GREEDY ALGORITHMS

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CS 140 – Fall 2022

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Admin

Assignment 7

Learning communities: may attend any group going forward

Grading update
- 5 graded
- 4.1 soon!
- Checkpoint revisions graded

2

Checkpoint 2

2 pages of notes

From hashtables (9/2) through dynamic programming (10/13)

Practice problems available in group assignment 7

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A problem

Input: a number k

Output: $\{n_p, n_n, n_d, n_q\}$, where $n_p+5n_n+10n_d+25n_q=k$ and $n_p+n_n+n_d+n_q$ is minimized

What is this problem?
How would you state it in English?

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Making change!

Input: a number k

Output: $\{n_p, n_n, n_d, n_q\}$, where $n_p+5n_n+10n_d+25n_q=k$ and $n_p+n_n+n_d+n_q$ is minimized

Provide (U.S.) coins that sum up to *k* such that we minimize the number of coins

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Making change!

Input: a number k

Output: $\{n_p, n_n, n_d, n_q\}$, where $n_p+5n_n+10n_d+25n_q=k$ and $n_p+n_n+n_d+n_q$ is minimized

Algorithm to solve it?

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Making change!

Input: a number k

Output: $\{n_p, n_n, n_d, n_q\}$, where $n_p+5n_n+10n_d+25n_q=k$ and $n_p+n_n+n_d+n_q$ is minimized

$n_q = \lfloor k \,/\, 25 \rfloor$ pick as many quarters as we can

Solve:
$n_p + 5nn + 10nd = \mathrm{k} - 25\lfloor k \,/\, 25 \rfloor$ recurse

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Algorithms vs heuristics

What is the difference between an algorithm and a heuristic?

Algorithm: a set of steps for arriving at the *correct* solution

Heuristic: a set of steps that will arrive at some solution

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Making change!

$n_q = \lfloor k / 25 \rfloor$ pick as many quarters as we can

Solve:
$n_p + 5nn + 10nd = \mathrm{k} - 25\lfloor k / 25 \rfloor$ recurse

Algorithm or heuristic?

Need to prove its correct!

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Greedy algorithms

What is a greedy algorithm?

Algorithm that makes a local decision with the goal of creating a globally optimal solution

Method for solving problems where optimal solutions can be defined in terms of optimal solutions to sub-problems

What does this mean? Where have we seen this before?

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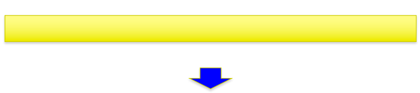
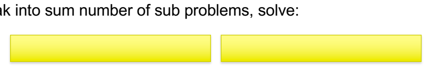
Greedy vs. divide and conquer

Divide and conquer

To solve the general problem:



Break into sum number of sub problems, solve:



then possibly do a little work

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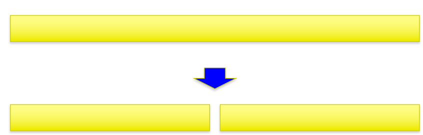
Divide and conquer

Divide and conquer

To solve the general problem:



The solution to the general problem is solved with respect to solutions to sub-problems!

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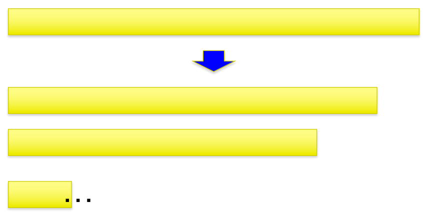
Dynamic programming

Dynamic programming

To solve the general problem:



The solution to the general problem is solved with respect to solutions to sub-problems **that overlap**!

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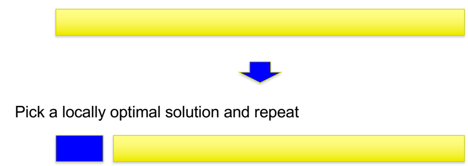
Greedy vs. divide and conquer

Greedy

To solve the general problem:



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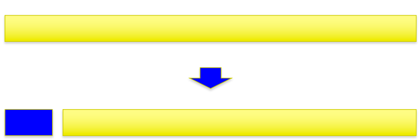
Greedy vs. divide and conquer

Greedy

To solve the general problem:



The solution to the general problem is solved with respect to solutions to sub-problems!

