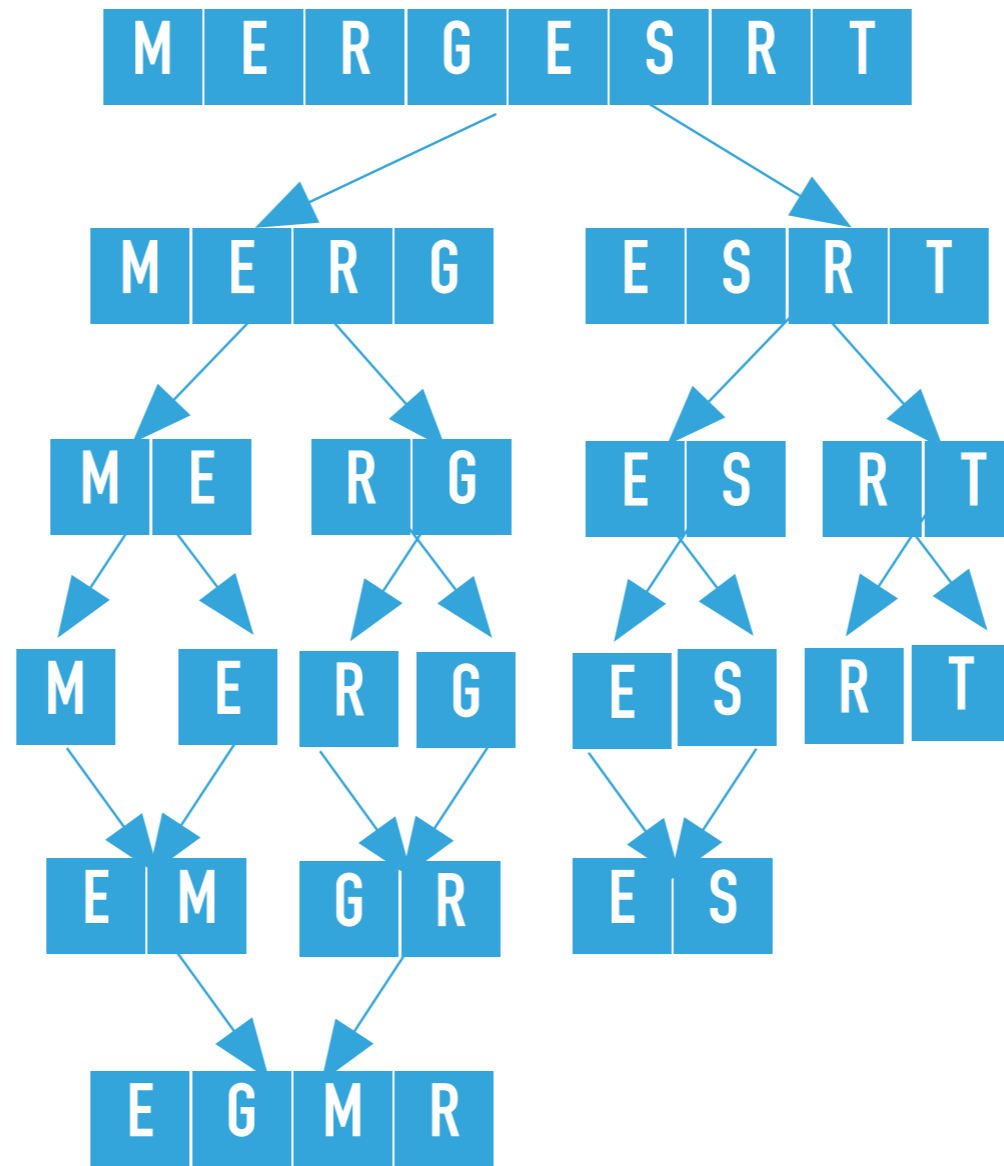


```

private static void sort(Comparable[] a, Comparable[] aux, int lo, int hi) {
    if (hi <= lo)
        return;
    int mid = lo + (hi - lo) / 2;
    sort(a, aux, lo, mid);
    sort(a, aux, mid+1, hi);
    merge(a, aux, lo, mid, hi);
}

```

`sort([E, G, M, R, E, S, R, T], [E, M, G, R, E, S, null, null], 6, 6)` finds `hi <= lo` and returns.

