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csci54 – discrete math & functional programming  
propositional logic continued, predicate logic

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# last time

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- ▶ introduction to propositional logic:

- ▶ Boole

- ▶ proposition

- ▶ well-formed propositional logic formulas (wff)

$$\phi ::= T|F|(\neg\phi)|(\phi \wedge \phi)|(\phi \vee \phi)|(\phi \Rightarrow \phi)$$

- ▶ truth tables for operators

- ▶ tautology/satisfiable/contingency (falsifiable)/contradiction

- ▶ implication

- ▶ logical equivalence



## converse, inverse, contrapositive

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Given an implication  $p \Rightarrow q$ , we can derive three other implications:

- converse:  $q \rightarrow p$
  - inverse:  $\neg p \rightarrow \neg q$
  - contrapositive:  $\neg q \rightarrow \neg p$
- Which, if any, of the converse, inverse, and contrapositive is logically equivalent to the original implication?



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consider the following statements . . .

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- ▶ If 2 is an even number then 3 is an odd number.
- ▶ If  $x$  is an even number, then  $x+1$  is an odd number.
  
- ▶ How would you express these two statements in propositional logic?



# predicate logic

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- ▶ A predicate  $P$  is function that assigns the value True or False to each element of a set  $U$ .
  - ▶ The set  $U$  is called the universe or domain of discourse
  - ▶  $P$  is a predicate over  $U$
- ▶ Examples:
  - ▶ the predicate "is an even number" over the positive integers.
  - ▶ the predicate "last name has at least 6 characters" over the set of people currently in this room.
- ▶ Once you specify the element of  $U$ , then you have a proposition with a truth value.



# quantifiers

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- ▶ quantifiers are another way to form propositions from a predicate

**Definition 3.21: Universal quantifier [for all,  $\forall$ ].**

Let  $P$  be a predicate over  $S$ . The proposition  $\forall x \in S : P(x)$  is true if, for *every* possible  $x \in S$ , we have that  $P(x)$  is true.

**Definition 3.22: Existential quantifier [there exists,  $\exists$ ].**

Let  $P$  be a predicate over  $S$ . The proposition  $\exists x \in S : P(x)$  is true if, for *at least one* possible  $x \in S$ , we have that  $P(x)$  is true.



## quantifiers - example

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- ▶ Imagine these predicates
  - ▶ "rested(n)" means "n got at least 8 hours of sleep in the past 24 hours"
  - ▶ "bornMA(n)" means "n was born in Massachusetts"
- ▶ Which, if any, of the following propositions is true? Justify your answer.
  - ▶  $\forall n$  in this room : rested(n)
  - ▶  $\forall n$  in this room : (rested(n)  $\wedge$  bornMA(n))
  - ▶  $\exists n$  currently enrolled at Pomona College : (rested(n)  $\vee$  bornMA(n))
  - ▶  $\exists n$  currently enrolled at Pomona College : (rested(n)  $\wedge$  bornMA(n))





## free and bound variables (an aside)

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- ▶ In an expression variables can be free/unbound or bound
  - ▶ With a free variable the value is not fixed by the expression
  - ▶ With a bound variable the value is defined within the expression

$$\forall x \in \mathbb{Z} : x^2 \geq y$$

- ▶ An expression of predicate logic with no free variables is called fully quantified



# theorems in predicate logic

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- ▶ A fully quantified expression of predicate logic is a theorem if and only if it is true for every possible meaning of each of its predicates.
- ▶ Is the following a theorem?  
$$[\forall x \in S : P(x)] \vee [\forall x \in S : \neg P(x)]$$
- ▶ What is an example of a predicate for which the statement is false? is true?



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## practice question

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- ▶ Exactly one of the following two propositions is a theorem.  
Which one?

$$(1) \quad [\forall x \in S : P(x) \vee Q(x)] \Leftrightarrow [\forall x \in S : P(x)] \vee [\forall x \in S : Q(x)]$$

$$(2) \quad [\exists x \in S : P(x) \vee Q(x)] \Leftrightarrow [\exists x \in S : P(x)] \vee [\exists x \in S : Q(x)]$$

- ▶ Justify your answer