

CS52 MACHINE:  
RECURSION

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CS 52 – Fall 2022

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Admin

Grading

2

Getting involved in the department

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Get on the colloquium mailing list:

- ▣ Go to: <https://listserv.pomona.edu/scripts/wa.exe?A0=CSCOLLOQ>
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Questions... reach out to the liaisons:  
[liaisons@cs.pomona.edu](mailto:liaisons@cs.pomona.edu)

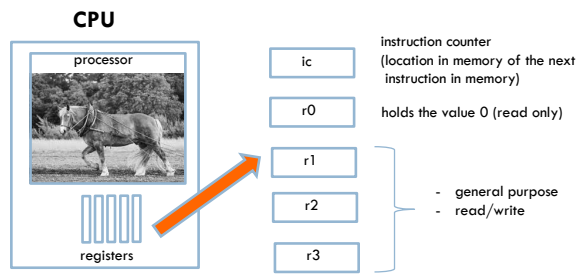
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Examples from this lecture

<https://cs.pomona.edu/classes/cs54/examples/cs52machine/>

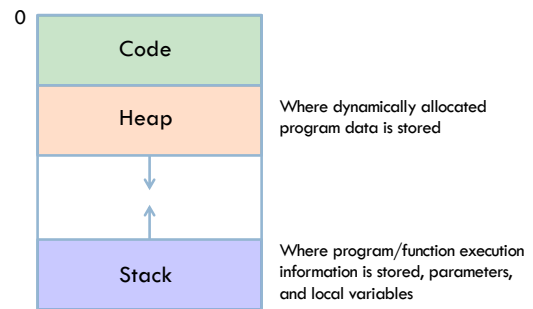
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## CS52 machine



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## Memory layout



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## Stack frame

- Key unit for keeping track of a function call
- return address (where to go when we're done executing)
  - parameters
  - local variables

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## CS52 function call conventions

- r1: reserved for the stack pointer
  - r2: contains the return address
  - r3: contains the first parameter
- additional parameters go on the stack (more on this)
- the result (i.e. the return value) should go in r3

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### Real structure of CS52 program

```

; great comments at the top!
;
  lw r1 stack           Save address of highest end
                        (highest address) of the stack in r1

  instruction1         ; comment
  instruction2         ; comment
  ...
  hlt

;
; stack area: 50 words
;
  dat 100              Reserve 50 words for the stack
stack
  
```

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### Revisit increment example

10

### Sum revisited

```

fun sum x =
  if x <= 0 then
    0
  else
    x + sum (x-1);
  
```

Note to future Dave from past Dave: write the function up on the board 😊

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```

sum
  psh r2           ; save the return address on the stack
  bgt r3 r0 recurse ; check base case
  add r3 r0 0      ; if x <= 0, result is 0
  brs done

recurse
  psh r3           ; save x on the stack
  sub r3 r3 1      ; x = x-1

  lw r2 sum        ; make recursive call
  cal r2 r2        ; sum (x-1), answer should be in r3

  pop r2           ; get x into r2
  add r3 r3 r2     ; r3 = x + sum (x-1)

done
  pop r2           ; get the return address
  jmp r2           ; go back to where we were called from
  
```

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```

sum
  psh r2 ; save the return address on the stack
  bgt r3 r0 recurse ; check base case
  add r3 r0 0 ; if x <= 0, result is 0
  brs done

recurse
  psh r3 ; save x on the stack
  sub r3 r3 1 ; x = x-1

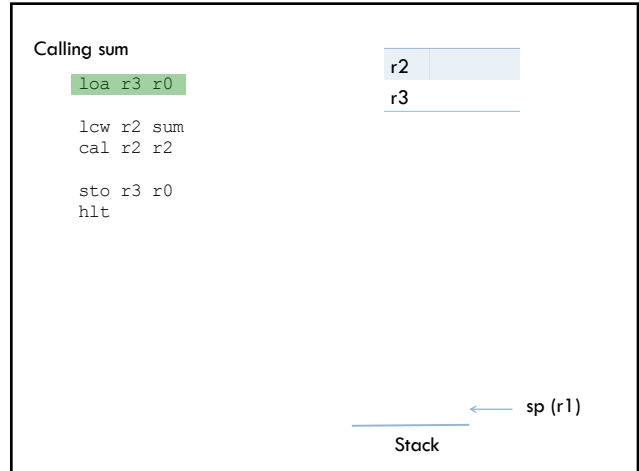
  lclw r2 sum ; make recursive call
  cal r2 r2 ; sum (x-1), answer should be in r3

  pop r2 ; get x into r2
  add r3 r3 r2 ; r3 = x + sum (x-1)

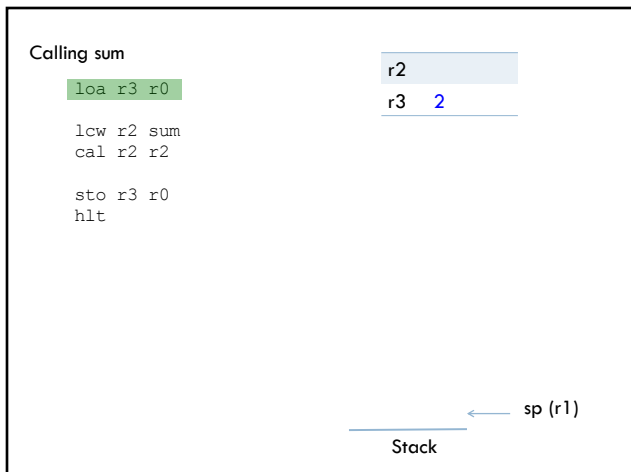
done
  pop r2 ; get the return address
  jmp r2 ; go back to where we were called from
  
```

Notice symmetry of psh and pop

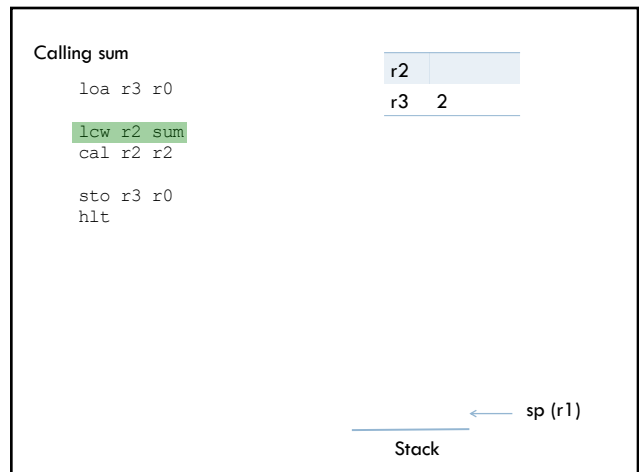
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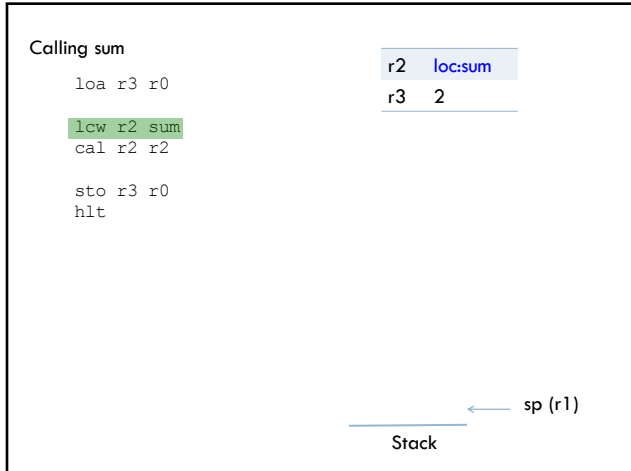
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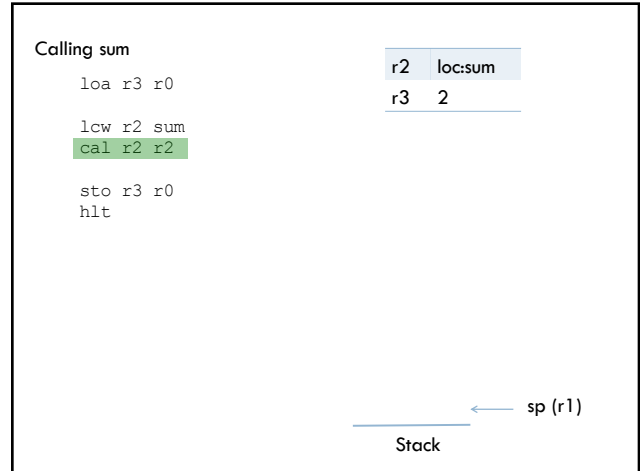
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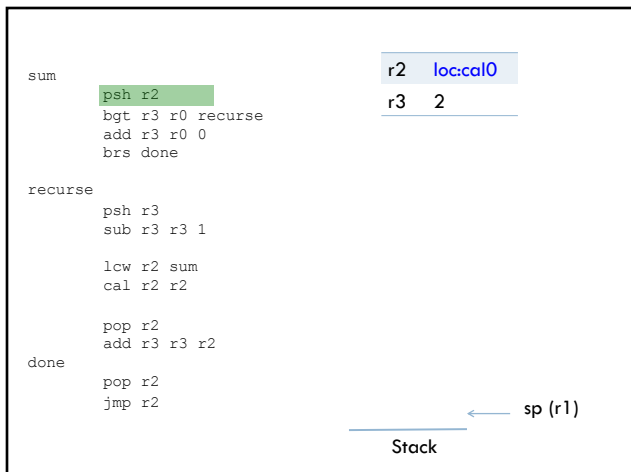
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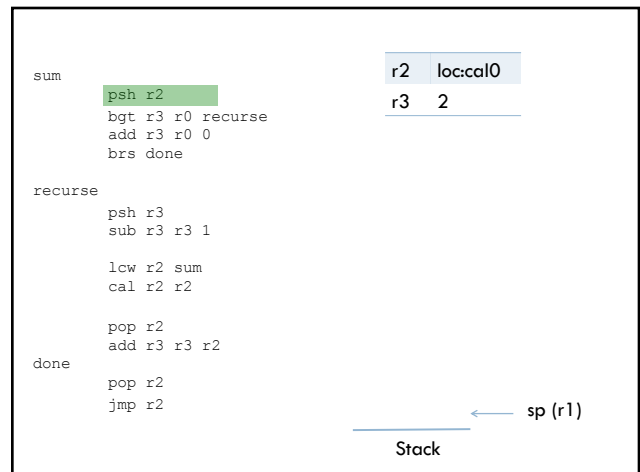
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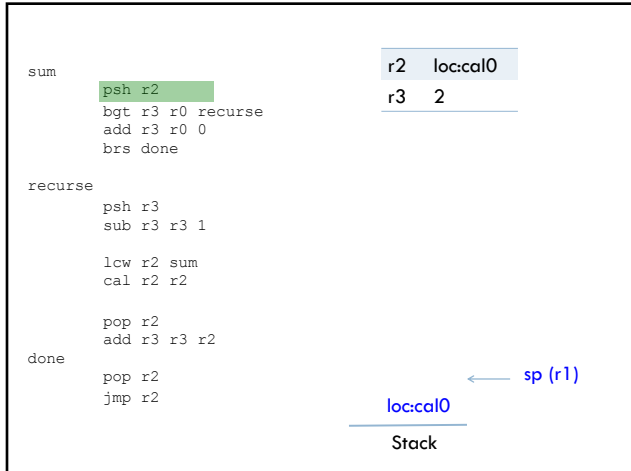
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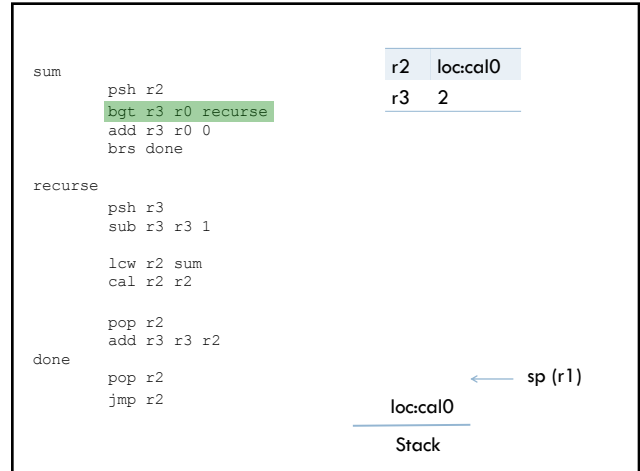
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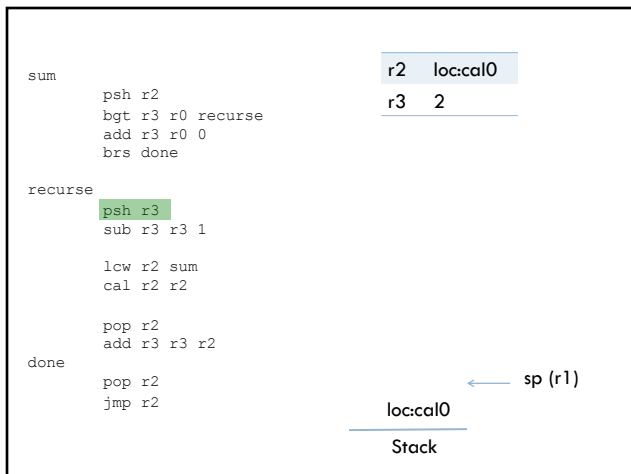
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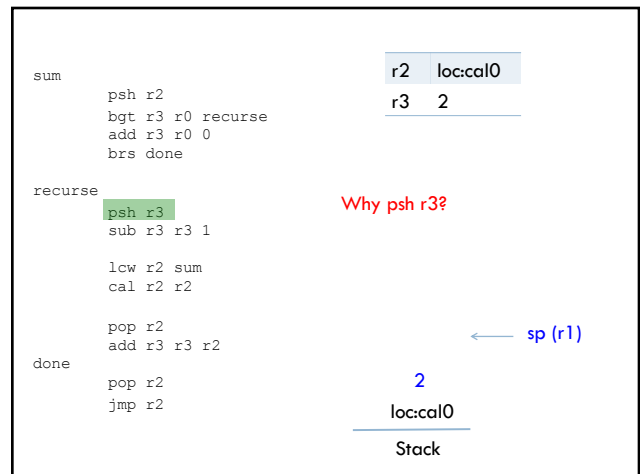
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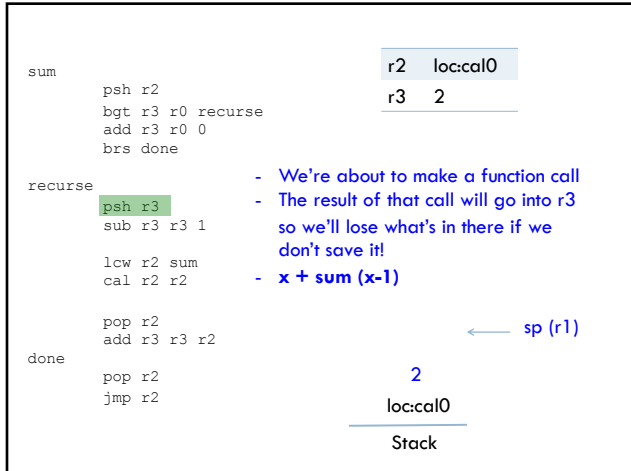
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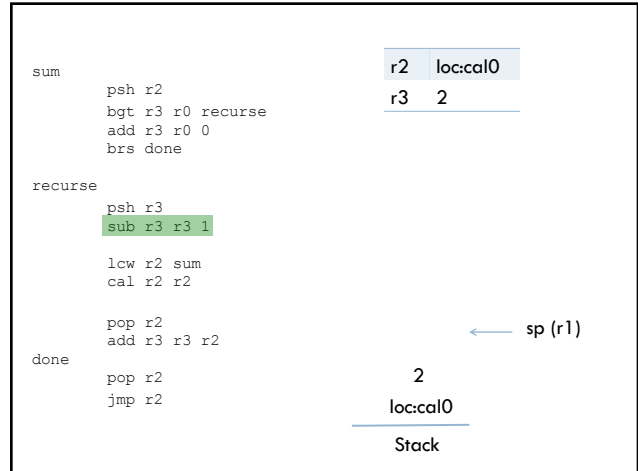
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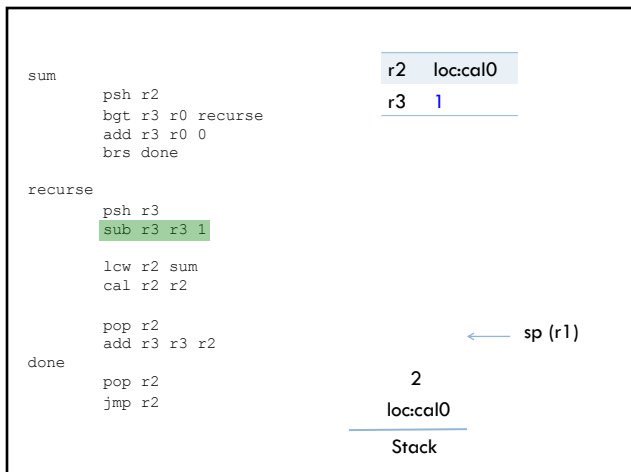
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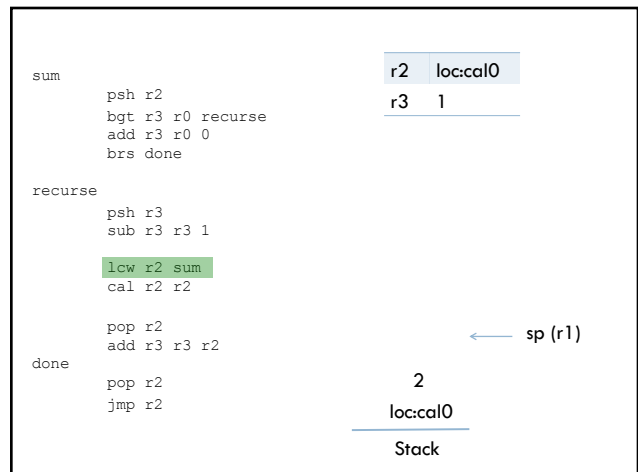
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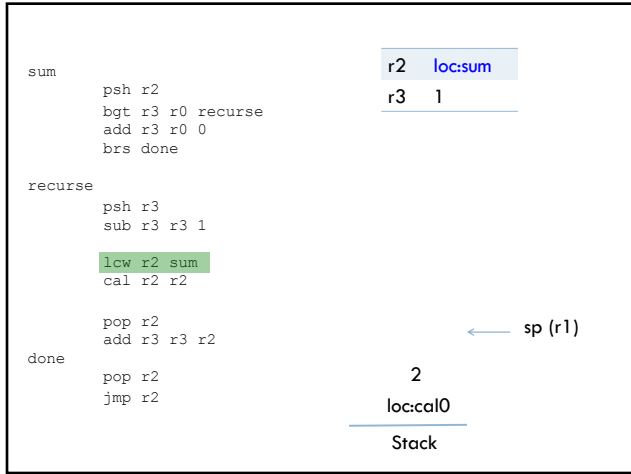
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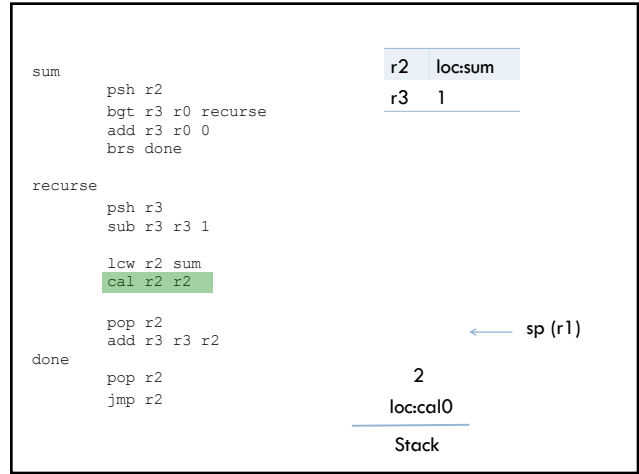
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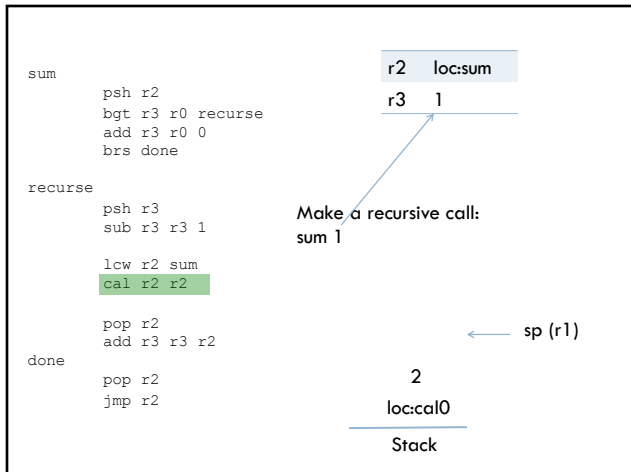
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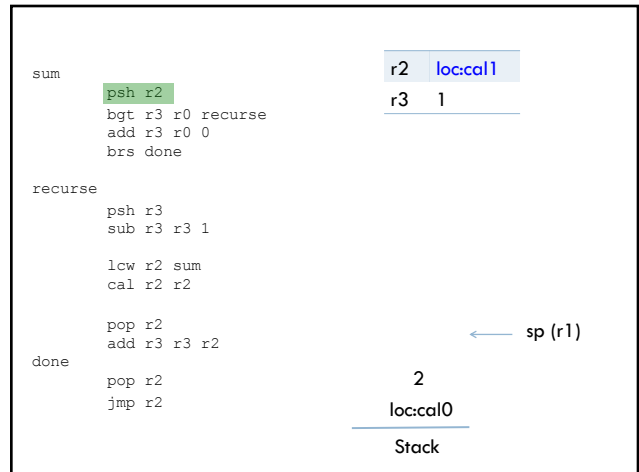
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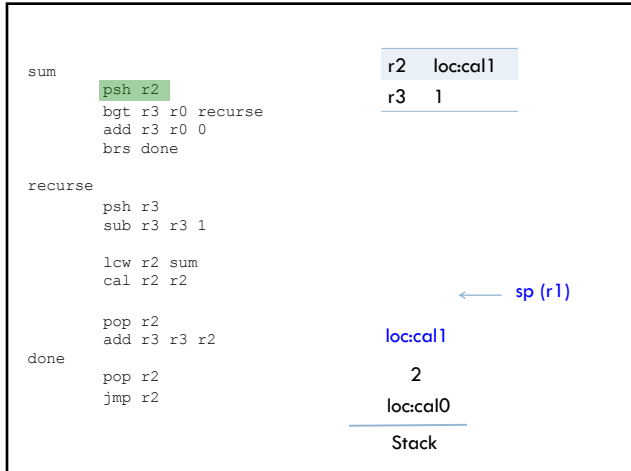


