

# Lecture 36: Graphs IV

CS 62

Fall 2018

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# Spanning Trees

- A spanning tree  $T$  of a graph  $G$  is a subset of the edges of  $G$  such that:
  - $T$  contains no cycles and
  - Every vertex in  $G$  is connected to every other vertex using just the edges in  $T$
- An unconnected graph has no spanning trees.
- A connected graph will have at least one spanning tree; it may have many

# Minimum Spanning Trees

- A weighted graph is a graph that has a weight associated with each edge.
- If  $G$  is a weighted graph, the cost of a tree is the sum of the costs (weights) of its edges.
- A tree  $T$  is a minimum spanning tree of  $G$  iff:
  - it is a spanning tree and
  - there is no other spanning tree whose cost is lower than that of  $T$ .

# Minimum Spanning Trees

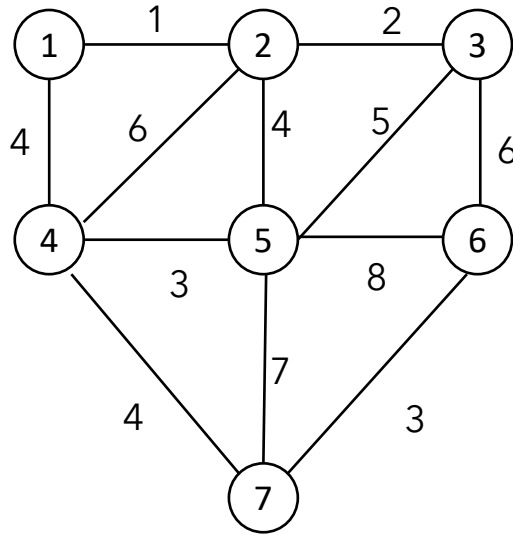
- Application:
  - The cheapest way to lay cable that connects a set of points is along a minimum spanning tree that connects those points.
- Many algorithms exist to find minimum spanning trees, most run in  $O(m \log m)$  time.
- In 1995 Karger, Klein & Tarjan found a linear time randomized algorithm, but there is no known linear time deterministic algorithm

