

# Breadth First Search

<https://cs.pomona.edu/classes/cs140/>

# Outline

## Topics and Learning Objectives

- Discuss breadth first search for graphs

## Exercises

- Continued from previous lecture slides
- Compute distance with Breadth-first search

# Extra Resources

- Introduction to Algorithms, 3<sup>rd</sup>, Chapter 22
- Algorithms Illuminated Part 2: Chapter 8

# General Algorithm

```
FUNCTION Connectivity(G, start_vertex)
  found = {v: FALSE FOR v IN G.vertices}
  found[start_vertex] = TRUE
```

**LOOP**

```
  (vFound, vNotFound) = get_valid_edge(G.edges, found)
```

```
  IF vFound == NONE || vNotFound == NONE
```

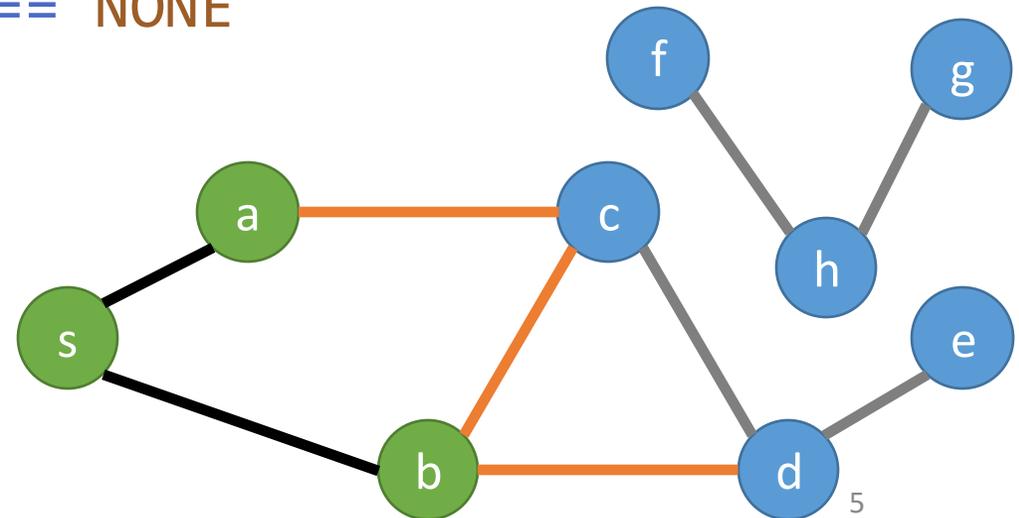
```
    BREAK
```

```
  ELSE
```

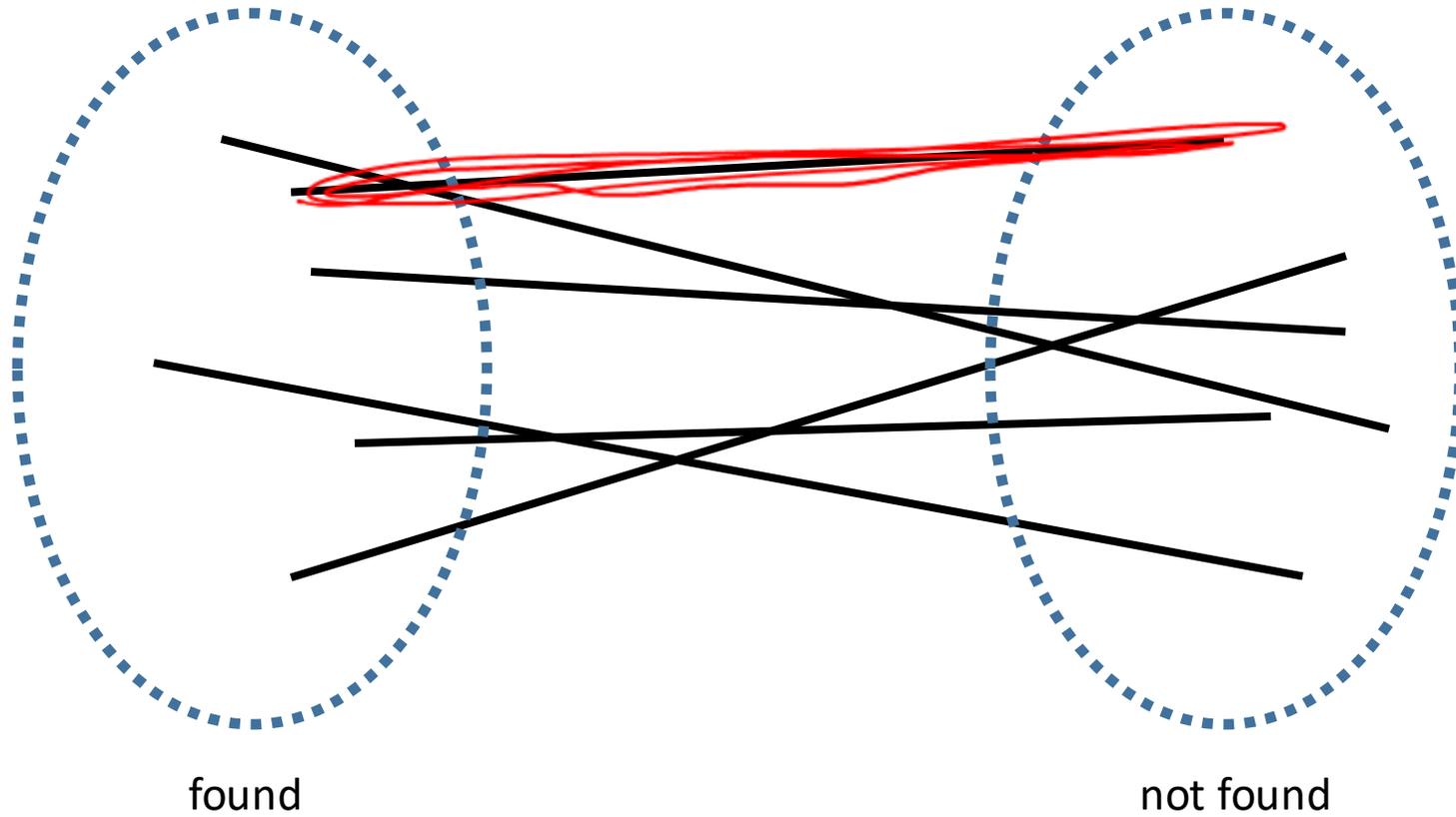
```
    found[vNotFound] = TRUE
```

```
  RETURN found
```

Find an edge where one vertex has been found and the other vertex has not been found.



How do we choose the next edge?



# Two common (and well studied) options

## Breadth-First Search

- Explore the graph in **layers**
- “*Cautious*” exploration
- Use a FIFO data structure (can you think of an example?)

## Depth-First Search

- Explore recursively
- A more “*aggressive*” exploration (we backtrack if necessary)
- Use a LIFO data structure (or recursion)

```
FUNCTION BFS(G, start_vertex)
```

```
found = {v: FALSE FOR v IN G.vertices}
```

```
found[start_vertex] = TRUE
```

```
visit_queue = [start_vertex]
```

```
WHILE visit_queue.length != 0
```

```
    vFound = visit_queue.pop() ↓
```

```
    FOR vOther IN G.edges[vFound]
```

```
        IF found[vOther] == FALSE ←
```