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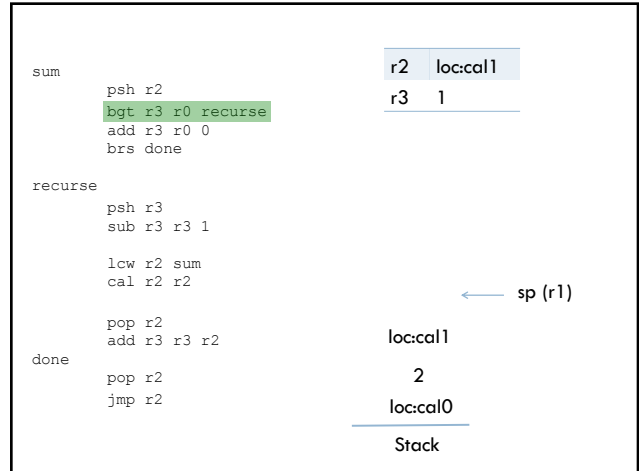


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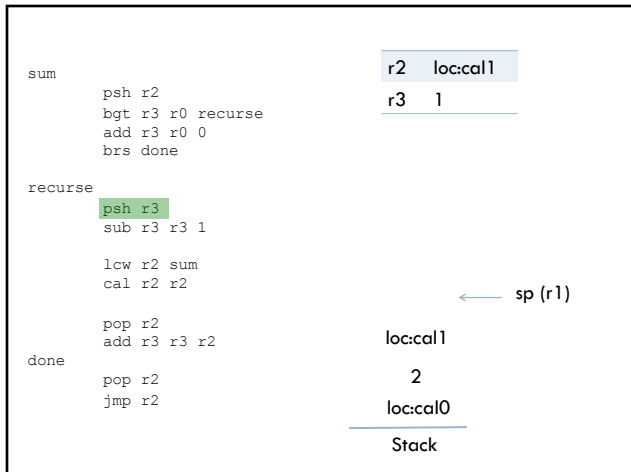




