Lecture 6: ArrayList implementation & Complexity

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Lab This Week

- Timing ArrayList operations
 - Encourage working in pairs
 - Stopwatch class: start(), stop(), getTime(), reset()
- Java has Just-In-Time compiler
 - Must "warm-up" before you get accurate timing
 - What can mess up timing?
- Uses Vector from Bailey rather than ArrayList from Java libraries because can change way it increases in size.

Programming Assignment This Week

- Weak AI/Natural Language Processing:
 - Generate text by building frequency lists based on pairs of words. ArrayList of Associations of String (words) and Integer (count of that word).



- Generated javadoc for fun
- See how ArrayList used in methods for PostItApplication
 - findWindowInList, moveToTop, removeWindow
 - Used in mouse-event-handling methods

ArrayList

- Not using Bailey implementation
 - see code on-line for implementation by Tomassia & Goodrich
- Standard Java libraries have lots of extra methods not in our implementation:
 - Many involve working on other collections
 - irrelevant for us at this point.
 - addAll, clear, contains, containsAll, listIterator, removeAll, replaceAll, retainAll, sort, spliterator, sublist, toArray

Back to ArrayList

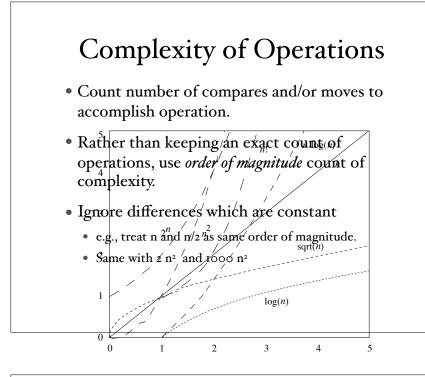
- Interface is IndexList<E>
- See ArrayIndexList<E>
 - Similar to ArrayList
 - Instance variables:
 - elts: array instance variable,
 - eltsFilled: number of slots filled.
- Creating new ArrayList is weird
 - Can't construct array of variable type!
 - Create array of Object, but coerce to believe array of E.

ArrayList Implementation

- Some operations very cheap:
 - size, isEmpty, get, set take constant time (no search)
- Others more expensive

Adding Elts in Slot i

- Easy if there is space:
 - At end, just add it
 - If before end, must move all elements at i and beyond to right before inserting
 - Delete similar
- What if run out of space
 - Create new array twice as big and copy old elements over before adding.
- How expensive is this?



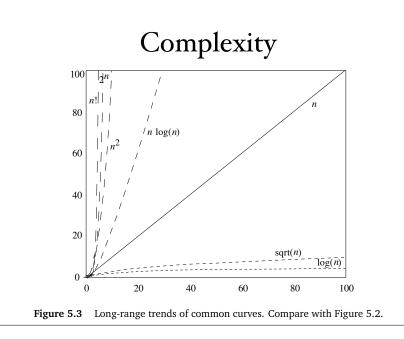
Order of Magnitude

- <u>Definition</u>: We say that g(n) is O(f(n)) if there exist two constants C and k such that |g(n)| <= C |f(n)| for all n > k.
- Examples: 2n+1, n³-n²+83, 2ⁿ+n²
- Used to measure time and space complexity of algorithms on data structures of size n.
- Most common are

• O(1) - for any constant

Use simplest version in O(...)

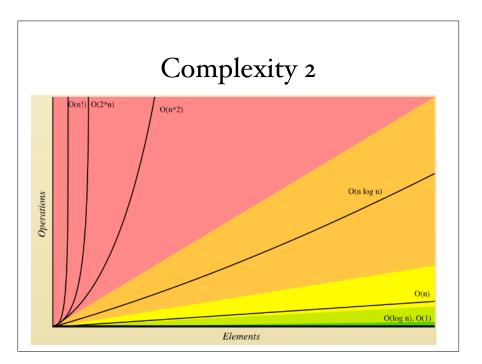
• O(log n), O(n), O(n log n), O(n²), ..., O(2ⁿ)



Comparing Orders of Magnitude

- Suppose have ops w/complexities given & problem of size n taking time t.
- How long if increase size of problem?

Problem Size:	10 n	100n	1000n
O(log n)	3+t	7 + t	10+ t
O(n)	10 t	100 t	1000 t
O(n log n)	> 10 t	> 100 t	> 1000 t
$O(n^2)$	100 t	10,000 t	1,000,000 t
$O(2^n)$	~ t ^{io}	~ t ¹⁰⁰	~ t ¹⁰⁰⁰



Adding to ArrayList

- Suppose n elements in ArrayList and add 1.
- If space:
 - Add to end is O(1)
 - Add to beginning is O(n)
- If not space,
 - What is cost of ensureCapacity?
 - O(n) because n elements in array

EnsureCapacity

- What if only increase in size by I each time?
 - Adding n elements one at a time to end
 - Total cost of copying over arrays: 1+2+3+...+(n-1) = n(n-1)/2
 - Total cost of $O(n^2)$
 - Average cost of each is O(n)
- What if double in size each time?
 - Suppose add $n = 2^m$ new elts to end
 - Total cost of copying over arrays: 1+2+4+...+n/2 = n-1, $\mathrm{O}(n)$
 - Average cost of O(1), but "lumpy"

ArrayList Ops

- Worst case
 - O(1): size, isEmpty, get, set
 - O(n): remove, add
- Add to end, on average O(I)