

Breadth First Search

<https://cs.pomona.edu/classes/cs140/>

Outline

Topics and Learning Objectives

- Discuss breadth first search for graphs

Exercises

- Continued from previous lecture slides
- Compute distance with Breadth-first search

Extra Resources

- Introduction to Algorithms, 3rd, Chapter 22
- Algorithms Illuminated Part 2: Chapter 8

General Algorithm

```
FUNCTION Connectivity(G, start_vertex)
    found = {v: FALSE FOR v IN G.vertices}
    found[start_vertex] = TRUE
```

LOOP

```
    (vFound, vNotFound) = get_valid_edge(G.edges, found)
```

```
    IF vFound == NONE || vNotFound == NONE
```

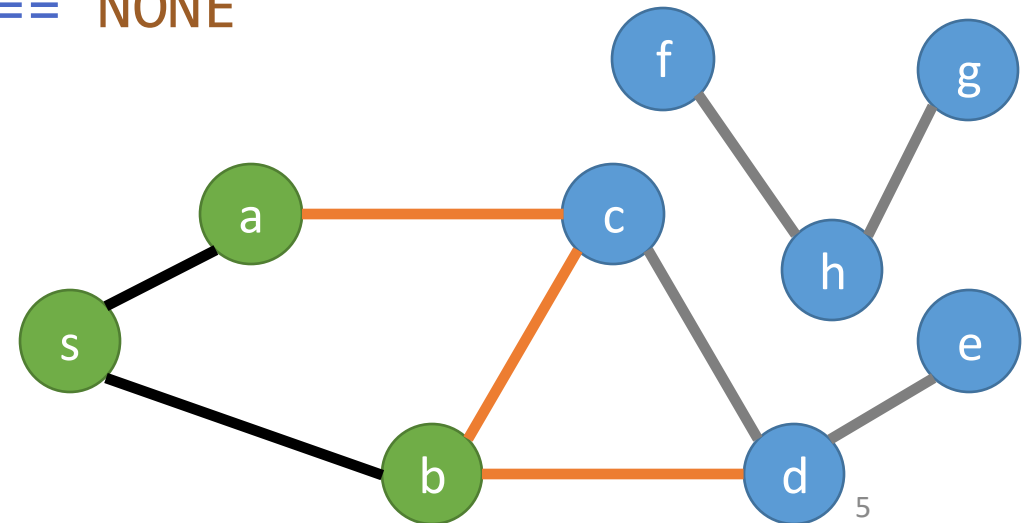
```
        BREAK
```

```
    ELSE
```

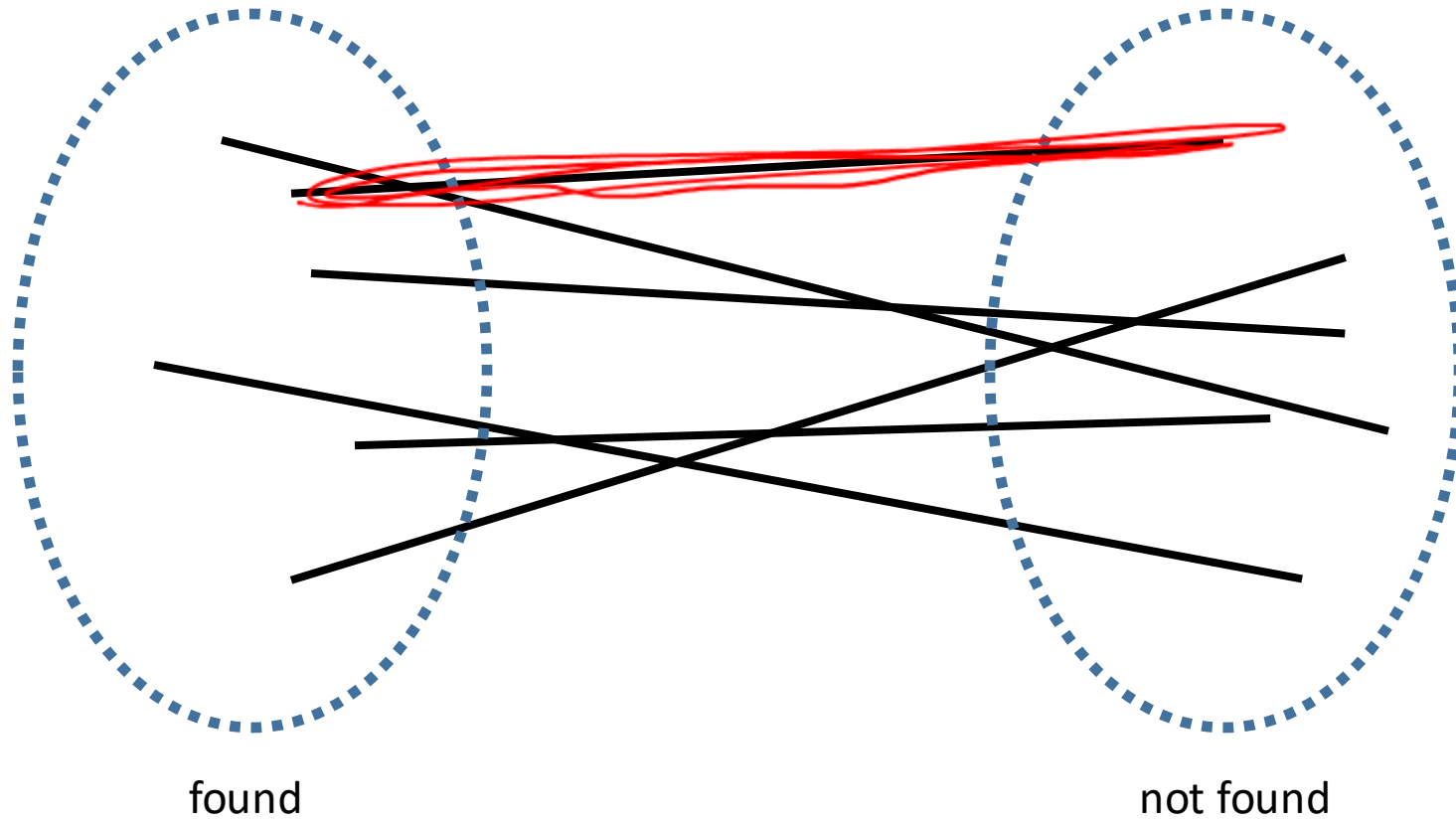
```
        found[vNotFound] = TRUE
```

```
RETURN found
```

Find an edge where one vertex has been found and the other vertex has not been found.



How do we choose the next edge?



Two common (and well studied) options

Breadth-First Search

- Explore the graph in **layers**
- “*Cautious*” exploration
- Use a FIFO data structure (can you think of an example?)

Depth-First Search

- Explore recursively
- A more “*aggressive*” exploration (we backtrack if necessary)
- Use a LIFO data structure (or recursion)

```
FUNCTION BFS(G, start_vertex)
```

```
    found = {v: FALSE FOR v IN G.vertices}
```

```
    found[start_vertex] = TRUE
```

```
    visit_queue = [start_vertex]
```

```
WHILE visit_queue.length != 0
```

```
    vFound = visit_queue.pop() ↓
```

```
    FOR vOther IN G.edges[vFound]
```

```
        IF found[vOther] == FALSE ←
```

```
            found[vOther] = TRUE ←
```

```
            visit_queue.add(vOther) ←
```

```
RETURN found
```

```
FUNCTION Connectivity(G, start_vertex)
```

```
    found = {v: FALSE FOR v IN G.vertices}
```

```
    found[start_vertex] = TRUE
```

```
    LOOP
```

```
        (vFound, vNotFound) =
```

```
            get_valid_edge(G.edges, found)
```

```
    IF vFound == NONE || vNotFound == NONE
```