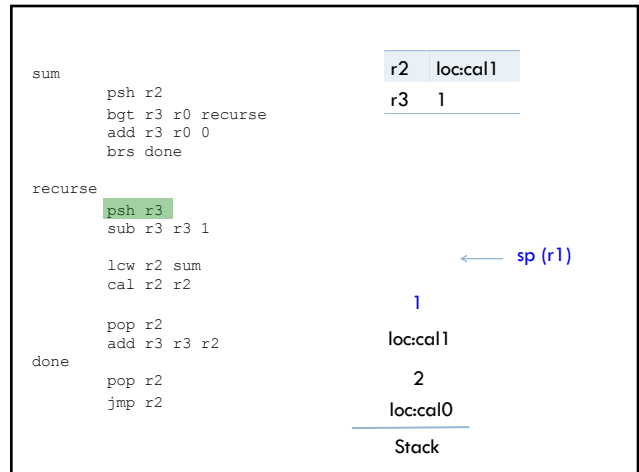
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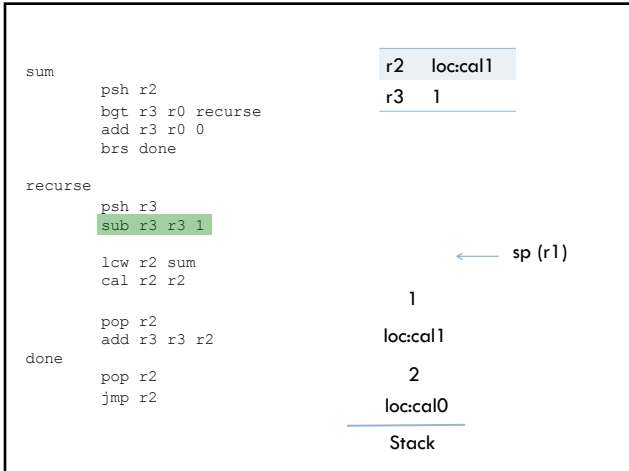
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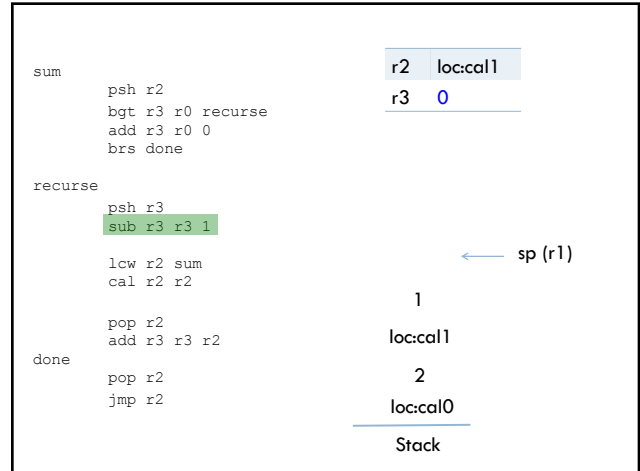
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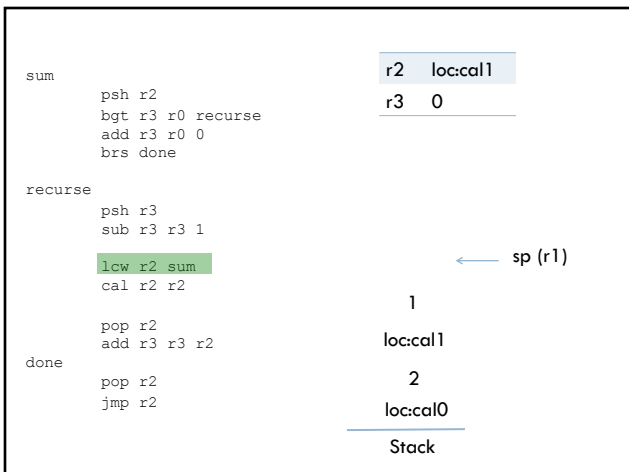
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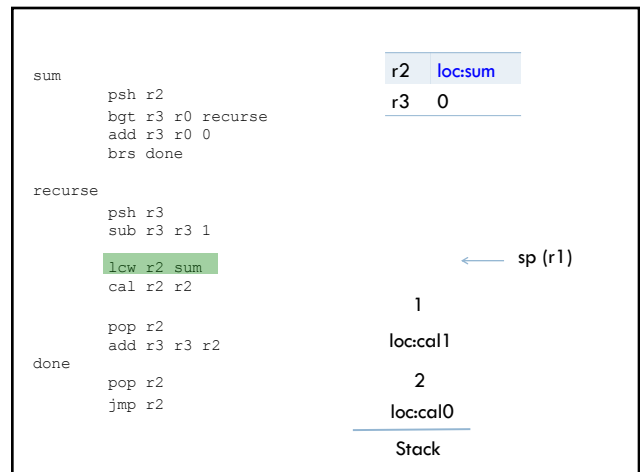
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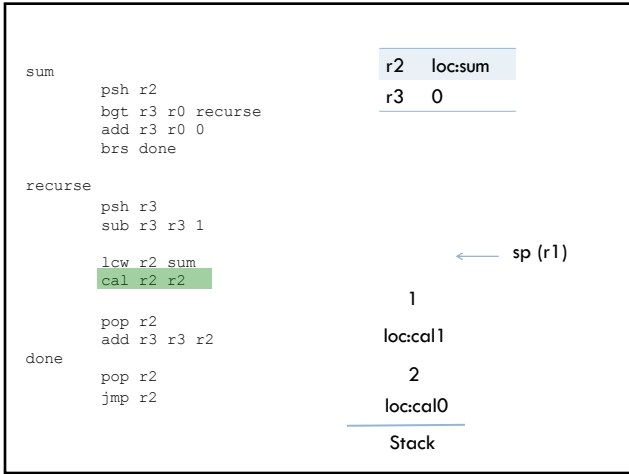
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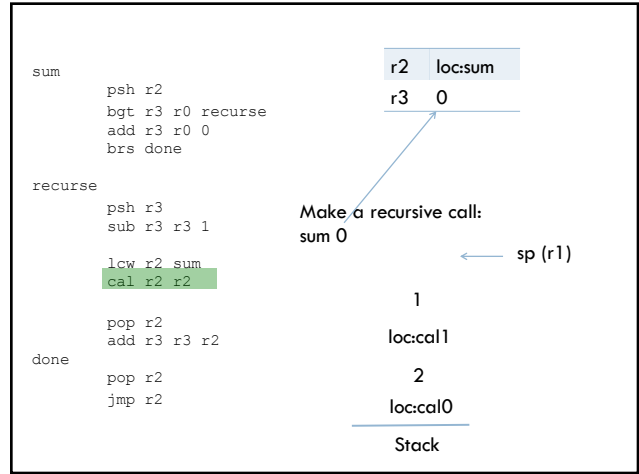
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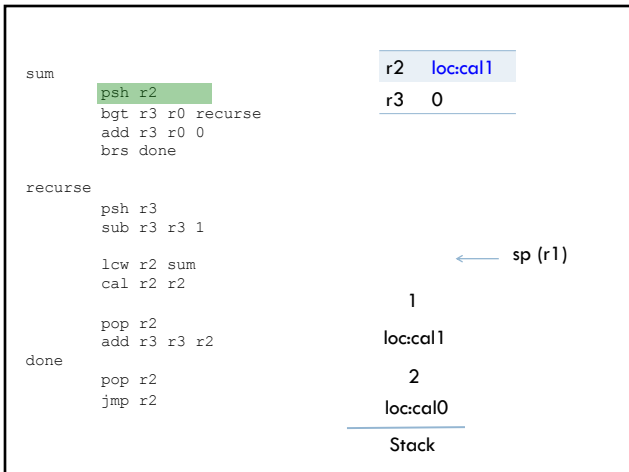
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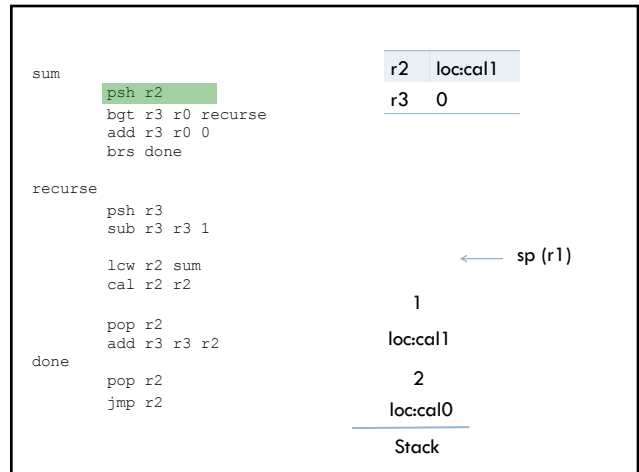
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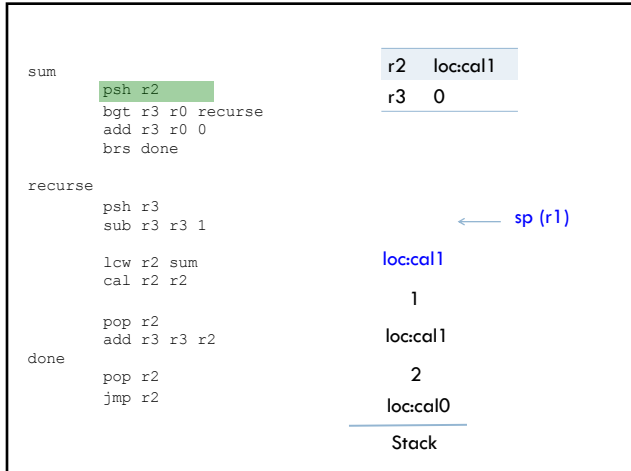
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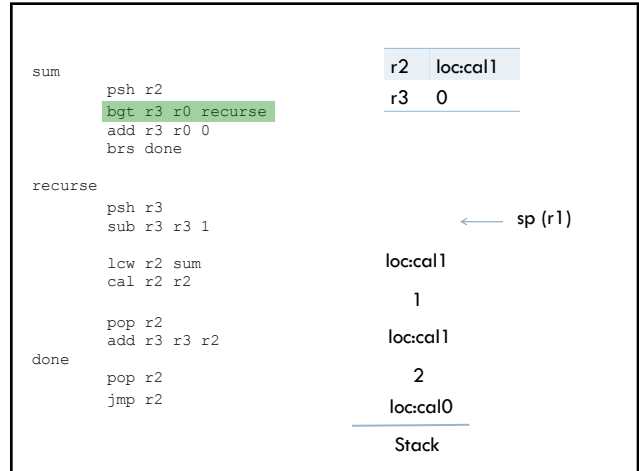
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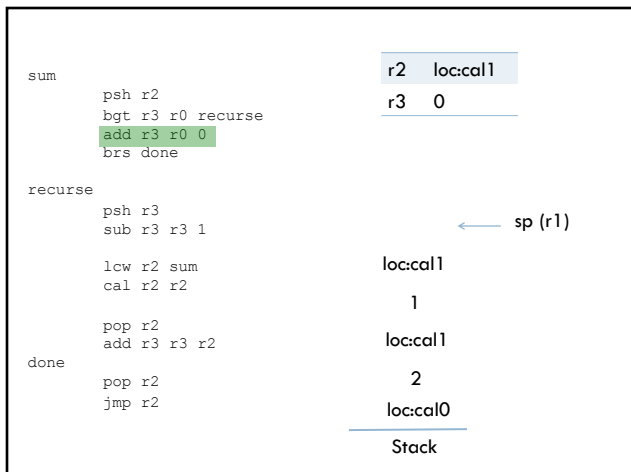
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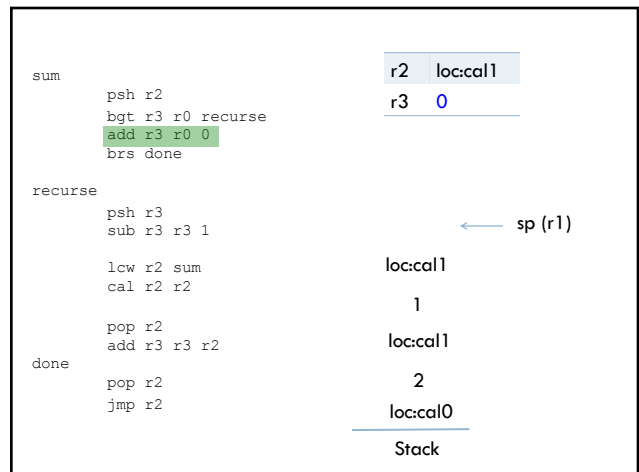
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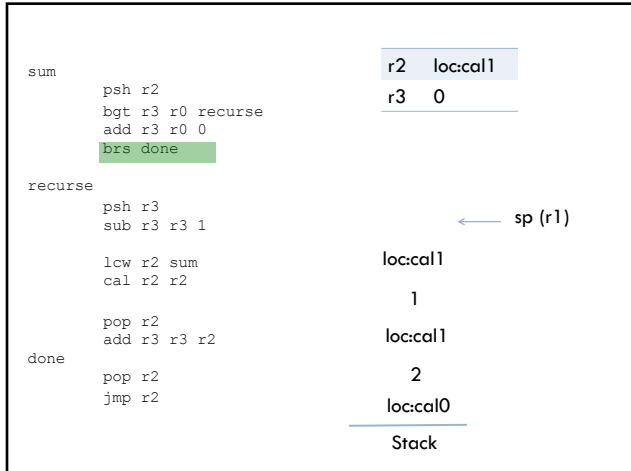
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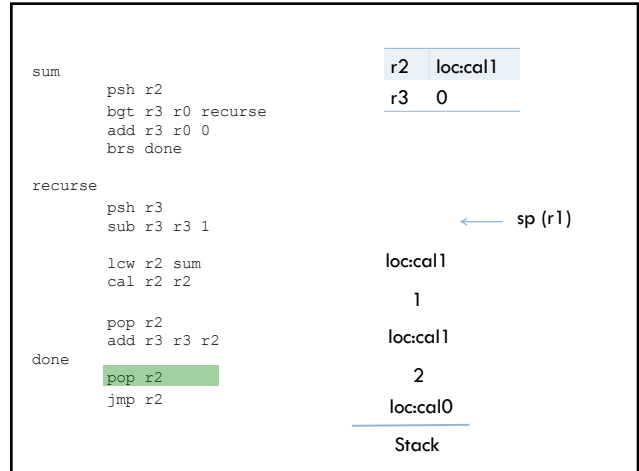
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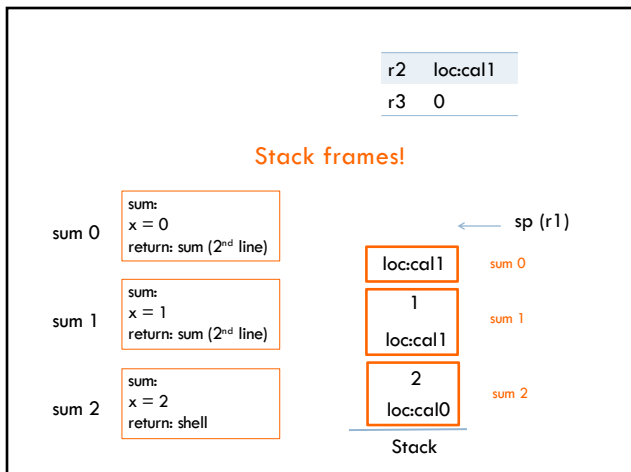
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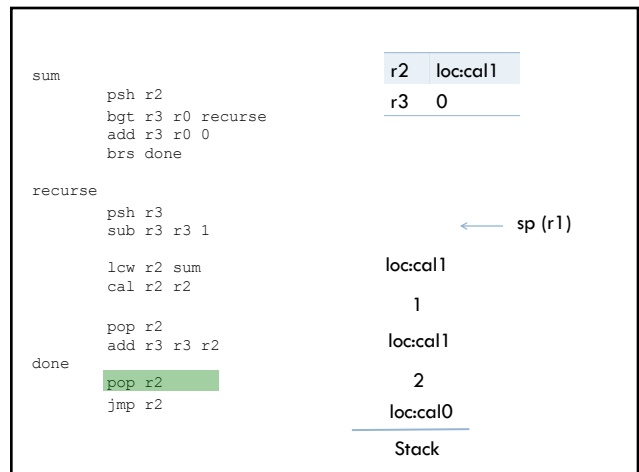
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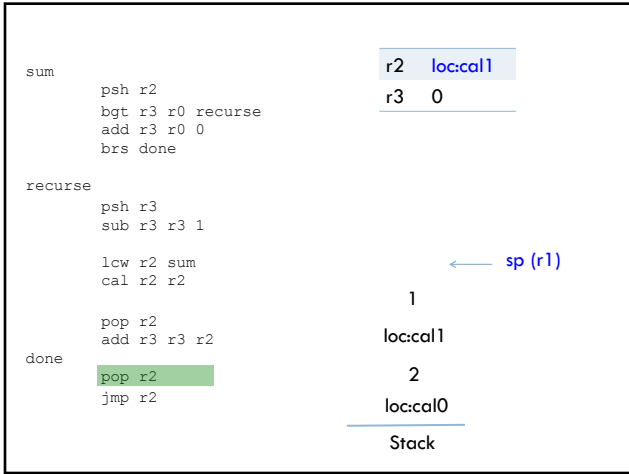
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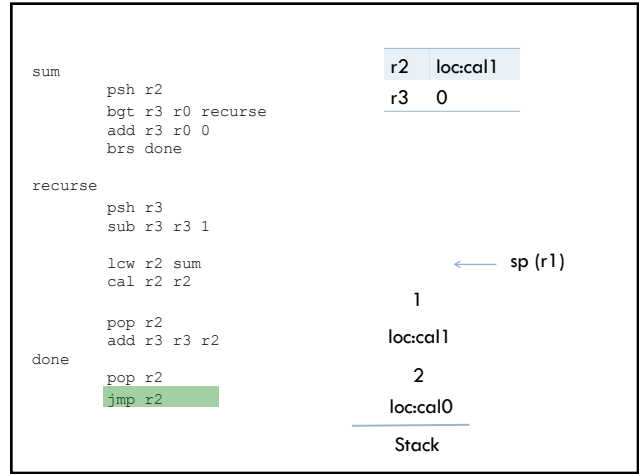
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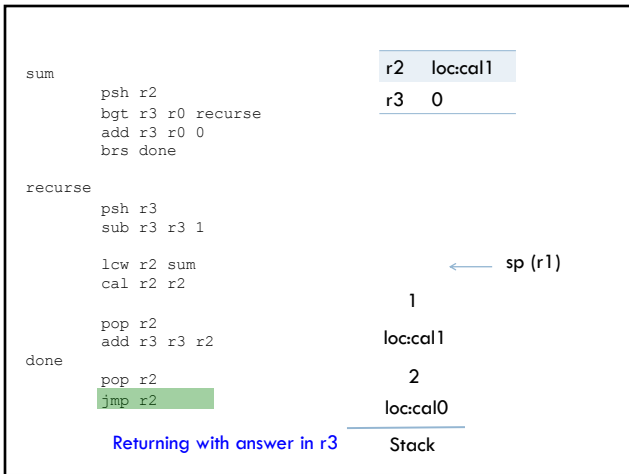
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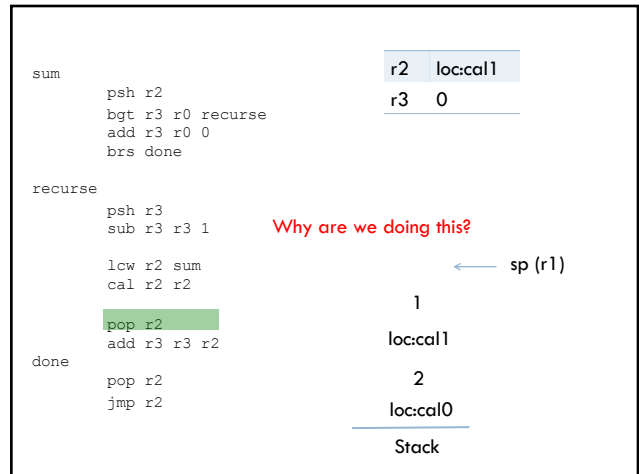
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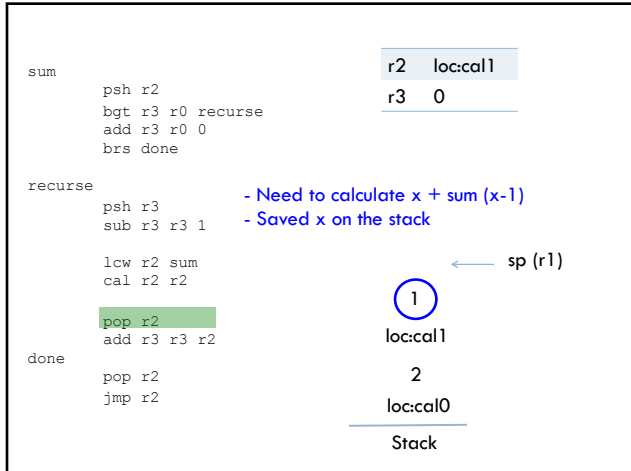
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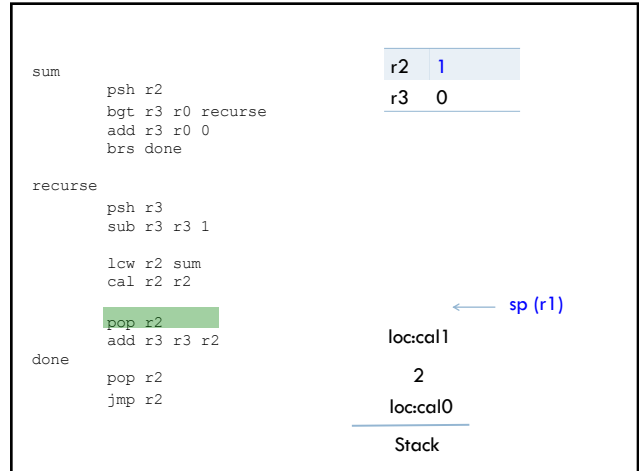
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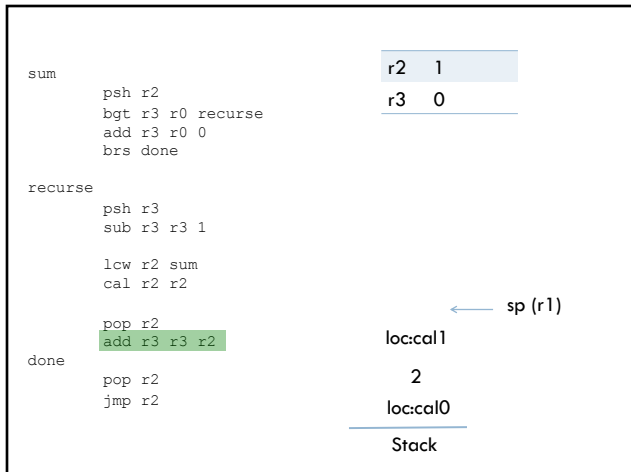
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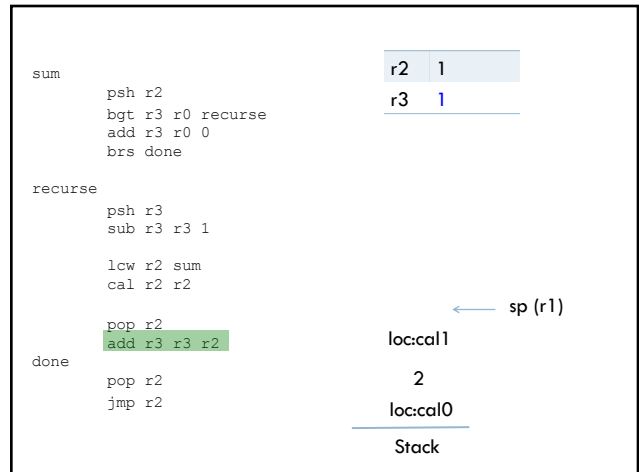
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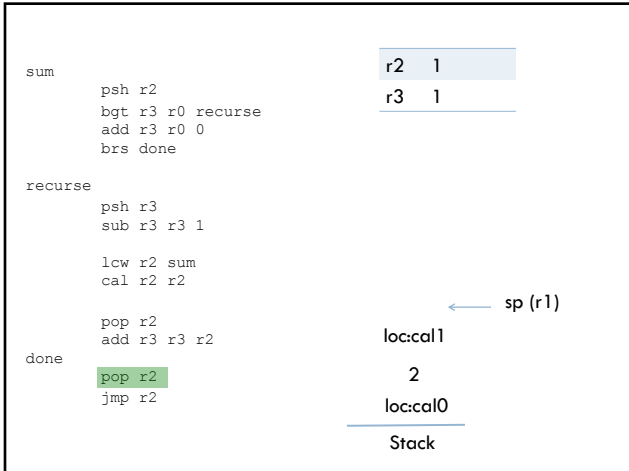
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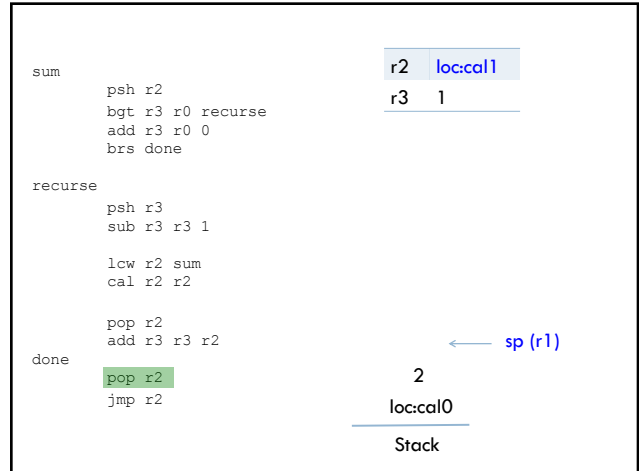
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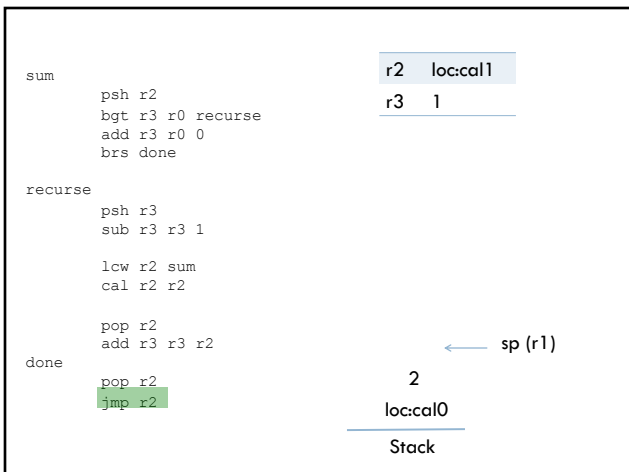
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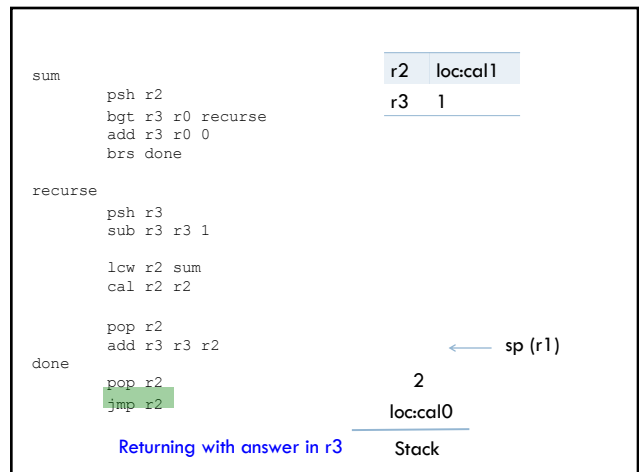
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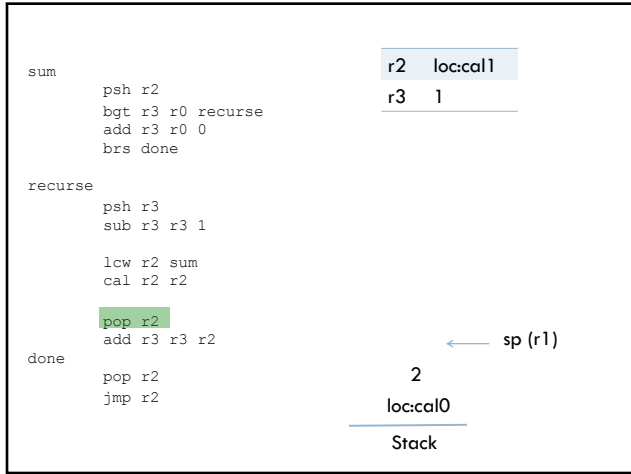


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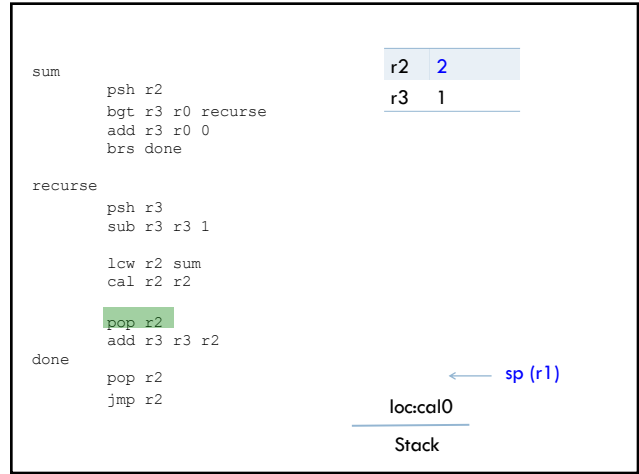


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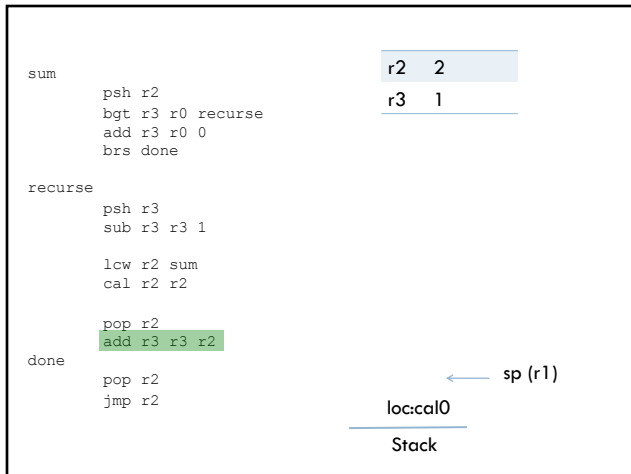