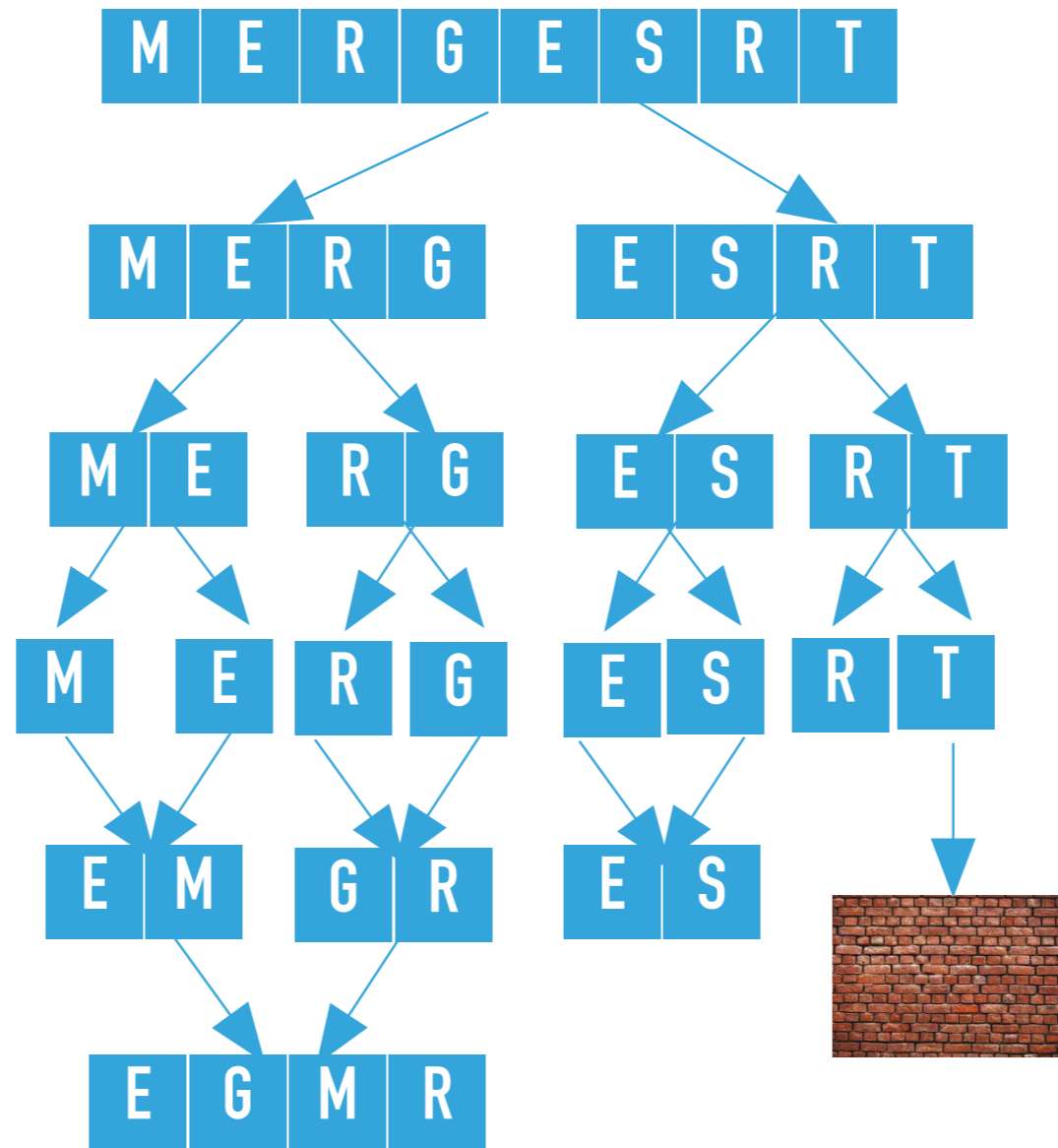


```

private static void sort(Comparable[] a, Comparable[] aux, int lo, int hi) {
    if (hi <= lo)
        return;
