

Drawing: Function Call

- Take three minutes to draw something related to a “Function Call”
- Include the following
 - A diagram of the stack
 - A few of the relevant assembly instructions

Assembly

B:

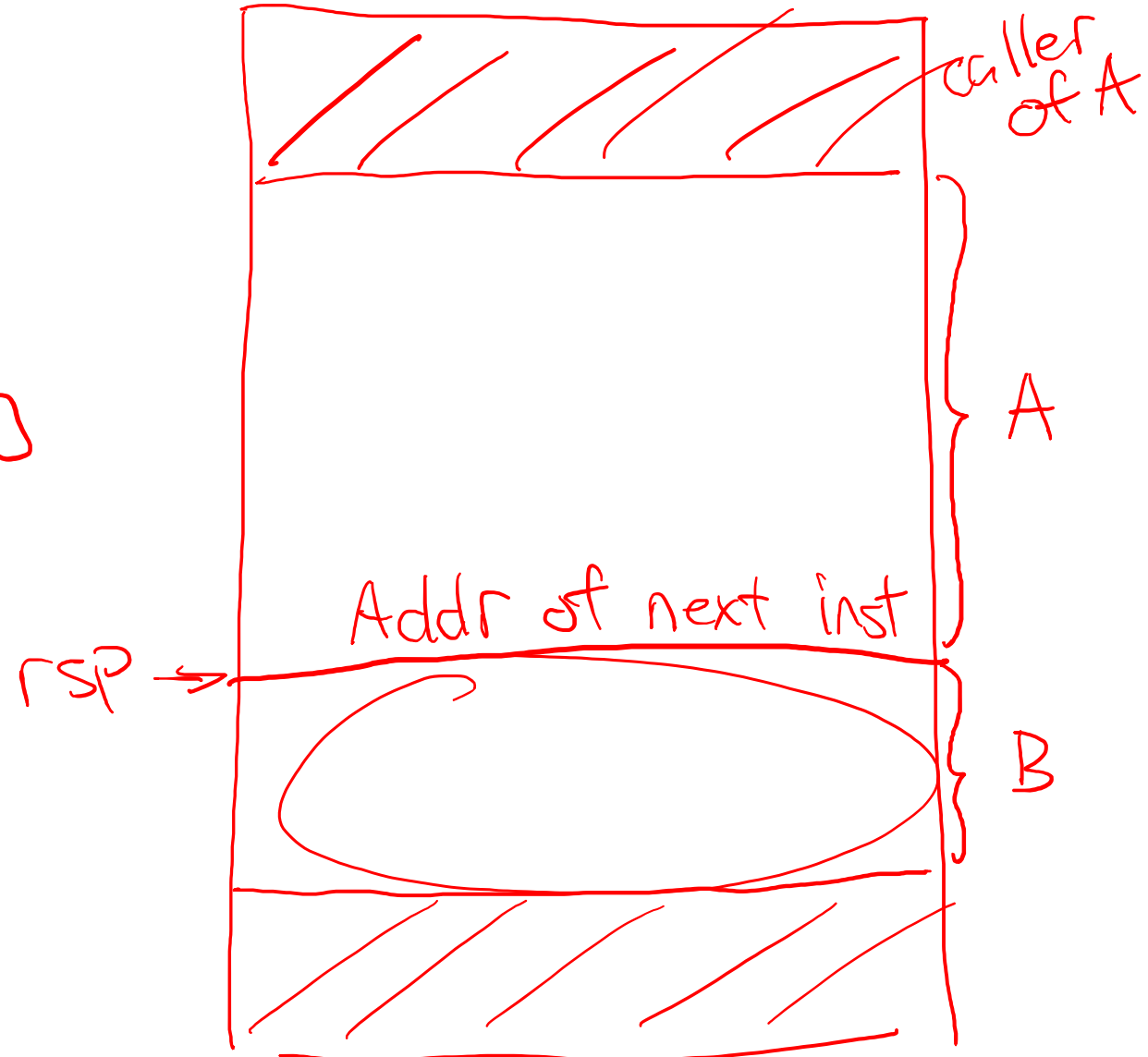
```
sub    rsp, 16 0x10
add
sub
mov    rax, [rsp + 0x8]
add   rsp, 0x10
ret
```

A:

```
mov    rdx, 6
mov    rsi, [rsp + 8]
call   B
```

// do something w/
rax

Stack



Security and Overflows

Stack Smashing and Return Oriented Programming

```

#include <stdio.h>

char *read_string() {
    char char_buffer[10];
    gets(char_buffer);
    return char_buffer;
}

void print_string(char *str) {
    printf("%s\n", str);
}

int main() {
    char *user_input = read_string();
    print_string(user_input);
}

```

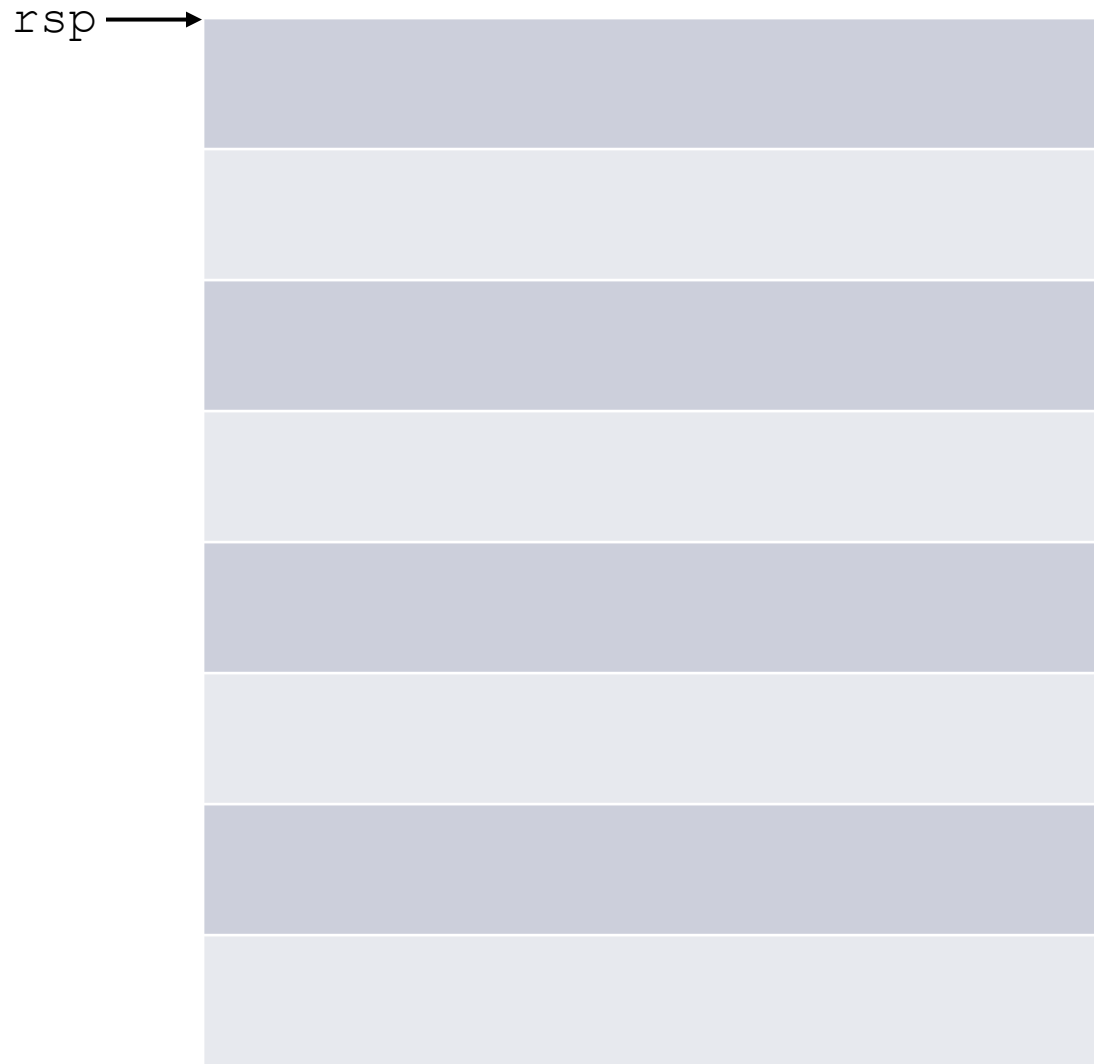
```

read_string:
    sub     rsp, 24
    lea    rdi, [rsp+6]
    mov    eax, 0
    call   gets
    mov    eax, 0
    add    rsp, 24
    ret

print_string:
    sub    rsp, 8
    call   puts
    add    rsp, 8
    ret

main:
    sub    rsp, 8
    mov    eax, 0
    call   read_string
    movsx rdi, eax
    mov    eax, 0
    call   print_string
    mov    eax, 0
    add    rsp, 8
    ret

```



```

read_string:
    sub     rsp, 24
    lea    rdi, [rsp+6]
    mov    eax, 0
    call   gets
    mov    eax, 0
    add    rsp, 24
    ret

```

```

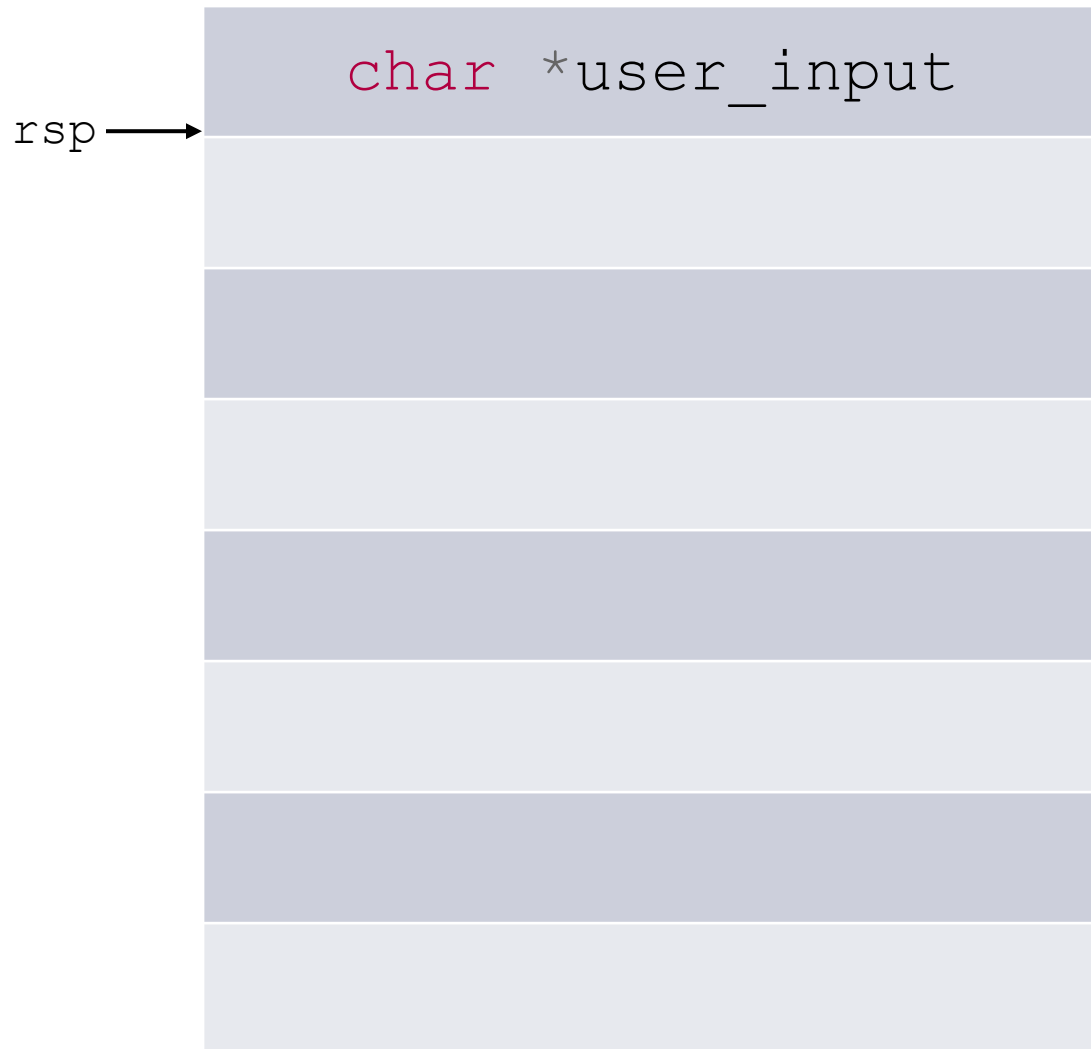
print_string:
    sub     rsp, 8
    call   puts
    add    rsp, 8
    ret

```

```

main:
    rip → sub     rsp, 8
          mov    eax, 0
          call   read_string
          movsx  rdi, eax
          mov    eax, 0
          call   print_string
          mov    eax, 0
          add    rsp, 8
          ret

```



```

read_string:
    sub    rsp, 24
    lea   rdi, [rsp+6]
    mov   eax, 0
    call  gets
    mov   eax, 0
    add   rsp, 24
    ret

```

```

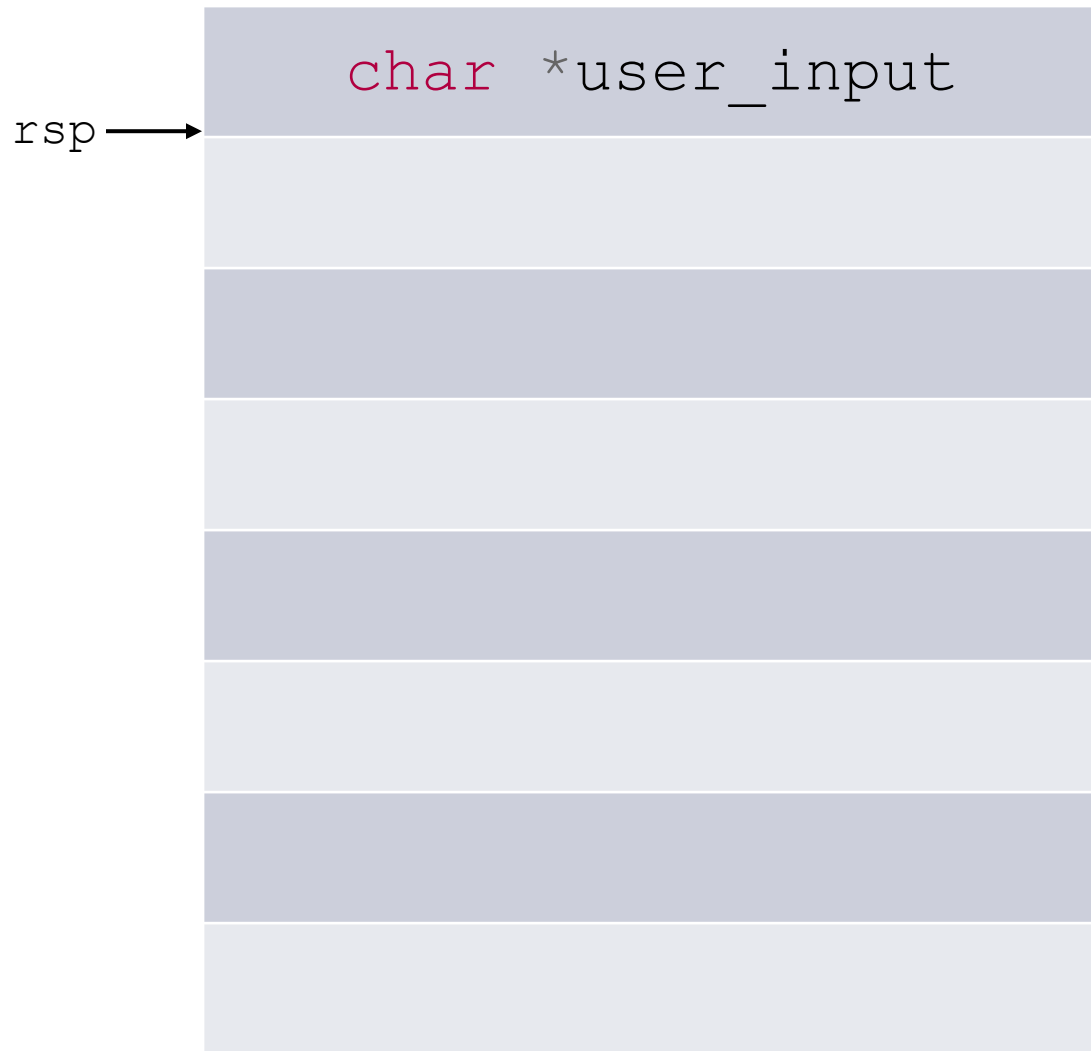
print_string:
    sub    rsp, 8
    call  puts
    add   rsp, 8
    ret

```

```

main:
    sub    rsp, 8
    rip → mov   eax, 0
    call  read_string
    movsx rdi, eax
    mov   eax, 0
    call  print_string
    mov   eax, 0
    add   rsp, 8
    ret

```



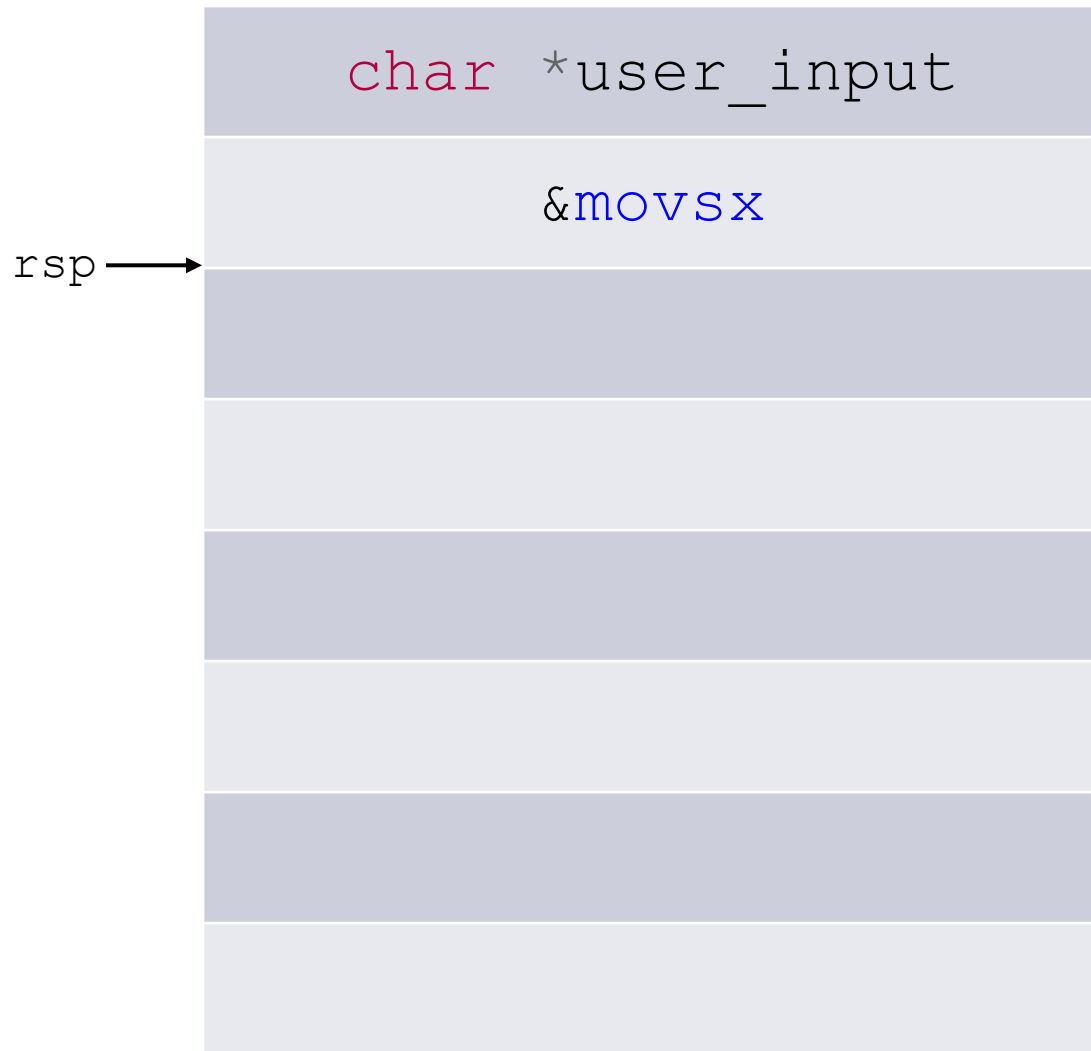
rax
0

rdi

```
read_string:
    sub    rsp, 24
    lea   rdi, [rsp+6]
    mov   eax, 0
    call  gets
    mov   eax, 0
    add   rsp, 24
    ret
```

```
print_string:
    sub    rsp, 8
    call  puts
    add   rsp, 8
    ret
```

```
main:
    sub    rsp, 8
    mov   eax, 0
    rip → call  read_string
    movsx rdi, eax
    mov   eax, 0
    call  print_string
    mov   eax, 0
    add   rsp, 8
    ret
```



```

read_string:
rip → sub    rsp, 24
      lea   rdi, [rsp+6]
      mov   eax, 0
      call gets
      mov   eax, 0
      add   rsp, 24
      ret

```

```

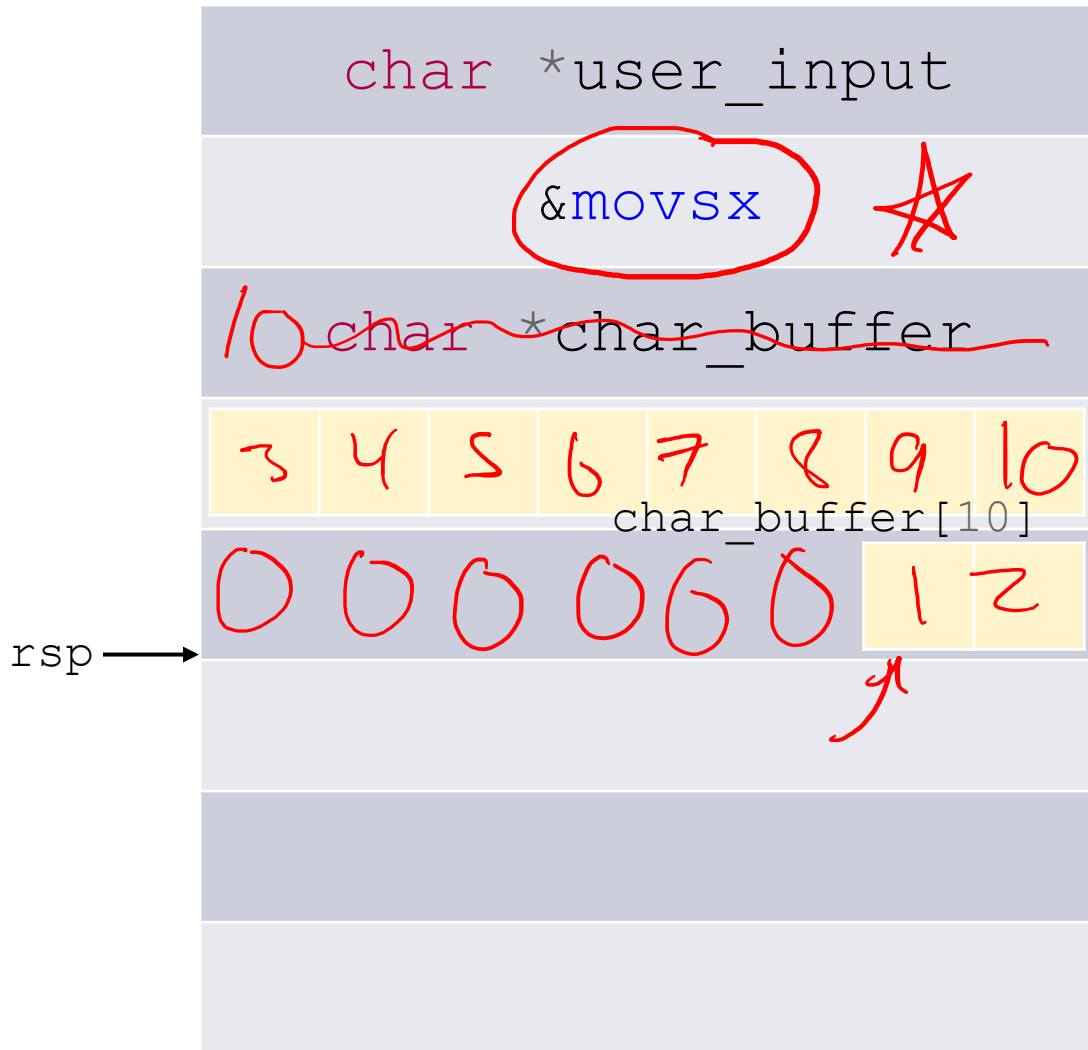
print_string:
      sub   rsp, 8
      call puts
      add   rsp, 8
      ret

```

```

main:
      sub   rsp, 8
      mov   eax, 0
      call read_string
      movsx rdi, eax
      mov   eax, 0
      call print_string
      mov   eax, 0
      add   rsp, 8
      ret

```

read_string:

```

sub     rsp, 24
rip → lea   rdi, [rsp+6]
mov     eax, 0
call   gets
mov     eax, 0
add    rsp, 24
ret

```

print_string:

```

sub     rsp, 8
call   puts
add    rsp, 8
ret

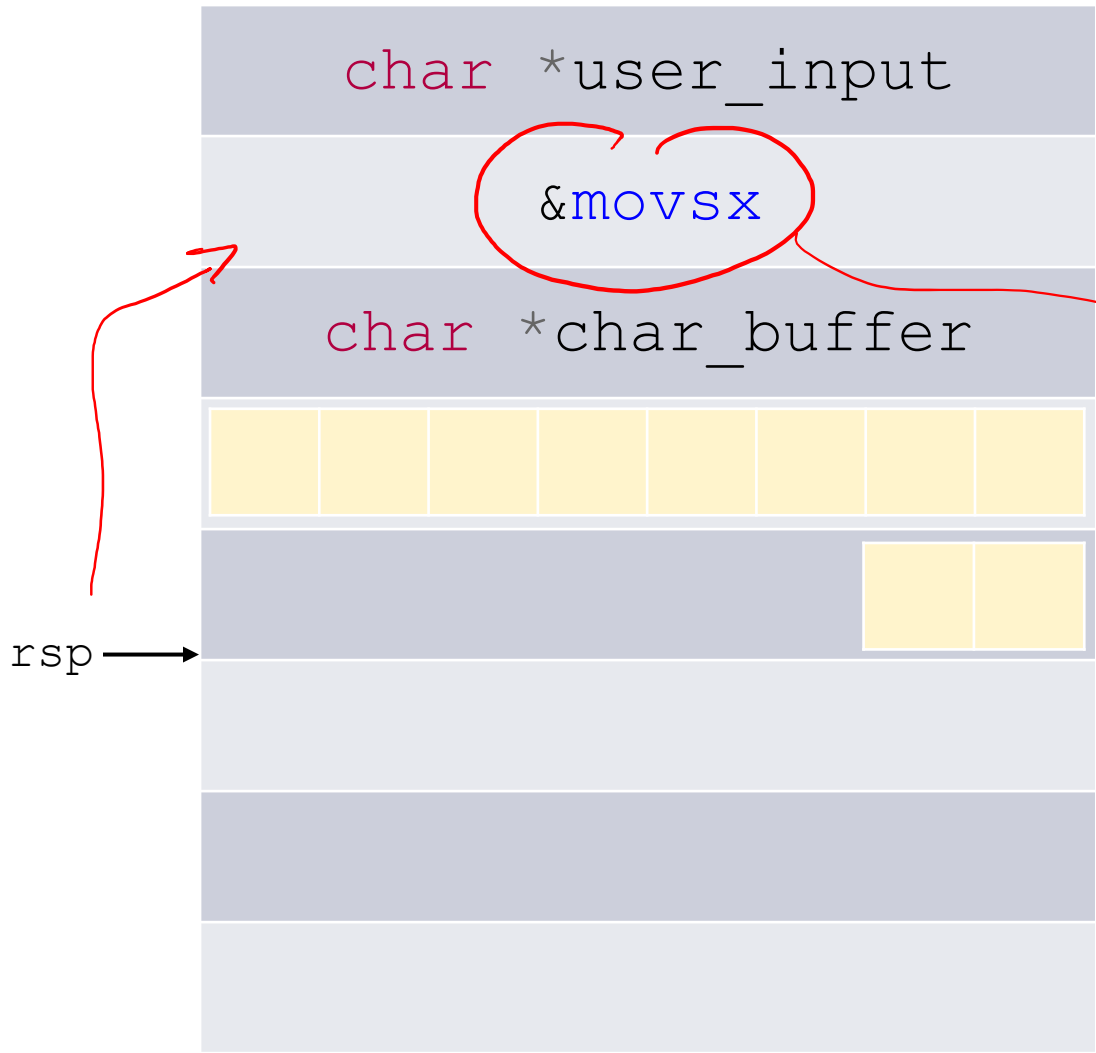
```

main:

```

sub     rsp, 8
mov     eax, 0
call   read_string
movsx  rdi, eax
mov     eax, 0
call   print_string
mov     eax, 0
add    rsp, 8
ret

```



```

read_string:
    sub    rsp, 24
    lea   rdi, [rsp+6]
rip → mov   eax, 0
    call  gets
    mov   eax, 0
    add   rsp, 24
    ret

