

# Kosaraju's Algorithm for Strongly Connected Components

<https://cs.pomona.edu/classes/cs140/>

# Outline

## Topics and Learning Objectives

- Review topological orderings
- Discuss strongly connected components
- Cover Kosaraju's Algorithm

## Exercise

- Work through Kosaraju's Algorithm

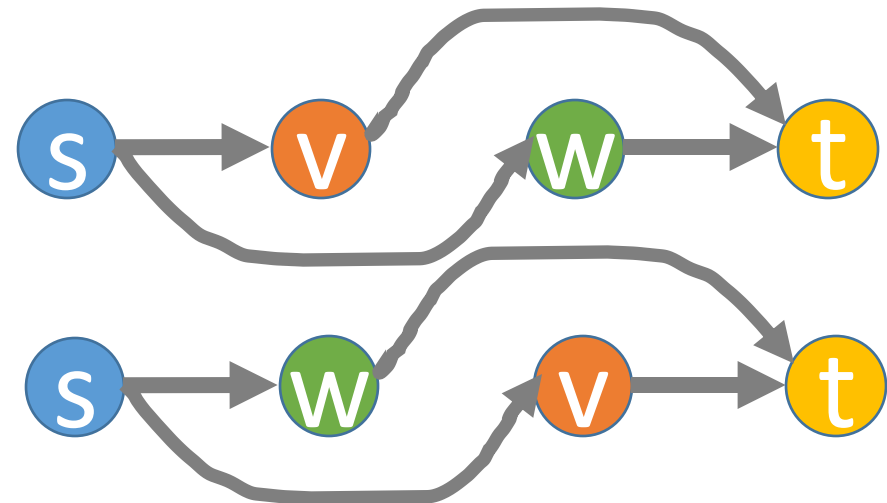
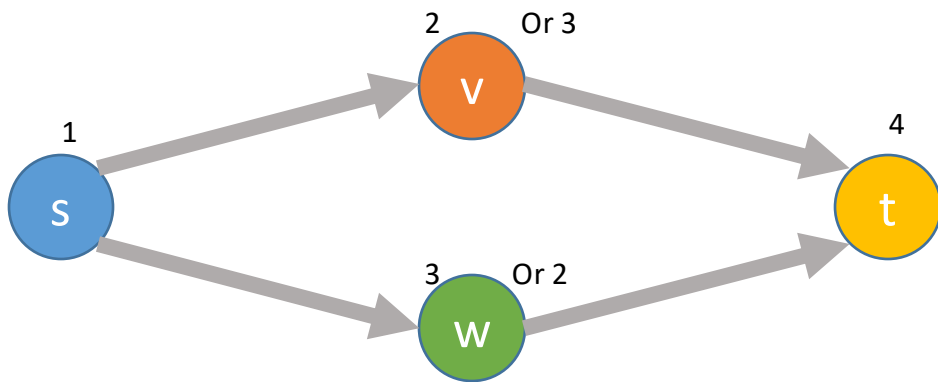
# Extra Resources

- Introduction to Algorithms, 3rd, chapter 22
- Algorithms Illuminated Part 2: Chapter 8

# Topological Orderings

Definition: a topological ordering of a **directed acyclic** graph is a labelling **f** of the graph's vertices such that:

1. The f-values are of the set  $\{1, 2, \dots, n\}$
2. For an edge  $(u, v)$  of  $G$ ,  $f(u) < f(v)$



# Solve with DFS

```
FUNCTION TopologicalOrdering(G)
```

```
    found = {v: FALSE FOR v IN G.vertices}
```

```
    fValues = {v: INFINITY FOR v IN G.vertices}
```

```
    f = G.vertices.length
```

```
    FOR v IN G.vertices
```

```
        IF found[v] == FALSE
```

```
            DFSTopological(G, v, found, f, fValues)
```

```
    RETURN fValues
```

```
FUNCTION DFSTopological(G, v, found, f, fValues)
```

```
    found[v] = TRUE
```

```
    FOR vOther IN G.edges[v]
```

```
        IF found[vOther] == FALSE
```

```
            DFSTopological(G, vOther, found, f, fValues)
```

```
    fValues[v] = f
```

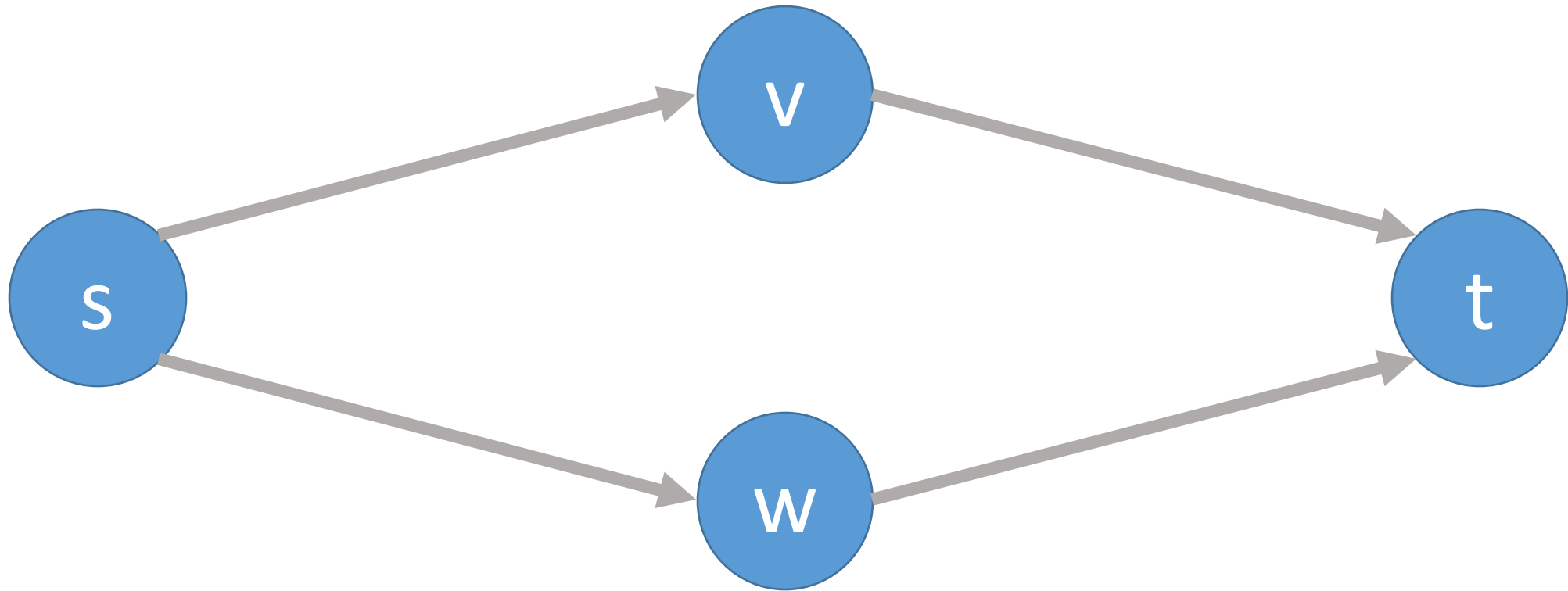
```
    f = f - 1
```

# Strongly Connected Components

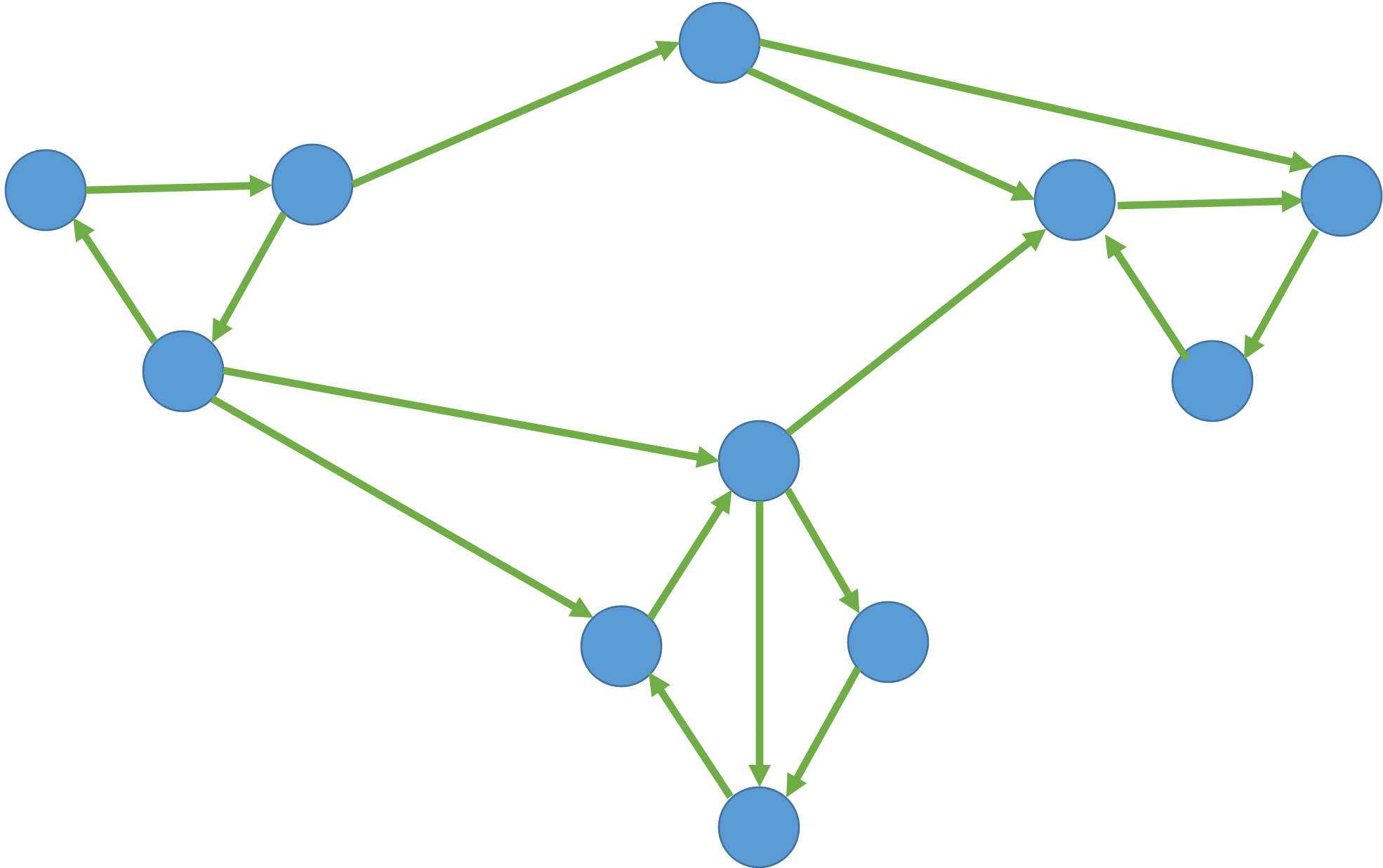
- Topological orderings are useful in their own right, but they also let us **efficiently** calculate the **strongly connected components (SCCs)** of a graph
- A **component (set of vertices)** of a graph is strongly connected if we can find a path from any vertex to any other vertex
- This is a concept for **directed** graphs only
- (just ***connected components*** for undirected graphs)

Why are SCCs useful?

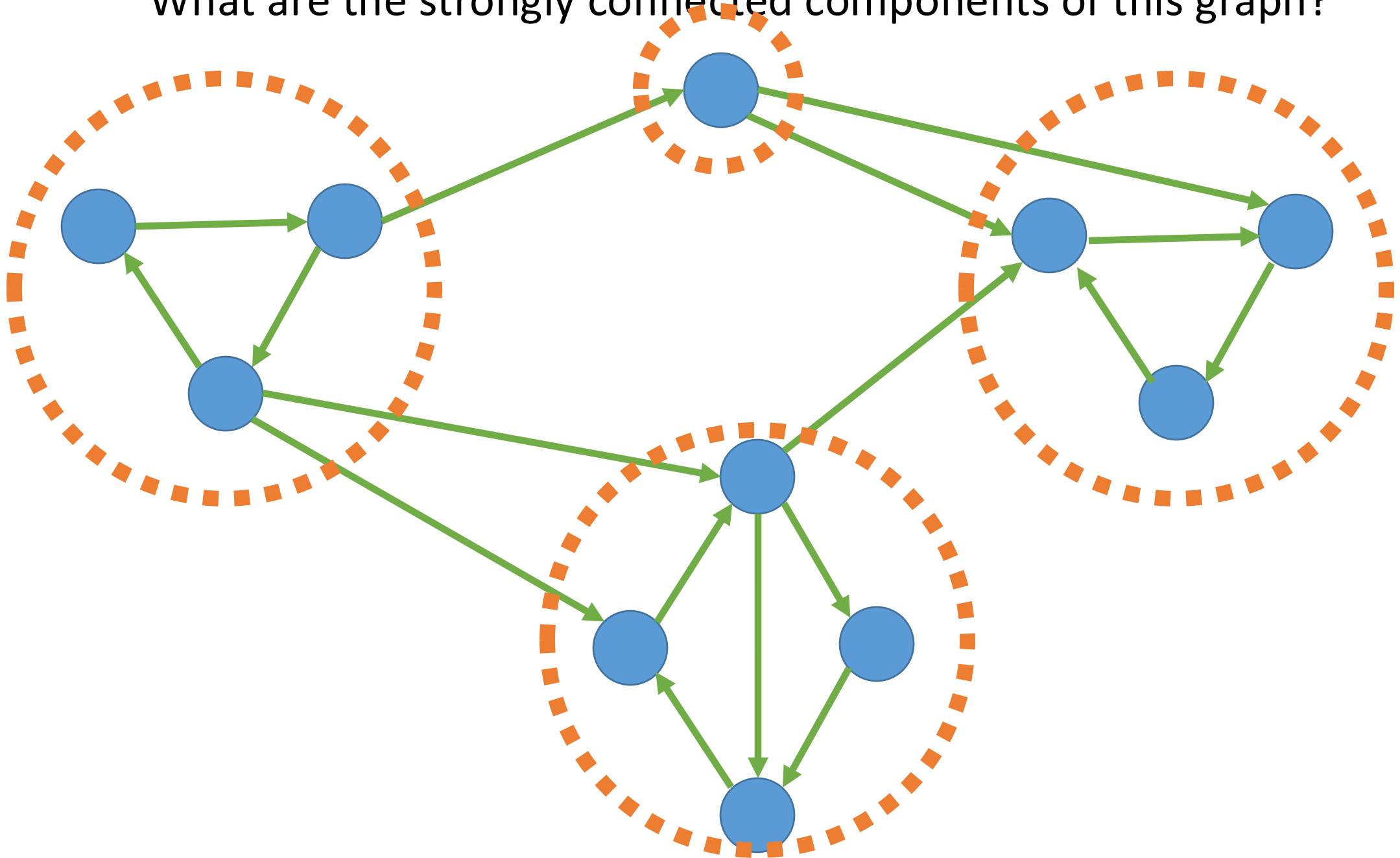
What are the strongly connected components of this graph?