    int mid = lo + (hi - lo) / 2;
    sort(a, aux, lo, mid);
    sort(a, aux, mid+1, hi);
    merge(a, aux, lo, mid, hi);
}

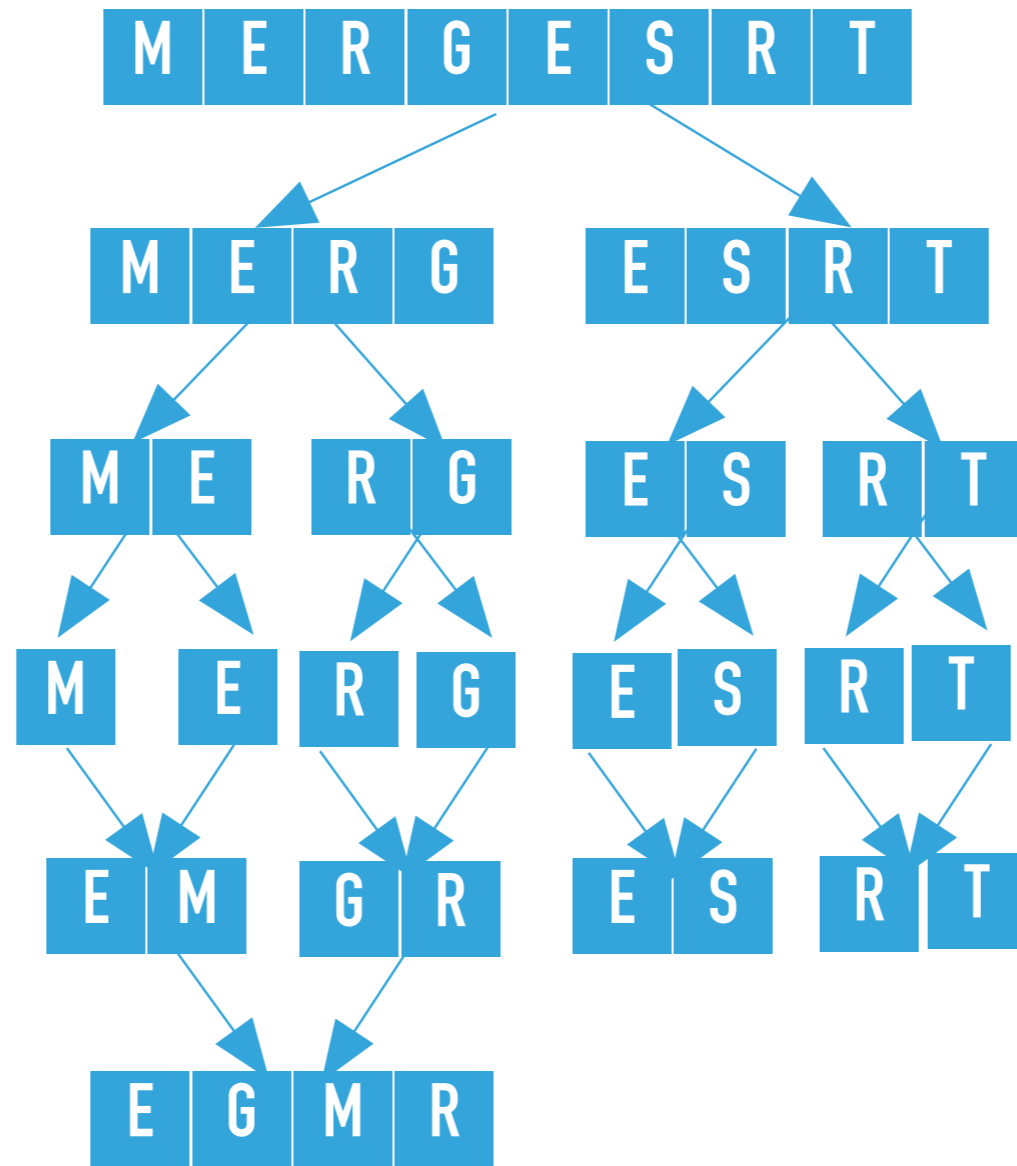
```