# Kruskal's Algorithm

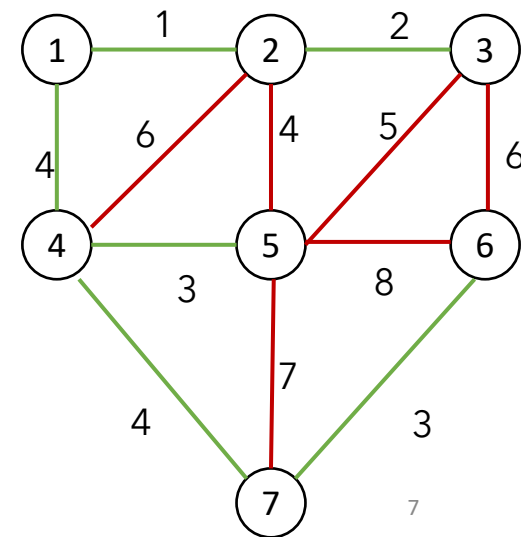
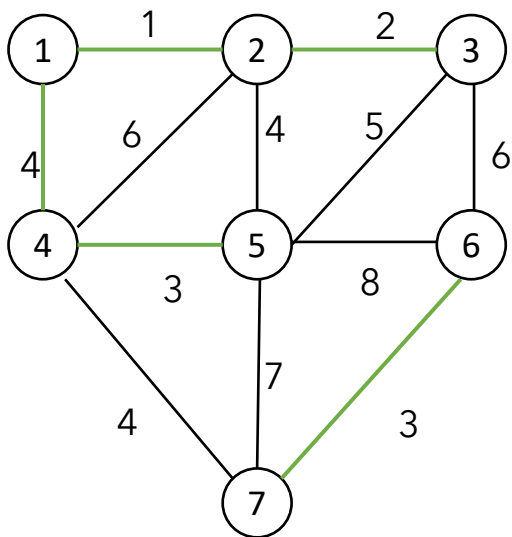
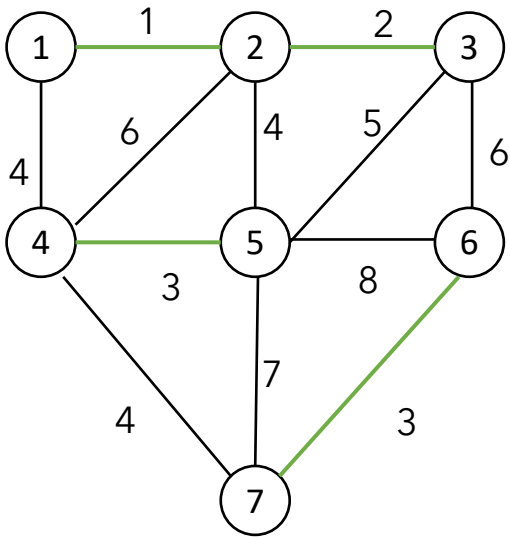
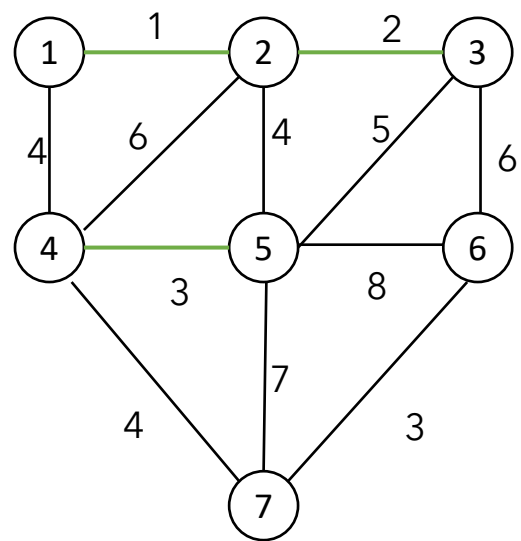
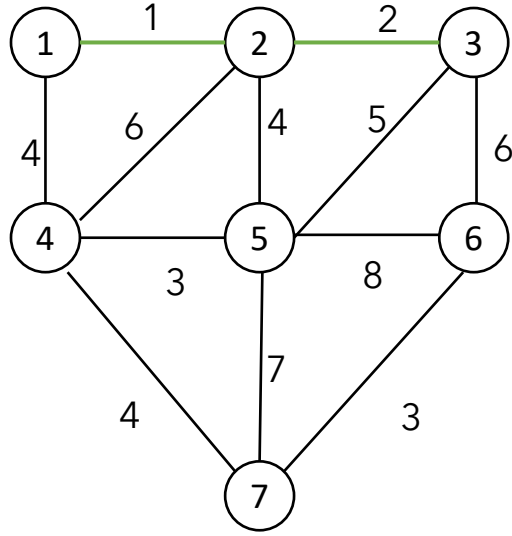
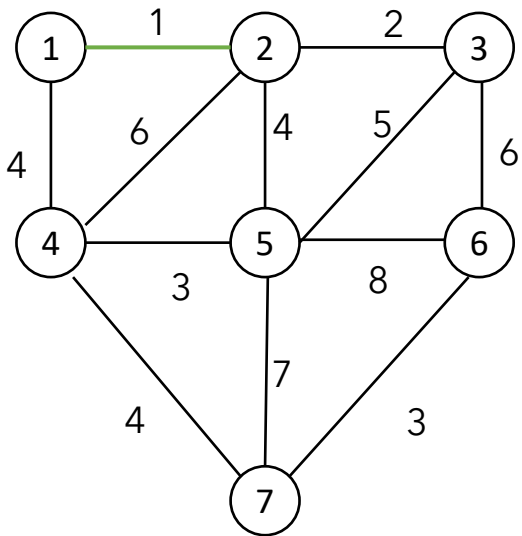
- Create forest  $F$  with no edges, using vertices in  $V$
- Sort the edges in the graph by their weight (smallest to largest)
- For each edge  $e$  in sorted order:
  - if  $e$  connects two different trees in  $F$ , then add  $e$  to  $F$