What are the strongly connected components of this graph?



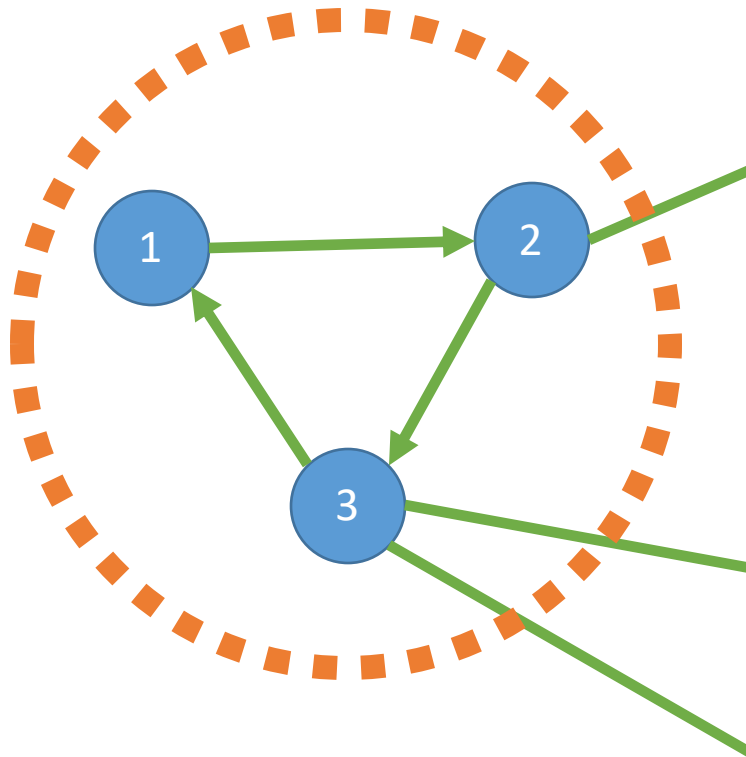
What are the strongly connected components of this graph?



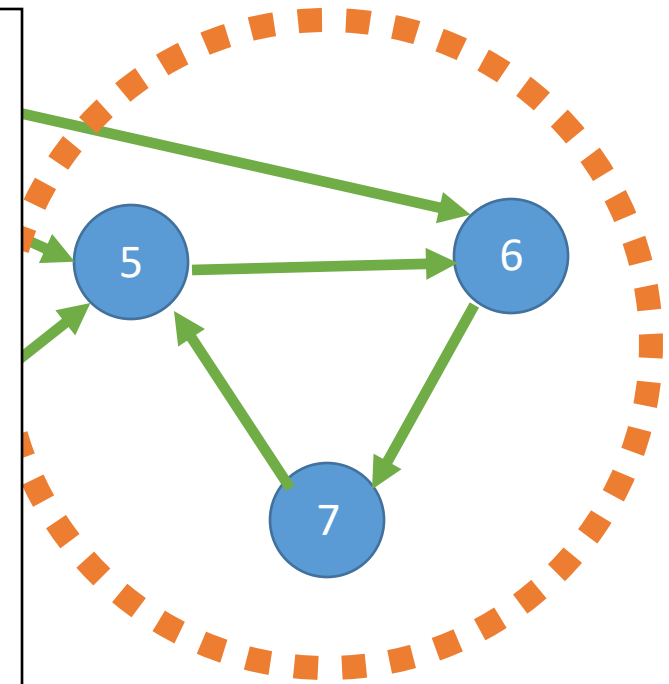
What are the strongly connected components of this graph?

```
G = {  
  1: [2],  
  2: [3, 4],  
  3: [1, 8, 9],  
  4: [5, 6],  
  5: [6],  
  6: [7],  
  7: [5],  
  8: [5, 10, 11],  
  9: [8],  
 10: [11],  
 11: [9]  
}
```

What are the strongly connected components of this graph?



```
G = {  
  1: [2],  
  2: [3, 4],  
  3: [1, 8, 9],  
  4: [5, 6],  
  5: [6],  
  6: [7],  
  7: [5],  
  8: [5, 10, 11],  
  9: [8],  
 10: [11],  
 11: [9]  
}
```



# Can we use DFS?

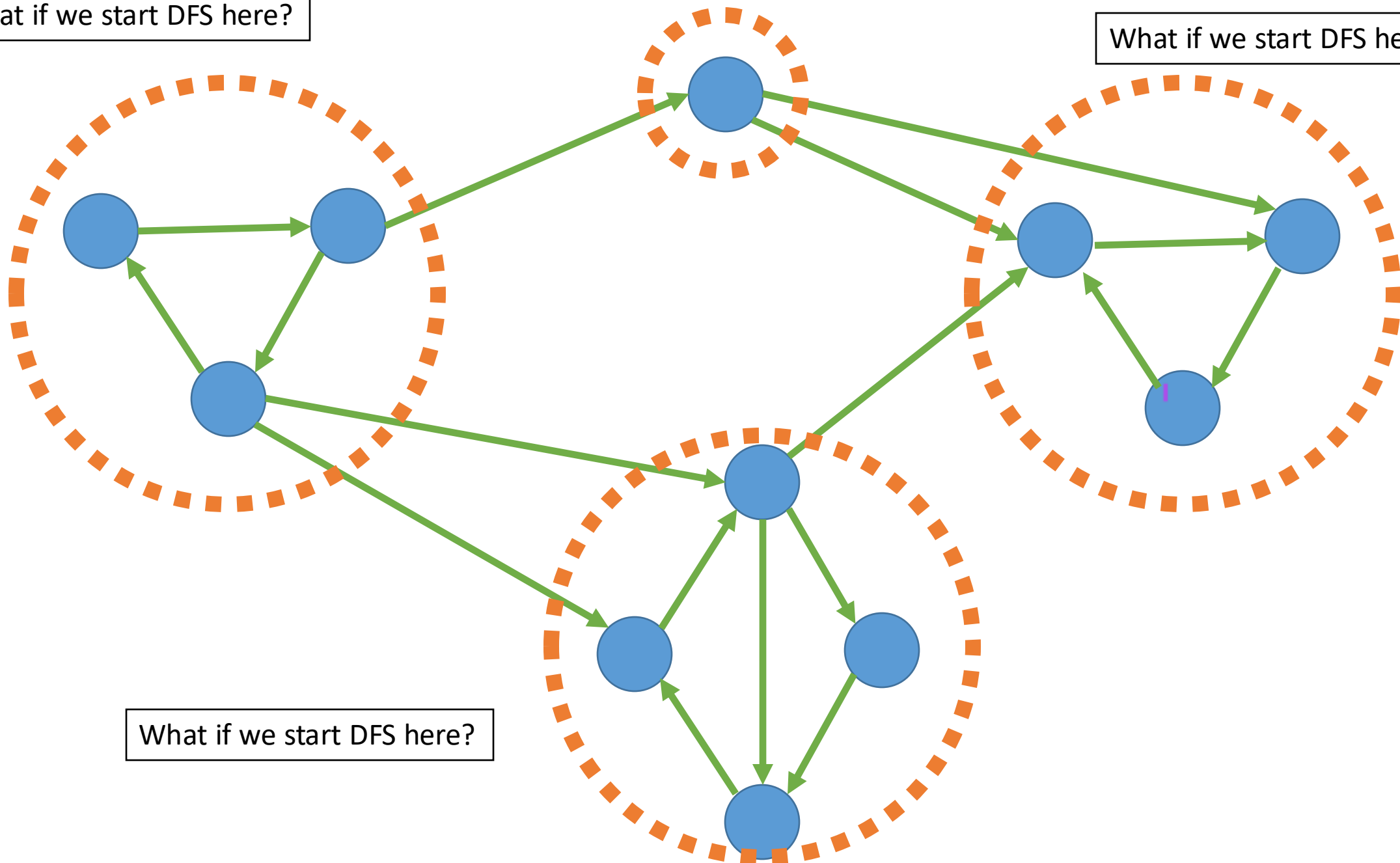
What does a DFS do?

- Finds everything that is findable
- Does not visit any vertex more than once

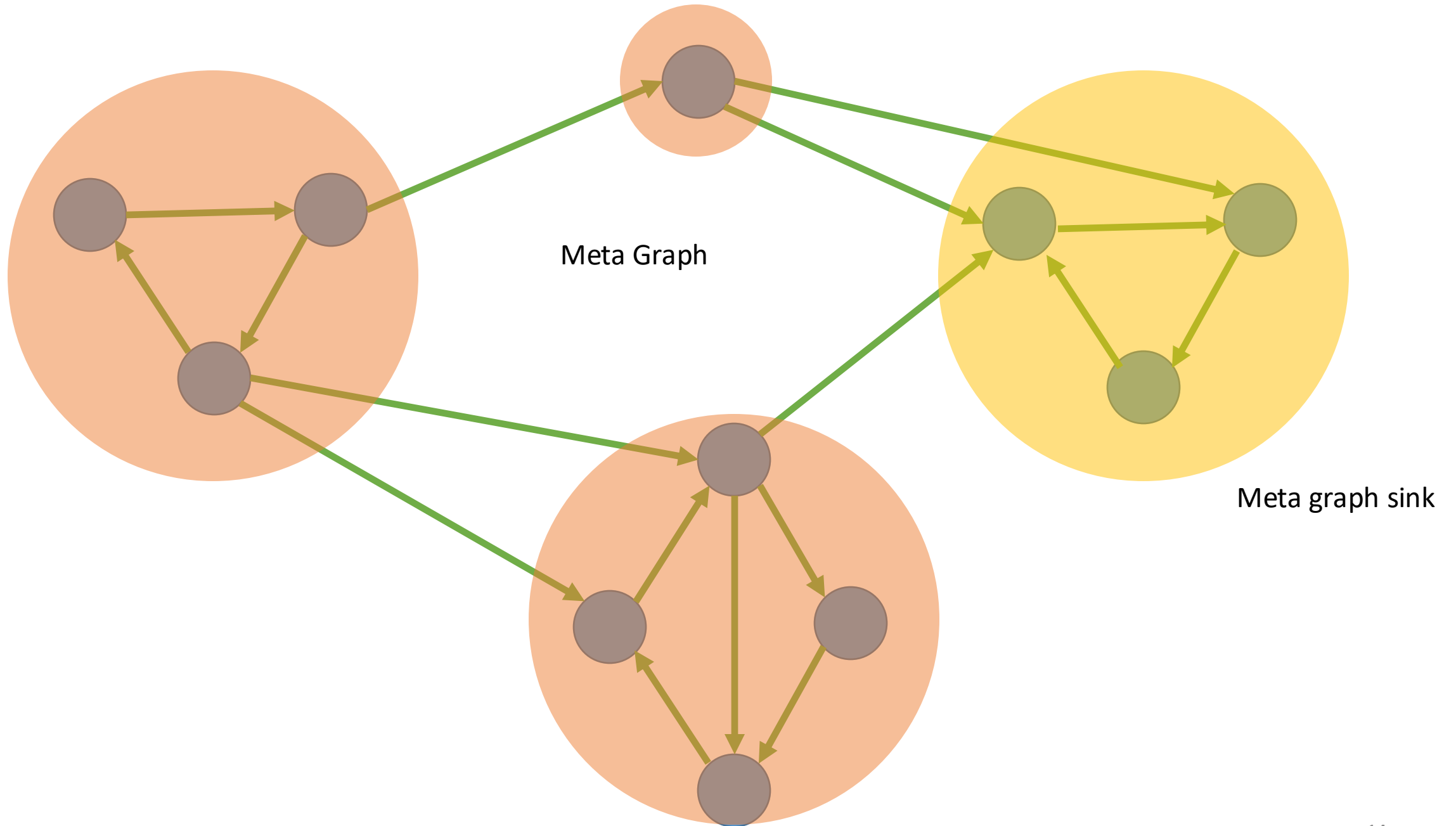
So, what can we find from each of the different nodes?