sort([E, G, M, R, E, S, R, T], [E, M, G, R, E, S, null, null], 6, 7) calls recursively sort on the right subarray, that is sort([E, G, M, R, E, S, R, T], [E, M, G, R, E, S, null, null], 7, 7), where lo = 7, hi = 7



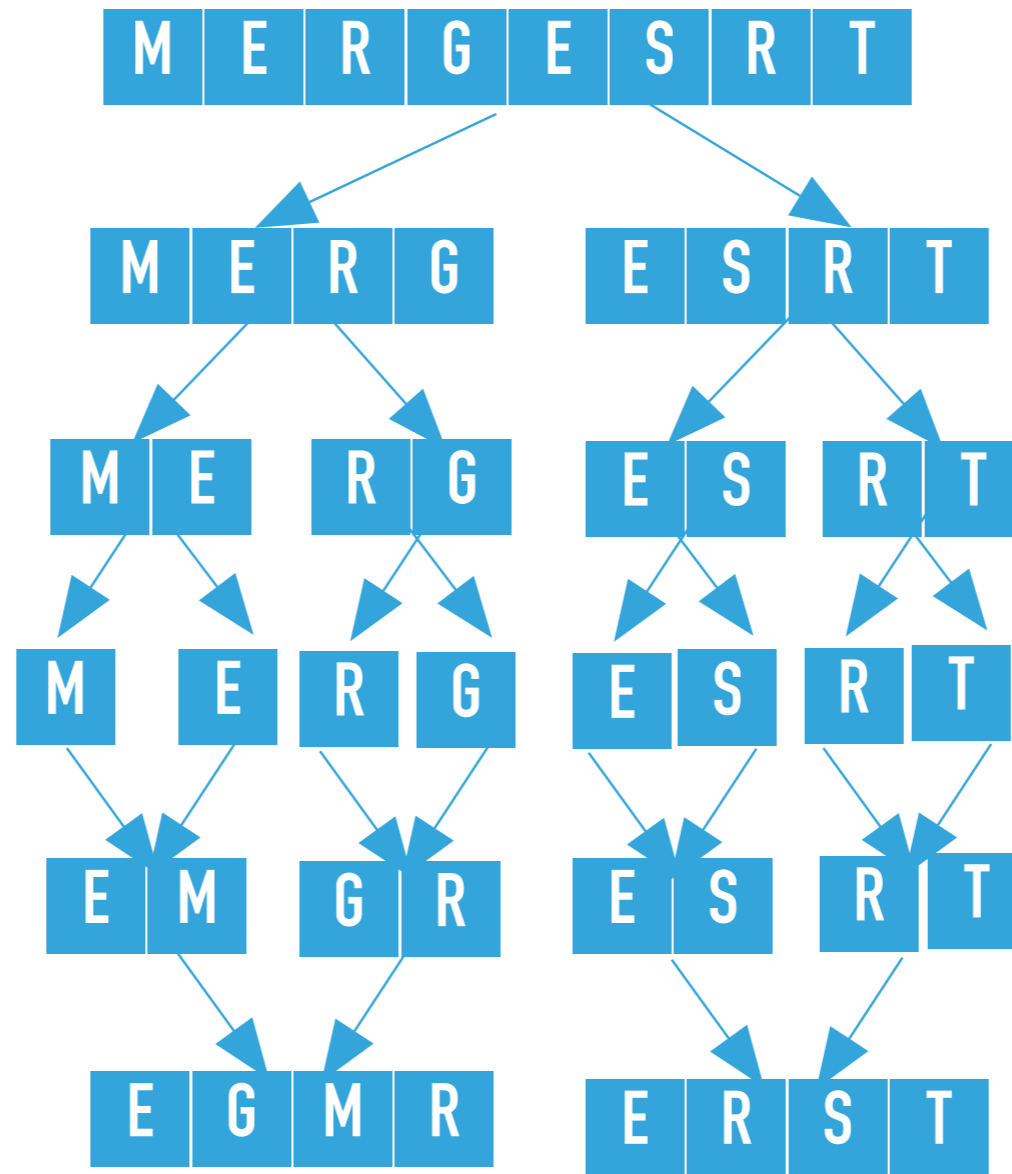
```
private static void sort(Comparable[] a, Comparable[] aux, int lo, int hi) {
    if (hi <= lo)
        return;
    int mid = lo + (hi - lo) / 2;
    sort(a, aux, lo, mid);
    sort(a, aux, mid+1, hi);
    merge(a, aux, lo, mid, hi);
}
```