Slightly different than divide and conquer

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Proving greedy algorithms correct

One approach, prove:

1) Optimal substructure: The optimal solution contains within it the optimal solution to subproblems

2) Greedy choice property: The greedy choice is contained within some optimal solution

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Making change!

$n_q = \lfloor k / 25 \rfloor$ pick as many quarters as we can

Solve:
$n_p + 5nn + 10nd = \text{k} - 25\lfloor k / 25 \rfloor$ recurse

$\{c_1, c_2, c_3, \ldots, cm\}$ solution: individual coins selected

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Optimal substructure

If $\{c_1, c_2, c_3, \ldots, cm\}$ is optimal for k, then $\{ c_2, c_3, \ldots, cm\}$ is optimal for k-c_1

We can combine a greedy choice with the optimal solution for the remaining problem and get a solution to the general problem

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Optimal substructure

Proof by contradiction:

Assume $\{c_1, c_2, c_3, \ldots, cm\}$ is optimal for k, *but* $\{ c_2, c_3, \ldots, cm\}$ is not optimal for k-c_1

What does that imply?

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Optimal substructure

Proof by contradiction:

Assume $\{c_1, c_2, c_3, \ldots, cm\}$ is optimal for k, *but* $\{ c_2, c_3, \ldots, cm\}$ is not optimal for k-c_1

There is some other set of coins $\{ c'2, c'3, \ldots, c'_n\}$ where n < m that add up to k-c_1

Any problem contradiction?

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Optimal substructure

Proof by contradiction:

Assume $\{c_1, c_2, c_3, \dots, cm\}$ is optimal for k,
but $\{c_2, c_3, \dots, cm\}$ is not optimal for k-c_1

There is some other set of coins $\{c'2, c'3, \dots, c'_n\}$ where n < m that add up to k-c_1

$\{c_1, c'2, c'3, \dots, c'_n\}$ would be a solution, but since n < m this implies that our original solution wasn't optimal!

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Optimal substructure

If $\{c_1, c_2, c_3, \dots, cm\}$ is optimal for

$\{c_2, c_3, \dots, cm\}$ is optimal for k-c_1

We can make greedy decisions

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Greedy choice property

Greedy choice property: The greedy choice is contained within some optimal solution

The greedy choice results in an optimal solution

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Greedy choice property

Proof by contradiction:

Let $\{c_1, c_2, c_3, \dots, cm\}$ be an optimal solution

Assume it is ordered from largest to smallest

Assume that it does not make the greedy choice at some coin c_i

Any problem contradiction?

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Greedy choice property

Proof by contradiction:

Let $\{c_1, c_2, c_3, \ldots, cm\}$ be an optimal solution
Assume it is ordered from largest to smallest

Assume that it does not make the greedy choice at some coin c_i

$g_i > c_i$. We need at least one more lower denomination coin because g_i can be made up of c_i and one or more of the other denominations

but that would mean that the solution is longer than the greedy!

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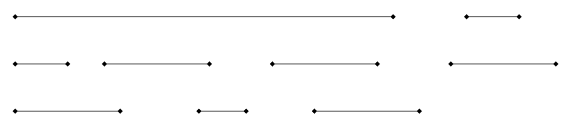
Interval scheduling

Given *n* activities A = [a_1,a_2, .., a_n] where each activity has start time s_i and a finish time f_i. Schedule as many as possible of these activities such that they don't conflict.



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Interval scheduling

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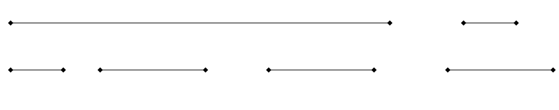

Which activities conflict?

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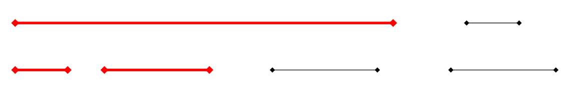
Interval scheduling

Given *n* activities A = [a_1,a_2, .., a_n] where each activity has start time s_i and a finish time f_i. Schedule as many as possible of these activities such that they don't conflict.



Which activities conflict?

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Simple recursive solution

Enumerate all possible solutions and find which schedules the most activities

```
IntervalSchedule-Recursive(A)
if A = {}
        return 0
else
        max = −∞
        for all a ∈ A
                A′ ← A minus a and all conflicting activites with a
                s = IntervalSchedule-Recursive(A′)
                if s > max
                        max = s
        return 1 + max
```

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Simple recursive solution

Is it correct?