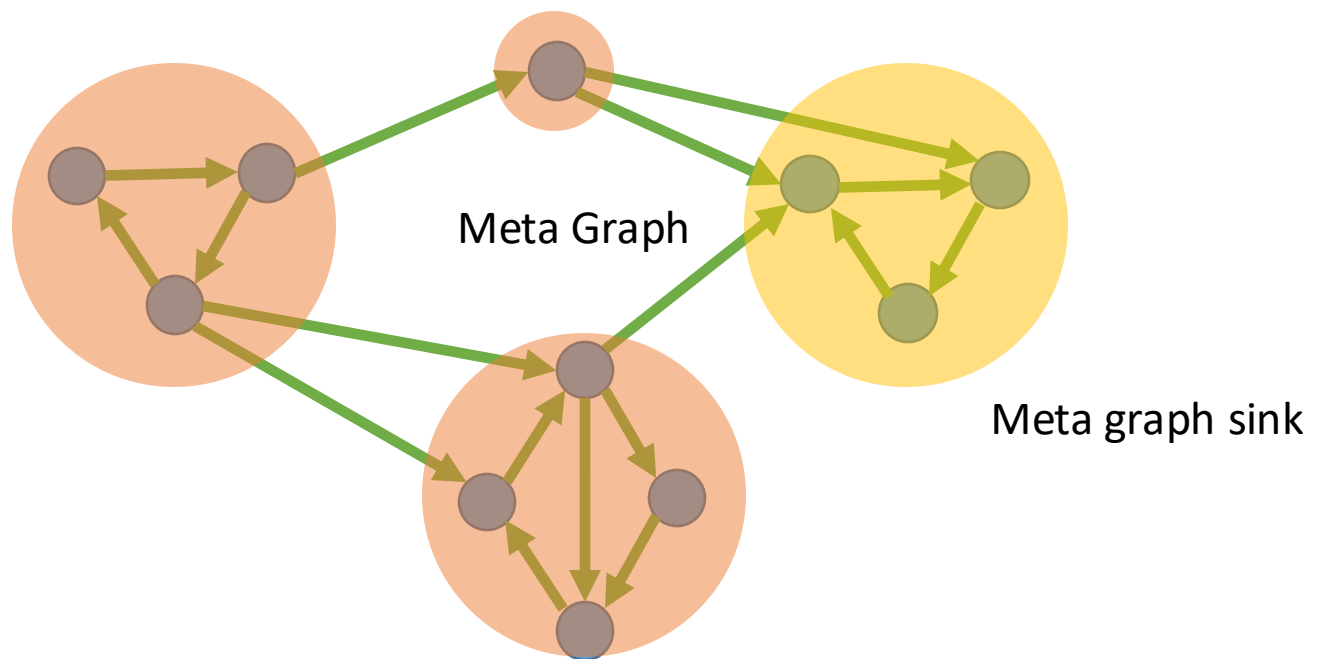
What if we start DFS here?

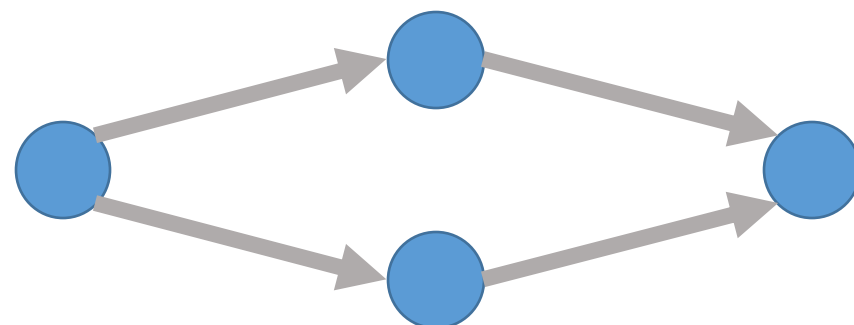
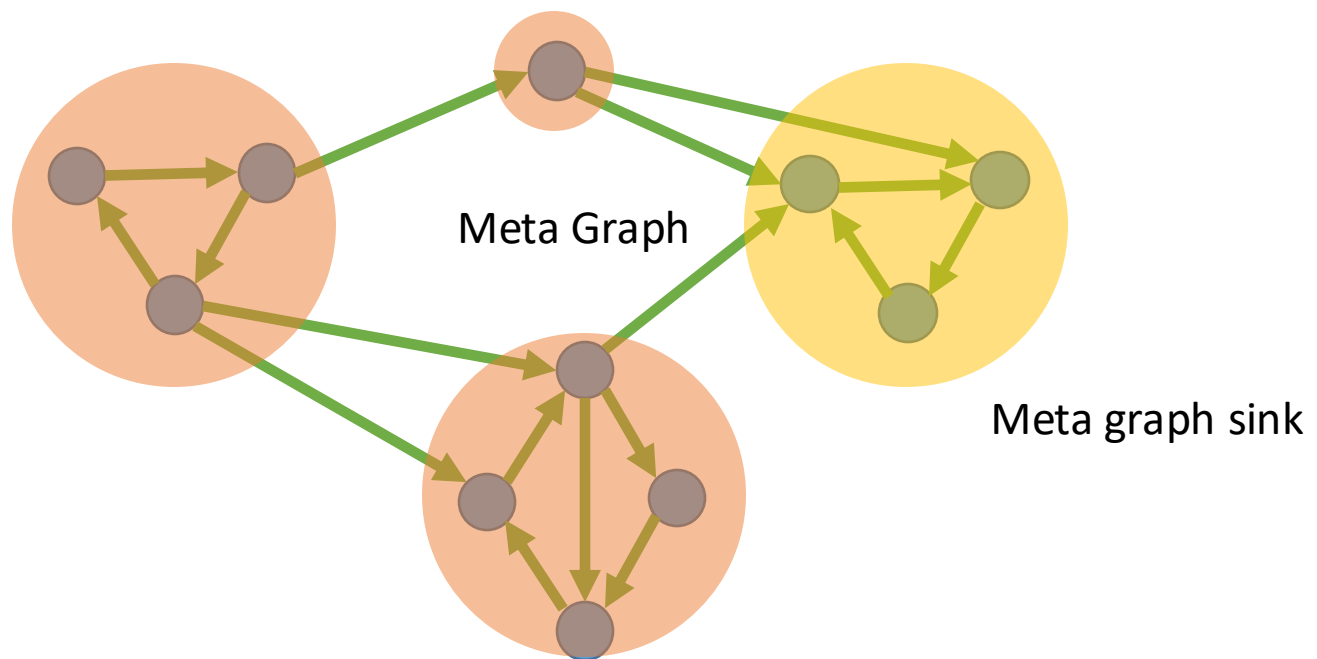
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What if we start DFS here?





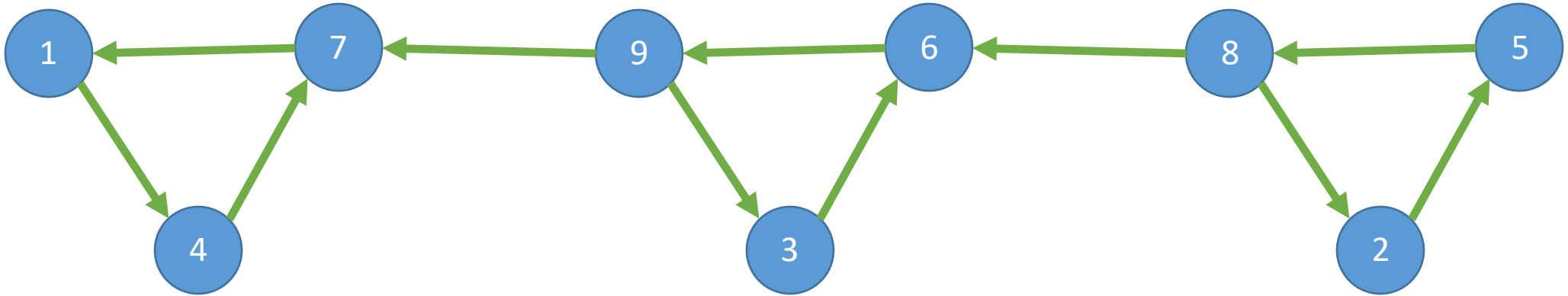


# Kosaraju

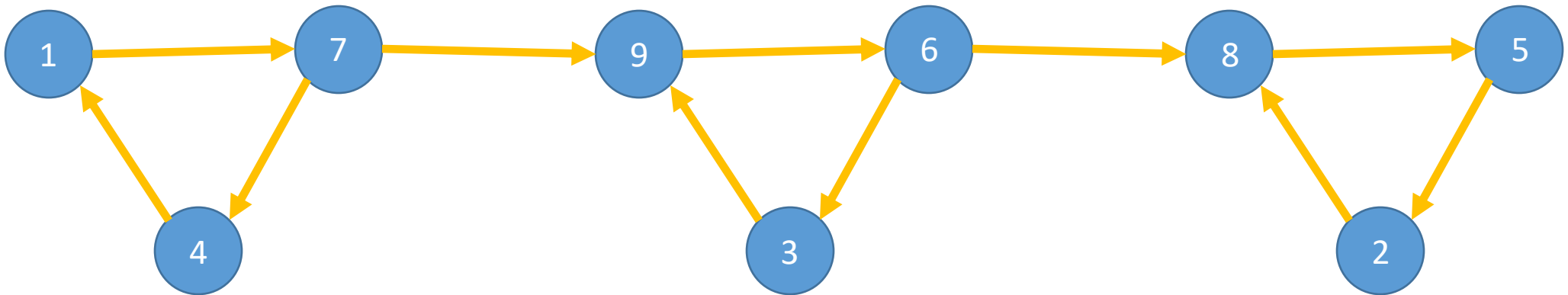
Computes the SCCs in  $O(m + n)$  time **(linear!)**

1. Create a reverse version of the  $G$  called  $G_{\text{reversed}}$

G



G\_reversed



# Kosaraju

Computes the SCCs in  $O(m + n)$  time **(linear!)**

1. Create a reverse version of the **G** called **G\_reversed**

2. Run **KosarajuLabels** on **G\_reversed**

Compute a topological order of the meta graph

3. Create a relabeled version of the **G** called **G\_relabeled**

4. Run **KosarajuLeaders** on **G\_relabeled**

Explore vertices in the new order

**FUNCTION** Kosaraju(G)

    G\_reversed = reverse\_graph(G)

    new\_labels = KosarajuLabels(G\_reversed)

    G\_relabeled = relabel\_graph(G, new\_labels)

    leaders = KosarajuLeaders(G\_relabeled)

**RETURN** leaders

**FUNCTION** Kosaraju(G)

    G\_reversed = reverse\_graph(G)

    new\_labels = KosarajuLabels(G\_reversed)

    G\_relabeled = relabel\_graph(G, new\_labels)

    leaders = KosarajuLeaders(G\_relabeled)

**RETURN** leaders

```

FUNCTION Kosaraju(G)
    G_reversed = reverse_graph(G)
    new_labels = KosarajuLabels(G_reversed)

    G_relabeled = relabel_graph(G, new_labels)
    leaders = KosarajuLeaders(G_relabeled)

    RETURN leaders

```

```

FUNCTION KosarajuLabels(G)
    found = {v: FALSE FOR v IN G.vertices}
    label = 0
    labels = {v: NONE FOR v IN G.vertices}

    FOR v IN G.vertices
        IF found[v] == FALSE
            DFSLabels(G, v, found, label, labels)

    RETURN labels

FUNCTION DFSLabels(G, v, found, label, labels)
    found[v] = TRUE
    FOR vOther IN G.edges[v]
        IF found[vOther] == FALSE
            DFSLabels(G, vOther, found, label, labels)
    label = label + 1
    labels[v] = label

```

```

FUNCTION Kosaraju(G)
    G_reversed = reverse_graph(G)
    new_labels = KosarajuLabels(G_reversed)

    G_relabeled = relabel_graph(G, new_labels)
    leaders = KosarajuLeaders(G_relabeled)

    RETURN leaders

```

```

FUNCTION KosarajuLeaders(G)
    found = {v: FALSE FOR v IN G.vertices}
    leaders = {v: NONE FOR v IN G.vertices}

    FOR v IN G.vertices.reverse_order
        IF found[v] == FALSE
            leader = v
            DFSLeaders(G, v, found, leader, leaders)

    RETURN leaders

FUNCTION DFSLeaders(G, v, found, leader, leaders)
    found[v] = TRUE
    leaders[v] = leader
    FOR vOther IN G.edges[v]
        IF found[vOther] == FALSE
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```

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```

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    RETURN leaders

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```

These are typically implemented in a single function

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FUNCTION KosarajuLabels(G)
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    RETURN labels

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    found[v] = TRUE
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    labels[v] = label

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    label = label + 1
    labels[v] = label

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    found[v] = TRUE
    leaders[v] = leader
    FOR vOther IN G.edges[v]
        IF found[vOther] == FALSE
            DFSLeaders(G, vOther, found, leader, leaders)
  
```

These are typically implemented in a single function

Does both labels and leaders.

```
FUNCTION KosarajuLoop(G)
    found = {v: FALSE FOR v IN G.vertices}
    label = 0
    labels = {v: NONE FOR v IN G.vertices}
    leaders = {v: NONE FOR v IN G.vertices}

    FOR v IN G.vertices.reverse_order
        IF found[v] == FALSE
            leader = v
            KosarajuDFS(G, v, found, label, labels, leader, leaders)

    RETURN labels, leaders

FUNCTION KosarajuDFS(G, v, found, label, labels, leader, leaders)
    found[v] = TRUE
    leaders[v] = leader
    FOR vOther IN G.edges[v]
        IF found[vOther] == FALSE
            KosarajuDFS(G, v, found, label, labels, leader, leaders)
    label = label + 1
    labels[v] = label
```

**FUNCTION** Kosaraju(G)

G\_reversed = reverse\_graph(G)

new\_labels = KosarajuLabels(G\_reversed)

G\_relabeled = relabel\_graph(G, new\_labels)

leaders = KosarajuLeaders(G\_relabeled)

**RETURN** leaders

**FUNCTION** Kosaraju(G)

G\_reversed = reverse\_graph(G)

new\_labels, \_ = KosarajuLoop(G\_reversed)

G\_relabeled = relabel\_graph(G, new\_labels)

\_, leaders = KosarajuLoop(G\_relabeled)

**RETURN** leaders

# Kosaraju

Computes the SCCs in  $O(m + n)$  time **(linear!)**

1. Create a reverse version of the **G** called **G\_reversed**

2. Run **KosarajuLoop** on **G\_reversed**

Compute a topological order of the meta graph

3. Create a relabeled version of the **G** called **G\_relabeled**

4. Run **KosarajuLoop** on **G\_relabeled**

Explore vertices in the new order

What are  
the SCCs?