```

```

print_string:
    sub    rsp, 8
    call   puts
    add    rsp, 8
    ret

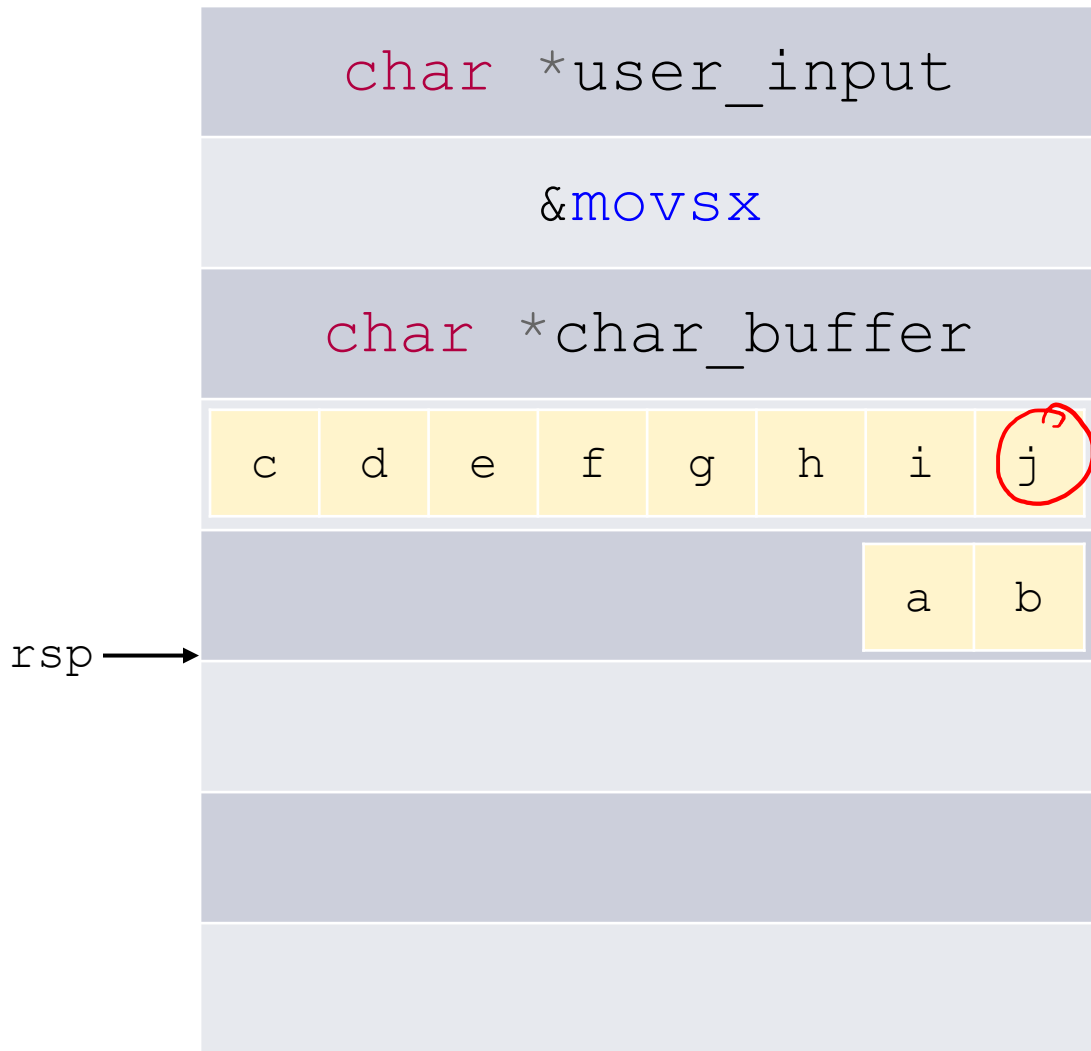
```

```

main:
    sub    rsp, 8
    mov   eax, 0
    call  read_string
    movsx rdi, eax
    mov   eax, 0
    call  print_string
    mov   eax, 0
    add   rsp, 8
    ret

```

rax	rdi
0	&char_buffer



10

```

read_string:
    sub     rsp, 24
    lea    rdi, [rsp+6]
    mov    eax, 0
    call   gets
    mov    eax, 0
rip → add   rsp, 24
    ret

```

```

print_string:
    sub     rsp, 8
    call   puts
    add    rsp, 8
    ret

```

```

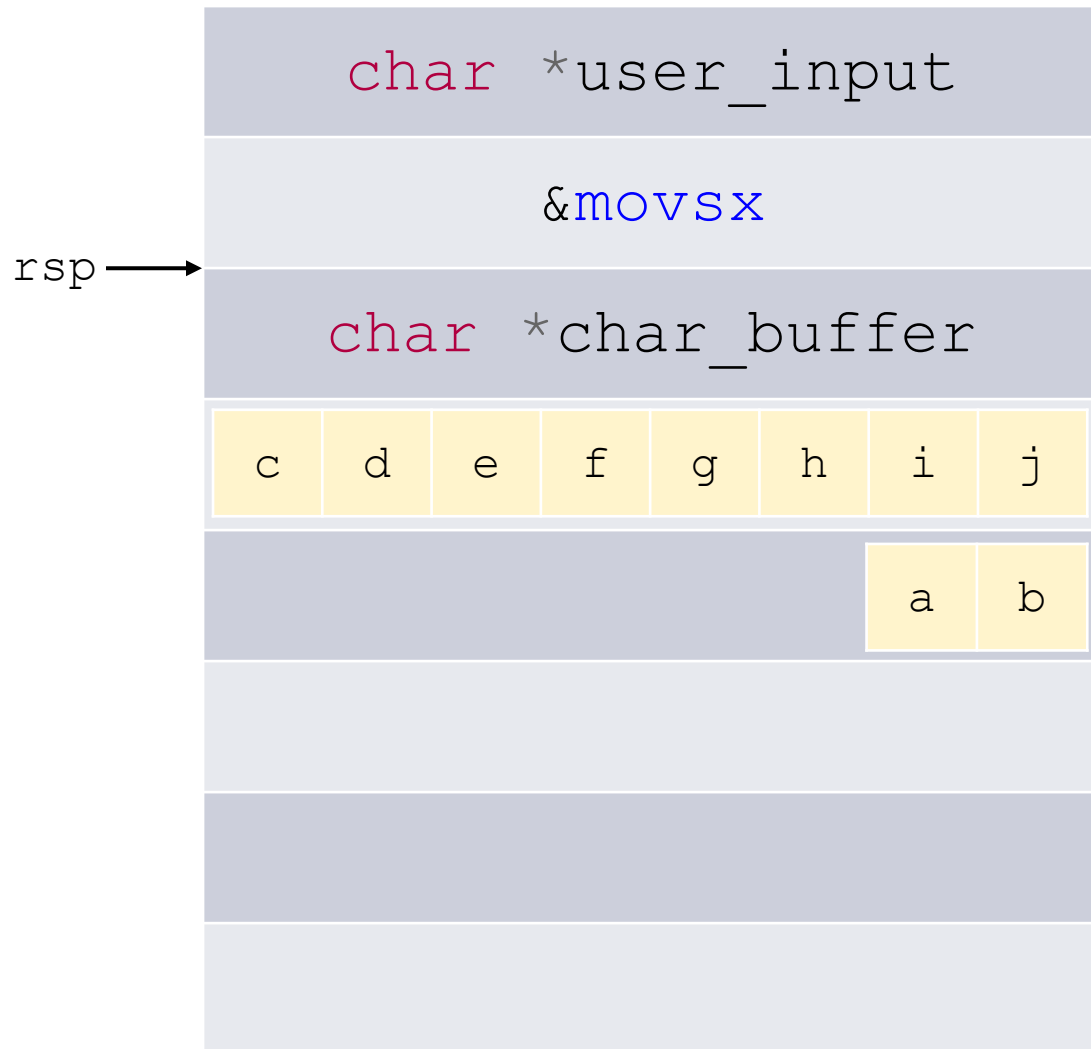
main:
    sub     rsp, 8
    mov    eax, 0
    call   read_string
    movsx  rdi, eax
    mov    eax, 0
    call   print_string
    mov    eax, 0
    add    rsp, 8
    ret

```

```

    rax    0
    rdi    &char_buffer

```



rax

0

rdi

&char_buffer

read_string:

```

sub    rsp, 24
lea   rdi, [rsp+6]
mov   eax, 0
call  gets
mov   eax, 0
add   rsp, 24

```

rip → ret

print_string:

```

sub    rsp, 8
call  puts
add   rsp, 8
ret

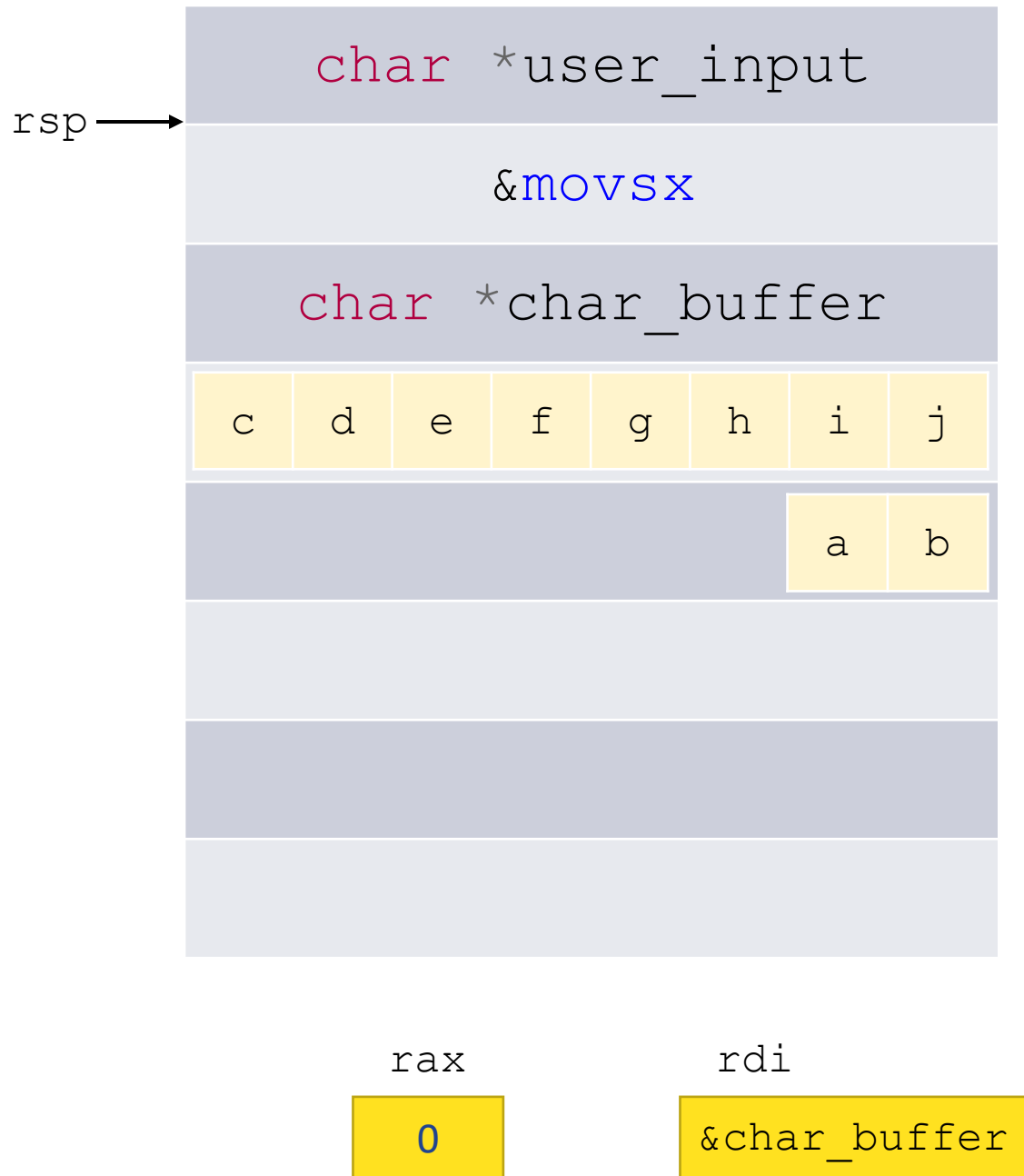
```

main:

```

sub    rsp, 8
mov   eax, 0
call  read_string
movsx rdi, eax
mov   eax, 0
call  print_string
mov   eax, 0
add   rsp, 8
ret

```



```

read_string:
    sub     rsp, 24
    lea    rdi, [rsp+6]
    mov    eax, 0
    call   gets
    mov    eax, 0
    add    rsp, 24
    ret

```

```

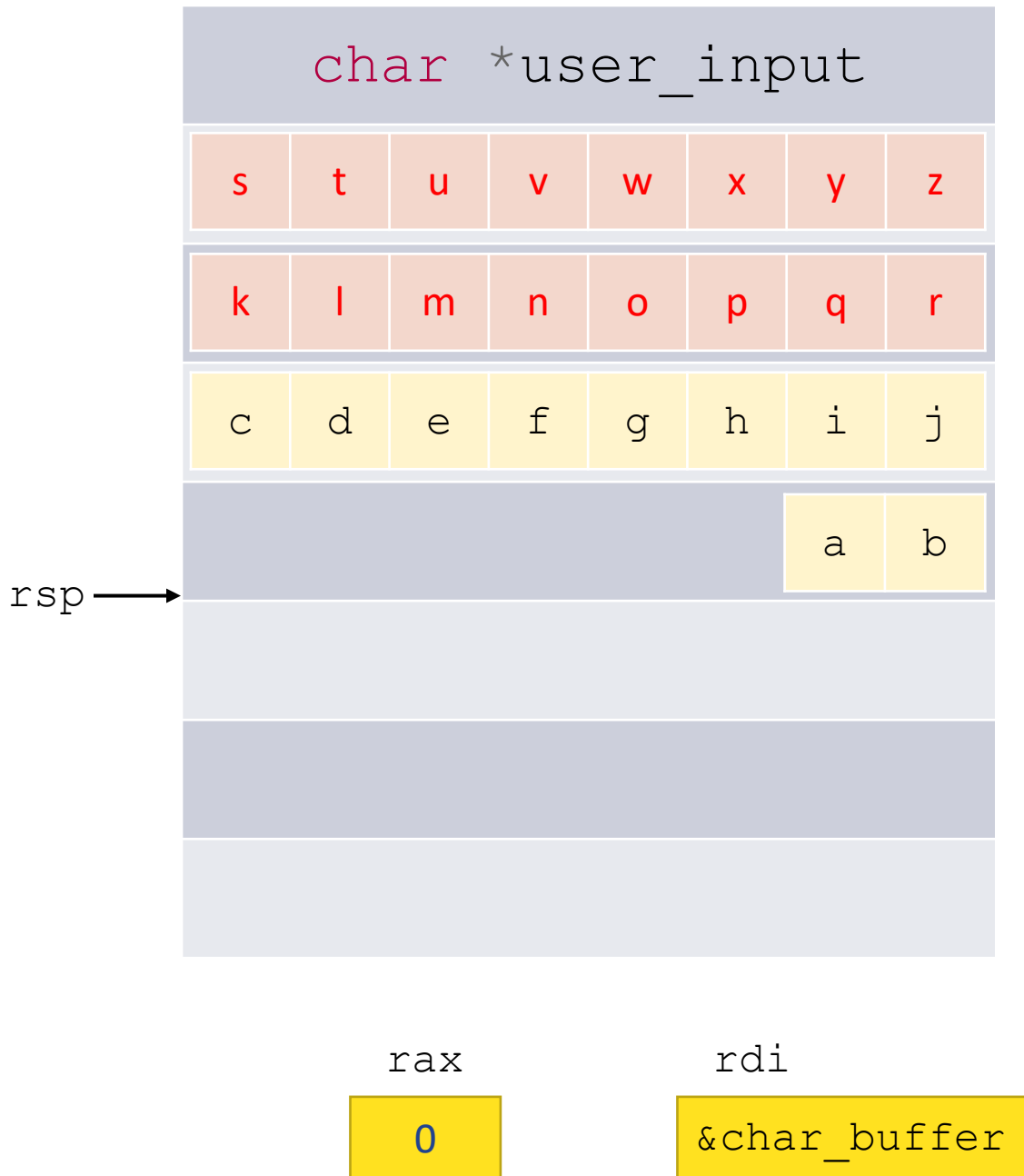
print_string:
    sub     rsp, 8
    call   puts
    add    rsp, 8
    ret

```

```

main:
    sub     rsp, 8
    mov    eax, 0
    call   read_string
    rip → movsx rdi, eax
    mov    eax, 0
    call   print_string
    mov    eax, 0
    add    rsp, 8
    ret

```



```

read_string:
    sub     rsp, 24
    lea    rdi, [rsp+6]
    mov    eax, 0
    call   gets
    mov    eax, 0
rip → add   rsp, 24
    ret

```

```

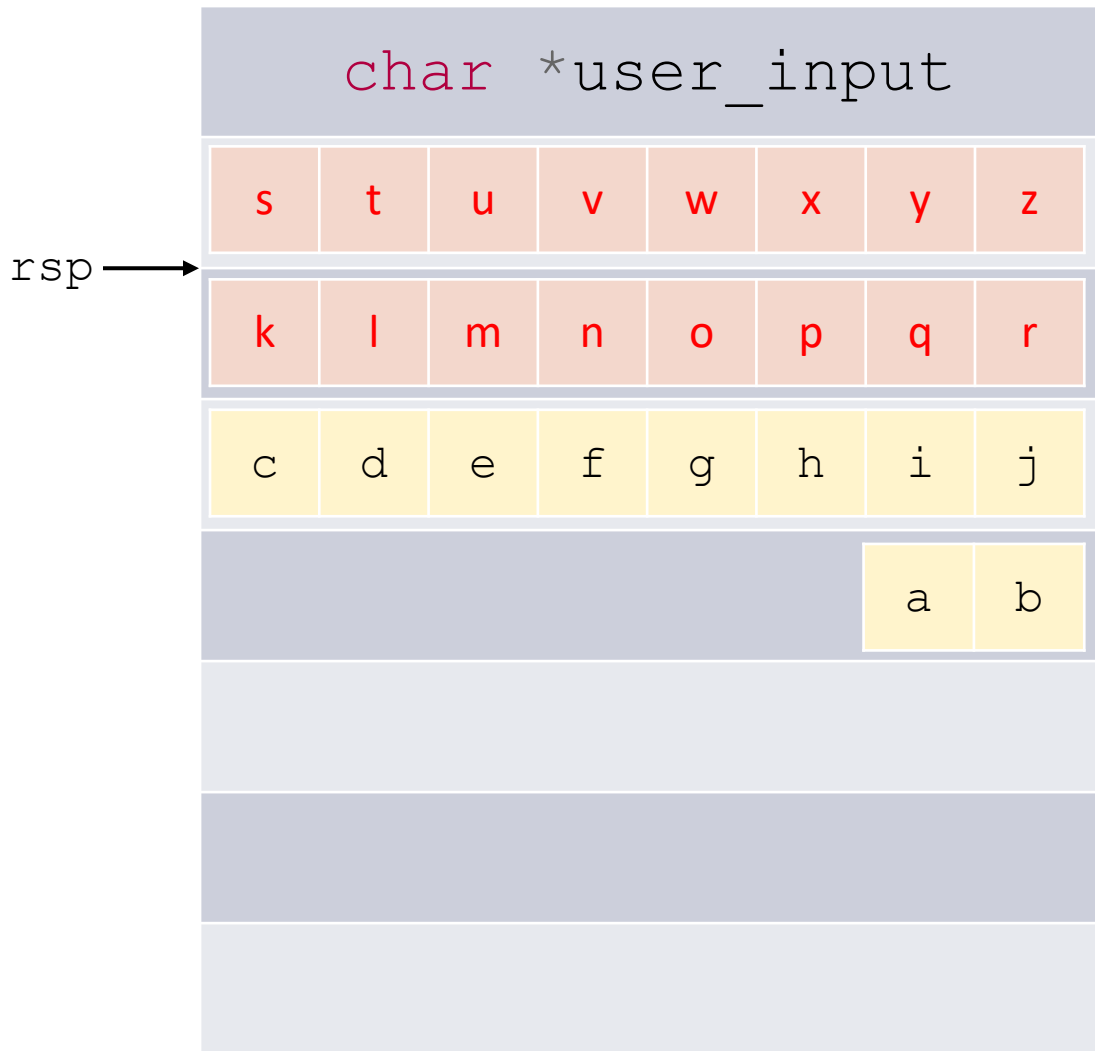
print_string:
    sub     rsp, 8
    call   puts
    add    rsp, 8
    ret

```

```

main:
    sub     rsp, 8
    mov    eax, 0
    call   read_string
    movsx  rdi, eax
    mov    eax, 0
    call   print_string
    mov    eax, 0
    add    rsp, 8
    ret

```



```

rax      rdi
0        &char_buffer

```

```

read_string:
    sub     rsp, 24
    lea    rdi, [rsp+6]
    mov    eax, 0
    call   gets
    mov    eax, 0
    rip ← add    rsp, 24
          ret

```

```

print_string:
    sub     rsp, 8
    call   puts
    add    rsp, 8
    ret

```

```

main:
    sub     rsp, 8
    mov    eax, 0
    call   read_string
    movsx  rdi, eax
    mov    eax, 0
    call   print_string
    mov    eax, 0
    add    rsp, 8
    ret

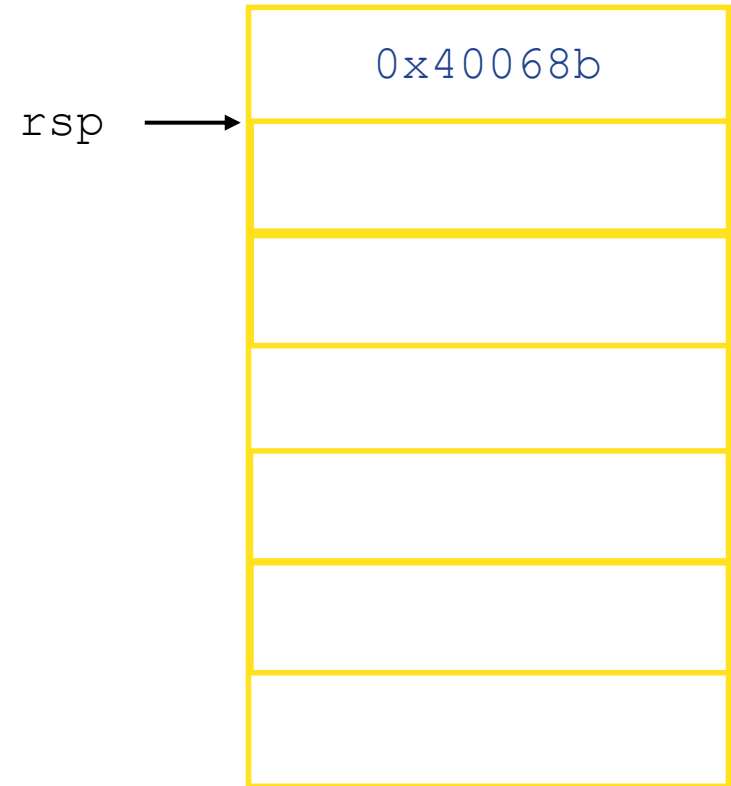
```

Practice with Memory Bugs

- What is the state of the stack immediately before the program returns from f2?
- What will happen immediately after f2 returns?

```
int f2(){  
    int a1[4] = {1,2,3,4};  
    a1[6] = 47;  
}
```

```
rip → f2:  
sub    rsp, 0x18  
mov    [rsp], 0x1  
mov    [rsp + 0x4], 0x2  
mov    [rsp + 0x8], 0x3  
mov    [rsp + 0xC], 0x4  
mov    [rsp + 0x18], 0x2F  
add    rsp, 0x18  
ret
```



Practice with Memory Bugs

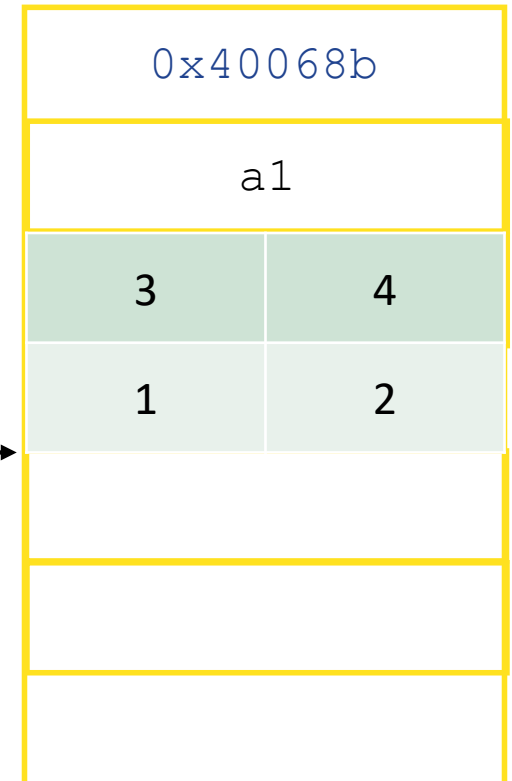
- What is the state of the stack immediately before the program returns from f2?
- What will happen immediately after f2 returns?

```
int f2(){  
    int a1[4] = {1,2,3,4};  
    a1[6] = 47;  
}
```

rip →

```
f2:  
    sub    rsp, 0x18  
    mov    [rsp], 0x1  
    mov    [rsp + 0x4], 0x2  
    mov    [rsp + 0x8], 0x3  
    mov    [rsp + 0xC], 0x4  
    mov    [rsp + 0x18], 0x2F  
    add    rsp, 0x18  
    ret
```

rsp →

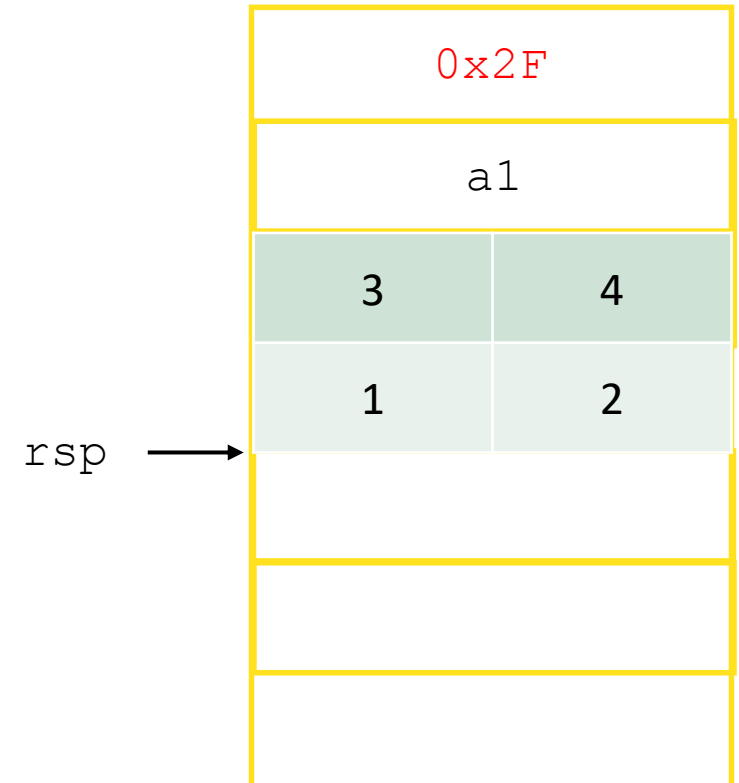


Practice with Memory Bugs

- What is the state of the stack immediately before the program returns from f2?
- What will happen immediately after f2 returns?

```
int f2(){  
    int a1[4] = {1,2,3,4};  
    a1[6] = 47;  
}
```

```
f2:  
    sub    rsp, 0x18  
    mov    [rsp], 0x1  
    mov    [rsp + 0x4], 0x2  
    mov    [rsp + 0x8], 0x3  
    mov    [rsp + 0xC], 0x4  
    mov    [rsp + 0x18], 0x2F  
    add    rsp, 0x18  
    ret
```

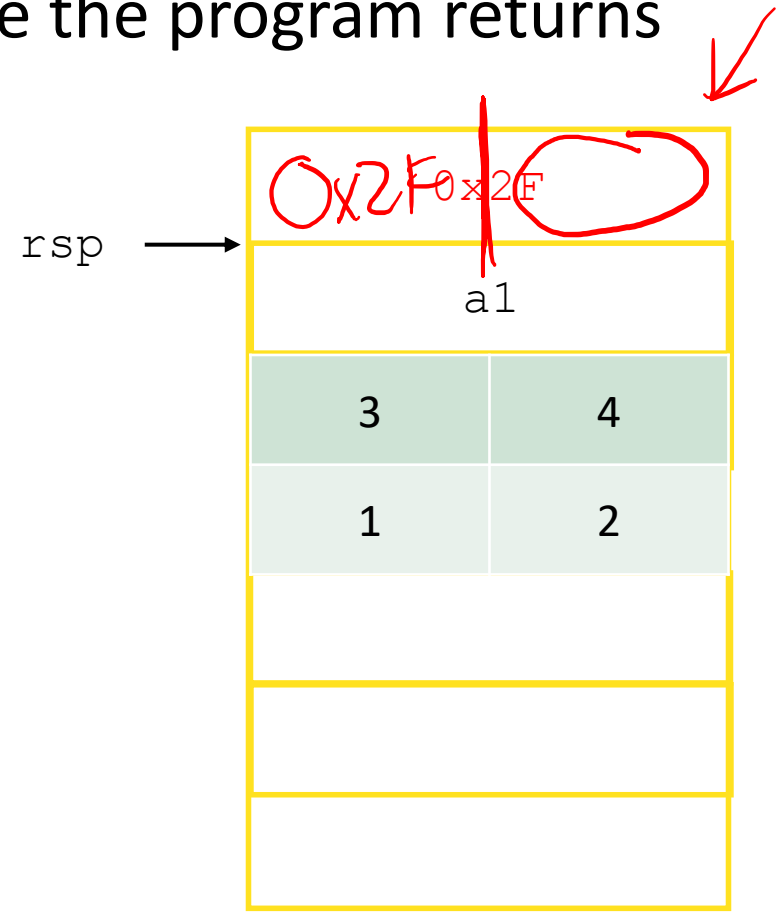


Practice with Memory Bugs

- What is the state of the stack immediately before the program returns from f2?
- What will happen immediately after f2 returns?

```
int f2(){  
    int a1[4] = {1,2,3,4};  
    a1[6] = 47;  
}
```

```
f2:  
    sub    rsp, 0x18  
    mov    [rsp], 0x1  
    mov    [rsp + 0x4], 0x2  
    mov    [rsp + 0x8], 0x3  
    mov    [rsp + 0xC], 0x4  
    mov    [rsp + 0x18], 0x2F  
    add    rsp, 0x18  
    rip → ret
```



Practice with Buffer Overflow Exploits

Construct an exploit string that will successfully cause the program to print "You are now logged in" without knowing the correct password

1. How many bytes of padding are in this exploit string?
2. What value will you overwrite the return address with?

```
int authenticate(char *password) {
    char buf[4];
    gets(buf);
    int correct = !strcmp(password, buf);
    return correct;
}

int main(int argc, char ** argv) {
    char * pw = "123456";
    printf("Enter your password: ");
    while(!authenticate(pw)) {
        printf("Incorrect. Try again: ");
    }
    printf("You are now logged in\n");
    return 0;
}
```

Practice Exploits

Construct an exploit that successfully prints "You are now logged in" without knowing the password.