- max{all possible solutions}

Running time?

- O(n!)

```
IntervalSchedule-Recursive(A)
if A = {}
        return 0
else
        max = −∞
        for all a ∈ A
                A′ ← A minus a and all conflicting activites with a
                s = IntervalSchedule-Recursive(A′)
                if s > max
                        max = s
        return 1 + max
```

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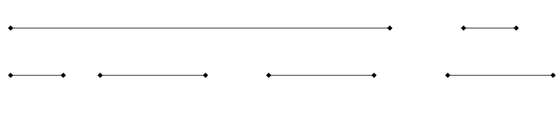
Can we do better?

Dynamic programming

- $O(n^2)$

Greedy solution – Is there a way to repeatedly make local decisions?

- Key: we'd still like to end up with the *optimal* solution



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Overview of a greedy approach

Greedily pick an activity to schedule

Add that activity to the answer

Remove that activity and all conflicting activities. Call this A'.

Repeat on A' until A' is empty

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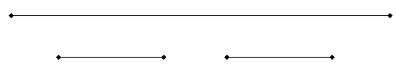
Greedy options

Select the activity that starts the earliest, i.e. argmin{s_1, s_2, s_3, …, s_n}?



33

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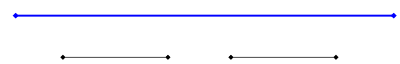

non-optimal

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Greedy options

Select the shortest activity, i.e. argmin{f_1-s_1, f_2-s_2, f_3-s_3, …, f_n-s_n}

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Greedy options

Select the shortest activity, i.e. argmin{f_1-s_1, f_2-s_2, f_3-s_3, …, f_n-s_n}

non-optimal

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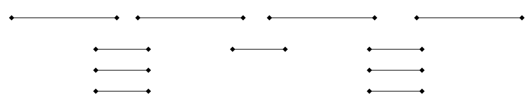
Greedy options

Select the activity with the smallest number of conflicts



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Greedy options

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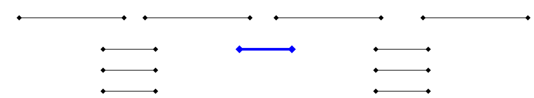

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Greedy options

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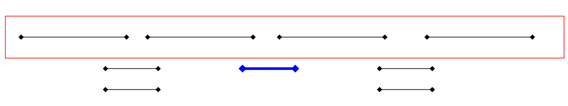

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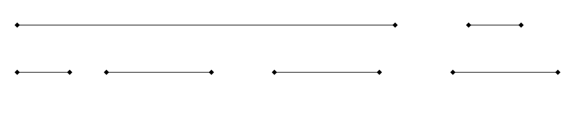
Greedy options

Select the activity that ends the earliest, i.e. argmin{f_1, f_2, f_3, …, f_n}?



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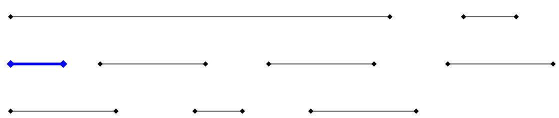
Greedy options

Select the activity that ends the earliest, i.e. argmin$\{f_1, f_2, f_3, \ldots, f_n\}$?



remove the conflicts

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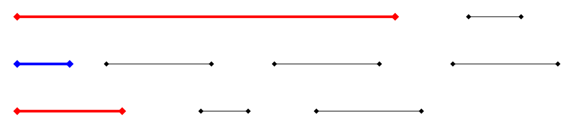
Greedy options

Select the activity that ends the earliest, i.e. argmin$\{f_1, f_2, f_3, \ldots, f_n\}$?



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Greedy options

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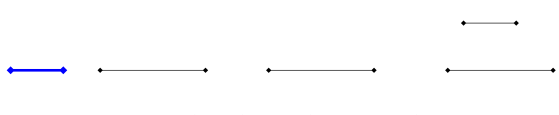

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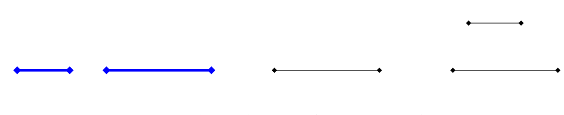

remove the conflicts

44

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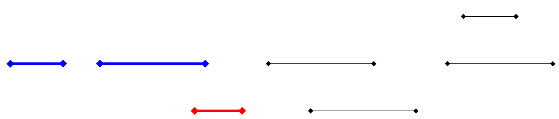

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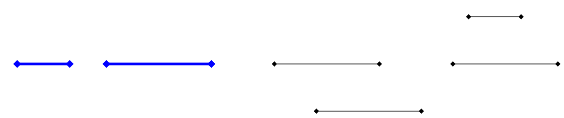
Greedy options

Select the activity that ends the earliest, i.e. argmin$\{f_1, f_2, f_3, \ldots, f_n\}$?



