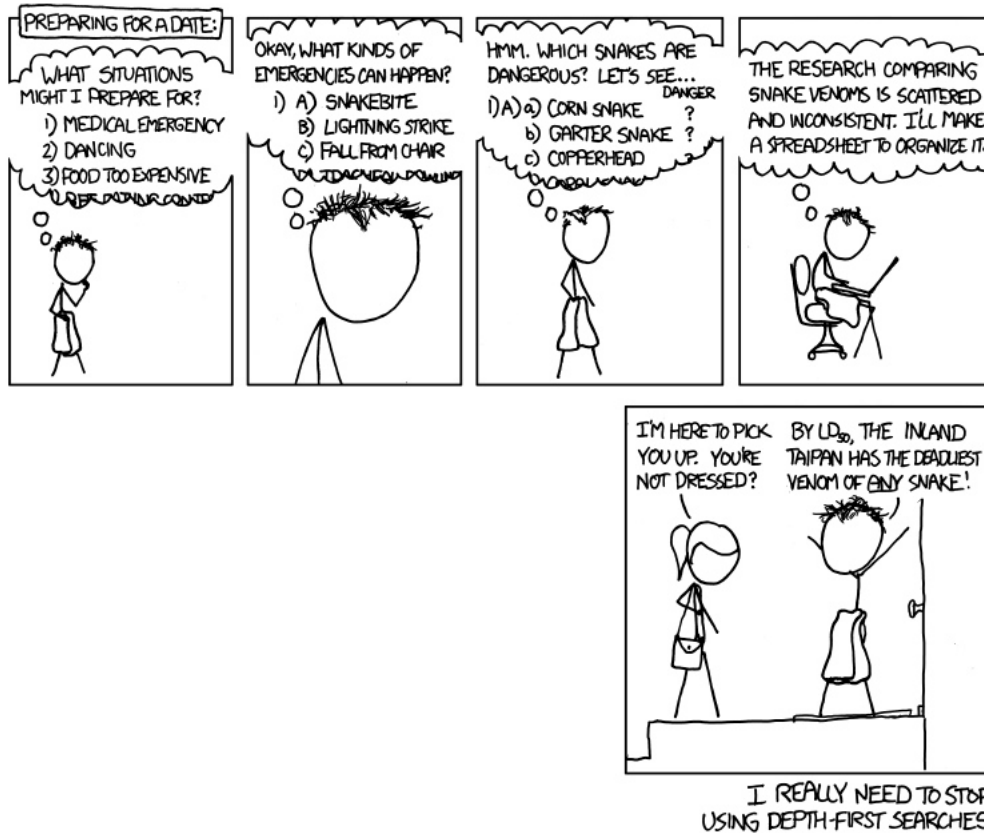


# CS140 - Assignment 8

Due: Sunday, Mar. 31st at 11:59pm



<http://xkcd.com/761/>

Notes:

- Many of the algorithms below can be accomplished by either modifying the graph and applying a known algorithm or slightly modifying a known algorithm. Try thinking of these *first* as they will save you a lot of work, and writing :) I don't expect long answers, but be precise.
- You will be graded on efficiency!
- If not specified in the problem, you may assume whatever graph representation makes your algorithm more efficient (adjacency list or adjacency matrix). State which one you are using.

1. [8 points] State whether the following statements about a graph  $G$  that is undirected and connected are true or false and justify your answer.
  - (a) Prim's algorithm works correctly if  $G$  has negative edge weights.
  - (b) The shortest path between two nodes is always part of some MST.
2. [5 points] Write pseudocode for an algorithm which, given an undirected graph  $G$  and a particular edge  $e$  in it, determines whether  $G$  has a cycle containing  $e$ . What is the runtime of this algorithm?
3. [8 points] Often there are multiple shortest paths between nodes of a graph. Write pseudocode for an algorithm that given an undirected, unweighted graph  $G$  and nodes  $u, v \in V$ , outputs the number of distinct shortest paths from  $u$  to  $v$ . What is the running time?
4. [5 points] Given a directed graph  $G = (V, E)$  with positive edge weights and a particular node  $v_i \in V$ , give an efficient algorithm for finding the shortest paths between **all pairs of nodes**, with the one restriction that these paths must all pass through  $v_i$ . Give the runtime of your algorithm. Points will be deducted for an inefficient algorithm.

*Hints:*

- Any path in this problem can be seen as two parts, the part to  $v_i$  and the part from  $v_i$ .
  - Look at how we determined if a graph was strongly connected.
5. [5 points] If a graph does not have a negative cycle, when calculating the shortest paths from a given vertex using the Bellman-Ford algorithm, we can stop early and do not need to do all  $|V| - 1$  iterations and will still have a correct answer for all the shortest paths from that vertex. Describe how to modify the Bellman-Ford algorithm to stop early when all of the distances are already correct.
  6. [6 points] Given an undirected graph  $G$  with nonnegative edge weights  $w_e \geq 0$ . Suppose you have calculated the minimum spanning tree of  $G$  and also the shortest paths to all nodes from a particular node  $s \in V$ . Now, suppose that each edge weight is increased by 1, i.e. the new weights are  $w'_e = w_e + 1$ .
    - (a) (3 points) Does the minimum spanning tree change? Give an example where it does or prove that it cannot change.
    - (b) (3 points) Do the shortest paths from  $s$  change? Given an example where it does or prove that it cannot change.