**FUNCTION** Kosaraju(G)

G\_reversed = reverse\_graph(G)

new\_labels, \_ = KosarajuLoop(G\_reversed)

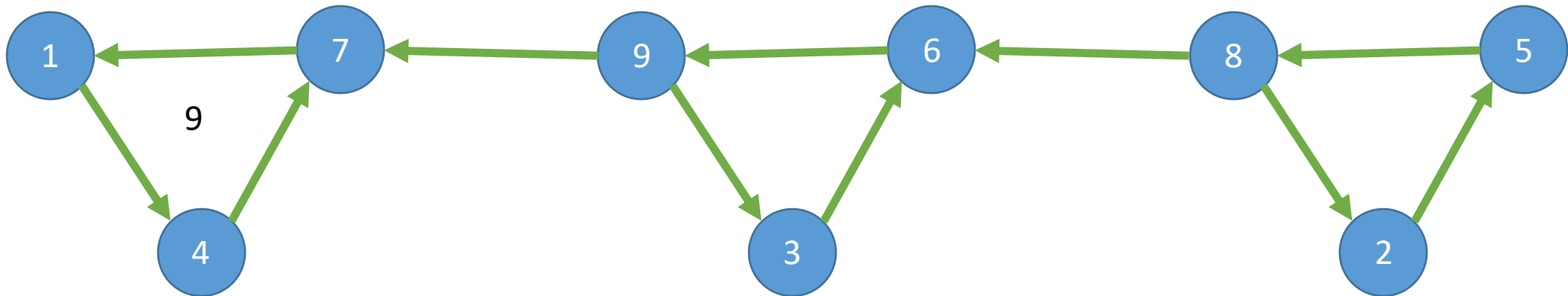
G\_relabeled = relabel\_graph(G, new\_labels)

\_, leaders = KosarajuLoop(G\_relabeled)

**RETURN** leaders

Where do we want to start DFS if  
we are looking for SCCs? (Which  
SCC is best to find first?)

G



**FUNCTION** Kosaraju(G)

`G_reversed = reverse_graph(G)`

`new_labels, _ = KosarajuLoop(G_reversed)`

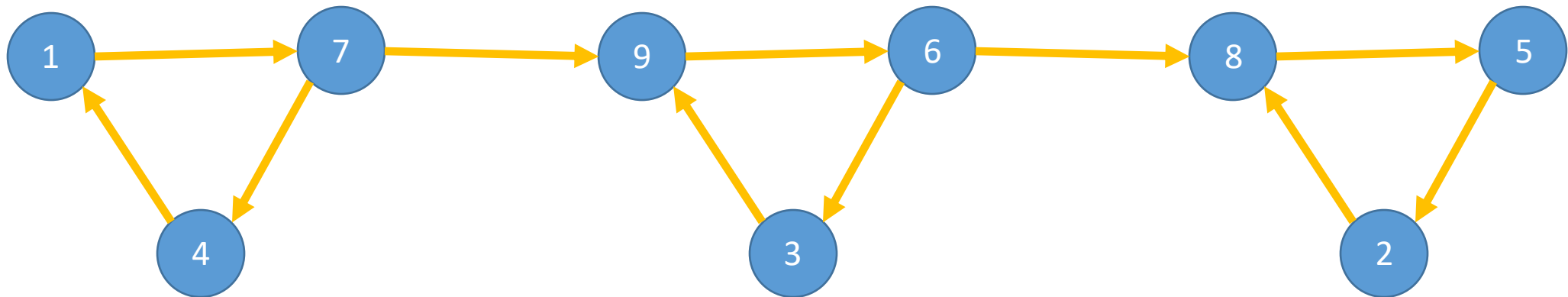
`G_relabeled = relabel_graph(G, new_labels)`

`_, leaders = KosarajuLoop(G_relabeled)`

**RETURN** leaders

Where do we want to start DFS if we are looking for SCCs? (Which SCC is best to find first?)

**G\_reversed**



**FUNCTION** Kosaraju(G)

G\_reversed = reverse\_graph(G)

new\_labels, \_ = KosarajuLoop(G\_reversed)

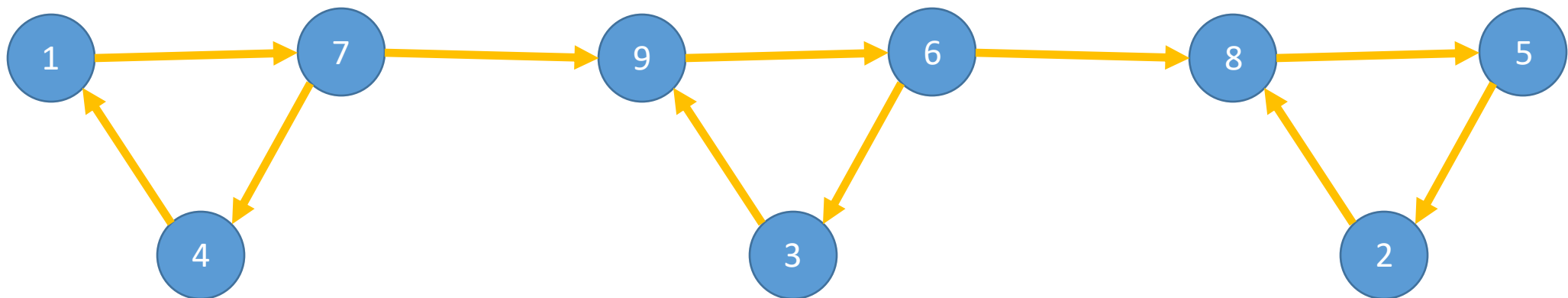
G\_relabeled = relabel\_graph(G, new\_labels)

\_, leaders = KosarajuLoop(G\_relabeled)

**RETURN** leaders

Where do we want to start DFS if we are looking for SCCs? (Which SCC is best to find first?)

G\_reversed



```

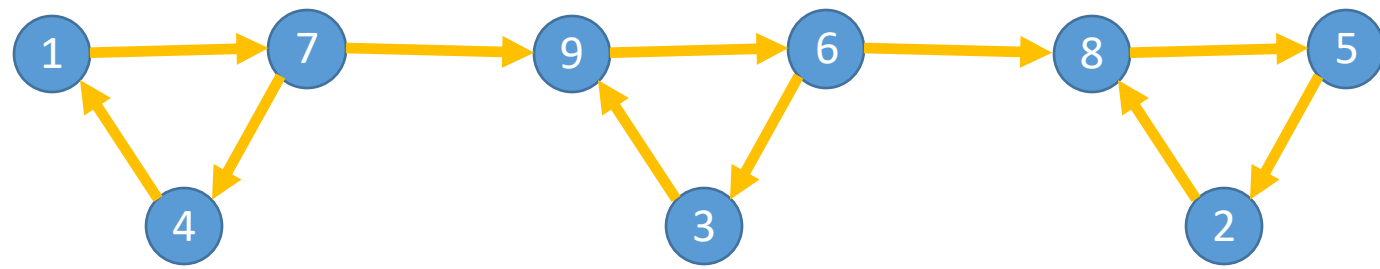
FUNCTION KosarajuLoop(G)
    found = {v: FALSE FOR v IN G.vertices}
    label = 0
    labels = {v: NONE FOR v IN G.vertices}
    leaders = {v: NONE FOR v IN G.vertices}

    FOR v IN G.vertices.reverse_order
        IF found[v] == FALSE
            leader = v
            KosarajuDFS(...)

    RETURN labels, leaders

```

Can we start anywhere?



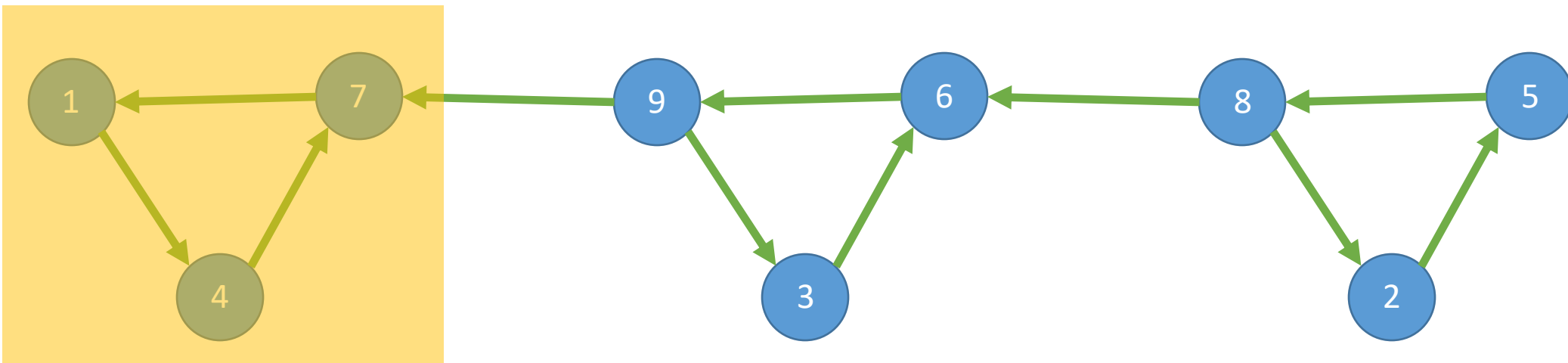
```

FUNCTION KosarajuDFS(...)
    found[v] = TRUE
    leaders[v] = leader
    FOR v0ther IN G.edges[v]
        IF found[v0ther] == FALSE
            KosarajuDFS(...)
    label = label + 1
    labels[v] = label

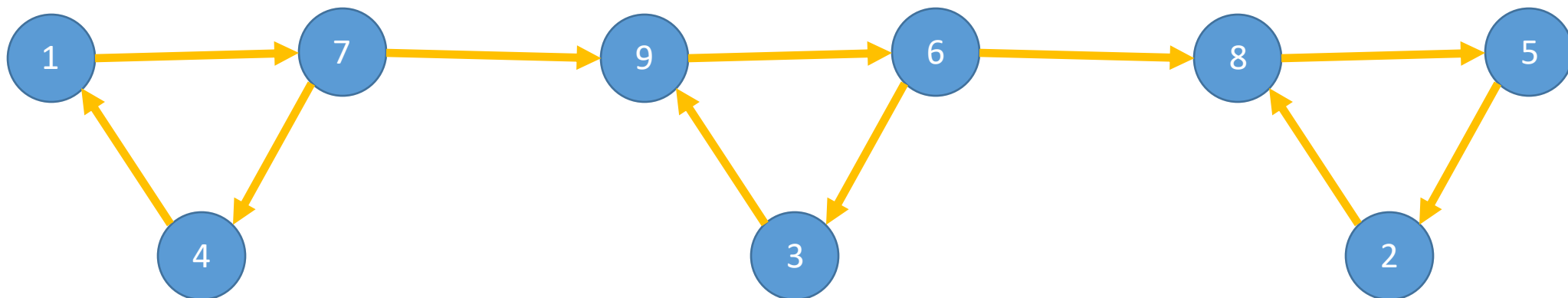
```

Ignore leaders the first pass  
Ignore labels the second pass

G

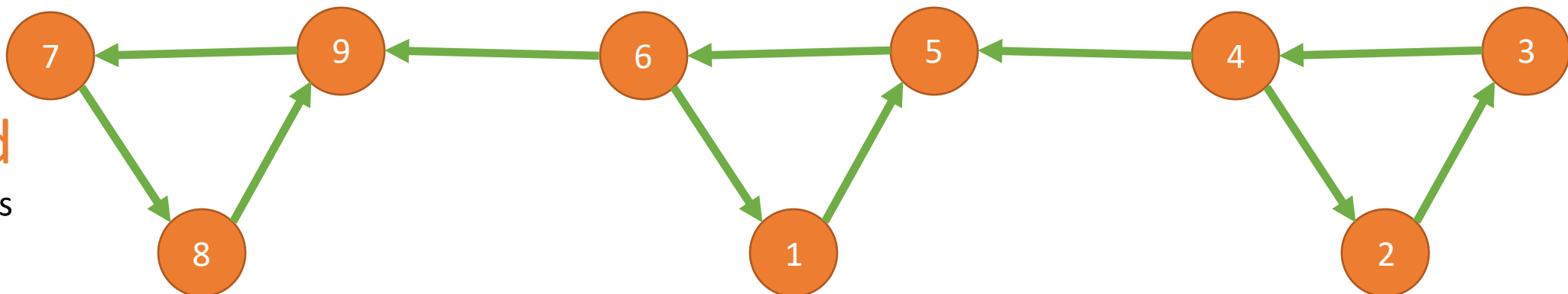


G\_reversed



G\_relabeled

Multiple possibilities



**FUNCTION** Kosaraju(G)

    G\_reversed = reverse\_graph(G)