# Kruskal on sample graph

- (1,2):1
- (2,3):2
- (4,5):3
- (6,7):3
- (1,4):4
- (2,5):4
- (4,7):4
- (3,5):5
- (2,4):6
- (3,6):6
- (5,7):7
- (5,6):8



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- (2,3):2
- (4,5):3
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- (1,4):4
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- (3,5):5
- (2,4):6
- (3,6):6
- (5,7):7
- (5,6):8



# Kruskal's Algorithm pseudocode

```
A = {};  
for(every vertex v in V) {  
  make-set(v)  
  for(every edge (u, v) ordered by increasing weight) {  
    if(find (u) != find (v)) {  
      A.add((u, v));  
      union(u, v);  
    }  
  }  
}  
return A;
```

make-set(v) - makes a set from a single vertex v

find(v) - finds the set that v belongs to

union(u, v) - makes the union of the sets containing u and v

} Union-find structure



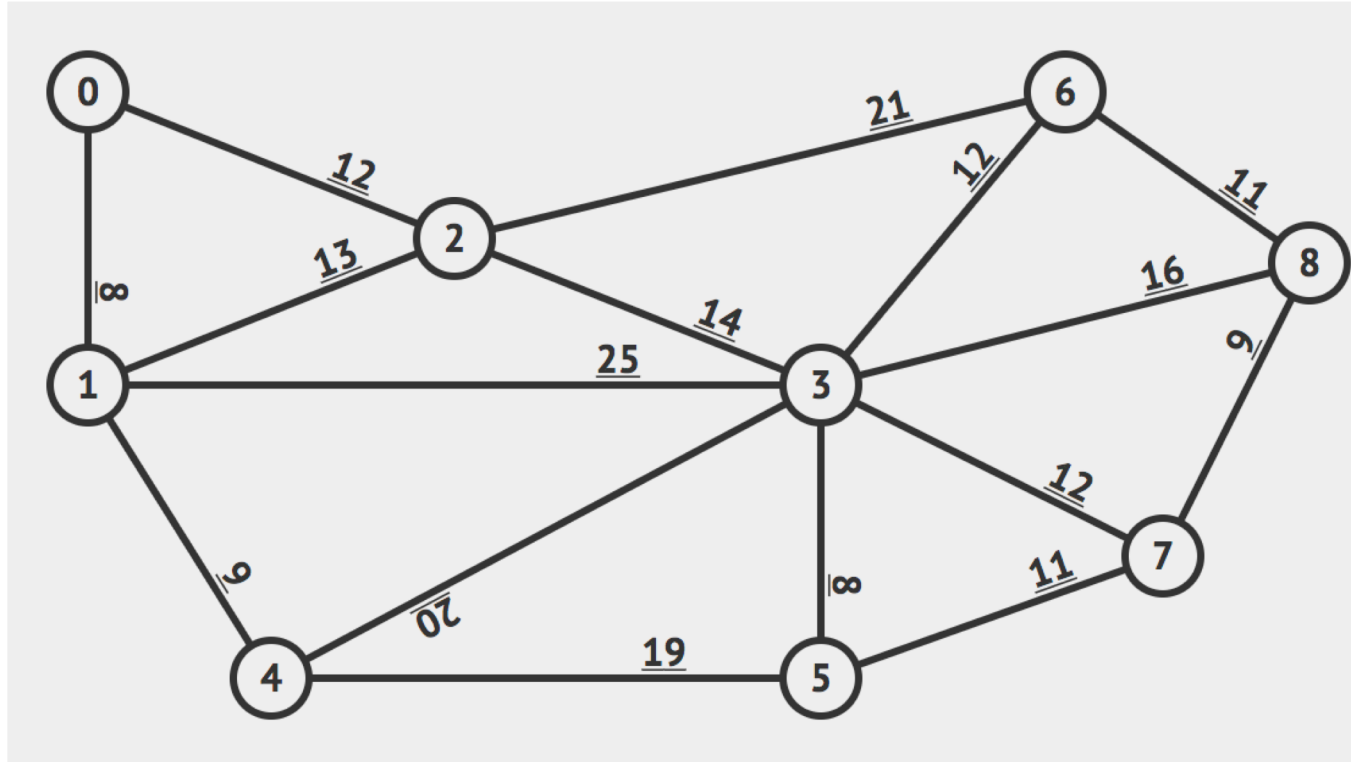
# Union-Find Data Structure

keeps track of a set of elements partitioned into a number of disjoint subsets

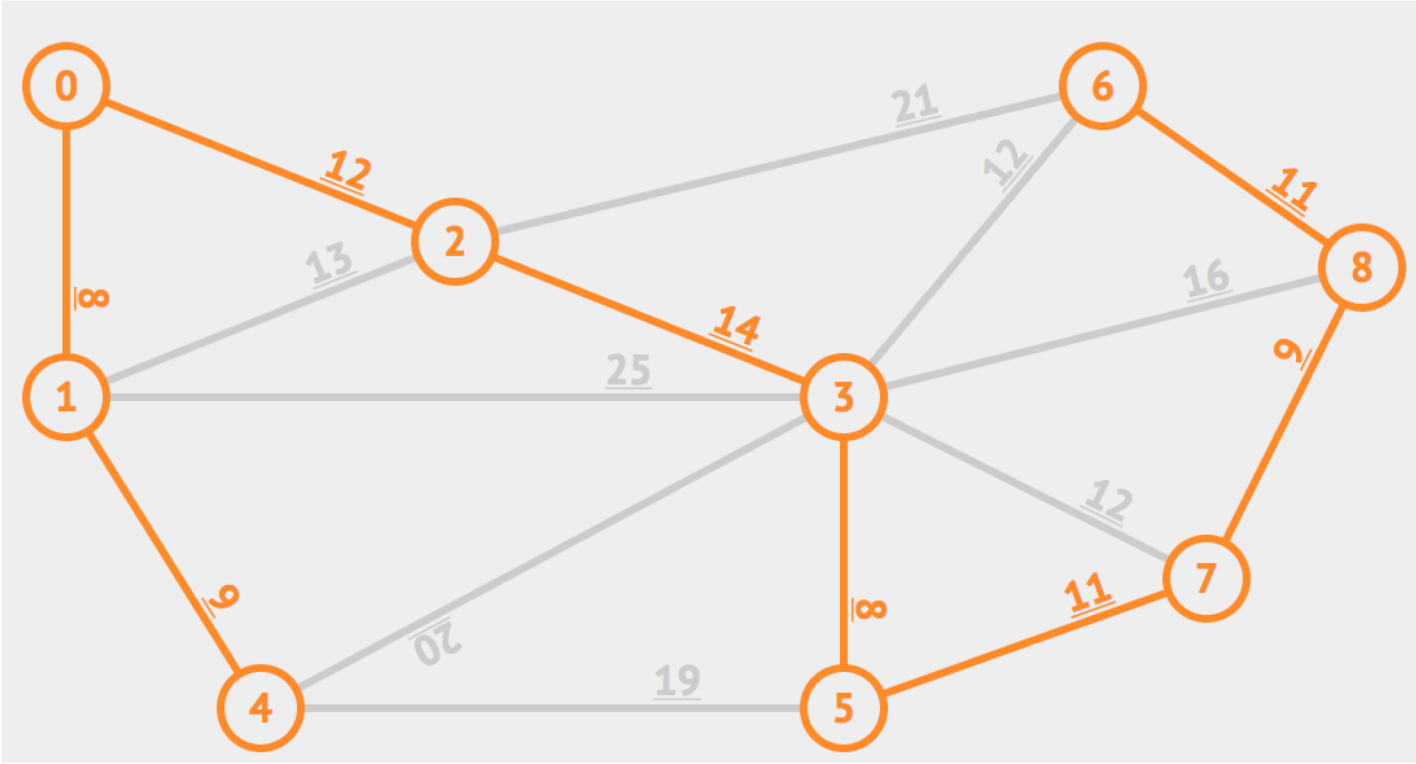
*Find*: Find what subset an element belongs. Use to find if two elements belong in the same subset

*Union*: Create a single subset out of two subsets

# Practice Time



# Answer



# Graph Algorithms

- Very important in practice!
- Sophisticated data structures
- Careful analysis of correctness and complexity
- CS 140: Algorithms