1. How many bytes does this exploit need?

2. What value should we return after the exploit?

```
.LC0: .string "Enter your password: "  
.LC1: .string "Incorrect. Try again: "  
.LC2: .string "123456"  
.LC3: .string "You are now logged in"  
authenticate:  
    push    rbx  
    sub     rsp, 16  
    mov     rbx, rdi  
    lea    rdi, [rsp+12]  
    mov     eax, 0  
    call   gets  
    lea    rsi, [rsp+12]  
    mov     rdi, rbx  
    call   strcmp  
    test   eax, eax  
    sete   al  
    movzx  eax, al  
    add    rsp, 16  
    pop    rbx  
    ret  
  
main:  
    sub     rsp, 8  
    mov     edi, OFFSET FLAT:.LC0  
    mov     eax, 0  
    call   printf  
    jmp    .L4  
  
.L5:  
    mov     edi, OFFSET FLAT:.LC1  
    mov     eax, 0  
    call   printf  
  
.L4:  
    mov     edi, OFFSET FLAT:.LC2  
    call   authenticate  
    test   eax, eax  
    je     .L5  
    mov     edi, OFFSET FLAT:.LC3  
    call   puts  
    mov     eax, 0  
    add    rsp, 8  
    ret
```

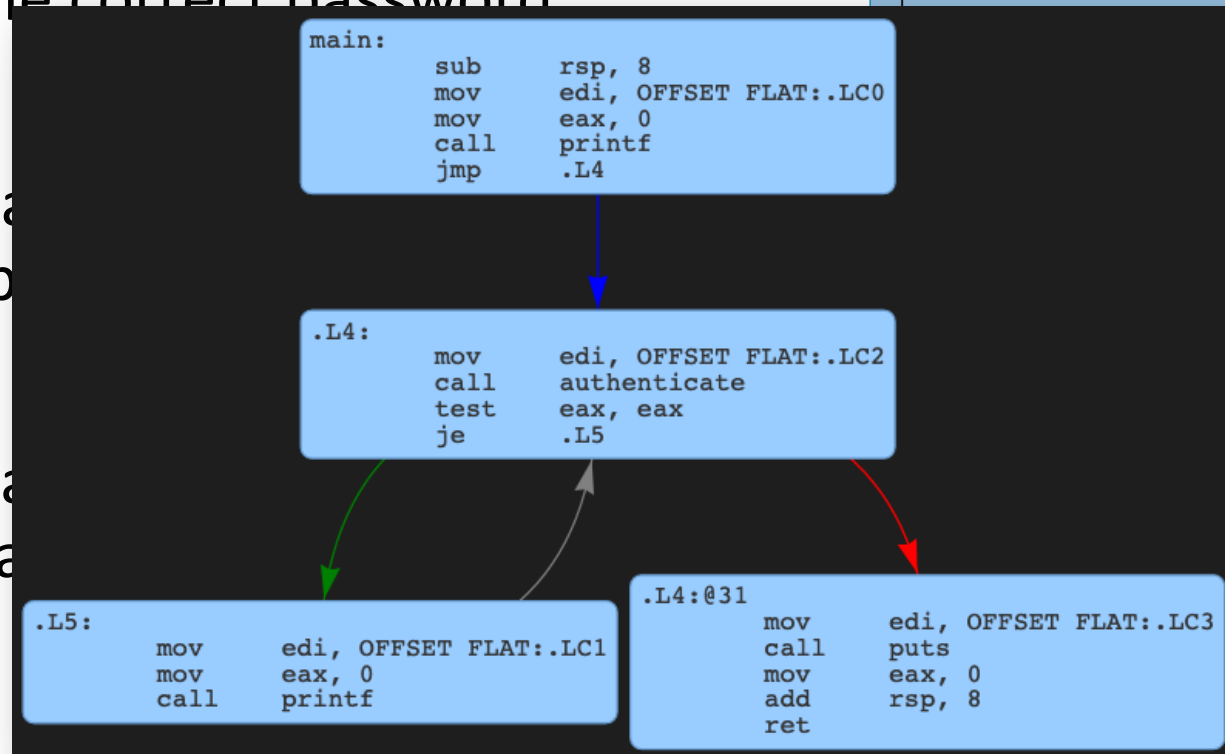
Workflow

```
int authenticate(char *password) {  
    char buf[4];  
    gets(buf);  
    int correct = !strcmp(password, buf);  
    return correct;  
}  
  
int main(int argc, char ** argv) {  
    char * pw = "123456";  
    printf("Enter your password: ");  
    while(!authenticate(pw)) {  
        printf("Incorrect. Try again: ");  
    }  
    printf("You are now logged in\n");  
    return 0;  
}
```

Practice with Buffer Overflow Exploits

Construct an exploit string that will successfully cause the program to print "You are now logged in" without knowing the correct password

1. How many characters does this exploit need?
2. What value should we return after the exploit?



```
int authenticate()
{
    char buf[4];
    gets(buf);
    int correct = 0;
    return correct;
}
```

```
.LC0: .string "Enter your password: "
.LC1: .string "Incorrect. Try again: "
.LC2: .string "123456"
.LC3: .string "You are now logged in"
authenticate:
    push    rbx
    sub     rsp, 16
    mov     rbx, rdi
    lea    rdi, [rsp+12]
    mov     eax, 0
    call   gets
    lea    rsi, [rsp+12]
    mov     rdi, rbx
    call   strcmp
    test   eax, eax
    sete   al
    movzx  eax, al
    add    rsp, 16
    pop    rbx
    ret

main:
    sub     rsp, 8
    mov     edi, OFFSET FLAT:.LC0
    mov     eax, 0
    call   printf
    jmp    .L4

.L5:
    mov     edi, OFFSET FLAT:.LC1
    mov     eax, 0
    call   printf

.L4:
    mov     edi, OFFSET FLAT:.LC2
    call   authenticate
    test   eax, eax
    je     .L5
    mov     edi, OFFSET FLAT:.LC3
    call   puts
    mov     eax, 0
    add    rsp, 8
    ret
```

Practice with Buffer Overflow Exploits

Construct an exploit string that will successfully cause the program to print "You are now logged in" without knowing the correct password

1. How many bytes of padding are in this exploit string?
2. What value will you overwrite the return address with?

```
int authenticate()  
{  
    char buf[4];  
    gets(buf);  
    int correct = 0;  
    return correct;  
}  
  
int main(int argc, char * argv[])  
{  
    char * pw = "123456";  
    printf("Enter your password: ");  
    while(!authenticate())  
        printf("Incorrect. Try again: ");  
    printf("You are now logged in");  
    return 0;  
}
```

```
.LC0:    .string "Enter your password: "  
.LC1:    .string "Incorrect. Try again: "  
.LC2:    .string "123456"  
.LC3:    .string "You are now logged in"  
authenticate:  
    push    rbx  
    sub     rsp, 16  
    mov     rbx, rdi  
    lea    rdi, [rsp+12]  
    mov     eax, 0  
    call   gets  
    lea    rsi, [rsp+12]  
    mov     rdi, rbx  
    call   strcmp  
    test   eax, eax  
    sete   al  
    movzx  eax, al  
    add    rsp, 16  
    pop    rbx  
    ret  
  
main:  
401188  sub     rsp, 8  
40118c  mov     edi, OFFSET FLAT:.LC0  
401191  mov     eax, 0  
401196  call   printf  
40119b  jmp     .L4  
.L5:  
40119d  mov     edi, OFFSET FLAT:.LC1  
4011a2  mov     eax, 0  
4011a7  call   printf  
.L4:  
4011ac  mov     edi, OFFSET FLAT:.LC2  
4011b1  call   authenticate  
4011b6  test   eax, eax  
4011b8  je     .L5  
4011ba  mov     edi, OFFSET FLAT:.LC3  
4011bf  call   puts  
4011c4  mov     eax, 0  
4011c9  add    rsp, 8  
4011cd  ret
```

Practice with Buffer Overflow Exploits

Construct an exploit string that will successfully cause the program to print "You are now logged in" without knowing the correct password

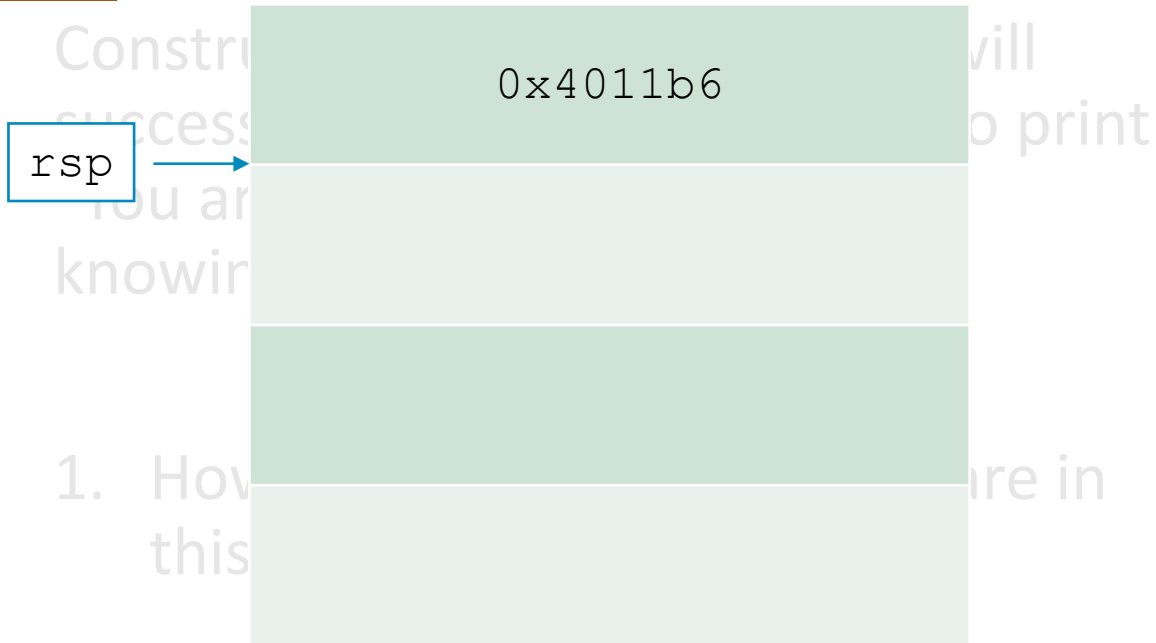
1. How many bytes of padding are in this exploit string?
2. What value will you overwrite the return address with?

```
int authenticate()  
{  
    char buf[4];  
    gets(buf);  
    int correct = 0;  
    return correct;  
}  
  
int main(int argc, char * argv[])  
{  
    char * pw = "123456";  
    printf("Enter your password: ");  
    while(!authenticate())  
        printf("Incorrect. Try again: ");  
    printf("You are now logged in");  
    return 0;  
}
```

```
.LC0:    .string "Enter your password: "  
.LC1:    .string "Incorrect. Try again: "  
.LC2:    .string "123456"  
.LC3:    .string "You are now logged in"  
authenticate:  
    push    rbx  
    sub     rsp, 16  
    mov     rbx, rdi  
    lea    rdi, [rsp+12]  
    mov     eax, 0  
    call   gets  
    lea    rsi, [rsp+12]  
    mov     rdi, rbx  
    call   strcmp  
    test   eax, eax  
    sete   al  
    movzx  eax, al  
    add    rsp, 16  
    pop    rbx  
    ret  
  
main:  
401188  sub     rsp, 8  
40118c  mov     edi, OFFSET FLAT:.LC0  
401191  mov     eax, 0  
401196  call   printf  
40119b  jmp     .L4  
.L5:  
40119d  mov     edi, OFFSET FLAT:.LC1  
4011a2  mov     eax, 0  
4011a7  call   printf  
.L4:  
4011ac  mov     edi, OFFSET FLAT:.LC2  
4011b1  call   authenticate  
4011b6  test   eax, eax  
4011b8  je     .L5  
4011ba  mov     edi, OFFSET FLAT:.LC3  
4011bf  call   puts  
4011c4  mov     eax, 0  
4011c9  add    rsp, 8  
4011cd  ret
```


Practice with Buffer Overflow Exploits

PUSH RBX



1. How will you overwrite the return address in this

2. What value will you overwrite the return address with?

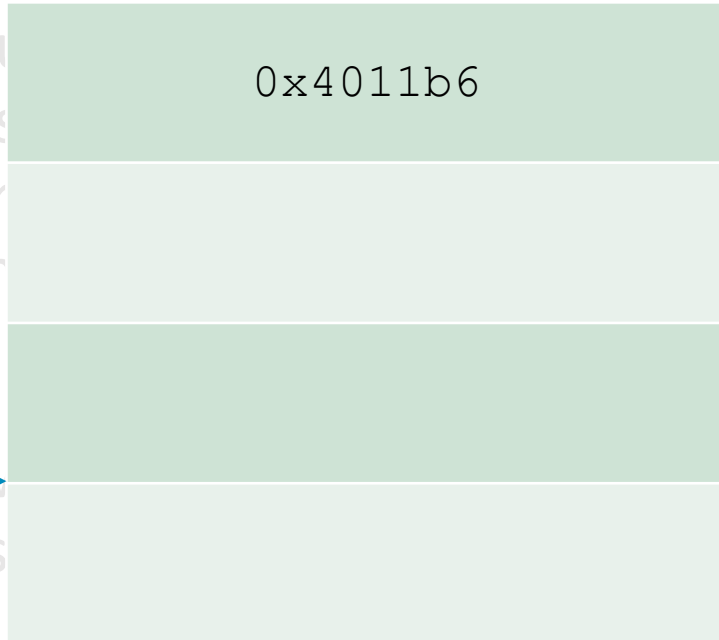
```
int authenticate()  
{  
    char buf[4];  
    gets(buf);  
    int correct = 0;  
    return correct;  
}  
  
int main(int argc, char * argv[])  
{  
    char * pw = "123456";  
    printf("Enter your password: ");  
    while(!authenticate())  
        printf("Incorrect. Try again: ");  
    printf("You are now logged in");  
    return 0;  
}
```

```
.LC0:    .string "Enter your password: "  
.LC1:    .string "Incorrect. Try again: "  
.LC2:    .string "123456"  
.LC3:    .string "You are now logged in"  
authenticate:  
    push    rbx  
    sub     rsp, 16  
    mov     rbx, rdi  
    lea    rdi, [rsp+12]  
    mov     eax, 0  
    call   gets  
    lea    rsi, [rsp+12]  
    mov     rdi, rbx  
    call   strcmp  
    test   eax, eax  
    sete   al  
    movzx  eax, al  
    add    rsp, 16  
    pop    rbx  
    ret  
  
main:  
401188  sub     rsp, 8  
40118c  mov     edi, OFFSET FLAT:.LC0  
401191  mov     eax, 0  
401196  call   printf  
40119b  jmp     .L4  
.L5:  
40119d  mov     edi, OFFSET FLAT:.LC1  
4011a2  mov     eax, 0  
4011a7  call   printf  
.L4:  
4011ac  mov     edi, OFFSET FLAT:.LC2  
4011b1  call   authenticate  
4011b6  test   eax, eax  
4011b8  je     .L5  
4011ba  mov     edi, OFFSET FLAT:.LC3  
4011bf  call   puts  
4011c4  mov     eax, 0  
4011c9  add    rsp, 8  
4011cd  ret
```

Practice with Buffer Overflow Exploits

Construct a buffer overflow exploit that will successfully overwrite the return address with "You are now logged in" knowing the address of the return address.

rsp



2. What value will you overwrite the return address with?

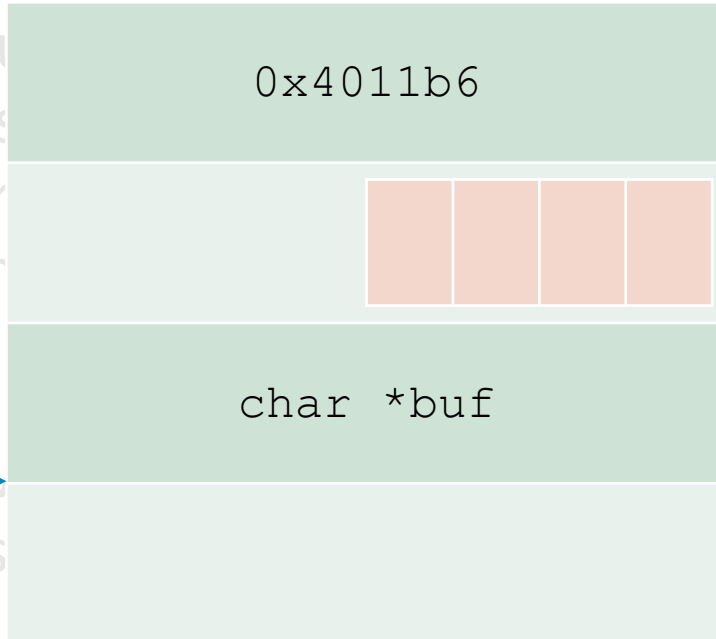
```
int authenticate()  
    char buf[4];  
    gets(buf);  
    int correct = 0;  
    return correct;  
}
```

```
int main(int argc, char * argv[])  
    char * pw = "123456";  
    printf("Enter password: ");  
    while(!authenticate())  
        printf("Incorrect. Try again: ");  
    printf("You are now logged in");  
    return 0;  
}
```

```
.LC0:  .string "Enter your password: "  
.LC1:  .string "Incorrect. Try again: "  
.LC2:  .string "123456"  
.LC3:  .string "You are now logged in"  
authenticate:  
    push    rbx  
    sub     rsp, 16  
    mov     rbx, rdi  
    lea    rdi, [rsp+12]  
    mov     eax, 0  
    call   gets  
    lea    rsi, [rsp+12]  
    mov     rdi, rbx  
    call   strcmp  
    test   eax, eax  
    sete   al  
    movzx  eax, al  
    add    rsp, 16  
    pop    rbx  
    ret  
  
main:  
401188  sub     rsp, 8  
40118c  mov     edi, OFFSET FLAT:.LC0  
401191  mov     eax, 0  
401196  call   printf  
40119b  jmp     .L4  
.L5:  
40119d  mov     edi, OFFSET FLAT:.LC1  
4011a2  mov     eax, 0  
4011a7  call   printf  
.L4:  
4011ac  mov     edi, OFFSET FLAT:.LC2  
4011b1  call   authenticate  
4011b6  test   eax, eax  
4011b8  je     .L5  
4011ba  mov     edi, OFFSET FLAT:.LC3  
4011bf  call   puts  
4011c4  mov     eax, 0  
4011c9  add    rsp, 8  
4011cd  ret
```

Practice with Buffer Overflow Exploits

Construct a buffer overflow exploit that will successfully overwrite the return address of the authenticate function with the address of the .LC0 label.



```
int authenticate(char *buf) {
    char buf[4];
    gets(buf);
    int correct = strcmp(buf, "123456");
    return correct;
}

int main(int argc, char *argv[]) {
    char *pw = "123456";
    printf("Enter password: ");
    while (!authenticate(pw)) {
        printf("Incorrect password. Try again: ");
    }
    printf("You are now logged in\n");
    return 0;
}
```

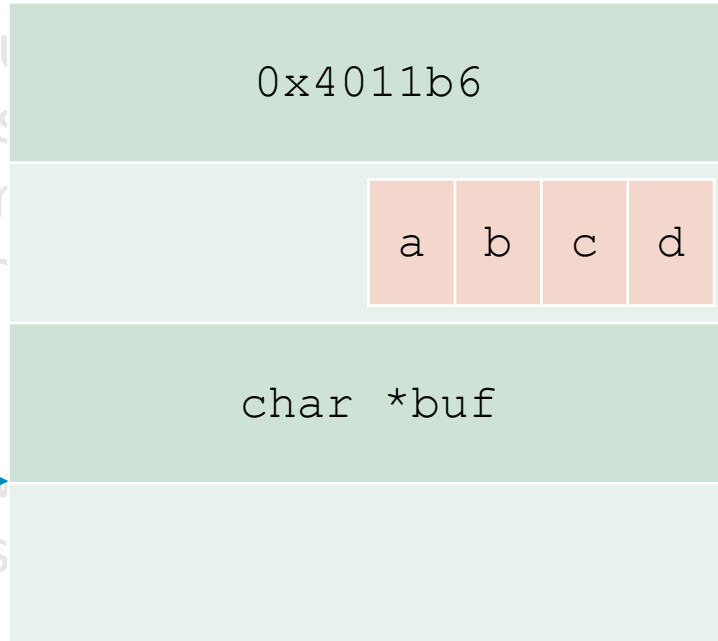
```
.LC0: .string "Enter your password: "
.LC1: .string "Incorrect. Try again: "
.LC2: .string "123456"
.LC3: .string "You are now logged in"
authenticate:
    push    rbx
    sub     rsp, 16
    mov     rbx, rdi
    lea    rdi, [rsp+12]
    mov     eax, 0
    call   gets
    lea    rsi, [rsp+12]
    mov     rdi, rbx
    call   strcmp
    test   eax, eax
    sete   al
    movzx  eax, al
    add    rsp, 16
    pop    rbx
    ret

main:
401188 sub     rsp, 8
40118c mov     edi, OFFSET FLAT:.LC0
401191 mov     eax, 0
401196 call   printf
40119b jmp    .L4
.L5:
40119d mov     edi, OFFSET FLAT:.LC1
4011a2 mov     eax, 0
4011a7 call   printf
.L4:
4011ac mov     edi, OFFSET FLAT:.LC2
4011b1 call   authenticate
4011b6 test   eax, eax
4011b8 je     .L5
4011ba mov     edi, OFFSET FLAT:.LC3
4011bf call   puts
4011c4 mov     eax, 0
4011c9 add    rsp, 8
4011cd ret
```

2. What value will you overwrite the return address with?

Practice with Buffer Overflow Exploits

Construct a buffer overflow exploit that will successfully overwrite the return address of the authenticate function with the address of the main function.



```
int authenticate(char *buf) {
    char buf[4];
    gets(buf);
    int correct = strcmp(buf, "123456");
    return correct;
}

int main(int argc, char *argv[]) {
    char *pw = "123456";
    printf("Enter your password: ");
    while (!authenticate(pw)) {
        printf("Incorrect. Try again: ");
    }
    printf("You are now logged in\n");
    return 0;
}
```

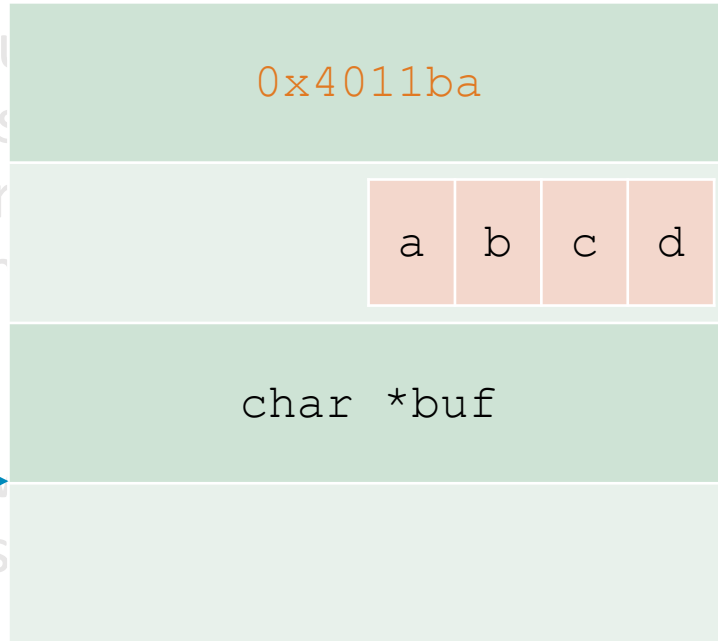
```
.LC0: .string "Enter your password: "
.LC1: .string "Incorrect. Try again: "
.LC2: .string "123456"
.LC3: .string "You are now logged in"
authenticate:
    push    rbx
    sub     rsp, 16
    mov     rbx, rdi
    lea    rdi, [rsp+12]
    mov     eax, 0
    call   gets
    lea    rsi, [rsp+12]
    mov     rdi, rbx
    call   strcmp
    test   eax, eax
    sete   al
    movzx  eax, al
    add    rsp, 16
    pop    rbx
    ret

main:
401188 sub     rsp, 8
40118c mov     edi, OFFSET FLAT:.LC0
401191 mov     eax, 0
401196 call   printf
40119b jmp    .L4
.L5:
40119d mov     edi, OFFSET FLAT:.LC1
4011a2 mov     eax, 0
4011a7 call   printf
.L4:
4011ac mov     edi, OFFSET FLAT:.LC2
4011b1 call   authenticate
4011b6 test   eax, eax
4011b8 je     .L5
4011ba mov     edi, OFFSET FLAT:.LC3
4011bf call   puts
4011c4 mov     eax, 0
4011c9 add    rsp, 8
4011cd ret
```

2. What value will you overwrite the return address with?

Practice with Buffer Overflow Exploits

Construct a buffer overflow exploit that will successfully overwrite the return address of the authenticate function with the address of the .LC3 label.



```
int authenticate(char *buf) {
    char buf[4];
    gets(buf);
    int correct = strcmp(buf, "123456");
    return correct;
}

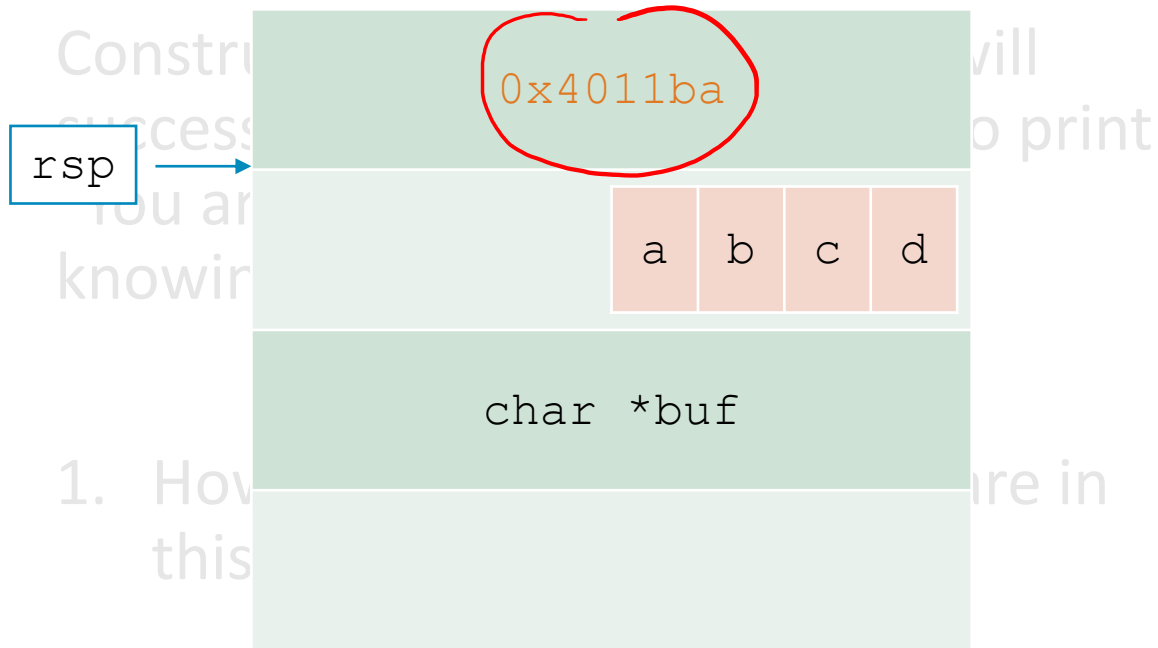
int main(int argc, char *argv[]) {
    char *pw = "123456";
    printf("Enter password: ");
    while (!authenticate(pw)) {
        printf("Incorrect password. Try again: ");
    }
    printf("You are now logged in\n");
    return 0;
}
```

```
.LC0: .string "Enter your password: "
.LC1: .string "Incorrect. Try again: "
.LC2: .string "123456"
.LC3: .string "You are now logged in"
authenticate:
    push    rbx
    sub     rsp, 16
    mov     rbx, rdi
    lea    rdi, [rsp+12]
    mov     eax, 0
    call   gets
    lea    rsi, [rsp+12]
    mov     rdi, rbx
    call   strcmp
    test   eax, eax
    sete   al
    movzx  eax, al
    add    rsp, 16
    pop    rbx
    ret

main:
401188 sub     rsp, 8
40118c mov     edi, OFFSET FLAT:.LC0
401191 mov     eax, 0
401196 call   printf
40119b jmp    .L4
.L5:
40119d mov     edi, OFFSET FLAT:.LC1
4011a2 mov     eax, 0
4011a7 call   printf
.L4:
4011ac mov     edi, OFFSET FLAT:.LC2
4011b1 call   authenticate
4011b6 test   eax, eax
4011b8 je     .L5
4011ba mov     edi, OFFSET FLAT:.LC3
4011bf call   puts
4011c4 mov     eax, 0
4011c9 add    rsp, 8
4011cd ret
```