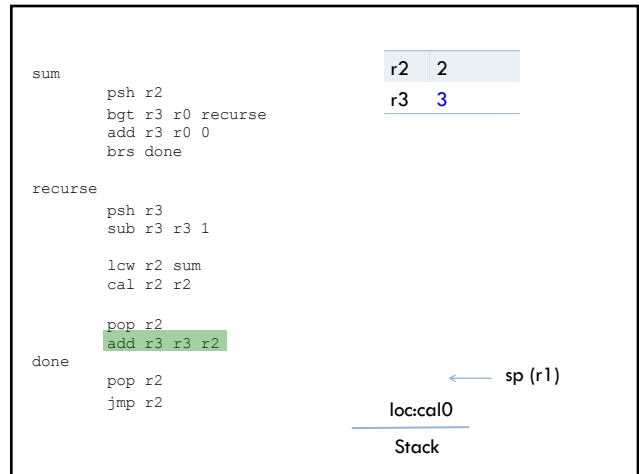
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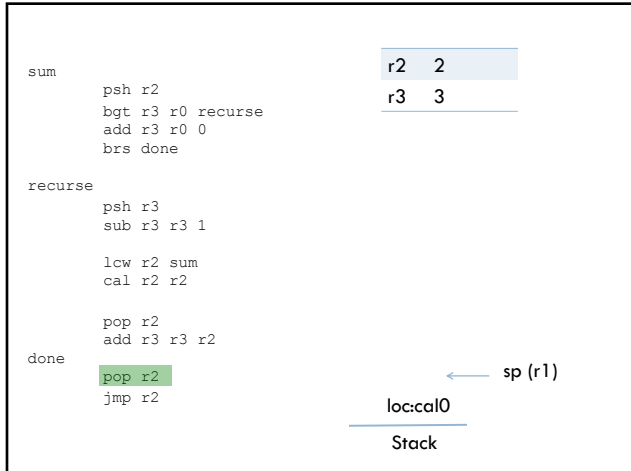
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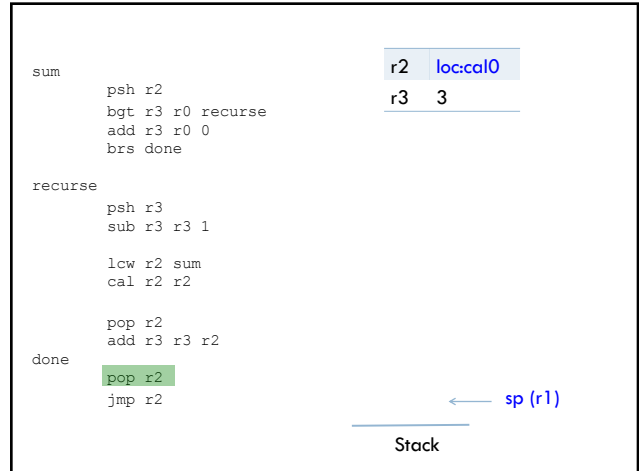
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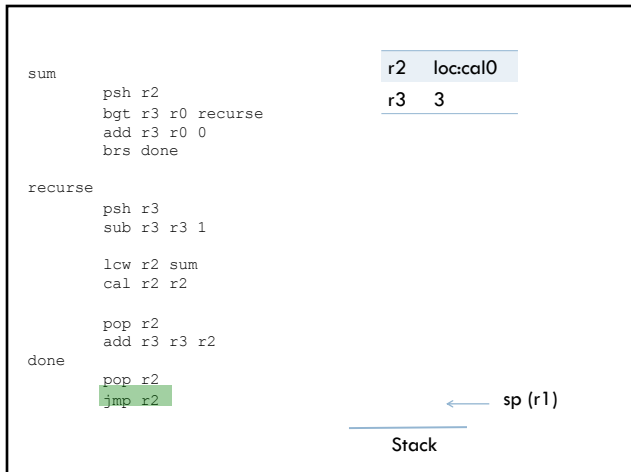
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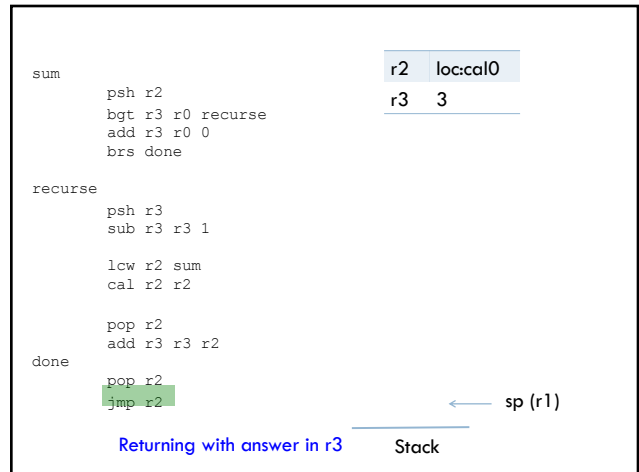
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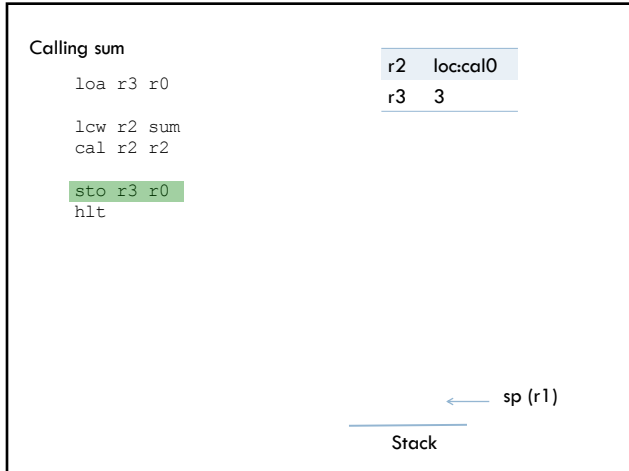
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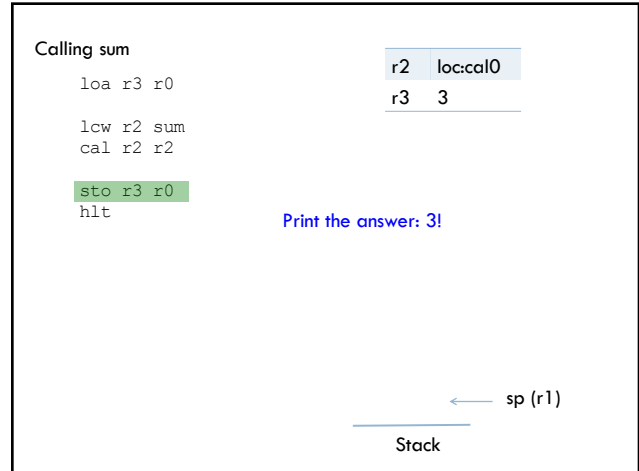
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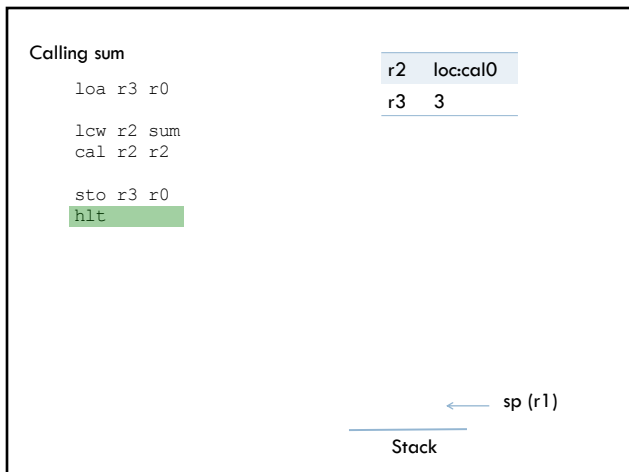
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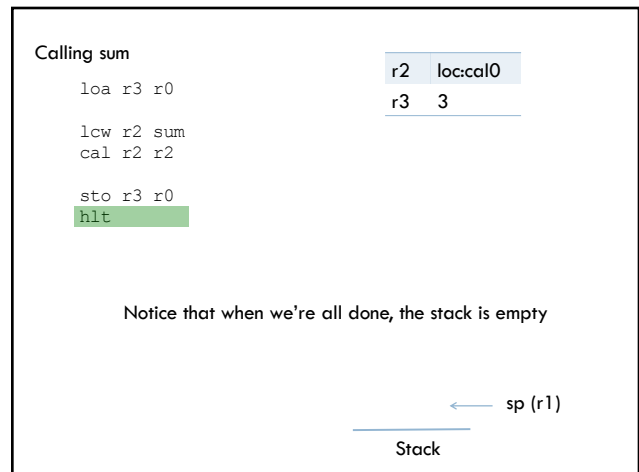
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## Real structure of CS52 program

```

; great comments at the top!
;
    lw r1 stack           Save address of highest end
                          (highest address) of the stack in r1

    instruction1         ; comment
    instruction2         ; comment
    ...

;
; stack area: 50 words
;
    dat 100              } Reserve 50 words for the stack
stack

```

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## Look at sum example

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## Structure of a single parameter function

```

fname
    psh r2               ; save return address on stack
    ...
                          ; do work using r3 as argument
                          ; put result in r3
    pop r2              ; restore return address from stack
    jmp r2              ; return to caller

```

### conventions:

- argument is in r3
- r1 is off-limits since it's used for the stack pointer
- return value goes in r3

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## Functions with multiple arguments

```

fname
    psh r2               ; save return address on stack
    loa r2 r1 4         ; load the second parameter into r2
    ...
                          ; do work using r3 and r2 as arguments
                          ; put result in r3
    pop r2              ; restore return address from stack
    jmp r2              ; return to caller

```

### conventions:

- first argument is in r3
- r1 is off-limits since it's used for the stack pointer
- additional arguments are put on the stack!
- return value goes in r3

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## Functions with multiple arguments

```
fname
    psh r2          ; save return address on stack
    loa r2 r1 4    ; load the second parameter into r2
    ...           ; do work using r3 and r2 as arguments
                ; put result in r3
    pop r2         ; restore return address from stack
    jmp r2        ; return to caller
```

$\text{loa } R_a R_b: R_a = \text{mem}[R_b]$

$\text{loa } R_a R_b S: R_a = \text{mem}[R_b + S]$

What does this operation do? What is the 4?

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## Functions with multiple arguments

```
fname
    psh r2          ; save return address on stack
    loa r2 r1 4    ; load the second parameter into r2
    ...           ; do work using r3 and r2 as arguments
                ; put result in r3
    pop r2         ; restore return address from stack
    jmp r2        ; return to caller
```

$\text{loa } R_a R_b: R_a = \text{mem}[R_b]$

$\text{loa } R_a R_b S: R_a = \text{mem}[R_b + S]$

- r1 is the stack pointer and points at the top (next) slot
- stacks grow towards smaller memory values

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## Functions with multiple arguments

```
fname
    psh r2          ; save return address on stack
    loa r2 r1 4    ; load the second parameter into r2
    ...           ; do work using r3 and r2 as arguments
                ; put result in r3
    pop r2         ; restore return address from stack
    jmp r2        ; return to caller
```

- r1 is the stack pointer and points at the top (next) slot
- stacks grow towards smaller memory values
- r1+2 is then the top value of the stack
- r1+4 is the 2<sup>nd</sup> value of the stack

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## Calling functions with multiple arguments

```
; put the first argument into r3
; put the second argument into r2
psh r2          ; push the second argument onto the stack
lcw r2 func    ; setup a call to the two-parameter function "func"
cal r2 r2      ; call the function
pop r0         ; cleanup the stack!
                ; if you wanted the value, you could also do pop r2
                ; which puts the value in r2
```

- Push the second argument onto the stack
- Make sure that after the function returns, you pop off the second argument
  - If you don't need the value, you can pop it to r0
  - If you want the value, you can pop it to r2 (r3 will have the return value of the function)

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## multiply

look at `mult_easy.a52` code

we can import libraries (really just functions in other files) using the “inc” command

the included files must be in the same directory as the .a52 file running

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## CS52 programming advice

1. Match your psh and pops
2. Follow the register conventions
3. Develop code incrementally
4. Debugging: write out stack, registers, etc. on paper and compare against system execution

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## More examples

I didn't have time to cover the next two examples in class, but left it in the notes as examples of functions that take two parameters

The first example is `multiply` (which is recursive) and the second example is `max` (which is not recursive).

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## Another recursive example

```
int mystery(int a, int b){
    if( b <= 0 ){
        return 0
    }
    else
        return a + mystery(a, b-1)
}
```

What does this function do?

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## Recursion

```

int mystery(int a, int b){
    if( b <= 0 ){
        return 0
    }
    else
        return a + mystery(a, b-1)
}
    
```

Multiplication...  $a*b$  (assuming  $b$  is positive)

Note to future Dave from past Dave: write the function up on the board 😊

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```

mult
  psh r2      ; save the return address
  loa r2 r1 4 ; get at the 2nd argument, b
              ; a = r3, b = r2

  bgt r2 r0 else ; r2 > 0, i.e. recursive case
  add r3 r0 0   ; return 0
  brs endif

else
  sub r2 r2 1 ; r2 = b-1

  psh r3      ; save first argument, a, on stack
              ; (it's going to get overwritten by the return!)
  psh r2      ; add r2 as 2nd argument
  lwc r2 mult ; call mult recursively
  cal r2 r2
  pop r0      ; pop 2nd argument off stack and throw away

  pop r2      ; pop 'a' into r2 off of the stack
  add r3 r3 r2 ; r3 = a + mult(a, b-1)
endif
pop r2      ; get the return address
jmp r2      ; return
    
```

Function startup

Base case

Recursive call

Recursive case

answer calculation

Function cleanup and return

90

```

mult
  psh r2      ; save the return address
  loa r2 r1 4 ; get at the 2nd argument, b
              ; a = r3, b = r2

  bgt r2 r0 else ; r2 > 0, i.e. recursive case
  add r3 r0 0   ; return 0
  brs endif

else
  sub r2 r2 1 ; r2 = b-1

  psh r3      ; save first argument, a, on stack
              ; (it's going to get overwritten by the return!)
  psh r2      ; add r2 as 2nd argument
  lwc r2 mult ; call mult recursively
  cal r2 r2
  pop r0      ; pop 2nd argument off stack and throw away

  pop r2      ; pop 'a' into r2 off of the stack
  add r3 r3 r2 ; r3 = a + mult(a, b-1)
endif
pop r2      ; get the return address
jmp r2      ; return
    
```

Function startup

if( b <= 0 )  
return 0

mystery(a, b-1)

a + mystery(a, b-1)

Function cleanup and return

91

```

mult
  psh r2      ; save the return address
  loa r2 r1 4 ; get at the 2nd argument, b
              ; a = r3, b = r2

  bgt r2 r0 else ; r2 > 0, i.e. recursive case
  add r3 r0 0   ; return 0
  brs endif

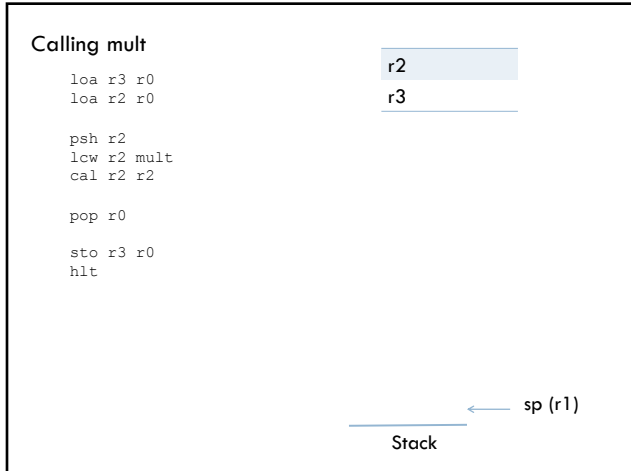
else
  sub r2 r2 1 ; r2 = a-1

  psh r3      ; save first argument, a, on stack
              ; (it's going to get overwritten by the return!)
  psh r2      ; add r2 as 2nd argument
  lwc r2 mult ; call mult recursively
  cal r2 r2
  pop r0      ; pop 2nd argument off stack and throw away

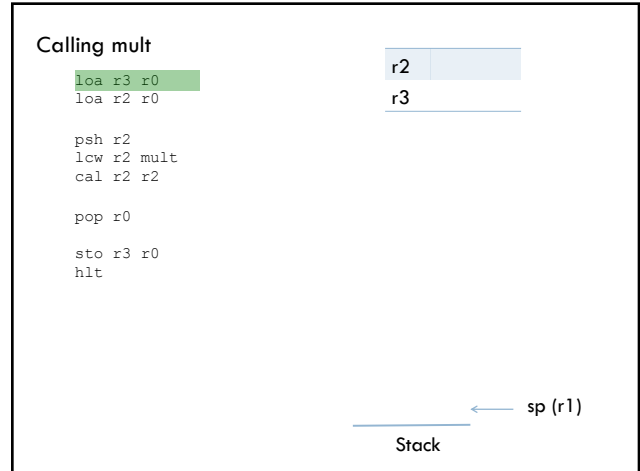
  pop r2      ; load a into r2 off of the stack
  add r3 r3 r2 ; r3 = a + mult(a, b-1)
endif
pop r2      ; get the return address
jmp r2      ; return
    