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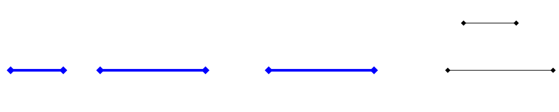
Greedy options

Select the activity that ends the earliest, i.e. argmin$\{f_1, f_2, f_3, \ldots, f_n\}$?



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Greedy options

Select the activity that ends the earliest, i.e. argmin$\{f_1, f_2, f_3, \ldots, f_n\}$?

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Greedy options

Select the activity that ends the earliest, i.e. argmin$\{f_1, f_2, f_3, \ldots, f_n\}$?

Multiple optimal solutions

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Greedy options

Select the activity that ends the earliest, i.e. argmin$\{f_1, f_2, f_3, \ldots, f_n\}$?

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Greedy options

Select the activity that ends the earliest, i.e. argmin$\{f_1, f_2, f_3, \ldots, f_n\}$?

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Efficient greedy algorithm

Once you've identified a reasonable greedy heuristic:

- Prove that it always gives the correct answer
- Develop an efficient solution

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Is our greedy approach correct?

"Stays ahead" argument:

show that no matter what other solution someone provides you, the solution provided by your algorithm always "stays ahead", in that no other choice could do better

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Is our greedy approach correct?

"Stays ahead" argument

Let $r_1, r_2, r_3, \ldots, r_k$ be the solution found by our approach

Let $o_1, o_2, o_3, \ldots, o_k$ be another optimal solution

Show our approach "stays ahead" of any other solution

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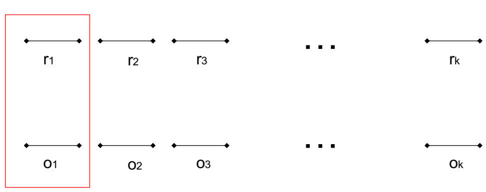
Stays ahead



Compare first activities of each solution

what do we know?

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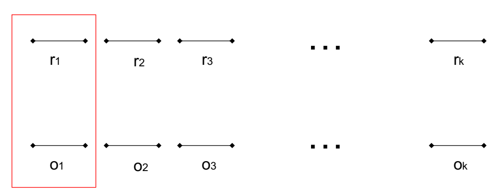
Stays ahead



$\text{finish}(r_1) \leq \text{finish}(o_1)$

what does this imply?

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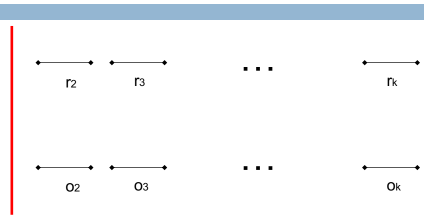
Stays ahead



We have **at least** as much time as any other solution to schedule the remaining 2…k tasks

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An efficient solution

```
IntervalSchedule-Greedy(A)
sort A based on finish times f_i
for i ← 1 to n
        add a_i to R
        finish ← f_i
        while s_i < finish
                i ← i + 1
return R
```

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Running time?

```
IntervalSchedule-Greedy(A)
sort A based on finish times f_i        Θ(n log n)
for i ← 1 to n                          ⎫
        add a_i to R                    ⎪
        finish ← f_i                    ⎬ Θ(n)
        while s_i < finish              ⎪
                i ← i + 1               ⎭
return R
```

Overall: Θ(n log n)

Better than:

$O(n!)$
$O(n^2)$

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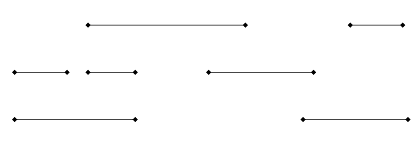
Scheduling *all* intervals

Given *n* activities, we need to schedule **all** activities.
Goal: minimize the number of resources required.