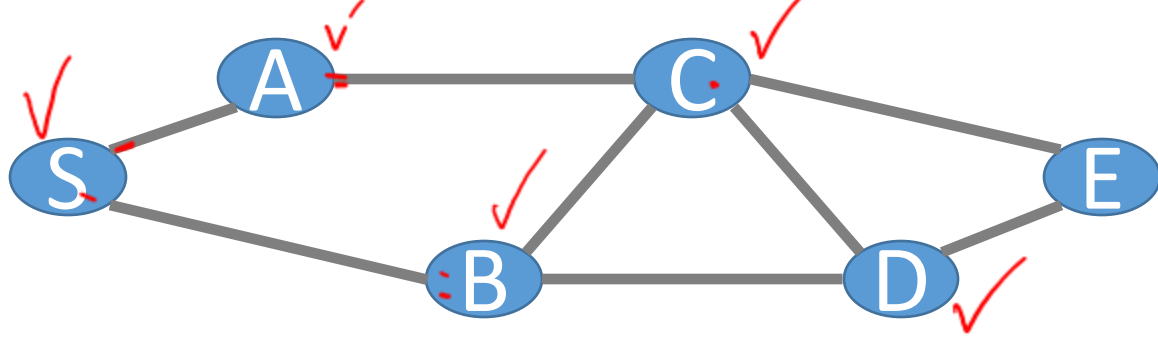
```
        BREAK
```

```
    ELSE
```

```
        found[vNotFound] = TRUE
```

```
    RETURN found
```

Exercise questions 2 and 3



$v_{\text{Found}} = S$
 $vO = A$
 $vO = B$
 $VF = A$
 $vO = C$
 $vO = S$
 $VF = B$
 $vO = C$
 $vO = D$
 $vO = S$

$VF = C$
 $vO = A$
 $vO = B$
 $vO = D$
 $vO = E$
 $VF = D$

S A B C D E

```

FUNCTION BFS(G, start_vertex)
    found = {v: FALSE FOR v IN G.vertices}
    found[start_vertex] = TRUE
    visit_queue = [start_vertex]
    = [S, A, B, C, D, E]
    WHILE visit_queue.length != 0
        → vFound = visit_queue.pop()
        FOR vOther IN G.edges[vFound]
            { IF found[vOther] == FALSE
                found[vOther] = TRUE
                visit_queue.add(vOther)
            }
    RETURN found
  
```

Given a tie, visit edges are in alphabetical order

Running Time

What is the running time?

```
FUNCTION BFS(G, start_vertex)
  found = {v: FALSE FOR v IN G.vertices}
  found[start_vertex] = TRUE
  visit_queue = [start_vertex]
```



How many times do we consider each edge?

```
  WHILE visit_queue.length != 0
    vFound = visit_queue.pop()
    FOR vOther IN G.edges[vFound]
      IF found[vOther] == FALSE
        found[vOther] = TRUE
        visit_queue.add(vOther)
```

```
  RETURN found
```

~~$O(m) \cdot O(n)$~~

$O(n) + O(m) = O(n+m)$

twice

Running Time

```
FUNCTION BFS(G, start_vertex)
    found = {v: FALSE FOR v IN G.vertices}
    found[start_vertex] = TRUE
    visit_queue = [start_vertex]

    WHILE visit_queue.length != 0
        vFound = visit_queue.pop()
        FOR vOther IN G.edges[vFound]
            IF found[vOther] == FALSE
                found[vOther] = TRUE
                visit_queue.add(vOther)

    RETURN found
```

What is the running time?

How many times do we consider each edge?

$$T_{BFS}(n, m) = O(\underbrace{n_s + m_s})$$

where n_s and m_s are the nodes and edges **findable/connected** from/to the start vertex

Proof: BFS

Claim: BFS finds all nodes connected to the start node.

At the end of the BFS algorithm, v is marked found if there exists a path from s to v

- Note: this is just a special case of the general algorithm that we proved by contradiction