```

Notice symmetry of psh and pop

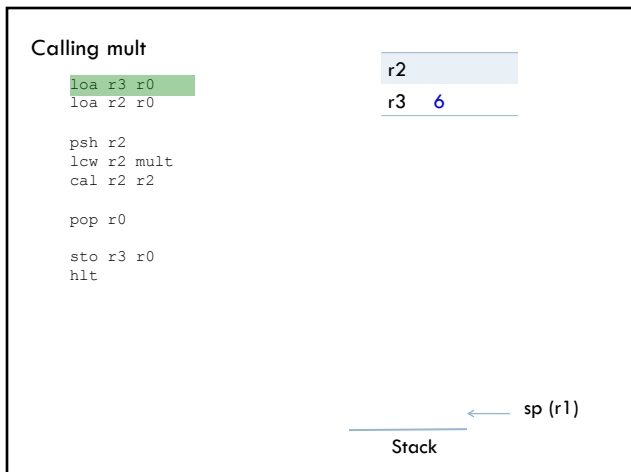
92



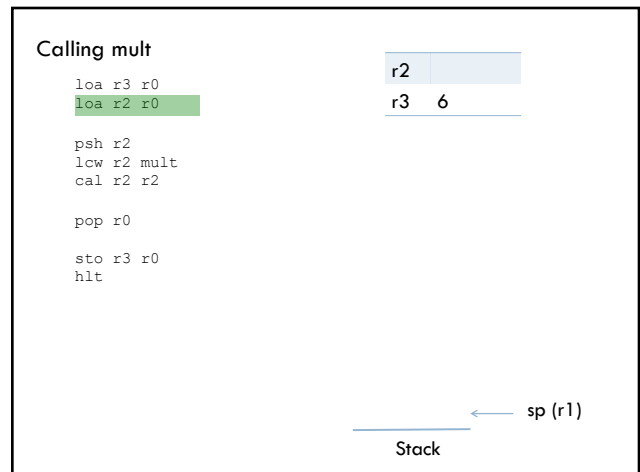
93



94

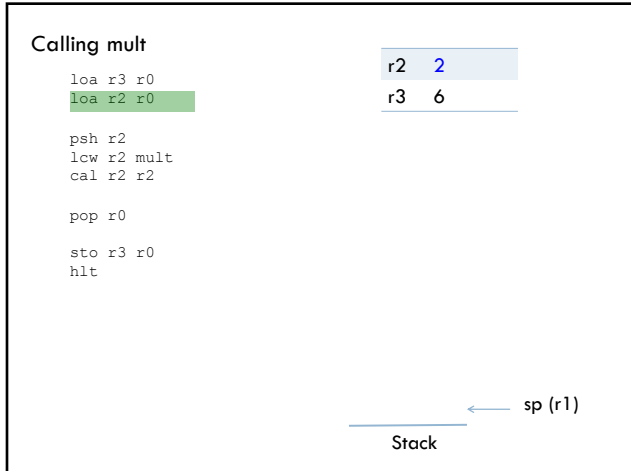


95

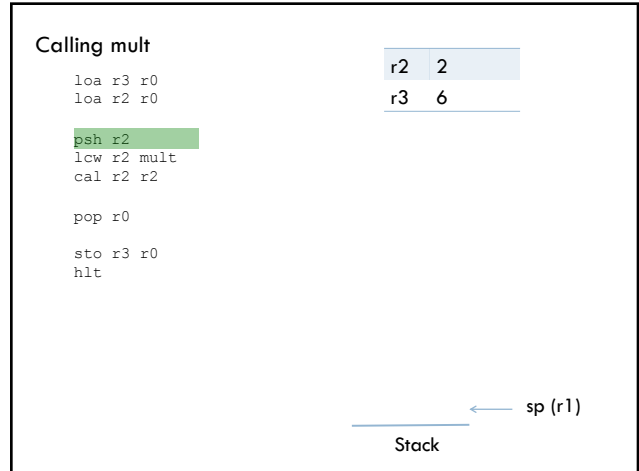


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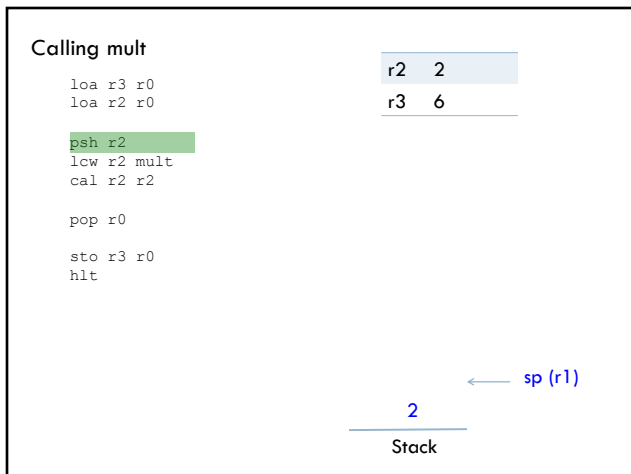




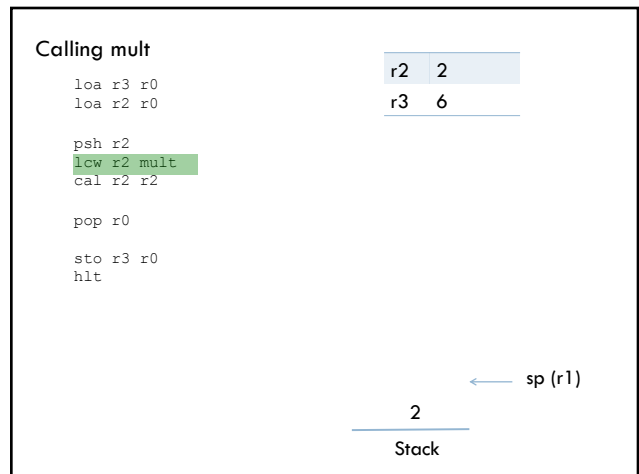
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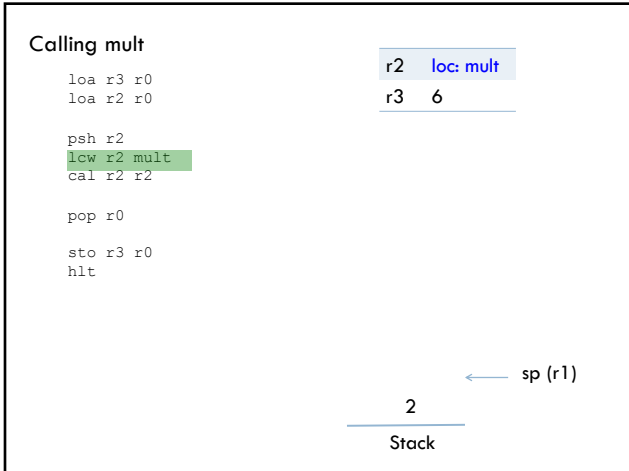
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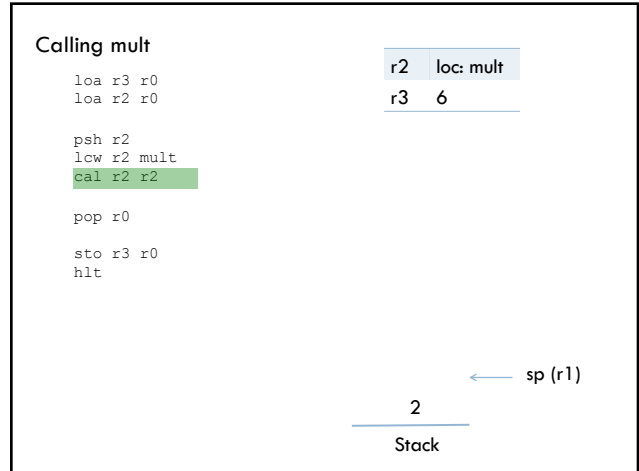
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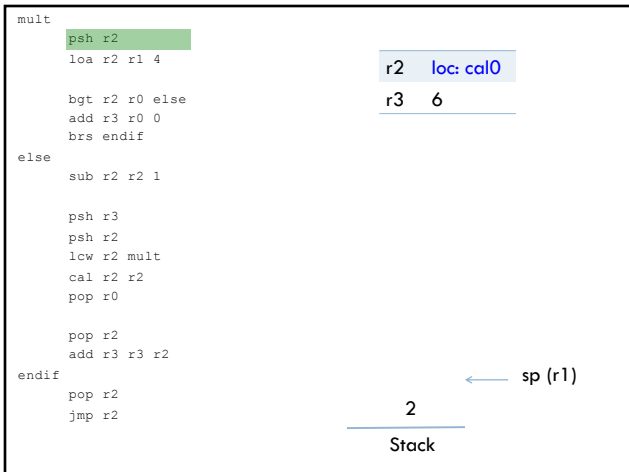
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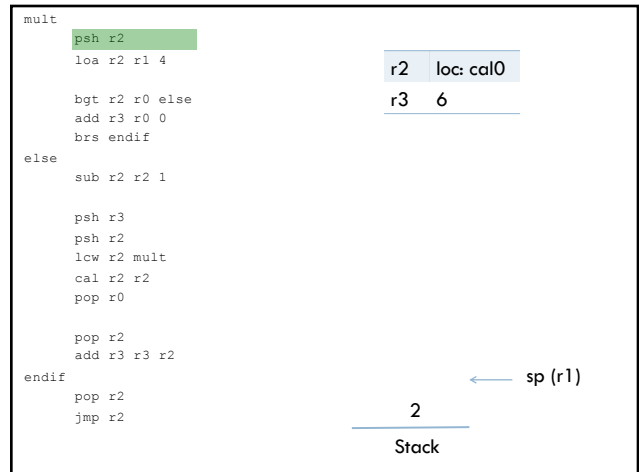
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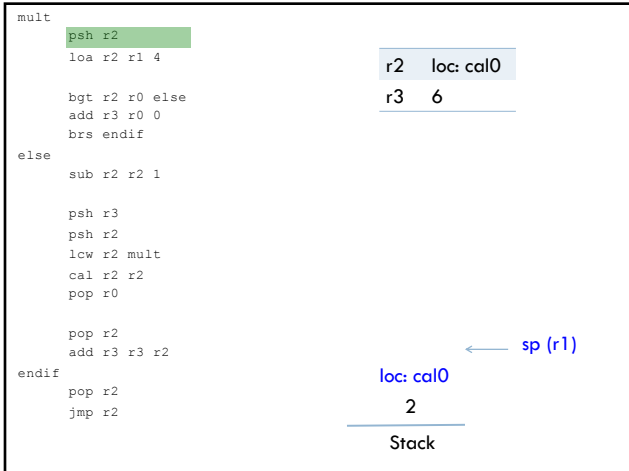
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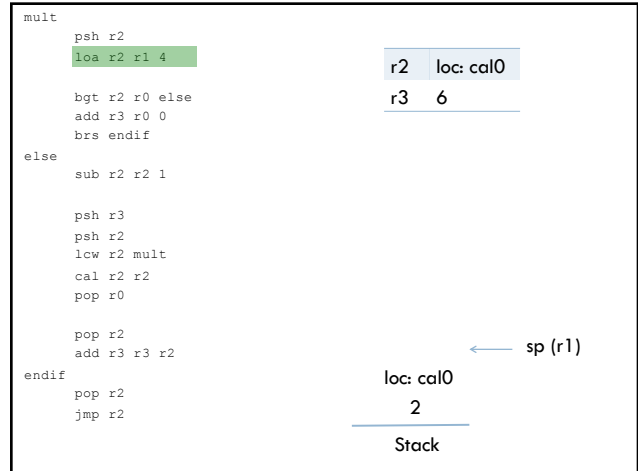
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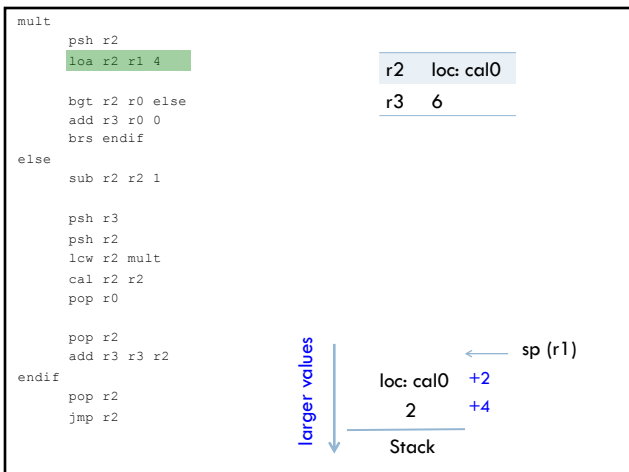
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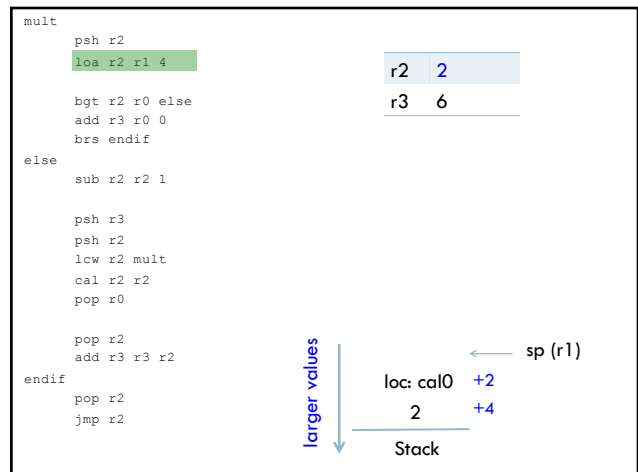
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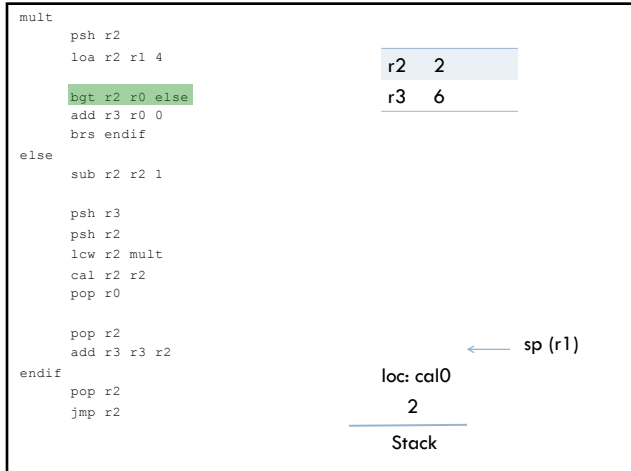
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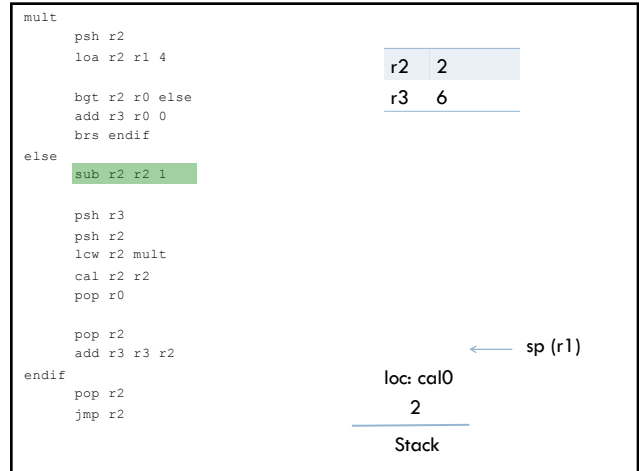
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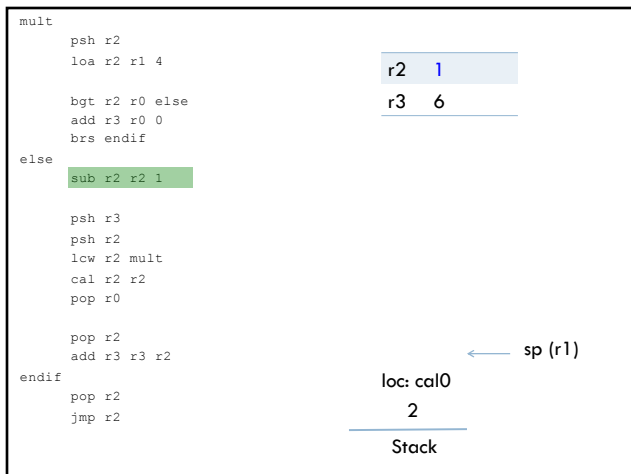
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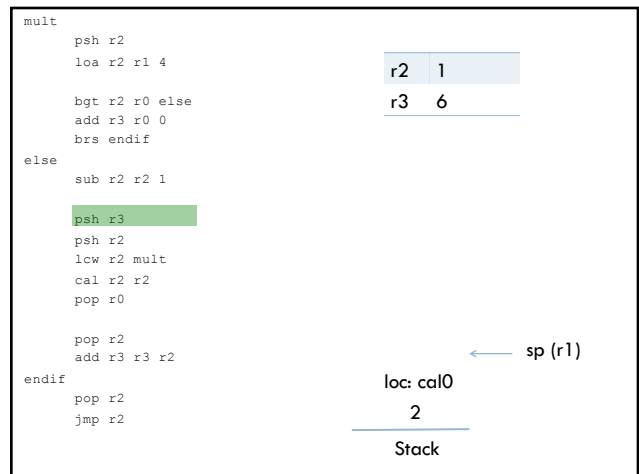
109



110



111



112

```

mult
  psh r2
  loa r2 r1 4
  bgt r2 r0 else
  add r3 r0 0
  brs endif
else
  sub r2 r2 1
  psh r3
  psh r2
  lcw r2 mult
  cal r2 r2
  pop r0

  pop r2
  add r3 r3 r2
endif
pop r2
jmp r2
    
```

r2	1
r3	6

Why psh r3?

← sp (r1)

6

loc: cal0

2

---

Stack

113

```

mult
  psh r2
  loa r2 r1 4
  bgt r2 r0 else
  add r3 r0 0
  brs endif
else
  sub r2 r2 1
  psh r3
  psh r2
  lcw r2 mult
  cal r2 r2
  pop r0

  pop r2
  add r3 r3 r2
endif
pop r2
jmp r2
    
```

r2	1
r3	6

← sp (r1)

6

loc: cal0

2

---

Stack

- We're about to make a function call
- The result of that call will go into r3 so we'll lose what's in there if we don't save it!

114

```

mult
  psh r2
  loa r2 r1 4
  bgt r2 r0 else
  add r3 r0 0
  brs endif
else
  sub r2 r2 1
  psh r3
  psh r2
  lcw r2 mult
  cal r2 r2
  pop r0

  pop r2
  add r3 r3 r2
endif
pop r2
jmp r2
    
```

r2	1
r3	6

← sp (r1)

6

loc: cal0

2

---

Stack

115

```

mult
  psh r2
  loa r2 r1 4
  bgt r2 r0 else
  add r3 r0 0
  brs endif
else
  sub r2 r2 1
  psh r3
  psh r2
  lcw r2 mult
  cal r2 r2
  pop r0

