# Loop Invariants

https://cs.pomona.edu/classes/cs140/

## Outline

**Topics and Learning Objectives** 

• Practice writing loop invariants

### Exercise

• Loop Invariant

## Extra Materials

- Chapter 2 of Introduction to Algorithms, Third Edition
- <u>https://www.win.tue.nl/~kbuchin/teaching/JBP030/notebooks/loop-invariants.html</u>

## Loop Invariant Proofs

- A procedural way to prove the correctness of some code with a loop
- Very similar to inductive proofs for recursive algorithms

```
FUNCTION SumArray(array)
sum = 0
i = 0
WHILE i < array.length
 sum = sum + array[i]
 i = i + 1</pre>
```

How do we prove that this code sums all values in the given array?

Some useful syntax:

- array[start ..= end] is the subarray
  - Including array [start], array [end], and everything in between
  - <u>Inclusive</u> lower and upper bounds
- array[start ... < end] is the subarray</li>
  - Including array [start], excluding array [end], and including everything in between
  - Inclusive lower bound, exclusive upper bound

## Loop Invariants

A loop invariant is a <u>predicate</u> (a statement that is either true or false) with the following properties:

1. It is true upon entering the loop the first time.

Initialization

2. If it is true upon starting an iteration of the loop, it remains true upon starting the next iteration. Maintenance

3. The loop terminates, and the loop invariant plus the reason that the loop terminates gives you the property that you want.

## Relation to Induction Proofs

Loop Invariant

- True before entering first loop
- True after executing any iteration
- True when loop terminates

#### Induction

- Prove for the base case
- State/assume inductive hypothesis
- Prove inductive step

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- Prove for the base case
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# How to perform a proof by loop invariant

- 1. State the loop invariant
  - 1. A statement that can be easily proven true or false
  - 2. The statement must reference the purpose of the loop
  - 3. The statement must reference variables that change each iteration

Initialization

2. Show that the loop invariant is true before the loop starts

Maintenance

- 3. Show that the loop invariant holds when executing any iteration
- 4. Show that the loop invariant holds once the loop ends Ter



## Loop Invariant

- 1. At the start of the iteration with <reference the looping variable>, ...
  - For example: "index j"
- 2. The <reference to partial solution> ...
  - For example: "subarray array[0 ..= j-1] consists of the elements originally in array[0 ..= j-1]..."
- 3. <something about why the partial solution is correct>.
  - For example: "In sorted order."

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What would be a good loop invariant for proving this procedure?

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  - 1. A statement that can be easily proven true or false
  - 2. The statement must reference the purpose of the loop
  - 3. The statement must reference variables that change each iteration

# FUNCTION SumArray(array) sum = 0 i = 0 WHILE i < array.length sum = sum + array[i] i = i + 1</pre>

At the start of the iteration with index i, the variable sum is the sum of all values in the subarray array[0 . . < i].

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Initialization
 Maintenance
 Termination

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## Example

# FUNCTION SumArray(array) sum = 0 i = 0 WHILE i < array.length sum = sum + array[i] i = i + 1</pre>

## Initialization:

Upon entering the first iteration, i = 0. There are no numbers in the subarray array[0 ... < i]. The sum of no terms is the identity for addition (0).

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WHILE i < array.length
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At the start of the iteration with index i, the variable sum is the sum of all values in the subarray array[0 ... < i].

### Maintenance:

Upon entering an iteration with index i, assume that sum is equal to the sum of all values in the subarray array[0 ... < i]:  $sum = \sum_{i=0}^{i-1} array[i]$ 

The current iteration adds array[i] to sum and then increments i, so that the loop invariant holds upon entering the next iteration.

At the start of the iteration with index i, the variable sum is the sum of all values in the subarray array [0 ... < i].

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## Termination:

The loop terminates with i = n. According to the loop invariant, sum is equal to the sum of all values in the subarray array  $[0 \dots < i]$ :  $sum = \sum_{i=0}^{i-1} array[i] = \sum_{i=0}^{n-1} array[i]$ 

which is the sum of all values in the array.

## A more complex example: Dijkstra's Algorithm

DIJKSTRA (G, W, S) S = nullQ = G.Vwhile Q is not null u = EXTRACT-MIN(Q)S = S union  $\{u\}$ for each vertex v adjacent to u RELAX(U, V, W)

**Loop Invariant:** 

At the start of each iteration of the while loop, v.d = delta(s, v) for each vertex v in S.

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#### **Initialization:**

Initially, S = null and so the invariant is trivially true

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#### Maintenance:

<long proof by contradiction on page 661 of the textbook>

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#### **Termination:**

At termination, Q = null which, along with our earlier invariant that Q = V - S, implies that S = V. Thus, u.d = delta(s, u) for all vertices in G.V.