`sort([E, G, M, R, E, S, R, T], [E, M, G, R, E, S, null, null], 7, 7)` finds `hi <= lo` and returns.



```
private static void sort(Comparable[] a, Comparable[] aux, int lo, int hi) {
    if (hi <= lo)
        return;
    int mid = lo + (hi - lo) / 2;
    sort(a, aux, lo, mid);
    sort(a, aux, mid+1, hi);
    merge(a, aux, lo, mid, hi);
}
```

`sort([E, G, M, R, E, S, R, T], [E, M, G, R, E, S, null, null], 6, 7)` merges the two subarrays that is calls `merge([E, G, M, R, E, S, R, T], [E, M, G, R, E, S, null, null], 6, 6, 7)`, where `lo = 6, mid = 6, and hi = 7`. The resulting partially sorted array is `[E, G, M, R, E, S, R, T]`.

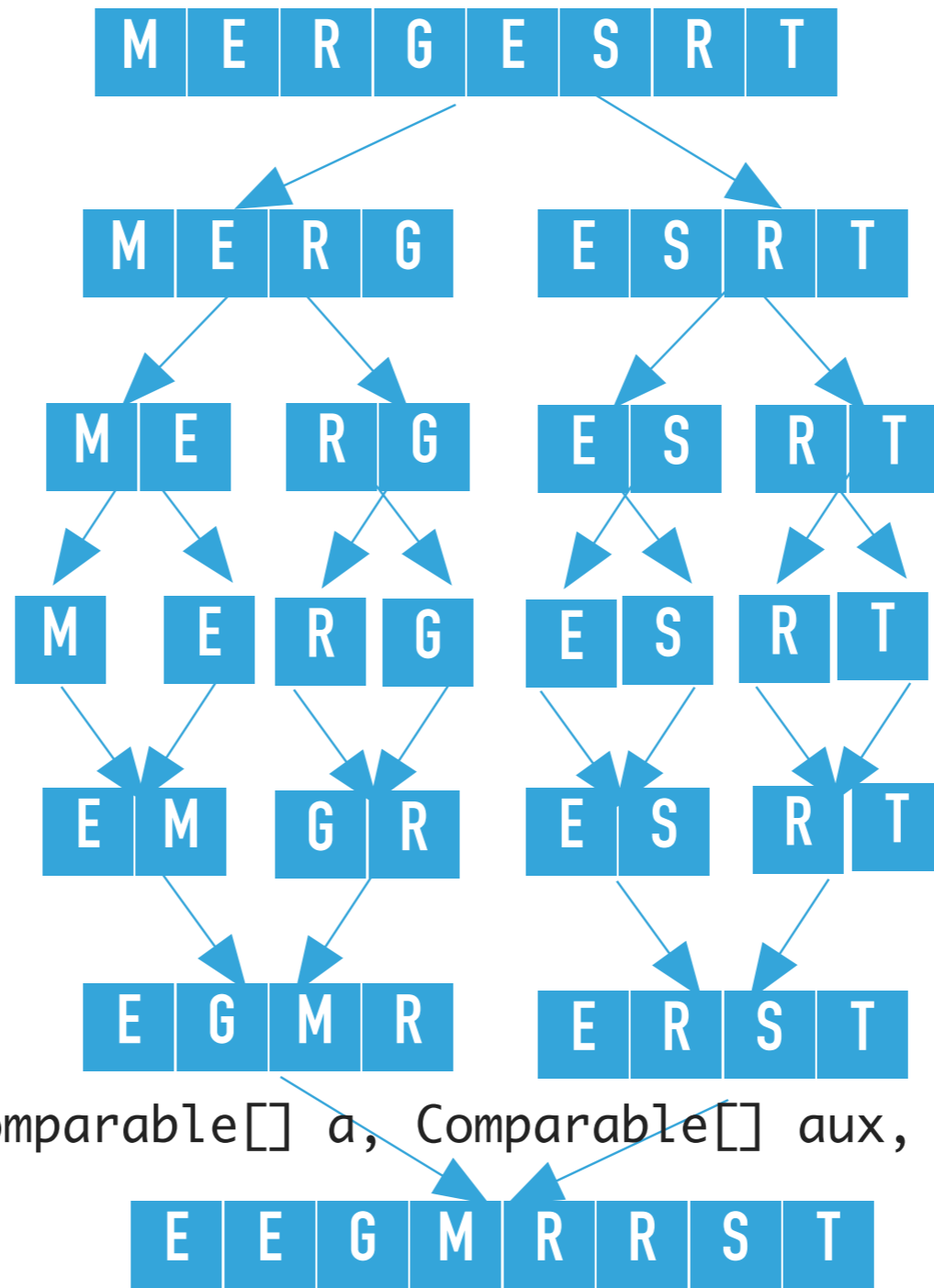


```

private static void sort(Comparable[] a, Comparable[] aux, int lo, int hi) {
    if (hi <= lo)
        return;
    int mid = lo + (hi - lo) / 2;
    sort(a, aux, lo, mid);
    sort(a, aux, mid+1, hi);
    merge(a, aux, lo, mid, hi);
}