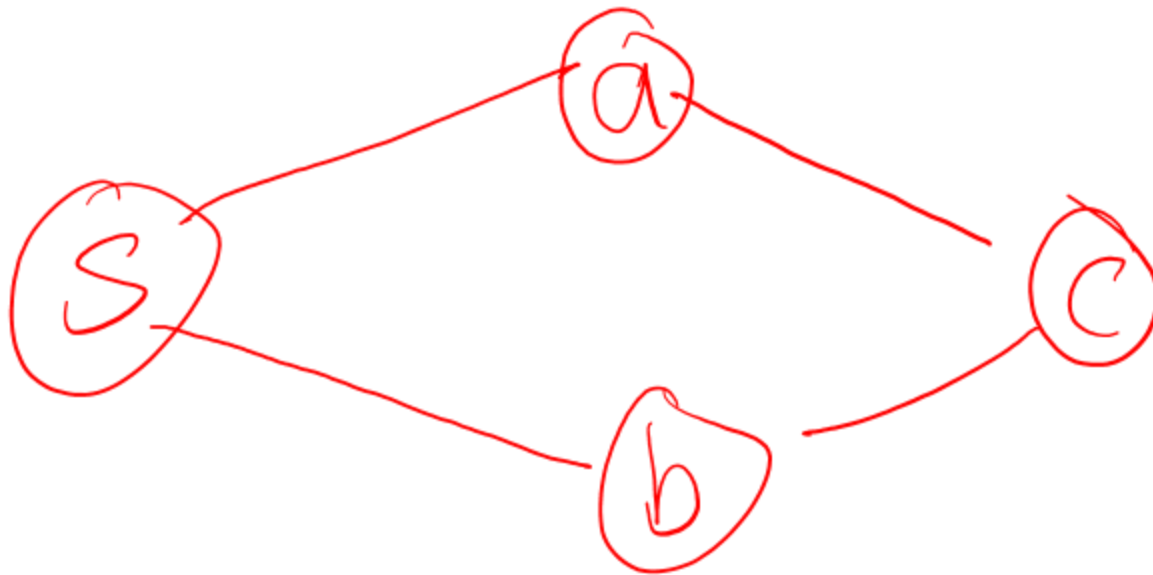
Practice for a loop invariant

Homework question

Question

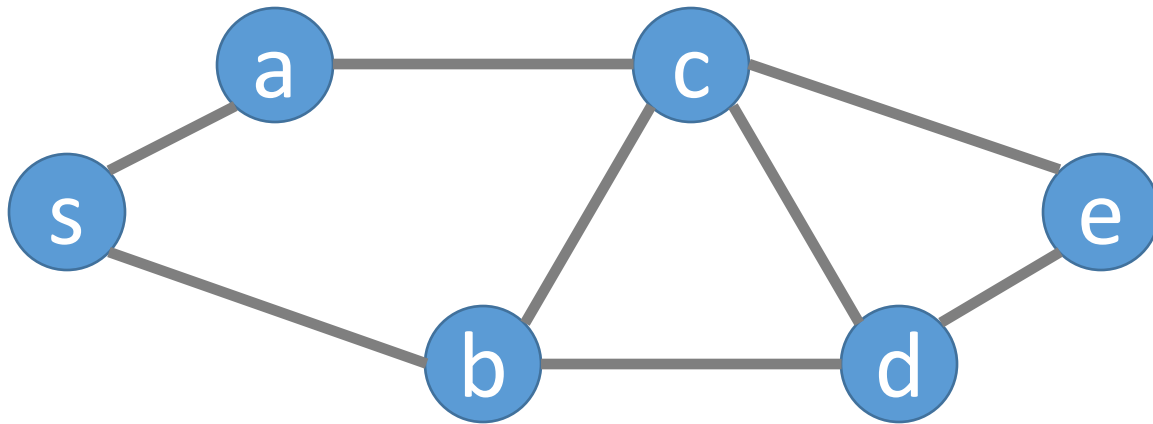
The Shortest Path Problem

- How can we determine the fewest number of hops between the start vertex and all other connected vertices?



BFS Exercise Question 1

How can we determine the fewest number of hops between the start vertex and all other connected vertices?



```
FUNCTION BFS(G, start_vertex)
    found = {v: FALSE FOR v IN G.vertices}
    found[start_vertex] = TRUE
    visit_queue = [start_vertex]

    WHILE visit_queue.length != 0
        vFound = visit_queue.pop()
        FOR vOther IN G.edges[vFound]
            IF found[vOther] == FALSE
                found[vOther] = TRUE
                visit_queue.add(vOther)

    RETURN found
```