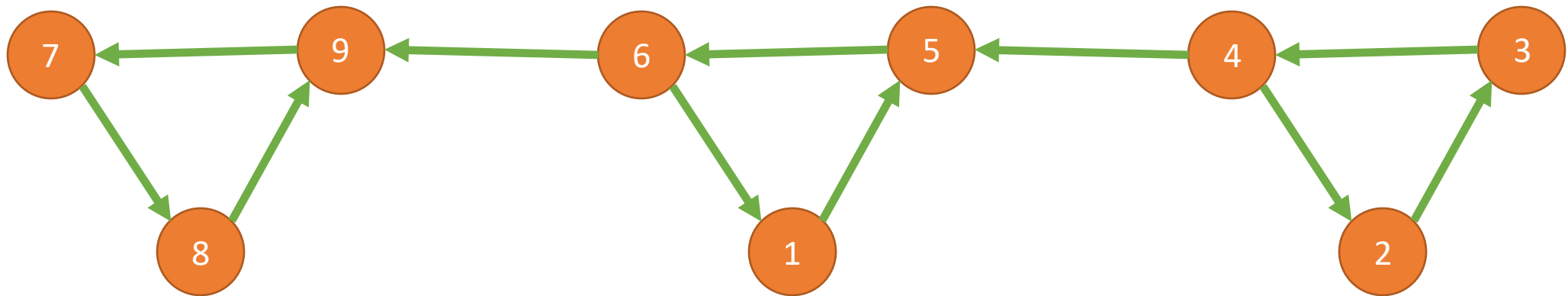
    new\_labels, \_ = KosarajuLoop(G\_reversed)

    G\_relabeled = relabel\_graph(G, new\_labels)

    \_, leaders = KosarajuLoop(G\_relabeled)

**RETURN** leaders

G\_relabeled



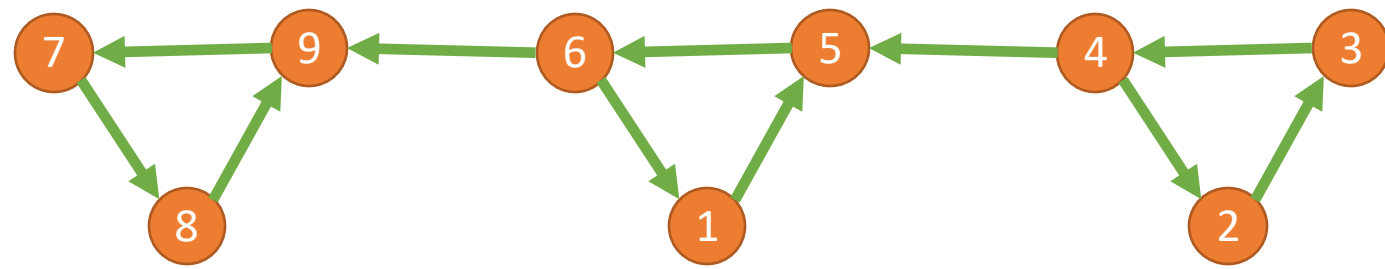
```

FUNCTION KosarajuLoop(G)
    found = {v: FALSE FOR v IN G.vertices}
    label = 0
    labels = {v: NONE FOR v IN G.vertices}
    leaders = {v: NONE FOR v IN G.vertices}

    FOR v IN G.vertices.reverse_order
        IF found[v] == FALSE
            leader = v
            KosarajuDFS(...)

    RETURN labels, leaders

```



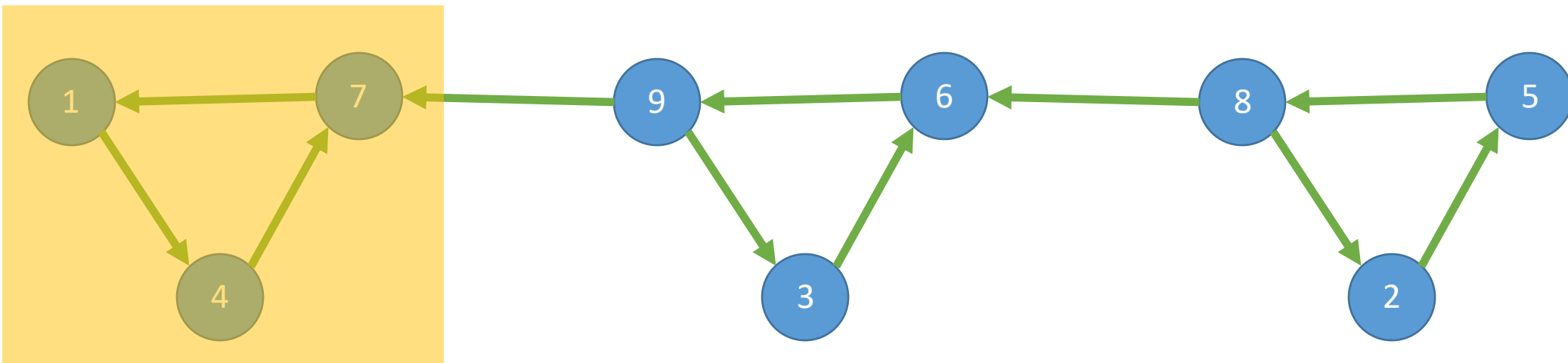
```

FUNCTION KosarajuDFS(...)
    found[v] = TRUE
    leaders[v] = leader
    FOR v0ther IN G.edges[v]
        IF found[v0ther] == FALSE
            KosarajuDFS(...)
    label = label + 1
    labels[v] = label

```

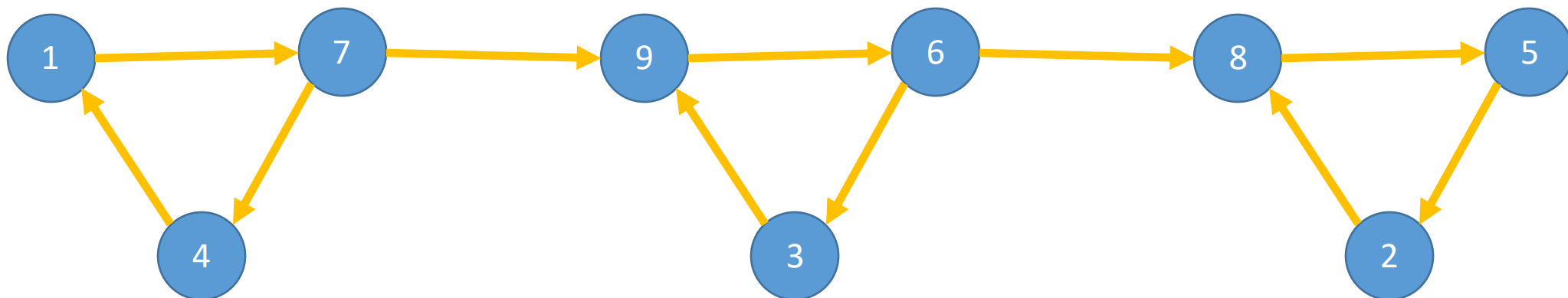
Ignore leaders the first pass  
Ignore labels the second pass

G

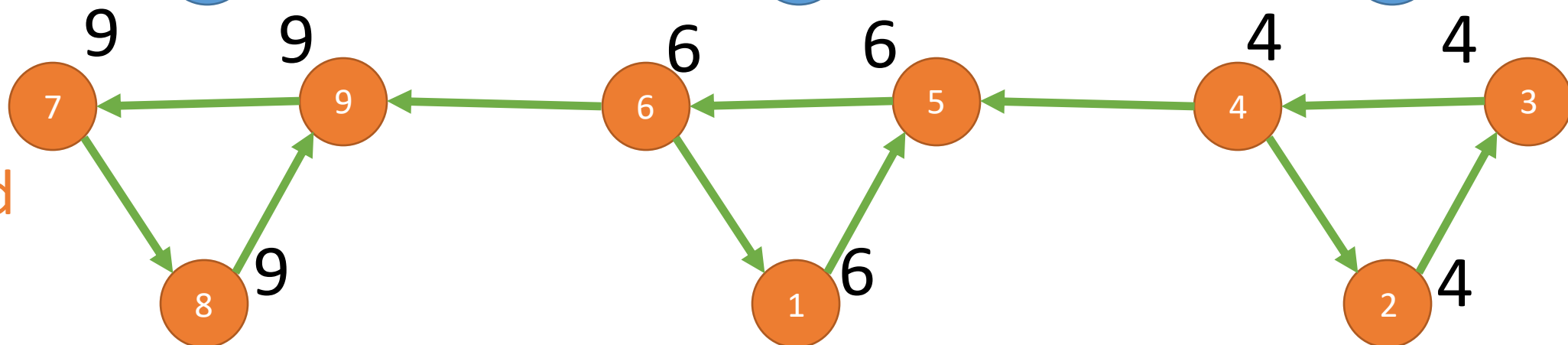


Sink SCC in  
Meta Graph

G\_reversed



G\_relabeled



**FUNCTION** Kosaraju(G)

G\_reversed = reverse\_graph(G)

new\_labels, \_ = KosarajuLoop(G\_reversed)

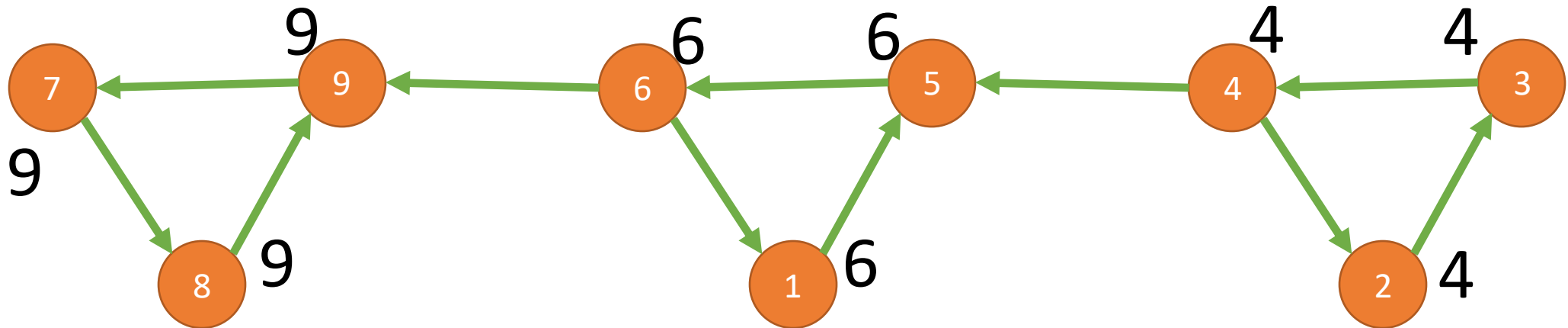
G\_relabeled = relabel\_graph(G, new\_labels)

\_, leaders = KosarajuLoop(G\_relabeled)

**RETURN** leaders

What could you do to make this API a bit nicer?

G\_relabeled



## Exercise

### FUNCTION KosarajuLoop(G)

```
found = {v: FALSE FOR v IN G.vertices}
label = 0
labels = {v: NONE FOR v IN G.vertices}
leaders = {v: NONE FOR v IN G.vertices}
```

```
FOR v IN G.vertices.reverse_order
```

```
  IF found[v] == FALSE
```

```
    leader = v
```

```
    KosarajuDFS(G, v, found, label, labels, leader, leaders)
```

```
RETURN labels, leaders
```

### FUNCTION KosarajuDFS(G, v, found, label, labels, leader, leaders)

```
found[v] = TRUE
```

```
leaders[v] = leader
```

```
FOR vOther IN G.edges[v]
```

```
  IF found[vOther] == FALSE
```

```
    KosarajuDFS(G, vOther, found, label, labels, leader, leaders)
```

```
label = label + 1
```

```
labels[v] = label
```

### FUNCTION Kosaraju(G)

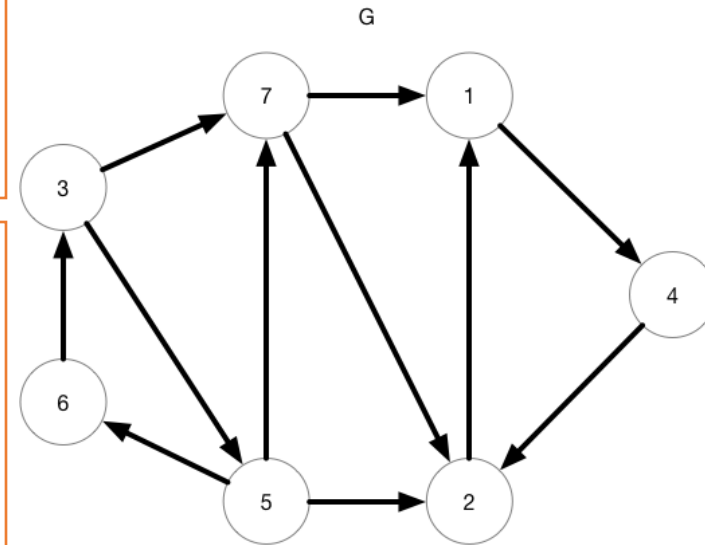
```
G_reversed = reverse_graph(G)
```

```
new_labels, _ = KosarajuLoop(G_reversed)
```

```
G_relabeled = relabel_graph(G, new_labels)
```

