csci54 – discrete math & functional programming relations, RSA

Relations recap

- A (binary) relation on a set A is a subset of AxA
- A relation can be any, or none, of the following:
 - reflexive
 - symmetric
 - transitive
- A relation that is reflexive, symmetric, and transitive is called an equivalence relation
- An equivalence relation on a set A partitions A into a set of equivalence classes

- For each of the following relations, indicate if it is reflexive, symmetric, and/or transitive. If it's all three and therefore an equivalence relation, describe the equivalence classes.
- 1. S = all juniors and seniors at Pomona. (x,y) in R_1 iff x and y share a major.

2.
$$S = Z$$
. (x,y) in R_2 iff x=y.

- 3. $S = \{1,2,3,4,5\}$. $R_3 = \{\langle 1, 5 \rangle, \langle 2, 2 \rangle, \langle 2, 4 \rangle, \langle 4, 1 \rangle, \langle 4, 2 \rangle\}$.
- S = all students at Pomona. Define an equivalence relation on S that isn't one of the ones discussed in lecture last time.

Closures

Definition 8.11: Reflexive, symmetric, and transitive closures.

Let $R \subseteq A \times A$ be a relation. Then:

The *reflexive closure of* R is the smallest relation $R' \supseteq R$ such that R' is reflexive.

The symmetric closure of R is the smallest relation $R'' \supseteq R$ such that R'' is symmetric.

The *transitive closure of* R is the smallest relation $R^+ \supseteq R$ such that R^+ is transitive.

- Consider the relation $R = \{(1, 5), (2, 2), (2, 4), (4, 1), (4, 2)\}$ on $\{1, 2, 3, 4, 5\}$
 - What is the reflexive closure?
 - What is the symmetric closure?
 - What is the transitive closure?
- S = all students at Pomona. (x,y) in R_1 if x and y share a major.

Transmitting information



- cryptography
- error correction
- compression



- goal is to keep someone with access to the channel from finding out information about the message.
- assumptions (for now)
 - message = message'
 - codeword = codeword'
- why?
- how?

Private key cryptography



- Symmetric-key algorithms
- The communicating parties share a piece of secret information (the key k)
- Examples?
- Challenges?

Public key cryptography



- asymmetric-key algorithm
- Everyone publishes their public key.
- If you want to send a message to someone, you encrypt it with their public key.
- When you get a message you can decrypt it with your private key.

Public key encryption in practice

Richard Adams

Α

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PUBLIC KEY

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-----END PGP PUBLIC KEY BLOCK-----

Encrypted Email

If you use PGP encryption, here is our fingerprint and link to our public key. If you use our public key with a mail encryption plugin, for example Mailvelope or Enigmail, this encrypts the contents of your message but not the subject line or the name of the sender.

Fingerprint:

EC6C 2905 F0F9 3C03 7394 6CA1 0642 427A 5FF7 80BE

Email: lockbox@washpost.com

The Post's public key

Сору

-BEGIN PGP PUBLIC KEY BLOCK-

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https://www.theguardian.com/pgp https://www.washingtonpost.com/anonymous-news-tips/

RSA algorithm

- A very widely used public key encryption algorithm
- Three algorithmic components
 - key generation
 - encryption
 - decryption

Our plan

- What is the algorithm?
- Why does it work?
- How to implement it efficiently?



Modular arithmetic – definitions and properties

 $a \equiv b \pmod{m}$ if and only if a mod $m = b \mod m$

i.e., exists x, y, z in Z:
$$a = x * m + z \land b = y * m + z$$

Some useful facts from Wikipedia:

(if a1=a2 (mod m) and b1=b2(mod m):

 $a + k \equiv b + k \pmod{m}$ for any integer k (compatibility with translation)

 $k a \equiv k b \pmod{m}$ for any integer k (compatibility with scaling)

 $k a \equiv k b \pmod{k * m}$ for any integer k

 $a1 + a2 \equiv b1 + b2 \pmod{m}$ (compatibility with addition)

 $a1 - a2 \equiv b1 - b2 \pmod{m}$ (compatibility with subtraction)

 $a1 * a2 \equiv b1 * b2 \pmod{m}$ (compatibility with multiplication)

 $a^{k} \equiv b^{k} \pmod{m}$ for any non-negative integer k (compatibility with exponentiation)