2. What value will you overwrite the return address with?

Practice with Buffer Overflow Exploits



1. How can we overflow this buffer?

2. What value will you overwrite the return address with?

```
int authenticate(char buf[4]) {
    gets(buf);
    int correct = strcmp(buf, "123456");
    return correct;
}

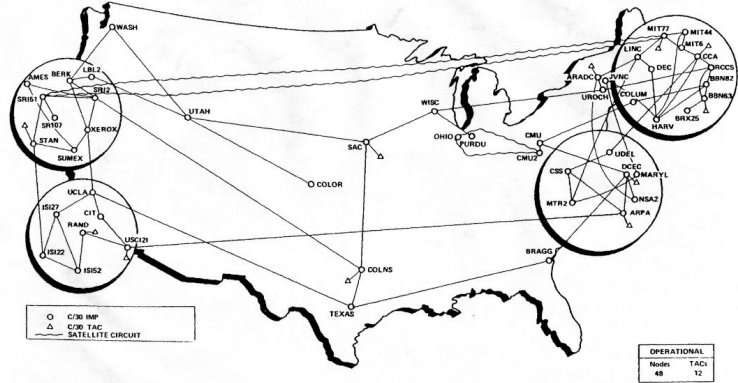
int main(int argc, char * argv[]) {
    char * pw = "123456";
    printf("Enter your password: ");
    while (!authenticate(pw)) {
        printf("Incorrect. Try again: ");
    }
    printf("You are now logged in\n");
    return 0;
}
```

```
.LC0: .string "Enter your password: "
.LC1: .string "Incorrect. Try again: "
.LC2: .string "123456"
.LC3: .string "You are now logged in"
authenticate:
    push    rbp
    sub     rsp, 16
    mov     rbp, rdi
    lea    rdi, [rsp+12]
    mov     eax, 0
    call   gets
    lea    rsi, [rsp+12]
    mov     rdi, rbp
    call   strcmp
    test   eax, eax
    sete   al
    movzx  eax, al
    add    rsp, 16
    pop    rbp
    ret

main:
401188 sub     rsp, 8
40118c mov     edi, OFFSET FLAT:.LC0
401191 mov     eax, 0
401196 call   printf
40119b jmp    .L4
.L5:
40119d mov     edi, OFFSET FLAT:.LC1
4011a2 mov     eax, 0
4011a7 call   printf
.L4:
4011ac mov     edi, OFFSET FLAT:.LC2
4011b1 call   authenticate
4011b6 test   eax, eax
4011b8 je     .L5
4011ba mov     edi, OFFSET FLAT:.LC3
4011bf call   puts
4011c4 mov     eax, 0
4011c9 add    rsp, 8
4011cd ret
```

Buffer Overflow Examples

ARPANET Geographic Map, 31 October 1988



Morris Worm (November 2, 1988), \$100,000-\$10mil, bugs in finger and sendmail

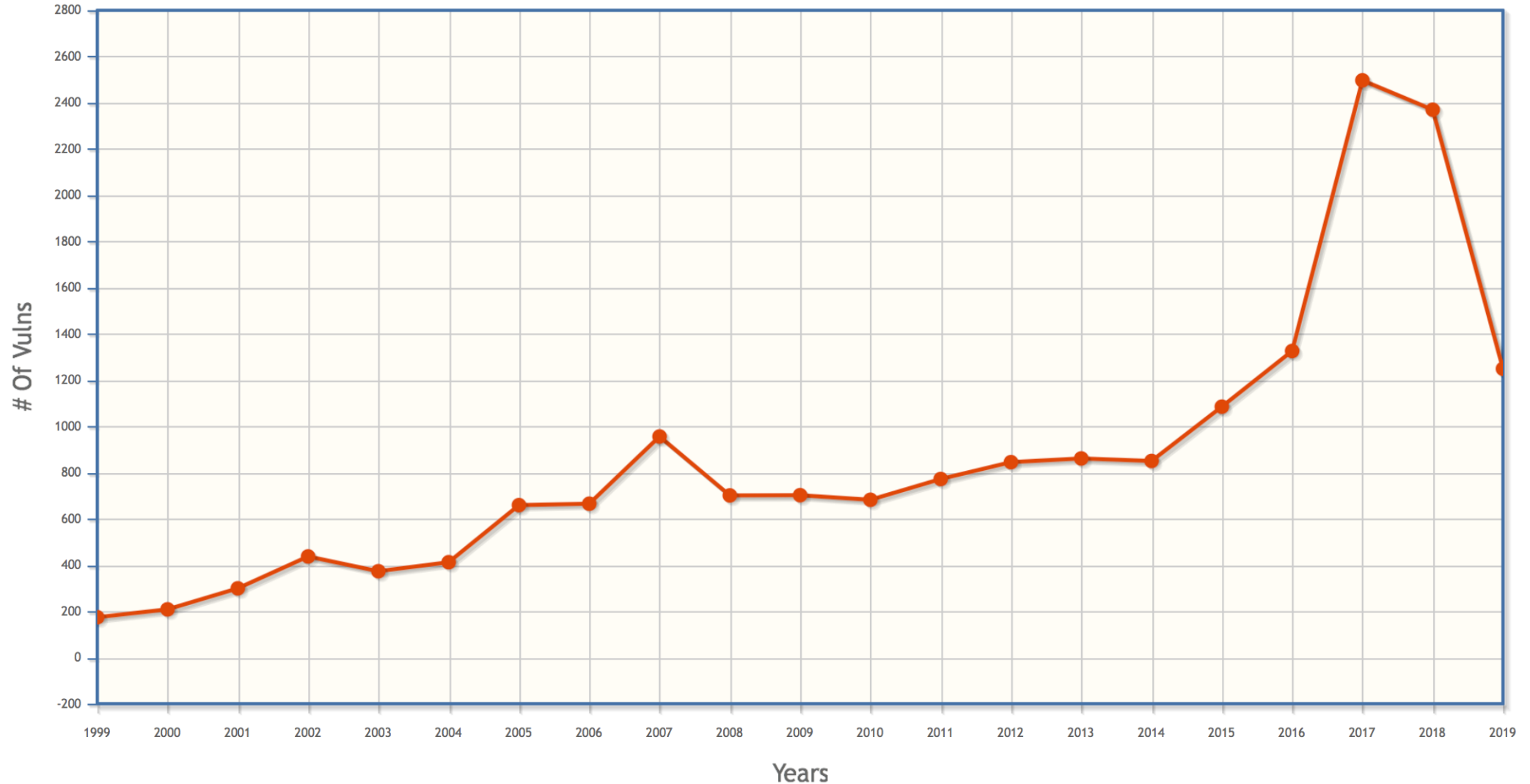


WhatsApp (2019)

Stuxnet (discovered 2010): bug in functions that process files to display icons when USB connected to PC



Buffer Overflow Vulnerabilities



Linus Torvalds: Rust will go into Linux 6.1

Sept. 19, 2022

Defense #1: Avoid Overflow Vulnerabilities

```
// Echo Line  
void echo() {  
    char buf[4]; // Way too small!  
    fgets(buf, 4, stdin);  
    puts(buf);  
}
```

- For example, use library routines that limit string lengths
 - **fgets** instead of **gets**
 - Don't use **scanf** with **%s** conversion specification (use **fgets** to read the string or use **%ns** where **n** is a suitable integer)

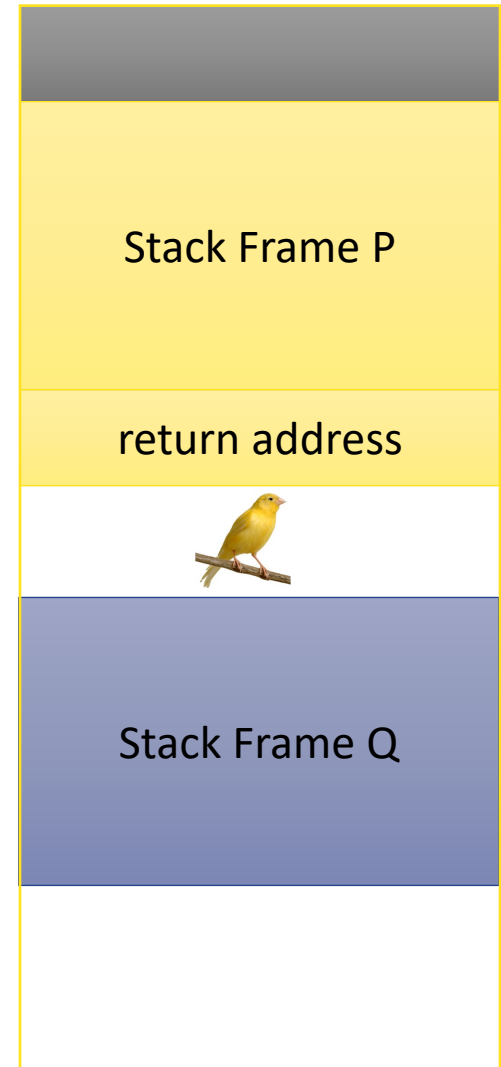
<https://github.com/leafsr/gcc-poison>

- Or use a better high-level language

Defense #2: Compiler checks

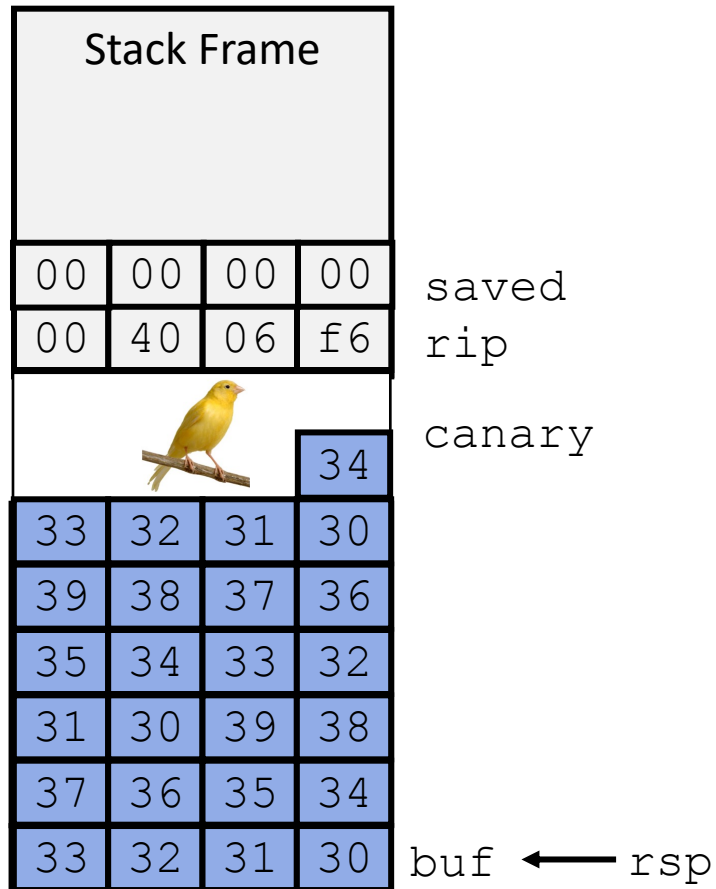
- Idea
 - Place special value (“canary”) on stack just beyond buffer
 - Check for corruption before exiting function
- GCC Implementation
 - -fstack-protector
 - Now the default (disabled earlier)

0x7FFFFFFF



0x00000000

Stack Canaries



authenticate:

```

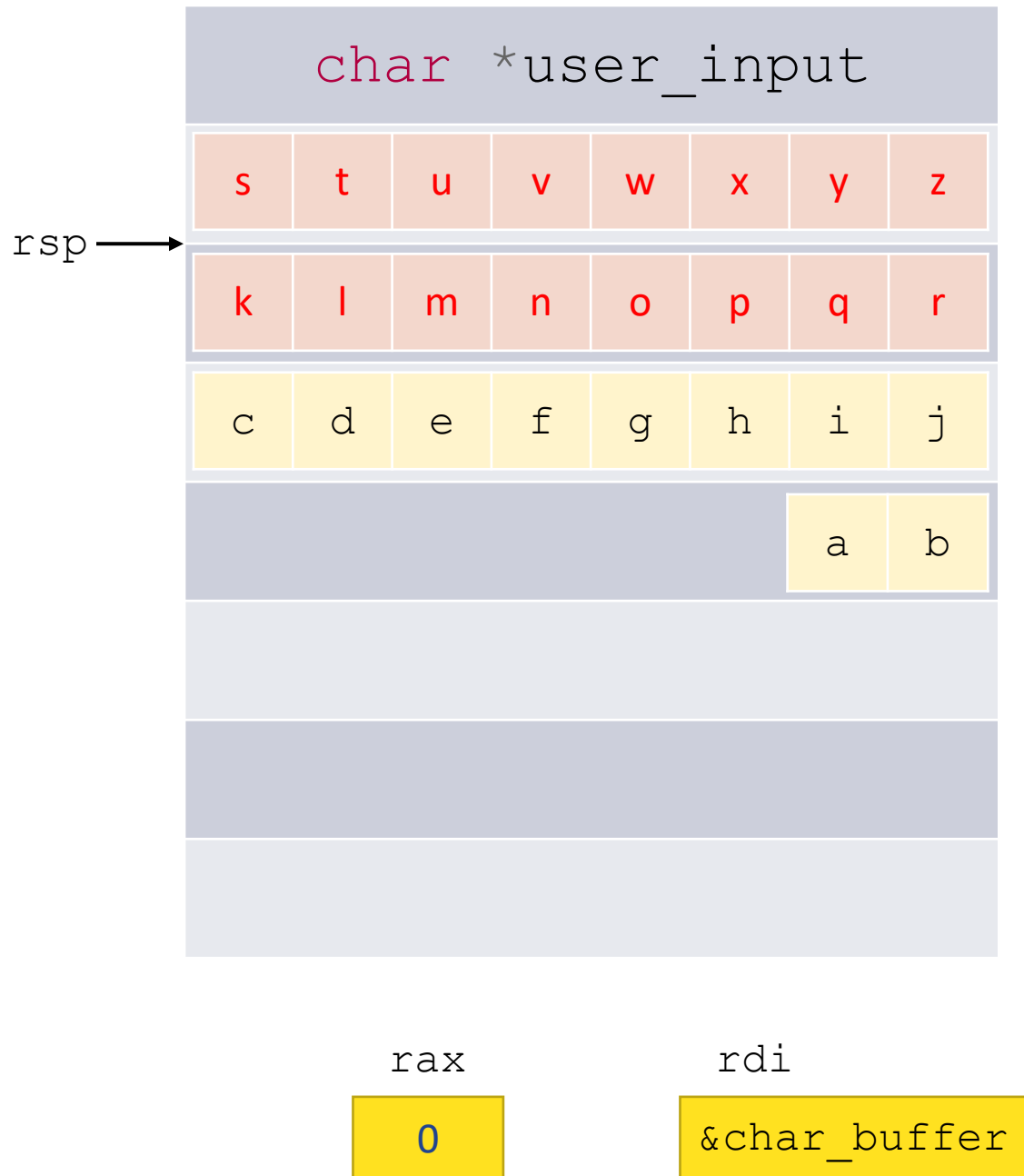
push    rbx
sub     rsp, 16
mov     rbx, rdi
mov     rax, fs:40
mov     [rsp+8], rax
xor     eax, eax
mov     rdi, rsp
call   gets
mov     rsi, rsp
mov     rdi, rbx
call   strcmp
test   eax, eax
sete   al
mov     rdx, [rsp+8]
xor     rdx, fs:40
je     .L2
call   __stack_chk_fail
.L2:
movzb  eax, al
add    rsp, 16
pop    rbx
ret

```

Stack Canaries

Which of the following would make a good stack canary?

1. A secret, constant value
2. A fixed sequence of common terminators (`\0`, EOF, etc.)
3. A random number chosen each time the program is run



```

read_string:
    sub     rsp, 24
    lea    rdi, [rsp+6]
    mov    eax, 0
    call   gets
    mov    eax, 0
    add    rsp, 24
rip → ret

```

```

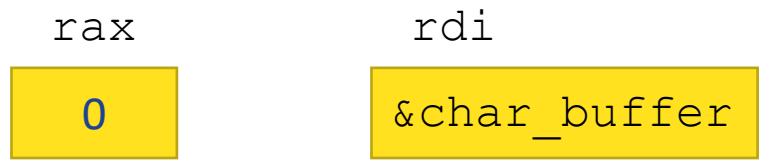
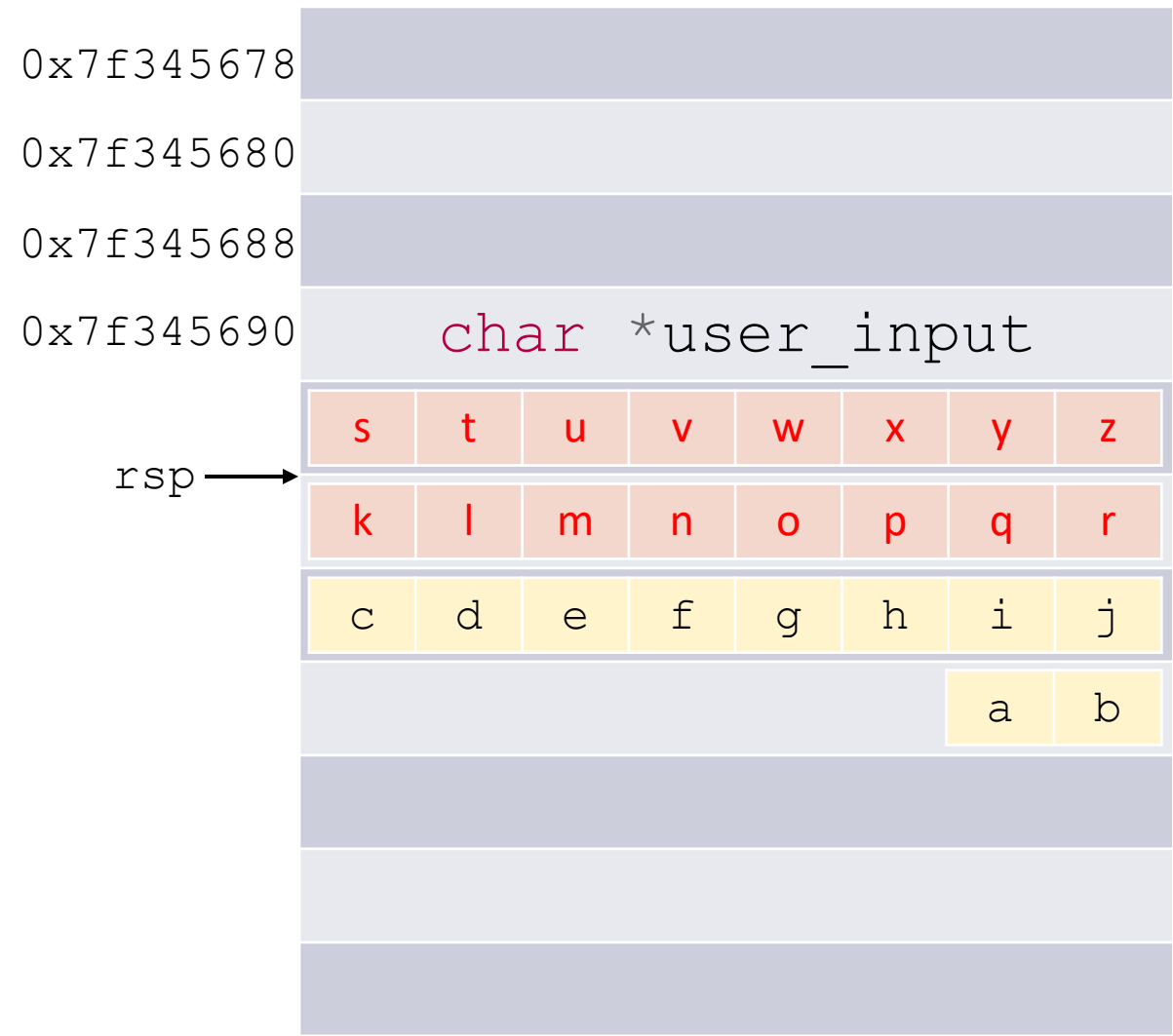
print_string:
    sub     rsp, 8
    call   puts
    add    rsp, 8
    ret

```

```

main:
    sub     rsp, 8
    mov    eax, 0
    call   read_string
    movsx  rdi, eax
    mov    eax, 0
    call   print_string
    mov    eax, 0
    add    rsp, 8
    ret

```



```

read_string:
    sub     rsp, 24
    lea    rdi, [rsp+6]
    mov    eax, 0
    call   gets
    mov    eax, 0
    add    rsp, 24
rip → ret

```

```

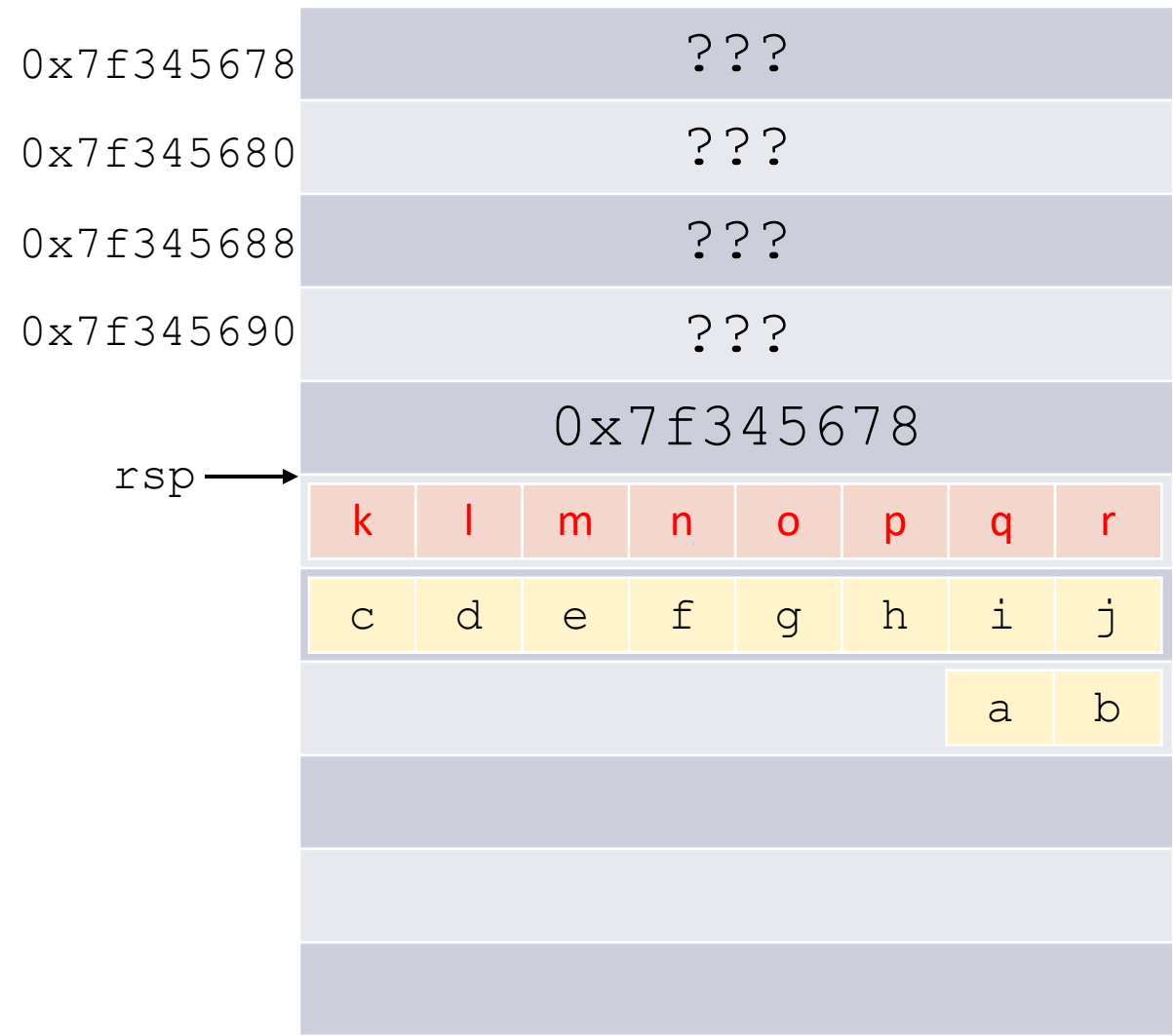
print_string:
    sub     rsp, 8
    call   puts
    add    rsp, 8
    ret

```

```

main:
    sub     rsp, 8
    mov    eax, 0
    call   read_string
    movsx  rdi, eax
    mov    eax, 0
    call   print_string
    mov    eax, 0
    add    rsp, 8
    ret

```



```

read_string:
    sub    rsp, 24
    lea   rdi, [rsp+6]
    mov   eax, 0
    call  gets
    mov   eax, 0
    add   rsp, 24
rip → ret

```

```

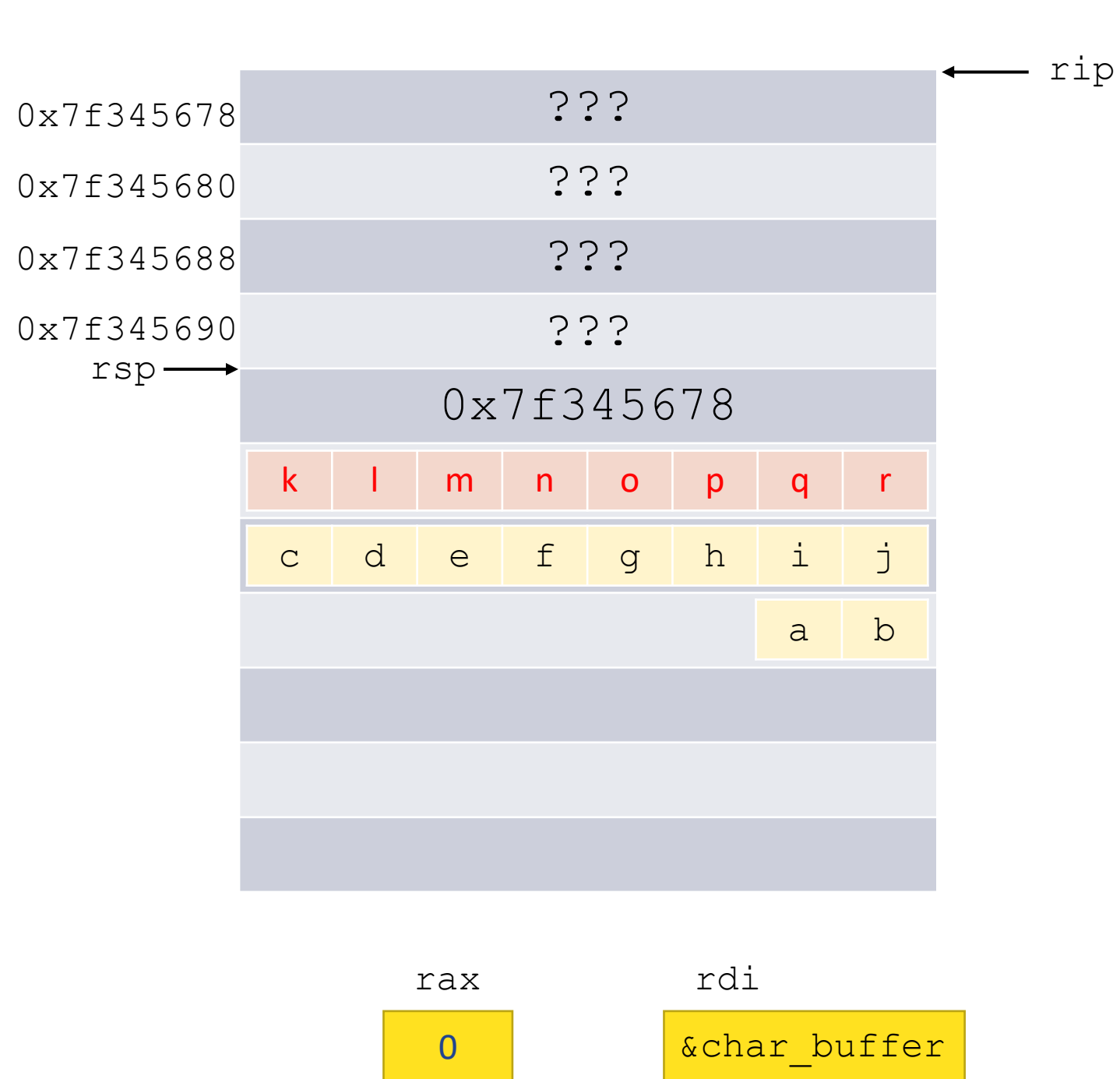
print_string:
    sub    rsp, 8
    call  puts
    add   rsp, 8
    ret

