

REVIEW

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CS 1.40 – Spring 2024

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Admin

Final

- posted on Gradescope
- due Wednesday (5/8) at 11:59pm (seniors: 5/2 at noon)
- time-limited (3 hours to take – 3.5 hours upload back to Gradescope)
- You may use:
 - the book
 - your notes
 - the class notes
 - the assignments
 - ONLY these things
- Do NOT discuss it with anyone until after Wednesday (5/8) at 11:59pm

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Test taking advice

- Read the questions carefully!
- Don't spend too much time on any problem
 - if you get stuck, move on and come back
- When you finish answering a question, reread the question and make sure that you answered everything the question asked
- Think about how you might be able to reuse an existing algorithm/approach
- Show your work (I can't give you partial credit if I can't figure out what went wrong)
- Don't rely on the book/notes for conceptual things
 - Do rely on the notes for a run-time you may not remember, etc.

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High-level approaches

Algorithm tools

- Divide and conquer
 - assume that we have a solver, but that can only solve sub-problems
 - define the current problem with respect to smaller problems
 - Key: sub-problems should be non-overlapping
- Dynamic programming
 - Same as above
 - Key difference: sub-problems are **overlapping**
 - Once you have this recursive relationship:
 - figure out the data structure to store sub-problem solutions
 - work from bottom up

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High-level approaches

Algorithm tools cont.

- Greedy
 - Same idea: most greedy problems can be solve using dynamic programming (but generally slower)
 - Key difference: Can decide between overlapping sub-problems without having to calculate them (i.e., we can make a local decision)
- Flow
 - Min-capacity cut
 - Bottleneck edge
 - Matching problems
 - Numerical maximization/minimization problems

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Data structures

A data structure

- Stores data
 - Supports access to/questions about data efficiently
 - the different bias towards different actions
 - No single best data structure
- Fast access/lookup?**
- If keys are sequential: array
 - If keys are non-sequential or non-numerical: hashtable
 - Guaranteed run-time/ordered: balanced binary search tree

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Data structures

Min/max?

- heap
- binomial heap

Fast insert/delete at positions?

- linked list

Others

- stacks/queues
- extensible data structures
- balanced BSTs
- disjoint sets

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Graphs

Graph types

- directed/undirected
- weighted/unweighted
- trees, DAGs
- cyclic
- connected

Algorithms

- connectedness
- contains a cycle
- traversal
 - dfs
 - bfs

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Graphs

Graph algorithms cont.

- ▣ minimum spanning trees (Prim's, Kruskal's)
- ▣ shortest paths
 - single source (BFS, Dijkstra's, Bellman-Ford)
 - all pairs (Johnson's, Floyd-Warshall)
- ▣ topological sort
- ▣ flow

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Other topics...

Analysis tools

- ▣ recurrences (master method, recurrence trees)
- ▣ big-O
- ▣ amortized analysis

NP-completeness

- ▣ proving NP-completeness
- ▣ reductions

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Proofs: general

Be clear and concise

Make sure you state assumptions and justify each step

Make sure when you're done you've shown what you need to show

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Proof by induction

1. State what you're trying to prove
We show that XXX using proof by induction
2. Prove base case
3. State the inductive hypothesis
4. Inductive proof
 - a. State what you want to show (may include a variable change, e.g., k in instead of n)
 - b. Show a step by step derivation from the left hand side resulting in the right hand side. Give justifications for steps that are non-trivial

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Prof by induction: structural

Use induction to prove that the number of degree-2 nodes in a non-empty binary tree is 1 less than the number of leaves. Recall that the degree of a vertex in a tree is the number of children that it has.

1. State what you're trying to prove
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Other (important) places we saw proofs

Recurrences (substitution method)

Big O (needed find constants c n_0)

Greedy algorithm correctness (proof by contradiction or stays ahead—induction —)

Proof of algorithm correctness (MSTs, Flow)

NP-completeness (proving correctness of reductions)

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Recurrences

Three ways to solve:

- Substitution
- Recurrence tree (may still have to use substitution to verify)
- Master method

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Recurrences

$$T(n) = 2T(n/3) + d$$

$$T(n) = aT(n/b) + f(n)$$

if $f(n) = O(n^{\log_b a - \epsilon})$ for $\epsilon > 0$, then $T(n) = \Theta(n^{\log_b a})$

if $f(n) = \Theta(n^{\log_b a})$, then $T(n) = \Theta(n^{\log_b a} \log n)$

if $f(n) = \Omega(n^{\log_b a + \epsilon})$ for $\epsilon > 0$ and $af(n/b) \leq cf(n)$ for $c < 1$
then $T(n) = \Theta(f(n))$

$$T(n) = T(n-1) + \log n$$

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Dynamic programming

Method for solving problems where optimal solutions can be defined in terms of optimal solutions to subproblems

AND

the subproblems are overlapping

Local decisions result in *different* subproblems. Not obvious how to make the first choice.

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DP advice

Write the recursive definition

- What is the input/output to the problem?
- What would a solution look like? What are the options for picking the first component of a solution?
- Assume you have a solver for subproblems. How can you combine the first decision with answer to subproblem.

Define DP structure: what are subproblems indexed by?

State how to fill in the table (including base cases and where the answer is)

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