Breadth First Search

https://cs.pomona.edu/classes/cs140/

Outline

Topics and Learning Objectives

• Discuss breadth first search for graphs

Exercises

- Continued from previous lecture slides
- Compute distance with Breadth-first search

Extra Resources

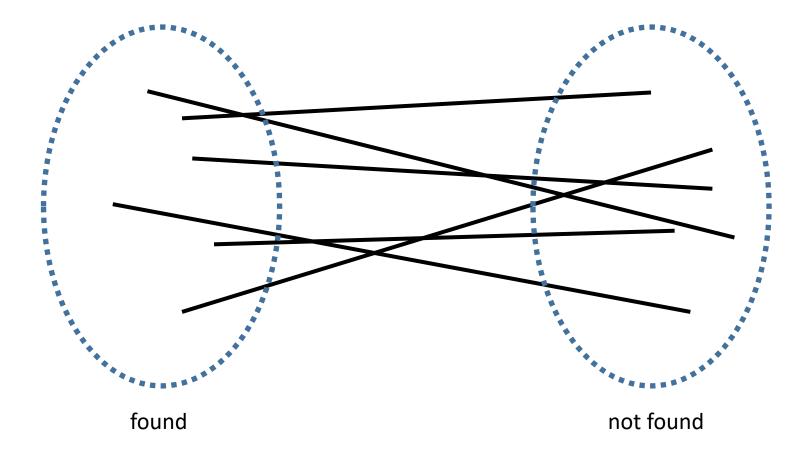
• Introduction to Algorithms, 3rd, Chapter 22

General Algorithm

```
has been found and the other
FUNCTION Connectivity(G, start vertex)
                                               vertex has not been found.
  found = {v: FALSE FOR v IN G.vertices}
  found[start_vertex] = TRUE
  LOOP
      (vFound, vNotFound) = get_valid_edge(G.edges, found)
      IF vFound == NONE | vNotFound == NONE
                                                                         g
         BREAK
      ELSE
         found[vNotFound] = TRUE
                                                                    h
  RETURN found
                                                       n
```

Find an edge where one vertex

How do we choose the <u>next</u> edge?



Two common (and well studied) options

Breadth-First Search

- Explore the graph in layers
- "Cautious" exploration
- Use a FIFO data structure (can you think of an example?)

Depth-First Search

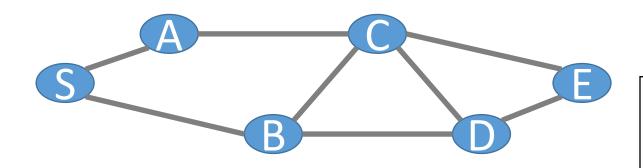
- Explore recursively
- A more "aggressive" exploration (we backtrack if necessary)
- Use a LIFO data structure (or recursion)

```
FUNCTION BFS(G, start_vertex)
found = {v: FALSE FOR v IN G.vertices}
found[start_vertex] = TRUE
visit_queue = [start_vertex]
```

```
WHILE visit_queue.length != 0
vFound = visit_queue.pop()
FOR vOther IN G.edges[vFound]
IF found[vOther] == FALSE
found[vOther] = TRUE
visit_queue.add(vOther)
```

```
FUNCTION Connectivity(G, start vertex)
  found = {v: FALSE FOR v IN G.vertices}
  found[start vertex] = TRUE
  LOOP
     (vFound, vNotFound) =
         get valid edge(G.edges, found)
     IF vFound == NONE | vNotFound == NONE
         BREAK
     ELSE
        found[vNotFound] = TRUE
  RETURN found
```

RETURN found



Exercise questions 2 and 3

FUNCTION BFS(G, start_vertex)

```
found = {v: FALSE FOR v IN G.vertices}
```

```
found[start_vertex] = TRUE
visit_queue = [start_vertex]
```

WHILE visit_queue.length != 0
vFound = visit_queue.pop()
FOR vOther IN G.edges[vFound]
IF found[vOther] == FALSE
found[vOther] = TRUE
visit_queue.add(vOther)

RETURN found

Given a tie, visit edges are in alphabetical order

Running Time

What is the running time?

```
FUNCTION BFS(G, start_vertex)
found = {v: FALSE FOR v IN G.vertices}
found[start_vertex] = TRUE
visit_queue = [start_vertex]
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WHILE visit_queue.length != 0
vFound = visit_queue.pop()
FOR vOther IN G.edges[vFound]
IF found[vOther] == FALSE
found[vOther] = TRUE
visit_queue.add(vOther)
```

How many times to we consider each edge?

RETURN found

Running Time

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FUNCTION BFS(G, start_vertex)
found = {v: FALSE FOR v IN G.vertices}
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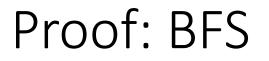
```
WHILE visit_queue.length != 0
vFound = visit_queue.pop()
FOR vOther IN G.edges[vFound]
IF found[vOther] == FALSE
found[vOther] = TRUE
visit_queue.add(vOther)
```

RETURN found

How many times to we consider each edge?

$$T_{BFS}(n,m) = O(n_s + m_s)$$

where n_s and m_s are the nodes and edges **findable/connected** from/to the start vertex



Claim: BFS finds all nodes connected to the start node.