  pop r2
  add r3 r3 r2
endif
pop r2
jmp r2
    
```

r2	1
r3	6

← sp (r1)

1

6

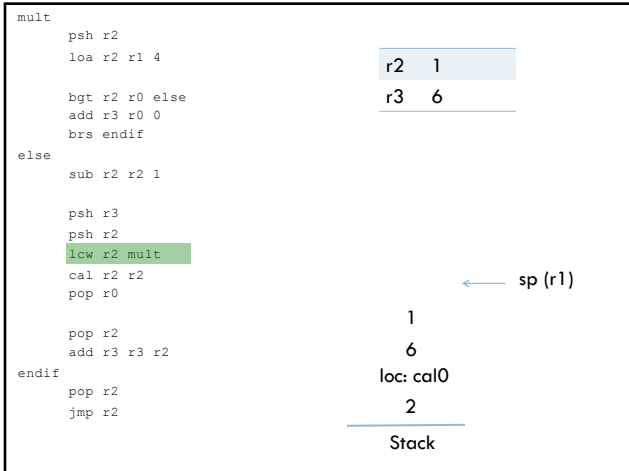
loc: cal0

2

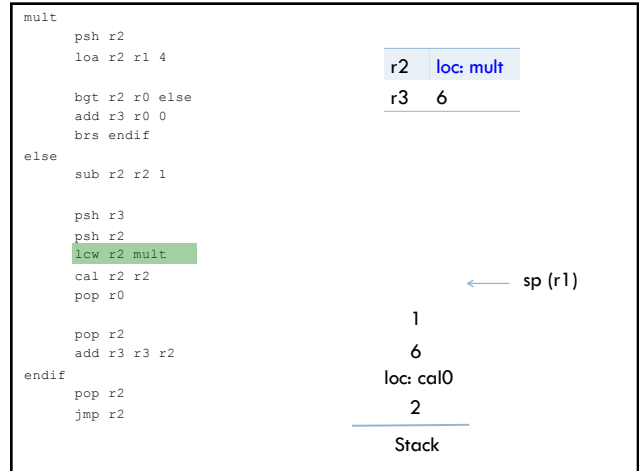
---

Stack

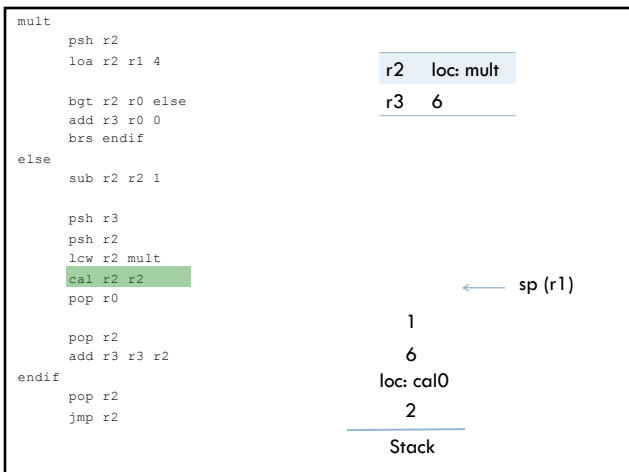
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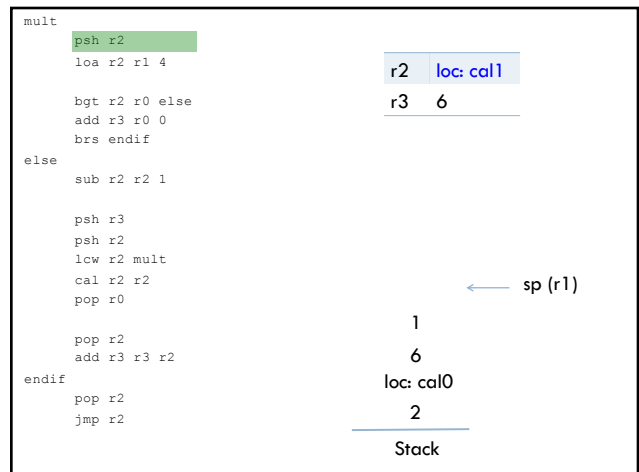
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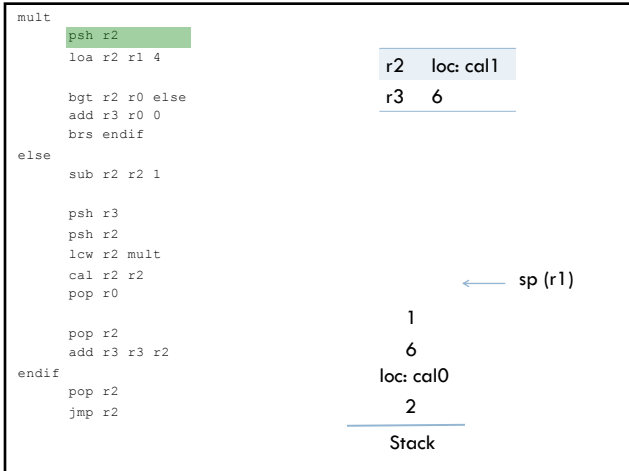
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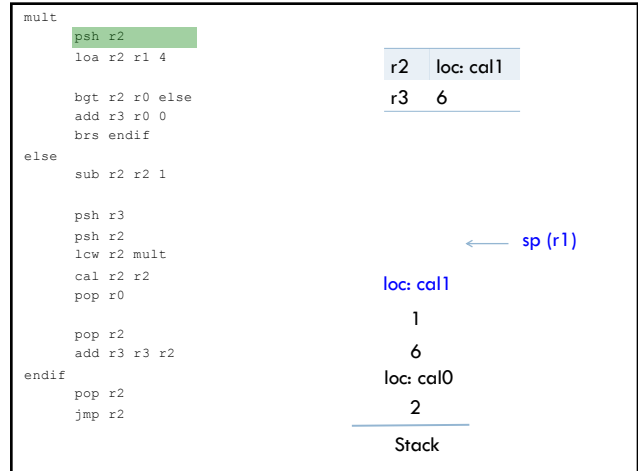
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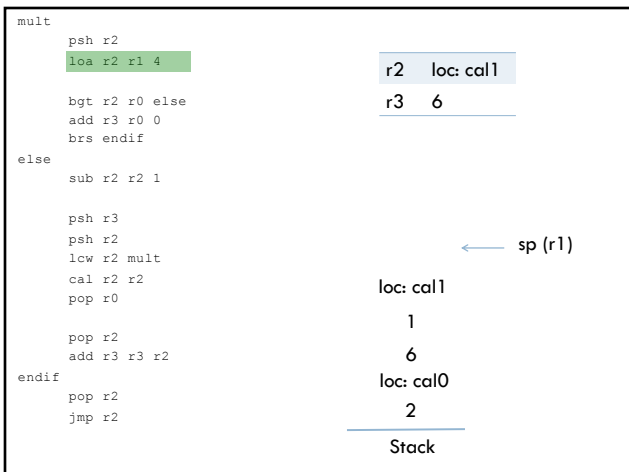
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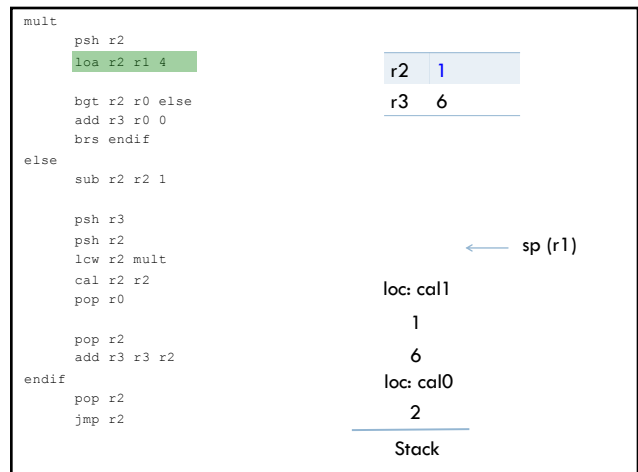
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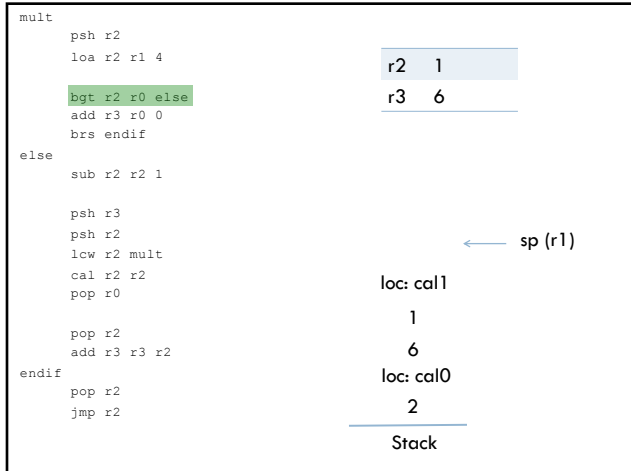
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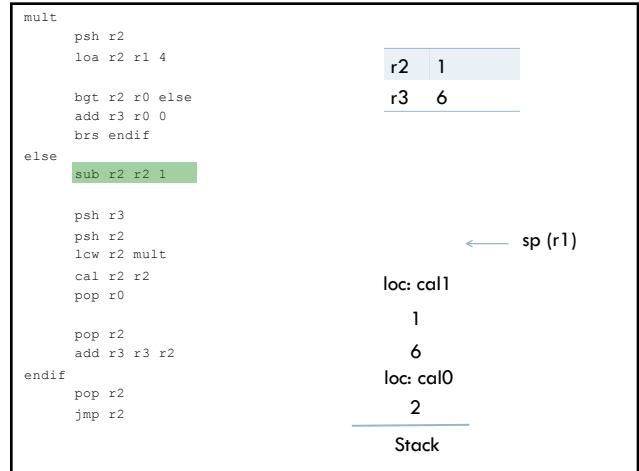
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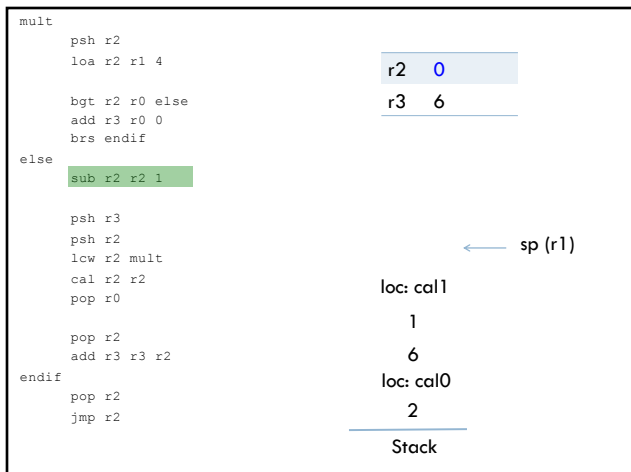
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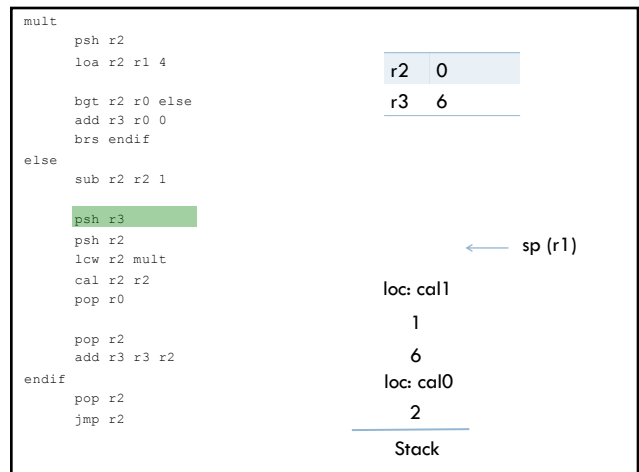
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126

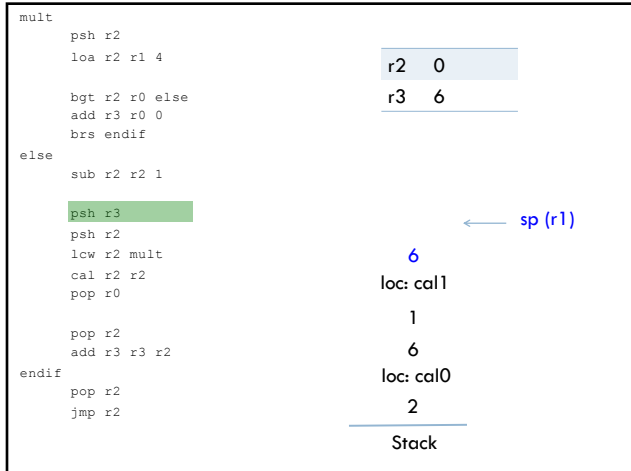


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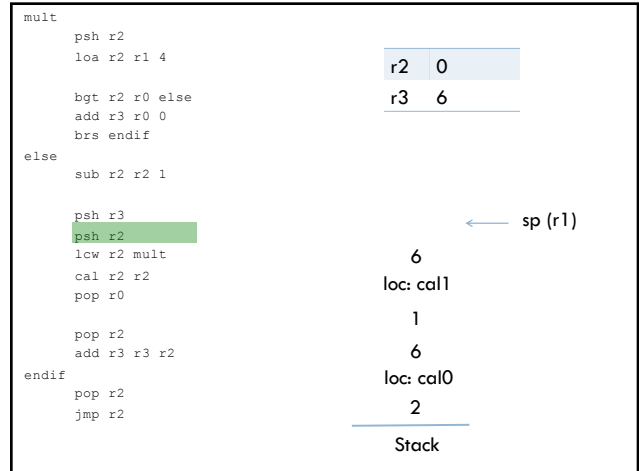


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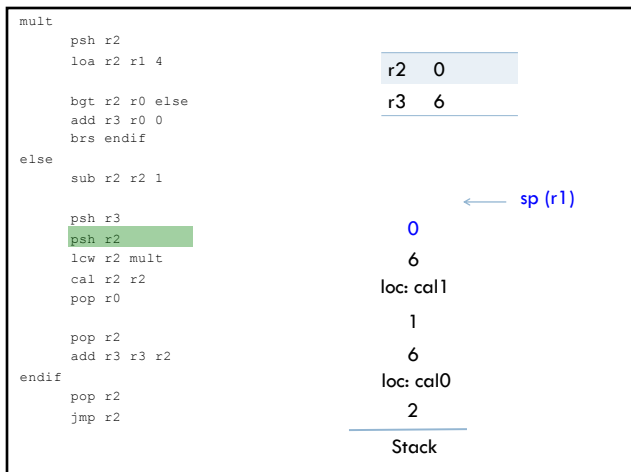




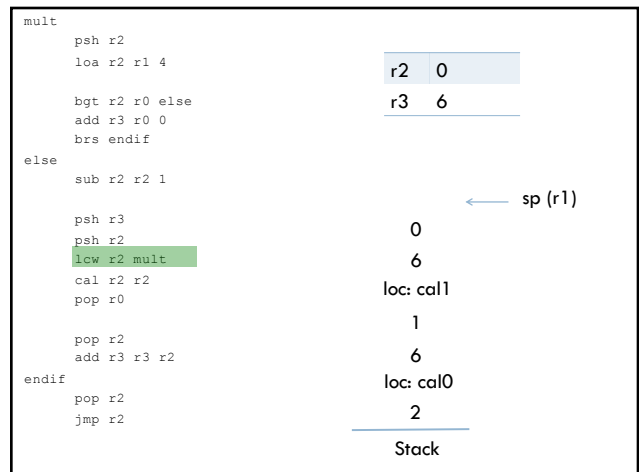
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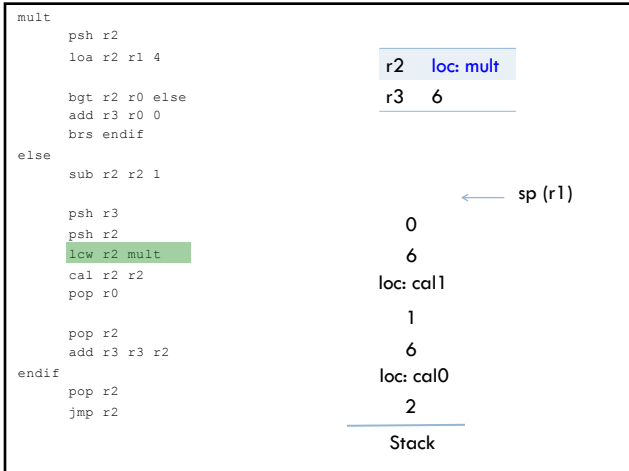
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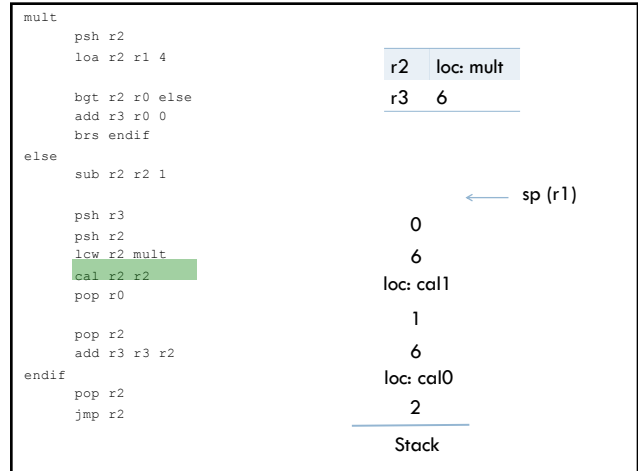
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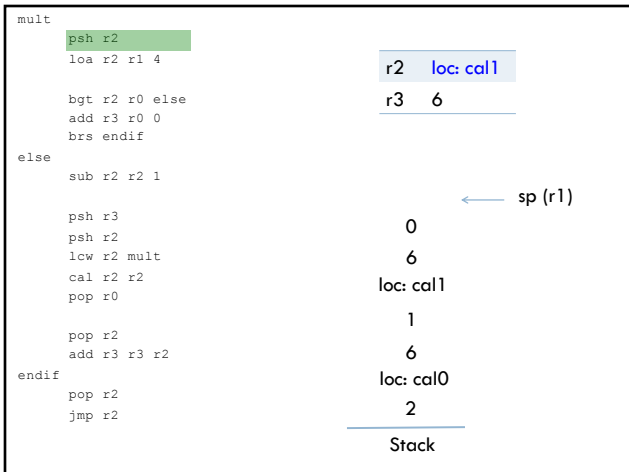
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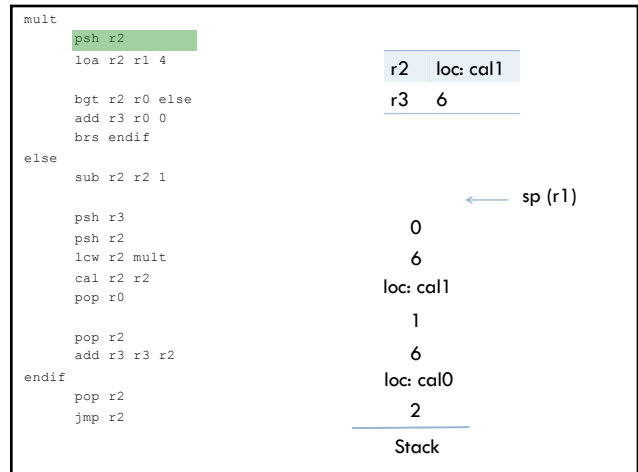
133



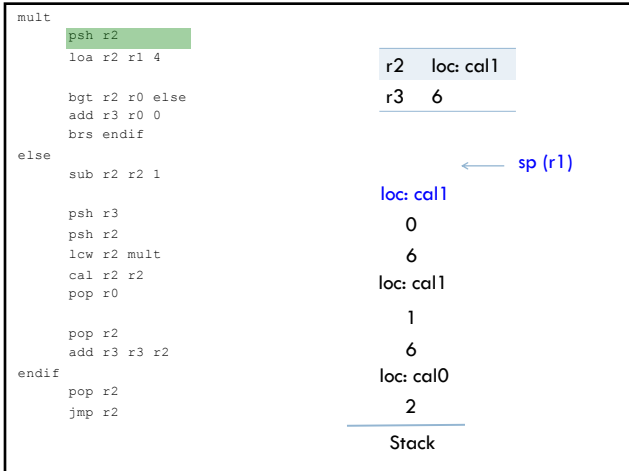
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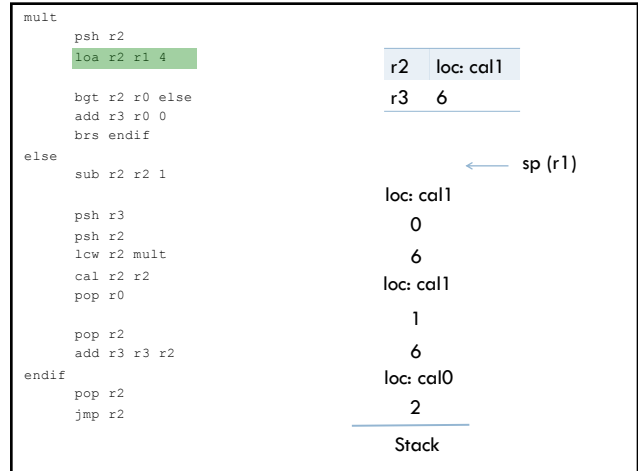
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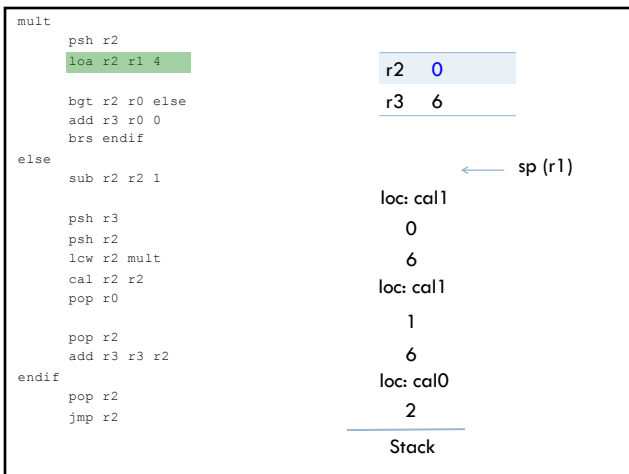
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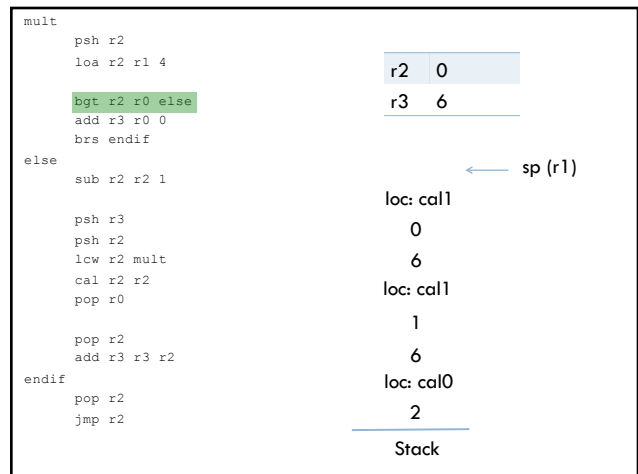