```

```

main:
    sub    rsp, 8
    mov   eax, 0
    call  read_string
    movsx rdi, eax
    mov   eax, 0
    call  print_string
    mov   eax, 0
    add   rsp, 8
    ret

```

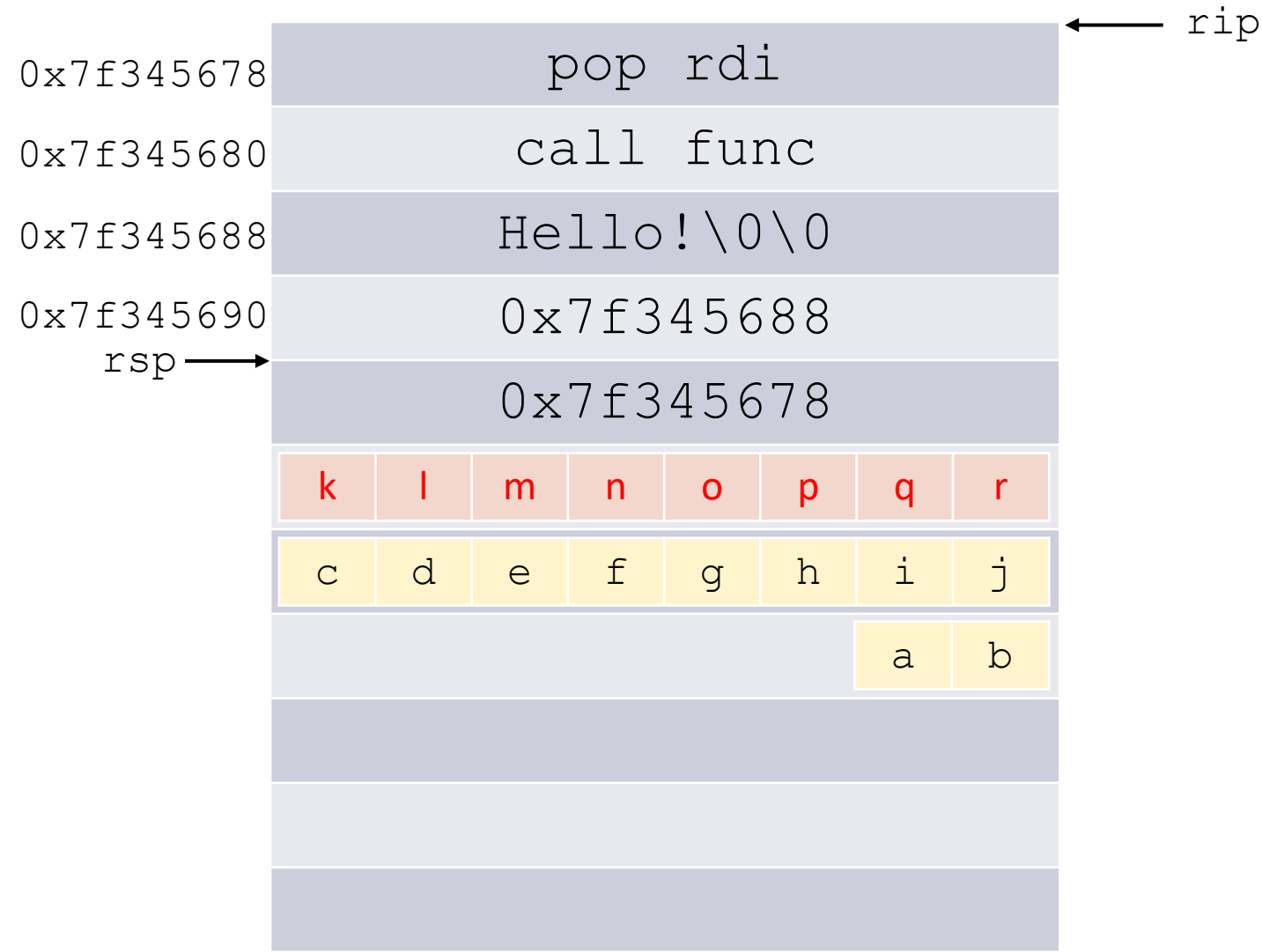
```

read_string:
    sub    rsp, 24
    lea   rdi, [rsp+6]
    mov   eax, 0
    call  gets
    mov   eax, 0
    add   rsp, 24
    ret

print_string:
    sub   rsp, 8
    call  puts
    add   rsp, 8
    ret

main:
    sub   rsp, 8
    mov   eax, 0
    call  read_string
    movsx rdi, eax
    mov   eax, 0
    call  print_string
    mov   eax, 0
    add   rsp, 8
    ret

```



```

read_string:
    sub    rsp, 24
    lea   rdi, [rsp+6]
    mov   eax, 0
    call  gets
    mov   eax, 0
    add   rsp, 24
    ret
  
```

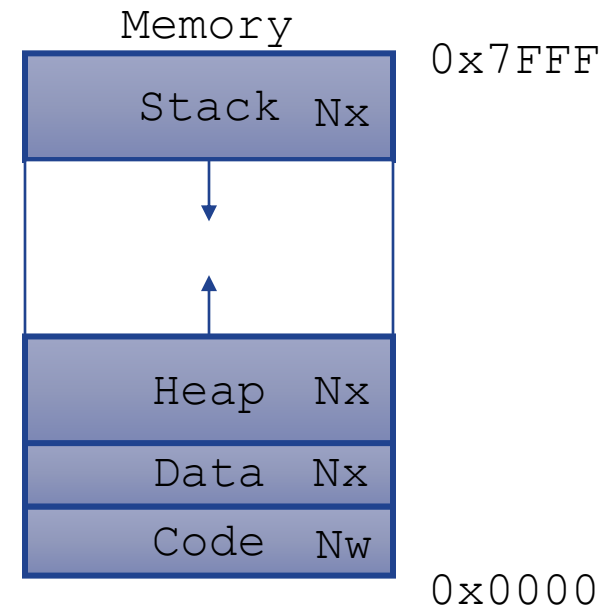
```

print_string:
    sub    rsp, 8
    call  puts
    add   rsp, 8
    ret
  
```

```

main:
    sub    rsp, 8
    mov   eax, 0
    call  read_string
    movsx rdi, eax
    mov   eax, 0
    call  print_string
    mov   eax, 0
    add   rsp, 8
    ret
  
```













Defense #3: Memory Tagging



Code Reuse Attacks

- Key idea: execute instructions that already exist
- Defeats memory tagging defenses
- Examples:
 1. return to a function or line in the current program
 2. return to a library function (e.g., return-into-libc)
 3. return to some other instruction (return-oriented programming)

Return-into-libc

Sr.No.	Function & Description
1	double atof(const char *str)  Converts the string pointed to, by the argument <i>str</i> to a floating-point number (type double).
2	int atoi(const char *str)  Converts the string pointed to, by the argument <i>str</i> to an integer (type int).
3	long int atol(const char *str)  Converts the string pointed to, by the argument <i>str</i> to a long integer (type long int).
8	void free(void *ptr)  Deallocates the memory previously allocated by a call to <i>calloc</i> , <i>malloc</i> , or <i>realloc</i> .
9	void *malloc(size_t size)  Allocates the requested memory and returns a pointer to it.
10	void *realloc(void *ptr, size_t size)  Attempts to resize the memory block pointed to by <i>ptr</i> that was previously allocated with a call to <i>malloc</i> or <i>calloc</i> .
15	int system(const char *string)  The command specified by <i>string</i> is passed to the host environment to be executed by the command processor.
16	void *bsearch(const void *key, const void *base, size_t nitems, size_t size, int (*compar)(const void *, const void *))  Performs a binary search.
17	void qsort(void *base, size_t nitems, size_t size, int (*compar)(const void *, const void*))  Sorts an array.
18	int abs(int x)  Returns the absolute value of <i>x</i> .
22	int rand(void)  Returns a pseudo-random number in the range of 0 to <i>RAND_MAX</i> .
23	void srand(unsigned int seed)  This function seeds the random number generator used by the function rand .

Defense #4: ASCII Armoring

- Make sure all system library addresses contain a null byte (0x00).
- Can be done by placing this code in the first 0x01010101 bytes of memory

Properties of x86-64 Assembly

- Lots of instructions
- Variable length instructions
- Not word aligned
- Dense instruction set → most bytes encode an actual instruction

Gadgets

```
void setval(unsigned *p)
{
    *p = 3347663060u; // 0xC78948D4
}
```

```
<setval>:
4004d9: c7 07 d4 48 89 c7 mov [rdi], 0xC78948D4
4004df: c3 ret
```

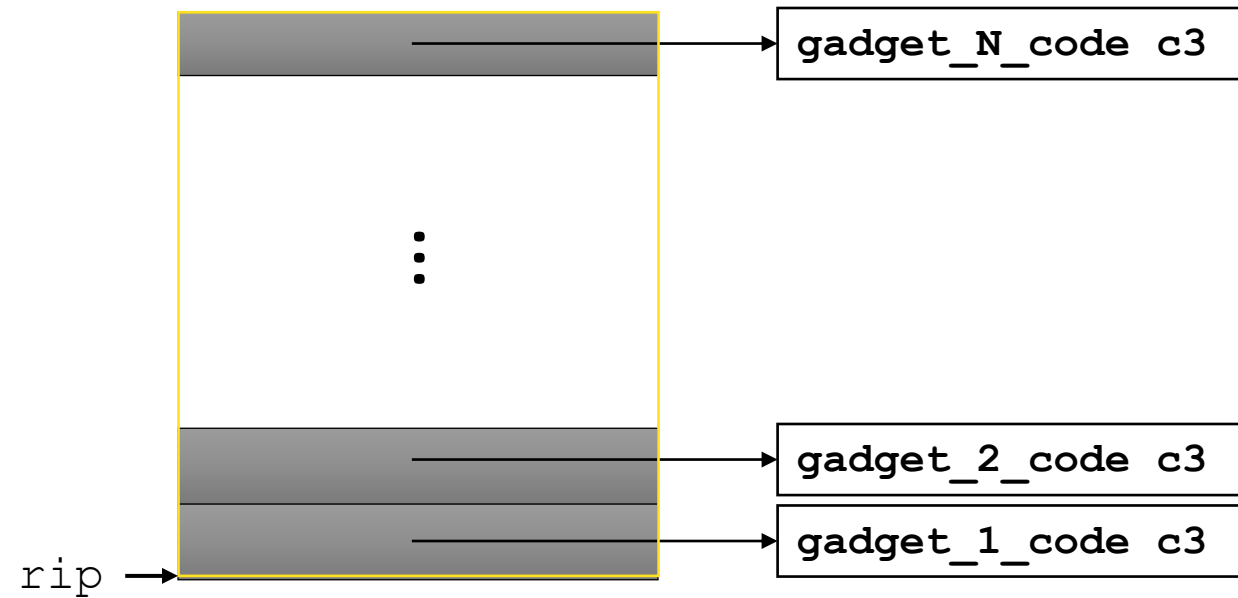
gadget address: 0x4004dc

encodes: mov rdi, rax
ret

Return-oriented Programming

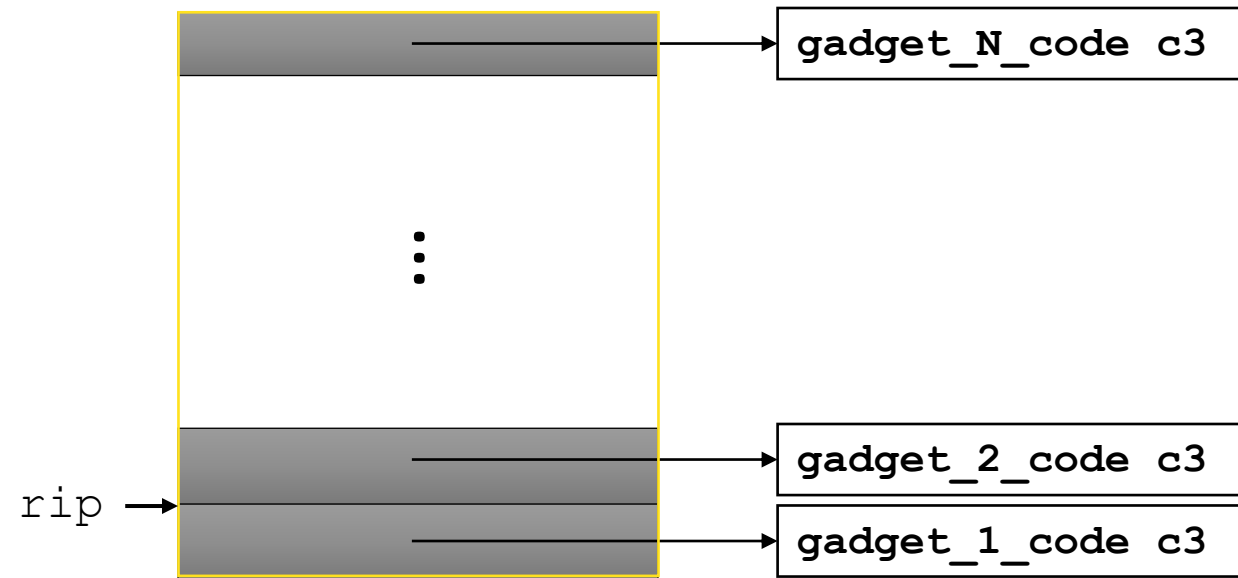


Return-oriented Programming



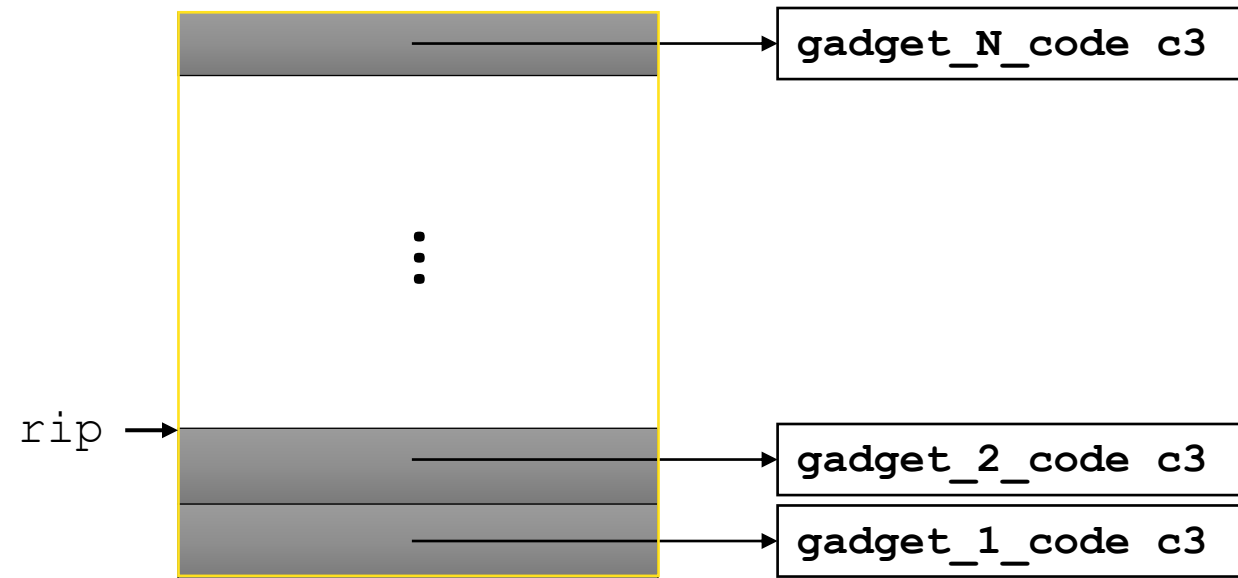
Final ret in each gadget sets pc (`rip`) to beginning of next gadget code

Return-oriented Programming



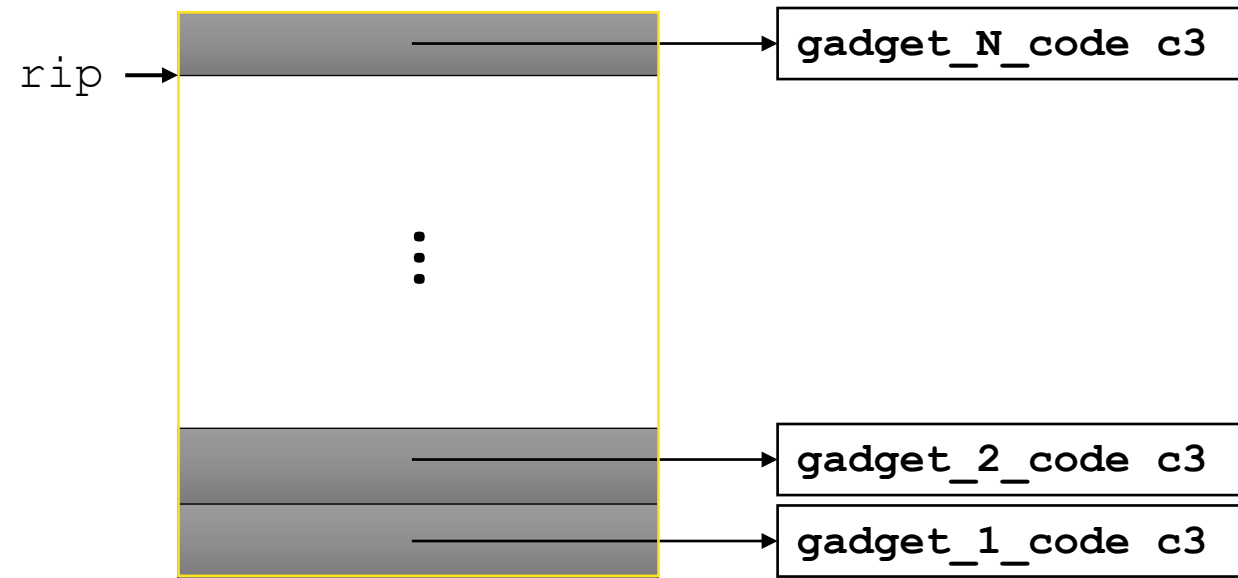
Final ret in each gadget sets pc (`rip`) to beginning of next gadget code

Return-oriented Programming



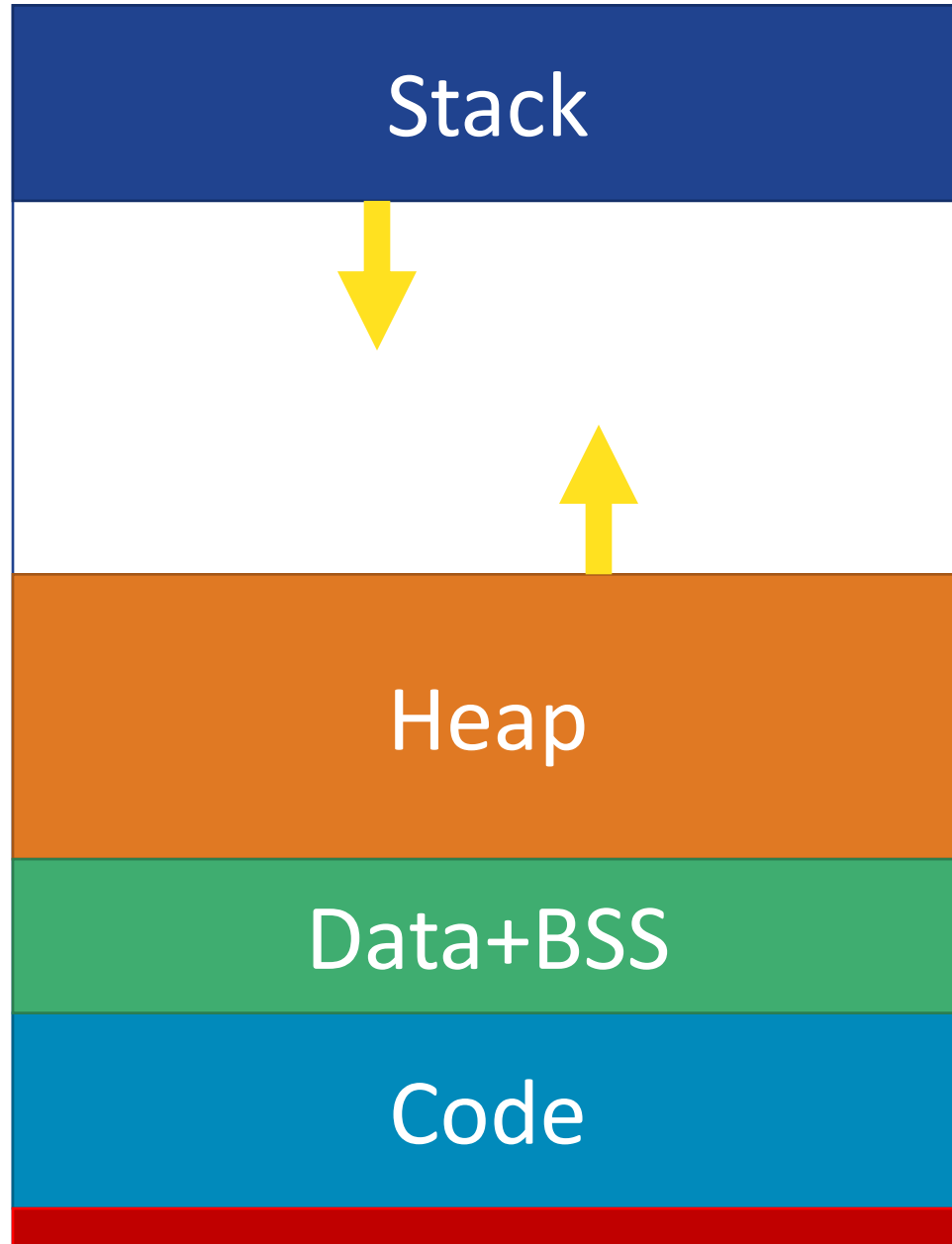
Final ret in each gadget sets pc (`rip`) to beginning of next gadget code

Return-oriented Programming



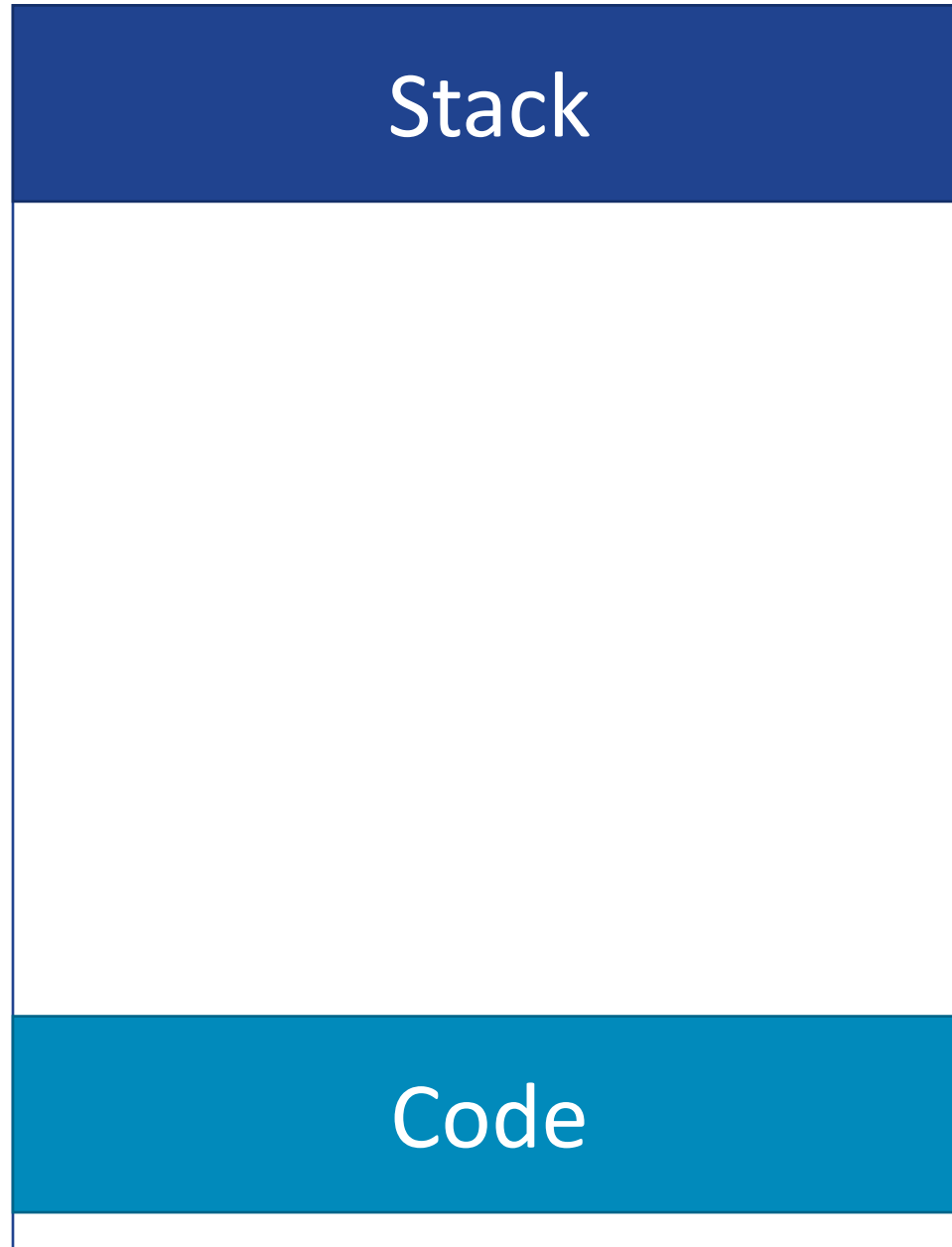
Final ret in each gadget sets pc (`rip`) to beginning of next gadget code

0x7fffffff



0x00000000

0x7fffffffffffffff



0x000000000000

Stack

Code

Stack

7fffffffefa88	...
7fffffffefa80	"\bin\sh\0"
7fffffffefa78	0x1122334455667788
7fffffffefa70	0x7fffffffefa80
7fffffffefa68	0x40042a
7fffffffefa60	0x7fffffffefa80
7fffffffefa58	0x7fffffffefa70
7fffffffefa50	0x4004b8
7fffffffefa48	0x4002a3
7fffffffefa40	0x400660
7fffffffefa38	0x7fffffffefa78
7fffffffefa30	0x3b3b3b3b3b3b3b3b
7fffffffefa28	0x400420
7fffffffefa20	0x40090b
7fffffffefa18	...

Code

Take a guess at the significance of each value on the stack.

Suppose that I manually crafted this smashed stack.

Stack

7fffffffefa88	...
7fffffffefa80	"\bin\sh\0"
7fffffffefa78	0x1122334455667788
7fffffffefa70	0x7fffffffefa80
7fffffffefa68	0x40042a
7fffffffefa60	0x7fffffffefa80
7fffffffefa58	0x7fffffffefa70
7fffffffefa50	0x4004b8
7fffffffefa48	0x4002a3
7fffffffefa40	0x400660
7fffffffefa38	0x7fffffffefa78
7fffffffefa30	0x3b3b3b3b3b3b3b3b
7fffffffefa28	0x400420
7fffffffefa20	0x40090b
7fffffffefa18	...

Code

Take a guess at the significance of each value on the stack.

Stack

7fffffffefa88	...
7fffffffefa80	"\bin\sh\0"
7fffffffefa78	0x1122334455667788
7fffffffefa70	0x7fffffffefa80
7fffffffefa68	0x40042a
7fffffffefa60	0x7fffffffefa80
7fffffffefa58	0x7fffffffefa70
7fffffffefa50	...
7fffffffefa48	...
7fffffffefa40	0x400660
7fffffffefa38	0x7fffffffefa78
7fffffffefa30	0x3b3b3b3b3b3b3b3b
7fffffffefa28	0x400420
7fffffffefa20	0x40090b
7fffffffefa18	...

What do you notice about the code section?

Code

0x4002a3	40 00 F8 C3
...	
0x400420	5F 5E C3
...	
0x40042a	0F 05 C3
...	
0x4004b8	5E 5F C3
...	
...	
0x400660	48 89 06 48 89 C2 C3
...	
...	
0x40090b	48 31 C0 C3
...	