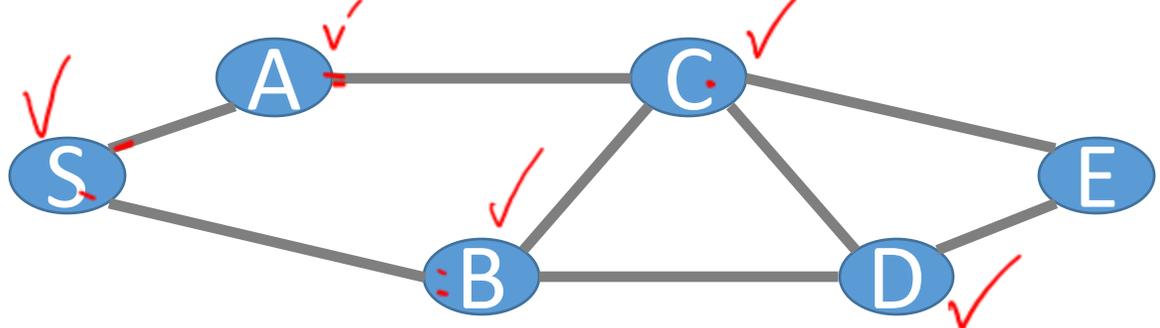
```
            found[vOther] = TRUE ←
```

```
            visit_queue.add(vOther) ←
```

```
RETURN found
```

```
FUNCTION Connectivity(G, start_vertex)
    found = {v: FALSE FOR v IN G.vertices}
    found[start_vertex] = TRUE
    LOOP
        (vFound, vNotFound) =
            get_valid_edge(G.edges, found)
        IF vFound == NONE || vNotFound == NONE
            BREAK
        ELSE
            found[vNotFound] = TRUE
    RETURN found
```

# Exercise questions 2 and 3



vFound = S  
 vO = A  
 vO = B  
 VF = A  
 vO = C  
 vO = S  
 VF = B  
 vO = C  
 vO = S

VF = C  
 vO = A  
 vO = B  
 vO = D  
 vO = E  
 VF = D

S A B C D E

```

FUNCTION BFS(G, start_vertex)
  found = {v: FALSE FOR v IN G.vertices}
  found[start_vertex] = TRUE
  visit_queue = [start_vertex]
  = [S A B C D E]
  WHILE visit_queue.length != 0
    → vFound = visit_queue.pop()
    FOR vOther IN G.edges[vFound]
      { IF found[vOther] == FALSE
        found[vOther] = TRUE
        visit_queue.add(vOther)
      }
  RETURN found
  
```

Given a tie, visit edges are in alphabetical order

# Running Time

```
FUNCTION BFS(G, start_vertex)
  found = {v: FALSE FOR v IN G.vertices}
  found[start_vertex] = TRUE
  visit_queue = [start_vertex]
```

```
  WHILE visit_queue.length != 0
    vFound = visit_queue.pop()
    FOR vOther IN G.edges[vFound]
      IF found[vOther] == FALSE
        found[vOther] = TRUE
        visit_queue.add(vOther)
```

```
  RETURN found
```

What is the running time?



How many times do we consider each edge?

twice

$$\cancel{O(m) \cdot O(n)}$$

$$O(n) + O(m) = O(n+m)$$

# Running Time

```
FUNCTION BFS(G, start_vertex)
  found = {v: FALSE FOR v IN G.vertices}
  found[start_vertex] = TRUE
  visit_queue = [start_vertex]

  WHILE visit_queue.length != 0
    vFound = visit_queue.pop()
    FOR vOther IN G.edges[vFound]
      IF found[vOther] == FALSE
        found[vOther] = TRUE
        visit_queue.add(vOther)

  RETURN found
```

What is the running time?

How many times do we consider each edge?

$$T_{BFS}(n, m) = O(n_s + m_s)$$

where  $n_s$  and  $m_s$  are the nodes and edges **findable/connected** from/to the start vertex

# Proof: BFS

**Claim:** BFS finds all nodes connected to the start node.

At the end of the BFS algorithm,  $v$  is marked found if there exists a path from  $s$  to  $v$

- Note: this is just a special case of the general algorithm that we proved by contradiction