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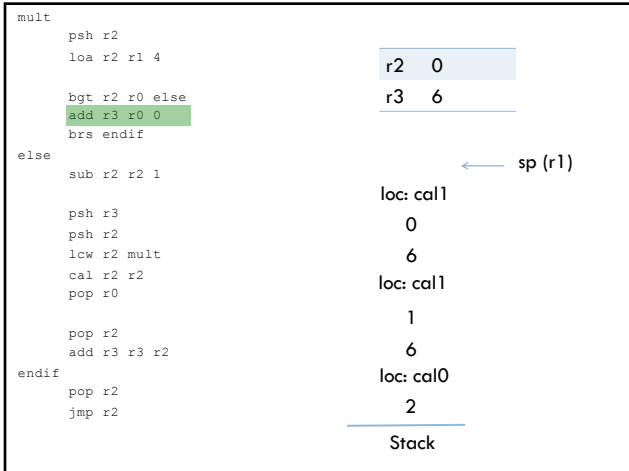
138



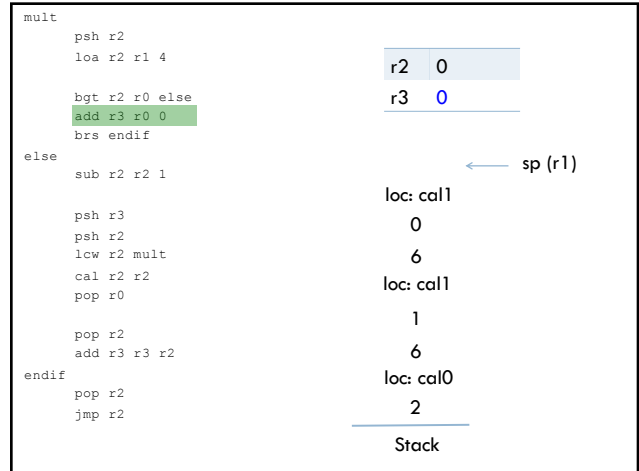
139



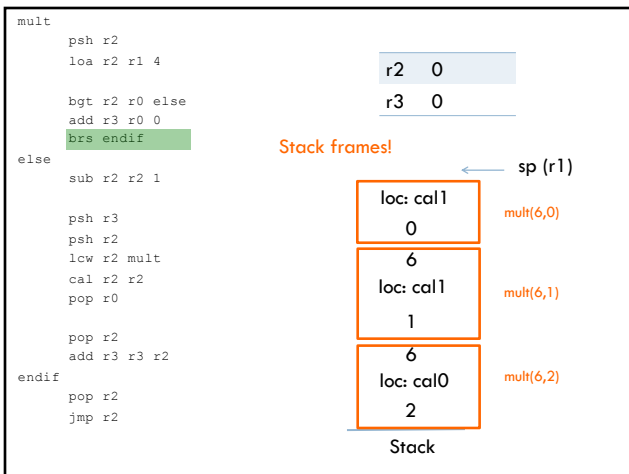
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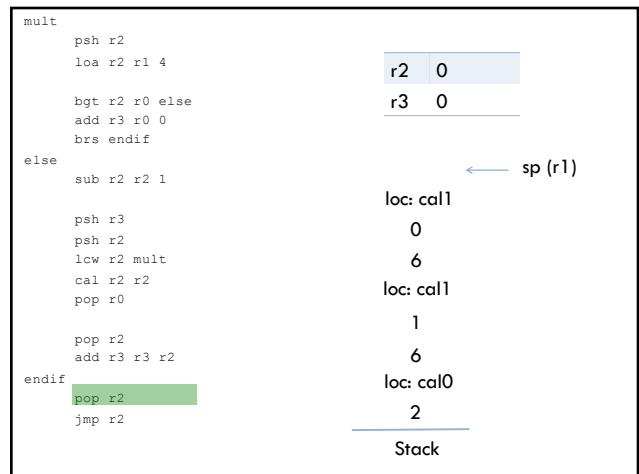
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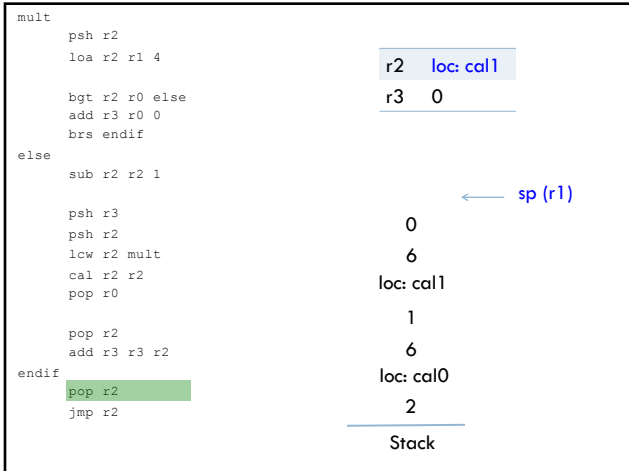
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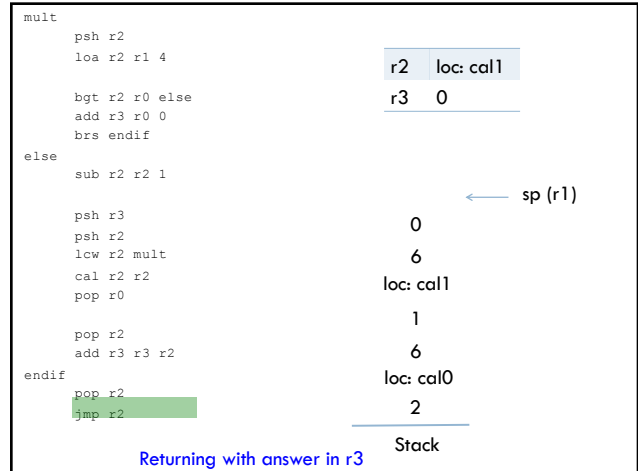
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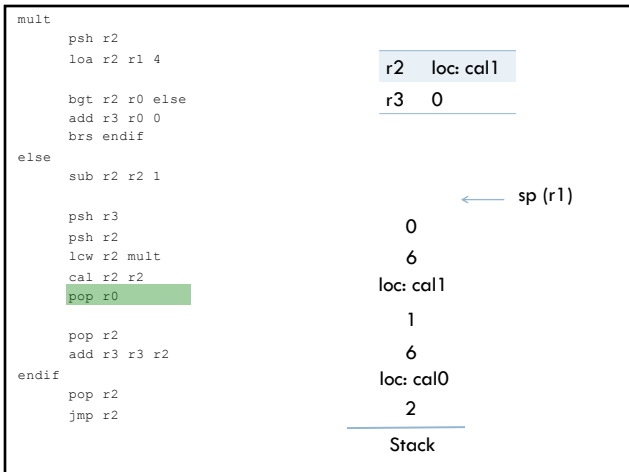
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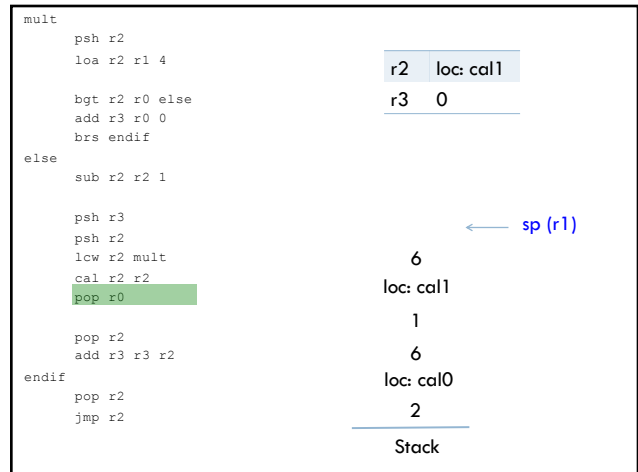
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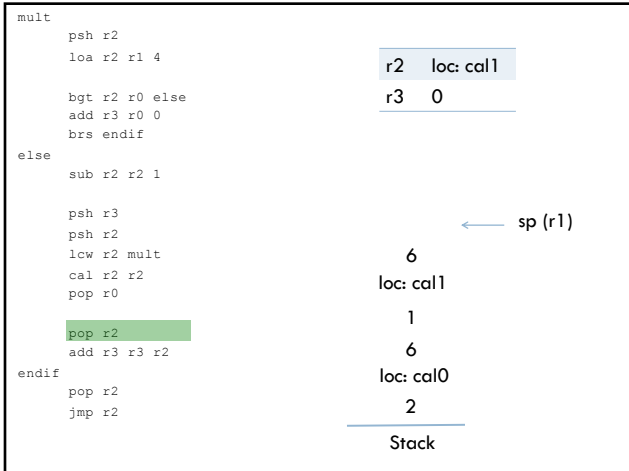
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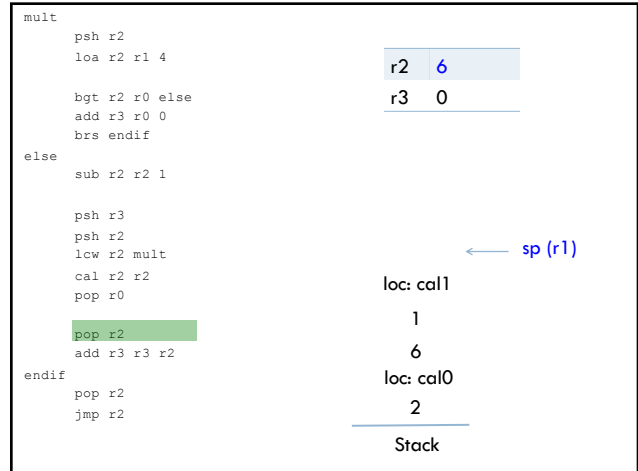
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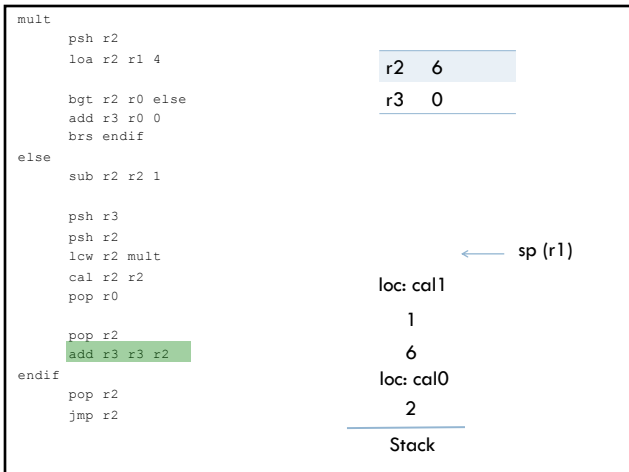
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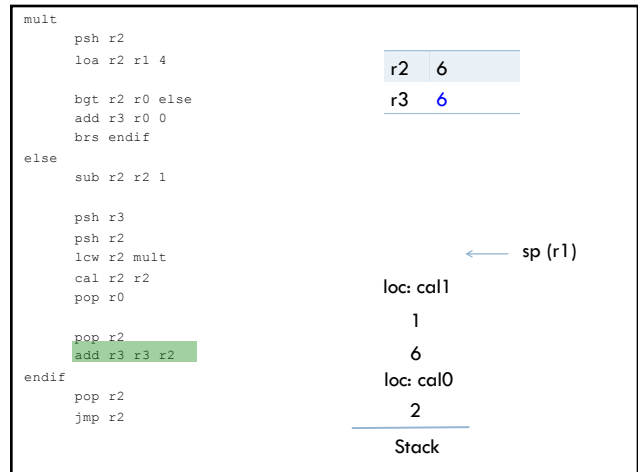
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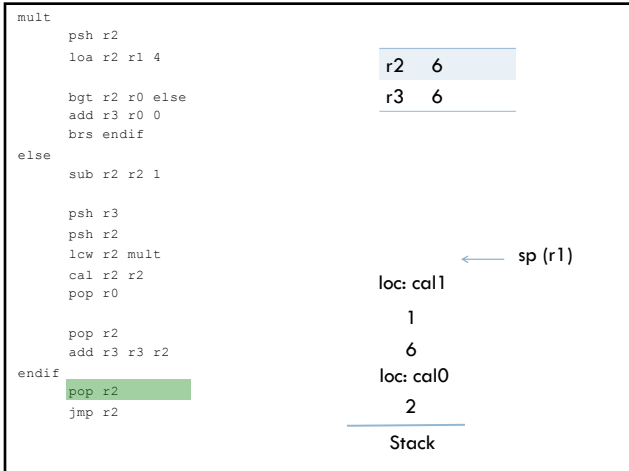
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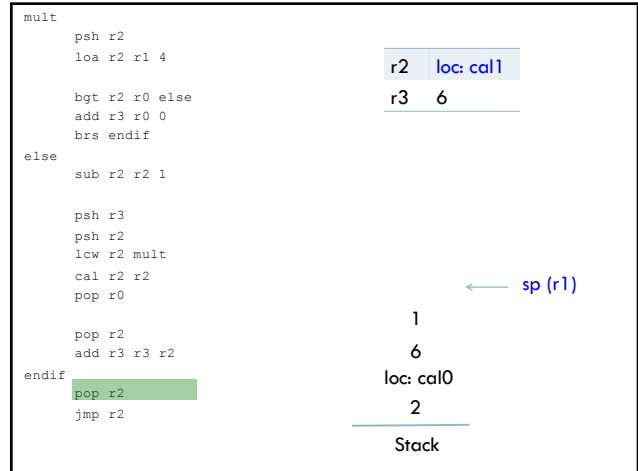
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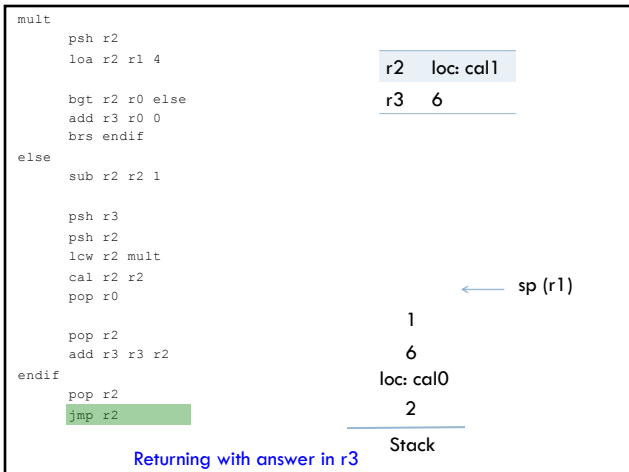
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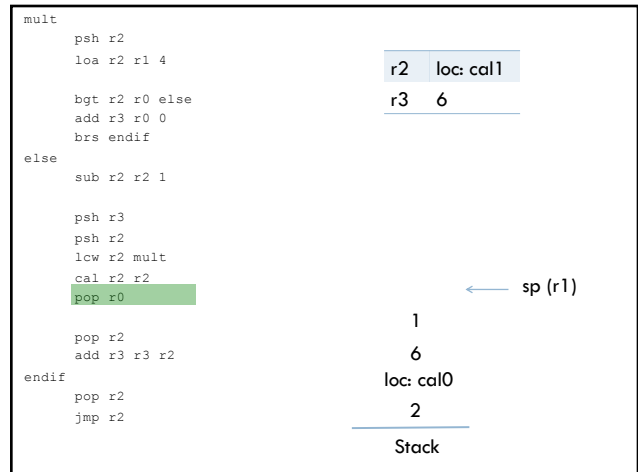
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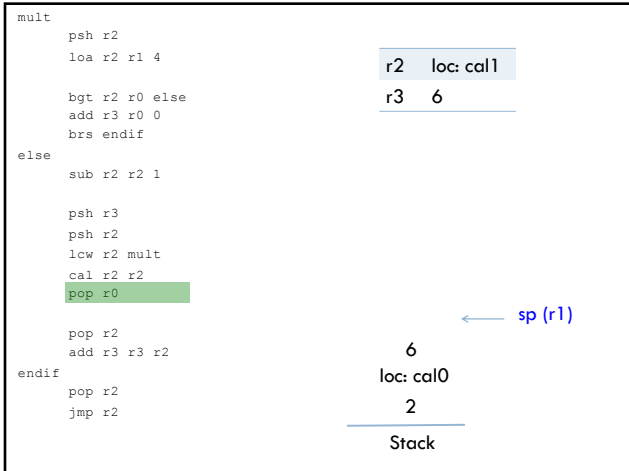
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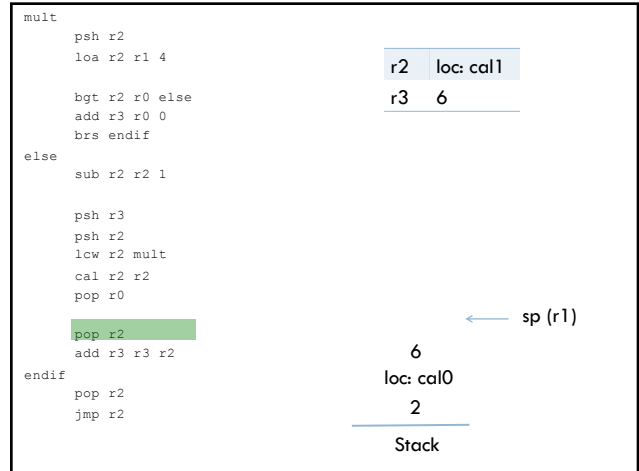
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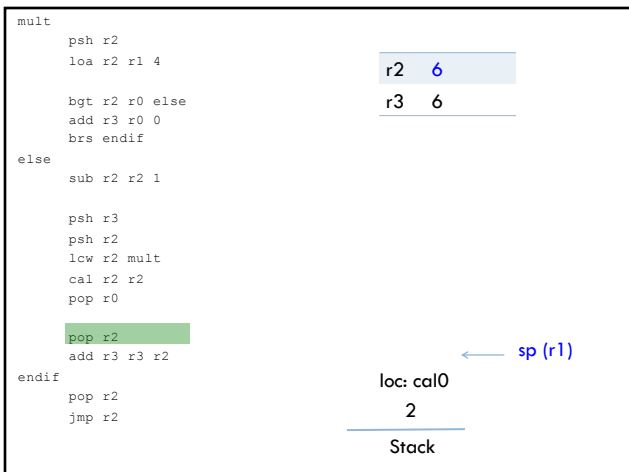
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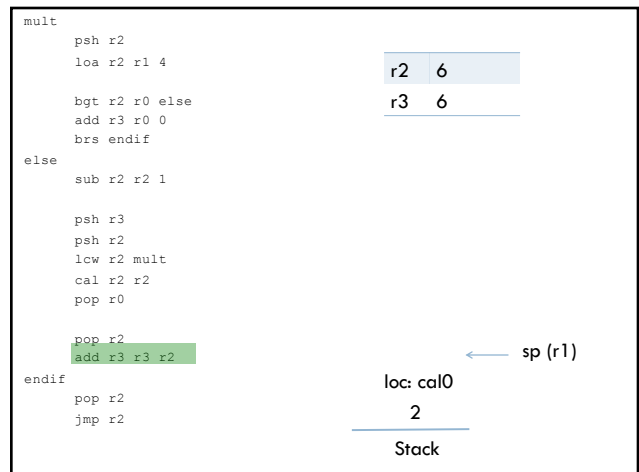
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158

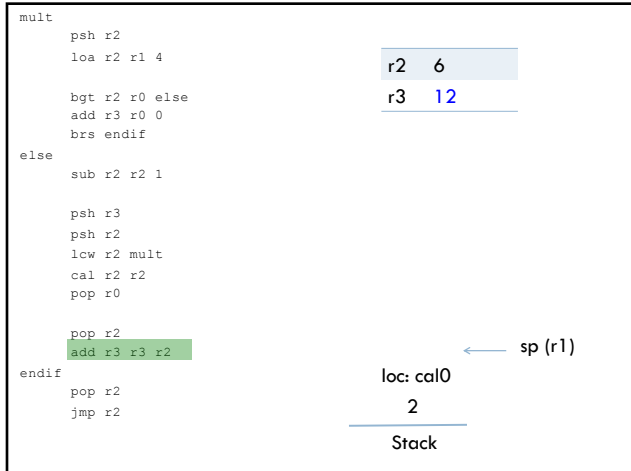


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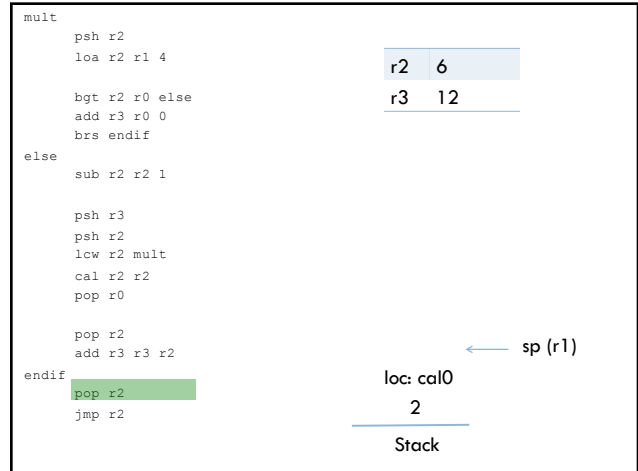


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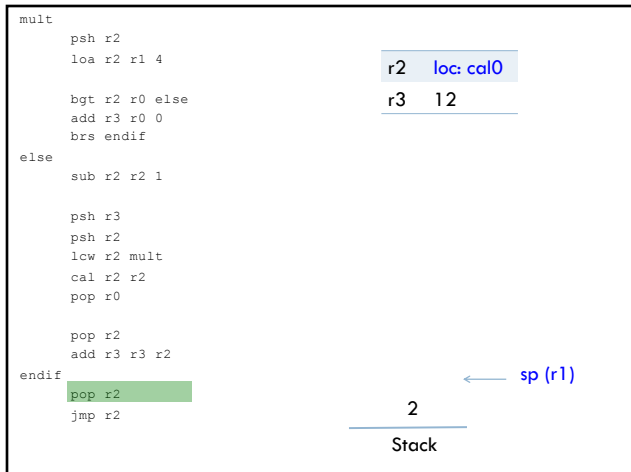




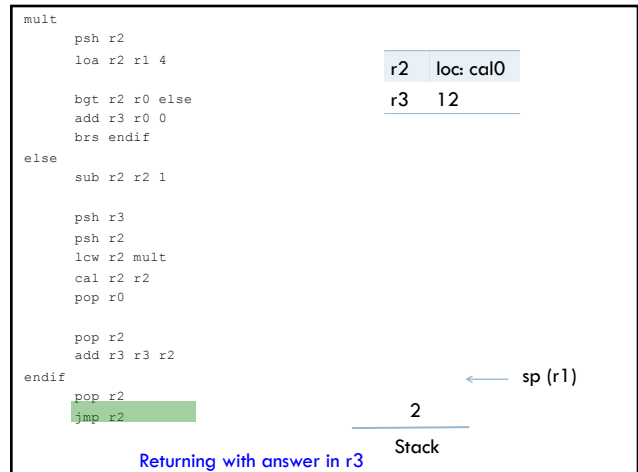
161



162



163



164

Calling mult