8-Byte Alignment Stack

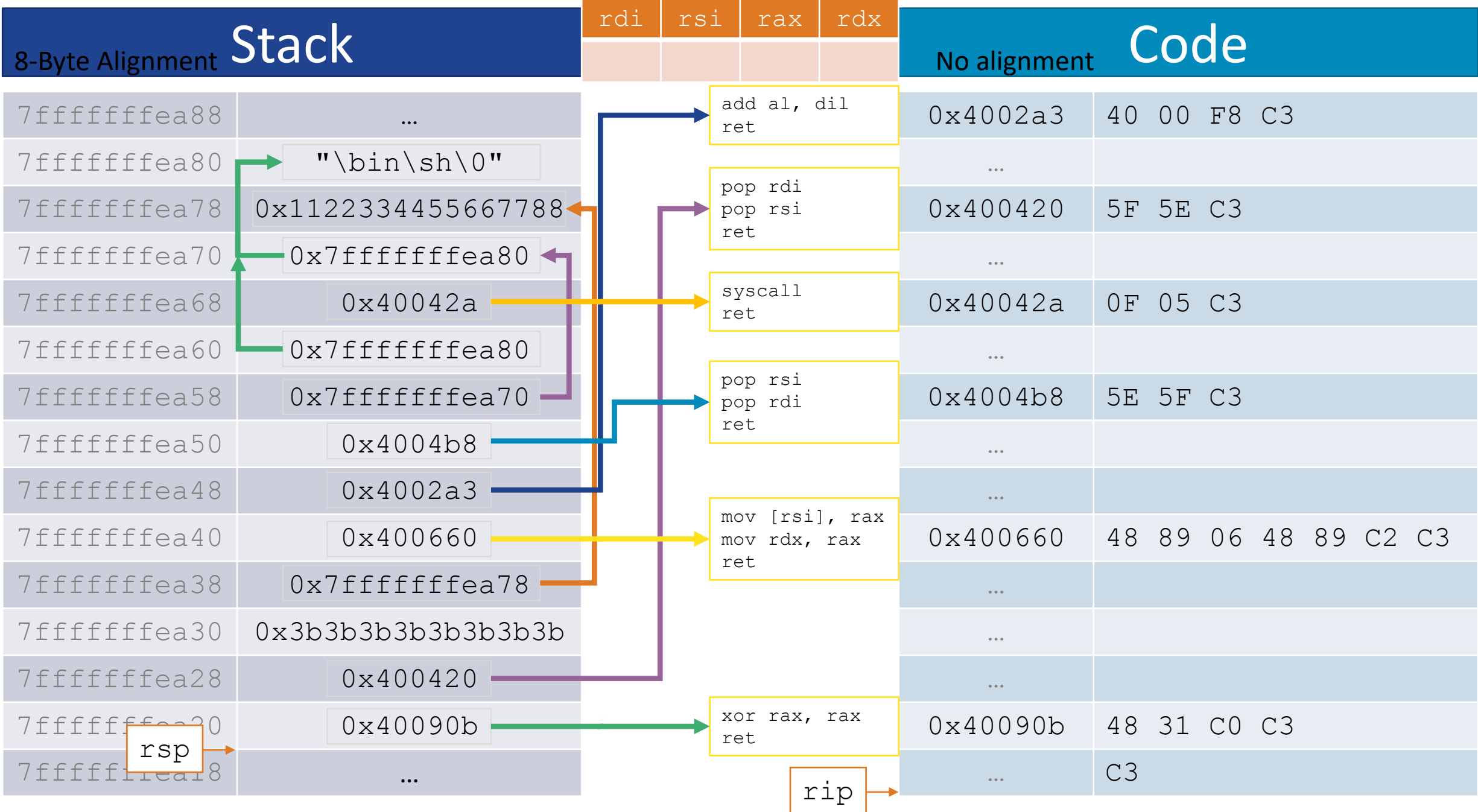
7fffffffefa88	...
7fffffffefa80	"\bin\sh\0"
7fffffffefa78	0x1122334455667788
7fffffffefa70	0x7fffffffefa80
7fffffffefa68	0x40042a
7fffffffefa60	0x7fffffffefa80
7fffffffefa58	0x7fffffffefa70
7fffffffefa50	0x4004b8
7fffffffefa48	0x4002a3
7fffffffefa40	0x400660
7fffffffefa38	0x7fffffffefa78
7fffffffefa30	0x3b3b3b3b3b3b3b3b
7fffffffefa28	0x400420
7fffffffefa20	0x40090b
7fffffffefa18	...

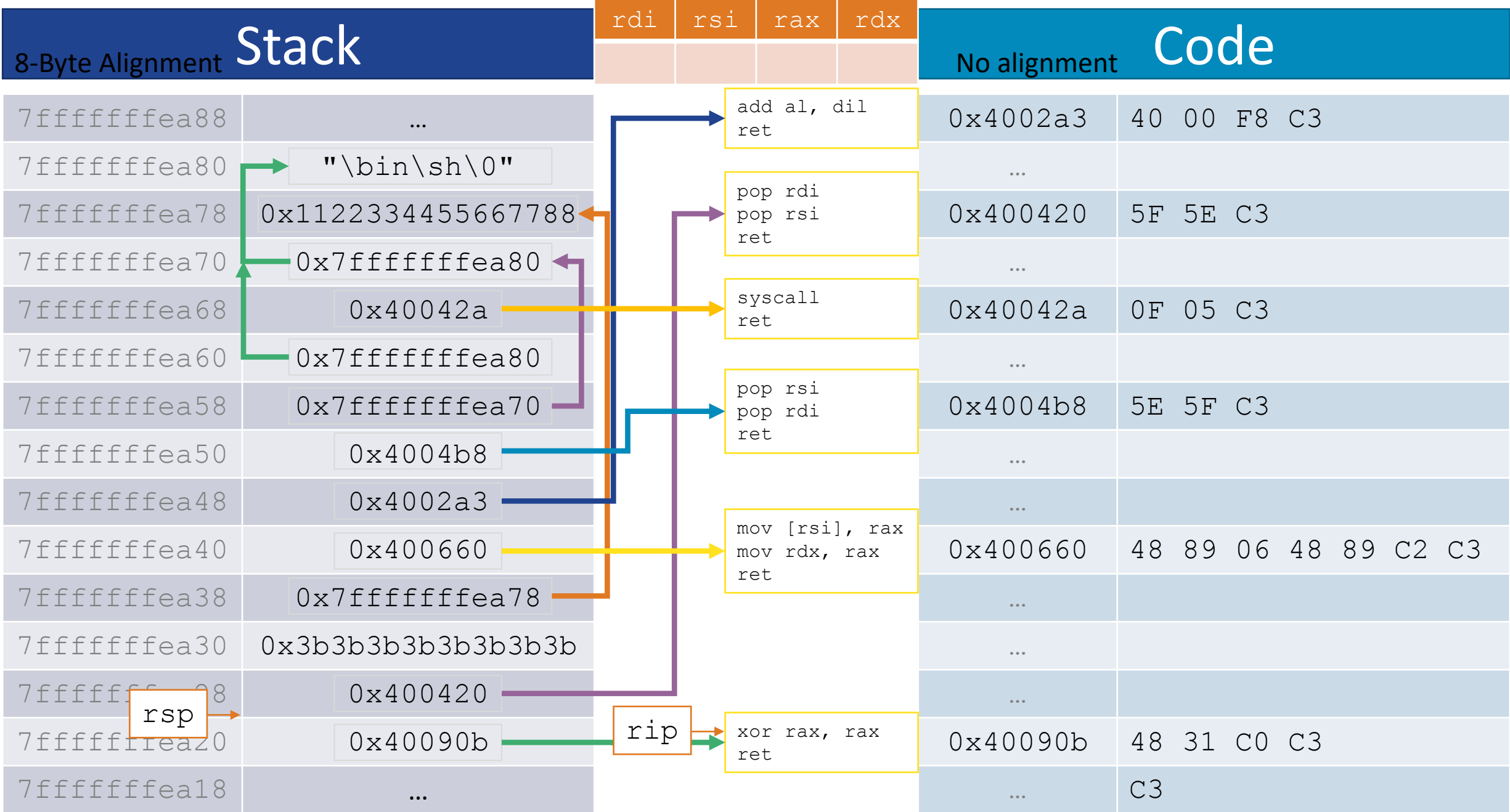
Gadgets

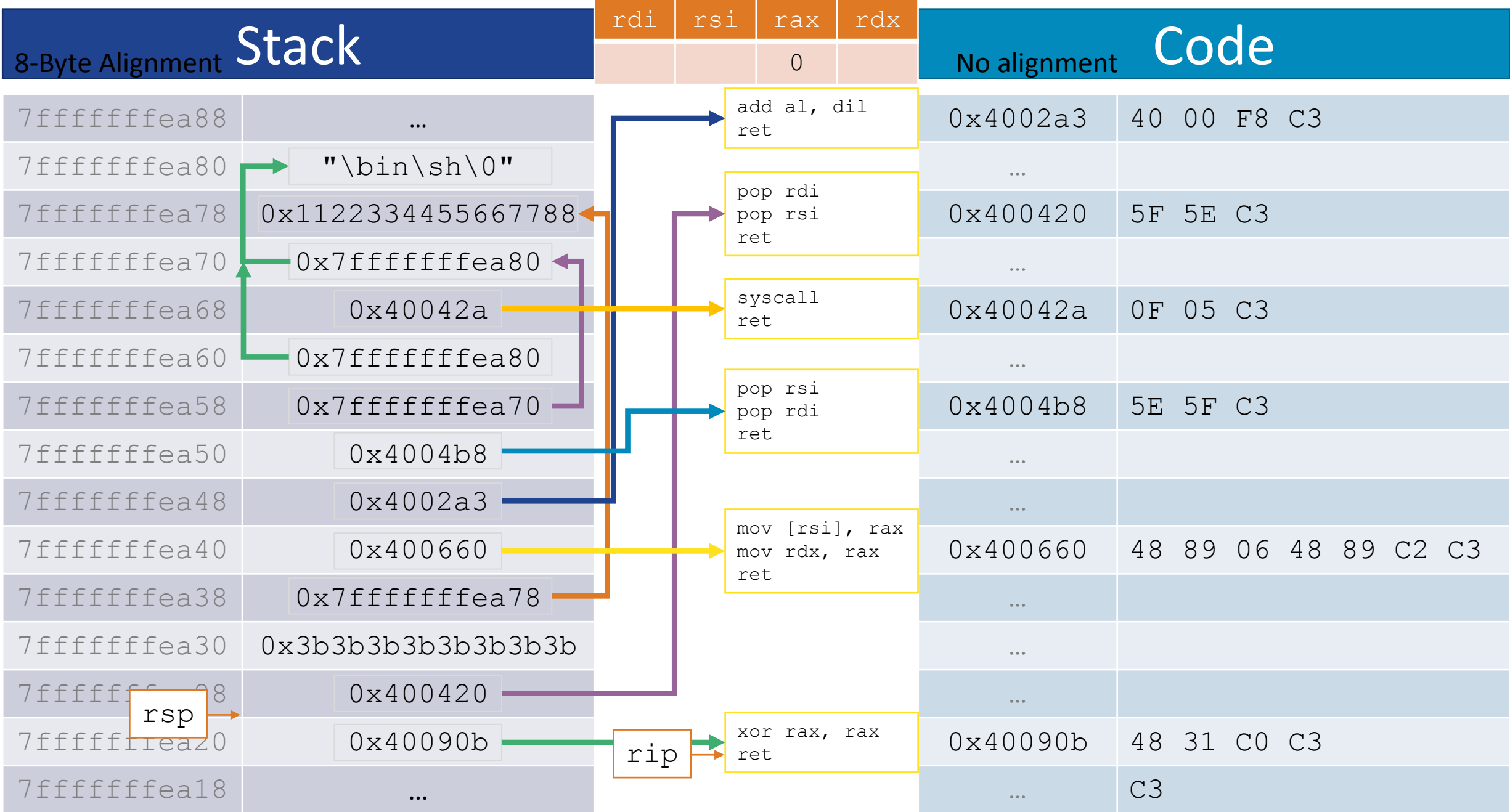
add al, dil ret
pop rdi pop rsi ret
syscall ret
pop rsi pop rdi ret
mov [rsi], rax mov rdx, rax ret
xor rax, rax ret

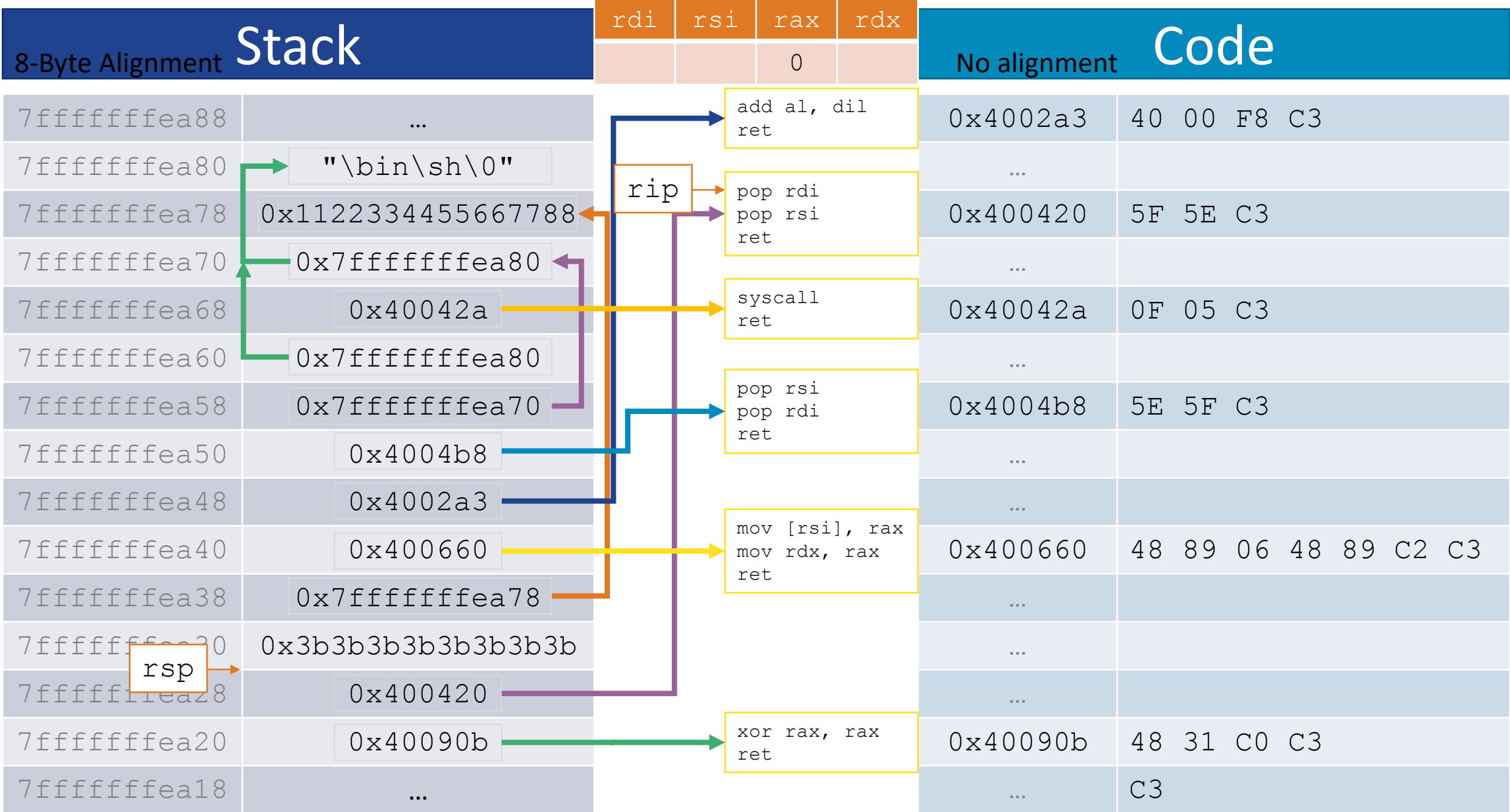
No alignment Code

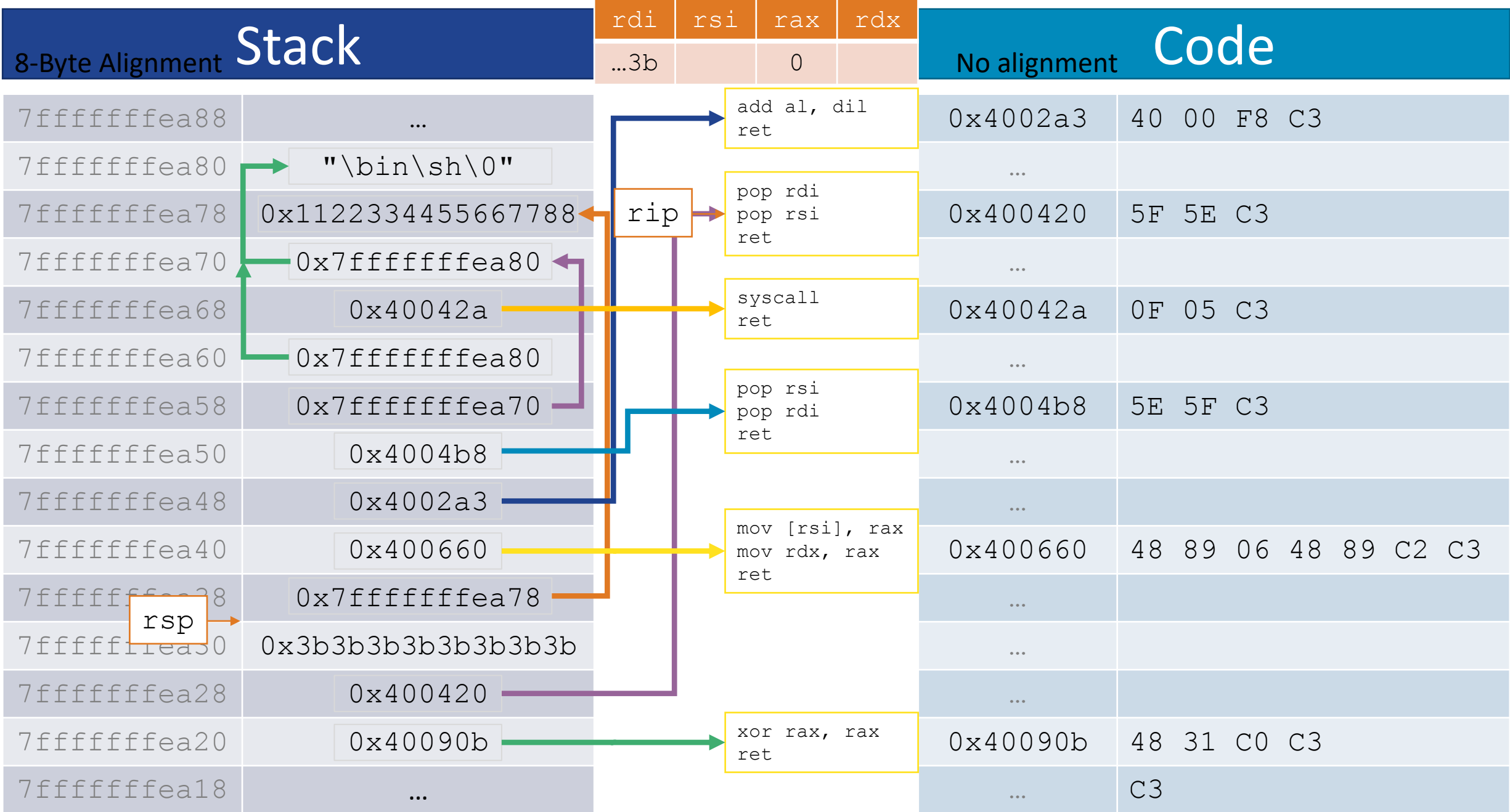
0x4002a3	40 00 F8 C3
...	
0x400420	5F 5E C3
...	
0x40042a	0F 05 C3
...	
0x4004b8	5E 5F C3
...	
...	
0x400660	48 89 06 48 89 C2 C3
...	
...	
0x40090b	48 31 C0 C3
...	

