Binary Search Trees

https://cs.pomona.edu/classes/cs140/

Outline

Topics and Learning Objectives

- Compare binary search trees with sorted arrays
- Discuss the importance of a binary search tree's height
- Discuss common search tree algorithms

Exercise

• Search tree exercise

Extra Resources

• Introduction to Algorithms, 3rd, chapter 12

Sorted Arrays

Operation	<u>Running Time</u>
Access	O(1)
Search	O(lg n)
Selection	O(1)
Predecessor	O(1)
Successor	O(1)
Output (print)	O(n)
Insert	O(n)
Delete	O(n)
Extract-Min	O(n)

Given a set of key values, is a BST unique? (ignore ties)



Binary Search Tree

Each node has:

- A pointer to a left subtree
- A pointer to a right subtree
- A pointer to a parent node
- A piece of data (the key value)

Search tree property:

- All keys found in a left subtree must be less than the key of the current node
- All keys found in a right subtree must be greater than the key of the current node

Trees and Graphs

- Trees are a special type of graph
- Trees cannot contain cycles (acyclic)
- Trees always have directed edges
- Trees have a single source (no incoming edges) vertex called root
- All tree vertices have one parent (except root, which has no parents)
- Trees always have n-1 edges
- BST compared to Heap?
 - Heap is always balanced, BSTs are not necessarily balanced
 - They have different properties (where are lesser values?)

Balanced Binary Search Tree (vs Sorted Array)

Operation

Access Search Selection Predecessor Successor Output (print) Insert Delete Extract Min



Running Time $O(1) \rightarrow O(\lg n)$ $O(\lg n)$ $O(1) \rightarrow O(\lg n)$ $O(1) \rightarrow O(\lg n)$ $O(1) \rightarrow O(\lg n)$ O(n) $O(n) \rightarrow O(\lg n)$ $O(n) \rightarrow O(\lg n)$ $O(n) \rightarrow O(\lg n)$











- Given a set of keys, we have many different choices for creating a binary search tree (we just have to satisfy the search tree properties)
- If we have nodes, what is the maximum height of the tree?
- If we have n nodes, what is the minimum height of the tree?



Searching a BST

Search the tree T for the key k

8

6

- 1. Start at the root node
- 2. Recursively:
 - 1. Traverse left if k < current key
 - 2. Traverse right if k > current key
- 3. Return the node when found or return NULL

Inserting into a BST

Insert the key k into the tree T

8

6

- 1. Start at the root node
- 2. Search for the key k (probably won't find it)
- 3. Create a new node and setup the correct pointer



Given a binary search tree that is not necessarily balanced or unbalanced, what is the maximum number of hops needed to search the tree or insert a new node?

Options:

- a. 1
- b. lg n
- c. tree height

d. n



1 3 4 6 7 8 10 13 14

- Min
- Max
- Predecessor (k)
- Successor (k)
- What is the running time?







How do you find:

• <mark>Min</mark>

O(tree M) NetoM)

- Max
- Predecessor (k)
- Successor (k)
- What is the running time?





- Max GO GO + ree M Predecessor (k) O + rete M
 - Successor (k)
 - What is the running time?

- Predecessor (k) Successor (k) Successor (k)

 - What is the running time?

$\begin{array}{c} & & & \\ & & & \\ & & & \\ 1 & & & \\ & & & \\ 4 & & & \\ \end{array}$

1 3 4 6 7 8 10

13

14

- Min
- Max
- Predecessor (k)
- Successor (k)
- What is the running time?



How would you print all nodes in order?

• In-order traversal:

- Recursively visit nodes on the left
- Print out the current node
- Recursively visit nodes on the right
- What is the running time?



Post-Order Traversal

- Recursively visit nodes on the left
- Recursively visit nodes on the right
- "Visit" the current node



Pre-Order Traversal

- "Visit" the current node
- Recursively visit nodes on the left
- Recursively visit nodes on the right





Deleting a node from a BST

Deletion is often the most difficult task for treelike structures

- Search for the key
 - Case 1: If the node has no children then just delete
 - Case 2: If the node has one child then splice it out
 - Case 3: if the node has both children
 - Find the node's predecessor
 - Swap the node with its predecessor
 - Delete the node



Selection and Rank with a BST

How would you compute the ith order statistic using a BST?

Idea: store some metadata at each node

 Let size(x) = the number of nodes rooted at x (the number of nodes that can be reached via the left and right children pointers

How would you calculate size(x)?

- What kind of traversal would this use (in order, pre, or post)?



FUNCTION UpdateSizes(bst_node)
IF bst node != NONE





Selection and Rank with a BST

```
FUNCTION GetIthOrderStatistic(bst_node, i)
   left_child_size = bst_node.left.size
  \frac{5}{\text{IF left_child_size}} = (i - 1)
                                                                    8
      RETURN bst node.value
   ELSE IF left_child_size ≥ i
      RETURN GetIthOrderStatistic(bst_node.left, i)
                                                                 6
   ELSE
      new_i = i - left_child_size - 1
      RETURN GetIthOrderStatistic(bst node.right, new i)
```

Balanced Binary Search Trees All ops depend on the height of

- Why is balancing important? \mathcal{V}
- What is the worst-case height for a binary tree?
- Balanced tree: the height of a balanced tree stays O(lg n) after insertions and deletions
- Many different types of balanced search trees:
 - AVL Tree, Splay Tree, B Tree, Red-Black Tree