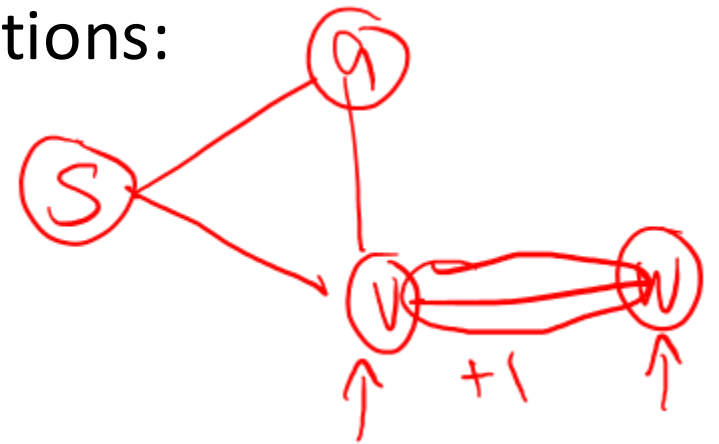
Given a tie, visit edges are in alphabetical order

The Shortest Path Problem

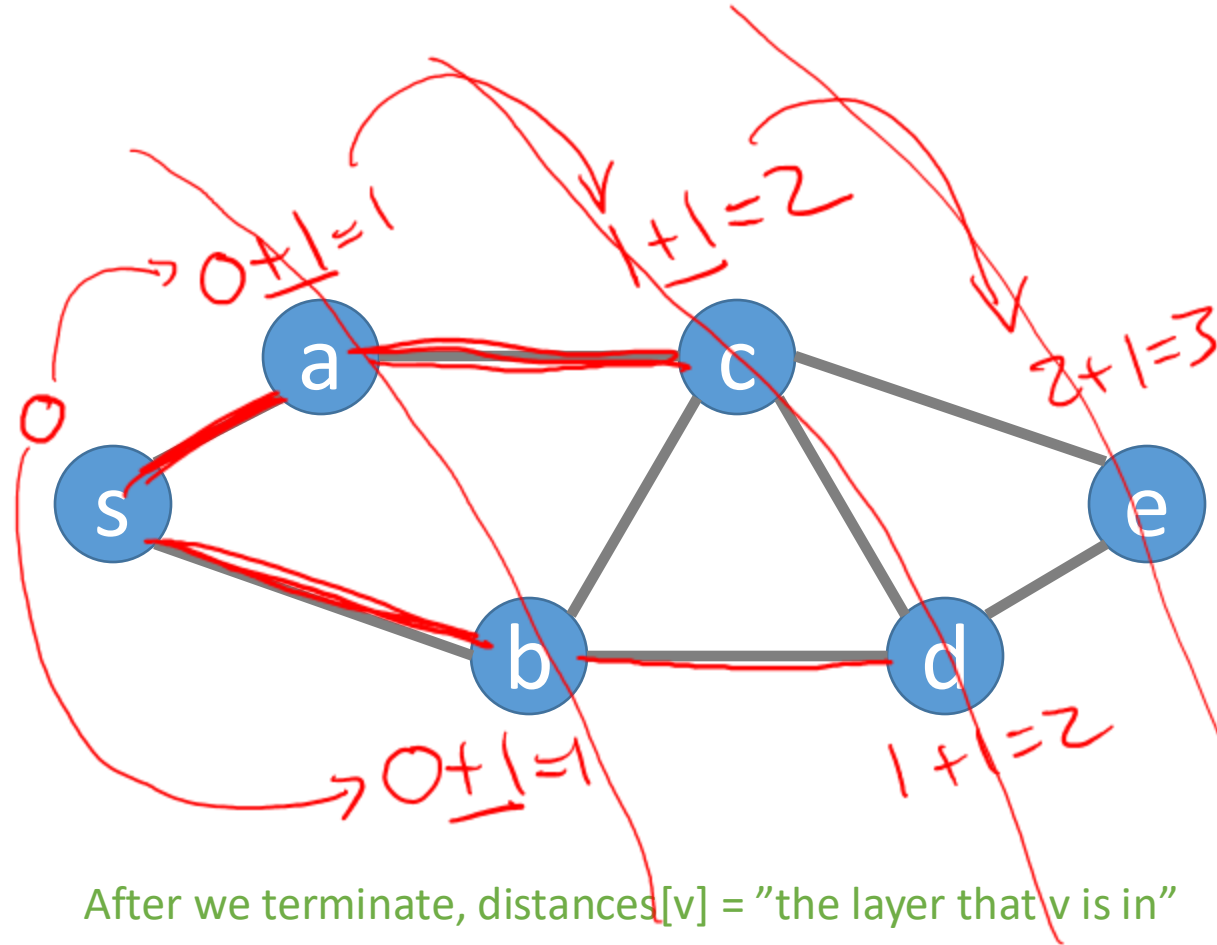
Determine the fewest number of hops between the start vertex and all other vertices