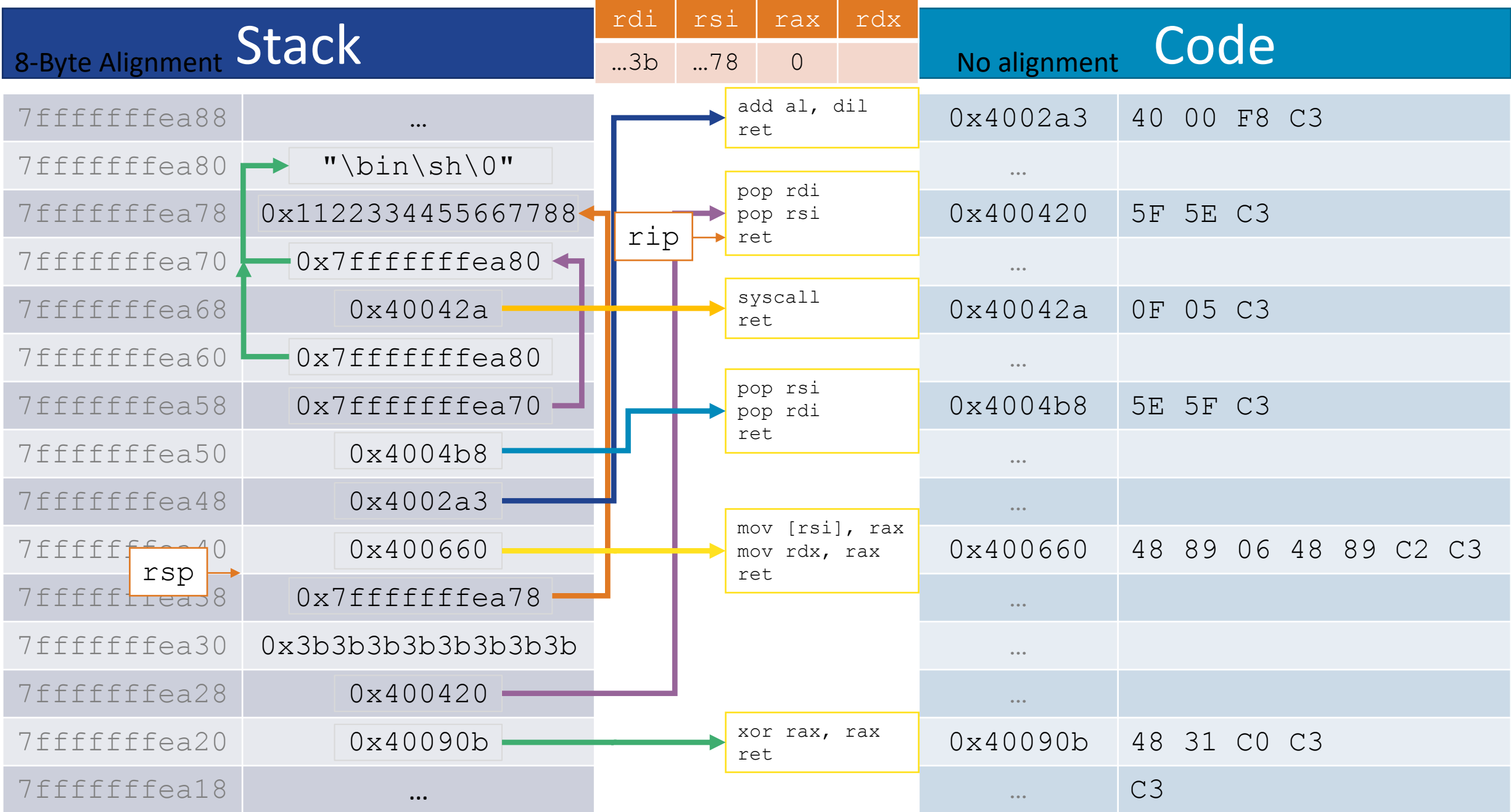


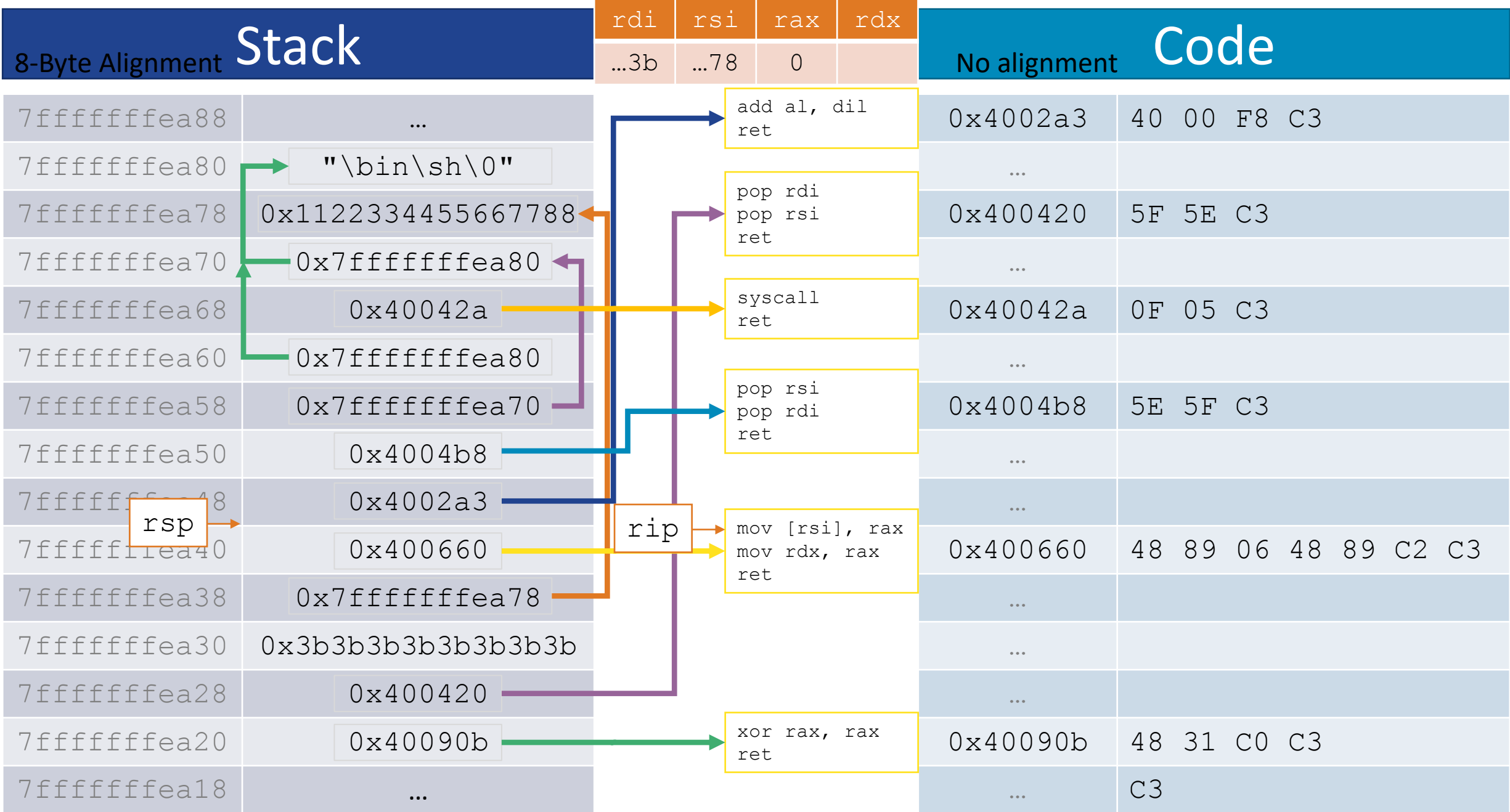


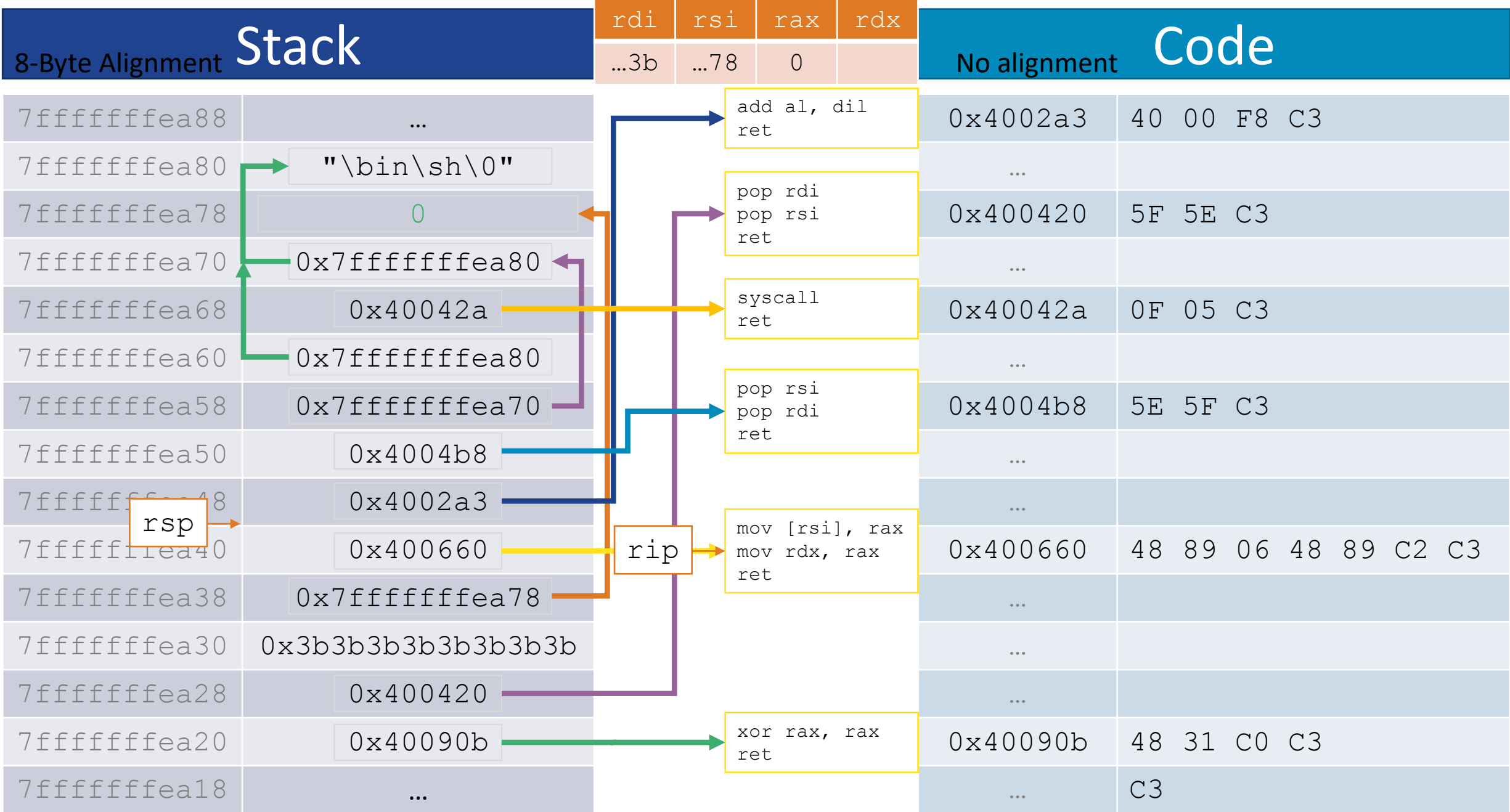


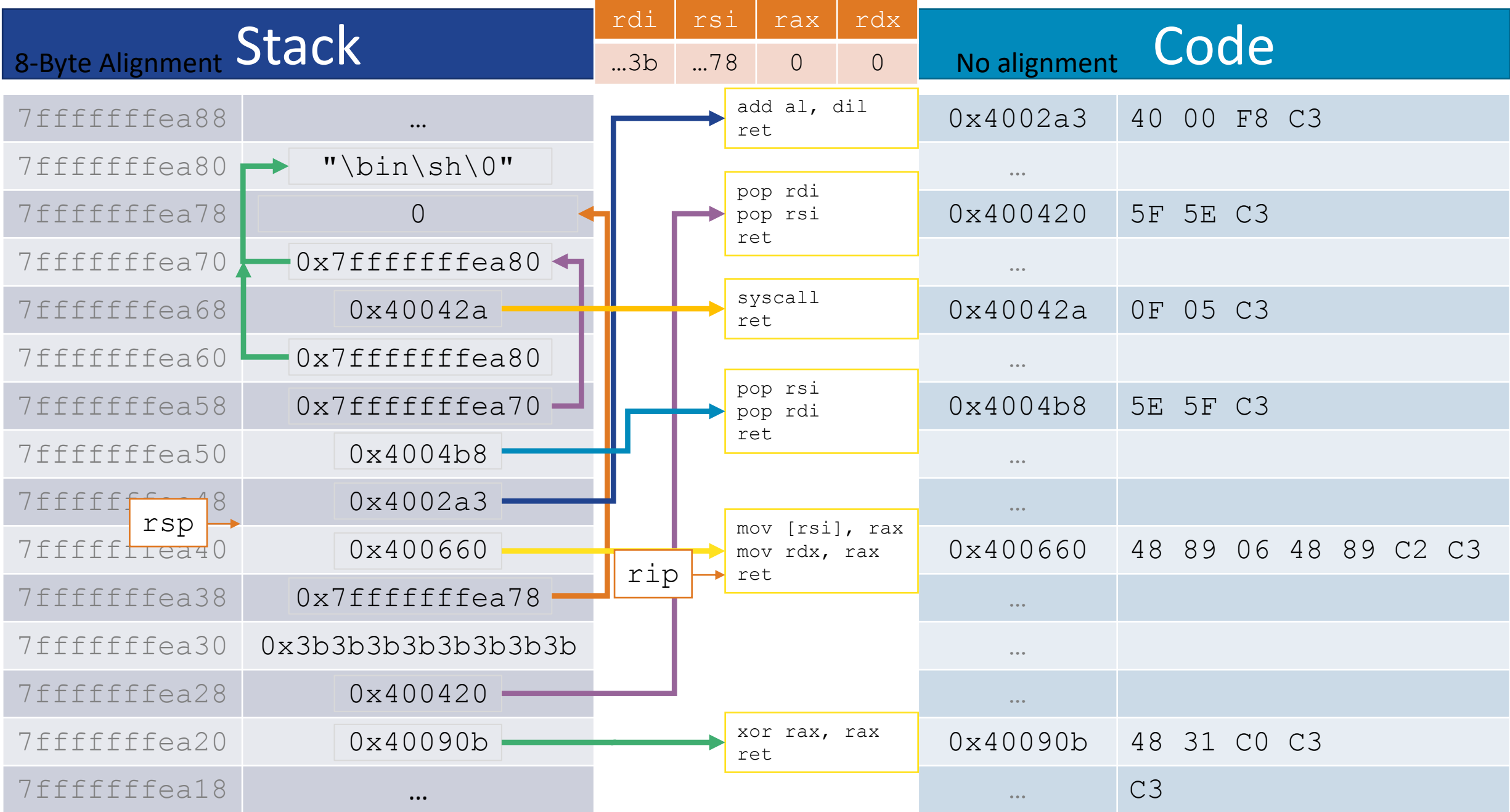


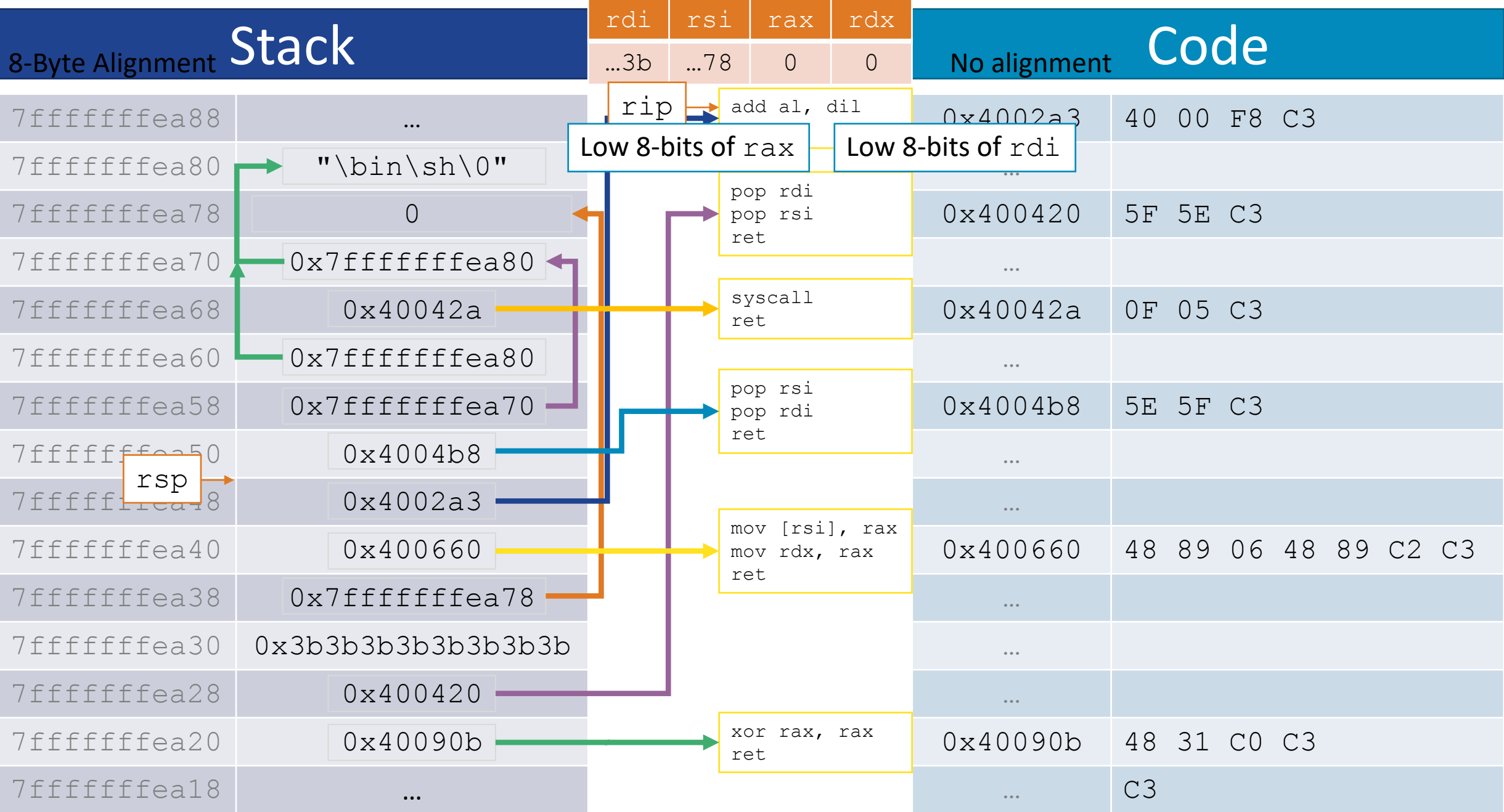


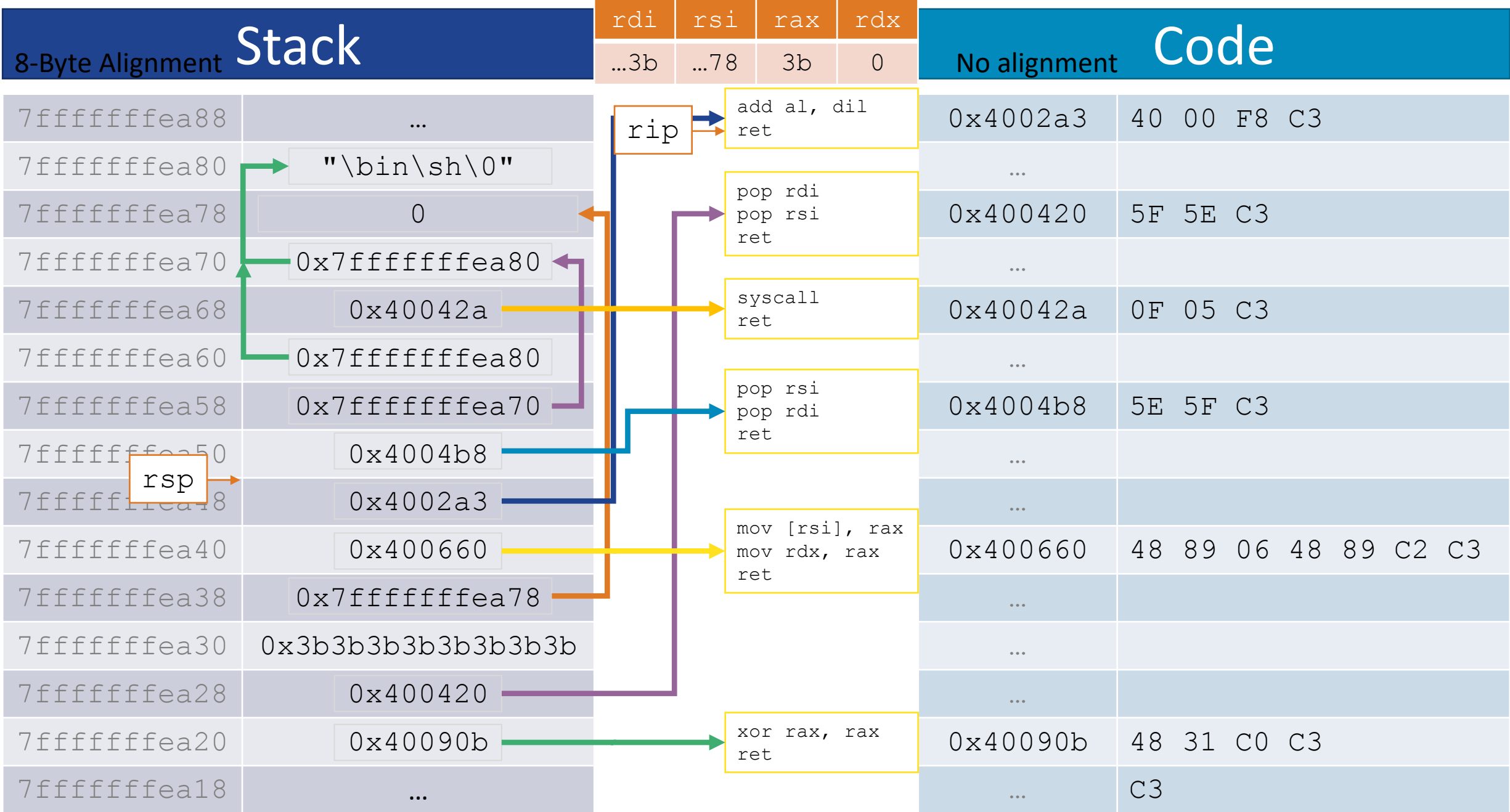


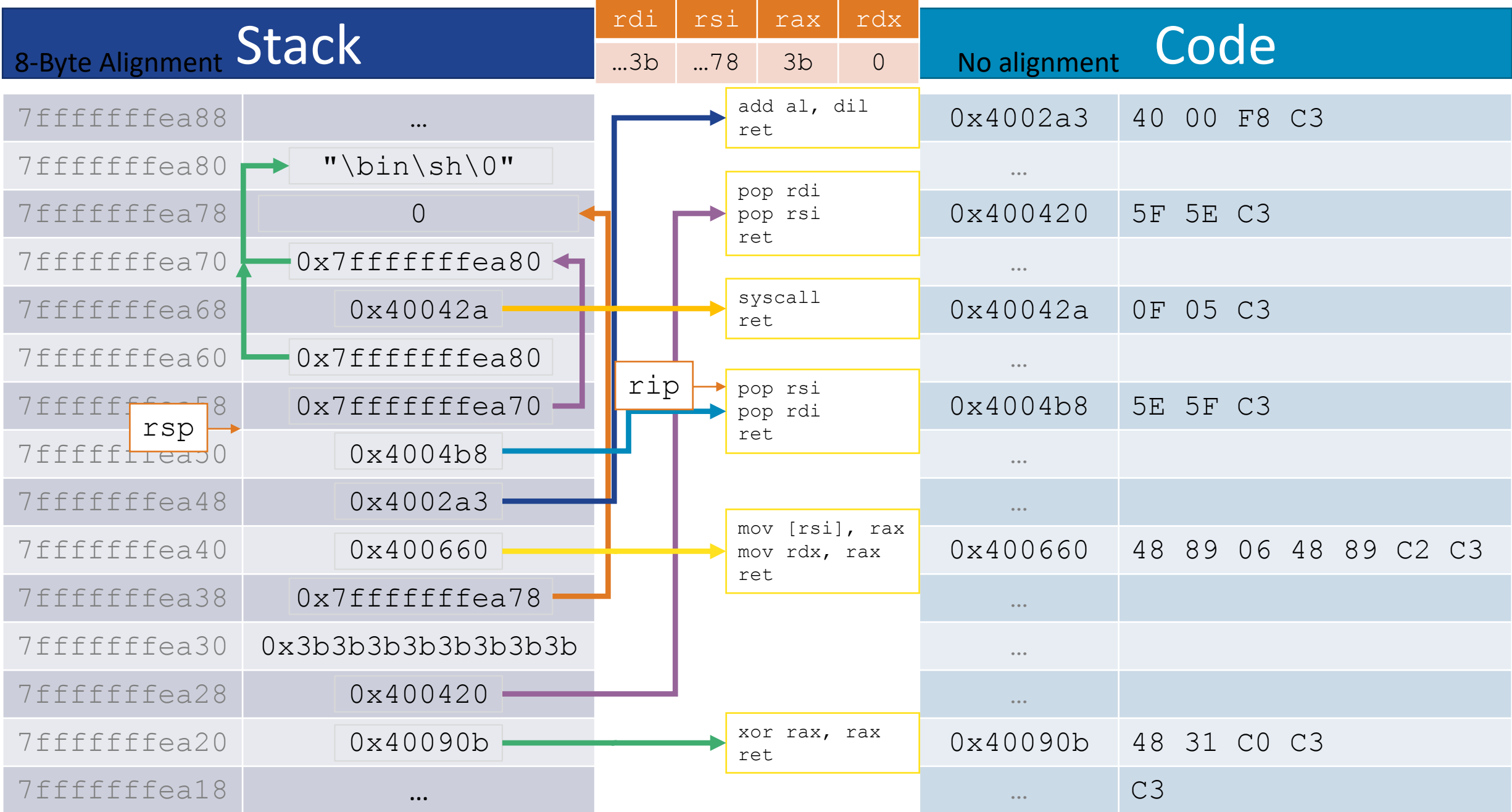


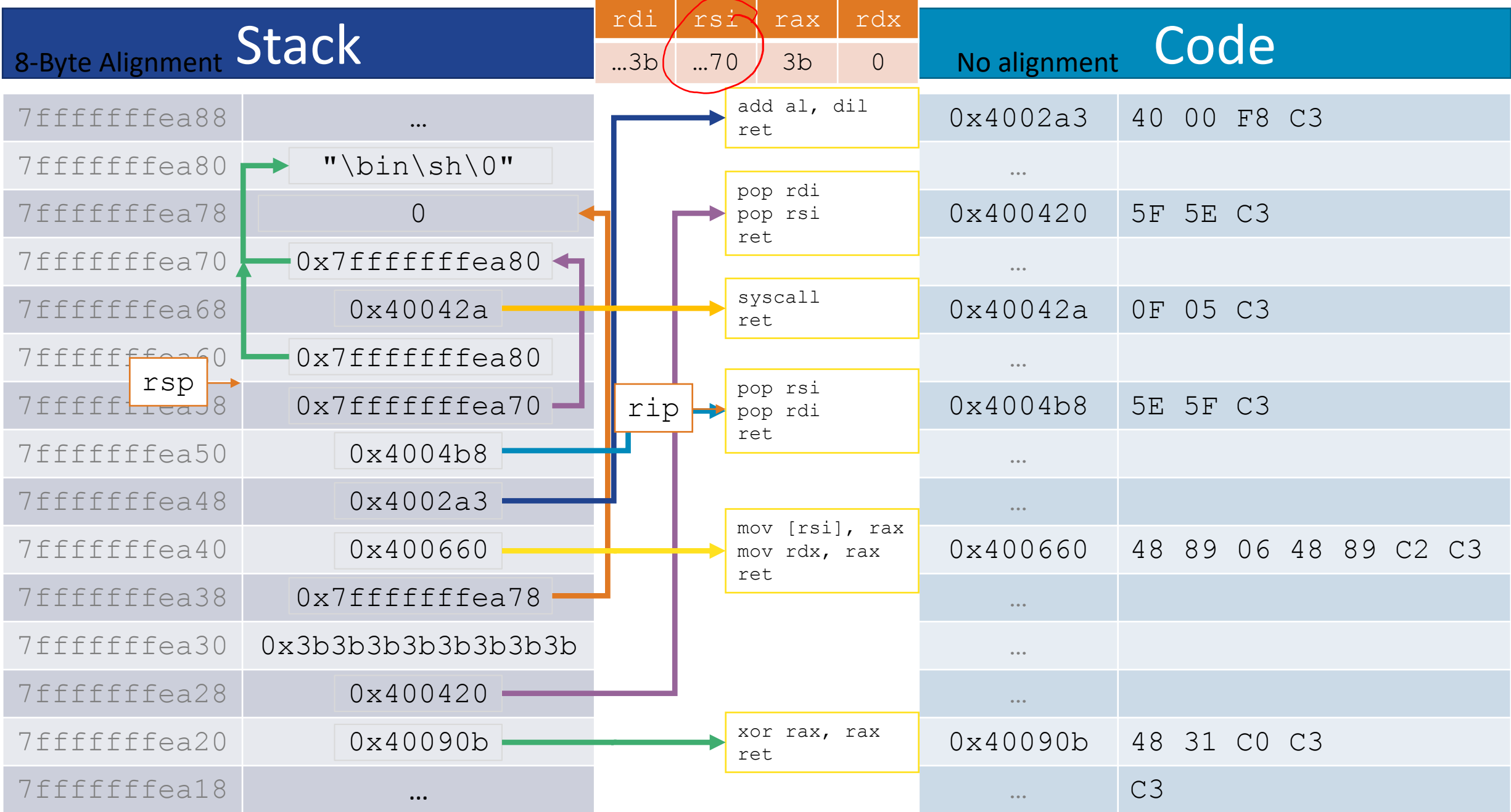


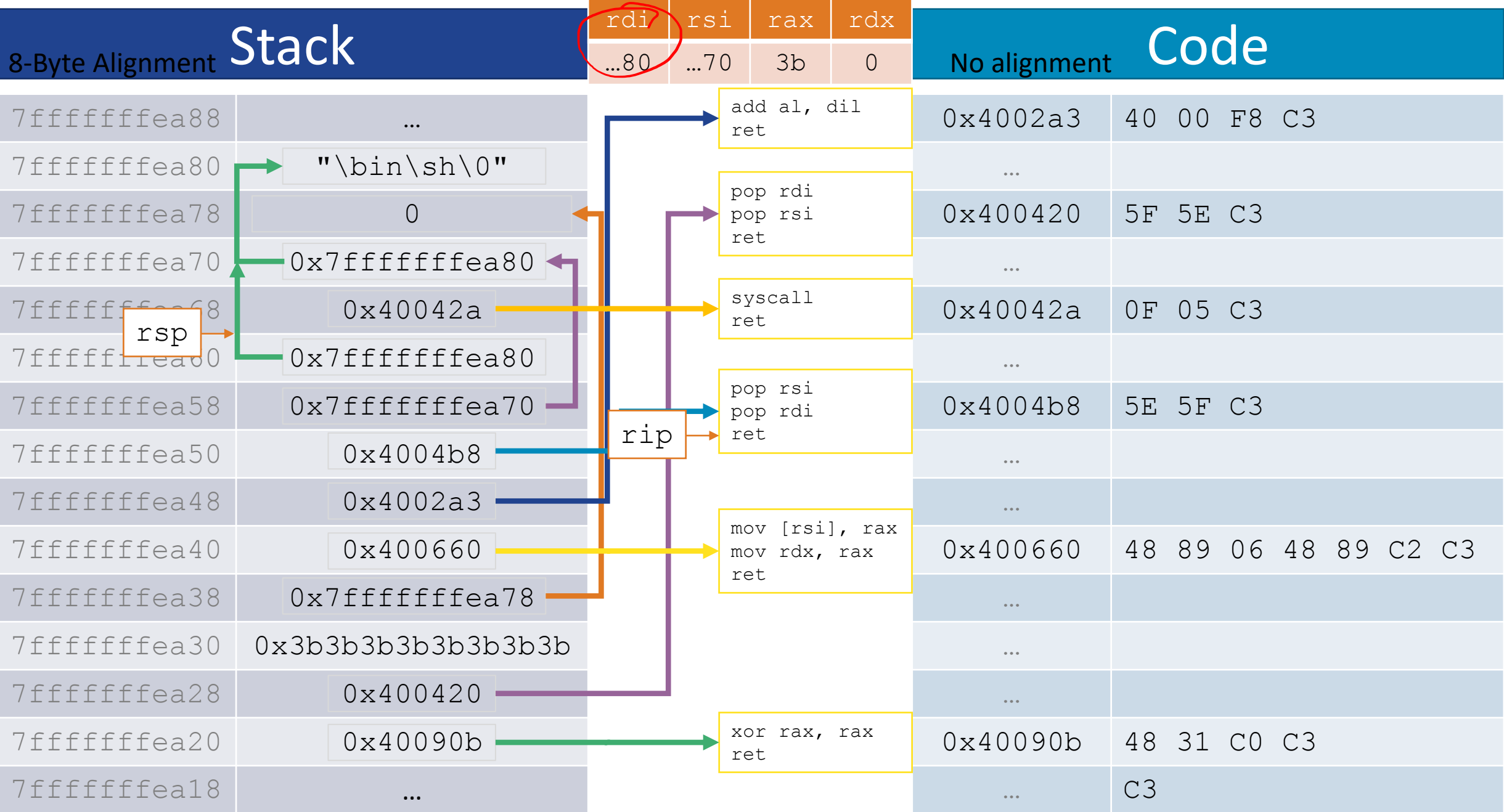


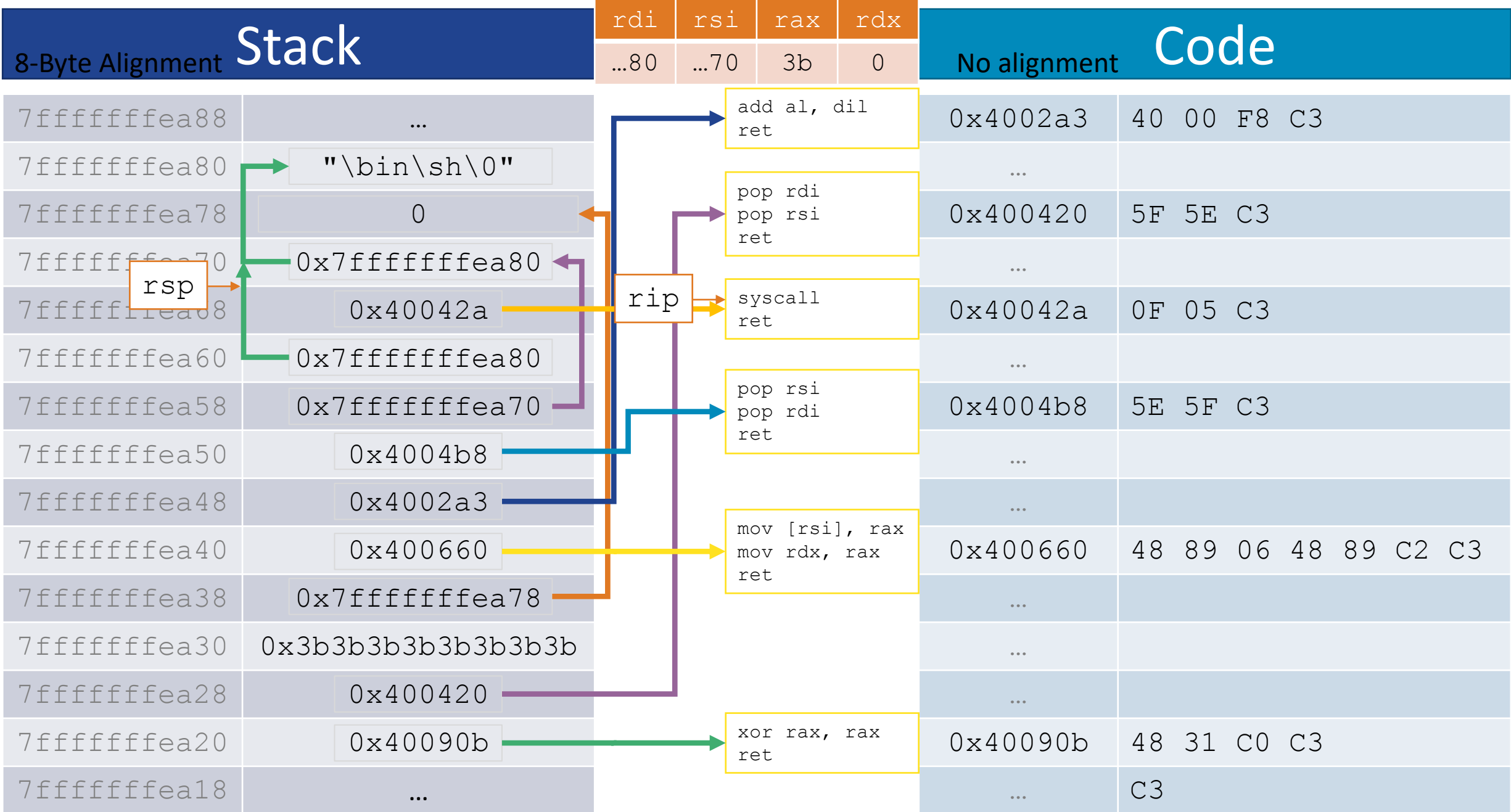


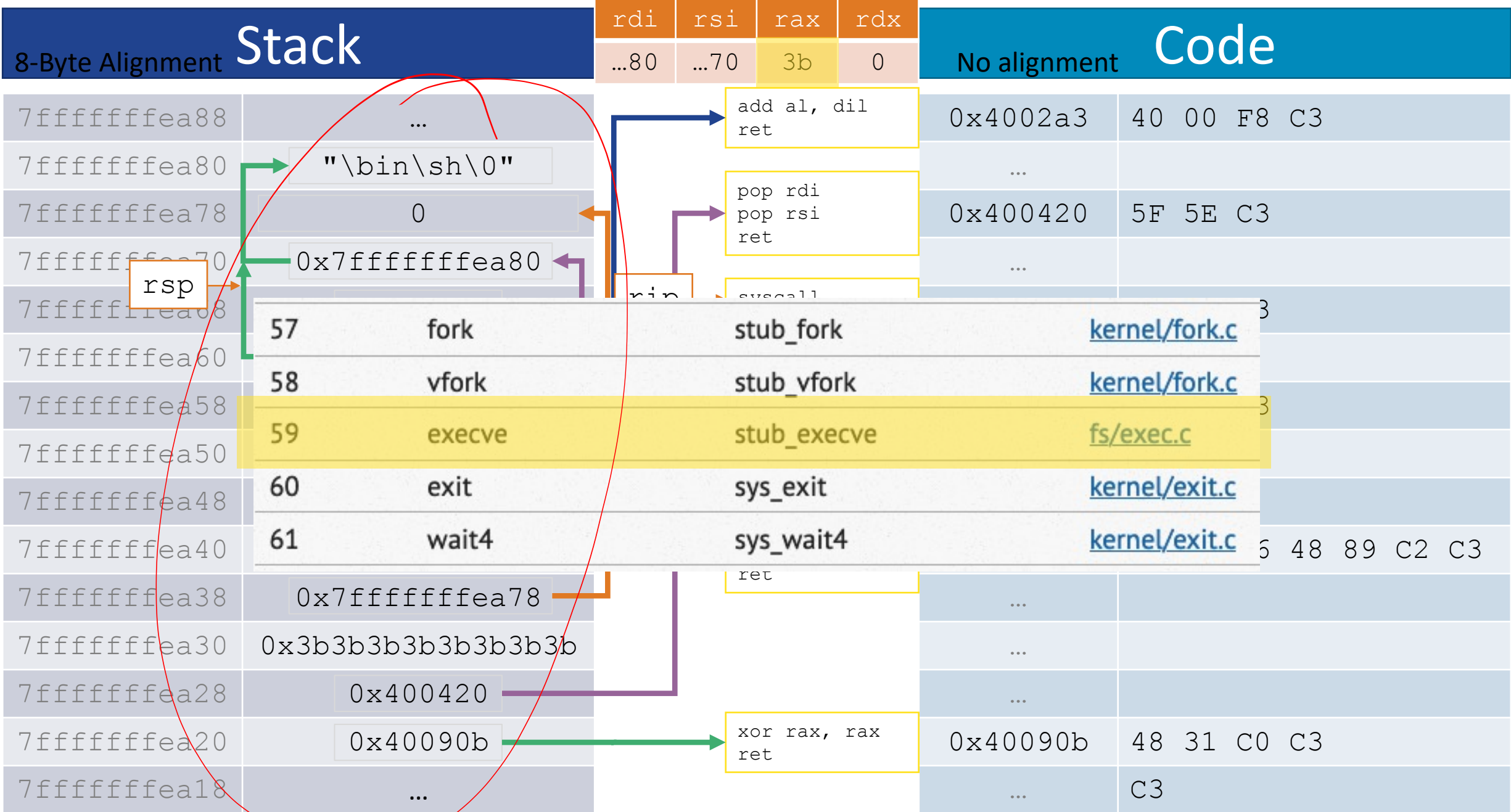












8-Byte Alignment Stack

rdi	rsi	rax	rdx
...80	...70	3b	0

No alignment Code

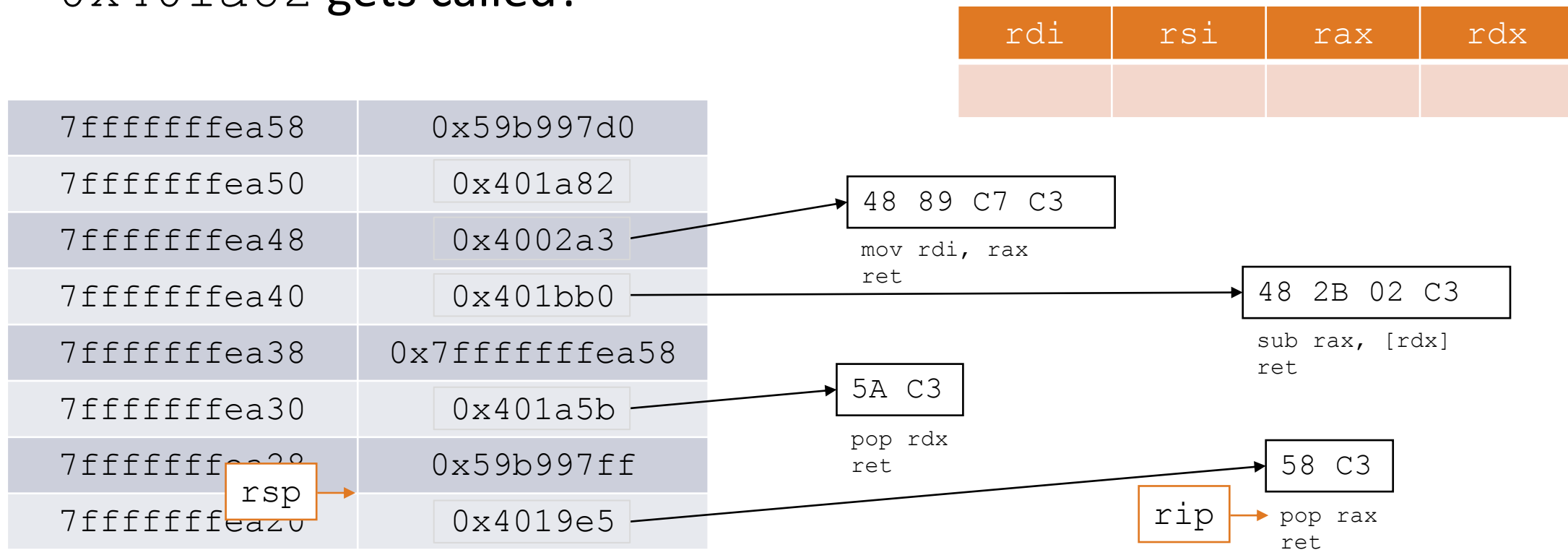
7fffffffefa88	...	add al, dil ret	0x4002a3	40 00 F8 C3
7fffffffefa80	"\bin\sh\0"		...	
7fffffffefa78	0	pop rdi pop rsi	0x400420	5F 5E C3

```
EXECVE(2)                                     System Calls Manual                                     EXECVE(2)
NAME
  execve - execute a file
SYNOPSIS
  #include <unistd.h>
  int
  execve(const char *path, char *const argv[], char *const envp[]);
DESCRIPTION
  execve() transforms the calling process into a new process. The new process is constructed from an ordinary file, whose name is pointed to by path, called the new process file. This file is either an executable object file, or a file of data for an interpreter. An executable object file consists of an identifying header, followed by pages of data representing the initial program (text) and initialized data pages. Additional pages may be specified by the header to be initialized with zero data; see a.out(5).
  An interpreter file begins with a line of the form:
      #! interpreter [arg ...]
  When an interpreter file is execve()'d, the system runs the specified interpreter. If any optional args are specified, they become the first (second, ...) argument to the interpreter. The name of the originally execve()'d file becomes the subsequent argument; otherwise, the name of the originally execve()'d file is the first argument. The original arguments to the invocation of the interpreter are shifted over to become the final arguments. The zeroth argument, normally the name of the execve()'d file, is left unchanged.