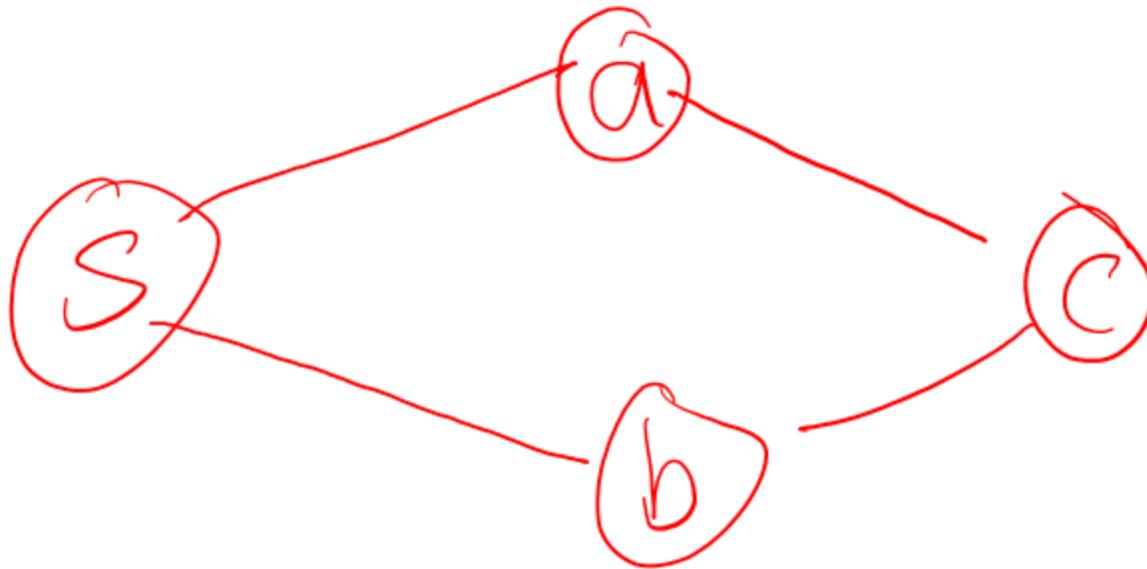
Practice for a loop invariant

Homework question

# Question

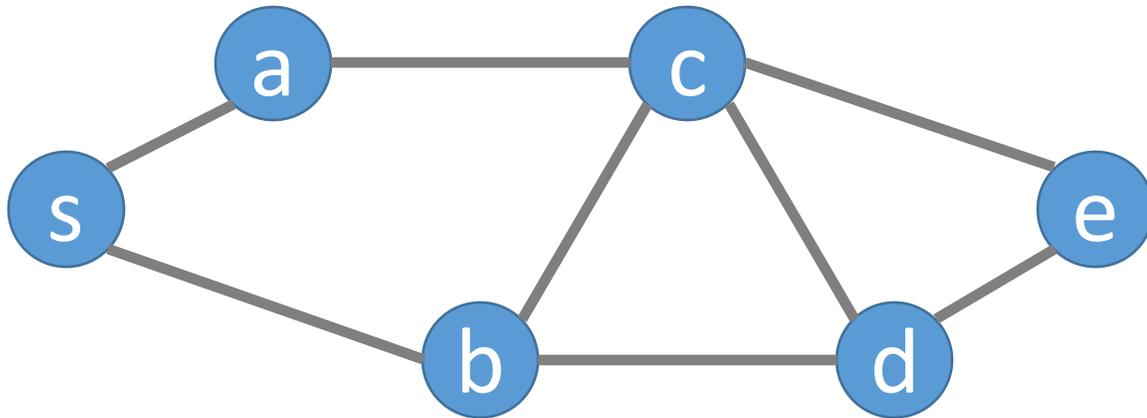
## The Shortest Path Problem

- How can we determine the fewest number of hops between the start vertex and all other connected vertices?



## BFS Exercise Question 1

How can we determine the fewest number of hops between the start vertex and all other connected vertices?



```
FUNCTION BFS(G, start_vertex)
    found = {v: FALSE FOR v IN G.vertices}
    found[start_vertex] = TRUE
    visit_queue = [start_vertex]

    WHILE visit_queue.length != 0
        vFound = visit_queue.pop()
        FOR vOther IN G.edges[vFound]
            IF found[vOther] == FALSE
                found[vOther] = TRUE
                visit_queue.add(vOther)

    RETURN found
```