At the end of the BFS algorithm, v is marked found if there exists a path from s to v

 Note: this is just a special case of the general algorithm that we proved by contradiction

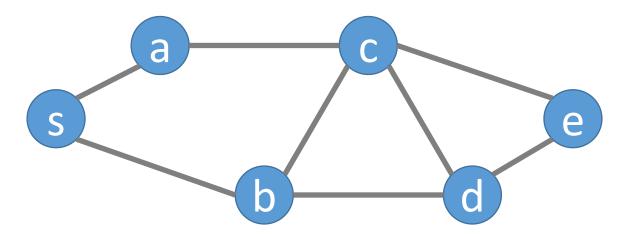


The Shortest Path Problem

• How can we determine the fewest number of hops between the start vertex and all other connected vertices?

BFS Exercise Question 1

How can we determine the fewest number of hops between the start vertex and all other connected vertices?



FUNCTION BFS(G, start_vertex) found = {v: FALSE FOR v IN G.vertices} found[start_vertex] = TRUE visit_queue = [start_vertex] WHILE visit queue.length != 0 vFound = visit_queue.pop() **FOR** vOther **IN** G.edges vFound **IF** found[vOther] == FALSE found[vOther] = TRUE visit_queue.add(v0ther)

RETURN found

Given a tie, visit edges are in alphabetical order

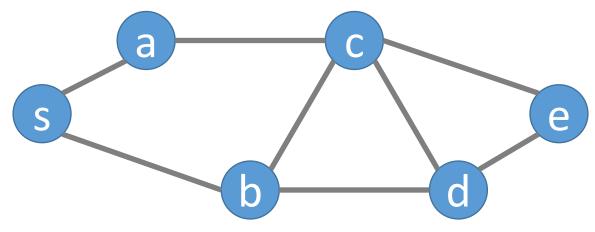
The Shortest Path Problem

Determine the fewest number of hops between the start vertex and all other vertices

Same algorithm as before with the following additions:

- Initialize the distances[s] as 0
- Initialize all other distances to **infinity**
- When considering an edge (v, w)
 - If w is not found, then set dist(w) to dist(v) + 1

The Shortest Path Problem



After we terminate, distances[v] = "the layer that v is in"

FUNCTION DistanceBFS(G, start_vertex) found = {v: FALSE FOR v IN G.vertices} found[start_vertex] = TRUE

```
distances = {v: INFINITY FOR v IN G.vertices}
distances[start_vertex] = 0
```

```
visit_queue = [start_vertex]
WHILE visit_queue.length != 0
vFound = visit_queue.pop()
FOR vOther IN G.edges[vFound]
IF found[vOther] == FALSE
found[vOther] = TRUE
visit_queue.add(vOther)
distances[vOther] = distances[vFound] + 1
```

```
RETURN distances
```

Given a tie, visit edges are in alphabetical order

Connected Components

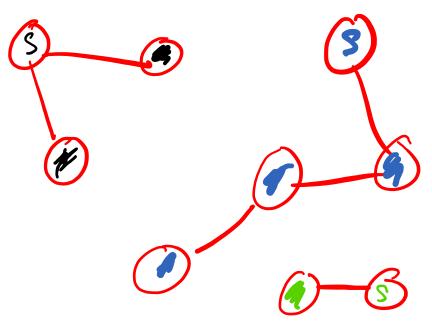
Let's only consider undirected graphs for now

Let G = (V,E) be an undirected graph

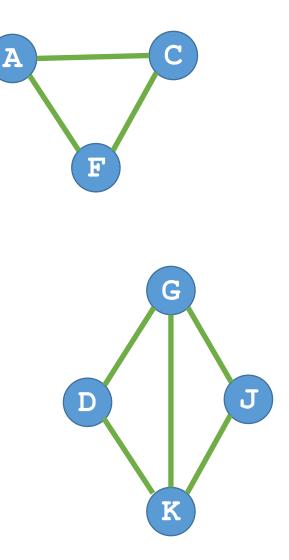
Goal: compute all connected components in O(m + n)

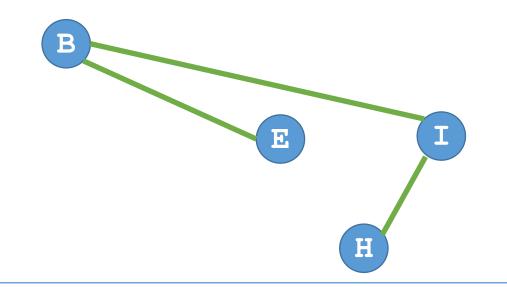
- A component is any group of vertices that can reach one another
- For example, if we are trying to see if a network has become disconnected

Exercise question 2: How would you do this using our BFS procedure from before?



BFS Exercise Question 2





```
FUNCTION FindComponents(G)
   components = []
   found = {v: FALSE FOR v IN G.vertices}
   FOR v IN G.vertices
      IF NOT found[v]
         newly found = BFS(G, v)
         new component = {
            w FOR w, w is found IN newly found
            IF w is found
         component.append(new component)
         FOR w IN new component:
            found[w] = TRUE
   RETURN components
```