```

loa r3 r0
loa r2 r0

psh r2
lcw r2 mult
cal r2 r2

pop r0

sto r3 r0
hlt
    
```

r2	loc: cal0
r3	12

← sp (r1)

2

---

Stack

165

Calling mult

```

loa r3 r0
loa r2 r0

psh r2
lcw r2 mult
cal r2 r2

pop r0

sto r3 r0
hlt
    
```

r2	loc: cal0
r3	12

← sp (r1)

Stack

166

Calling mult

```

loa r3 r0
loa r2 r0

psh r2
lcw r2 mult
cal r2 r2

pop r0

sto r3 r0
hlt
    
```

r2	loc: cal0
r3	12

← sp (r1)

Stack

167

Calling mult

```

loa r3 r0
loa r2 r0

psh r2
lcw r2 mult
cal r2 r2

pop r0

sto r3 r0
hlt
    
```

Print the answer: 12!

r2	loc: cal0
r3	12

← sp (r1)

Stack

168

Calling mult

```

loa r3 r0
loa r2 r0

psh r2
lcw r2 mult
cal r2 r2

pop r0

sto r3 r0
hlt

```

Print the answer: 12!

Stack ← sp (r1)

r2	loc: cal0
r3	12

169

## multiply the easy way

look at mult\_easy.a52 code

we can import libraries (really just functions in other files) using the “inc” command

the included files must be in the same directory as the .a52 file running

170

## Another example

```

max
    psh r2
    loa r2 r1 4

    bge r3 r2 endif
    add r3 r2 0
endif
    pop r2
    jmp r2

```

What does this code do?

186

## Another example

```

max
    psh r2
    loa r2 r1 4

    bge r3 r2 endif
    add r3 r2 0
endif
    pop r2
    jmp r2

```

max, as a function!

187

### Calling max

```

    loa r3 r0
    loa r2 r0

    psh r2
    lcw r2 max
    cal r2 r2
    pop r0

    sto r3 r0
    hlt
    
```

Anything different?

188

### Calling max

```

    loa r3 r0
    loa r2 r0

    psh r2
    lcw r2 max
    cal r2 r2
    pop r0

    sto r3 r0
    hlt
    
```

For the second argument,  
psh it on the stack

189

```

    loa r3 r0
    loa r2 r0

    psh r2
    lcw r2 max
    cal r2 r2
    pop r2

    sto r3 r0
    hlt
max
    psh r2
    loa r2 r1 4

    bge r3 r2 endif
    add r3 r2 0
endif
    pop r2
    jmp r2
    
```

r2
r3

← sp (r1)

Stack

190

```

    loa r3 r0
    loa r2 r0

    psh r2
    lcw r2 max
    cal r2 r2
    pop r2

    sto r3 r0
    hlt
max
    psh r2
    loa r2 r1 4

    bge r3 r2 endif
    add r3 r2 0
endif
    pop r2
    jmp r2
    
```

r2
r3 10

← sp (r1)

Stack

191

```

    loa r3 r0
    loa r2 r0
    psh r2
    lcw r2 max
    cal r2 r2
    pop r2

    sto r3 r0
    hlt
max
    psh r2
    loa r2 r1 4

    bge r3 r2 endif
    add r3 r2 0
endif
    pop r2
    jmp r2

```

r2	2
r3	10

← sp (r1)

Stack

192

```

    loa r3 r0
    loa r2 r0
    psh r2
    lcw r2 max
    cal r2 r2
    pop r2

    sto r3 r0
    hlt
max
    psh r2
    loa r2 r1 4

    bge r3 r2 endif
    add r3 r2 0
endif
    pop r2
    jmp r2

```

r2	2
r3	10

← sp (r1)

2

Stack

193

```

    loa r3 r0
    loa r2 r0
    psh r2
    lcw r2 max
    cal r2 r2
    pop r2

    sto r3 r0
    hlt
max
    psh r2
    loa r2 r1 4

    bge r3 r2 endif
    add r3 r2 0
endif
    pop r2
    jmp r2

```

r2	2
r3	10

← sp (r1)

2

Stack

194

```

    loa r3 r0
    loa r2 r0
    psh r2
    lcw r2 max
    cal r2 r2
    pop r2

    sto r3 r0
    hlt
max
    psh r2
    loa r2 r1 4

    bge r3 r2 endif
    add r3 r2 0
endif
    pop r2
    jmp r2

```

r2	max
r3	10

← sp (r1)

2

Stack

Notice that we overwrote the value in r2

If we hadn't saved it on the stack, it would have been lost

195

```

    loa r3 r0
    loa r2 r0

    psh r2
    lcw r2 max
    cal r2 r2
    pop r2

    sto r3 r0
    hlt
max
    psh r2
    loa r2 r1 4

    bge r3 r2 endif
    add r3 r2 0
endif
    pop r2
    jmp r2

```

r2	max
r3	10

← sp (r1)

2

Stack

196

```

    loa r3 r0
    loa r2 r0

    psh r2
    lcw r2 max
    cal r2 r2
    pop r2

    sto r3 r0
    hlt
max
    psh r2
    loa r2 r1 4

    bge r3 r2 endif
    add r3 r2 0
endif
    pop r2
    jmp r2

```

r2	loc: cal
r3	10

← sp (r1)

2

Stack

197

```

    loa r3 r0
    loa r2 r0

    psh r2
    lcw r2 max
    cal r2 r2
    pop r2

    sto r3 r0
    hlt
max
    psh r2
    loa r2 r1 4

    bge r3 r2 endif
    add r3 r2 0
endif
    pop r2
    jmp r2

```

r2	loc: cal
r3	10

← sp (r1)

2

Stack

198

```

    loa r3 r0
    loa r2 r0

    psh r2
    lcw r2 max
    cal r2 r2
    pop r2

    sto r3 r0
    hlt
max
    psh r2
    loa r2 r1 4

    bge r3 r2 endif
    add r3 r2 0
endif
    pop r2
    jmp r2

```

r2	loc: cal
r3	10

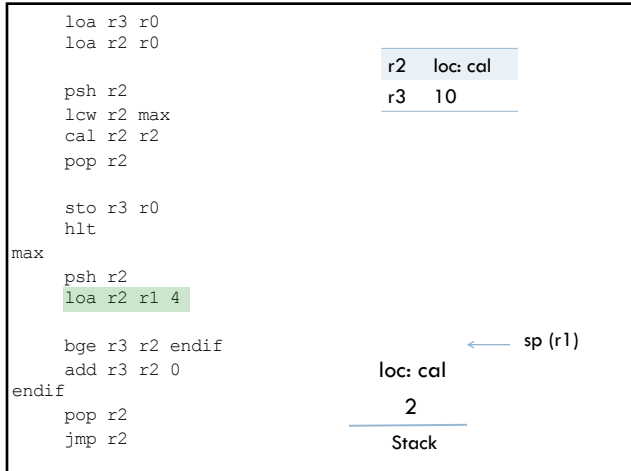
← sp (r1)

loc: cal

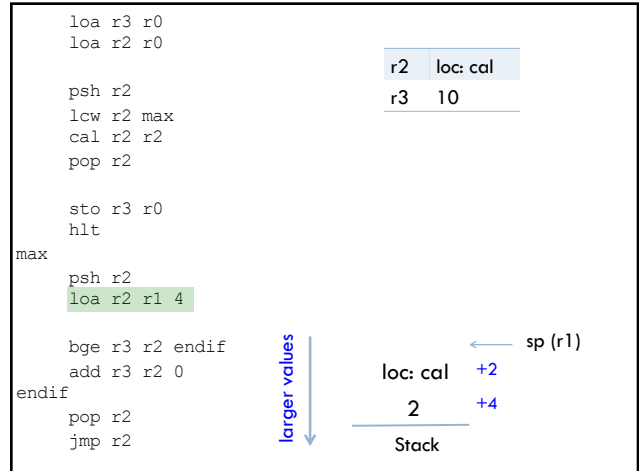
2

Stack

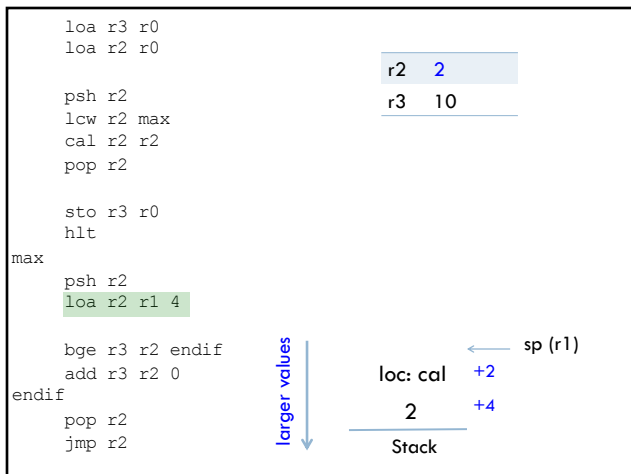
199



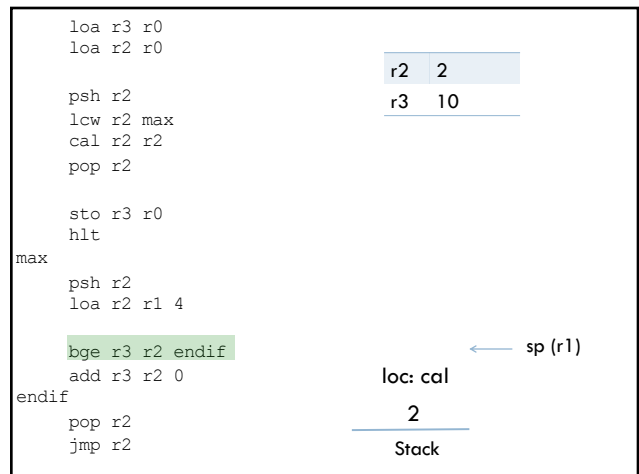
200



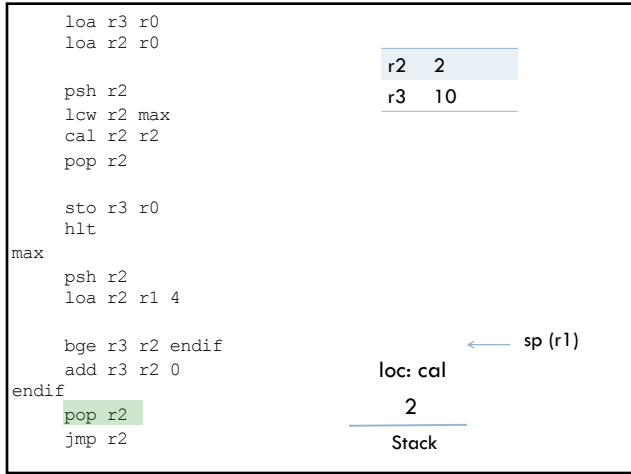
201



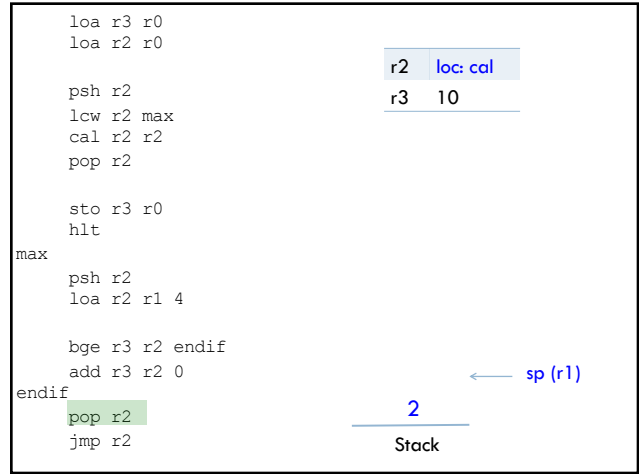
202



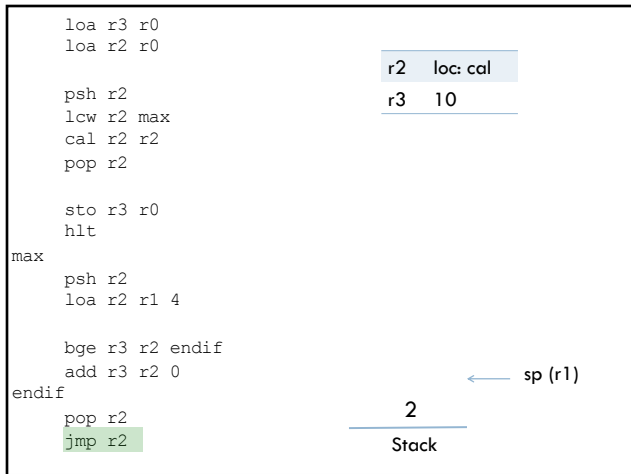
203



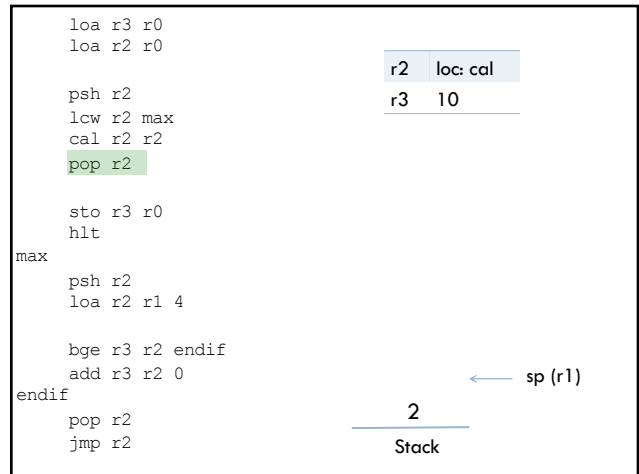
204



205



206



207



```

    loa r3 r0
    loa r2 r0

    psh r2
    lcw r2 max
    cal r2 r2
    pop r2

    sto r3 r0
    hlt
max
    psh r2
    loa r2 r1 4

    bge r3 r2 endif
    add r3 r2 0
endif
    pop r2
    jmp r2

```

r2 2  
r3 10

← sp (r1)

Stack

208

```

    loa r3 r0
    loa r2 r0

    psh r2
    lcw r2 max
    cal r2 r2
    pop r2

    sto r3 r0
    hlt
max
    psh r2
    loa r2 r1 4

    bge r3 r2 endif
    add r3 r2 0
endif
    pop r2
    jmp r2

```

r2 2  
r3 10

← sp (r1)

Stack

209

```

    loa r3 r0
    loa r2 r0

    psh r2
    lcw r2 max
    cal r2 r2
    pop r2

    sto r3 r0
    hlt
max
    psh r2
    loa r2 r1 4

    bge r3 r2 endif
    add r3 r2 0
endif
    pop r2
    jmp r2

```

r2 2  
r3 10

10!

← sp (r1)

Stack

210

```

    loa r3 r0
    loa r2 r0

    psh r2
    lcw r2 max
    cal r2 r2
    pop r2

    sto r3 r0
    hlt
max
    psh r2
    loa r2 r1 4

    bge r3 r2 endif
    add r3 r2 0
endif
    pop r2
    jmp r2

```

r2 2  
r3 10

← sp (r1)

Stack

211