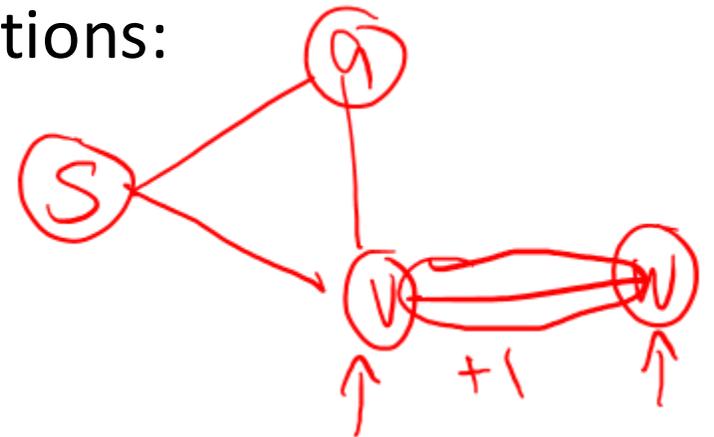
*Given a tie, visit edges are in alphabetical order*

# The Shortest Path Problem

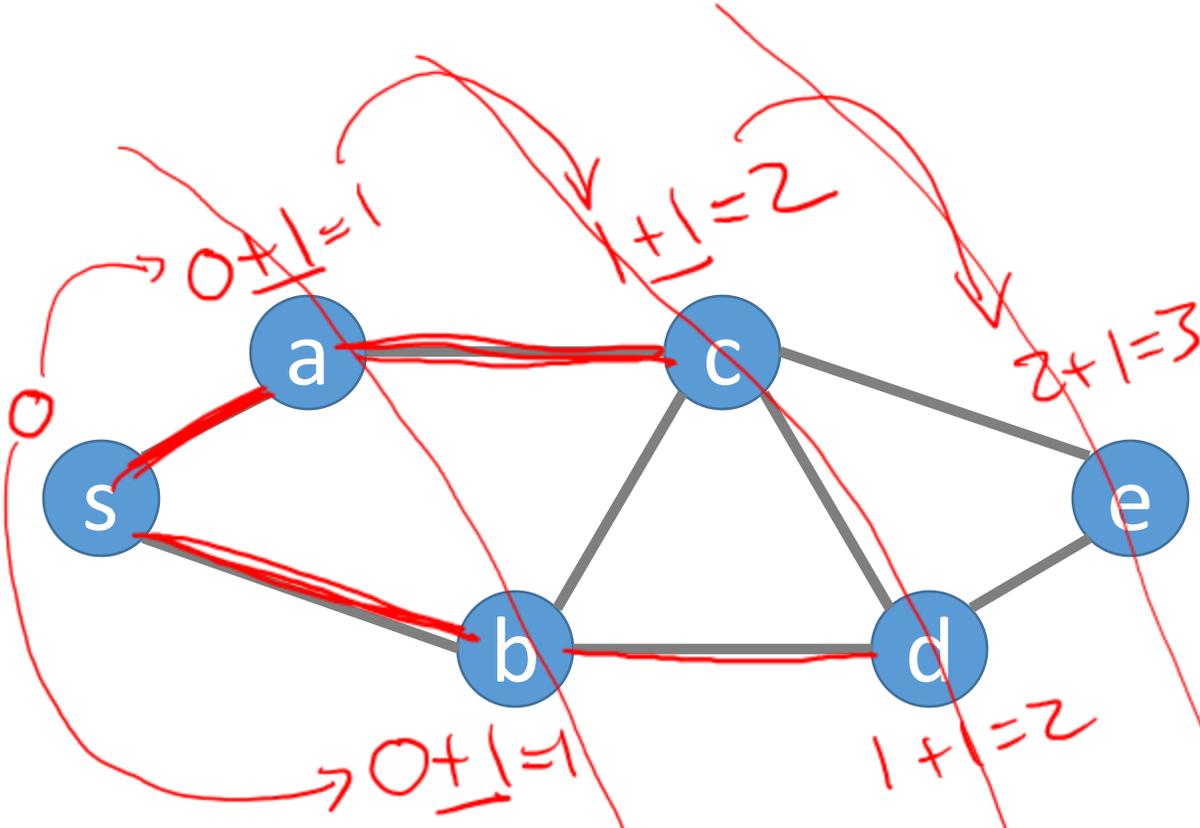
Determine the fewest number of hops between the start vertex and all other vertices