  The argument argv is a pointer to a null-terminated array of character pointers to null-terminated character strings. These strings construct the argument list to be made available to the new process. At least one argument must be present in the array; by custom, the first element should be the name of the executed program (for example, the last component of path).
man://execve(2) [R0]                                     1,1                                     Top
```

C3

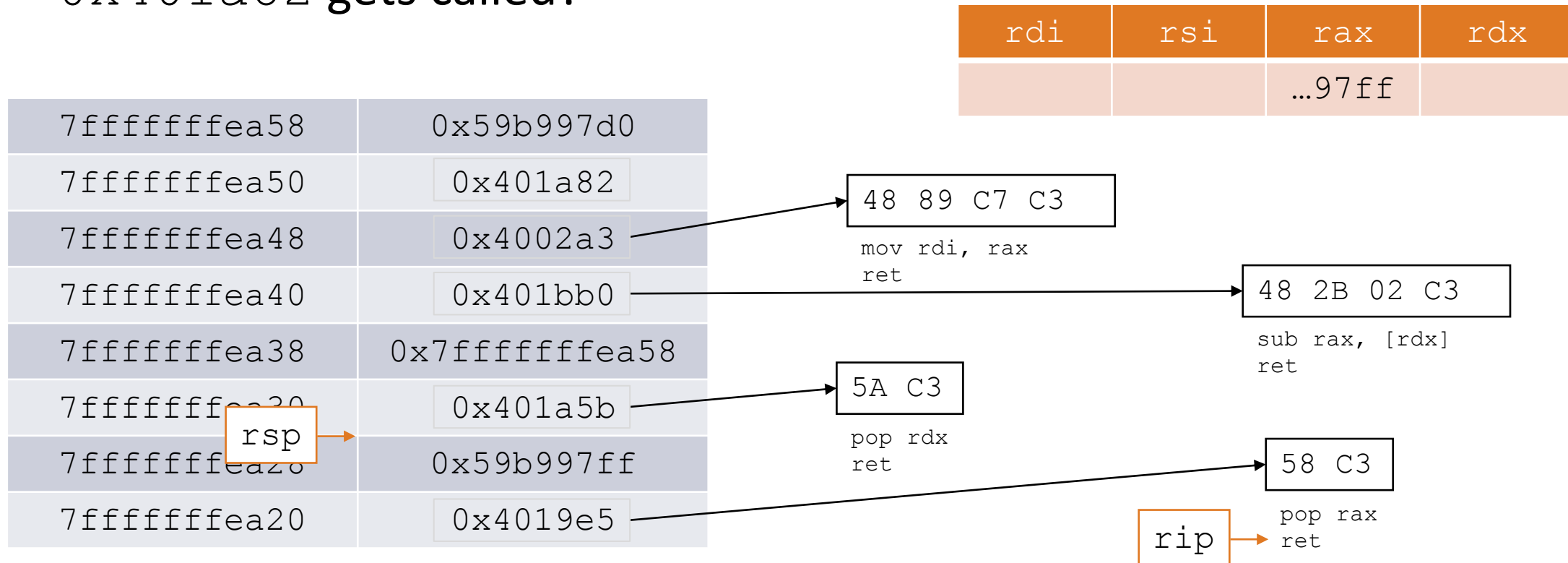
Practice with Return-Oriented Programming (ROP)

What are the values in the registers when the function at address 0x401a82 gets called?



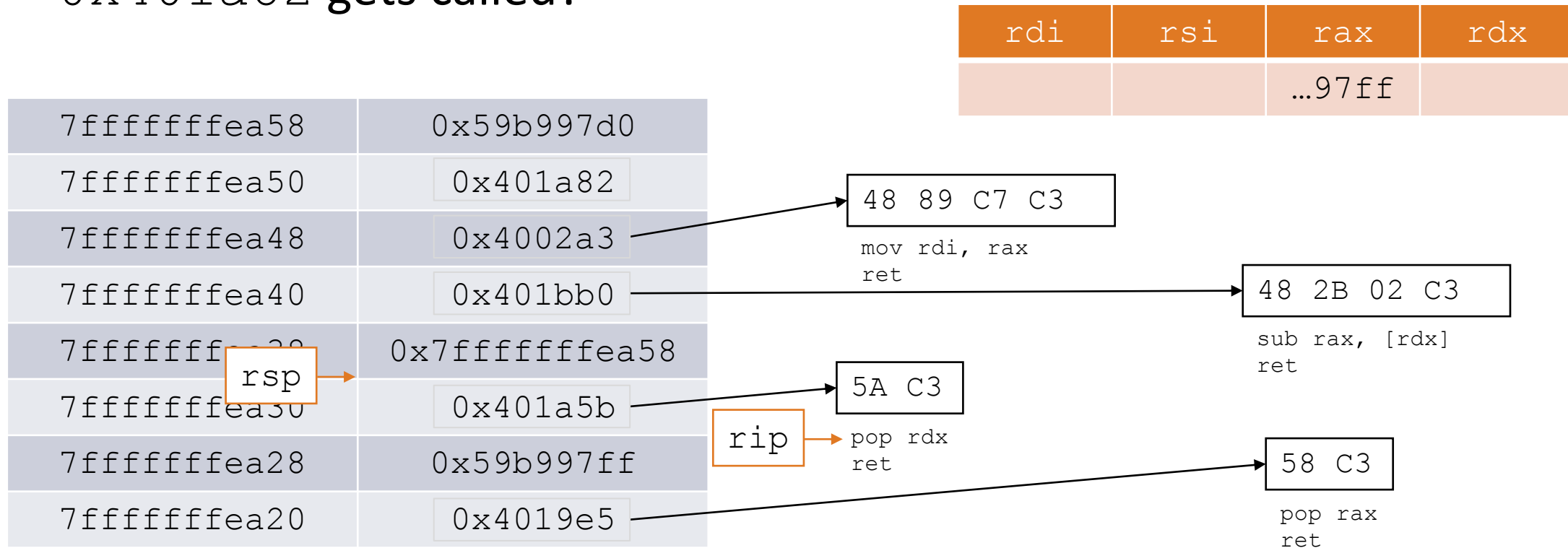
Practice with Return-Oriented Programming (ROP)

What are the values in the registers when the function at address 0x401a82 gets called?



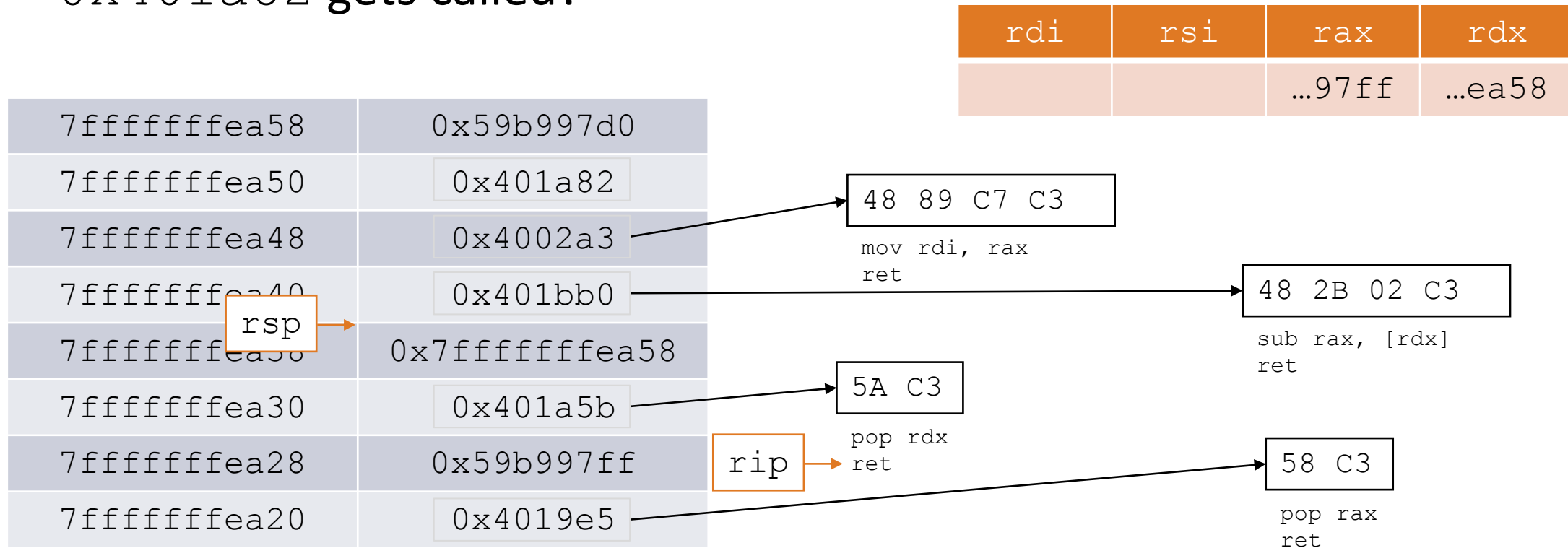
Practice with Return-Oriented Programming (ROP)

What are the values in the registers when the function at address 0x401a82 gets called?



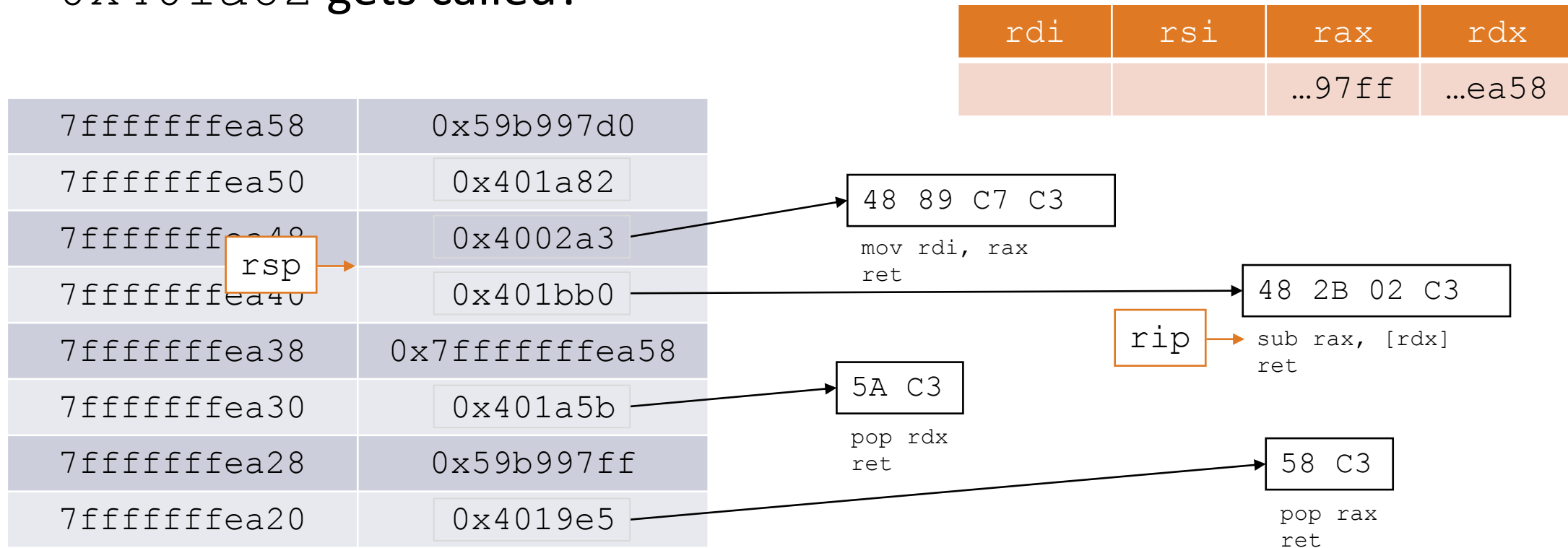
Practice with Return-Oriented Programming (ROP)

What are the values in the registers when the function at address 0x401a82 gets called?



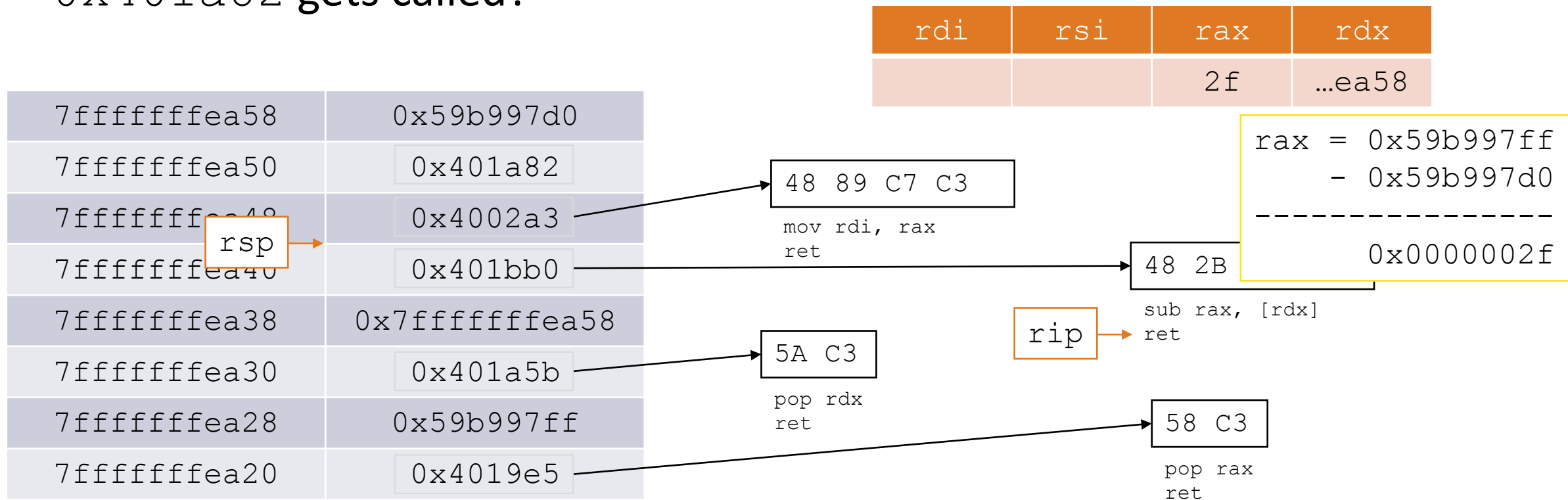
Practice with Return-Oriented Programming (ROP)

What are the values in the registers when the function at address 0x401a82 gets called?



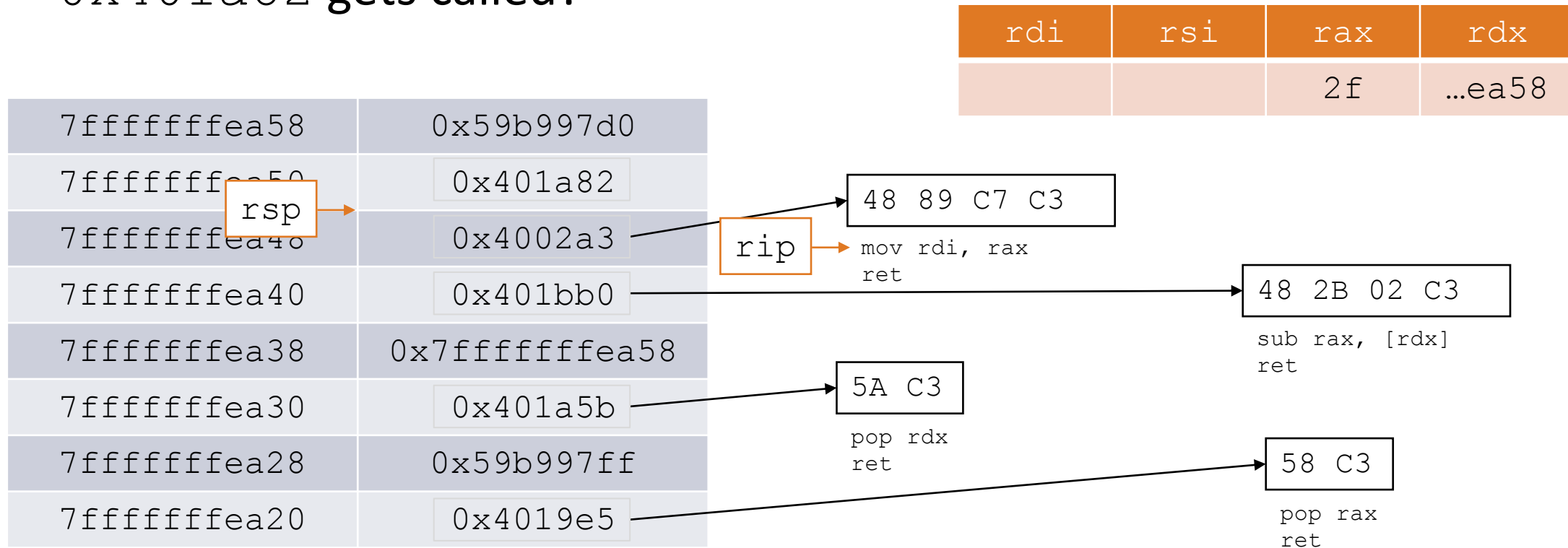
Practice with Return-Oriented Programming (ROP)

What are the values in the registers when the function at address 0x401a82 gets called?



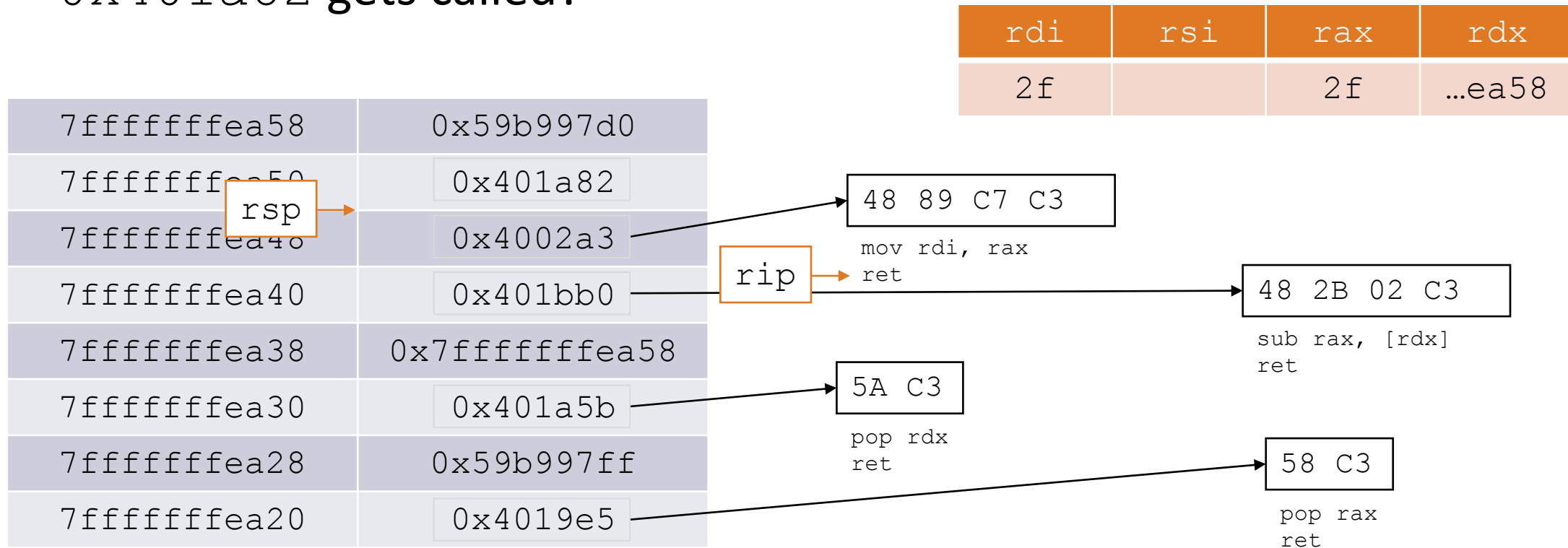
Practice with Return-Oriented Programming (ROP)

What are the values in the registers when the function at address 0x401a82 gets called?



Practice with Return-Oriented Programming (ROP)

What are the values in the registers when the function at address 0x401a82 gets called?



Defense #5: Address Space Layout Randomization

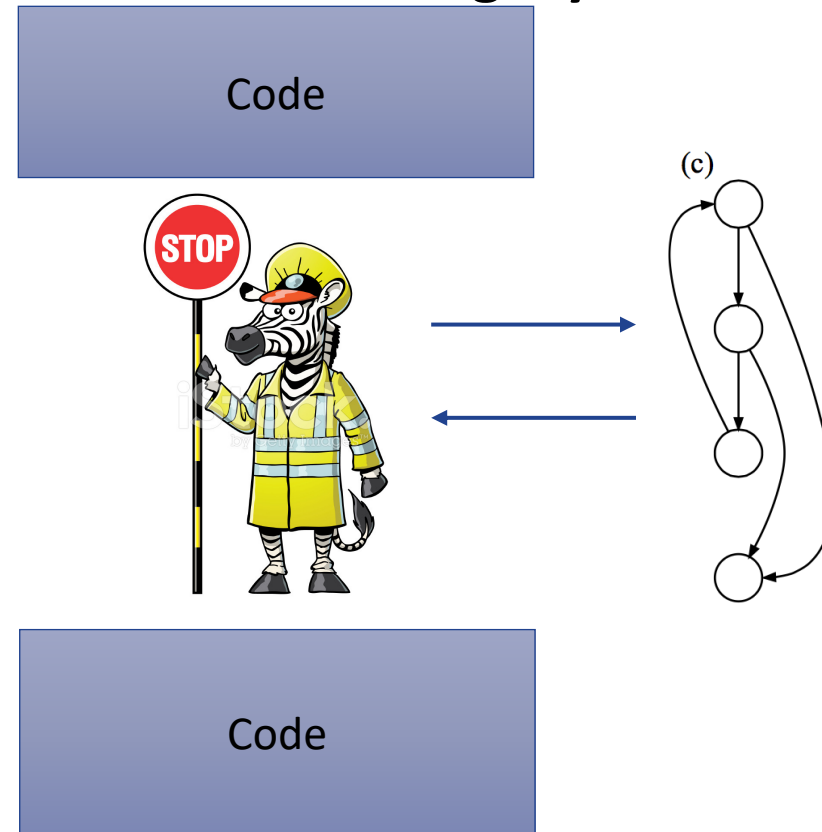


Other defenses

Gadget Elimination



Control Flow Integrity



The state of the world

Defenses:

- Use high-level languages
- Stack Canaries
- Memory tagging
- Address Space Layout Randomization
- Continuing research and development.

But they aren't perfect!