Same algorithm as before with the following additions:

- Initialize the distances[s] as 0
- Initialize all other distances to **infinity**
- When considering an edge (v , w)
 - If w is not found, then set $\text{dist}(w)$ to $\text{dist}(v) + 1$



The Shortest Path Problem



```
FUNCTION DistanceBFS(G, start_vertex)
    found = {v: FALSE FOR v IN G.vertices}
    found[start_vertex] = TRUE
```

```
    distances = {v: INFINITY FOR v IN G.vertices}
    distances[start_vertex] = 0
```

```
    visit_queue = [start_vertex]
```

```
    WHILE visit_queue.length != 0
```

```
        vFound = visit_queue.pop()
```

```
        FOR vOther IN G.edges[vFound]
```

```
            IF found[vOther] == FALSE
```

```
                found[vOther] = TRUE
```

```
                visit_queue.add(vOther)
```

```
                distances[vOther] = distances[vFound] + 1
```

```
    RETURN distances
```

Given a tie, visit edges are in alphabetical order

Connected Components

Let's only consider undirected graphs for now

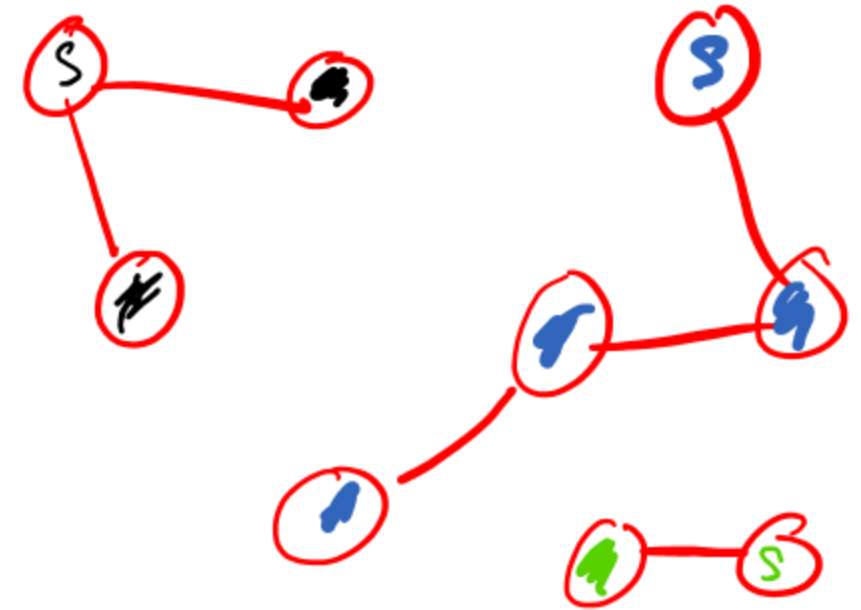
Let $G = (V, E)$ be an undirected graph

Goal: compute all **connected components** in $O(m + n)$

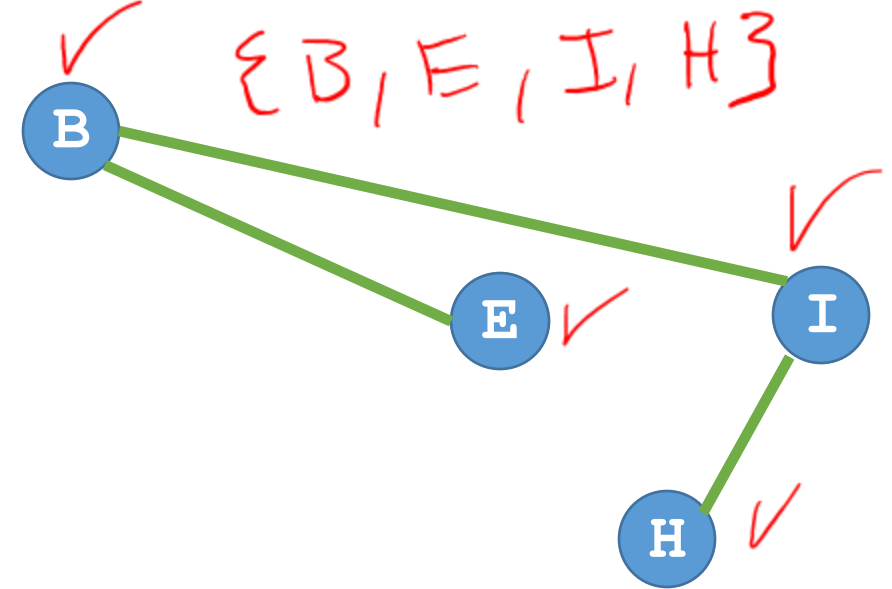
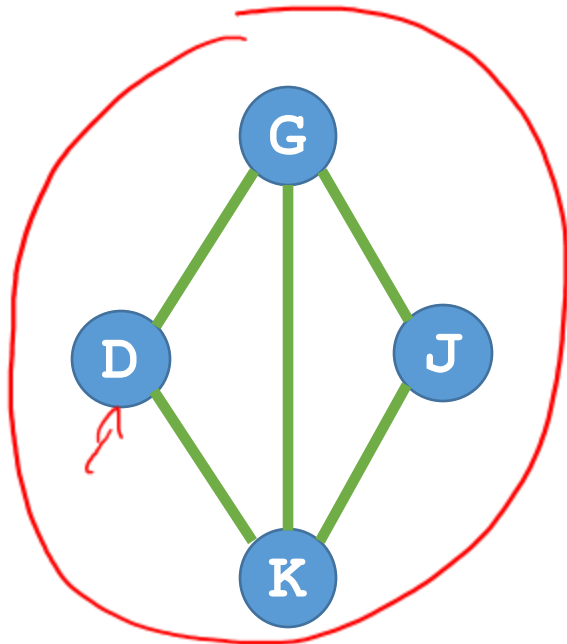
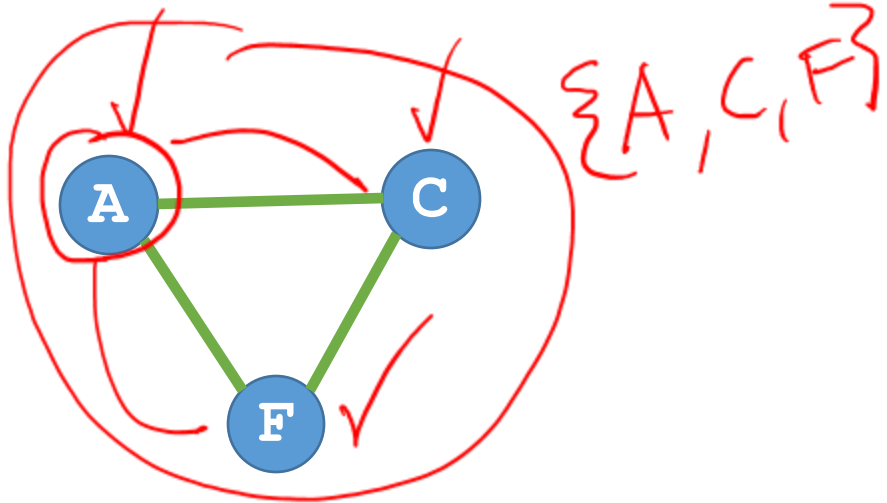
- A component is any group of vertices that can reach one another
- For example, if we are trying to see if a network has become disconnected

Exercise question 2:

How would you do this using our BFS procedure from before?



BFS Exercise Question 2



```

FUNCTION FindComponents (G)
    components = []
    found = {v: FALSE FOR v IN G.vertices}
    FOR v IN G.vertices
        IF NOT found[v]
            newly_found = BFS(G, v)
            new_component = {
                w FOR w, w_is_found IN newly_found
                IF w_is_found
            }
            component.append(new_component)
            FOR w IN new_component:
                found[w] = TRUE
    RETURN components
    
```