Same algorithm as before with the following additions:

- Initialize the distances[s] as 0
- Initialize all other distances to **infinity**
- When considering an edge (v, w)
  - If w is not found, then set  $\text{dist}(w)$  to  $\text{dist}(v) + 1$



# The Shortest Path Problem



After we terminate, distances[v] = "the layer that v is in"

```

FUNCTION DistanceBFS(G, start_vertex)
    found = {v: FALSE FOR v IN G.vertices}
    found[start_vertex] = TRUE

    distances = {v: INFINITY FOR v IN G.vertices}
    distances[start_vertex] = 0

    visit_queue = [start_vertex]
    WHILE visit_queue.length != 0
        vFound = visit_queue.pop()
        FOR vOther IN G.edges[vFound]
            IF found[vOther] == FALSE
                found[vOther] = TRUE
                visit_queue.add(vOther)
                distances[vOther] = distances[vFound] + 1

    RETURN distances
    
```

Given a tie, visit edges are in alphabetical order

# Connected Components

Let's only consider undirected graphs for now

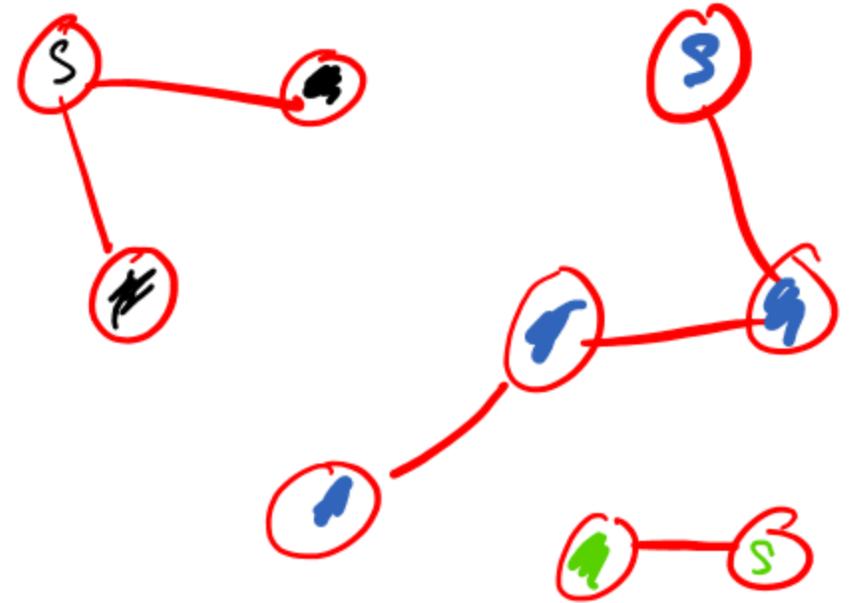
Let  $G = (V, E)$  be an undirected graph