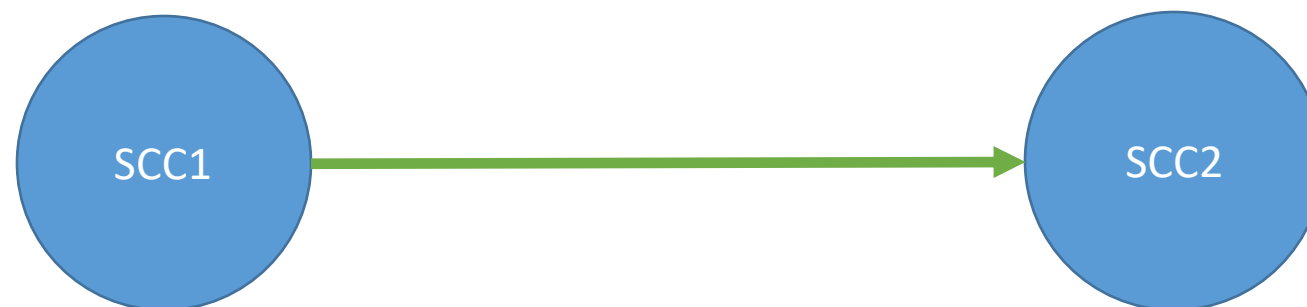
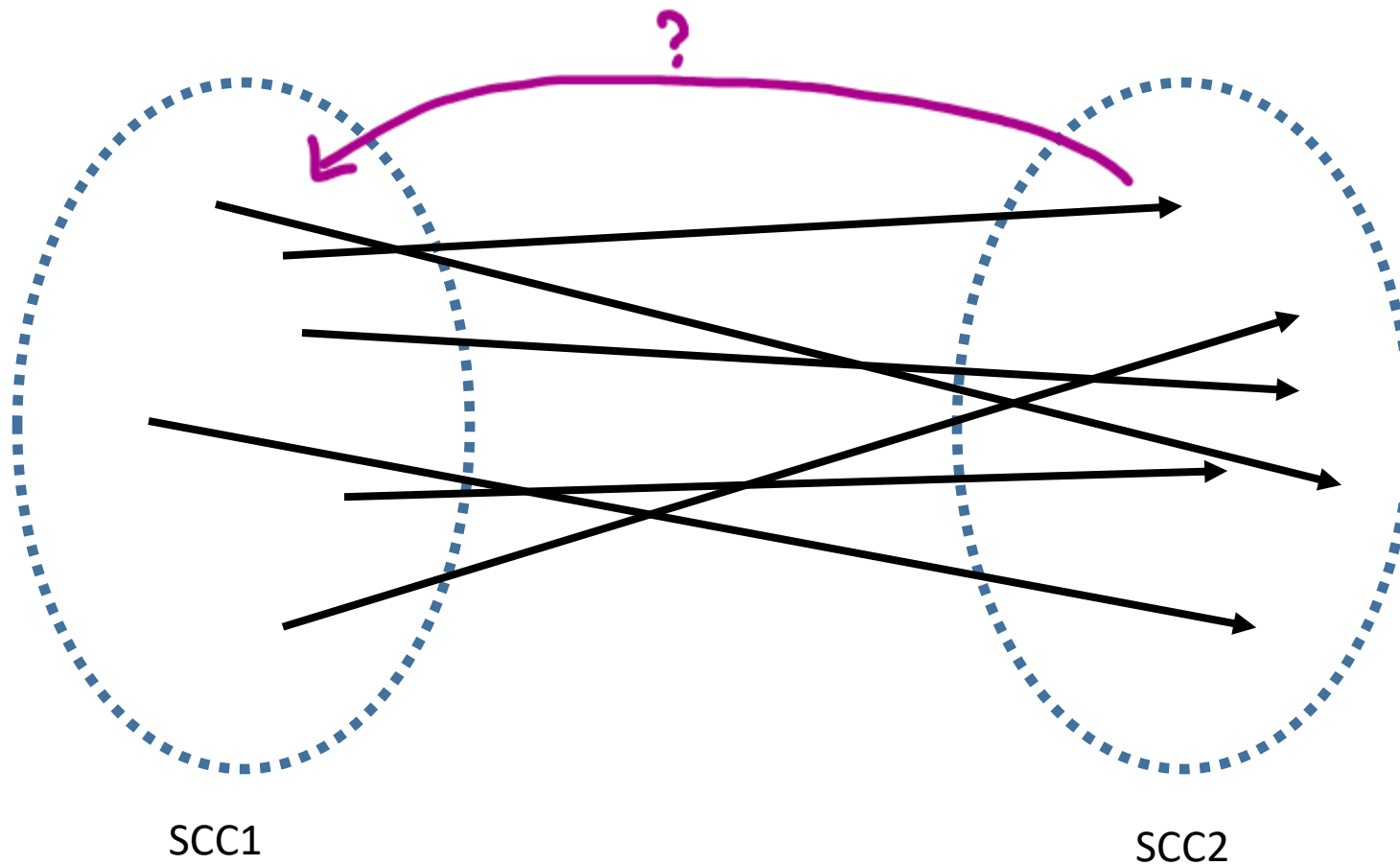
```
_, leaders = KosarajuLoop(G_relabeled)
```

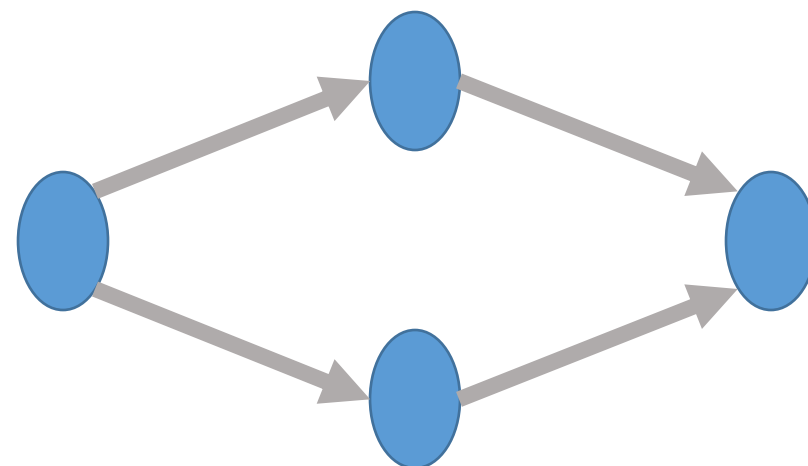
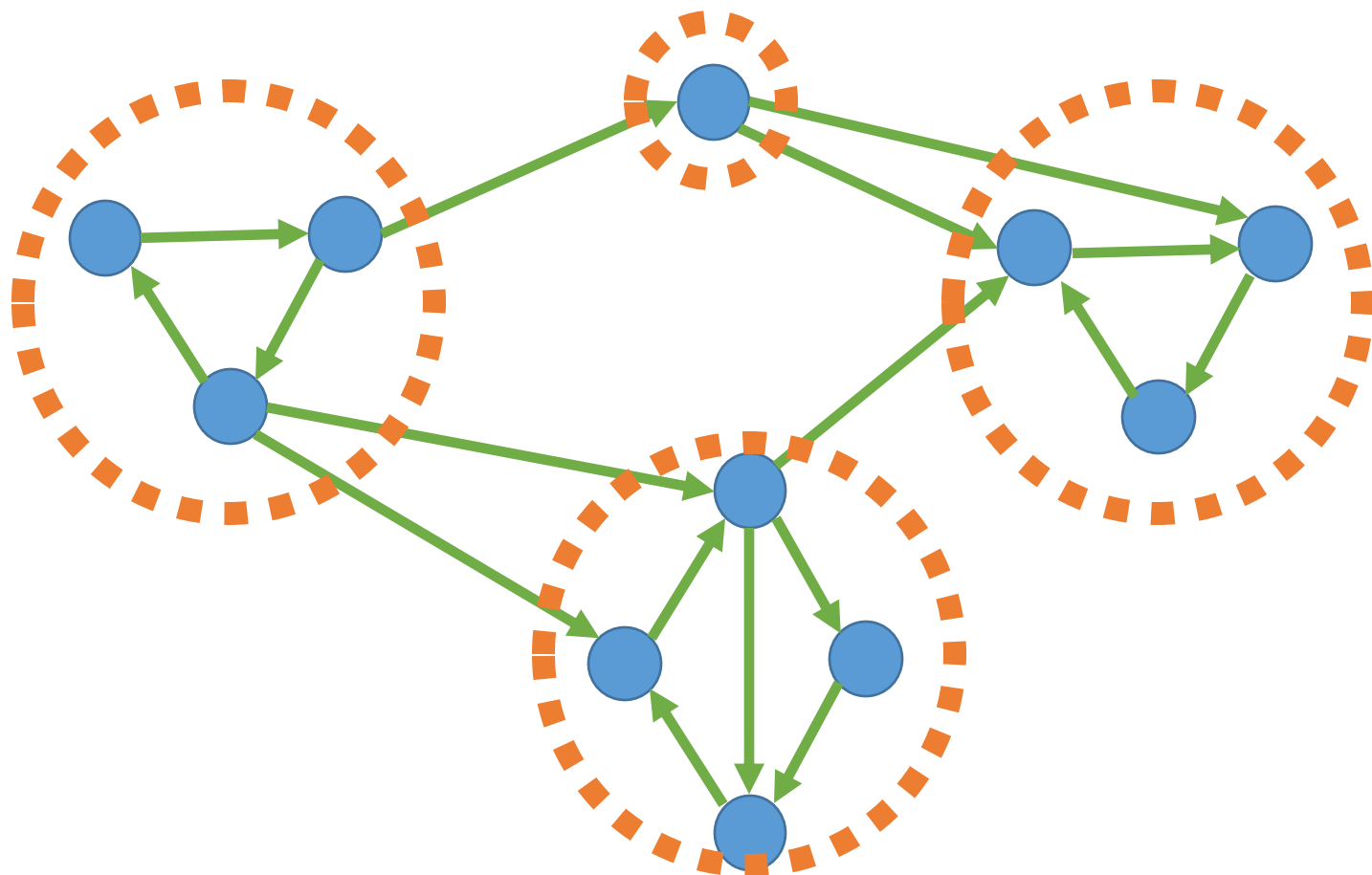
```
RETURN leaders
```

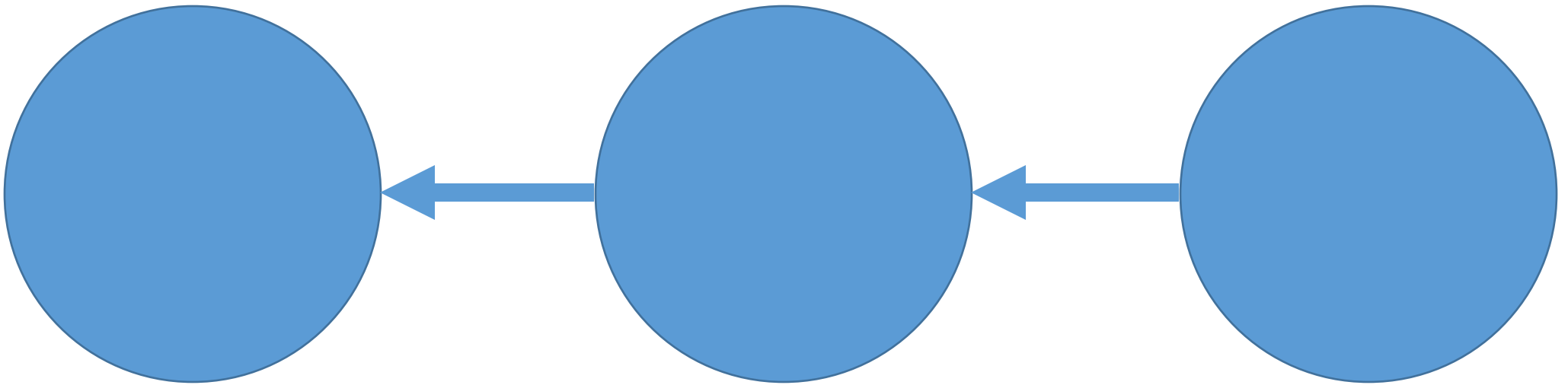
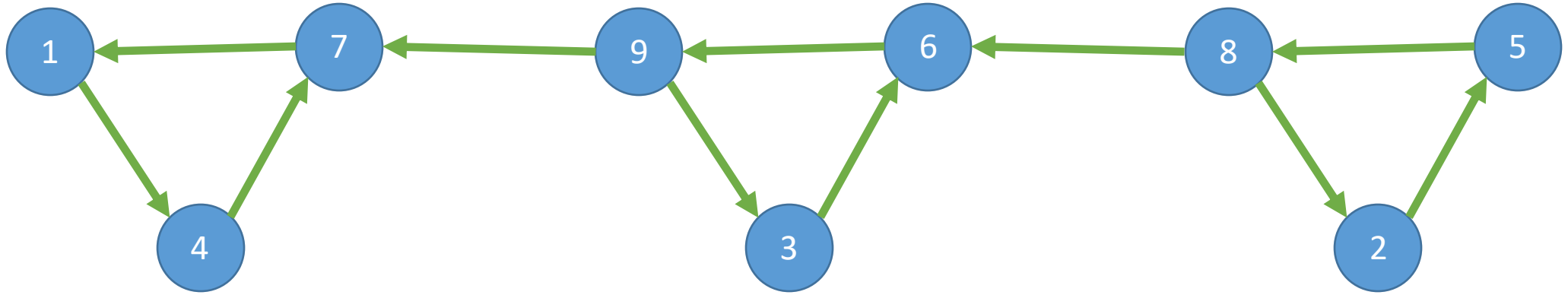