```

`sort([E, G, M, R, E, S, R, T], [E, M, G, R, E, S, R, T], 4, 7)` merges the two subarrays that is calls `merge([E, G, M, R, E, S, R, T], [E, M, G, R, E, S, R, T], 4, 5, 7)`, where `lo = 4`, `mid = 5`, and `hi = 7`. The resulting partially sorted array is `[E, G, M, R, E, R, S, T]`.



```
private static void sort(Comparable[] a, Comparable[] aux, int lo, int hi) {
    if (hi <= lo)
        return;
    int mid = lo + (hi - lo) / 2;
    sort(a, aux, lo, mid);
    sort(a, aux, mid+1, hi);
    merge(a, aux, lo, mid, hi);
}
```

`sort([E, G, M, R, E, R, S, T], [E, M, G, R, E, S, R, T], 0, 7)` merges the two subarrays that is calls `merge([E, G, M, R, E, R, S, T], [E, M, G, R, E, S, R, T], 0, 3, 7)`, where `lo = 0`, `mid = 3`, and `hi = 7`. The resulting sorted array is `[E, E, G, M, R, R, S, T]`.

Practice time

Which of the following subarray lengths will occur when running mergesort on an array of length 10?

A. { 1, 2, 3, 5, 10 }

B. { 2, 4, 6, 8, 10 }

C. { 1, 2, 5, 10 }

D. { 1, 2, 3, 4, 5, 10 }

Answer

Which of the following subarray lengths will occur when running mergesort on an array of length 10?