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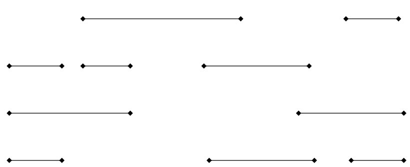
Greedy approach?

The best we could ever do is the maximum number of conflicts for any time period



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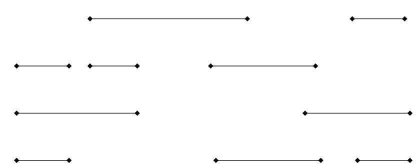
Calculating max conflicts efficiently



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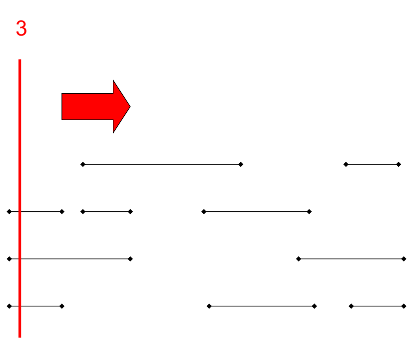
Calculating max conflicts efficiently



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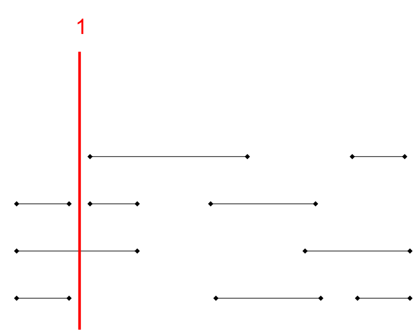
Calculating max conflicts efficiently



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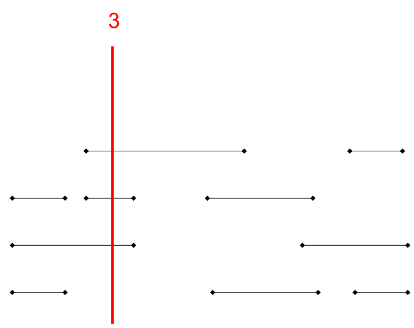

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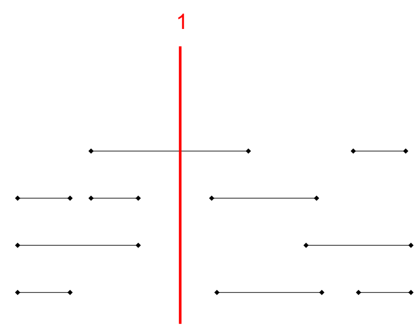

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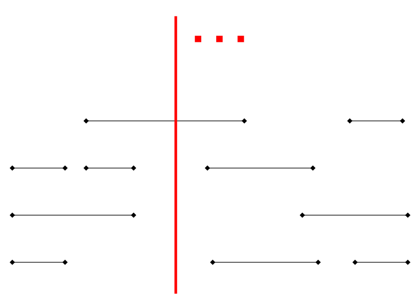
Calculating max conflicts efficiently



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Calculating max conflicts

```
AllIntervalScheduleCount(A)
Sort the start and end times, call this X
current ← 0
max ← 0
for i ← 1 to length[X]
        if x_i is a start node
                current++
        else
                current--
        if current > max
                max ← current
return max
```

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Correctness?

We can do no better then the max number of conflicts. This exactly counts the max number of conflicts.

```
AllIntervalScheduleCount(A)
Sort the start and end times, call this X
current ← 0
max ← 0
for i ← 1 to length[X]
        if x_i is a start node
                current++
        else
                current--
        if current > max
                max ← current
return max
```

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Runtime?

O(2n log 2n + n) = O(n log n)

```
AllIntervalScheduleCount(A)
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for i ← 1 to length[X]
        if x_i is a start node
                current++
        else
                current--
        if current > max
                max ← current
return max
```

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Knapsack problems: Greedy or not?

0-1 Knapsack – A thief robbing a store finds n items worth v_1, v_2, .., v_n dollars and weight w_1, w_2, …, w_n pounds, where v_i and w_i are integers. The thief can carry at most W pounds in the knapsack. Which items should the thief take if he wants to maximize value.

Fractional knapsack problem – Same as above, but the thief happens to be at the bulk section of the store and can carry fractional portions of the items. For example, the thief could take 20% of item i for a weight of $0.2w_i$ and a value of $0.2v_i$.

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