# Why does this work?

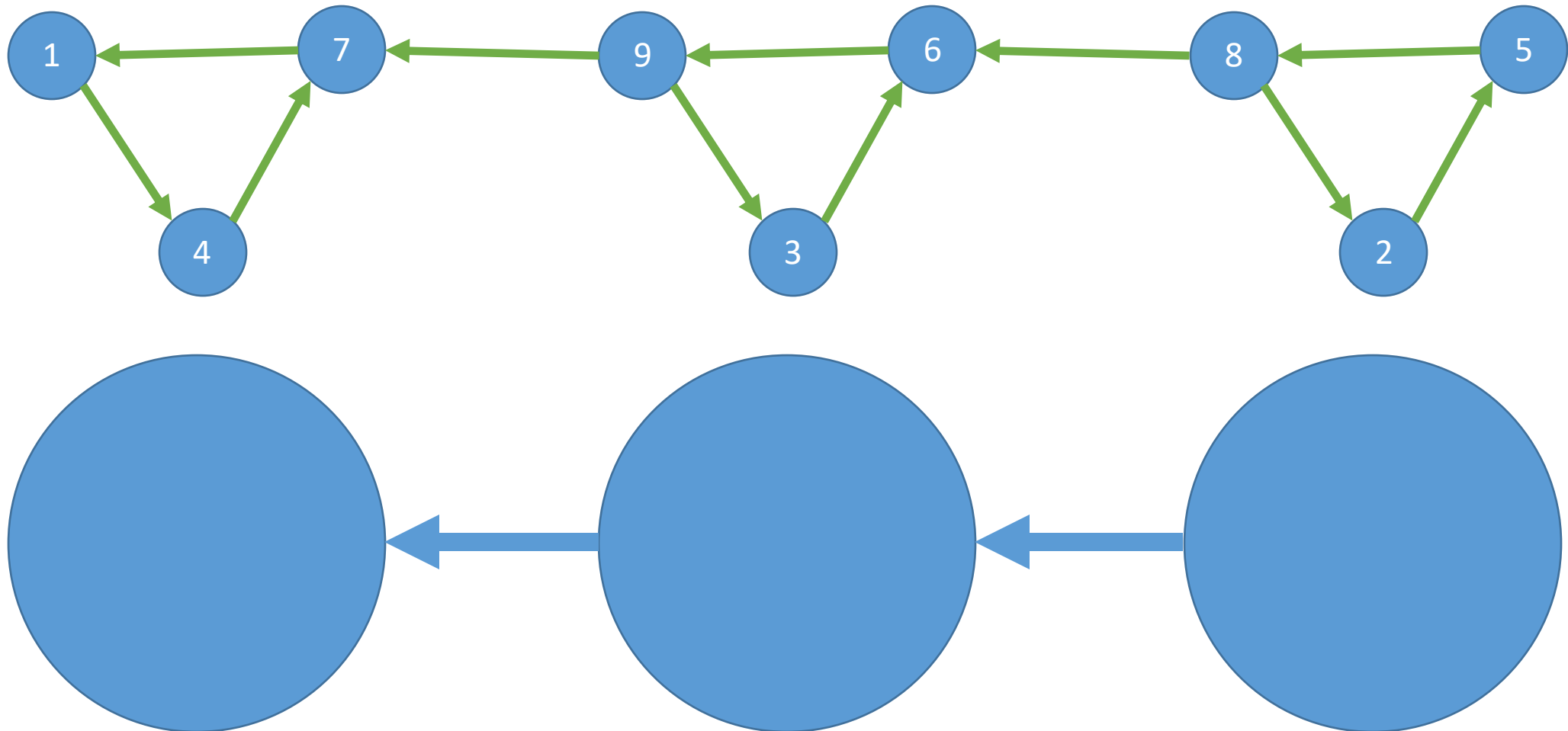
- Does this work for all graphs, or just this example?
- The SCCs of  $G$  create an **acyclic** “meta-graph”
- For the “meta-graph”
  - Vertices correspond to the SCCs
  - Edges correspond to paths among the SCCs



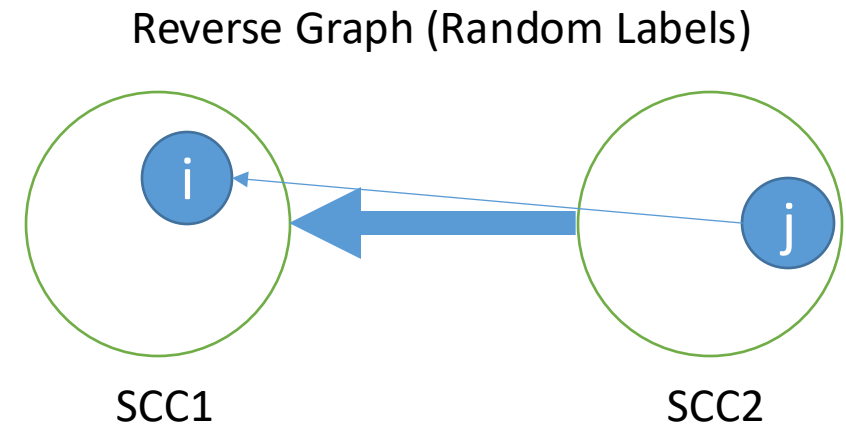
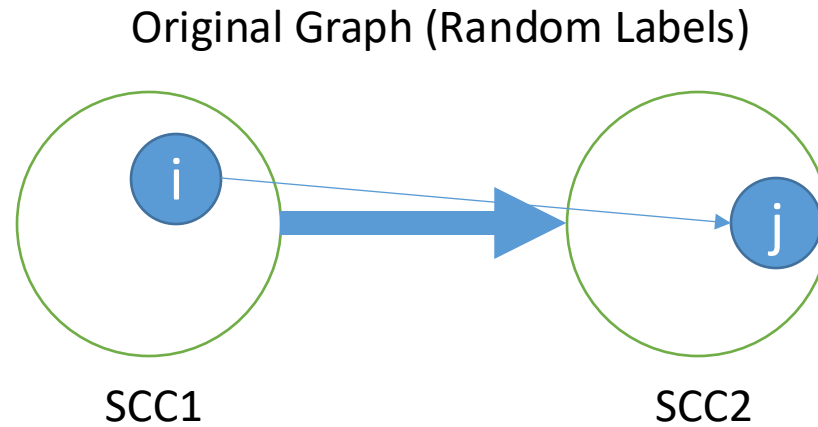




How do we know that the SCC based meta-graph is acyclic?



# Key Lemma

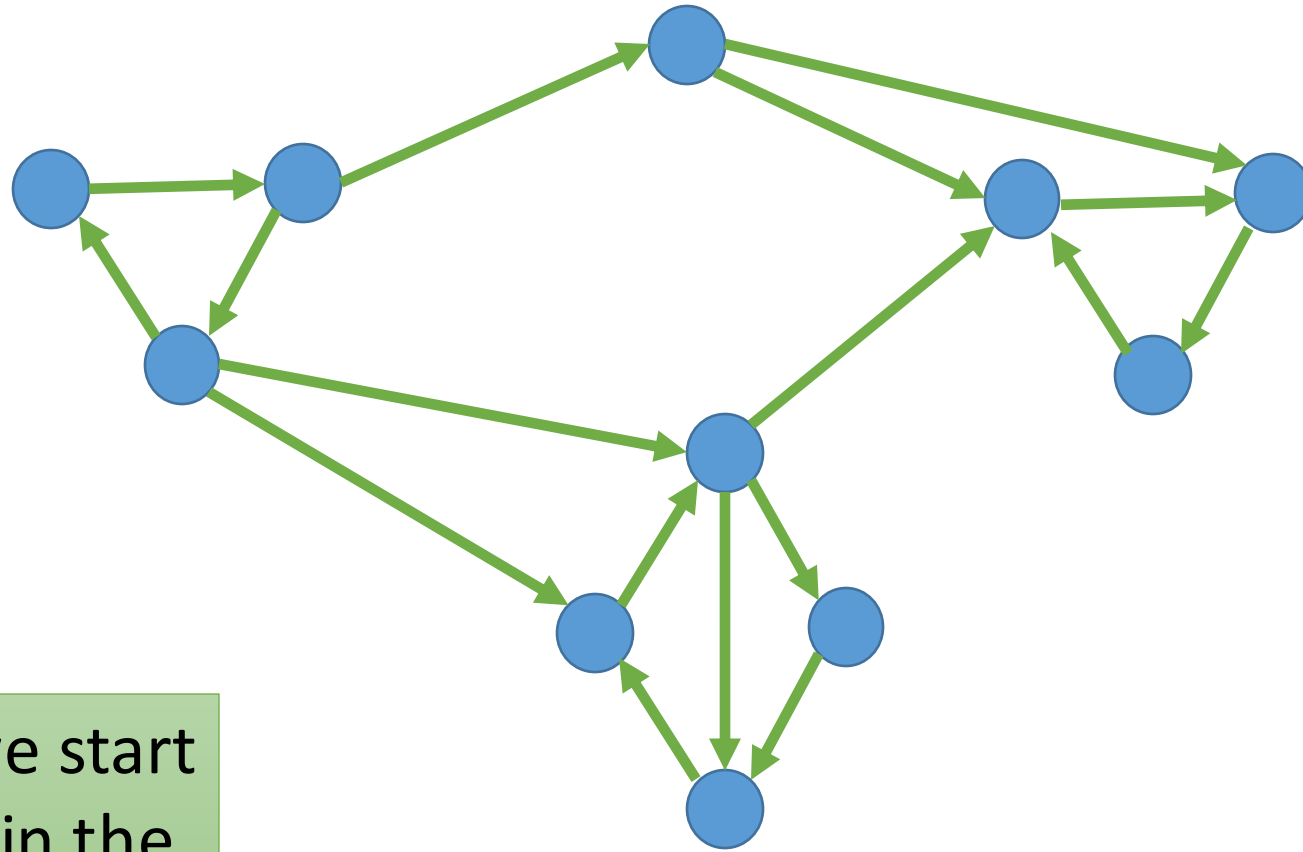


- Consider the two adjacent SCCs in the meta-graph above
- Now consider the re-labeling found from the reverse graph

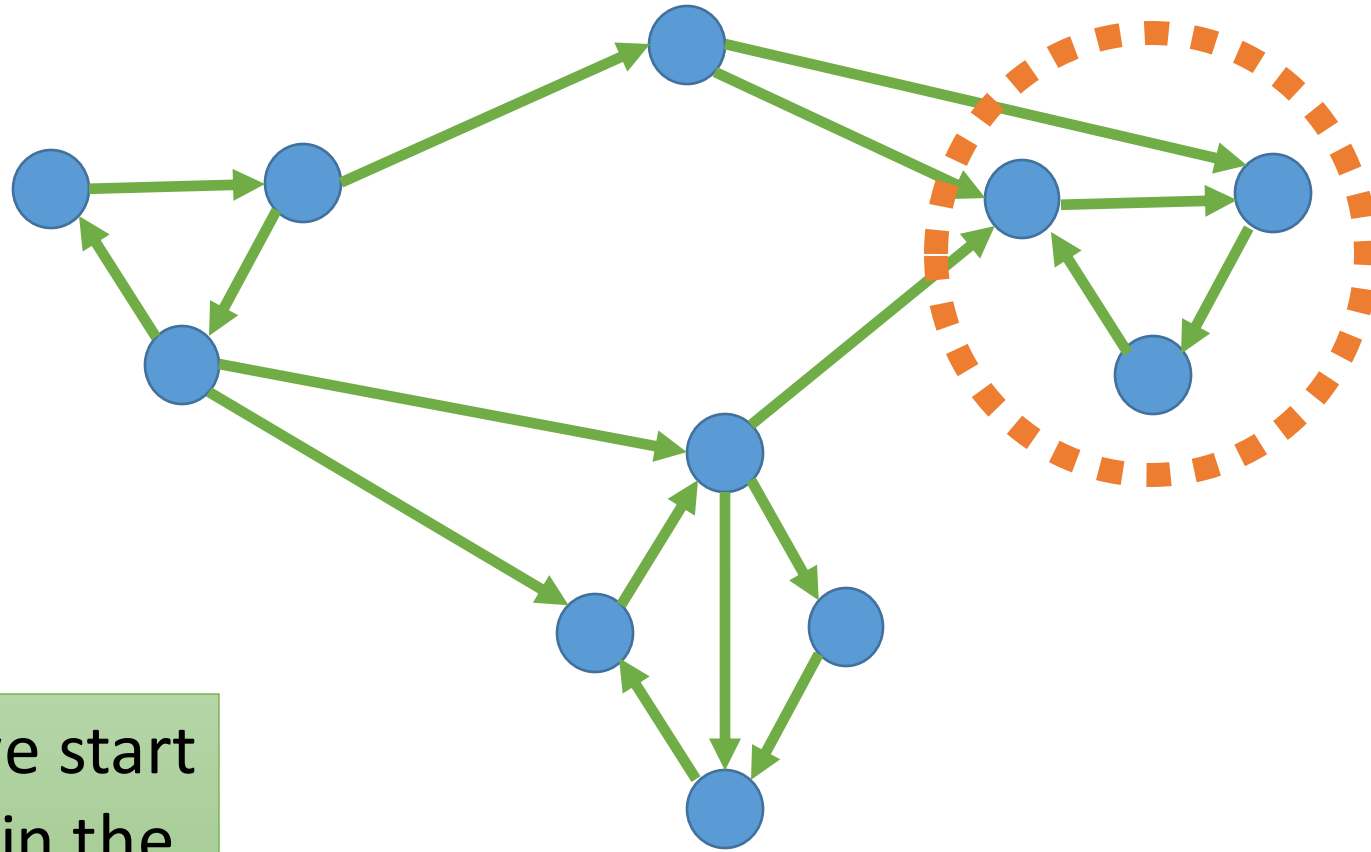
Where do we want to start DFS in the leaders pass?