A. { 1, 2, 3, 5, 10 }

Good algorithms are better than supercomputers

- ▶ Your laptop executes 10^8 comparisons per second
- ▶ A supercomputer executes 10^{12} comparisons per second

	Insertion sort			Mergesort		
Computer	Thousand inputs	Million inputs	Billion inputs	Thousand inputs	Million inputs	Billion inputs
Home	Instant	2 hours	300 years	instant	1 sec	15 min
Supercomputer	Instant	1 second	1 week	instant	instant	instant

Analysis of comparisons

- ▶ We will assume that n is a power of 2 ($n = 2^k$, where $k = \log_2 n$) and the number of comparisons $T(n)$ to sort an array of length n with merge sort satisfies the recurrence:
 - ▶ $T(n) = T(n/2) + T(n/2) + (n - 1) = O(n \log n)$
- ▶ Number of array accesses (rather than exchanges, here) is also $O(n \log n)$.

Mergesort uses $\leq n \log n$ compares to sort an array of length n

If $n = 4$, 2 levels

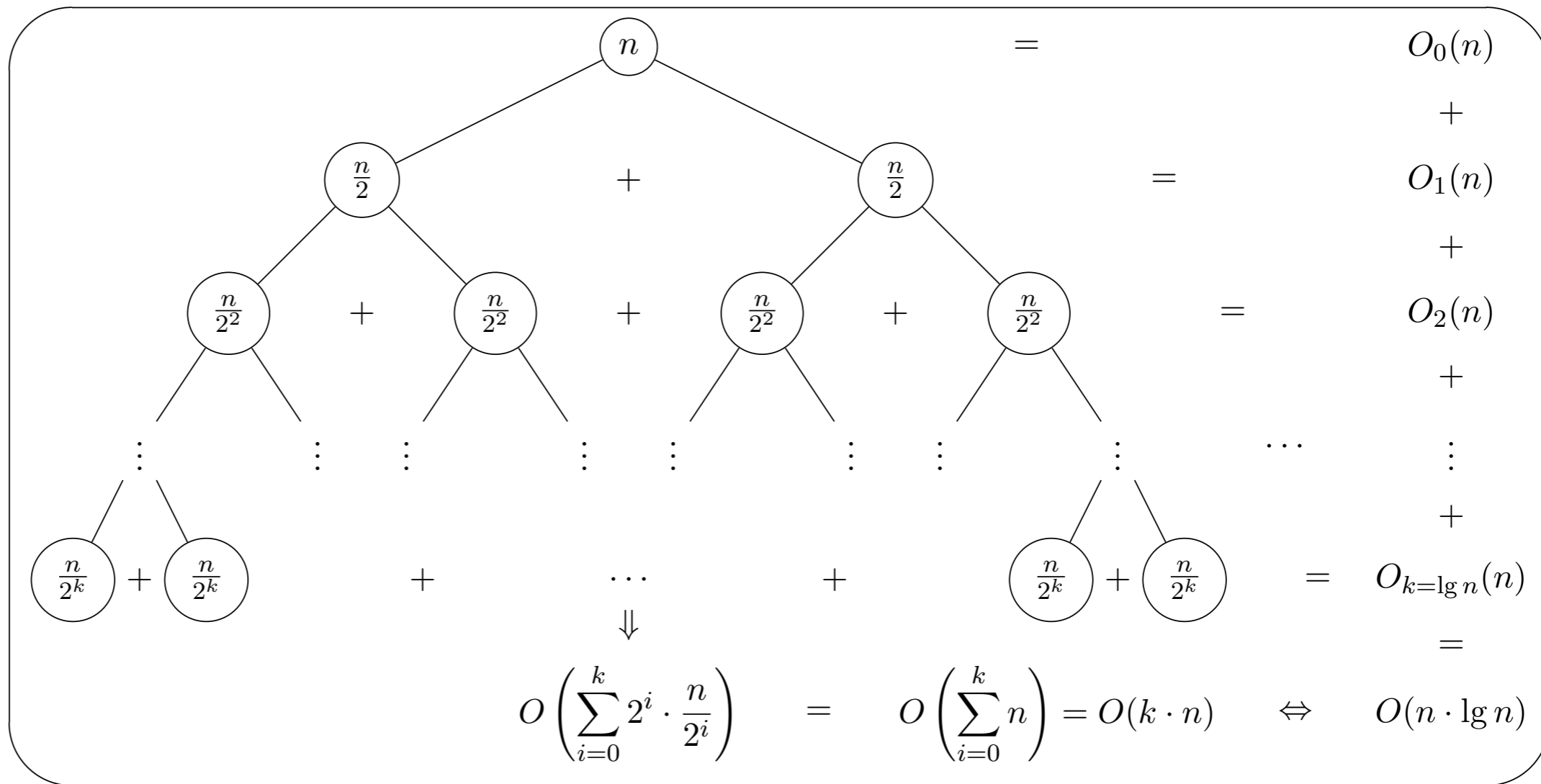
If $n = 8$, 3 levels

If $n = 16$, 4 levels

...

If $n = 2^k$, k levels,

or $k = \log_2 n$



(log n levels) x (n comparisons) = O(n log n)

Any algorithm with the same structure takes $n \log n$ time

```
public static void f(int n) {  
    if (n == 0)  
        return;  
    f(n/2);  
    f(n/2);  
    linear(n);  
}
```

Mergesort basics

- ▶ Auxiliary memory is proportional to n , as `aux[]` needs to be of length n for the last merge.
- ▶ At its simplest form, merge sort is **not an in-place algorithm**.
- ▶ There are modifications for halting the size of the auxiliary array but in-place merge is very hard.
- ▶ **Stable**: Look into `merge()`, if equal keys, it takes them from the left subarray.
 - ▶ So is insertion sort, but not selection sort.

Practical improvements for Mergesort

- ▶ Use insertion sort for small subarrays.
- ▶ Stop if already sorted.