Goal: compute all **connected components** in  $O(m + n)$

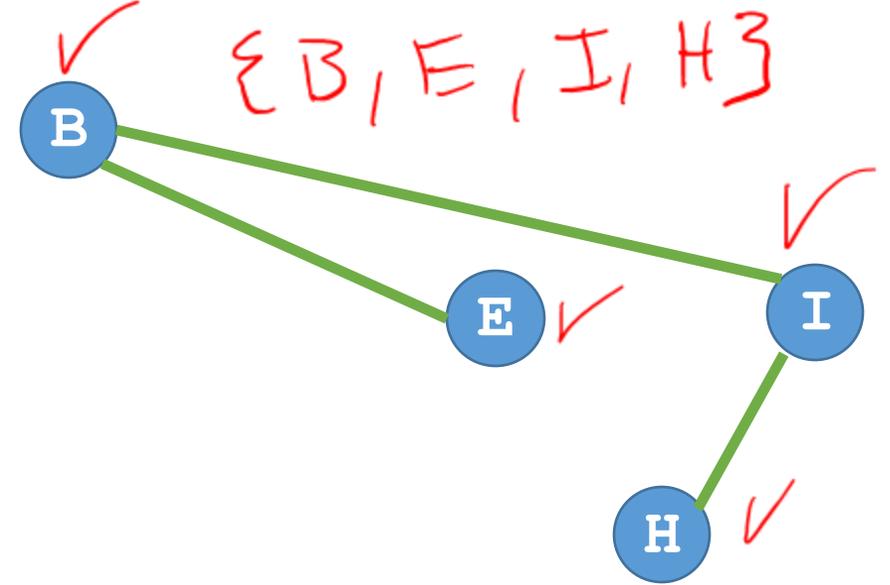
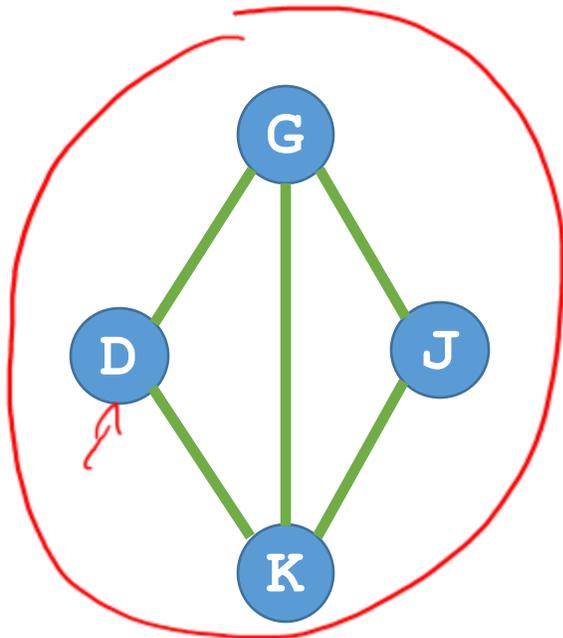
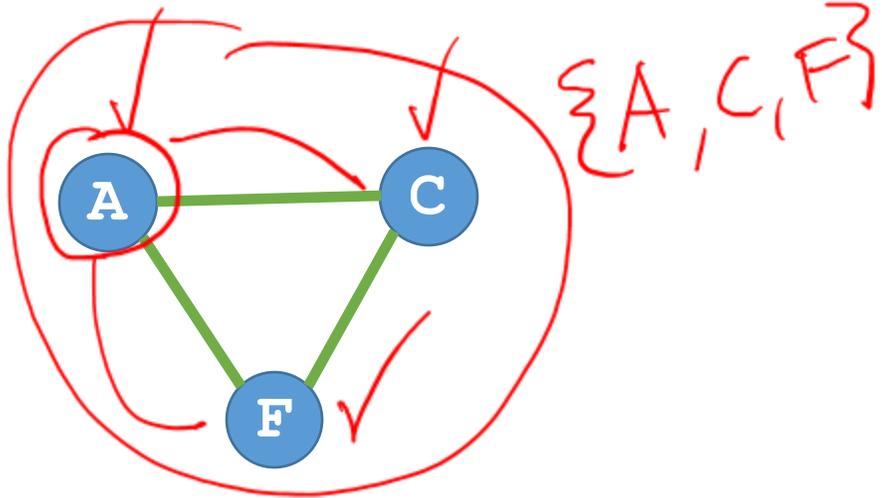
- A component is any group of vertices that can reach one another
- For example, if we are trying to see if a network has become disconnected

Exercise question 2:

How would you do this using our BFS procedure from before?



# BFS Exercise Question 2



```
FUNCTION FindComponents (G)
  components = []
  found = {v: FALSE FOR v IN G.vertices}
  FOR v IN G.vertices
    IF NOT found[v]
      newly_found = BFS(G, v)
      new_component = {
        w FOR w, w_is_found IN newly_found
        IF w_is_found
      }
      component.append(new_component)
      FOR w IN new_component:
        found[w] = TRUE
  RETURN components
```