- Let  $f(v)$  = the re-labeling resulting from  $\text{KosarajuLoop}(G\_reversed)$
- Then  $\max[f(.) \text{ in } SCC1] < \max[f(.) \text{ in } SCC2]$
- Corollary: the maximum label must lie in a “sink SCC” of the original graph

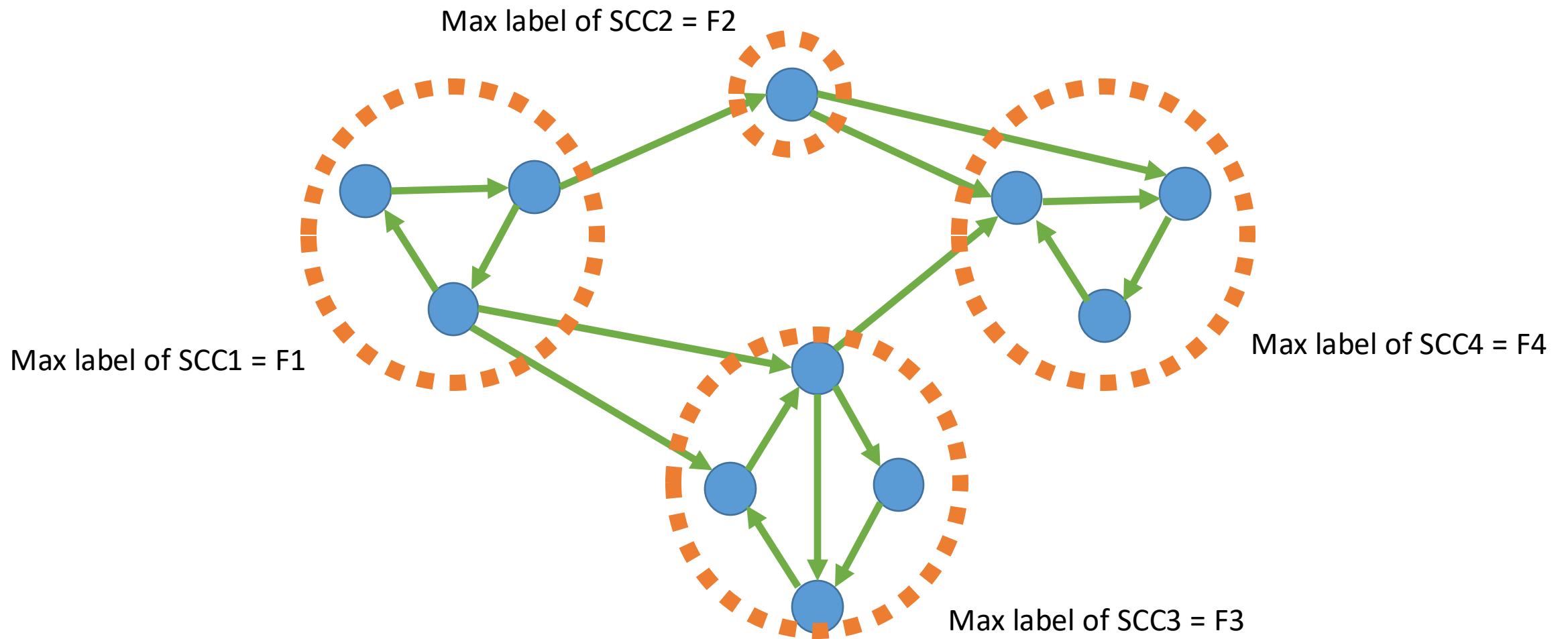
```
FUNCTION KosarajuDFS(...)
    found[v] = TRUE
    leaders[v] = leader
    FOR vOther IN G.edges[v]
        IF found[vOther] == FALSE
            KosarajuDFS(...)
    label = label + 1
    labels[v] = label
```



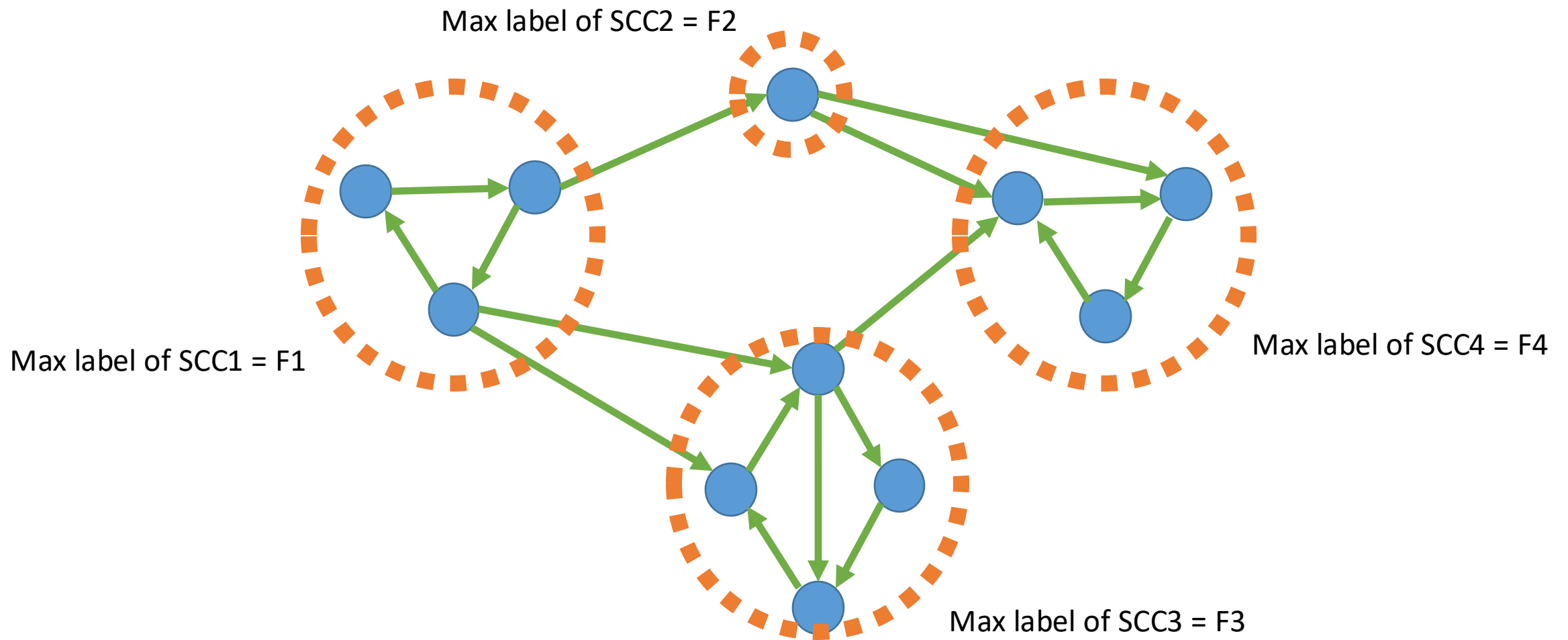
Where should we start labeling leaders in the second pass?



Where should we start labeling leaders in the second pass?



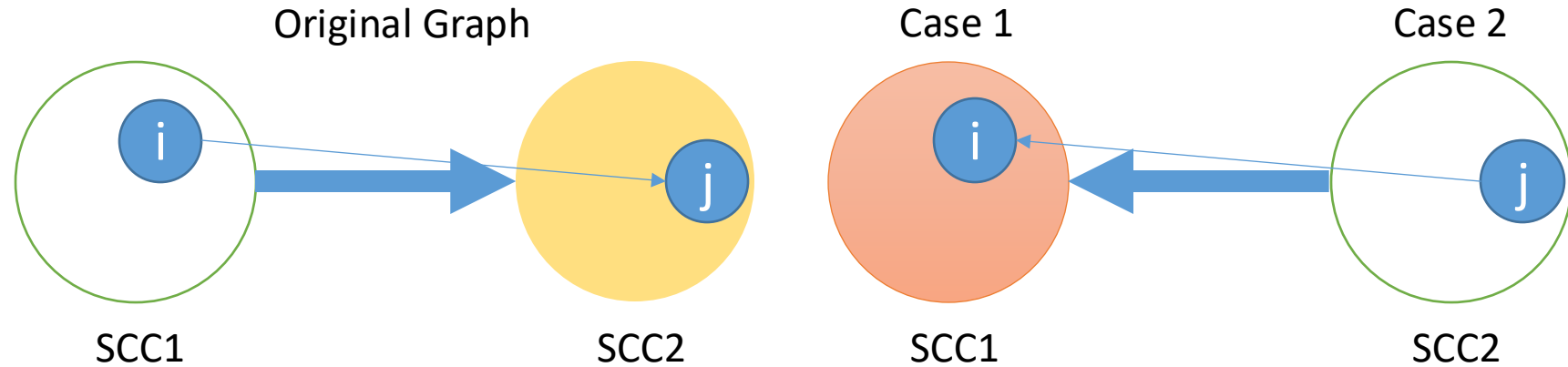
Then  $F1 < \{F2, F3\} < F4$



Then  $F1 < \{F2, F3\} < F4$

What would happen if SCC4 had a link back to SCC3?

# Proof of Lemma

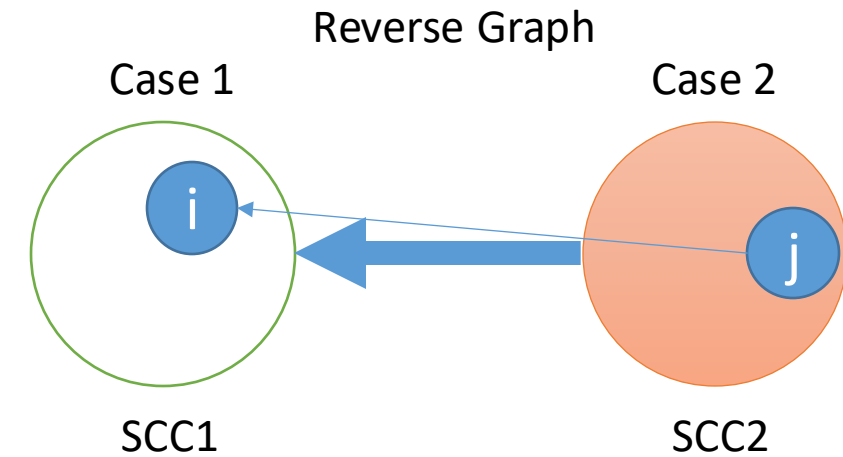
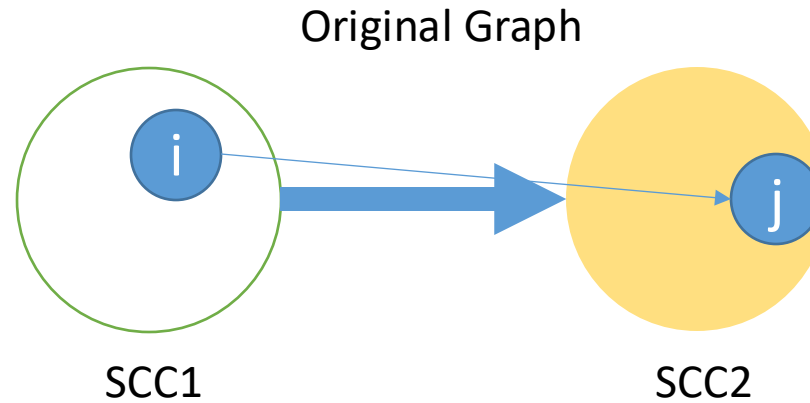


Case 1: consider the case when the first vertex that we explore is in SCC1

- Then all SCC1 is explored before SCC2
- Therefore, all labels in SCC1 are less than all labels in SCC2
- So, in the original graph we will start in SCC2 (the sink)

```
FUNCTION KosarajuDFS(...)
    found[v] = TRUE
    leaders[v] = leader
    FOR vOther IN G.edges[v]
        IF found[vOther] == FALSE
            KosarajuDFS(...)
    label = label + 1
    labels[v] = label
```

# Proof of Lemma



Case 2: consider the case when the first vertex that we explore is in SCC2

- All other vertices in SCC2 are explored before vertex j
- All vertices in SCC1 are explored before vertex j
- Therefore, all labels in SCC1 and SCC2 are less than the label of vertex j
- So, in the original graph we will start at vertex j in SCC2 (the sink)

```
FUNCTION KosarajuDFS(...)
    found[v] = TRUE
    leaders[v] = leader
    FOR vOther IN G.edges[v]
        IF found[vOther] == FALSE
            KosarajuDFS(...)
    label = label + 1
    labels[v] = label
```

# What does this mean?

- We'll start the second KosarajuLoop at an “SCC sink”
- That sink will then be *removed* (by marking all vertices in the SCC as explored) and we'll next move to the newly created sink
- And so on

# Kosaraju's Algorithm Summary

Computes the SCCs in  $O(m + n)$  time **(linear!)**

1. Create a reverse version of the  $G$  called  $G_{\text{reversed}}$
2. Run  $\text{KosarajuLoop}$  on  $G_{\text{reversed}}$ 
  - Create a topological ordering on the meta graph
3. Create a relabeled version of the  $G$  called  $G_{\text{relabelled}}$
4. Run  $\text{KosarajuLoop}$  on  $G_{\text{relabelled}}$ 
  - Find all nodes with the same "leader"