- ▶ Eliminate the copy to the auxiliary array by saving time (not space).

```
private static void sort(Comparable[] src, Comparable[] dst, int lo, int hi) {  
    if (hi <= lo + 7) {  
        insertionSort(dst, lo, hi);  
        return;  
    }  
  
    int mid = lo + (hi - lo) / 2;  
    sort(dst, src, lo, mid);  
    sort(dst, src, mid+1, hi);  
  
    if (!less(src[mid+1], src[mid])) {  
        for (int i = lo; i <= hi; i++) dst[i] = src[i];  
        return;  
    }  
  
    merge(src, dst, lo, mid, hi);  
}
```

- ▶ For years, Java used this version to sort Collections of objects.

Sorting: the story so far

	In place	Stable	Best	Average	Worst	Remarks
Selection	X		$O(n^2)$	$O(n^2)$	$O(n^2)$	n exchanges
Insertion	X	X	$O(n)$	$O(n^2)$	$O(n^2)$	Use for small arrays or partially
Merge sort		X	$O(n \log n)$	$O(n \log n)$	$O(n \log n)$	Guaranteed performance; stable

Lecture 13: Mergesort

- ▶ Mergesort

Readings:

- ▶ Textbook:
 - ▶ Chapter 2.2 (pages 270-277)
- ▶ Website:
 - ▶ Mergesort: <https://algs4.cs.princeton.edu/22mergesort/>
 - ▶ Code: <https://algs4.cs.princeton.edu/22mergesort/Merge.java.html>

Practice Problems:

- ▶ 2.2.1-2.2.2, 2.2.11

Readings:

- ▶ Textbook:
 - ▶ Chapter 2.1 (pages 244-262), Chapter 2.1 (Page 247), Chapter 2.5 (Pages 338-339)
- ▶ Website:
 - ▶ Elementary sorts: <https://algs4.cs.princeton.edu/21elementary/>
 - ▶ Code: <https://algs4.cs.princeton.edu/21elementary/Selection.java.html> and <https://algs4.cs.princeton.edu/21elementary/Insertion.java.html>
- ▶ Oracle documentation:
 - ▶ Collections: <https://docs.oracle.com/javase/tutorial/collections/intro/index.html>
 - ▶ Comparable: <https://docs.oracle.com/javase/8/docs/api/java/lang/Comparable.html>
 - ▶ Comparator: <https://docs.oracle.com/javase/8/docs/api/java/util/Comparator.html>

Practice Problems:

- ▶ 2.1.1-2.1.8